

# Notes for CIE 6002 Matrix Analysis

Binghao He

March 17, 2021

## Contents

<b>Lecture 2. Fundamental and Basic Operations</b>	<b>3</b>
Subspace . . . . .	3
Definition . . . . .	3
Some important subspaces . . . . .	3
Linear independence . . . . .	3
Basis . . . . .	3
Rank . . . . .	3
Nonsingularity/Invertibility . . . . .	3
Determinant . . . . .	3
Vector Norm . . . . .	3
Inner Product . . . . .	4
Projection onto a Subspace . . . . .	4
Projection . . . . .	4
Orthogonal complement . . . . .	4
Sum of subspaces . . . . .	4
Orthogonal Set . . . . .	4
Orthogonal/Unitary Matrix . . . . .	4
Dimensional Theorem . . . . .	5
Complexity of Matrix Computation . . . . .	5
Flops . . . . .	5
Big O notation . . . . .	5
<b>Lecture 3. Least Square Problem (LSP)</b>	<b>6</b>
Formulation . . . . .	6
LS problem . . . . .	6
Linear representation of signals . . . . .	6
AR and Polynomial Models . . . . .	6
Autoregressive (AR) model . . . . .	6
Polynomial model . . . . .	6
Basis Representation . . . . .	6
Basis matrix . . . . .	6
Fourier basis . . . . .	6
LTI System . . . . .	7
Lienar time invariant system . . . . .	7
System identification . . . . .	7
Deconvolution . . . . .	7
Toeplitz matrix . . . . .	7

Circulant Matrix . . . . .	7
<b>Lecture 4. Eigenvalues and Eigenvectors</b>	<b>8</b>
Diagonalizability . . . . .	8
Similar matrices . . . . .	8
Diagonalizable matrix . . . . .	8
When Diagonalizable . . . . .	8
Distinct eigenvalues . . . . .	8
Repeated eigenvalues . . . . .	9
Symmetric/Hermitian Matrices . . . . .	9
Unitarily diagonalizable . . . . .	10
Schur triangularization . . . . .	10
Variational Characterization of Eigenvalues . . . . .	11
Rayleigh quotient . . . . .	11
Courant-Fischer theorem . . . . .	11
Weyl's inequality . . . . .	12
The Power Method . . . . .	12
<b>Lecture 5. PageRank</b>	<b>14</b>
PageRank Model . . . . .	14
Importance score . . . . .	14
The model . . . . .	14
PageRank Problem . . . . .	14
Answer to Question 1 . . . . .	14
Answer to Question 2 . . . . .	15
Answer to Question 3 . . . . .	15
Answer to Question 4 . . . . .	15
<b>Lecture 6. Positive Semidefinite Matrix</b>	<b>16</b>
Properties of PSD Matrices . . . . .	16
Propeties . . . . .	16
Eigenvalues of PSD matrices . . . . .	16
Matrix Inequality . . . . .	17

# Lecture 2. Fundamental and Basic Operations

## Subspace

### Definition

The set  $S \subseteq \mathbb{R}^m$  is a subspace, if for  $x, y \in S$   
 $\alpha x + \beta y \in S, \quad \alpha, \beta \in \mathbb{R}^m$

This indicates

$$x_i \in S, \quad i = 1, \dots, N \implies \sum_{i=1}^N \alpha_i x_i \in S$$

### Some important subspaces

1. Span: a collection of vectors

$$\text{span}\{a_1, \dots, a_n\} = \{y \in \mathbb{R}^m | y = \sum_{i=1}^n \alpha_i a_i, \alpha_i \in \mathbb{R}\}$$

2. Orthogonal complement: given a subspace  $S \subseteq \mathbb{R}^m$

$$S_{\perp} = \{y \in \mathbb{R}^m | y^T x = 0, \forall x \in S\}$$

3. Range space of  $A$  (linear comb. of col. vectors)

$$\begin{aligned} R(A) &= \text{span}\{a_1, a_2, \dots, a_n\} \\ &= \{y \in \mathbb{R}^m | y = Ax, x \in \mathbb{R}^n\} \end{aligned}$$

4. Null space of  $A$  (linear comb. of col. vectors)

$$\text{Null}(A) = \{x \in \mathbb{R}^n | Ax = 0\}$$

### Linear independence

Given a set  $\{a_1, a_2, \dots, a_n\}$ , we say they are linear independent if

$$\sum_{i=1}^n \alpha_i a_i = 0 \implies \alpha_i = 0, \forall i = 1, \dots, n$$

### Basis

$\{b_1, b_2, \dots, b_k\} \subseteq \mathbb{R}^m$  is a basis for the subspace  $S \subseteq \mathbb{R}^m$  if  $S = \text{span}\{b_1, b_2, \dots, b_k\}$  and  $\{b_1, b_2, \dots, b_k\}$  is linear independent.

The number of the vectors of the basis for  $S$

$$k = \dim(S)$$

1.  $k$ : max number of linear indep. vectors in  $S$
2.  $k$ : min number of vectors that can span  $S$

### Rank

For a matrix  $A \in \mathbb{R}^{m \times n}$

$$\dim(R(A)) \leq \min\{m, n\}$$

Rank of a matrix (?)

$$\begin{aligned} \text{rank}(A) &= \dim(R(A)) \\ &= \text{max number of linear indep.} \\ &\quad \text{columns(rows) of } A \\ &= \text{rank}(A^T) \end{aligned}$$

And

$$\text{Null}(A) = (R(A^T))_{\perp}$$

### Nonsingularity/Invertibility

A symmetric matrix  $A \in \mathbb{R}^{m \times m}$  is non-singular if

$$Ax = 0 \implies x = 0$$

$\implies$  columns of  $A$  are linearly indep.

$\implies \exists A^{-1} \in \mathbb{R}^{m \times m}$  that  $AA^{-1} = A^{-1}A = I_m$

Some properties

1.  $(AB)^{-1} = B^{-1}A^{-1}$
2.  $(A^{-1})^T = (A^T)^{-1}$

### Determinant

For a matrix  $A \in \mathbb{R}^{m \times m}$ ,

Inductive

1. If  $m = 1$ ,  $\det(A) = A$
2. If  $m > 1$ , Define:  $A_{ij} \in \mathbb{R}^{(m-1) \times (m-1)}$  is a sub-matrix that removes  $i$ th row and  $j$ th column of  $A$ .

Define:  $c_{ij} = (-1)^{i+j} \det(A_{ij})$ .

$$\begin{aligned} \det(A) &= \sum_{j=1}^m a_{ij} c_{ij} \quad \forall i = 1, \dots, m \\ &= \sum_{i=1}^m a_{ij} c_{ij} \quad \forall j = 1, \dots, m \end{aligned}$$

For  $A, B \in \mathbb{R}^{m \times m}$

- $\det(AB) = \det(A)\det(B)$
- $\det(A) = \det(A^T)$
- $\det(\alpha A) = \alpha^m \det(A)$
- $\det(A) = 0 \iff A$  is singular
- $\det(A^{-1}) = 1/\det(A)$  if  $A$  is invertible

### Vector Norm

A mapping  $f : \mathbb{R}^> \rightarrow \mathbb{R}$  is a norm if it satisfies that  $\forall x, y \in \mathbb{R}^m$

- $f(x) \geq 0$
- $f(x) = 0$  iff.  $x = 0$
- $f(x + y) \leq f(x) + f(y)$
- $f(\alpha x) = |\alpha|f(x)$

Examples

- $\|x\|_2 = \sqrt{\sum_{i=1}^m |x_i|^2}$  (Euclidean norm)
- $\|x\|_1 = \sum_{i=1}^m |x_i|$
- $\|x\|_{\infty} = \max_{i=1, \dots, m} |x_i|$
- $\|x\|_p = \sqrt[p]{\sum_{i=1}^m |x_i|^p}$

## Inner Product

$$\langle x, y \rangle = x^T y$$

$$\langle x, x \rangle = \|x\|_2^2$$

Cauchy-Schwarz Inequality

$$|x^T y| \leq \|x\|_2 \cdot \|y\|_2$$

Holder Inequality

$$|x^T y| \leq \|x\|_p \cdot \|y\|_q$$

where

$$1/p + 1/q = 1 \quad , \quad p \cdot q \geq 1$$

## Projection onto a Subspace

### Projection

Given a set  $S \subseteq \mathbb{R}^m$  and a point  $y \in \mathbb{R}^m$ , the projection of  $y$  onto  $S$  is

$$y_S = \arg \min_{z \in S} \|z - y\|_2^2$$

#### Theorem 1

Suppose  $S$  is a subspace,

1.  $y_S$  exists and is unique
2. Let  $y_S = \Pi_S(y)$  iff.  
 $z^T(y_S - y) = 0 \quad \forall z \in S$

### Orthogonal complement

$S_\perp = \{y \in \mathbb{R}^m | y^T x = 0, \forall x \in S\}$  is the orthogonal complement of subspace  $S$ .

