
MODULE *TPaxos*

Copied from my junior.
<https://github.com/Starydark/PaxosStore-tla/blob/master/specification/TPaxos.tla>

EXTENDS *Integers*, *FiniteSets*

$Max(m, n) \triangleq \text{IF } m > n \text{ THEN } m \text{ ELSE } n$
 $Injective(f) \triangleq \forall a, b \in \text{DOMAIN } f : (a \neq b) \Rightarrow (f[a] \neq f[b])$

CONSTANTS *Participant*, the set of participants
Value the set of possible input values for *Participant* to propose

$None \triangleq \text{CHOOSE } b : b \notin \text{Value}$

$Quorum \triangleq \{Q \in \text{SUBSET } Participant : \text{Cardinality}(Q) * 2 > \text{Cardinality}(Participant)\}$

ASSUME $QuorumAssumption \triangleq \wedge \forall Q \in Quorum : Q \subseteq Participant$
 $\wedge \forall Q1, Q2 \in Quorum : Q1 \cap Q2 \neq \{\}$

$Ballot \triangleq Nat$

$MaxBallot \triangleq \text{Cardinality}(Ballot) - 1$

$PIndex \triangleq \text{CHOOSE } f \in [Participant \rightarrow 1.. \text{Cardinality}(Participant)] : Injective(f)$
 $Bals(p) \triangleq \{b \in Ballot : b \% \text{Cardinality}(Participant) = PIndex[p] - 1\}$ allocate ballots for each $p \in Participant$

$State \triangleq [maxBal : Ballot \cup \{-1\}, maxVVal : Ballot \cup \{-1\}, maxVVal : Value \cup \{None\}]$

$InitState \triangleq [maxBal \mapsto -1, maxVVal \mapsto -1, maxVVal \mapsto None]$

For simplicity, in this specification, we choose to send the complete state of a participant each time. When receiving such a message, the participant processes only the “partial” state it needs.

$Message \triangleq [from : Participant, to : \text{SUBSET } Participant, state : [Participant \rightarrow State]]$

VARIABLES *state*, $state[p][q]$: the state of $q \in Participant$ from the view of $p \in Participant$
msgs the set of messages that have been sent

$vars \triangleq \langle state, msgs \rangle$

$TypeOK \triangleq \wedge state \in [Participant \rightarrow [Participant \rightarrow State]]$
 $\wedge msgs \subseteq Message$

$Send(m) \triangleq msgs' = msgs \cup \{m\}$

$Init \triangleq \wedge state = [p \in Participant \mapsto [q \in Participant \mapsto InitState]]$
 $\wedge msgs = \{\}$

$p \in Participant$ starts the prepare phase by issuing a ballot $b \in Ballot$.

$Prepare(p, b) \triangleq \wedge b \in Bals(p)$
 $\wedge state[p][p].maxBal < b$

$$\begin{aligned} & \wedge state' = [state \text{ EXCEPT } ![p][p].maxBal = b] \\ & \wedge Send([from \mapsto p, to \mapsto Participant \setminus \{p\}, state \mapsto state'[p]]) \end{aligned}$$

$q \in Participant$ updates its own state $state[q]$ according to the actual state pp of $p \in Participant$ extracted from a message $m \in Message$ it receives. This is called by $OnMessage(q)$. Note: pp is $m.state[p]$; it may not be equal to $state[p][p]$ at the time $UpdateState$ is called.

$$\begin{aligned} UpdateState(q, p, pp) \triangleq & \text{ LET } maxB \triangleq Max(state[q][q].maxBal, pp.maxBal) \\ & \text{ IN } state' = [state \text{ EXCEPT } ![q][p].maxBal = Max(@, pp.maxBal), \\ & \quad ![q][p].maxVVal = Max(@, pp.maxVVal), \\ & \quad ![q][p].maxVVal = \text{ IF } state[q][p].maxVVal < pp.maxVVal \\ & \quad \quad \text{ THEN } pp.maxVVal \text{ ELSE } @, \\ & \quad ![q][q].maxBal = maxB, \quad \text{make prom} \\ & \quad ![q][q].maxVVal = \text{ IF } maxB \leq pp.maxVVal \quad \text{accept} \\ & \quad \quad \text{ THEN } pp.maxVVal \text{ ELSE } @, \\ & \quad ![q][q].maxVVal = \text{ IF } maxB \leq pp.maxVVal \quad \text{accept} \\ & \quad \quad \text{ THEN } pp.maxVVal \text{ ELSE } @] \end{aligned}$$

