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1  ┌────────────────── MODULE XJupiter ───────────────────┐
    Specification of the Jupiter protocol described in CSCW'2014 by Yi Xu, Chengzheng Sun, and
    Mo Li. We call it XJupiter, with 'X' for "Xu".
7  EXTENDS JupiterInterface
8  ┌──────────────────┐
    Direction flags for edges in 2D state spaces and OT.
12 Local  $\triangleq$  0
13 Remote  $\triangleq$  1
14 ┌──────────────────┐
    Cop: operation of type Op with context
18 Oid  $\triangleq$  [c : Client, seq : Nat] operation identifier
19 Cop  $\triangleq$  [op : Op  $\cup$  {Nop}, oid : Oid, ctx : SUBSET Oid]

    OT of two operations of type Cop.
24 COT(lcop, rcop)  $\triangleq$  [lcop EXCEPT !.op = Xform(lcop.op, rcop.op), !.ctx = @  $\cup$  {rcop.oid}]
25 ┌──────────────────┐
26 VARIABLES
27   cseq,      cseq[c]: local sequence number at client c  $\in$  Client
    The 2D state spaces (ss, for short). Each client maintains one 2D state space. The server
    maintains n 2D state spaces, one for each client.
33   c2ss,      c2ss[c]: the 2D state space at client c  $\in$  Client
34   s2ss,      s2ss[c]: the 2D state space maintained by the Server for client c  $\in$  Client
35   cur       cur[r]: the current node of the 2D state space at replica r  $\in$  Replica
37   vars  $\triangleq$  {chins, cseq, cur, cincoming, sincoming, c2ss, s2ss, state}
38 ┌──────────────────┐
    A 2D state space is a directed graph with labeled edges. It is represented by a record with node
    field and edge field. Each node is characterized by its context, a set of operations. Each edge is
    labeled with an operation and a direction flag indicating whether this edge is LOCAL or REMOTE.
    For clarity, we denote edges by records instead of tuples.
47 IsSS(G)  $\triangleq$ 
48    $\wedge$  G = [node  $\mapsto$  G.node, edge  $\mapsto$  G.edge]
49    $\wedge$  G.node  $\subseteq$  (SUBSET Oid)
50    $\wedge$  G.edge  $\subseteq$  [from : G.node, to : G.node, cop : Cop, lr : {Local, Remote}]
52 EmptySS  $\triangleq$  [node  $\mapsto$  {}, edge  $\mapsto$  {}]
    Take union of two state spaces ss1 and ss2.
56 ss1  $\oplus$  ss2  $\triangleq$  [node  $\mapsto$  ss1.node  $\cup$  ss2.node, edge  $\mapsto$  ss1.edge  $\cup$  ss2.edge]
58 TypeOK  $\triangleq$ 
59    $\wedge$  TypeOKInt
60    $\wedge$  cseq  $\in$  [Client  $\rightarrow$  Nat]
    For the 2D state spaces:
64    $\wedge \forall c \in \text{Client} : \text{IsSS}(c2ss[c]) \wedge \text{IsSS}(s2ss[c])$ 
65    $\wedge cur \in [Replica \rightarrow \text{SUBSET } Oid]$ 

