

Notes for CIE 6002 Matrix Analysis

Binghao He

March 17, 2021

Contents

Lecture 2. Fundamental and Basic Operations	3
Subspace	3
Definition	3
Some important subspaces	3
Linear independence	3
Basis	3
Rank	3
Nonsingularity/Invertibility	3
Determinant	3
Vector Norm	3
Inner Product	4
Projection onto a Subspace	4
Projection	4
Orthogonal complement	4
Sum of subspaces	4
Orthogonal Set	4
Orthogonal/Unitary Matrix	4
Dimensional Theorem	5
Complexity of Matrix Computation	5
Flops	5
Big O notation	5
Lecture 3. Least Square Problem (LSP)	6
Formulation	6
LS problem	6
Linear representation of signals	6
AR and Polynomial Models	6
Autoregressive (AR) model	6
Polynomial model	6
Basis Representation	6
Basis matrix	6
Fourier basis	6
LTI System	7
Lienar time invariant system	7
System identification	7
Deconvolution	7
Toeplitz matrix	7

Circulant Matrix	7
Lecture 4. Eigenvalues and Eigenvectors	8
Diagonalizability	8
Similar matrices	8
Diagonalizable matrix	8
When Diagonalizable	8
Distinct eigenvalues	8
Repeated eigenvalues	9
Symmetric/Hermitian Matrices	9
Unitarily diagonalizable	10
Schur triangularization	10
Variational Characterization of Eigenvalues	11
Rayleigh quotient	11
Courant-Fischer theorem	11
Weyl's inequality	12
The Power Method	12
Lecture 5. PageRank	14
PageRank Model	14
Importance score	14
The model	14
PageRank Problem	14
Answer to Question 1	14
Answer to Question 2	15
Answer to Question 3	15
Answer to Question 4	15
Lecture 6. Positive Semidefinite Matrix	16
Properties of PSD Matrices	16
Propeties	16
Eigenvalues of PSD matrices	16
Matrix Inequality	17
Lecture 7. Singular Value Decomposition	18

Lecture 2. Fundamental and Basic Operations

Subspace

Definition

The set $S \subseteq \mathbb{R}^m$ is a subspace, if for $x, y \in S$
 $\alpha x + \beta y \in S, \quad \alpha, \beta \in \mathbb{R}^m$

This indicates

$$x_i \in S, \quad i = 1, \dots, N \implies \sum_{i=1}^N \alpha_i x_i \in S$$

Some important subspaces

1. Span: a collection of vectors

$$\text{span}\{a_1, \dots, a_n\} = \{y \in \mathbb{R}^m | y = \sum_{i=1}^n \alpha_i a_i, \alpha_i \in \mathbb{R}\}$$

2. Orthogonal complement: given a subspace $S \subseteq \mathbb{R}^m$

$$S_{\perp} = \{y \in \mathbb{R}^m | y^T x = 0, \forall x \in S\}$$

3. Range space of A (linear comb. of col. vectors)

$$\begin{aligned} R(A) &= \text{span}\{a_1, a_2, \dots, a_n\} \\ &= \{y \in \mathbb{R}^m | y = Ax, x \in \mathbb{R}^n\} \end{aligned}$$

4. Null space of A (linear comb. of col. vectors)

$$\text{Null}(A) = \{x \in \mathbb{R}^n | Ax = 0\}$$

Linear independence

Given a set $\{a_1, a_2, \dots, a_n\}$, we say they are linear independent if

$$\sum_{i=1}^n \alpha_i a_i = 0 \implies \alpha_i = 0, \forall i = 1, \dots, n$$

Basis

$\{b_1, b_2, \dots, b_k\} \subseteq \mathbb{R}^m$ is a basis for the subspace $S \subseteq \mathbb{R}^m$ if $S = \text{span}\{b_1, b_2, \dots, b_k\}$ and $\{b_1, b_2, \dots, b_k\}$ is linear independent.

The number of the vectors of the basis for S

$$k = \dim(S)$$

1. k : max number of linear indep. vectors in S
2. k : min number of vectors that can span S

Rank

For a matrix $A \in \mathbb{R}^{m \times n}$

$$\dim(R(A)) \leq \min\{m, n\}$$

Rank of a matrix (?)

$$\begin{aligned} \text{rank}(A) &= \dim(R(A)) \\ &= \text{max number of linear indep.} \\ &\quad \text{columns(rows) of } A \\ &= \text{rank}(A^T) \end{aligned}$$

And

$$\text{Null}(A) = (R(A^T))_{\perp}$$

Nonsingularity/Invertibility

A symmetric matrix $A \in \mathbb{R}^{m \times m}$ is non-singular if

$$Ax = 0 \implies x = 0$$

\implies columns of A are linearly indep.

$\implies \exists A^{-1} \in \mathbb{R}^{m \times m}$ that $AA^{-1} = A^{-1}A = I_m$

Some properties

1. $(AB)^{-1} = B^{-1}A^{-1}$
2. $(A^{-1})^T = (A^T)^{-1}$

Determinant

For a matrix $A \in \mathbb{R}^{m \times m}$,

Inductive

1. If $m = 1$, $\det(A) = A$
2. If $m > 1$, Define: $A_{ij} \in \mathbb{R}^{(m-1) \times (m-1)}$ is a sub-matrix that removes i th row and j th column of A .

Define: $c_{ij} = (-1)^{i+j} \det(A_{ij})$.

$$\begin{aligned} \det(A) &= \sum_{j=1}^m a_{ij} c_{ij} \quad \forall i = 1, \dots, m \\ &= \sum_{i=1}^m a_{ij} c_{ij} \quad \forall j = 1, \dots, m \end{aligned}$$

For $A, B \in \mathbb{R}^{m \times m}$

- $\det(AB) = \det(A)\det(B)$
- $\det(A) = \det(A^T)$
- $\det(\alpha A) = \alpha^m \det(A)$
- $\det(A) = 0 \iff A$ is singular
- $\det(A^{-1}) = 1/\det(A)$ if A is invertible

Vector Norm

A mapping $f : \mathbb{R}^> \rightarrow \mathbb{R}$ is a norm if it satisfies that $\forall x, y \in \mathbb{R}^m$

- $f(x) \geq 0$
- $f(x) = 0$ iff. $x = 0$
- $f(x + y) \leq f(x) + f(y)$
- $f(\alpha x) = |\alpha|f(x)$

Examples

- $\|x\|_2 = \sqrt{\sum_{i=1}^m |x_i|^2}$ (Euclidean norm)
- $\|x\|_1 = \sum_{i=1}^m |x_i|$
- $\|x\|_{\infty} = \max_{i=1, \dots, m} |x_i|$
- $\|x\|_p = \sqrt[p]{\sum_{i=1}^m |x_i|^p}$

Inner Product

$$\langle x, y \rangle = x^T y$$

$$\langle x, x \rangle = \|x\|_2^2$$

Cauchy-Schwarz Inequality

$$|x^T y| \leq \|x\|_2 \cdot \|y\|_2$$

Holder Inequality

$$|x^T y| \leq \|x\|_p \cdot \|y\|_q$$

where

$$1/p + 1/q = 1 \quad , \quad p \cdot q \geq 1$$

Projection onto a Subspace

Projection

Given a set $S \subseteq \mathbb{R}^m$ and a point $y \in \mathbb{R}^m$, the projection of y onto S is

$$y_S = \arg \min_{z \in S} \|z - y\|_2^2$$

Theorem 1

Suppose S is a subspace,

1. y_S exists and is unique
2. Let $y_S = \Pi_S(y)$ iff.
 $z^T(y_S - y) = 0 \quad \forall z \in S$

Orthogonal complement

$S_\perp = \{y \in \mathbb{R}^m | y^T x = 0, \forall x \in S\}$ is the orthogonal complement of subspace S .

- Orthogonal complement is a subspace
- $S \cap S_\perp = \{0\}$

Theorem 2

Let $S \subseteq \mathbb{R}^m$ be a subspace. For any $y \in \mathbb{R}^m$

$$y = \Pi_S(y) + \Pi_{S_\perp}(y)$$

Proof

$$y = \Pi_S(y) + \Pi_{S_\perp}(y)$$

$$\begin{aligned} \iff y - \Pi_S(y) &= \Pi_{S_\perp}(y) \\ \iff z^T(y - \Pi_S(y)) &= 0 \quad \forall z \in S_\perp \\ \iff z^T \Pi_S(y) &= 0 \quad \forall z \in S_\perp \end{aligned}$$

$$\therefore \Pi_S(y) \in S$$

$$\therefore z^T \Pi_S(y) = 0 \quad \forall z \in S_\perp$$

□

Sum of subspaces

For two sets $X, Y \subseteq \mathbb{R}^m$,

$$X + Y = \{x + y \in \mathbb{R}^m\}$$

For a subspace $S \subseteq \mathbb{R}^m$

1. $S + S_\perp = \mathbb{R}^m$
2. $\dim(S) + \dim(S_\perp) = m$
3. $(S_\perp)_\perp = S$

Proof

1. Comes from Theorem 2
2. Let $\{a_1, a_2, \dots, a_k\}$ be a basis for S ($\dim(S) = k$) and $\{b_1, b_2, \dots, b_\ell\}$ be a basis for S_\perp ($\dim(S_\perp) = \ell$).
 $\{a_1, a_2, \dots, a_k\} \cup \{b_1, b_2, \dots, b_\ell\}$ is
 (a) linear independent
 (b) Span $S + S_\perp = \mathbb{R}^m$
 Therefore, $\dim(S) + \dim(S_\perp) = m$
3. Obvious

□

Orthogonal Set

$\{a_1, a_2, \dots, a_n\}$ is orthogonal/orthonormal set if

1. $\|a_i\|_2 = 1 \quad \forall i = 1, \dots, n$
2. $a_i^T a_j = 0 \quad \forall i \neq j$

The orthogonal set is linear independent.

$$y = \sum_{i=1}^n \alpha_i a_i \implies \alpha_i = a_i^T y = \alpha \|a_i\|_2^2$$

Given a linear independent set $\{a_1, a_2, \dots, a_n\}$. The procedure of finding an orthogonal set $\{q_1, q_2, \dots, q_n\}$ so that $\text{span}\{a_1, a_2, \dots, a_n\} = \text{span}\{q_1, q_2, \dots, q_n\}$ is called "Gram-Schmidt" procedure.