- Orthogonal complement is a subspace
- $S \cap S_\perp = \{0\}$

#### Theorem 2

Let  $S \subseteq \mathbb{R}^m$  be a subspace. For any  $y \in \mathbb{R}^m$

$$y = \Pi_S(y) + \Pi_{S_\perp}(y)$$

*Proof*

$$y = \Pi_S(y) + \Pi_{S_\perp}(y)$$

$$\begin{aligned} \iff y - \Pi_S(y) &= \Pi_{S_\perp}(y) \\ \iff z^T(y - \Pi_S(y)) &= 0 \quad \forall z \in S_\perp \\ \iff z^T \Pi_S(y) &= 0 \quad \forall z \in S_\perp \end{aligned}$$

$$\therefore \Pi_S(y) \in S$$

$$\therefore z^T \Pi_S(y) = 0 \quad \forall z \in S_\perp$$

□

## Sum of subspaces

For two sets  $X, Y \subseteq \mathbb{R}^m$ ,

$$X + Y = \{x + y \in \mathbb{R}^m\}$$

For a subspace  $S \subseteq \mathbb{R}^m$

1.  $S + S_\perp = \mathbb{R}^m$
2.  $\dim(S) + \dim(S_\perp) = m$
3.  $(S_\perp)_\perp = S$

*Proof*

1. Comes from Theorem 2
2. Let  $\{a_1, a_2, \dots, a_k\}$  be a basis for  $S$  ( $\dim(S) = k$ ) and  $\{b_1, b_2, \dots, b_\ell\}$  be a basis for  $S_\perp$  ( $\dim(S_\perp) = \ell$ ).  
 $\{a_1, a_2, \dots, a_k\} \cup \{b_1, b_2, \dots, b_\ell\}$  is  
 (a) linear independent  
 (b) Span  $S + S_\perp = \mathbb{R}^m$   
 Therefore,  $\dim(S) + \dim(S_\perp) = m$
3. Obvious

□

## Orthogonal Set

$\{a_1, a_2, \dots, a_n\}$  is orthogonal/orthonormal set if

1.  $\|a_i\|_2 = 1 \quad \forall i = 1, \dots, n$
2.  $a_i^T a_j = 0 \quad \forall i \neq j$

The orthogonal set is linear independent.

$$y = \sum_{i=1}^n \alpha_i a_i \implies \alpha_i = a_i^T y = \alpha \|a_i\|_2^2$$

Given a linear independent set  $\{a_1, a_2, \dots, a_n\}$ . The procedure of finding an orthogonal set  $\{q_1, q_2, \dots, q_n\}$  so that  $\text{span}\{a_1, a_2, \dots, a_n\} = \text{span}\{q_1, q_2, \dots, q_n\}$  is called "Gram-Schmidt" procedure.

## Orthogonal/Unitary Matrix

$Q = [q_1, q_2, \dots, q_m] \in \mathbb{R}^{n \times m}$ ,  $q_i, i = 1, \dots, m$  are orthogonal.

- If  $m = n$ ,  
 $Q$  is orthogonal matrix  $\implies Q^T Q = I_m = Q Q^T$ .
- If  $n > m$ ,  
 semi-orthogonal  $\implies Q^T Q = I_m \neq Q Q^T$ .

$Q = [q_1, q_2, \dots, q_m] \in \mathbb{C}^{n \times m}$ ,  $q_i, i = 1, \dots, m$  are orthogonal.

- If  $m = n$ ,  
 $Q$  is unitary matrix  $\implies Q^H Q = I_m = Q Q^H$ .
- If  $n > m$ ,  
 semi-unitary  $\implies Q^H Q = I_m \neq Q Q^H$ .

□

If  $Q \in \mathbb{R}^{n \times m}$  ( $n > m$ ) is semi-orthogonal,  $\exists \tilde{Q} \in \mathbb{R}^{n \times (n-m)}$  so that  $\begin{bmatrix} Q & \tilde{Q} \end{bmatrix}$  is orthogonal.

### Dimensional Theorem

For a matrix  $A \in \mathbb{R}^{m \times n}$ , it holds

$$\dim(R(A)) + \dim(\text{Null}(A)) = n$$

*Proof*

$$\begin{aligned} \dim(R(A^T)) + \dim(R(A^T)_{\perp}) &= n \\ \dim(R(A^T)) &= \text{rank}(A^T) = \text{rank}(A) = \dim(R(A)) \\ \dim(R(A^T)_{\perp}) &= \text{Null}(A) = n \end{aligned}$$

□

## Complexity of Matrix Computation

### Flops

Arithmetic operations including addition, subtraction, multiplication, division.

### Big O notation

Big O notation: "order" of the complexity.

$$f(n) = O(g(n))$$

for some  $g(n)$  if  $\exists$  a constant  $C > 0$  and  $n_0$  so that for  $n > n_0$

$$|f(n)| \leq C \cdot |g(n)|$$

Examples: for  $x \in \mathbb{R}^m$ ,  $y \in \mathbb{R}^m$ ,  $A \in \mathbb{R}^{n \times m}$ ,  $B \in \mathbb{R}^{m \times p}$

$x + y$	$m$ flops	$O(m)$
$x^T y$	$2m - 1$ flops	$O(m)$
$Ax$	$n \cdot (2m - 1)$ flops	$O(nm)$
$AB$	$p \cdot n \cdot (2m - 1)$ flops	$O(pnm)$

# Lecture 3. Least Square Problem (LSP)

## Formulation

### LS problem

Given  $y \in \mathbb{R}^m$ ,  $A \in \mathbb{R}^{m \times n}$

$$\min_{x \in \mathbb{R}^n} \|y - Ax\|_2^2$$

### Linear representation of signals

$$y = Ax$$

In practice,

$$y = Ax + v$$

where

- $y$  is the observation (Given)
- $A$  is the model (Given)
- $x$  is the parameter to be estimated/determined
- $v$  is the noise follow multi-variate Gaussian distribution

$$\text{pdf}(v) \propto e^{-\|v\|^2}$$

$\Rightarrow$  Maximum likelihood estimator

$$\max_x P(y|x) \propto \min_x \|y - Ax\|_2^2$$

## AR and Polynomial Models

### Autoregressive (AR) model

Given a time series  $y_t$ ,  $t = 0, 1, 2, \dots$ ,

$$\underbrace{y_t = \alpha_1 y_{t-1} + \alpha_2 y_{t-2} + \dots + \alpha_p y_{t-p}}_{\text{AR model}} + \underbrace{v_t}_{\text{noise}}$$

If  $\alpha_1, \alpha_2, \dots, \alpha_p$  are known, estimate  $y_t$  by

$$\hat{y}_t = \underbrace{\alpha_1 y_{t-1} + \alpha_2 y_{t-2} + \dots + \alpha_p y_{t-p}}_{\text{historical observation}}$$

Collect

$$\underbrace{\begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_T \end{bmatrix}}_{y \in \mathbb{R}^T} = \underbrace{\begin{bmatrix} y_0 & & & \\ y_1 & y_0 & & \\ \vdots & \ddots & \ddots & \\ y_{p-1} & \cdots & \cdots & y_0 \\ \vdots & \ddots & \ddots & \vdots \\ y_{T-1} & \cdots & \cdots & y_T \end{bmatrix}}_{A \in \mathbb{R}^{T \times p}} \underbrace{\begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_p \end{bmatrix}}_{x \in \mathbb{R}^p} + \underbrace{\begin{bmatrix} v_0 \\ v_1 \\ \vdots \\ v_T \end{bmatrix}}_{v \in \mathbb{R}^T}$$

$\Rightarrow$  Solve LS problem to obtain  $\alpha_i$ ,  $i = 1, \dots, p$

### Polynomial model

$$y_t = \underbrace{\alpha_0 + \alpha_1 t + \alpha_2 t^2 + \dots + \alpha_p t^p}_{\text{polynomial model}} + \underbrace{v_t}_{\text{noise}}$$

Collect

$$\underbrace{\begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_T \end{bmatrix}}_{y \in \mathbb{R}^T} = \underbrace{\begin{bmatrix} 1 & 0 & \cdots & 0 \\ 1 & 1 & \cdots & 1 \\ & \vdots & & \\ 1 & t & \cdots & t^p \\ & \vdots & & \\ 1 & T-1 & \cdots & (T-1)^p \end{bmatrix}}_{A \in \mathbb{R}^{T \times p}} \underbrace{\begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_p \end{bmatrix}}_{x \in \mathbb{R}^p} + \underbrace{\begin{bmatrix} v_0 \\ v_1 \\ \vdots \\ v_T \end{bmatrix}}_{v \in \mathbb{R}^T}$$

$\Rightarrow$  Solve LS problem to obtain  $\alpha_i$ ,  $i = 1, \dots, p$

## Basis Representation

### Basis matrix

For  $y \in \mathbb{R}^m$ , given a set of basis vectors (can be given or learned)

$$\Phi = \{\phi_1, \phi_2, \dots, \phi_n\} \in \mathbb{R}^{m \times n}$$

We model

$$y = \sum_{i=1}^n \phi_i x_i = \Phi x$$

- $m \gg n$ , assume data lies in a low-dimensional subspace
- $m \ll n$ , dictionary  $\Rightarrow x$  is sparse

### Fourier basis

Discrete Fourier Transform (DFT)

Inverse Discrete Fourier Transform (IDFT)

$$\Phi = [\phi_1, \dots, \phi_n] \in \mathbb{C}^{n \times n}$$

$$\phi_i = \frac{1}{\sqrt{n}} \begin{bmatrix} 1 \\ e^{j \frac{2\pi}{n}(i-1)} \\ e^{j \frac{2\pi}{n}(i-1) \times 2} \\ \vdots \\ e^{j \frac{2\pi}{n}(i-1) \times n-1} \end{bmatrix} \in \mathbb{R}^n$$

- $\Phi$  is IDFT matrix,  $\Phi^H$  is DFT matrix.
- $\Phi$  is unitary matrix.

