MODULE *Stuttering*

This module is explained in Section 5 of the paper "Auxiliary Variables in TLA+". It defines operators used to add a stuttering variable s to a specification *Spec* to form a specification *SpecS*. It is mean to be instantiated with s replaced by the stuttering variable to be added and *vars* replaced by the tuple of all variables in the original specification.

If *Init* is the initial predicate of *Spec* , then the initial predicate of *SpecS* is $Init \land (s = top)$, where top is defined below.

The next-state action of *SpecS* is obtained by replacing each subaction A of a disjunctive representation of the next-state action Next of *Spec* with an action AS written in terms of operators defined below. (Disjunctive representations are described in Section 3.2 of "Auxiliary Variables in TLA+".) Action AS executes A and stuttering steps added either before or after an A step. The basic idea is that s equals top except while stuttering steps are being taken, when it is a record with the following fields:

id: A value that identifies the action A .

ctxt: A value identifying the context under which A is executed. For example, if A appears in a formula $\exists\, i, j \in S : A$, this would equal the value of the pair $\langle i, j \rangle$ for which A is being executed.

val: A value that is decremented by each stuttering step, until it reaches a minimum value.

The arguments of the operators defined in this module have the following meanings.

A
The subaction A of Next.

id
A string identifying action A .

Sigma
A set of values ordered by some "less-than" relation. This is the set of possible values of $s.val$ when executing stuttering steps before or after subaction A .

bot
The minimum element of *Sigma* under its less-than relation.

initVal
The initial value to which $s.val$ is set for executing stuttering steps before or after A .

decr
An operator such that each stuttering step changes $s.val$ to $decr(s.val)$.

context
The context in which A appears. It is the expression that is evaluated to determine the value to which $s.ctxt$ is set.

enabled
A formula equivalent to ENABLED A . You can always take it to be ENABLED A , but you can usually find an expression that equals ENABLED A in every reachable state of *Spec* but is easier for *TLC* to compute.

EXTENDS *Naturals*, *TLC*

$top \triangleq [top \mapsto \text{"top"}]$

VARIABLES s, *vars*

$NoStutter(A) \triangleq (s = top) \land A \land (s' = s)$

$PostStutter(A,\ actionId,\ context,\ bot,\ initVal,\ decr(\_)) \triangleq$
IF $s = top$ THEN $\land A$
$\land s' = [id \mapsto actionId,\ ctxt \mapsto context,\ val \mapsto initVal]$
ELSE $\land s.id = actionId$
$\land$ UNCHANGED $vars$
$\land s' =$ IF $s.val = bot$ THEN top
ELSE $[s$ EXCEPT $!.val = decr(s.val)]$

$PreStutter(A,\ enabled,\ actionId,\ context,\ bot,\ initVal,\ decr(\_)) \triangleq$
IF $s = top$
THEN $\land enabled$
$\land$ UNCHANGED $vars$
$\land s' = [id \mapsto actionId,\ ctxt \mapsto context,\ val \mapsto initVal]$
ELSE $\land s.id = actionId$
$\land$ IF $s.val = bot$ THEN $\land s.ctxt = context$
$\land A$
$\land s' = top$
ELSE $\land$ UNCHANGED $vars$
$\land s' = [s$ EXCEPT $!.val = decr(s.val)]$

$MayPostStutter(A,\ actionId,\ context,\ bot,\ initVal,\ decr(\_)) \triangleq$
IF $s = top$ THEN $\land A$
$\land s' =$ IF $initVal = bot$
THEN s
ELSE $[id \mapsto actionId,\ ctxt \mapsto context,$
$val \mapsto initVal]$
ELSE $\land s.id = actionId$
$\land$ UNCHANGED $vars$
$\land s' =$ IF $decr(s.val) = bot$
THEN top
ELSE $[s$ EXCEPT $!.val = decr(s.val)]$

$MayPreStutter(A,\ enabled,\ actionId,\ context,\ bot,\ initVal,\ decr(\_)) \triangleq$
IF $s = top$
THEN $\land enabled$
$\land$ UNCHANGED $vars$
$\land$ IF $initVal = bot$
THEN $A \land (s' = s)$
ELSE $s' = [id \mapsto actionId,\ ctxt \mapsto context,\ val \mapsto initVal]$
ELSE $\land s.id = actionId$
$\land$ IF $s.val = bot$ THEN $\land s.ctxt = context$
$\land A$
$\land s' = top$
ELSE $\land$ UNCHANGED $vars$

$\land\ s' = [s \text{ EXCEPT } !.val = decr(s.val)]$

$StutterConstantCondition(Sigma,\ bot,\ decr(\_)) \triangleq$
LET $InverseDecr(S) \triangleq \{sig \in Sigma \setminus S : decr(sig) \in S\}$
$R[n \in Nat] \triangleq$ IF $n = 0$ THEN $\{bot\}$
ELSE LET $T \triangleq R[n-1]$
IN $T \cup InverseDecr(T)$

IN $Sigma =$ UNION $\{R[n] : n \in Nat\}$

LET $InverseDecr(S) \triangleq \{sig \in Sigma \setminus S : decr(sig) \in S\}$
$ReachBot[S \in \text{SUBSET } Sigma] \triangleq$
IF $InverseDecr(S) = \{\}$ THEN S
ELSE $ReachBot[S \cup InverseDecr(S)]$
IN $ReachBot[\{bot\}] = Sigma$

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* Modification History
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