Orthogonal/Unitary Matrix

$Q = [q_1, q_2, \dots, q_m] \in \mathbb{R}^{n \times m}$, $q_i, i = 1, \dots, m$ are orthogonal.

- If $m = n$,
 Q is orthogonal matrix $\implies Q^T Q = I_m = Q Q^T$.
- If $n > m$,
 semi-orthogonal $\implies Q^T Q = I_m \neq Q Q^T$.

$Q = [q_1, q_2, \dots, q_m] \in \mathbb{C}^{n \times m}$, $q_i, i = 1, \dots, m$ are orthogonal.

- If $m = n$,
 Q is unitary matrix $\implies Q^H Q = I_m = Q Q^H$.
- If $n > m$,
 semi-unitary $\implies Q^H Q = I_m \neq Q Q^H$.

□

If $Q \in \mathbb{R}^{n \times m}$ ($n > m$) is semi-orthogonal, $\exists \tilde{Q} \in \mathbb{R}^{n \times (n-m)}$ so that $\begin{bmatrix} Q & \tilde{Q} \end{bmatrix}$ is orthogonal.

Dimensional Theorem

For a matrix $A \in \mathbb{R}^{m \times n}$, it holds

$$\dim(R(A)) + \dim(\text{Null}(A)) = n$$

Proof

$$\begin{aligned} \dim(R(A^T)) + \dim(R(A^T)_{\perp}) &= n \\ \dim(R(A^T)) = \text{rank}(A^T) = \text{rank}(A) &= \dim(R(A)) \\ \dim(R(A^T)_{\perp}) = \text{Null}(A) &= n \end{aligned}$$

□

Complexity of Matrix Computation

Flops

Arithmetic operations including addition, subtraction, multiplication, division.

Big O notation

Big O notation: "order" of the complexity.

$$f(n) = O(g(n))$$

for some $g(n)$ if \exists a constant $C > 0$ and n_0 so that for $n > n_0$

$$|f(n)| \leq C \cdot |g(n)|$$

Examples: for $x \in \mathbb{R}^m$, $y \in \mathbb{R}^m$, $A \in \mathbb{R}^{n \times m}$, $B \in \mathbb{R}^{m \times p}$

$x + y$	m flops	$O(m)$
$x^T y$	$2m - 1$ flops	$O(m)$
Ax	$n \cdot (2m - 1)$ flops	$O(nm)$
AB	$p \cdot n \cdot (2m - 1)$ flops	$O(pnm)$

Lecture 3. Least Square Problem (LSP)

Formulation

LS problem

Given $y \in \mathbb{R}^m$, $A \in \mathbb{R}^{m \times n}$

$$\min_{x \in \mathbb{R}^n} \|y - Ax\|_2^2$$

Linear representation of signals

$$y = Ax$$

In practice,

$$y = Ax + v$$

where

- y is the observation (Given)
- A is the model (Given)
- x is the parameter to be estimated/determined
- v is the noise follow multi-variate Gaussian distribution

$$\text{pdf}(v) \propto e^{-\|v\|^2}$$

\Rightarrow Maximum likelihood estimator

$$\max_x P(y|x) \propto \min_x \|y - Ax\|_2^2$$

AR and Polynomial Models

Autoregressive (AR) model

Given a time series y_t , $t = 0, 1, 2, \dots$,

$$y_t = \underbrace{\alpha_1 y_{t-1} + \alpha_2 y_{t-2} + \dots + \alpha_p y_{t-p}}_{\text{AR model}} + \underbrace{v_t}_{\text{noise}}$$

If $\alpha_1, \alpha_2, \dots, \alpha_p$ are known, estimate y_t by

$$\hat{y}_t = \underbrace{\alpha_1 y_{t-1} + \alpha_2 y_{t-2} + \dots + \alpha_p y_{t-p}}_{\text{historical observation}}$$

Collect

$$\underbrace{\begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_T \end{bmatrix}}_{y \in \mathbb{R}^T} = \underbrace{\begin{bmatrix} y_0 & & & \\ y_1 & y_0 & & \\ \vdots & \ddots & \ddots & \\ y_{p-1} & \cdots & \cdots & y_0 \\ \vdots & \ddots & \ddots & \vdots \\ y_{T-1} & \cdots & \cdots & y_T \end{bmatrix}}_{A \in \mathbb{R}^{T \times p}} \underbrace{\begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_p \end{bmatrix}}_{x \in \mathbb{R}^p} + \underbrace{\begin{bmatrix} v_0 \\ v_1 \\ \vdots \\ v_T \end{bmatrix}}_{v \in \mathbb{R}^T}$$

\Rightarrow Solve LS problem to obtain α_i , $i = 1, \dots, p$

Polynomial model

$$y_t = \underbrace{\alpha_0 + \alpha_1 t + \alpha_2 t^2 + \dots + \alpha_p t^p}_{\text{polynomial model}} + \underbrace{v_t}_{\text{noise}}$$

Collect

$$\underbrace{\begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_T \end{bmatrix}}_{y \in \mathbb{R}^T} = \underbrace{\begin{bmatrix} 1 & 0 & \cdots & 0 \\ 1 & 1 & \cdots & 1 \\ & \vdots & & \\ 1 & t & \cdots & t^p \\ & \vdots & & \\ 1 & T-1 & \cdots & (T-1)^p \end{bmatrix}}_{A \in \mathbb{R}^{T \times p}} \underbrace{\begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_p \end{bmatrix}}_{x \in \mathbb{R}^p} + \underbrace{\begin{bmatrix} v_0 \\ v_1 \\ \vdots \\ v_T \end{bmatrix}}_{v \in \mathbb{R}^T}$$

\Rightarrow Solve LS problem to obtain α_i , $i = 1, \dots, p$

Basis Representation

Basis matrix

For $y \in \mathbb{R}^m$, given a set of basis vectors (can be given or learned)

$$\Phi = \{\phi_1, \phi_2, \dots, \phi_n\} \in \mathbb{R}^{m \times n}$$

We model

$$y = \sum_{i=1}^n \phi_i x_i = \Phi x$$

- $m \gg n$, assume data lies in a low-dimensional subspace
- $m \ll n$, dictionary $\Rightarrow x$ is sparse

Fourier basis

Discrete Fourier Transform (DFT)

Inverse Discrete Fourier Transform (IDFT)

$$\Phi = [\phi_1, \dots, \phi_n] \in \mathbb{C}^{n \times n}$$

$$\phi_i = \frac{1}{\sqrt{n}} \begin{bmatrix} 1 \\ e^{j \frac{2\pi}{n}(i-1)} \\ e^{j \frac{2\pi}{n}(i-1) \times 2} \\ \vdots \\ e^{j \frac{2\pi}{n}(i-1) \times n-1} \end{bmatrix} \in \mathbb{R}^n$$

- Φ is IDFT matrix, Φ^H is DFT matrix.
- Φ is unitary matrix.

LTI System

Lienar time invariant system

$$\begin{aligned} - \text{Linearity} & \quad \begin{cases} x_t^1 \rightarrow \boxed{\text{LTI}} \rightarrow y_t^1 \\ x_t^2 \rightarrow \boxed{\text{LTI}} \rightarrow y_t^2 \\ \alpha x_t^1 + \beta x_t^2 \rightarrow \boxed{\text{LTI}} \rightarrow \alpha y_t^1 + \beta y_t^2 \end{cases} \\ - \text{Time Invariant} & \quad x_{t-d} \rightarrow \boxed{\text{LTI}} \rightarrow y_{t-d} \end{aligned}$$

Impulse $\delta(t)$

$$\begin{cases} \delta(t) = 0, t \neq 0 \\ \int_{-\infty}^{+\infty} \delta(t) dt = 1 \end{cases}$$

LTI system is characterized by an impulse response

$$h_t, \quad t = 0, \dots, p-1$$

$\Rightarrow x_t$ and y_t is related through "convolution"

$$\begin{aligned} y_t &= \sum_{i=1}^p h_i x_{t-i} \\ &= h_t * x_t \end{aligned}$$

$$x_t \rightarrow \overbrace{\boxed{h_t}}^{\text{LTI}} \rightarrow y_t$$

System identification

Given $\{y_t\}_{t=0}^{T-1}$ and $\{x_t\}_{t=0}^{T-1}$, find the impulse response $\{h_t\}_{t=0}^p$

$$\begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_p \\ \vdots \\ y_T \end{bmatrix} = \begin{bmatrix} x_0 & & & & \\ x_1 & x_0 & & & \\ \vdots & \vdots & \ddots & \ddots & \\ x_p & \cdots & \cdots & x_0 & \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ x_{T-1} & \cdots & \cdots & x_{T-p} & \end{bmatrix} \begin{bmatrix} h_0 \\ h_1 \\ \vdots \\ h_p \end{bmatrix} + \text{noise}$$