## LTI System

### Lienar time invariant system

$$\begin{aligned} - \text{Linearity} & \quad \begin{cases} x_t^1 \rightarrow \boxed{\text{LTI}} \rightarrow y_t^1 \\ x_t^2 \rightarrow \boxed{\text{LTI}} \rightarrow y_t^2 \\ \alpha x_t^1 + \beta x_t^2 \rightarrow \boxed{\text{LTI}} \rightarrow \alpha y_t^1 + \beta y_t^2 \end{cases} \\ - \text{Time Invariant} & \quad x_{t-d} \rightarrow \boxed{\text{LTI}} \rightarrow y_{t-d} \end{aligned}$$

Impulse  $\delta(t)$

$$\begin{cases} \delta(t) = 0, t \neq 0 \\ \int_{-\infty}^{+\infty} \delta(t) dt = 1 \end{cases}$$

LTI system is characterized by an impulse response

$$h_t, \quad t = 0, \dots, p-1$$

$\Rightarrow x_t$  and  $y_t$  is related through "convolution"

$$\begin{aligned} y_t &= \sum_{i=1}^p h_i x_{t-i} \\ &= h_t * x_t \end{aligned}$$

$$x_t \rightarrow \overbrace{\boxed{h_t}}^{\text{LTI}} \rightarrow y_t$$

### System identification

Given  $\{y_t\}_{t=0}^{T-1}$  and  $\{x_t\}_{t=0}^{T-1}$ , find the impulse response  $\{h_t\}_{t=0}^p$

$$\begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_p \\ \vdots \\ y_T \end{bmatrix} = \begin{bmatrix} x_0 & & & & \\ x_1 & x_0 & & & \\ \vdots & \vdots & \ddots & \ddots & \\ x_p & \cdots & \cdots & \cdots & x_0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ x_{T-1} & \cdots & \cdots & \cdots & x_{T-p} \end{bmatrix} \begin{bmatrix} h_0 \\ h_1 \\ \vdots \\ h_p \end{bmatrix} + \text{noise}$$

$\Rightarrow$  Solve LS problem to obtain  $\{h_t\}_{t=0}^p$

### Deconvolution

Given  $\{y_t\}_{t=0}^{T-1}$  and  $h_{t=0}^p$ , find out  $\{x_t\}_{t=0}^{T-1}$

Put them into a linear system of equations

$$\begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_p \\ \vdots \\ y_T \end{bmatrix} = \underbrace{\begin{bmatrix} h_0 & 0 & \cdots & & 0 \\ h_1 & h_0 & 0 & \cdots & \ddots \\ \vdots & \vdots & & \ddots & \\ h_p & h_{p-1} & \cdots & h_1 & h_0 \\ \vdots & \ddots & \ddots & \vdots & \vdots \\ 0 & \cdots & h_p & h_{p-1} & \cdots & h_1 & h_0 \end{bmatrix}}_{A \in \mathbb{R}^{T \times T}} \begin{bmatrix} x_0 \\ x_1 \\ \vdots \\ x_p \\ \vdots \\ x_T \end{bmatrix} + \text{noise}$$

$\Rightarrow$  Solve LS problem to obtain  $\{x_t\}_{t=0}^{T-1}$

### Toeplitz matrix

$$A = \begin{bmatrix} a_0 & a_{-1} & \cdots & \cdots & a_{-n+1} \\ a_1 & a_0 & a_{-1} & \cdots & a_{-n+2} \\ a_2 & a_1 & a_0 & a_{-1} & \ddots \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ a_{n-1} & \cdots & a_2 & a_1 & a_0 \end{bmatrix} \in \mathbb{R}^{n \times n}$$

### Circulant Matrix

Circulant matrix is a special case of Toeplitz matrix

$$H = \begin{bmatrix} h_0 & h_{n-1} & \cdots & \cdots & h_1 \\ h_1 & h_0 & h_{n-1} & \cdots & h_2 \\ h_2 & h_1 & h_0 & h_{n-1} & \ddots \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ h_{n-1} & \cdots & h_2 & h_1 & h_0 \end{bmatrix} \in \mathbb{R}^{n \times n}$$

$$H \xLeftrightarrow{O(n \log n)} H^{-1}$$

The key property:

$$H \cdot \underbrace{\phi_i}_{\text{IDFT}} = d_i \cdot \phi_i, \quad d_i \in \mathbb{R}$$

Let

$$D = \begin{bmatrix} d_1 & & \\ & \ddots & \\ & & d_n \end{bmatrix}$$

Then

$$\begin{aligned} H \cdot \Phi &= H \cdot [\phi_1, \dots, \phi_n] \\ &= [d_1 \phi_1, \dots, d_n \phi_n] \\ &= \Phi \cdot D \end{aligned}$$

$$\underbrace{H = \Phi \cdot D \cdot \Phi^H}_{\text{eigen-decomposition of } H} \iff \Phi^H H \Phi = D$$

# Lecture 4. Eigenvalues and Eigenvectors

## Definition (Eigenvalues and Eigenvector)

Let  $A \in \mathbb{R}^{m \times m}$ , if a scalar  $\lambda \in \mathbb{C}$  and a nonzero vector  $v \in \mathbb{C}^m$  satisfy the equation

$$Av = \lambda v$$

then  $\lambda$  is called the eigenvalue of  $A$  and  $v$  is called the eigenvector of  $A$  associated with  $\lambda$ .

$\Rightarrow$

$$\underbrace{(A - \lambda I)}_{\text{singular}} v = 0$$

$P_A(\lambda) \triangleq \det(A - \lambda I) = \prod_{i=1}^m (\lambda - \lambda_i)$  is a polynomial of  $\lambda$  with degree  $m$ . The equation  $P_A(\lambda) \triangleq \det(A - \lambda I) = 0$  is called the characteristic equation of  $A$ .

For each  $\lambda_i$  find  $v_i$  such that

$$(A - \lambda_i I)v_i = 0$$

$v_i \in \text{Null}(A - \lambda_i I)$  and  $\text{Null}(A - \lambda_i I)$  is the eigenspace for  $\lambda_i$

- Even if  $A \in \mathbb{R}^{m \times m}$  is real-valued,  $\lambda_i$  and  $v_i$  could be complex. *e.g.*

$$A = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix} \Rightarrow \lambda = \pm j$$

- But if  $\lambda_i \in \mathbb{R}$  is real-valued,  $v_i \in \mathbb{R}^m$  should be in  $\mathbb{R}^m$ .

## $\Rightarrow$ (Eigenvalue decomposition)

There exists an invertible matrix  $V \in \mathbb{R}^{m \times m}$  and a diagonal matrix  $D \in \mathbb{R}^{m \times m}$  s.t.

$$A = \underbrace{VDV^{-1}}$$

eigenvalue decomposition (EVD)

$\Rightarrow$

$$AV = VD$$

$$V = [v_1, v_2, \dots, v_m] \quad , \quad D = \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_m \end{bmatrix}$$

$\Leftrightarrow$

$$Av_i = \lambda_i v_i, \quad \forall i = 1, \dots, m$$

$\Rightarrow$

$A$  is diagonalizable *iff.* there is a set of linear indep. vectors which are eigenvectors of  $A$ . **Properties**

If a matrix  $A \in \mathbb{R}^{m \times m}$  is diagonalizable

1.  $\det(A) = \prod_{i=1}^m \lambda_i$
2.  $\text{Tr}(A) = \sum_{i=1}^m \lambda_i$
3.  $\text{rank}(A)$  = the number of non-zero eigenvalues

*Proof*

1.  $\det(A) = \det(VDV^{-1}) = \det(V)\det(D)\det(V^{-1}) = \det(D) = \prod_{i=1}^m \lambda_i$
2.  $\text{Tr}(A) = \text{Tr}(VDV^{-1}) = \text{Tr}(DV^{-1}V) = \text{Tr}(D) = \sum_{i=1}^m \lambda_i$

□

## Diagonalizability

### Similar matrices

#### Definition (Similar matrices)

$A \in \mathbb{R}^{m \times m}$  and  $B \in \mathbb{R}^{m \times m}$  are similar to each other if there exists a nonsingular  $S \in \mathbb{R}^{m \times m}$  such that

$$B = SAS^{-1}$$

If  $A$  and  $B$  are similar to each other, then they have the same eigenvalues.

*Proof*

$$\begin{aligned} P_B(\lambda) &= \det(B - \lambda I) \\ &= \det(SAS^{-1} - \lambda I) \\ &= \det(S(A - \lambda I)S^{-1}) \\ &= \det(S)\det(A - \lambda I)\det(S^{-1}) \\ &= \det(A - \lambda I) \\ &= P_A(\lambda) \end{aligned}$$

□

### Diagonalizable matrix

#### Definition (Diagonalizable)

A matrix  $A \in \mathbb{R}^{m \times m}$  is diagonalizable if  $A$  is similar to a diagonal matrix.

## When Diagonalizable

### Distinct eigenvalues

Let  $\{\lambda_1, \lambda_2, \dots, \lambda_k\}_{k \in M}$  ( $k \geq 2$ ) be a subset of eigenvalues of  $A$  and  $\lambda_i \neq \lambda_j \quad \forall i \neq j$ . Then the corresponding eigenvectors, denoted with  $v_1, \dots, v_k$  respectively, are linearly independent.

If  $A$  has  $\lambda_1 \neq \lambda_2 \neq \dots \neq \lambda_k$ , then the corresponding eigenvectors  $v_1, \dots, v_k$  are linearly independent.

