
MODULE *Peterson*

EXTENDS *Integers*, *TLAPS*

$Not(i) \triangleq \text{IF } i = 0 \text{ THEN } 1 \text{ ELSE } 0$

```
--algorithm Peterson{
  variables flag = [i ∈ {0, 1} ↦ FALSE], turn = 0;
  process ( proc ∈ {0, 1} ) {
    a0: while ( TRUE ) {
      a1:   flag[self] := TRUE;
      a2:   turn := Not(self);
      a3a:  if ( flag[Not(self)] ) { goto a3b } else { goto cs };
      a3b:  if ( turn = Not(self) ) { goto a3a } else { goto cs };
      cs:   skip;
      a4:   flag[self] := FALSE;
    }
  }
}
```

BEGIN TRANSLATION

VARIABLES *flag*, *turn*, *pc*

$vars \triangleq \langle flag, turn, pc \rangle$

$ProcSet \triangleq (\{0, 1\})$

$Init \triangleq$ Global variables
 $\wedge flag = [i \in \{0, 1\} \mapsto FALSE]$
 $\wedge turn = 0$
 $\wedge pc = [self \in ProcSet \mapsto "a0"]$

$a0(self) \triangleq$ $\wedge pc[self] = "a0"$
 $\wedge pc' = [pc \text{ EXCEPT } ![self] = "a1"]$
 $\wedge \text{UNCHANGED } \langle flag, turn \rangle$

$a1(self) \triangleq$ $\wedge pc[self] = "a1"$
 $\wedge flag' = [flag \text{ EXCEPT } ![self] = TRUE]$
 $\wedge pc' = [pc \text{ EXCEPT } ![self] = "a2"]$
 $\wedge turn' = turn$

$a2(self) \triangleq$ $\wedge pc[self] = "a2"$
 $\wedge turn' = Not(self)$
 $\wedge pc' = [pc \text{ EXCEPT } ![self] = "a3a"]$
 $\wedge flag' = flag$

$a3a(self) \triangleq$ $\wedge pc[self] = "a3a"$

$$\begin{aligned}
& \wedge \text{IF } flag[Not(self)] \\
& \quad \text{THEN } \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"a3b"}] \\
& \quad \text{ELSE } \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"cs"}] \\
& \wedge \text{UNCHANGED } \langle flag, turn \rangle \\
a3b(self) & \triangleq \wedge pc[self] = \text{"a3b"} \\
& \wedge \text{IF } turn = Not(self) \\
& \quad \text{THEN } \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"a3a"}] \\
& \quad \text{ELSE } \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"cs"}] \\
& \wedge \text{UNCHANGED } \langle flag, turn \rangle \\
cs(self) & \triangleq \wedge pc[self] = \text{"cs"} \\
& \wedge \text{TRUE} \\
& \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"a4"}] \\
& \wedge \text{UNCHANGED } \langle flag, turn \rangle \\
a4(self) & \triangleq \wedge pc[self] = \text{"a4"} \\
& \wedge flag' = [flag \text{ EXCEPT } ![self] = \text{FALSE}] \\
& \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"a0"}] \\
& \wedge turn' = turn \\
proc(self) & \triangleq a0(self) \vee a1(self) \vee a2(self) \vee a3a(self) \vee a3b(self) \\
& \vee cs(self) \vee a4(self) \\
Next & \triangleq (\exists self \in \{0, 1\} : proc(self)) \\
Spec & \triangleq Init \wedge \Box [Next]_{vars} \\
\text{END TRANSLATION} & \\
MutualExclusion & \triangleq (pc[0] \neq \text{"cs"}) \vee (pc[1] \neq \text{"cs"}) \\
TypeOK & \triangleq \wedge pc \in [\{0, 1\} \rightarrow \{\text{"a0"}, \text{"a1"}, \text{"a2"}, \text{"a3a"}, \text{"a3b"}, \text{"cs"}, \text{"a4"}\}] \\
& \wedge turn \in \{0, 1\} \\
& \wedge flag \in [\{0, 1\} \rightarrow \text{BOOLEAN}] \\
I & \triangleq \forall i \in \{0, 1\} : \\
& \quad \wedge pc[i] \in \{\text{"a2"}, \text{"a3a"}, \text{"a3b"}, \text{"cs"}, \text{"a4"}\} \Rightarrow flag[i] \\
& \quad \wedge pc[i] \in \{\text{"cs"}, \text{"a4"}\} \Rightarrow \wedge pc[Not(i)] \notin \{\text{"cs"}, \text{"a4"}\} \\
& \quad \wedge pc[Not(i)] \in \{\text{"a3a"}, \text{"a3b"}\} \Rightarrow turn = i \\
Inv & \triangleq TypeOK \wedge I
\end{aligned}$$