\Rightarrow Solve LS problem to obtain $\{h_t\}_{t=0}^p$

Deconvolution

Given $\{y_t\}_{t=0}^{T-1}$ and $h_{t=0}^p$, find out $\{x_t\}_{t=0}^{T-1}$

Put them into a linear system of equations

$$\begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_p \\ \vdots \\ y_T \end{bmatrix} = \underbrace{\begin{bmatrix} h_0 & 0 & \cdots & & 0 \\ h_1 & h_0 & 0 & \cdots & \ddots \\ \vdots & \vdots & & \ddots & \\ h_p & h_{p-1} & \cdots & h_1 & h_0 \\ \vdots & \ddots & \ddots & \vdots & \vdots \\ 0 & \cdots & h_p & h_{p-1} & \cdots & h_1 & h_0 \end{bmatrix}}_{A \in \mathbb{R}^{T \times T}} \begin{bmatrix} x_0 \\ x_1 \\ \vdots \\ x_p \\ \vdots \\ x_T \end{bmatrix} + \text{noise}$$

\Rightarrow Solve LS problem to obtain $\{x_t\}_{t=0}^{T-1}$

Toeplitz matrix

$$A = \begin{bmatrix} a_0 & a_{-1} & \cdots & \cdots & a_{-n+1} \\ a_1 & a_0 & a_{-1} & \cdots & a_{-n+2} \\ a_2 & a_1 & a_0 & a_{-1} & \ddots \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ a_{n-1} & \cdots & a_2 & a_1 & a_0 \end{bmatrix} \in \mathbb{R}^{n \times n}$$

Circulant Matrix

Circulant matrix is a special case of Toeplitz matrix

$$H = \begin{bmatrix} h_0 & h_{n-1} & \cdots & \cdots & h_1 \\ h_1 & h_0 & h_{n-1} & \cdots & h_2 \\ h_2 & h_1 & h_0 & h_{n-1} & \ddots \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ h_{n-1} & \cdots & h_2 & h_1 & h_0 \end{bmatrix} \in \mathbb{R}^{n \times n}$$

$$H \xLeftrightarrow{O(n \log n)} H^{-1}$$

The key property:

$$H \cdot \underbrace{\phi_i}_{\text{IDFT}} = d_i \cdot \phi_i, \quad d_i \in \mathbb{R}$$

Let

$$D = \begin{bmatrix} d_1 & & \\ & \ddots & \\ & & d_n \end{bmatrix}$$

Then

$$\begin{aligned} H \cdot \Phi &= H \cdot [\phi_1, \dots, \phi_n] \\ &= [d_1 \phi_1, \dots, d_n \phi_n] \\ &= \Phi \cdot D \end{aligned}$$

$$\underbrace{H = \Phi \cdot D \cdot \Phi^H}_{\text{eigen-decomposition of } H} \iff \Phi^H H \Phi = D$$

Lecture 4. Eigenvalues and Eigenvectors

Definition (Eigenvalues and Eigenvector)

Let $A \in \mathbb{R}^{m \times m}$, if a scalar $\lambda \in \mathbb{C}$ and a nonzero vector $v \in \mathbb{C}^m$ satisfy the equation

$$Av = \lambda v$$

then λ is called the eigenvalue of A and v is called the eigenvector of A associated with λ .

\Rightarrow

$$\underbrace{(A - \lambda I)}_{\text{singular}} v = 0$$

$P_A(\lambda) \triangleq \det(A - \lambda I) = \prod_{i=1}^m (\lambda - \lambda_i)$ is a polynomial of λ with degree m . The equation $P_A(\lambda) \triangleq \det(A - \lambda I) = 0$ is called the characteristic equation of A .

For each λ_i find v_i such that

$$(A - \lambda_i I)v_i = 0$$

$v_i \in \text{Null}(A - \lambda_i I)$ and $\text{Null}(A - \lambda_i I)$ is the eigenspace for λ_i

- Even if $A \in \mathbb{R}^{m \times m}$ is real-valued, λ_i and v_i could be complex. *e.g.*

$$A = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix} \Rightarrow \lambda = \pm j$$

- But if $\lambda_i \in \mathbb{R}$ is real-valued, $v_i \in \mathbb{R}^m$ should be in \mathbb{R}^m .

\Rightarrow (Eigenvalue decomposition)

There exists an invertible matrix $V \in \mathbb{R}^{m \times m}$ and a diagonal matrix $D \in \mathbb{R}^{m \times m}$ s.t.

$$A = \underbrace{VDV^{-1}}$$

eigenvalue decomposition (EVD)

\Rightarrow

$$AV = VD$$

$$V = [v_1, v_2, \dots, v_m] \quad , \quad D = \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_m \end{bmatrix}$$

\Leftrightarrow

$$Av_i = \lambda_i v_i, \quad \forall i = 1, \dots, m$$

\Rightarrow

A is diagonalizable *iff.* there is a set of linear indep. vectors which are eigenvectors of A . **Properties**

If a matrix $A \in \mathbb{R}^{m \times m}$ is diagonalizable

1. $\det(A) = \prod_{i=1}^m \lambda_i$
2. $\text{Tr}(A) = \sum_{i=1}^m \lambda_i$
3. $\text{rank}(A)$ = the number of non-zero eigenvalues

Proof

1. $\det(A) = \det(VDV^{-1}) = \det(V)\det(D)\det(V^{-1}) = \det(D) = \prod_{i=1}^m \lambda_i$
2. $\text{Tr}(A) = \text{Tr}(VDV^{-1}) = \text{Tr}(DV^{-1}V) = \text{Tr}(D) = \sum_{i=1}^m \lambda_i$

□

Diagonalizability

Similar matrices

Definition (Similar matrices)

$A \in \mathbb{R}^{m \times m}$ and $B \in \mathbb{R}^{m \times m}$ are similar to each other if there exists a nonsingular $S \in \mathbb{R}^{m \times m}$ such that

$$B = SAS^{-1}$$

If A and B are similar to each other, then they have the same eigenvalues.

Proof

$$\begin{aligned} P_B(\lambda) &= \det(B - \lambda I) \\ &= \det(SAS^{-1} - \lambda I) \\ &= \det(S(A - \lambda I)S^{-1}) \\ &= \det(S)\det(A - \lambda I)\det(S^{-1}) \\ &= \det(A - \lambda I) \\ &= P_A(\lambda) \end{aligned}$$

□

Diagonalizable matrix

Definition (Diagonalizable)

A matrix $A \in \mathbb{R}^{m \times m}$ is diagonalizable if A is similar to a diagonal matrix.

When Diagonalizable

Distinct eigenvalues

Let $\{\lambda_1, \lambda_2, \dots, \lambda_k\}_{k \in M}$ ($k \geq 2$) be a subset of eigenvalues of A and $\lambda_i \neq \lambda_j \quad \forall i \neq j$. Then the corresponding eigenvectors, denoted with v_1, \dots, v_k respectively, are linearly independent.

If A has $\lambda_1 \neq \lambda_2 \neq \dots \neq \lambda_k$, then the corresponding eigenvectors v_1, \dots, v_k are linearly independent.

Proof

Suppose v_1, \dots, v_k are linearly **dependent**. Then there exists $\alpha_1, \alpha_2, \dots, \alpha_k$ not all zero (without losing generality, assume $\alpha_1 \neq 0$) s.t.

$$\sum_{i=1}^k \alpha_i v_i = 0$$

Then

$$A \cdot \left(\sum_{i=1}^k \alpha_i v_i \right) = \sum_{i=1}^k \alpha_i Av_i = \sum_{i=1}^k \alpha_i \lambda_i v_i = 0$$

$$\sum_{i=1}^{k-1} \alpha_i \lambda_i v_i + \alpha_k \lambda_k v_k = 0$$

On the other hand,

$$\lambda_k \cdot \left(\sum_{i=1}^k \alpha_i v_i \right) = \sum_{i=1}^k \alpha_i \lambda_k v_i = 0$$

$$\sum_{i=1}^{k-1} \alpha_i \lambda_k v_i + \alpha_k \lambda_k v_k = 0$$

(1) - (2):

$$\sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) v_i = 0$$

Then

$$\begin{aligned} A \cdot \left(\sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) v_i \right) &= \sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) A v_i \\ &= \sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) \lambda_i v_i \\ &= 0 \end{aligned}$$

$$\sum_{i=1}^{k-2} \alpha_i (\lambda_i - \lambda_k) \lambda_i v_i + \alpha_{k-1} (\lambda_{k-1} - \lambda_k) \lambda_{k-1} v_{k-1} = 0 \quad (3)$$

On the other hand

$$\lambda_{k-1} \cdot \left(\sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) v_i \right) = \sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) \lambda_{k-1} v_i = 0$$

$$\sum_{i=1}^{k-2} \alpha_i (\lambda_i - \lambda_k) \lambda_{k-1} v_i + \alpha_{k-1} (\lambda_{k-1} - \lambda_k) \lambda_{k-1} v_{k-1} = 0$$

(4)

(3) - (4):

$$\sum_{i=1}^{k-2} \alpha_i (\lambda_i - \lambda_k) (\lambda_i - \lambda_{k-1}) v_i = 0$$

Repeat until $i = 1$

$$\alpha_1 \left(\prod_{i=1}^k (\lambda_1 - \lambda_i) \right) v_1 = 0$$

Since $\alpha_1 \neq 0$ and $\prod_{i=1}^k (\lambda_1 - \lambda_i) \neq 0$, so $v_1 = 0$ (Contradict!)