*Proof*

Suppose  $v_1, \dots, v_k$  are linearly **dependent**. Then there exists  $\alpha_1, \alpha_2, \dots, \alpha_k$  not all zero (without losing generality, assume  $\alpha_1 \neq 0$ ) s.t.

$$\sum_{i=1}^k \alpha_i v_i = 0$$

Then

$$A \cdot \left( \sum_{i=1}^k \alpha_i v_i \right) = \sum_{i=1}^k \alpha_i Av_i = \sum_{i=1}^k \alpha_i \lambda_i v_i = 0$$



$$\sum_{i=1}^{k-1} \alpha_i \lambda_i v_i + \alpha_k \lambda_k v_k = 0$$

On the other hand,

$$\lambda_k \cdot \left( \sum_{i=1}^k \alpha_i v_i \right) = \sum_{i=1}^k \alpha_i \lambda_k v_i = 0$$

$$\sum_{i=1}^{k-1} \alpha_i \lambda_k v_i + \alpha_k \lambda_k v_k = 0$$

(1) - (2):

$$\sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) v_i = 0$$

Then

$$\begin{aligned} A \cdot \left( \sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) v_i \right) &= \sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) A v_i \\ &= \sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) \lambda_i v_i \\ &= 0 \end{aligned}$$

$$\sum_{i=1}^{k-2} \alpha_i (\lambda_i - \lambda_k) \lambda_i v_i + \alpha_{k-1} (\lambda_{k-1} - \lambda_k) \lambda_{k-1} v_{k-1} = 0 \quad (3)$$

On the other hand

$$\lambda_{k-1} \cdot \left( \sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) v_i \right) = \sum_{i=1}^{k-1} \alpha_i (\lambda_i - \lambda_k) \lambda_{k-1} v_i = 0$$

$$\sum_{i=1}^{k-2} \alpha_i (\lambda_i - \lambda_k) \lambda_{k-1} v_i + \alpha_{k-1} (\lambda_{k-1} - \lambda_k) \lambda_{k-1} v_{k-1} = 0$$

(4)

(3) - (4):

$$\sum_{i=1}^{k-2} \alpha_i (\lambda_i - \lambda_k) (\lambda_i - \lambda_{k-1}) v_i = 0$$

Repeat until  $i = 1$

$$\alpha_1 \left( \prod_{i=1}^k (\lambda_1 - \lambda_i) \right) v_1 = 0$$

Since  $\alpha_1 \neq 0$  and  $\prod_{i=1}^k (\lambda_1 - \lambda_i) \neq 0$ , so  $v_1 = 0$  (Contradict!)

Therefore,  $v_1, \dots, v_k$  are linearly **independent**  $\square$

If  $A \in \mathbb{R}^{m \times m}$  has  $m$  distinct eigenvalues, then  $A$  is diagonalizable.

### Repeated eigenvalues

Let  $\lambda$  be an eigenvalue of  $A$  repeated  $r$  times. Define

- Algebra multiplicity:  $r$ .
- Geometric multiplicity:  $s = \dim(\text{Null}(A - \lambda I))$

Then

$$s \leq r$$

(1) *Proof*

Let  $\dim(\text{Null}(A - \bar{\lambda}I)) = s$  and  $\{q_1, \dots, q_s\}$  be an orthogonal basis for  $\text{Null}(A - \bar{\lambda}I)$ .

Then  $Q_1 = [q_1, \dots, q_s] \in \mathbb{R}^{m \times s}$  is semi-unitary and  $\exists Q_2 \in \mathbb{R}^{m \times (m-s)}$  s.t.  $Q \triangleq [Q_1, Q_2]$  is unitary.

Consider

$$Q^H A Q = \begin{bmatrix} Q_1^H A Q_1 & Q_1^H A Q_2 \\ Q_2^H A Q_1 & Q_2^H A Q_2 \end{bmatrix} = \begin{bmatrix} \bar{\lambda} I_s & Q_1^H A Q_2 \\ 0 & Q_2^H A Q_2 \end{bmatrix}$$

Then

$$\begin{aligned} P_A(\lambda) &= P_{Q^H A Q}(\lambda) \\ &= \det(Q^H A Q - \lambda I) \\ &= \det \left( \begin{bmatrix} (\bar{\lambda} - \lambda) I_s & Q_1^H A Q_2 \\ 0 & Q_2^H A Q_2 - \lambda I_{m-s} \end{bmatrix} \right) \\ &= \underbrace{(\bar{\lambda} - \lambda)^s}_{\bar{\lambda} \text{ repeats } s \text{ times}} \det(Q_2^H A Q_2 - \lambda I_{m-s}) \\ &= 0 \end{aligned}$$

So

$$r \geq s$$

$\square$

If  $s = r$ ,  $A$  is diagonalizable.

## Symmetric/Hermitian Matrices

- **Symmetric**  $A = A^T$   $A \in \mathbb{R}^{m \times n}$

- **Hermitian**  $A = A^H$   $A \in \mathbb{C}^{m \times n}$

**Property 1:** The eigenvalues of a Hermitian matrix  $A \in \mathbb{C}^{m \times n}$  are real-valued.

*Proof*

Let  $\lambda$  be the eigenvalue of  $A$ ,  $v$  be the eigenvector of  $A$  and  $\|v\|_2 = 1$ .

$$\implies Av = \lambda v$$

$$\implies v^H Av = \lambda v^H v = \lambda$$

$$\stackrel{\text{take } H}{\implies} v^H A^H v = \lambda^*$$

Since  $A$  is Hermitian,  $A = A^H$  and hence  $\lambda = \lambda^*$

So  $\lambda$  is real-valued  $\square$

**Property 2:** If  $A \in \mathbb{C}^{m \times n}$  is Hermitian and that

- $\lambda_i \neq \lambda_j$  be two eigenvalues,
- $v_i, v_j$  be the corresponding eigenvectors.

Then the eigenvectors are mutually orthogonal.

*Proof*

$$\begin{cases} v_i^H A v_j = v_i^H (\lambda_j v_j) = \lambda_j v_i^H v_j \\ v_j^H A v_i = v_j^H (\lambda_i v_i) = \lambda_i v_j^H v_i \stackrel{\text{take } H}{\implies} v_i^H A^H v_j = \lambda_i v_i^H v_j \end{cases}$$

Since  $v_i^H A^H v_j = v_i^H A v_j$ , so we have

$$(\lambda_i - \lambda_j) v_i^H v_j = 0 \implies v_i^H v_j = 0$$

$\square$

### Unitarily diagonalizable

If Hermitian/symmetric  $A$  is diagonalizable

$$A = VDV^{-1}$$

where

- $V$  contains orthogonal eigenvectors and is unitary matrix,
- $D$  real-valued diagonal eigenvalues.

Therefore,  $VV^H = I$  and  $V^{-1} = V^H$ . Thus,

$$A = VDV^H$$

And we say  $A$  is unitarily diagonalizable.

### Schur triangularization

Any  $A \in \mathbb{C}^{m \times m}$  can be decomposed as an upper triangular matrix.

#### Theorem 3 (Schur triangularization)

Given  $A \in \mathbb{C}^{m \times m}$  with eigenvalues  $\lambda_1, \lambda_2, \dots, \lambda_m \in \mathbb{C}$  there exists a unitary matrix  $U \in \mathbb{C}^{m \times m}$  such that

$$U^H AU = T \iff A = UTU^H$$

where  $T$  is an upper triangular matrix with  $[T]_{ii} = \lambda_i, i = 1, 2, \dots, m$

If  $A \in \mathbb{C}^{m \times m}$ , by Schur triangularization, there exists unitary  $U \in \mathbb{C}^{m \times m}$  s.t.

$$U^H AU = T$$

Then by taking  $H$ , we have

$$U^H A^H U = T^H$$

$\implies$

$$T = T^H$$

Since  $T$  is an upper triangular matrix, so  $T^H$  is a lower triangular matrix. Therefore,  $T$  is diagonal.  $\implies$

$$U^H AU = T$$

Thus,  $A$  is diagonalizable.

*Proof* for Schur Triangularization (**Theorem 3**)

*This proof is also an algorithm.*

Let  $v_1 \in \mathbb{R}^m$  be the eigenvector of  $A$  associated with  $\lambda_1$ , i.e.,  $Av_1 = \lambda_1 v_1$ . Without losing generality, we can assume  $v_1^H v_1 = 1$  and  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_m$

Construct a unitary matrix  $U_1$

$$U_1 = [v_1 \quad V_1] \in \mathbb{C}^{m \times m}$$

where  $V_1 \in \mathbb{C}^{m \times (m-1)}$ .  $U_1$  being unitary indicates

$$v_1^H V_1 = \vec{0}^T$$

Then we have

$$U_1^H AU_1 = \begin{bmatrix} v_1^H \\ V_1^H \end{bmatrix} A [v_1 \quad V_1] = \begin{bmatrix} v_1^H Av_1 & v_1^H AV_1 \\ V_1^H Av_1 & V_1^H AV_1 \end{bmatrix}$$

where

$$\begin{cases} v_1^H Av_1 = \lambda_1 \\ v_1^H AV_1 \neq \vec{0} \\ V_1^H Av_1 = \vec{0} \end{cases}$$

So,

$$U_1^H AU_1 = \begin{bmatrix} \lambda_1 & \times \\ 0 & V_1^H AV_1 \end{bmatrix}$$

Define

$$A_1 \triangleq V_1^H AV_1 \in \mathbb{C}^{(m-1) \times (m-1)}$$

Then

$$\begin{aligned} P_A(\lambda) &= \det(A - \lambda I) \\ &= \det(U_1^H AU_1 - \lambda I)^\dagger \\ &= \det\left(\begin{bmatrix} \lambda_1 & \times \\ 0 & A_1 \end{bmatrix} - \lambda I\right) \\ &= \det\left(\begin{bmatrix} \lambda_1 - \lambda & \times \\ 0 & A_1 - \lambda I \end{bmatrix}\right) \\ &= (\lambda_1 - \lambda) \underbrace{\det(A_1 - \lambda I)}_{=P_{A_1}(\lambda)} \end{aligned}$$

$\dagger$ :  $A = U_1(U_1^H AU_1)U_1^{-1}$  so  $A$  is similar to  $U_1^H AU_1$  and hence  $A$  and  $U_1^H AU_1$  have same eigenvalues.

Therefore,  $A_1$  has eigenvalues  $\lambda_2, \lambda_3, \dots, \lambda_m$

Let  $v_2 \in \mathbb{C}^{m-1}$  be the eigenvector of  $A_1$  associated with  $\lambda_2$ .

Construct a unitary matrix  $U_2$

$$U_2 = [v_2 \quad V_2] \in \mathbb{C}^{(m-1) \times (m-1)}$$

where  $V_2 \in \mathbb{C}^{(m-1) \times (m-2)}$ .  $U_2$  being unitary indicates

$$v_2^H V_2 = \vec{0}^T$$

$\implies$

$$U_2^H A_1 U_2 = \begin{bmatrix} \lambda_2 & \times \\ 0 & V_2^H A_1 V_2 \end{bmatrix}$$

Let  $\tilde{U}_2 = \begin{bmatrix} 1 & 0^T \\ 0 & U_2 \end{bmatrix}$ , then  $\tilde{U}_2$  is unitary.