THEOREM $Spec \Rightarrow \Box MutualExclusion$

$\langle 1 \rangle 1. Init \Rightarrow Inv$

BY *Zenon*, *Isa*DEFS *Init*, *Inv*, *TypeOK*, *I*, *ProcSet*

$\langle 1 \rangle 2. Inv \wedge [Next]_{vars} \Rightarrow Inv'$

BY DEFS *Inv*, *Next*, *TypeOK*, *I*, *Not*, *proc*, *a0*, *a1*, *a2*, *a3a*, *a3b*, *cs*, *a4*, *vars*

$\langle 2 \rangle 1. \text{SUFFICES ASSUME } Inv, Next$

PROVE Inv'

BY *Zenon*, *Isa*DEFS *Inv*, *TypeOK*, *I*, *vars*

$\langle 2 \rangle 2. TypeOK'$

BY *Zenon*, *Isa*, $\langle 2 \rangle 1$

DEFS *Inv*, *TypeOK*, *Next*, *proc*, *a0*, *a1*, *a2*, *a3a*, *a3b*, *cs*, *a4*, *Not*

$\langle 2 \rangle 3. I'$

$\langle 3 \rangle 1. \text{SUFFICES ASSUME NEW } j \in \{0, 1\}$

PROVE $I!(j)'$

BY *Zenon*, *Isa*DEFS *I*

$\langle 3 \rangle 2. \text{PICK } i \in \{0, 1\} : proc(i)$

BY *Zenon*, *Isa*, $\langle 2 \rangle 1$

DEFS *Next*, *I*, *TypeOK*, *Next*, *proc*, *a0*, *a1*, *a2*, *a3a*, *a3b*, *cs*, *a4*, *Not*

$\langle 3 \rangle 3. \text{CASE } i = j$

BY *Zenon*, *Isa*, $\langle 2 \rangle 1$, $\langle 3 \rangle 2$, $\langle 3 \rangle 3$

DEFS *Inv*, *I*, *TypeOK*, *proc*, *a0*, *a1*, *a2*, *a3a*, *a3b*, *cs*, *a4*, *Not*

$\langle 3 \rangle 4. \text{CASE } i \neq j$

BY *Zenon*, *Isa*, $\langle 2 \rangle 1$, $\langle 3 \rangle 2$, $\langle 3 \rangle 4$

DEFS *Inv*, *I*, *TypeOK*, *proc*, *a0*, *a1*, *a2*, *a3a*, *a3b*, *cs*, *a4*, *Not*

$\langle 3 \rangle 5. \text{QED}$

BY *Zenon*, *Isa*, $\langle 3 \rangle 3$, $\langle 3 \rangle 4$

$\langle 2 \rangle 4. \text{QED}$

BY *Zenon*, *Isa*, $\langle 2 \rangle 2$, $\langle 2 \rangle 3$ DEFS *Inv*

$\langle 1 \rangle 3. Inv \Rightarrow MutualExclusion$

BY *Zenon*, *Isa*DEFS *MutualExclusion*, *Inv*, *I*, *Not*

$\langle 1 \rangle 4. \text{QED}$

BY *Zenon*, *Isa*, $\langle 1 \rangle 1$, $\langle 1 \rangle 2$, $\langle 1 \rangle 3$, *PTL* DEF *Spec*

\ * Modification History

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