Therefore, v_1, \dots, v_k are linearly **independent** \square

If $A \in \mathbb{R}^{m \times m}$ has m distinct eigenvalues, then A is diagonalizable.

Repeated eigenvalues

Let λ be an eigenvalue of A repeated r times. Define

- Algebra multiplicity: r .
- Geometric multiplicity: $s = \dim(\text{Null}(A - \lambda I))$

Then

$$s \leq r$$

(1) *Proof*

Let $\dim(\text{Null}(A - \bar{\lambda}I)) = s$ and $\{q_1, \dots, q_s\}$ be an orthogonal basis for $\text{Null}(A - \bar{\lambda}I)$.

Then $Q_1 = [q_1, \dots, q_s] \in \mathbb{R}^{m \times s}$ is semi-unitary and $\exists Q_2 \in \mathbb{R}^{m \times (m-s)}$ s.t. $Q \triangleq [Q_1, Q_2]$ is unitary.

Consider

$$Q^H A Q = \begin{bmatrix} Q_1^H A Q_1 & Q_1^H A Q_2 \\ Q_2^H A Q_1 & Q_2^H A Q_2 \end{bmatrix} = \begin{bmatrix} \bar{\lambda} I_s & Q_1^H A Q_2 \\ 0 & Q_2^H A Q_2 \end{bmatrix}$$

Then

$$\begin{aligned} P_A(\lambda) &= P_{Q^H A Q}(\lambda) \\ &= \det(Q^H A Q - \lambda I) \\ &= \det \left(\begin{bmatrix} (\bar{\lambda} - \lambda) I_s & Q_1^H A Q_2 \\ 0 & Q_2^H A Q_2 - \lambda I_{m-s} \end{bmatrix} \right) \\ &= \underbrace{(\bar{\lambda} - \lambda)^s}_{\bar{\lambda} \text{ repeats } s \text{ times}} \det(Q_2^H A Q_2 - \lambda I_{m-s}) \\ &= 0 \end{aligned}$$

So

$$r \geq s$$

\square

If $s = r$, A is diagonalizable.

Symmetric/Hermitian Matrices

- **Symmetric** $A = A^T$ $A \in \mathbb{R}^{m \times n}$

- **Hermitian** $A = A^H$ $A \in \mathbb{C}^{m \times n}$

Property 1: The eigenvalues of a Hermitian matrix $A \in \mathbb{C}^{m \times n}$ are real-valued.

Proof

Let λ be the eigenvalue of A , v be the eigenvector of A and $\|v\|_2 = 1$.

$$\implies Av = \lambda v$$

$$\implies v^H Av = \lambda v^H v = \lambda$$

$$\stackrel{\text{take } H}{\implies} v^H A^H v = \lambda^*$$

Since A is Hermitian, $A = A^H$ and hence $\lambda = \lambda^*$

So λ is real-valued \square

Property 2: If $A \in \mathbb{C}^{m \times n}$ is Hermitian and that

- $\lambda_i \neq \lambda_j$ be two eigenvalues,
- v_i, v_j be the corresponding eigenvectors.

Then the eigenvectors are mutually orthogonal.

Proof

$$\begin{cases} v_i^H A v_j = v_i^H (\lambda_j v_j) = \lambda_j v_i^H v_j \\ v_j^H A v_i = v_j^H (\lambda_i v_i) = \lambda_i v_j^H v_i \stackrel{\text{take } H}{\implies} v_i^H A^H v_j = \lambda_i v_i^H v_j \end{cases}$$

Since $v_i^H A^H v_j = v_i^H A v_j$, so we have

$$(\lambda_i - \lambda_j) v_i^H v_j = 0 \implies v_i^H v_j = 0$$

\square

Unitarily diagonalizable

If Hermitian/symmetric A is diagonalizable

$$A = VDV^{-1}$$

where

- V contains orthogonal eigenvectors and is unitary matrix,
- D real-valued diagonal eigenvalues.

Therefore, $VV^H = I$ and $V^{-1} = V^H$. Thus,

$$A = VDV^H$$

And we say A is unitarily diagonalizable.

Schur triangularization

Any $A \in \mathbb{C}^{m \times m}$ can be decomposed as an upper triangular matrix.

Theorem 3 (Schur triangularization)

Given $A \in \mathbb{C}^{m \times m}$ with eigenvalues $\lambda_1, \lambda_2, \dots, \lambda_m \in \mathbb{C}$ there exists a unitary matrix $U \in \mathbb{C}^{m \times m}$ such that

$$U^H AU = T \iff A = UTU^H$$

where T is an upper triangular matrix with $[T]_{ii} = \lambda_i, i = 1, 2, \dots, m$

If $A \in \mathbb{C}^{m \times m}$, by Schur triangularization, there exists unitary $U \in \mathbb{C}^{m \times m}$ s.t.

$$U^H AU = T$$

Then by taking H , we have

$$U^H A^H U = T^H$$

\implies

$$T = T^H$$

Since T is an upper triangular matrix, so T^H is a lower triangular matrix. Therefore, T is diagonal. \implies

$$U^H AU = T$$

Thus, A is diagonalizable.

Proof for Schur Triangularization (**Theorem 3**)

This proof is also an algorithm.

Let $v_1 \in \mathbb{R}^m$ be the eigenvector of A associated with λ_1 , i.e., $Av_1 = \lambda_1 v_1$. Without losing generality, we can assume $v_1^H v_1 = 1$ and $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_m$

Construct a unitary matrix U_1

$$U_1 = [v_1 \quad V_1] \in \mathbb{C}^{m \times m}$$

where $V_1 \in \mathbb{C}^{m \times (m-1)}$. U_1 being unitary indicates

$$v_1^H V_1 = \vec{0}^T$$

Then we have

$$U_1^H AU_1 = \begin{bmatrix} v_1^H \\ V_1^H \end{bmatrix} A [v_1 \quad V_1] = \begin{bmatrix} v_1^H A v_1 & v_1^H A V_1 \\ V_1^H A v_1 & V_1^H A V_1 \end{bmatrix}$$

where

$$\begin{cases} v_1^H A v_1 = \lambda_1 \\ v_1^H A V_1 \neq \vec{0} \\ V_1^H A v_1 = \vec{0} \end{cases}$$

So,

$$U_1^H AU_1 = \begin{bmatrix} \lambda_1 & \times \\ 0 & V_1^H A V_1 \end{bmatrix}$$

Define

$$A_1 \triangleq V_1^H A V_1 \in \mathbb{C}^{(m-1) \times (m-1)}$$

Then

$$\begin{aligned} P_A(\lambda) &= \det(A - \lambda I) \\ &= \det(U_1^H AU_1 - \lambda I)^\dagger \\ &= \det\left(\begin{bmatrix} \lambda_1 & \times \\ 0 & A_1 \end{bmatrix} - \lambda I\right) \\ &= \det\left(\begin{bmatrix} \lambda_1 - \lambda & \times \\ 0 & A_1 - \lambda I \end{bmatrix}\right) \\ &= (\lambda_1 - \lambda) \underbrace{\det(A_1 - \lambda I)}_{=P_{A_1}(\lambda)} \end{aligned}$$

\dagger : $A = U_1(U_1^H AU_1)U_1^{-1}$ so A is similar to $U_1^H AU_1$ and hence A and $U_1^H AU_1$ have same eigenvalues.

Therefore, A_1 has eigenvalues $\lambda_2, \lambda_3, \dots, \lambda_m$

Let $v_2 \in \mathbb{C}^{m-1}$ be the eigenvector of A_1 associated with λ_2 .

Construct a unitary matrix U_2

$$U_2 = [v_2 \quad V_2] \in \mathbb{C}^{(m-1) \times (m-1)}$$

where $V_2 \in \mathbb{C}^{(m-1) \times (m-2)}$. U_2 being unitary indicates

$$v_2^H V_2 = \vec{0}^T$$

\implies

$$U_2^H A_1 U_2 = \begin{bmatrix} \lambda_2 & \times \\ 0 & V_2^H A_1 V_2 \end{bmatrix}$$

Let $\tilde{U}_2 = \begin{bmatrix} 1 & 0^T \\ 0 & U_2 \end{bmatrix}$, then \tilde{U}_2 is unitary.