So,

$$\begin{aligned} \tilde{U}_2^H [U_1^H AU_1] \tilde{U}_2 &= \tilde{U}_2^H \begin{bmatrix} \lambda_1 & \times \\ 0 & A_1 \end{bmatrix} \tilde{U}_2 \\ &= \begin{bmatrix} \lambda_1 & \times \\ 0 & U_2^H A_1 U_2 \end{bmatrix} \\ &= \begin{bmatrix} \lambda_1 & \times \\ 0 & \begin{bmatrix} \lambda_2 & \times \\ 0 & V_2^H A_1 V_2 \end{bmatrix} \end{bmatrix} \\ &= \begin{bmatrix} \lambda_1 & \times & \times \\ 0 & \lambda_2 & \times \\ 0 & 0 & V_2^H A_1 V_2 \end{bmatrix} \end{aligned}$$

Repeat this operations and we can finally get

$$\begin{aligned} \tilde{U}_m^H \dots \tilde{U}_3^H \tilde{U}_2^H (U_1 A U_1) \tilde{U}_2 \tilde{U}_3 \dots \tilde{U}_m \\ \underbrace{\hspace{10em}}_{\triangleq U \text{ unitary}} \\ = \begin{bmatrix} \lambda_1 & \times & \times & \times \\ & \lambda_2 & \times & \times \\ & & \ddots & \times \\ & & & \lambda_m \end{bmatrix} \triangleq T \end{aligned}$$

Therefore,

$$U^H A U = T$$

□

## Variational Characterization of Eigenvalues

### Rayleigh quotient

**Definition (Rayleigh quotient)**

$$R(A, x) = \frac{x^H A x}{\|x\|_2^2} = \left( \frac{x}{\|x\|_2} \right)^H A \left( \frac{x}{\|x\|_2} \right)$$

**Theorem 4 (Rayleigh-Ritz theorem)**

Let  $A \in \mathbb{C}^{m \times m}$  be Hermitian, then

$$\begin{aligned} \lambda_{\max}(A) &= \max_{x \in \mathbb{C}^m} \frac{x^H A x}{\|x\|_2^2} \\ &= \max_{x \in \mathbb{C}^m} x^H A x \quad (\|x\|_2 = 1) \\ \lambda_{\min}(A) &= \min_{x \in \mathbb{C}^m} \frac{x^H A x}{\|x\|_2^2} \end{aligned}$$

and

$$\lambda_{\min}(A) \cdot \|x\|_2^2 \leq x^H A x \leq \lambda_{\max}(A) \cdot \|x\|_2^2$$

*Proof*

According to the EVD of  $A$ ,

$$A = V \Lambda V^H = \sum_{i=1}^m (\lambda_i \cdot v_i v_i^H)$$

$\Rightarrow$

$$\begin{aligned} x^H A x &= x^H V \Lambda V^H x \\ &= \sum_{i=1}^m (\lambda_i \cdot x^H v_i v_i^H x) \\ &= \sum_{i=1}^m (\lambda_i \cdot |v_i^H x|^2) \\ &\leq \lambda_{\max}(A) \cdot \sum_{i=1}^m |v_i^H x|^2 \end{aligned}$$

Since

$$\sum_{i=1}^m |v_i^H x|^2 = x^H V V^H x = \|x\|_2^2$$

So,

$$x^H A x \leq \lambda_{\max}(A) \cdot \|x\|_2^2$$

$\Rightarrow$

$$\frac{x^H A x}{\|x\|_2^2} \leq \lambda_{\max}(A)$$

The equality holds only when  $x = v_1$ , which is an eigenvector associated with  $\lambda_{\max}(A)$  and is called the principal eigenvector. □

### Courant-Fischer theorem

**Theorem 5 (Courant-Fischer theorem)**

Let  $A \in \mathbb{C}^{m \times m}$  be Hermitian and  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_m$  are ordered eigenvalues. Then

$$\begin{aligned} \lambda_k &= \min_{\omega_1, \omega_2, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \omega_2, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \right\} \\ &= \max_{\omega_{k+1}, \omega_{k+2}, \dots, \omega_m} \left\{ \min_{x \in \mathbb{C}^m, x \perp \omega_{k+1}, \omega_{k+2}, \dots, \omega_m} \frac{x^H A x}{\|x\|_2^2} \right\} \end{aligned}$$

**Proof ideas:**

1.

$$\max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{x^H A x}{\|x\|_2^2} = \lambda_k$$

where  $v_k$  is the eigenvector associated with  $\lambda_k$  of  $A$ .

2.

$$\max_{x \in \mathbb{C}^m, x \perp \omega_1, \omega_2, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \geq \lambda_k$$

1+2  $\Rightarrow$

$$\min_{\omega_1, \omega_2, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \omega_2, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \right\} = \lambda_k$$

attained when  $\omega_i = v_i$  ( $i = 1, \dots, k-1$ )

*Proof*

1.

Since  $A = V \Lambda V^H = \sum_{i=1}^m \lambda_i v_i v_i^H$ ,

$$\begin{aligned} &\max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{x^H A x}{\|x\|_2^2} \\ &= \max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{\sum_{i=1}^m \lambda_i |v_i^H x|^2}{\|x\|_2^2} \\ &= \max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{\sum_{i=k}^m \lambda_i |v_i^H x|^2}{\|x\|_2^2} \\ &\leq \max_{x \in \mathbb{C}^m} \frac{\sum_{i=k}^m \lambda_i |v_i^H x|^2}{\|x\|_2^2} \quad (\text{Relax constraint}) \\ &\leq \max_{x \in \mathbb{C}^m} \frac{\lambda_k \sum_{i=k}^m |v_i^H x|^2}{\|x\|_2^2} \quad \dagger \\ &\leq \lambda_k \quad (\text{equality attained if } x = v_k) \end{aligned}$$

$\dagger$ :  $\sum_{i=k}^m |v_i^H x|^2 = x^H V V^H x = \|x\|_2^2$ .

Therefore,

$$\max_{x \in \mathbb{C}^m, x \perp v_1, v_2, \dots, v_{k-1}} \frac{x^H A x}{\|x\|_2^2} = \lambda_k$$

2.

Let  $y = V^H x$

$$\begin{aligned} &\max_{x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \\ &= \max_{x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H V \Lambda V^H x}{\|x\|_2^2} \\ &= \max_{y \perp V^H \omega_1, \dots, V^H \omega_{k-1}} \frac{y^H \Lambda y}{\|y\|_2^2} \\ &\geq \max_{y \perp V^H \omega_1, \dots, V^H \omega_{k-1}, y_{k+1}=y_{k+2}=\dots=y_m=0, \|y\|_2=1} y^H \Lambda y \\ &= \max_{y \perp V^H \omega_1, \dots, V^H \omega_{k-1}, |y_1|^2 + |y_2|^2 + \dots + |y_k|^2 = 1} \sum_{i=1}^k \lambda_i |y_i|^2 \\ &\geq \lambda_k \end{aligned}$$

□

## Weyl's inequality

For  $A, B \in \mathbb{C}^{m \times m}$  be two Hermitian matrices.  
Then

1.  $\lambda_k(A) + \lambda_m(B) \leq \lambda_k(A+B) \leq \lambda_k(A) + \lambda_1(B)$
2.  $\lambda_{k+1}(A) \leq \lambda_k(X \pm z z^H) \leq \lambda_{k-1}(A)$
3.  $\lambda_{k+r}(A) \leq \lambda_k(A+B) \leq \lambda_{k-r}(A)$  where  $\text{rank}(B) \leq r$
4.  $\lambda_{j+k-1}(A+B) \leq \lambda_j(A) + \lambda_k(B) \leq \lambda_{j+k-m}(A+B)$
5. Let  $A_I$  be the principal submatrix of  $A$  associated with subset  $I = \{i_1, i_2, \dots, i_r\} \subseteq \{1, 2, \dots, m\}$   
 $\lambda_{k+m-r}(A) \leq \lambda_k(A_I) \leq \lambda_k(A)$
6. Let  $U \in \mathbb{C}^{m \times r}$  be semiunitary  
 $\lambda_{k+m-r} \leq \lambda_k(U^H A U) \leq \lambda_k(A)$

*Proof* 1.

$$\lambda_k(A+B) = \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H (A+B)x}{\|x\|_2^2} \right\}$$

where

$$\frac{x^H (A+B)x}{\|x\|_2^2} = \frac{x^H A x}{\|x\|_2^2} + \underbrace{\frac{x^H B x}{\|x\|_2^2}}_{\leq \lambda_1(B)}$$

So,

$$\begin{aligned} & \lambda_k(A+B) \\ & \leq \underbrace{\min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H A x}{\|x\|_2^2} \right\}}_{=\lambda_k(A)} + \lambda_1(B) \end{aligned}$$

Therefore,

$$\lambda_k(A+B) \leq \lambda_k(A) + \lambda_1(B)$$

*Proof* 2.

$$\begin{aligned} & \lambda_k(A \pm z z^H) \\ & = \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}} \frac{x^H (A \pm z z^H) x}{\|x\|_2^2} \right\} \\ & \geq \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}, \boxed{z}} \frac{x^H (A \pm z z^H) x}{\|x\|_2^2} \right\} \\ & = \min_{\omega_1, \dots, \omega_{k-1}, \omega_k = z} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}, \omega_k} \frac{x^H A x}{\|x\|_2^2} \right\} \\ & \geq \min_{\omega_1, \dots, \omega_{k-1}, \omega_k} \left\{ \max_{x \in \mathbb{C}^m, x \perp \omega_1, \dots, \omega_{k-1}, \omega_k} \frac{x^H A x}{\|x\|_2^2} \right\} \\ & = \lambda_{k+1}(A) \end{aligned}$$

*Proof* 3.

$$\begin{aligned} & \because B = \sum_{i=1}^r \lambda_i(B) u_i u_i^H \\ & \therefore A+B = A + \sum_{i=1}^{r-1} \lambda_i(B) u_i u_i^H + \lambda_r(B) u_r u_r^H \end{aligned}$$

Define  $B^{r-k} \triangleq \sum_{i=1}^{r-k} \lambda_i(B) u_i u_i^H$   $k = 0, 1, \dots, r$

$$\lambda_{k+1}(A+B^{r-1}) \leq \lambda_k(A+B^{r-1} + \lambda_r(B) u_r u_r^H) \leq \lambda_{k-1}(A+B^{r-1})$$

Do again and we get

$$\begin{aligned} \lambda_{k+2}(A+B^{r-2}) & \leq \lambda_{k+1}(A+B^{r-2} + \lambda_{r-1}(B) u_{r-1} u_{r-1}^H) \\ & = \lambda_{k+1}(A+B^{r-1}) \end{aligned}$$

Repeat  $r$  times, we can get

$$\lambda_{k+r}(A) \leq \lambda_k(A+B) \leq \lambda_{k-r}(A)$$

where  $\text{rank}(B) \leq r$  □

*Proof* 4.

