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1  |----- MODULE XJupiter -----|
   | Specification of the Jupiter protocol described in CSCW'2014 by Yi Xu, Chengzheng Sun, and |
   | Mo Li. We call it XJupiter, with 'X' for "Xu". |
7  | EXTENDS StateSpace |
8  |-----|
9  | VARIABLES |
10 |   c2ss,   c2ss[c]: the 2D state space (2ss, for short) at client c ∈ Client |
11 |   s2ss    s2ss[c]: the 2D state space maintained by the Server for client c ∈ Client |
13 | vars ≜ ⟨intVars, ctxVars, c2ss, s2ss⟩ |
14 |-----|
15 | TypeOK ≜ |
16 |   ∧ TypeOKInt |
17 |   ∧ TypeOKCtx |
18 |   ∧ Comm(Cop)! TypeOK |
19 |   ∧ ∀ c ∈ Client : IsSS(c2ss[c]) ∧ IsSS(s2ss[c]) |
20 |-----|
21 | Init ≜ |
22 |   ∧ InitInt |
23 |   ∧ InitCtx |
24 |   ∧ Comm(Cop)! Init |
25 |   ∧ c2ss = [c ∈ Client ↦ EmptySS] |
26 |   ∧ s2ss = [c ∈ Client ↦ EmptySS] |
27 |-----|
   | xForm: iteratively transform cop with a path through the 2D state space ss at some client. |
32 | xForm(cop, ss, cur) ≜ |
33 |   LET u ≜ Locate(cop, ss) |
34 |   v ≜ u ∪ {cop.oid} |
35 |   RECURSIVE xFormHelper(−, −, −, −) |
36 |   xFormHelper(uh, vh, coph, xss) ≜ xss: eXtra ss created during transformation |
37 |   IF uh = cur THEN [xss ↦ xss, xcop ↦ coph] |
38 |   ELSE LET e ≜ CHOOSE e ∈ ss.edge : e.from = uh ∧ ClientOf(e.cop) ≠ ClientOf(cop) |
39 |       copprime ≜ e.cop |
40 |       uprime ≜ e.to |
41 |       vprime ≜ vh ∪ {copprime.oid} |
42 |       coph2copprime ≜ COT(coph, copprime) |
43 |       copprime2coph ≜ COT(copprime, coph) |
44 |   IN xFormHelper(uprime, vprime, coph2copprime, |
45 |       xss ⊕ [node ↦ {vprime}, |
46 |       edge ↦ {[from ↦ vh, to ↦ vprime, cop ↦ copprime2coph], |
47 |       [from ↦ uprime, to ↦ vprime, cop ↦ coph2copprime]}]) |
48 |   IN xFormHelper(u, v, cop, [node ↦ {v}, edge ↦ {[from ↦ u, to ↦ v, cop ↦ cop]}]) |
49 |-----|
   | Client c ∈ Client perform operation cop. |

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53  $ClientPerform(cop, c) \triangleq$ 
54   LET  $xform \triangleq xForm(cop, c2ss[c], ds[c])$   $xform: [xss, xcop]$ 
55   IN    $\wedge c2ss' = [c2ss \text{ EXCEPT } ![c] = @ \oplus xform.xss]$ 
56        $\wedge state' = [state \text{ EXCEPT } ![c] = Apply(xform.xcop.op, @)]$ 
    Client  $c \in Client$  generates an operation  $op$ .
60  $DoOp(c, op) \triangleq$ 
61   LET  $cop \triangleq [op \mapsto op, oid \mapsto [c \mapsto c, seq \mapsto cseq'[c]], ctx \mapsto ds[c]]$ 
62   IN    $\wedge ClientPerform(cop, c)$ 
63        $\wedge Comm(Cop)!CSend(cop)$ 
65  $DoIns(c) \triangleq$ 
66    $\exists ins \in \{op \in Ins : op.pos \in 1 \dots (Len(state[c]) + 1) \wedge op.ch \in chins \wedge op.pr = Priority[c]\} :$ 
67      $\wedge DoOp(c, ins)$ 
68      $\wedge chins' = chins \setminus \{ins.ch\}$ 
70  $DoDel(c) \triangleq$ 
71    $\exists del \in \{op \in Del : op.pos \in 1 \dots Len(state[c])\} :$ 
72      $\wedge DoOp(c, del)$ 
73      $\wedge \text{UNCHANGED } chins$ 
75  $Do(c) \triangleq$ 
76    $\wedge DoCtx(c)$ 
77    $\wedge \vee DoIns(c)$ 
78    $\vee DoDel(c)$ 
79    $\wedge \text{UNCHANGED } s2ss$ 
    Client  $c \in Client$  receives a message from the Server.
83  $Rev(c) \triangleq$ 
84    $\wedge Comm(Cop)!CRev(c)$ 
85    $\wedge \text{LET } cop \triangleq Head(cincoming[c])$ 
86   IN    $ClientPerform(cop, c)$ 
87    $\wedge RevCtx(c)$ 
88    $\wedge \text{UNCHANGED } \langle chins, s2ss \rangle$ 
89 |-----|
    The Server performs operation  $cop$ .
93  $ServerPerform(cop) \triangleq$ 
94   LET  $c \triangleq ClientOf(cop)$ 
95    $scur \triangleq ds[Server]$ 
96    $xform \triangleq xForm(cop, s2ss[c], scur)$   $xform: [xss, xcop]$ 
97    $xcop \triangleq xform.xcop$ 
98    $xcur \triangleq scur \cup \{cop.oid\}$ 
99   IN    $\wedge s2ss' = [cl \in Client \mapsto$ 
100        IF  $cl = c$ 
101        THEN  $s2ss[cl] \oplus xform.xss$ 
102        ELSE  $s2ss[cl] \oplus [node \mapsto \{xcur\},$ 
103             $edge \mapsto \{[from \mapsto scur, to \mapsto xcur, cop \mapsto xcop]\}]$ 

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104         ]
105          $\wedge state' = [state \text{ EXCEPT } ![Server] = Apply(xcop.op, @)]$ 
106          $\wedge Comm(Cop)!SSendSame(c, xcop)$ 
107         The Server receives a message.
108
109      $SRev \triangleq$ 
110          $\wedge Comm(Cop)!SRev$ 
111          $\wedge LET \ cop \triangleq Head(sincoming)$ 
112         IN  $ServerPerform(cop)$ 
113          $\wedge SRevCtx$ 
114          $\wedge UNCHANGED \langle chins, c2ss \rangle$ 
115
116 |-----|
117      $Next \triangleq$ 
118          $\vee \exists c \in Client : Do(c) \vee Rev(c)$ 
119          $\vee SRev$ 
120
121      $Fairness \triangleq$  There is no requirement that the clients ever generate operations.
122      $WF_{vars}(SRev \vee \exists c \in Client : Rev(c))$ 
123
124      $Spec \triangleq Init \wedge \Box[Next]_{vars} \wedge Fairness$ 
125 |-----|
126      $CSSync \triangleq$  Each client  $c \in Client$  is synchronized with the Server.
127      $\forall c \in Client : (ds[c] = ds[Server]) \Rightarrow c2ss[c] = s2ss[c]$ 
128
129 THEOREM  $Spec \Rightarrow \Box CSSync$ 
130 |-----|
131
132 \ * Modification History
133 \ * Last modified Sat Dec 29 18:50:10 CST 2018 by hengxin
134 \ * Created Tue Oct 09 16:33:18 CST 2018 by hengxin

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