So,

$$\begin{aligned} \tilde{U}_2^H [U_1^H AU_1] \tilde{U}_2 &= \tilde{U}_2^H \begin{bmatrix} \lambda_1 & \times \\ 0 & A_1 \end{bmatrix} \tilde{U}_2 \\ &= \begin{bmatrix} \lambda_1 & \times \\ 0 & U_2^H A_1 U_2 \end{bmatrix} \\ &= \begin{bmatrix} \lambda_1 & \times \\ 0 & \begin{bmatrix} \lambda_2 & \times \\ 0 & V_2^H A_1 V_2 \end{bmatrix} \end{bmatrix} \\ &= \begin{bmatrix} \lambda_1 & \times & \times \\ 0 & \lambda_2 & \times \\ 0 & 0 & V_2^H A_1 V_2 \end{bmatrix} \end{aligned}$$

Repeat this operations and we can finally get

$$\begin{aligned} \tilde{U}_m^H \dots \tilde{U}_3^H \tilde{U}_2^H (U_1 A U_1) \tilde{U}_2 \tilde{U}_3 \dots \tilde{U}_m \\ \underbrace{\hspace{10em}}_{\triangleq U \text{ unitary}} \\ = \begin{bmatrix} \lambda_1 & \times & \times & \times \\ & \lambda_2 & \times & \times \\ & & \ddots & \times \\ & & & \lambda_m \end{bmatrix} \triangleq T \end{aligned}$$

Therefore,

$$U^H A U = T$$

□

Variational Characterization of Eigenvalues

Rayleigh quotient

Definition (Rayleigh quotient)

$$R(A, x) = \frac{x^H A x}{\|x\|_2^2} = \left(\frac{x}{\|x\|_2} \right)^H A \left(\frac{x}{\|x\|_2} \right)$$

Theorem 4 (Rayleigh-Ritz theorem)

Let $A \in \mathbb{C}^{m \times m}$ be Hermitian, then

$$\begin{aligned} \lambda_{\max}(A) &= \max_{x \in \mathbb{C}^m} \frac{x^H A x}{\|x\|_2^2} \\ &= \max_{x \in \mathbb{C}^m} x^H A x \quad (\|x\|_2 = 1) \\ \lambda_{\min}(A) &= \min_{x \in \mathbb{C}^m} \frac{x^H A x}{\|x\|_2^2} \end{aligned}$$

and

$$\lambda_{\min}(A) \cdot \|x\|_2^2 \leq x^H A x \leq \lambda_{\max}(A) \cdot \|x\|_2^2$$

Proof

According to the EVD of A ,

$$A = V \Lambda V^H = \sum_{i=1}^m (\lambda_i \cdot v_i v_i^H)$$

\Rightarrow

$$\begin{aligned} x^H A x &= x^H V \Lambda V^H x \\ &= \sum_{i=1}^m (\lambda_i \cdot x^H v_i v_i^H x) \\ &= \sum_{i=1}^m (\lambda_i \cdot |v_i^H x|^2) \\ &\leq \lambda_{\max}(A) \cdot \sum_{i=1}^m |v_i^H x|^2 \end{aligned}$$

Since

$$\sum_{i=1}^m |v_i^H x|^2 = x^H V V^H x = \|x\|_2^2$$

So,

$$x^H A x \leq \lambda_{\max}(A) \cdot \|x\|_2^2$$

\Rightarrow

$$\frac{x^H A x}{\|x\|_2^2} \leq \lambda_{\max}(A)$$

The equality holds only when $x = v_1$, which is an eigenvector associated with $\lambda_{\max}(A)$ and is called the principal eigenvector. □

Courant-Fischer theorem

Theorem 5 (Courant-Fischer theorem)

Let $A \in \mathbb{C}^{m \times m}$ be Hermitian and $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_m$ are ordered eigenvalues. Then

$$\begin{aligned} \lambda_k &= \min_{\omega_1, \omega_2, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \omega_2, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \right\} \\ &= \max_{\omega_{k+1}, \omega_{k+2}, \dots, \omega_m} \left\{ \min_{x \in \mathbb{C}^m, x \perp \omega_{k+1}, \omega_{k+2}, \dots, \omega_m} \frac{x^H A x}{\|x\|_2^2} \right\} \end{aligned}$$

Proof ideas:

1.

$$\max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{x^H A x}{\|x\|_2^2} = \lambda_k$$

where v_k is the eigenvector associated with λ_k of A .

2.

$$\max_{x \in \mathbb{C}^m, x \perp \omega_1, \omega_2, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \geq \lambda_k$$

1+2 \Rightarrow

$$\min_{\omega_1, \omega_2, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \omega_2, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \right\} = \lambda_k$$

attained when $\omega_i = v_i$ ($i = 1, \dots, k-1$)

Proof

1.

Since $A = V \Lambda V^H = \sum_{i=1}^m \lambda_i v_i v_i^H$,

$$\begin{aligned} &\max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{x^H A x}{\|x\|_2^2} \\ &= \max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{\sum_{i=1}^m \lambda_i |v_i^H x|^2}{\|x\|_2^2} \\ &= \max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{\sum_{i=k}^m \lambda_i |v_i^H x|^2}{\|x\|_2^2} \\ &\leq \max_{x \in \mathbb{C}^m} \frac{\sum_{i=k}^m \lambda_i |v_i^H x|^2}{\|x\|_2^2} \quad (\text{Relax constraint}) \\ &\leq \max_{x \in \mathbb{C}^m} \frac{\lambda_k \sum_{i=k}^m |v_i^H x|^2}{\|x\|_2^2} \quad \dagger \\ &\leq \lambda_k \quad (\text{equality attained if } x = v_k) \end{aligned}$$

\dagger : $\sum_{i=k}^m |v_i^H x|^2 = x^H V V^H x = \|x\|_2^2$.

Therefore,

$$\max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{x^H A x}{\|x\|_2^2} = \lambda_k$$

2.

Let $y = V^H x$

$$\begin{aligned} &\max_{x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \\ &= \max_{x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H V \Lambda V^H x}{\|x\|_2^2} \\ &= \max_{y \perp V^H \omega_1, \dots, V^H \omega_{k-1}} \frac{y^H \Lambda y}{\|y\|_2^2} \\ &\geq \max_{y \perp V^H \omega_1, \dots, V^H \omega_{k-1}, y_{k+1}=y_{k+2}=\dots=y_m=0, \|y\|_2=1} y^H \Lambda y \\ &= \max_{y \perp V^H \omega_1, \dots, V^H \omega_{k-1}, |y_1|^2 + |y_2|^2 + \dots + |y_k|^2 = 1} \sum_{i=1}^k \lambda_i |y_i|^2 \\ &\geq \lambda_k \end{aligned}$$

□

Weyl's inequality

For $A, B \in \mathbb{C}^{m \times m}$ be two Hermitian matrices.
Then

1. $\lambda_k(A) + \lambda_m(B) \leq \lambda_k(A+B) \leq \lambda_k(A) + \lambda_1(B)$
2. $\lambda_{k+1}(A) \leq \lambda_k(X \pm zz^H) \leq \lambda_{k-1}(A)$
3. $\lambda_{k+r}(A) \leq \lambda_k(A+B) \leq \lambda_{k-r}(A)$ where $\text{rank}(B) \leq r$
4. $\lambda_{j+k-1}(A+B) \leq \lambda_j(A) + \lambda_k(B) \leq \lambda_{j+k-m}(A+B)$
5. Let A_I be the principal submatrix of A associated with subset $I = \{i_1, i_2, \dots, i_r\} \subseteq \{1, 2, \dots, m\}$
 $\lambda_{k+m-r}(A) \leq \lambda_k(A_I) \leq \lambda_k(A)$
6. Let $U \in \mathbb{C}^{m \times r}$ be semiunitary
 $\lambda_{k+m-r} \leq \lambda_k(U^H A U) \leq \lambda_k(A)$

Proof 1.

$$\lambda_k(A+B) = \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H(A+B)x}{\|x\|_2^2} \right\}$$

where

$$\frac{x^H(A+B)x}{\|x\|_2^2} = \frac{x^H A x}{\|x\|_2^2} + \underbrace{\frac{x^H B x}{\|x\|_2^2}}_{\leq \lambda_1(B)}$$

So,

$$\begin{aligned} & \lambda_k(A+B) \\ & \leq \underbrace{\min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \right\}}_{=\lambda_k(A)} + \lambda_1(B) \end{aligned}$$

Therefore,

$$\lambda_k(A+B) \leq \lambda_k(A) + \lambda_1(B)$$

Proof 2.

$$\begin{aligned} & \lambda_k(A \pm zz^H) \\ & = \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H(A \pm zz^H)x}{\|x\|_2^2} \right\} \\ & \geq \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}, \boxed{z}} \frac{x^H(A \pm zz^H)x}{\|x\|_2^2} \right\} \\ & = \min_{\omega_1, \dots, \omega_{k-1}, \omega_k = z} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}, \omega_k} \frac{x^H A x}{\|x\|_2^2} \right\} \\ & \geq \min_{\omega_1, \dots, \omega_{k-1}, \omega_k} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}, \omega_k} \frac{x^H A x}{\|x\|_2^2} \right\} \\ & = \lambda_{k+1}(A) \end{aligned}$$

Proof 3.

$$\begin{aligned} & \because B = \sum_{i=1}^r \lambda_i(B) u_i u_i^H \\ & \therefore A+B = A + \sum_{i=1}^{r-1} \lambda_i(B) u_i u_i^H + \lambda_r(B) u_r u_r^H \end{aligned}$$

Define $B^{r-k} \triangleq \sum_{i=1}^{r-k} \lambda_i(B) u_i u_i^H$ $k = 0, 1, \dots, r$
 $\lambda_{k+1}(A+B^{r-1}) \leq \lambda_k(A+B^{r-1} + \lambda_r(B) u_r u_r^H) \leq \lambda_{k-1}(A+B^{r-1})$
Do again and we get

$$\begin{aligned} \lambda_{k+2}(A+B^{r-2}) & \leq \lambda_{k+1}(A+B^{r-2} + \lambda_{r-1}(B) u_{r-1} u_{r-1}^H) \\ & = \lambda_{k+1}(A+B^{r-1}) \end{aligned}$$

Repeat r times, we can get

$$\lambda_{k+r}(A) \leq \lambda_k(A+B) \leq \lambda_{k-r}(A)$$

where $\text{rank}(B) \leq r$ □

Proof 4.