Define  $A_{j-1} = \sum_{i=1}^{j-1} \lambda_i(A) v_i v_i^H$  (if  $A = V \Lambda V^H$ ).

Define  $B_{k-1} = \sum_{i=1}^{k-1} \lambda_i(B) u_i u_i^H$  (if  $B = U \Lambda U^H$ ).

Then

$$\begin{aligned} \lambda_j(A) & = \lambda_1(A - A_{j-1}) \\ \lambda_k(B) & = \lambda_1(B - B_{k-1}) \end{aligned}$$

$\Rightarrow$

$$\begin{aligned} \lambda_j(A) + \lambda_k(B) & = \lambda_1(A - A_{j-1}) + \lambda_1(B - B_{k-1}) \\ & \geq \lambda_1(A+B - (A_{j-1} + B_{k-1})) \\ & \geq \lambda_{j+k-1}(A+B)^\dagger \end{aligned}$$

$\dagger$ :  $\text{rank}((A_{j-1} + B_{k-1})) \leq j+k-2$  and from 3. □

*Proof* 5.

$$\begin{aligned} & \lambda_k(A) \\ & = \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \perp \omega_1, \dots, \omega_{k-1}, \|x\|_2=1} x^H A x \right\} \\ & \geq \min_{\omega_1, \dots, \omega_{k-1}} \left\{ \max_{x \perp \omega_1, \dots, \omega_{k-1}, \|x\|_2=1, x \in \text{span}\{e_{i_1}, \dots, e_{i_r}\}} x^H A x \right\}^\dagger \\ & = \min_{v_1, \dots, v_{k-1}} \left\{ \max_{y \perp v_1, \dots, v_{k-1}, \|y\|_2=1} y^H A y \right\}^\ddagger \\ & = \lambda_k(A_I) \end{aligned}$$

$\dagger$ :  $e_i = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}$ , only the  $i$ th entry equal to 1. □

$\ddagger$ : Define  $y \in \mathbb{C}^r$ ,  $y_\ell = x_{i_\ell}$  ( $\ell = 1, \dots, r$ ) and take  $v_i \in \mathbb{C}^r$ ,  $(v_i)_\ell = (\omega_i)_{i_\ell}$  □

*Proof* 6.

Let  $\tilde{U} = \left[ \underbrace{U}_{\in \mathbb{C}^{m \times r}}, \underbrace{\tilde{U}}_{\in \mathbb{C}^{m \times r}} \right]$  be unitary. Then

$$\tilde{U}^H A \tilde{U} = \begin{bmatrix} U^H A U & \times \\ \times & \times \end{bmatrix}$$

where  $U^H A U$  is the principal submatrix of  $\tilde{U}^H A \tilde{U}$  with  $I = 1, 2, \dots, r$

Therefore,

$$\lambda_k(A) = \lambda_k(\tilde{U}^H A \tilde{U}) \geq \lambda_k(U^H A U)$$

□

□

## The Power Method

The power method is a simple method to obtain eigenvector/eigenvalue of a matrix, particularly for large scale problem.

**Algorithm 1 (Power Method)**

Given a matrix  $A \in \mathbb{C}^{m \times m}$  and  $v^0 \in \mathbb{C}^m$

$$v^{(k+1)} \leftarrow \frac{A \cdot v^{(k)}}{\|A \cdot v^{(k)}\|_2}$$

until a stopping condition is met.

$\Rightarrow v^{(k+1)}$  takes as an estimate of the **principle eigenvector** of  $A$ .

Assumption:

- $A$  is diagonalizable (in particular,  $A$  is Hermitian (symmetric)).  $A = V\Lambda V^H$ .
- $|\lambda_1| \geq |\lambda_2| \geq \dots \geq |\lambda_m|$  and  $|\lambda_1| > |\lambda_2|$

$$v^{(k+1)} = \frac{A \cdot v^{(k)}}{\|A \cdot v^{(k)}\|_2} = \frac{A^{k+1} \cdot v^{(0)}}{\|A^{k+1} \cdot v^{(0)}\|_2}$$

Consider  $A^k x$ ,  $x$  satisfies  $[V^H x]_1 \neq 0$ . ( $x$  is  $v^{(0)}$ )

Let  $\alpha = V^H x$

$$A^k \cdot x = (V\Lambda^k V^H) \cdot x = V\Lambda^k \alpha$$

Let  $V = [v_1, v_2, \dots, v_m]$

$$\begin{aligned} A^k \cdot x &= V\Lambda^k \alpha \\ &= \sum_{i=1}^m v_i \cdot \lambda_i^k \cdot \alpha_i \\ &= v_1(\lambda_1 \alpha_1) + \sum_{i=2}^m \alpha_i \lambda_i^k v_i \\ &= \lambda_1^k \alpha_1 \underbrace{\left( v_1 + \sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \left( \frac{\lambda_i}{\lambda_1} \right)^k v_i \right)}_{\triangleq r_k(\text{residue})} \end{aligned}$$

$\Rightarrow$

$$\begin{aligned} \|r_k\|_2 &= \left\| \sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \left( \frac{\lambda_i}{\lambda_1} \right)^k v_i \right\|_2 \\ &\leq \sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \cdot \underbrace{\left( \frac{\lambda_i}{\lambda_1} \right)^k}_{=1} \cdot \|v_i\|_2 \\ &\leq \underbrace{\left| \frac{\lambda_2}{\lambda_1} \right|^k}_{<1} \cdot \sum_{i=2}^m \left| \frac{\alpha_i}{\alpha_1} \right| \rightarrow 0 \text{ as } k \rightarrow \infty \end{aligned}$$

Consider  $c_k = \frac{|\alpha_1| |\lambda_1^k|}{\alpha_1 \lambda_1^k}$  and  $\lim_{k \rightarrow \infty} c_k \cdot \frac{A^k \cdot x}{\|A^k \cdot x\|_2}$

$$\begin{aligned} c_k \cdot \frac{A^k \cdot x}{\|A^k \cdot x\|_2} &= \frac{|\alpha_1| |\lambda_1^k|}{\alpha_1 \lambda_1^k} \cdot \frac{1}{\|A^k \cdot x\|_2^\dagger} \cdot \lambda_1^k \alpha_1 \left( v_1 + \sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \left( \frac{\lambda_i}{\lambda_1} \right)^k v_i \right) \\ &= \frac{|\alpha_1| |\lambda_1^k|}{\sqrt{\sum_{i=1}^m (\lambda_i^k)^2 \alpha_i^2}}^\ddagger \cdot \left( v_1 + \underbrace{\sum_{i=2}^m \frac{\alpha_i}{\alpha_1} \left( \frac{\lambda_i}{\lambda_1} \right)^k v_i}_{\rightarrow 0 \text{ as } k \rightarrow \infty} \right) \end{aligned}$$

$$\dagger: \|A^k x\|_2 = \|V\Lambda^k V^H x\|_2 = \|V\Lambda^k \alpha\|_2 = \|\Lambda^k \alpha\|_2 = \sqrt{\sum_{i=1}^m (\lambda_i^k \alpha_i)^2}$$

$$\begin{aligned} \ddagger: \frac{|\alpha_1| |\lambda_1^k|}{\sqrt{\sum_{i=1}^m (\lambda_i^k)^2 \alpha_i^2}} &= \frac{|\alpha_1| |\lambda_1^k|}{\sqrt{(\lambda_1^k)^2 \alpha_1^2 + \sum_{i=2}^m (\lambda_i^k)^2 \alpha_i^2}} \\ &= \frac{1}{\sqrt{1 + \sum_{i=2}^m \left( \frac{\lambda_i}{\lambda_1} \right)^{2k} \left( \frac{\alpha_i}{\alpha_1} \right)^2}} \end{aligned}$$

Since  $\sum_{i=2}^m \left( \frac{\lambda_i}{\lambda_1} \right)^{2k} \left( \frac{\alpha_i}{\alpha_1} \right)^2 \rightarrow 0$  as  $k \rightarrow \infty$ , so  $\frac{|\alpha_1| |\lambda_1^k|}{\sqrt{\sum_{i=1}^m (\lambda_i^k)^2 \alpha_i^2}} \rightarrow 1$  as  $k \rightarrow \infty$ .

Therefore,  $\lim_{k \rightarrow \infty} c_k \cdot \frac{A^k \cdot x}{\|A^k \cdot x\|_2} \rightarrow v_1$ .