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69    $\wedge \text{Comm}(\text{Cop})! \text{TypeOK}$ 
70 |-----|
71    $\text{Init} \triangleq$ 
72      $\wedge \text{InitInt}$ 
73      $\wedge \text{cseq} = [c \in \text{Client} \mapsto 0]$ 
74   For the 2D state spaces:
75    $\wedge \text{c2ss} = [c \in \text{Client} \mapsto \text{EmptySS}]$ 
76    $\wedge \text{s2ss} = [c \in \text{Client} \mapsto \text{EmptySS}]$ 
77    $\wedge \text{cur} = [r \in \text{Replica} \mapsto \{\}]$ 
78   For communication between the server and the clients:
79    $\wedge \text{Comm}(\text{Cop})! \text{Init}$ 
80 |-----|
81   Locate the node in the 2D state space  $ss$  which matches the context  $ctx$  of  $\text{cop}$ .
82    $\text{Locate}(\text{cop}, ss) \triangleq \text{CHOOSE } n \in ss.\text{node} : n = \text{cop}.ctx$ 
83    $x\text{Form}$ : iteratively transform  $\text{cop}$  with a path through the 2D state space  $ss$  at some client,
84   following the edges with the direction flag  $d$ .
85    $x\text{Form}(\text{cop}, ss, \text{current}, d) \triangleq$ 
86     LET  $u \triangleq \text{Locate}(\text{cop}, ss)$ 
87      $v \triangleq u \cup \{\text{cop}.oid\}$ 
88     RECURSIVE  $x\text{FormHelper}(-, -, -, -, -, -)$ 
89     'h' stands for "helper";  $xss$ :  $eXtra$   $ss$  created during transformation
90      $x\text{FormHelper}(uh, vh, coph, xss, xcoph, xcurh) \triangleq$ 
91     IF  $uh = \text{current}$ 
92     THEN  $\langle xss, xcoph, xcurh \rangle$ 
93     ELSE LET  $e \triangleq \text{CHOOSE } e \in ss.\text{edge} : e.\text{from} = uh \wedge e.\text{lr} = d$ 
94      $u\text{prime} \triangleq e.\text{to}$ 
95      $cop\text{prime} \triangleq e.\text{cop}$ 
96      $coph2cop\text{prime} \triangleq \text{COT}(coph, cop\text{prime})$ 
97      $cop\text{prime}2coph \triangleq \text{COT}(cop\text{prime}, coph)$ 
98      $v\text{prime} \triangleq vh \cup \{cop\text{prime}.oid\}$ 
99     IN  $x\text{FormHelper}(u\text{prime}, v\text{prime}, coph2cop\text{prime},$ 
100        $[node \mapsto xss.\text{node} \cup \{v\text{prime}\},$ 
101        $edge \mapsto xss.\text{edge} \cup \{[from \mapsto vh, to \mapsto v\text{prime}, cop \mapsto cop\text{prime}2coph, lr \mapsto d],$ 
102        $[from \mapsto u\text{prime}, to \mapsto v\text{prime}, cop \mapsto coph2cop\text{prime}, lr \mapsto 1 - d]\},$ 
103        $coph2cop\text{prime}, v\text{prime})$ 
104     IN  $x\text{FormHelper}(u, v, cop, [node \mapsto \{v\}, edge \mapsto \{[from \mapsto u, to \mapsto v, cop \mapsto cop, lr \mapsto 1 - d]\}], cop, v)$ 
105 |-----|
106   Client  $c \in \text{Client}$  perform operation  $\text{cop}$  guided by the direction flag  $d$ .
107    $\text{ClientPerform}(\text{cop}, c, d) \triangleq$ 
108     LET  $x\text{form} \triangleq x\text{Form}(\text{cop}, \text{c2ss}[c], \text{cur}[c], d)$   $x\text{form}: \langle xss, xcop, xcur \rangle$ 
109      $xss \triangleq x\text{form}[1]$ 
110      $xcop \triangleq x\text{form}[2]$ 

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122       $xcur \triangleq xform[3]$ 
123      IN  $\wedge c2ss' = [c2ss \text{ EXCEPT } ![c] = @ \oplus xss]$ 
124       $\wedge cur' = [cur \text{ EXCEPT } ![c] = xcur]$ 
125       $\wedge state' = [state \text{ EXCEPT } ![c] = Apply(xcop.op, @)]$ 
    Client  $c \in Client$  generates an operation  $op$ .
129   $DoOp(c, op) \triangleq$ 
130       $\wedge cseq' = [cseq \text{ EXCEPT } ![c] = @ + 1]$ 
131       $\wedge \text{LET } cop \triangleq [op \mapsto op, oid \mapsto [c \mapsto c, seq \mapsto cseq'[c]], ctx \mapsto cur[c]]$ 
132      IN  $\wedge ClientPerform(cop, c, Remote)$ 
133       $\wedge Comm(Cop)!CSend(cop)$ 

135   $DoIns(c) \triangleq$ 
136       $\exists ins \in \{op \in Ins : op.pos \in 1 \dots (Len(state[c]) + 1) \wedge op.ch \in chins \wedge op.pr = Priority[c]\} :$ 
137       $\wedge DoOp(c, ins)$ 
138       $\wedge chins' = chins \setminus \{ins.ch\}$  We assume that all inserted elements are unique.