$q \in Participant$ receives and processes a message in $Message$.

$$\begin{aligned} OnMessage(q) \triangleq & \exists m \in msgs : \\ & \wedge q \in m.to \\ & \wedge \text{ LET } p \triangleq m.from \\ & \quad \text{ IN } UpdateState(q, p, m.state[p]) \\ & \wedge \text{ LET } qm \triangleq [from \mapsto m.from, to \mapsto m.to \setminus \{q\}, state \mapsto m.state] \quad \text{remove } q \text{ from to} \\ & \quad nm \triangleq [from \mapsto q, to \mapsto \{m.from\}, state \mapsto state'[q]] \quad \text{new message to reply} \\ & \quad \text{ IN } \text{ IF } \vee m.state[q].maxBal < state'[q][q].maxBal \\ & \quad \quad \vee m.state[q].maxVVal < state'[q][q].maxVVal \\ & \quad \quad \text{ THEN } msgs' = (msgs \setminus \{m\}) \cup \{qm, nm\} \\ & \quad \quad \text{ ELSE } msgs' = (msgs \setminus \{m\}) \cup \{qm\} \end{aligned}$$

$$\begin{aligned} Accept(p, b, v) \triangleq & \wedge b \in Bals(p) \\ & \wedge state[p][p].maxBal \leq b \quad \text{corresponding the first conjunction in Voting} \\ & \wedge state[p][p].maxVVal \neq b \quad \text{cooresponding the first conjunction in Voting} \\ & \wedge \exists Q \in Quorum : \\ & \quad \wedge \forall q \in Q : state[p][q].maxBal = b \\ & \wedge \forall q \in Participant : state[p][q].maxVVal = -1 \\ & \quad \vee \exists q \in Participant : \\ & \quad \quad \wedge state[p][q].maxVVal = v \\ & \quad \quad \wedge \forall r \in Participant : state[p][q].maxVVal \geq state[p][r].maxVVal \\ & \wedge state' = [state \text{ EXCEPT } ![p][p].maxVVal = b, \\ & \quad \quad \quad ![p][p].maxVVal = v] \\ & \wedge Send([from \mapsto p, to \mapsto Participant \setminus \{p\}, state \mapsto state'[p]]) \end{aligned}$$

$$\begin{aligned} Next \triangleq & \exists p \in Participant : \vee OnMessage(p) \\ & \quad \vee \exists b \in Ballot : \vee Prepare(p, b) \\ & \quad \quad \vee \exists v \in Value : Accept(p, b, v) \end{aligned}$$

$$Spec \triangleq Init \wedge \Box[Next]_{vars}$$

$$\begin{aligned}
ChosenP(p) &\triangleq \{v \in Value : \exists b \in Ballot : \\
&\quad \exists Q \in Quorum : \forall q \in Q : \wedge state[p][q].maxVVal = b \\
&\quad \wedge state[p][q].maxVVal = v\} \\
chosen &\triangleq \text{UNION } \{ChosenP(p) : p \in Participant\} \\
Consistency &\triangleq Cardinality(chosen) \leq 1
\end{aligned}$$

THEOREM $Spec \Rightarrow \Box Consistency$

For checking *Liveness* $WF(A)$: if A ever becomes enabled, then an A step will eventually occur even if A remains enabled for only a fraction of a nanosecond and is never again enabled. Liveness in *TPaxos*: like paxos, there should be a single-leader to prapre and accept.

$$\begin{aligned}
LConstrain &\triangleq \wedge \exists p \in Participant : \\
&\quad \wedge MaxBallot \in Bals(p) \\
&\quad \wedge WF_{vars}(Prepare(p, MaxBallot)) \\
&\quad \wedge \forall v \in Value : WF_{vars}(Accept(p, MaxBallot, v)) \\
&\quad \wedge \exists Q \in Quorum : \\
&\quad \quad \wedge p \in Q \\
&\quad \quad \wedge \forall q \in Q : WF_{vars}(OnMessage(q))
\end{aligned}$$

$$LSpec \triangleq Spec \wedge LConstrain$$

$$Liveness \triangleq \Diamond(chosen \neq \{\})$$

\ * Modification History
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