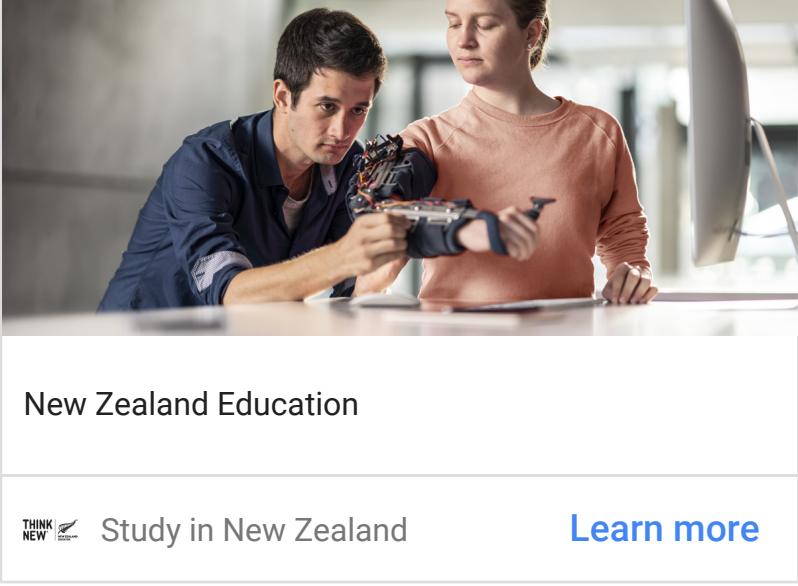


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LRU Cache Implementation

How to implement LRU caching scheme? What data structures should be used?

We are given total possible page numbers that can be referred. We are also given cache (or memory) size (Number of page frames that cache can hold at a time). The LRU caching scheme is to remove the least recently used frame when the cache is full and a new page is referenced which is not there in cache. Please see the Galvin book for more details (see the LRU page replacement slide [here](#)).

Recommended: Please solve it on "**PRACTICE**" first, before moving on to the solution.

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Got It !

We use two data structures to implement an LRU Cache.

1. **Queue** which is implemented using a doubly linked list. The maximum size of the queue will be equal to the total number of frames available (cache size). The most recently used pages will be near front end and least recently pages will be near the rear end.
2. **A Hash** with page number as key and address of the corresponding queue node as value.

When a page is referenced, the required page may be in the memory. If it is in the memory, we need to detach the node of the list and bring it to the front of the queue.

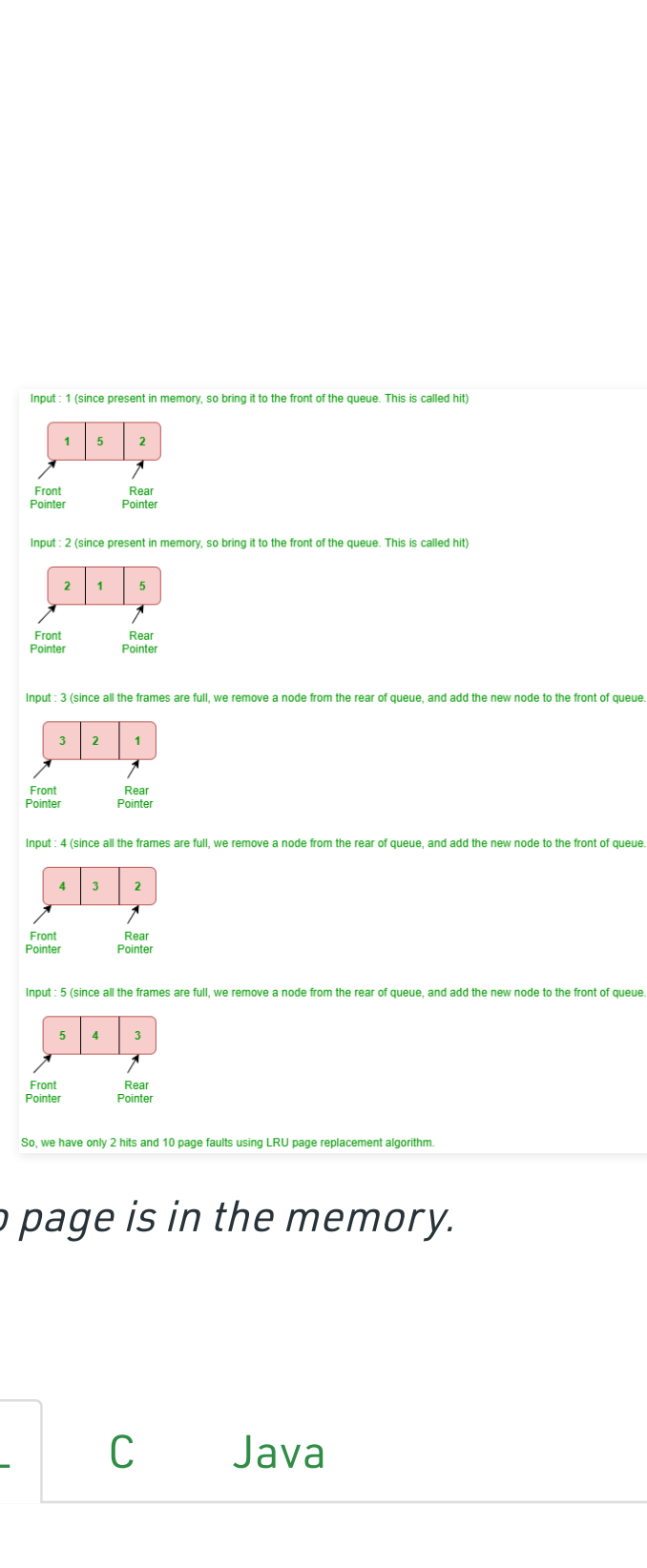
If the required page is not in memory, we bring that in memory. In simple words, we add a new node to the front of the queue and update the corresponding node address in the hash. If the queue is full, i.e. all the frames are full, we remove a node from the rear of the queue, and add the new node to the front of the queue.