**Deflation:**

$$\begin{aligned} A &= \sum_{i=1}^m \lambda_i v_i v_i^H \\ &= \lambda_1 v_1 v_1^H + \sum_{i=2}^m \lambda_i v_i v_i^H \end{aligned}$$

Obtain  $A^{(1)} = A - \lambda_1 v_1 v_1^H$  and apply power method to obtain  $v_2$

Obtain  $A^{(2)} = A - \lambda_2 v_2 v_2^H$  and apply power method to obtain  $v_3$

...

# Lecture 5. PageRank

## PageRank Model

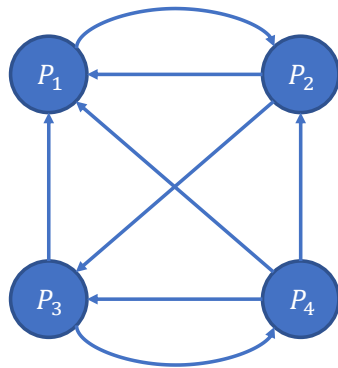
### Importance score

- $v_i$ : importance score of page  $i$
- $c_j$ : number of outgoing links of page  $j$
- $\mathcal{L}_i$ : set of pages that refers to page  $i$

$$\begin{aligned} v_i &= |\mathcal{L}_i| \\ &\Downarrow \\ v_i &= \sum_{j \in \mathcal{L}_i} v_j \\ &\Downarrow \\ v_i &= \sum_{j \in \mathcal{L}_i} \frac{v_j}{c_j} \end{aligned}$$

### The model

e.g.



$$\begin{aligned} v_1 &= 1 \cdot v_2 + \frac{1}{2}v_3 + \frac{1}{3}v_4 \\ v_2 &= 1 \cdot v_1 + \frac{1}{3}v_4 \end{aligned}$$

...

$\Rightarrow$

$$\underbrace{\begin{bmatrix} 0 & 1/2 & 1/2 & 1/3 \\ 1 & 0 & 0 & 1/3 \\ 0 & 1/2 & 0 & 1/3 \\ 0 & 0 & 1/2 & 0 \end{bmatrix}}_{\text{link matrix}} \begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{bmatrix} = \begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{bmatrix}$$

$\Rightarrow$

$$Av = v, \quad v \geq 0$$

## PageRank Problem

Find solution  $v \geq 0$  satisfying  $Av = v$

The four questions:

1. Does  $A$  admit a solution of  $Av = v$ ?  
(Does  $A$  has eigenvalue equal to 1?)
2. Does  $Av = v$  admit a non-negative solution?  
( $v \geq 0$  or  $v \leq 0$ )

3. Is the solution unique?  
(unique ranking)
4. Can we obtain the solution efficiently?  
(Power method applicable?)

### Answer to Question 1

Does  $A$  admit a solution of  $Av = v$ ? (Does  $A$  has eigenvalue equal to 1?)

Note that  $A \geq 0$  and column-sum equal to 1 (column-stochastic matrix).

$$\Rightarrow \mathbf{1}^T A = \mathbf{1}^T$$

$$\Rightarrow A^T \cdot \mathbf{1} = \mathbf{1}$$

$$\Rightarrow A^T \text{ has an eigenvalue equal to 1}$$

Since  $A$  and  $A^T$  have the same set of eigenvalues

$$\Rightarrow A \text{ has an eigenvalue equal to 1}$$

### Theorem 6

Let  $A$  be non-negative and column(row)-stochastic, then

1.  $\lambda = 1$  is an eigenvalue of  $A$
2.  $|\lambda| \leq 1$  for all eigenvalues of  $A$

*Proof* Let  $Av = \lambda v$  for some  $|\lambda| \neq 1$ .

Let  $|v_k| = \max\{|v_1|, |v_2|, \dots, |v_m|\}$ .

$$\lambda v_k = \sum_{i=1}^m [A]_{ki} v_i$$

$\Rightarrow$

$$\begin{aligned} |\lambda| |v_k| &= \left| \sum_{i=1}^m [A]_{ki} v_i \right| \\ &\leq \sum_{i=1}^m |[A]_{ki}| \cdot |v_i| \\ &\leq |v_k| \cdot \sum_{i=1}^m |[A]_{ki}| \\ &= |v_k| \end{aligned}$$

$\Rightarrow$

$$|\lambda| \leq 1$$

□

### Disjoint page networks

$$A = \begin{bmatrix} A_1 & 0 \\ 0 & A_2 \end{bmatrix}$$

$$A_1 v_1 = v_1$$

$$A_2 v_2 = v_2$$

$$A \left( \alpha_1 \begin{pmatrix} v_1 \\ 0 \end{pmatrix} + \alpha_2 \begin{pmatrix} 0 \\ v_2 \end{pmatrix} \right) = \left( \alpha_1 \begin{pmatrix} v_1 \\ 0 \end{pmatrix} + \alpha_2 \begin{pmatrix} 0 \\ v_2 \end{pmatrix} \right)$$

### Modified model

$$M \triangleq (1 - \beta)A + \beta \left( \frac{\mathbf{1}\mathbf{1}^T}{m} \right) > 0$$

This is positive column-stochastic

$$\mathbf{1}^T M = (1 - \beta)\mathbf{1}^T A + \beta \mathbf{1}^T \left( \frac{\mathbf{1}\mathbf{1}^T}{m} \right) = \mathbf{1}^T$$

### Answer to Question 2

Does  $Av = v$  admit a non-negative solution? ( $v \geq 0$  or  $v \leq 0$ )

*Proof*

1)  $v \in S \iff 1^T v = 0$   
 2)

□

#### Theorem 7

For positive and column-stochastic matrix  $M$ , any eigenvector associated with  $\lambda_1$  has either all positive or all negative components.

*Proof* Note that for  $y \in \mathbb{R}^m$ ,  $|\sum_{i=1}^m y_i| < \sum_{i=1}^m |y_i|$  if  $\{y_i\}$  have mixed signs.

Suppose the above theorem is not true. Then we have

$$|v_i| = \left| \sum_{j=1}^m [M]_{ij} v_j \right| < \sum_{j=1}^m [M]_{ij} |v_j|$$

$$\implies \sum_{i=1}^m |v_i| < \sum_{i=1}^m \sum_{j=1}^m [M]_{ij} v_j = \sum_{j=1}^m \sum_{i=1}^m [M]_{ij} v_j = \sum_{j=1}^m |v_j|$$

Contradict! So the above theorem is true. □

### Answer to Question 3

Is the solution unique? (unique ranking)

Since  $Mv = v \implies (M - I)v = 0$ , show

$$\dim(\text{Null}(M - I)) = 1$$

i.e., there is a unique ranking.

*Proof*

Simple lemma: Let  $a$  and  $b$  be linearly independent, then  $\exists x \in \text{span}\{a, b\}$  which has mixed sign.

*Proof for this lemma:*

If  $1^T q = 0 \implies q$  has mixed sign.

If  $1^T q \neq 0$ , set  $\alpha_1 = \frac{-1^T b}{1^T q}$ ,  $\alpha_2 = 1$  and let  $x = \alpha_1 \cdot a + \alpha_2 \cdot b$

$\implies x$  satisfies  $1^T x = 0$ , so  $x$  has mixed sign.

Suppose  $\dim(\text{Null}(M - I)) > 1$ , then for  $q_1, q_2 \in \text{Null}(M - I)$  that are linear independent.  $\implies \exists v = \alpha_1 q_1 + \alpha_2 q_2 \in \text{Null}(M - I)$

$\implies Mv = v$  which has mixed sign. (Contradict!)

$\implies \dim(\text{Null}(M - I)) = 1$  □

### Answer to Question 4

Can we obtain the solution efficiently? (Power method applicable?)

Property: Let  $M$  be positive and column-stochastic.

Let  $S = \{v \in \mathbb{R}^m | 1^T v = 0\}$ .

Then,

1.  $Mv \in S$  if  $v \in S$
2.  $\|Mv\|_1 \leq c\|v\|_1$  with  $c < 1 \forall$  any  $v \in S$
3. the eigenvector associated with  $\lambda = 1$ , denoted  $v_1$ , can be obtainable by the power method

## Lecture 6. Positive Semidefinite Matrix

For  $A \in \mathbb{C}^{m \times m}$ , if

- $A$  is symmetric, we denote  $A \in S^m$
- $A$  is Hermitian, we denote  $A \in H^m$

### Definition

We say  $A \in H^m$  is positive semidefinite (PSD) if

$$x^H Ax > 0 \quad \forall x \in \mathbb{C}^m$$

We say  $A \in H^m$  is positive definite (PD) if

$$x^H Ax > 0 \quad \forall x \in \mathbb{C}^m, x \neq 0$$

We say  $A \in H^m$  is indefinite if  $A$  is not PSD.

- $\det(A) \geq 0 (> 0)$

If  $A$  is PD, then  $A$  is invertible.

$$A = V \Lambda V^H$$

$$A^{-1} = V\Lambda^{-1}V^H$$

**Property 2:**  $A \in H^m$  can be factorized as  $A = BB^H$  for some matrix  $B$  if and only if  $A$  is PSD.

*Proof*

 $\Rightarrow$ 

$$x^H A x = x^H B B^H x = \|B^H x\|^2 \geq 0 \implies A \text{ is PSD}$$

 $\Leftarrow$ 

If  $A$  is PSD, define  $\Lambda^{1/2} = \text{diag}(\sqrt{\lambda_1}, \sqrt{\lambda_2}, \dots, \sqrt{\lambda_m})$

$$\begin{aligned} A &= V\lambda V^H \\ &= V\Lambda^{1/2}\Lambda^{1/2}V^H \\ &= V\Lambda^{1/2}Q^HQ\Lambda^{1/2}V^H \end{aligned}$$

where  $Q$  is a Hermitian matrix.