Define $A_{j-1} = \sum_{i=1}^{j-1} \lambda_i(A) v_i v_i^H$ (if $A = V \Lambda V^H$).
Define $B_{k-1} = \sum_{i=1}^{k-1} \lambda_i(B) u_i u_i^H$ (if $B = U \Lambda U^H$).
Then

$$\begin{aligned} \lambda_j(A) & = \lambda_1(A - A_{j-1}) \\ \lambda_k(B) & = \lambda_1(B - B_{k-1}) \end{aligned}$$

\Rightarrow

$$\begin{aligned} \lambda_j(A) + \lambda_k(B) & = \lambda_1(A - A_{j-1}) + \lambda_1(B - B_{k-1}) \\ & \geq \lambda_1(A+B - (A_{j-1} + B_{k-1})) \\ & \geq \lambda_{j+k-1}(A+B)^\dagger \end{aligned}$$

\dagger : $\text{rank}((A_{j-1} + B_{k-1})) \leq j+k-2$ and from 3. □

Proof 5.

$$\begin{aligned} & \lambda_k(A) \\ & = \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \perp \omega_1, \dots, \omega_{k-1}, \|x\|_2=1} x^H A x \right\} \\ & \geq \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \perp \omega_1, \dots, \omega_{k-1}, \|x\|_2=1, x \in \text{span}\{e_{i_1}, \dots, e_{i_r}\}} x^H A x \right\}^\dagger \\ & = \min_{v_1, \dots, v_{k-1}} \left\{ \max_{y \perp v_1, \dots, v_{k-1}, \|y\|_2=1} y^H A y \right\}^\ddagger \\ & = \lambda_k(A_I) \end{aligned}$$

\dagger : $e_i = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}$, only the i th entry equal to 1. □

\ddagger : Define $y \in \mathbb{C}^r$, $y_\ell = x_{i_\ell}$ ($\ell = 1, \dots, r$) and take $v_i \in \mathbb{C}^r$, $(v_i)_\ell = (\omega_i)_{i_\ell}$ □

Proof 6.

Let $\tilde{U} = \begin{bmatrix} U & \tilde{U} \end{bmatrix}$ be unitary. Then

$$\tilde{U}^H A \tilde{U} = \begin{bmatrix} U^H A U & \times \\ \times & \times \end{bmatrix}$$

where $U^H A U$ is the principal submatrix of $\tilde{U}^H A \tilde{U}$ with $I = 1, 2, \dots, r$

Therefore,

$$\lambda_k(A) = \lambda_k(\tilde{U}^H A \tilde{U}) \geq \lambda_k(U^H A U)$$

□

□

The Power Method

The power method is a simple method to obtain eigenvector/eigenvalue of a matrix, particularly for large scale problem.

Algorithm 1 (Power Method)

Given a matrix $A \in \mathbb{C}^{m \times m}$ and $v^0 \in \mathbb{C}^m$

$$v^{(k+1)} \leftarrow \frac{A \cdot v^{(k)}}{\|A \cdot v^{(k)}\|_2}$$

until a stopping condition is met.

$\Rightarrow v^{(k+1)}$ takes as an estimate of the **principle eigenvector** of A .

Assumption:

- A is diagonalizable (in particular, A is Hermitian (symmetric)). $A = V\Lambda V^H$.
- $|\lambda_1| \geq |\lambda_2| \geq \dots \geq |\lambda_m|$ and $|\lambda_1| > |\lambda_2|$

$$v^{(k+1)} = \frac{A \cdot v^{(k)}}{\|A \cdot v^{(k)}\|_2} = \frac{A^{k+1} \cdot v^{(0)}}{\|A^{k+1} \cdot v^{(0)}\|_2}$$

Consider $A^k x$, x satisfies $[V^H x]_1 \neq 0$. (x is $v^{(0)}$)

Let $\alpha = V^H x$

$$A^k \cdot x = (V\Lambda^k V^H) \cdot x = V\Lambda^k \alpha$$

Let $V = [v_1, v_2, \dots, v_m]$

$$\begin{aligned} A^k \cdot x &= V\Lambda^k \alpha \\ &= \sum_{i=1}^m v_i \cdot \lambda_i^k \cdot \alpha_i \\ &= v_1(\lambda_1 \alpha_1) + \sum_{i=2}^m \alpha_i \lambda_i^k v_i \\ &= \lambda_1^k \alpha_1 \underbrace{\left(v_1 + \sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \left(\frac{\lambda_i}{\lambda_1} \right)^k v_i \right)}_{\triangleq r_k(\text{residue})} \end{aligned}$$

\Rightarrow

$$\begin{aligned} \|r_k\|_2 &= \left\| \sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \left(\frac{\lambda_i}{\lambda_1} \right)^k v_i \right\|_2 \\ &\leq \sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \cdot \underbrace{\left(\frac{\lambda_i}{\lambda_1} \right)^k}_{=1} \cdot \|v_i\|_2 \\ &\leq \underbrace{\left| \frac{\lambda_2}{\lambda_1} \right|^k}_{<1} \cdot \sum_{i=2}^m \left| \frac{\alpha_i}{\alpha_1} \right| \rightarrow 0 \text{ as } k \rightarrow \infty \end{aligned}$$

Consider $c_k = \frac{|\alpha_1| |\lambda_1^k|}{\alpha_1 \lambda_1^k}$ and $\lim_{k \rightarrow \infty} c_k \cdot \frac{A^k \cdot x}{\|A^k \cdot x\|_2}$

$$\begin{aligned} c_k \cdot \frac{A^k \cdot x}{\|A^k \cdot x\|_2} &= \frac{|\alpha_1| |\lambda_1^k|}{\alpha_1 \lambda_1^k} \cdot \frac{1}{\|A^k \cdot x\|_2^\dagger} \cdot \lambda_1^k \alpha_1 \left(v_1 + \sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \left(\frac{\lambda_i}{\lambda_1} \right)^k v_i \right) \\ &= \frac{|\alpha_1| |\lambda_1^k|}{\sqrt{\sum_{i=1}^m (\lambda_i^k)^2 \alpha_i^2}}^\ddagger \cdot \left(v_1 + \underbrace{\sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \left(\frac{\lambda_i}{\lambda_1} \right)^k v_i}_{\rightarrow 0 \text{ as } k \rightarrow \infty} \right) \end{aligned}$$

$$\dagger: \|A^k x\|_2 = \|V\Lambda^k V^H x\|_2 = \|V\Lambda^k \alpha\|_2 = \|\Lambda^k \alpha\|_2 = \sqrt{\sum_{i=1}^m (\lambda_i^k \alpha_i)^2}$$

$$\begin{aligned} \ddagger: \frac{|\alpha_1| |\lambda_1^k|}{\sqrt{\sum_{i=1}^m (\lambda_i^k)^2 \alpha_i^2}} &= \frac{|\alpha_1| |\lambda_1^k|}{\sqrt{(\lambda_1^k)^2 \alpha_1^2 + \sum_{i=2}^m (\lambda_i^k)^2 \alpha_i^2}} \\ &= \frac{1}{\sqrt{1 + \sum_{i=2}^m \left(\frac{\lambda_i}{\lambda_1} \right)^{2k} \left(\frac{\alpha_i}{\alpha_1} \right)^2}} \end{aligned}$$

Since $\sum_{i=2}^m \left(\frac{\lambda_i}{\lambda_1} \right)^{2k} \left(\frac{\alpha_i}{\alpha_1} \right)^2 \rightarrow 0$ as $k \rightarrow \infty$, so $\frac{|\alpha_1| |\lambda_1^k|}{\sqrt{\sum_{i=1}^m (\lambda_i^k)^2 \alpha_i^2}} \rightarrow 1$ as $k \rightarrow \infty$.

Therefore, $\lim_{k \rightarrow \infty} c_k \cdot \frac{A^k \cdot x}{\|A^k \cdot x\|_2} \rightarrow v_1$.

Deflation:

$$\begin{aligned} A &= \sum_{i=1}^m \lambda_i v_i v_i^H \\ &= \lambda_1 v_1 v_1^H + \sum_{i=2}^m \lambda_i v_i v_i^H \end{aligned}$$

Obtain $A^{(1)} = A - \lambda_1 v_1 v_1^H$ and apply power method to obtain v_2

Obtain $A^{(2)} = A - \lambda_2 v_2 v_2^H$ and apply power method to obtain v_3

...