140   $DoDel(c) \triangleq$ 
141       $\exists del \in \{op \in Del : op.pos \in 1 \dots Len(state[c])\} :$ 
142       $\wedge DoOp(c, del)$ 
143       $\wedge \text{UNCHANGED } \langle chins \rangle$ 

145   $Do(c) \triangleq$ 
146       $\wedge \vee DoIns(c)$ 
147       $\vee DoDel(c)$ 
148       $\wedge \text{UNCHANGED } \langle s2ss \rangle$ 
    Client  $c \in Client$  receives a message from the Server.

152   $Rev(c) \triangleq$ 
153       $\wedge Comm(Cop)!CRev(c)$ 
154       $\wedge \text{LET } cop \triangleq Head(cincomig[c])$  the received (transformed) operation
155      IN  $ClientPerform(cop, c, Local)$ 
156       $\wedge \text{UNCHANGED } \langle chins, cseq, s2ss \rangle$ 
157  |-----|
    The Server performs operation  $cop$ .

161   $ServerPerform(cop) \triangleq$ 
162      LET  $c \triangleq cop.oid.c$ 
163       $scur \triangleq cur[Server]$ 
164       $xform \triangleq xForm(cop, s2ss[c], scur, Remote)$   $xform: \langle xss, xcop, xcur \rangle$ 
165       $xss \triangleq xform[1]$ 
166       $xcop \triangleq xform[2]$ 
167       $xcur \triangleq xform[3]$ 
168      IN  $\wedge s2ss' = [cl \in Client \mapsto$ 
169          IF  $cl = c$ 
170          THEN  $s2ss[cl] \oplus xss$ 
171          ELSE  $s2ss[cl] \oplus [node \mapsto \{xcur\},$ 
172               $edge \mapsto \{[from \mapsto scur, to \mapsto xcur,$ 

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173                                      $cop \mapsto xcop, lr \mapsto Remote\}}]$ 
174                                     ]
175          $\wedge cur' = [cur \text{ EXCEPT } ![Server] = xcur]$ 
176          $\wedge state' = [state \text{ EXCEPT } ![Server] = Apply(xcop.op, @)]$ 
177          $\wedge Comm(Cop)!SSendSame(c, xcop)$  broadcast the transformed operation
178         The Server receives a message.
179
180  $SRev \triangleq$ 
181      $\wedge Comm(Cop)!SRev$ 
182      $\wedge \text{LET } cop \triangleq Head(sincoming)$ 
183     IN  $ServerPerform(cop)$ 
184      $\wedge \text{UNCHANGED } \langle chins, cseq, c2ss \rangle$ 
185
186 |-----|
187  $Next \triangleq$ 
188      $\vee \exists c \in Client : Do(c) \vee Rev(c)$ 
189      $\vee SRev$ 
190
191  $Fairness \triangleq$ 
192      $WF_{vars}(SRev \vee \exists c \in Client : Rev(c))$ 
193
194  $Spec \triangleq Init \wedge \Box[Next]_{vars} \wedge Fairness$ 
195 |-----|
196 In Jupiter (not limited to XJupiter), each client synchronizes with the server. In XJupiter, this
197 is expressed as the following CSSync property.
198
199  $CSSync \triangleq$ 
200      $\forall c \in Client : (cur[c] = cur[Server]) \Rightarrow c2ss[c] = s2ss[c]$ 
201
202 |-----|
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204 \ * Last modified Tue Dec 04 21:12:08 CST 2018 by hengxin
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