Example – Consider the following reference string :

1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5

Find the number of page faults using least recently used (LRU) page replacement algorithm with 3 page frames.

Explanation –



Note: Initially no page is in the memory.

C++ using STL C Java

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```
// We can use stl container list as a double ended queue to store the cache keys, with the descending time of reference from front to back and a set container to check presence of a key. But to fetch the address of the key in the list using find(), it takes O(N) time. // This can be optimized by storing a reference (iterator) to each key in a hash map.
#include <bits/stdc++.h>
using namespace std;

class LRUCache {
    // store keys of cache
    list<int> dq;

    // store references of key in cache
    unordered_map<int, list<int>::iterator> ma;
    int csz; // maximum capacity of cache

public:
    LRUCache(int);
    void refer(int);
    void display();
};

// Declare the size
LRUCache::LRUCache(int n)
{
    csz = n;
}

// Refers key x with in the LRU cache
void LRUCache::refer(int x)
{
    // not present in cache
    if (ma.find(x) == ma.end()) {
        // cache is full
        if (dq.size() == csz) {
            // delete least recently used element
            int last = dq.back();

            // Pops the last element
            dq.pop_back();

            // Erase the last
            ma.erase(last);
        }

        // present in cache
        else
            dq.erase(ma[x]);

        // update reference
        dq.push_front(x);
        ma[x] = dq.begin();
    }

    // Function to display contents of cache
    void LRUCache::display()
    {
        // Iterate in the deque and print
        // all the elements in it
        for (auto it = dq.begin(); it != dq.end(); it++)
            cout << (*it) << " ";

        cout << endl;
    }

    // Driver Code
    int main()
    {
        LRUCache ca(4);

        ca.refer(1);
        ca.refer(2);
        ca.refer(3);
        ca.refer(1);
        ca.refer(4);
        ca.refer(5);
        ca.display();

        return 0;
    }
    // This code is contributed by Satish Srinivas
```

Output

5 4 1 3

Java Implementation using [LinkedHashMap](#).

The idea is to use a LinkedHashMap that maintains insertion order of elements. This way implementation becomes short and easy.

Java

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```
// Java program to implement LRU cache
// using LinkedHashMap
import java.util.*;

class LRUCache {

    Set<Integer> cache;
    int capacity;

    public LRUCache(int capacity)
    {
        this.cache = new LinkedHashMap<Integer>(capacity,
        this.capacity = capacity;
    }

    // This function returns false if key is not
    // present in cache. Else it moves the key to
    // front by first removing it and then adding
    // it, and returns true.
    public boolean get(int key)
    {
        if (!cache.contains(key))
            return false;
        cache.remove(key);
        cache.add(key);
        return true;
    }

    /* Refers key x with in the LRU cache */
    public void refer(int key)
    {
        if (get(key) == false)
            put(key);
    }

    // displays contents of cache in Reverse Order
    public void display()
    {
        LinkedList<Integer> list = new LinkedList<>(cache)

        // The descendingIterator() method of java.util.Li
        // class is used to return an iterator over the el
        // in this LinkedList in reverse sequential order
        Iterator<Integer> itr = list.descendingIterator();

        while (itr.hasNext())
            System.out.print(itr.next() + " ");
    }

    public void put(int key)
    {
        if (cache.size() == capacity) {
            int firstKey = cache.iterator().next();
            cache.remove(firstKey);
        }

        cache.add(key);
    }

    public static void main(String[] args)
    {
        LRUCache ca = new LRUCache(4);
        ca.refer(1);
        ca.refer(2);
        ca.refer(3);
        ca.refer(1);
        ca.refer(4);
        ca.refer(5);
        ca.display();
    }
}
```

Output

5 4 1 3

[Python implementation using \[OrderedDict\]\(#\)](#)

This article is compiled by [Aashish Barnwal](#) and reviewed by GeeksforGeeks team. Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.

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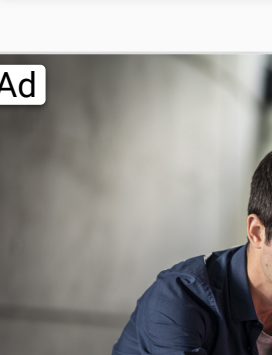
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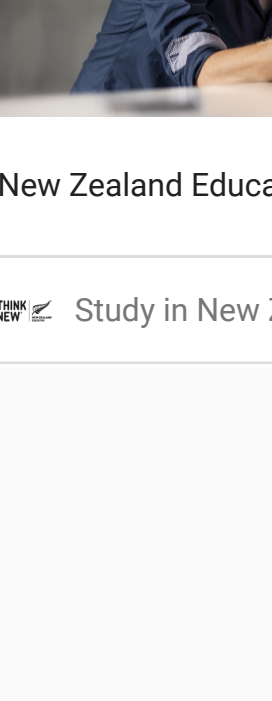
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