Lecture 5. PageRank

PageRank Model

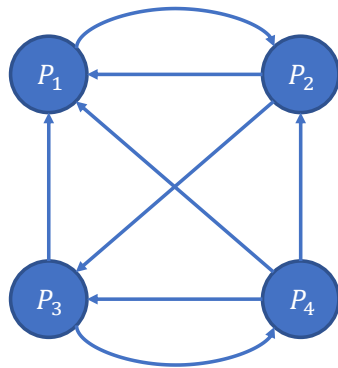
Importance score

- v_i : importance score of page i
- c_j : number of outgoing links of page j
- \mathcal{L}_i : set of pages that refers to page i

$$\begin{aligned} v_i &= |\mathcal{L}_i| \\ &\Downarrow \\ v_i &= \sum_{j \in \mathcal{L}_i} v_j \\ &\Downarrow \\ v_i &= \sum_{j \in \mathcal{L}_i} \frac{v_j}{c_j} \end{aligned}$$

The model

e.g.



$$\begin{aligned} v_1 &= 1 \cdot v_2 + \frac{1}{2}v_3 + \frac{1}{3}v_4 \\ v_2 &= 1 \cdot v_1 + \frac{1}{3}v_4 \end{aligned}$$

...

\Rightarrow

$$\underbrace{\begin{bmatrix} 0 & 1/2 & 1/2 & 1/3 \\ 1 & 0 & 0 & 1/3 \\ 0 & 1/2 & 0 & 1/3 \\ 0 & 0 & 1/2 & 0 \end{bmatrix}}_{\text{link matrix}} \begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{bmatrix} = \begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{bmatrix}$$

\Rightarrow

$$Av = v, \quad v \geq 0$$

PageRank Problem

Find solution $v \geq 0$ satisfying $Av = v$

The four questions:

1. Does A admit a solution of $Av = v$?
(Does A has eigenvalue equal to 1?)
2. Does $Av = v$ admit a non-negative solution?
($v \geq 0$ or $v \leq 0$)

3. Is the solution unique?
(unique ranking)
4. Can we obtain the solution efficiently?
(Power method applicable?)

Answer to Question 1

Does A admit a solution of $Av = v$? (Does A has eigenvalue equal to 1?)

Note that $A \geq 0$ and column-sum equal to 1 (column-stochastic matrix).

$$\Rightarrow \mathbf{1}^T A = \mathbf{1}^T$$

$$\Rightarrow A^T \cdot \mathbf{1} = \mathbf{1}$$

$$\Rightarrow A^T \text{ has an eigenvalue equal to 1}$$

Since A and A^T have the same set of eigenvalues

$$\Rightarrow A \text{ has an eigenvalue equal to 1}$$

Theorem 6

Let A be non-negative and column(row)-stochastic, then

1. $\lambda = 1$ is an eigenvalue of A
2. $|\lambda| \leq 1$ for all eigenvalues of A

Proof Let $Av = \lambda v$ for some $|\lambda| \neq 1$.

Let $|v_k| = \max\{|v_1|, |v_2|, \dots, |v_m|\}$.

$$\lambda v_k = \sum_{i=1}^m [A]_{ki} v_i$$

\Rightarrow

$$\begin{aligned} |\lambda| |v_k| &= \left| \sum_{i=1}^m [A]_{ki} v_i \right| \\ &\leq \sum_{i=1}^m |[A]_{ki}| \cdot |v_i| \\ &\leq |v_k| \cdot \sum_{i=1}^m |[A]_{ki}| \\ &= |v_k| \end{aligned}$$

\Rightarrow

$$|\lambda| \leq 1$$

□

Disjoint page networks

$$A = \begin{bmatrix} A_1 & 0 \\ 0 & A_2 \end{bmatrix}$$

$$A_1 v_1 = v_1$$

$$A_2 v_2 = v_2$$

$$A \left(\alpha_1 \begin{pmatrix} v_1 \\ 0 \end{pmatrix} + \alpha_2 \begin{pmatrix} 0 \\ v_2 \end{pmatrix} \right) = \left(\alpha_1 \begin{pmatrix} v_1 \\ 0 \end{pmatrix} + \alpha_2 \begin{pmatrix} 0 \\ v_2 \end{pmatrix} \right)$$

Modified model

$$M \triangleq (1 - \beta)A + \beta \left(\frac{\mathbf{1}\mathbf{1}^T}{m} \right) > 0$$

This is positive column-stochastic

$$\mathbf{1}^T M = (1 - \beta)\mathbf{1}^T A + \beta \mathbf{1}^T \left(\frac{\mathbf{1}\mathbf{1}^T}{m} \right) = \mathbf{1}^T$$

Answer to Question 2

Does $Av = v$ admit a non-negative solution? ($v \geq 0$ or $v \leq 0$)

Theorem 7

For positive and column-stochastic matrix M , any eigenvector associated with λ_1 has either all positive or all negative components.

Proof Note that for $y \in \mathbb{R}^m$, $|\sum_{i=1}^m y_i| < \sum_{i=1}^m |y_i|$ if $\{y_i\}$ have mixed signs.

Suppose the above theorem is not true. Then we have

$$|v_i| = \left| \sum_{j=1}^m [M]_{ij} v_j \right| < \sum_{j=1}^m [M]_{ij} |v_j|$$

$$\begin{aligned} \Rightarrow \sum_{i=1}^m |v_i| &< \sum_{i=1}^m \sum_{j=1}^m [M]_{ij} v_j = \sum_{j=1}^m \sum_{i=1}^m [M]_{ij} v_j = \sum_{j=1}^m |v_j| \end{aligned}$$

Contradict! So the above theorem is true. \square

Proof

$$1) \quad v \in S \iff 1^T v = 0$$

2)

$$w \triangleq Mv$$

$$\begin{aligned} \|w\|_1 &= \|Mv\|_1 \\ &= \sum_{i=1}^m |w_i| = \sum_{i=1}^m \text{sign}(w_i) \cdot w_i \\ &= \sum_{i=1}^m \left[\text{sign}(w_i) \cdot \sum_{j=1}^m [M]_{ij} v_j \right] \\ &= \sum_{j=1}^m \left[\left(\sum_{i=1}^m \text{sign}(w_i) [M]_{ij} \right) \cdot v_j \right] \end{aligned}$$

$$a_j \triangleq \left(\sum_{i=1}^m \text{sign}(w_i) [M]_{ij} \right)$$

$$|a_j| = \left| \sum_{i=1}^m \text{sign}(w_i) [M]_{ij} \right|$$

\square

Answer to Question 3

Is the solution unique? (unique ranking)

Since $Mv = v \implies (M - I)v = 0$, show

$$\dim(\text{Null}(M - I)) = 1$$

i.e., there is a unique ranking.

Proof

Simple lemma: Let a and b be linearly independent, then $\exists x \in \text{span}\{a, b\}$ which has mixed sign.

Proof for this lemma:

If $1^T q = 0 \implies q$ has mixed sign.

If $1^T q \neq 0$, set $\alpha_1 = \frac{-1^T b}{1^T q}$, $\alpha_2 = 1$ and let $x = \alpha_1 \cdot a + \alpha_2 \cdot b$

$\implies x$ satisfies $1^T x = 0$, so x has mixed sign.

Suppose $\dim(\text{Null}(M - I)) > 1$, then for $q_1, q_2 \in \text{Null}(M - I)$ that are linear independent. $\implies \exists v = \alpha_1 q_1 + \alpha_2 q_2 \in \text{Null}(M - I)$

$\implies Mv = v$ which has mixed sign. (Contradict!)

$\implies \dim(\text{Null}(M - I)) = 1$ \square

Answer to Question 4

Can we obtain the solution efficiently? (Power method applicable?)

Property: Let M be positive and column-stochastic.

Let $S = \{v \in \mathbb{R}^m | 1^T v = 0\}$.

Then,

1. $Mv \in S$ if $v \in S$
2. $\|Mv\|_1 \leq c \|v\|_1$ with $c < 1 \forall$ any $v \in S$
3. the eigenvector associated with $\lambda = 1$, denoted v_1 , can be obtainable by the power method

Lecture 6. Positive Semidefinite Matrix

For $A \in \mathbb{C}^{m \times m}$, if

- A is symmetric, we denote $A \in S^m$
- A is Hermitian, we denote $A \in H^m$

Definition

We say $A \in H^m$ is positive semidefinite (PSD) if

$$x^H Ax > 0 \quad \forall x \in \mathbb{C}^m$$

We say $A \in H^m$ is positive definite (PD) if

$$x^H Ax > 0 \quad \forall x \in \mathbb{C}^m, x \neq 0$$

We say $A \in H^m$ is indefinite if A is not PSD.

- $\det(A) \geq 0 (> 0)$

If A is PD, then A is invertible.

$$A = V \Lambda V^H$$

$$A^{-1} = V\Lambda^{-1}V^H$$

Property 2: $A \in H^m$ can be factorized as $A = BB^H$ for some matrix B if and only if A is PSD.

Proof

 \Rightarrow

$$x^H A x = x^H B B^H x = \|B^H x\|^2 \geq 0 \implies A \text{ is PSD}$$

 \Leftarrow

If A is PSD, define $\Lambda^{1/2} = \text{diag}(\sqrt{\lambda_1}, \sqrt{\lambda_2}, \dots, \sqrt{\lambda_m})$

$$\begin{aligned} A &= V\lambda V^H \\ &= V\Lambda^{1/2}\Lambda^{1/2}V^H \\ &= V\Lambda^{1/2}Q^HQ\Lambda^{1/2}V^H \end{aligned}$$

where Q is a Hermitian matrix.