Define  $B \triangleq V\Lambda^{1/2}Q^H$ . Then,  $A = BB^H$ .  $\square$

### Theorem 2

Consider  $A \in H^m$ ,  $B \in \mathbb{C}^{m \times n}$  and  $C = B^H A B$

- If  $A$  is PSD, then  $C$  is PSD.
- If  $A$  is PD, then  $C$  is PD if and only if  $rank(B) = n$
- If  $B$  is square and invertible, then  $A$  is PSD(PD) if and only if  $C$  is PSD(PD)

*Proof* of 2)

( $\implies$ ) Show  $C$  is PSD  $\implies \text{rank}(B) = n$

Assume  $B$  is not full column rank.

Consider  $x^H C x = x^H B^H A B x$ , then there exists  $x \neq 0$  s.t.  $Bx = 0 \implies x^H C x = 0$  (Contradict!) ( $\Leftarrow$ ) Show  $\text{rank}(B) = n \implies$  Show  $C$  is PSD  
 $\text{rank}(B) = n \implies Bx = 0$  iff  $x = 0$ . So  $\forall x \neq 0 \implies Bx \neq 0$ .

$\implies x^H C x = x^H B^H A B x > 0$  as  $A$  is PD. So  $C$  is PD.  $\square$

## Eigenvalues of PSD matrices

### Theorem 1

$A \in H^m$  is PSD(PD) if and only if  $\lambda_1, \lambda_2, \dots, \lambda_3 \geq 0(> 0)$

*Proof*

Consider EVD of  $A$  as

$$A = V\Lambda V^H, \quad \Lambda = diag(\lambda_1, \lambda_2, ..., \lambda_m)$$

Consider  $x^H A x = x^H V \Lambda V^H x = z^H \Lambda z$ .

$$A \text{ is PSD} \iff x^H A x \geq 0 \quad \forall x \in \mathbb{C}^m \iff z^H \Lambda z \geq 0 \quad \forall z \in \mathbb{C}^m.$$

$$z^H \Lambda z = \sum_{i=1}^m \lambda_i |z_i|^2 \geq 0 \quad \forall z \in \mathbb{C}^m \iff \lambda_i \geq 0 \quad \forall i = 1, \dots, m \quad \square$$

If  $A$  is PSD(PD),

- $Tr(A) \geq 0 (> 0)$

**Property 3:** Let  $A \in \mathbb{C}^{m \times k}$  and  $B \in \mathbb{C}^k$  and  $B$  has full row rank ( $rank(B) = k$ ). Then

$$Range(A) = Range(AB)$$

*Proof*

$Range(B) = \mathbb{C}^k$  as  $B$  is full row rank.

$$\begin{aligned} \text{Range}(AB) &= \{y|y = ABx, x \in \mathbb{C}^n\} \\ &= \{y|y = Az, \underbrace{z = Bx, x \in \mathbb{C}^n}_{z \in \text{Range}(B) = \mathbb{C}^k}\} \\ &= \{y|y = Az, z \in \mathbb{C}^k\} \\ &= \text{Range}(A) \end{aligned}$$



**Property 4:** Let  $B \in \mathbb{C}^{m \times k}$  and  $C \in \mathbb{C}^{m \times k}$  be two matrices with full column rank. Then

$$BB^H = CC^H$$

if and only if  $C = BQ$  for some unitary  $Q$ . *Proof*

( $\Leftarrow$ ): obvious.

( $\Rightarrow$ ):  $(C^H)^+ \triangleq C(C^H C)^{-1}$ . Since  $\text{rank}(C^H C) = \text{rank}(C)$ ,  $C^H C$  is invertible  $^\dagger$ .

$^\dagger$ : As  $C$  has full column rank, then  $Cx = 0 \iff x = 0$ .  
 $Cx = 0 \implies C^H Cx = 0 \implies x^H C^H Cx = 0 \iff \|Cx\|^2 = 0 \implies x = 0$ . Therefore,  $C^H Cx = 0 \implies x = 0$  and hence  $\text{rank}(C^H C) = \text{rank}(C)$ .

$$\begin{aligned} BB^H(C^H)^+ &= CC^H(C^H)^+ \\ &= CC^H C(C^H C)^{-1} \\ &= C \end{aligned}$$

$\implies$

$$\begin{aligned} C &= BB^H(C^H)^+ \\ &= \underbrace{BB^H C(C^H C)^{-1}}_{\triangleq Q(\text{square})} \\ Q^H Q &= \underbrace{(C^H C)^{-H} C^H}_{=I} \underbrace{BB^H C(C^H C)^{-1}}_{=I} \\ &= I \end{aligned}$$

Therefore,  $Q$  is unitary.  $\square$

## Matrix Inequality

For  $A, B \in H^{m \times m}$ , we write  $A \succeq (\succ) B$  if  $A - B$  is a PSD(PD) matrix.

So  $A \succeq 0$  if  $A$  is PSD.

**Basic Properties:**

- If  $A \succeq 0$ , then  $\alpha A \succeq 0 \forall \alpha \geq 0$
- If  $A \succeq 0, B \succeq 0$ , then  $\alpha A + B \succeq 0$
- If  $A \succeq B, B \succeq C$ , then  $\alpha A \succeq C$
- If  $A \not\succeq B$ , it doesn't imply  $B \succ A$

Let  $\lambda_1(A) \geq \lambda_2(A) \geq \dots \geq \lambda_m(A)$  for  $A \in H^{m \times m}$ , then

1.  $A \succeq I \iff \lambda_i(A) \geq 1 \forall i = 1, \dots, m$
2.  $I \succeq A \iff \lambda_i(A) \leq 1 \forall i = 1, \dots, m$
3. Suppose  $A, B \succ 0$ , then  $A \succeq B \iff B^{-1} \succeq A^{-1}$
4. If  $A \succeq B$ , then  $\lambda_i(A) \geq \lambda_i(B) \forall i = 1, \dots, m$

*Proof*

1)

$$\begin{aligned} A \succeq I &\iff A - I \succ 0 \iff V \Lambda V^H - I \succeq 0 \iff \\ V(\Lambda - I)V^H &\succeq 0 \iff \Lambda - I \succeq 0 \iff \lambda_i(A) \geq 1 \\ \forall i = 1, \dots, m \end{aligned}$$

2) similar to 1)

3)

$$\begin{aligned} A \succeq B &\iff A - B \succeq 0 \\ &\iff A^{\frac{1}{2}} A^{\frac{1}{2}} - B \succeq 0 \\ &\iff A^{\frac{1}{2}} (I - A^{-\frac{1}{2}} B A^{-\frac{1}{2}}) A^{\frac{1}{2}} \succeq 0 \\ &\iff I - A^{-\frac{1}{2}} B A^{-\frac{1}{2}} \succeq 0 \\ &\iff \lambda_i(A^{-\frac{1}{2}} B A^{-\frac{1}{2}}) \leq 1 \quad \forall i = 1, \dots, m \\ &\iff \frac{1}{\lambda_i(A^{\frac{1}{2}} B^{-1} A^{\frac{1}{2}})} \leq 1 \quad \forall i = 1, \dots, m^\dagger \\ &\iff \lambda_i(A^{\frac{1}{2}} B^{-1} A^{\frac{1}{2}}) \geq 1 \quad \forall i = 1, \dots, m \\ &\iff A^{\frac{1}{2}} B^{-1} A^{\frac{1}{2}} - I \succeq 0 \\ &\iff A^{-\frac{1}{2}} (I - A^{\frac{1}{2}} B^{-1} A^{\frac{1}{2}}) A^{-\frac{1}{2}} \succeq 0 \\ &\iff B^{-1} - A^{-1} \succeq 0 \\ &\iff B^{-1} \succeq A^{-1} \end{aligned}$$

$$^\dagger: \lambda_i(A^{-1}) = \frac{1}{\lambda_i(A)}$$

4)

$$\begin{aligned} \lambda_i(A) &= \lambda_i(A - B + B) \\ &\geq \lambda_i(B) + \underbrace{\lambda_m(A - B)}_{\geq 0 \text{ since } A - B \succeq 0} \\ &\geq \lambda_i(B) \quad \forall i = 1, \dots, m \end{aligned}$$

$\square$

From 4), if  $A \succeq B \implies \text{Tr}(A) \geq \text{Tr}(B)$ , if additionally  $A, B \succeq 0 \implies \det(A) \geq \det(B)$ , and if  $A, B \succ 0$  and  $A \succeq B$  then  $\text{Tr}(B^{-1}) \geq \text{Tr}(A^{-1})$

### Definition (Schur Complement)

Consider a matrix  $X \in H^{(m+n) \times (m+n)}$  to be partitioned

$$X = \begin{bmatrix} A & B \\ B^H C & \end{bmatrix}$$

where  $A \in H^{m \times m}, B \in \mathbb{C}^{m \times n}, C \in H^{n \times n} \succ 0$ . The term

$$S \triangleq A - B C^{-1} B^H \in H^{m \times m}$$

is called *Schur Complement* of  $X$ .

Then  $X \succeq (\succ) 0 \iff S \succeq (\succ) 0$  *Proof*

$$\text{Let } Y = \begin{bmatrix} I_m & 0 \\ -C^{-1} B^H & I_n \end{bmatrix} \in \mathbb{C}^{(m+n) \times (m+n)}.$$

Then  $Y$  is a lower triangular matrix and  $\det(Y) = 1 \implies Y$  is invertible.

Then

$$Y^H X Y = \begin{bmatrix} A - B C^{-1} B^H & 0 \\ 0 & C \end{bmatrix}$$

$$Y^H X Y \succeq (\succ) 0 \iff x \succeq (\succ) 0 \implies X \succeq (\succ) 0 \iff X \succeq (\succ) 0 \quad \square$$

$$\begin{aligned} \det(Y X^H Y) &= \det(X) \\ &= \det(S) \det(C) \end{aligned}$$

## Lecture 7. Singular Value Decomposition