Define $B \triangleq V\Lambda^{1/2}Q^H$. Then, $A = BB^H$. \square

Theorem 2

Consider $A \in H^m$, $B \in \mathbb{C}^{m \times n}$ and $C = B^H A B$

- If A is PSD, then C is PSD.
- If A is PD, then C is PD if and only if $rank(B) = n$
- If B is square and invertible, then A is PSD(PD) if and only if C is PSD(PD)

Proof of 2)

(\implies) Show C is PSD $\implies \text{rank}(B) = n$

Assume B is not full column rank.

Consider $x^H C x = x^H B^H A B x$, then there exists $x \neq 0$ s.t. $Bx = 0 \implies x^H C x = 0$ (Contradict!) (\Leftarrow) Show $\text{rank}(B) = n \implies$ Show C is PSD
 $\text{rank}(B) = n \implies Bx = 0$ iff $x = 0$. So $\forall x \neq 0 \implies Bx \neq 0$.

$\implies x^H C x = x^H B^H A B x > 0$ as A is PD. So C is PD. \square

Eigenvalues of PSD matrices

Theorem 1

$A \in H^m$ is PSD(PD) if and only if $\lambda_1, \lambda_2, \dots, \lambda_3 \geq 0(> 0)$

Proof

Consider EVD of A as

$$A = V\Lambda V^H, \quad \Lambda = diag(\lambda_1, \lambda_2, \dots, \lambda_m)$$

Consider $x^H A x = x^H V \Lambda V^H x = z^H \Lambda z$.

$$A \text{ is PSD} \iff x^H A x \geq 0 \quad \forall x \in \mathbb{C}^m \iff z^H \Lambda z \geq 0 \quad \forall z \in \mathbb{C}^m.$$

$$z^H \Lambda z = \sum_{i=1}^m \lambda_i |z_i|^2 \geq 0 \quad \forall z \in \mathbb{C}^m \iff \lambda_i \geq 0 \quad \forall i = 1, \dots, m \quad \square$$

If A is PSD(PD),

- $Tr(A) \geq 0 (> 0)$

Property 3: Let $A \in \mathbb{C}^{m \times k}$ and $B \in \mathbb{C}^k$ and B has full row rank ($rank(B) = k$). Then

$$Range(A) = Rnage(AB)$$

Proof

$Range(B) = \mathbb{C}^k$ as B is full row rank.

$$\begin{aligned} \text{Range}(AB) &= \{y|y = ABx, x \in \mathbb{C}^n\} \\ &= \{y|y = Az, \underbrace{z = Bx, x \in \mathbb{C}^n}_{z \in \text{Range}(B) = \mathbb{C}^k}\} \\ &= \{y|y = Az, z \in \mathbb{C}^k\} \\ &= \text{Range}(A) \end{aligned}$$

Property 4: Let $B \in \mathbb{C}^{m \times k}$ and $C \in \mathbb{C}^{m \times k}$ be two matrices with full column rank. Then

$$BB^H = CC^H$$

if and only if $C = BQ$ for some unitary Q . *Proof*

(\Leftarrow): obvious.

(\Rightarrow): $(C^H)^+ \triangleq C(C^H C)^{-1}$. Since $\text{rank}(C^H C) = \text{rank}(C)$, $C^H C$ is invertible † .

† : As C has full column rank, then $Cx = 0 \iff x = 0$.
 $Cx = 0 \implies C^H Cx = 0 \implies x^H C^H Cx = 0 \iff \|Cx\|^2 = 0 \implies x = 0$. Therefore, $C^H Cx = 0 \implies x = 0$ and hence $\text{rank}(C^H C) = \text{rank}(C)$.

$$\begin{aligned} BB^H(C^H)^+ &= CC^H(C^H)^+ \\ &= CC^H C(C^H C)^{-1} \\ &= C \end{aligned}$$

\implies

$$\begin{aligned} C &= BB^H(C^H)^+ \\ &= \underbrace{BB^H C(C^H C)^{-1}}_{\triangleq Q(\text{square})} \\ Q^H Q &= \underbrace{(C^H C)^{-H} C^H}_{=I} \underbrace{BB^H C(C^H C)^{-1}}_{=I} \\ &= I \end{aligned}$$

Therefore, Q is unitary. \square

Matrix Inequality

For $A, B \in H^{m \times m}$, we write $A \succeq (\succ) B$ if $A - B$ is a PSD(PD) matrix.

So $A \succeq 0$ if A is PSD.

Basic Properties:

- If $A \succeq 0$, then $\alpha A \succeq 0 \forall \alpha \geq 0$
- If $A \succeq 0, B \succeq 0$, then $\alpha A + B \succeq 0$
- If $A \succeq B, B \succeq C$, then $\alpha A \succeq C$
- If $A \not\succeq B$, it doesn't imply $B \succ A$

Let $\lambda_1(A) \geq \lambda_2(A) \geq \dots \geq \lambda_m(A)$ for $A \in H^{m \times m}$, then

1. $A \succeq I \iff \lambda_i(A) \geq 1 \forall i = 1, \dots, m$
2. $I \succeq A \iff \lambda_i(A) \leq 1 \forall i = 1, \dots, m$
3. Suppose $A, B \succ 0$, then $A \succeq B \iff B^{-1} \succeq A^{-1}$
4. If $A \succeq B$, then $\lambda_i(A) \geq \lambda_i(B) \forall i = 1, \dots, m$

Proof

1)

$$\begin{aligned} A \succeq I &\iff A - I \succ 0 \iff V \Lambda V^H - I \succeq 0 \iff \\ V(\Lambda - I)V^H &\succeq 0 \iff \Lambda - I \succeq 0 \iff \lambda_i(A) \geq 1 \\ \forall i = 1, \dots, m \end{aligned}$$

2) similar to 1)

3)

$$\begin{aligned} A \succeq B &\iff A - B \succeq 0 \\ &\iff A^{\frac{1}{2}} A^{\frac{1}{2}} - B \succeq 0 \\ &\iff A^{\frac{1}{2}} (I - A^{-\frac{1}{2}} B A^{-\frac{1}{2}}) A^{\frac{1}{2}} \succeq 0 \\ &\iff I - A^{-\frac{1}{2}} B A^{-\frac{1}{2}} \succeq 0 \\ &\iff \lambda_i(A^{-\frac{1}{2}} B A^{-\frac{1}{2}}) \leq 1 \quad \forall i = 1, \dots, m \\ &\iff \frac{1}{\lambda_i(A^{\frac{1}{2}} B^{-1} A^{\frac{1}{2}})} \leq 1 \quad \forall i = 1, \dots, m^\dagger \\ &\iff \lambda_i(A^{\frac{1}{2}} B^{-1} A^{\frac{1}{2}}) \geq 1 \quad \forall i = 1, \dots, m \\ &\iff A^{\frac{1}{2}} B^{-1} A^{\frac{1}{2}} - I \succeq 0 \\ &\iff A^{-\frac{1}{2}} (I - A^{\frac{1}{2}} B^{-1} A^{\frac{1}{2}}) A^{-\frac{1}{2}} \succeq 0 \\ &\iff B^{-1} - A^{-1} \succeq 0 \\ &\iff B^{-1} \succeq A^{-1} \end{aligned}$$

$$^\dagger: \lambda_i(A^{-1}) = \frac{1}{\lambda_i(A)}$$

4)

$$\begin{aligned} \lambda_i(A) &= \lambda_i(A - B + B) \\ &\geq \lambda_i(B) + \underbrace{\lambda_m(A - B)}_{\geq 0 \text{ since } A - B \succeq 0} \\ &\geq \lambda_i(B) \quad \forall i = 1, \dots, m \end{aligned}$$

\square

From 4), if $A \succeq B \implies \text{Tr}(A) \geq \text{Tr}(B)$, if additionally $A, B \succeq 0 \implies \det(A) \geq \det(B)$, and if $A, B \succ 0$ and $A \succeq B$ then $\text{Tr}(B^{-1}) \geq \text{Tr}(A^{-1})$

Definition (Schur Complement)

Consider a matrix $X \in H^{(m+n) \times (m+n)}$ to be partitioned

$$X = \begin{bmatrix} A & B \\ B^H C & \end{bmatrix}$$

where $A \in H^{m \times m}, B \in \mathbb{C}^{m \times n}, C \in H^{n \times n} \succ 0$. The term

$$S \triangleq A - B C^{-1} B^H \in H^{m \times m}$$

is called *Schur Complement* of X .

Then $X \succeq (\succ) 0 \iff S \succeq (\succ) 0$ *Proof*

$$\text{Let } Y = \begin{bmatrix} I_m & 0 \\ -C^{-1} B^H & I_n \end{bmatrix} \in \mathbb{C}^{(m+n) \times (m+n)}.$$

Then Y is a lower triangular matrix and $\det(Y) = 1 \implies Y$ is invertible.

Then

$$Y^H X Y = \begin{bmatrix} A - B C^{-1} B^H & 0 \\ 0 & C \end{bmatrix}$$

$$Y^H X Y \succeq (\succ) 0 \iff x \succeq (\succ) 0 \implies X \succeq (\succ) 0 \iff X \succeq (\succ) 0 \quad \square$$

$$\begin{aligned} \det(Y X^H Y) &= \det(X) \\ &= \det(S) \det(C) \end{aligned}$$

Lecture 7. Singular Value Decomposition