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- MODULE CJupiter
 1
     Model of our own CJupiter protocol.
 6 EXTENDS Integers, OT, TLC, AdditionalFunctionOperators, AdditionalSequenceOperators
     CONSTANTS
           Client,
                             the set of client replicas
           Server,
                             the (unique) server replica
10
           Char,
                              set of characters allowed
11
           InitState
                             the initial state of each replica
12
     Replica \stackrel{\Delta}{=} Client \cup \{Server\}
     List \stackrel{\triangle}{=} Seq(Char \cup Range(InitState)) all possible lists/strings
     MaxLen \triangleq Cardinality(Char) + Len(InitState) the max length of lists in any states;
17
             We assume that all inserted elements are unique.
18
     ClientNum \triangleq Cardinality(Client)
20
     Priority \triangleq \text{CHOOSE } f \in [Client \rightarrow 1 .. ClientNum] : Injective(f)
21
22
           \land Range(InitState) \cap Char = \{\} due to the uniqueness requirement
24
           \land Priority \in [Client \rightarrow 1 .. ClientNum]
26
     The set of all operations. Note: The positions are indexed from 1.
    Rd \stackrel{\triangle}{=} [type : \{ \text{"Rd"} \}]
     \begin{array}{l} \textit{Tu} = [\textit{type}: \{ \textit{``log} \}] \\ \textit{Del} \stackrel{\triangle}{=} [\textit{type}: \{ \textit{``log} \}, \; \textit{pos}: 1 \ldots \textit{MaxLen}] \\ \textit{Ins} \stackrel{\triangle}{=} [\textit{type}: \{ \textit{``log} \}, \; \textit{pos}: 1 \ldots (\textit{MaxLen} + 1), \; \textit{ch}: \textit{Char}, \; \textit{pr}: 1 \ldots \textit{ClientNum}] \; \; \textit{pr}: \; \textit{priority} \end{array}
     Op \triangleq Ins \cup Del
     Cop: operation of type Op with context
    Oid \stackrel{\triangle}{=} [c:Client, seq:Nat] operation identifier
    Cop \triangleq [op : Op \cup \{Nop\}, oid : Oid, ctx : SUBSET Oid, sctx : SUBSET Oid]
     tb: Is cop1 totally ordered before cop2?
     At a given replica r \in Replica, these can be determined in terms of sctx.
     tb(cop1, cop2, r) \stackrel{\triangle}{=}
48
           \lor cop1.oid \in cop2.sctx
49
            \lor \land cop1.oid \notin cop2.sctx
50
                 \land cop2.oid \notin cop1.sctx
51
                 \land cop1.oid.c \neq r
52
     OT of two operations of type Cop.
    COT(lcop, rcop) \triangleq [lcop \ EXCEPT \ !.op = Xform(lcop.op, rcop.op), \ !.ctx = @ \cup \{rcop.oid\}]
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VARIABLES
          For the client replicas:
                       cseq[c]: local sequence number at client c \in Client
 63
          For the server replica:
           soids.
 67
                      the set of operations the Server has executed
          For all replicas: the n-ary ordered state space
 71
           css,
                      css[r]: the n-ary ordered state space at replica r \in Replica
 72
           cur,
                      cur[r]: the current node of css at replica r \in Replica
           state,
                      state[r]: state (the list content) of replica r \in Replica
 73
          For communication between the Server and the Clients:
           cincoming,
                             cincoming[c]: incoming channel at the client c \in Client
                             incoming channel at the Server
           sincoming,
 78
          For model checking:
 82
           chins
                      a set of chars to insert
 84
      comm \stackrel{\Delta}{=} \text{Instance } CSComm \text{ with } Msg \leftarrow Cop
 86
      eVars \triangleq \langle chins \rangle
                               variables for the environment
      cVars \stackrel{\triangle}{=} \langle cseq \rangle
                               variables for the clients
     sVars \stackrel{\triangle}{=} \langle soids \rangle variables for the server
     dsVars \triangleq \langle css, cur, state \rangle
                                                              variables for the data structure: the n-ary ordered state space
      commVars \triangleq \langle cincoming, sincoming \rangle
                                                              variables for communication
     vars \triangleq \langle eVars, eVars, sVars, commVars, dsVars \rangle all variables
     An css is a directed graph with labeled edges.
     It is represented by a record with node field and edge field.
     Each node is characterized by its context, a set of operations.
     Each edge is labeled with an operation. For clarity, we denote edges by records instead of tuples.
104 \; IsCSS(G) \stackrel{\triangle}{=}
           \land G = [node \mapsto G.node, edge \mapsto G.edge]
105
           \land G.node \subseteq (SUBSET\ Oid)
106
           \land G.edge \subseteq [from : G.node, to : G.node, cop : Cop]
107
      TypeOK \triangleq
109
          For the client replicas:
           \land cseq \in [Client \rightarrow Nat]
113
          For the server replica:
           \land soids \subseteq Oid
117
          For all replicas: the n-ary ordered state space
           \land \forall r \in Replica : IsCSS(css[r])
121
122
           \land cur \in [Replica \rightarrow SUBSET\ Oid]
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\land state \in [Replica \rightarrow List]
123
          For communication between the server and the clients:
           \land comm! TypeOK
127
          For model checking:
           \land chins \subseteq Char
131
132 F
     The Init predicate.
     Init \stackrel{\triangle}{=}
136
          For the client replicas:
           \land cseq = [c \in Client \mapsto 0]
140
          For the server replica:
           \land soids = \{\}
144
          For all replicas: the n-ary ordered state space
           \land css = [r \in Replica \mapsto [node \mapsto \{\{\}\}\}, edge \mapsto \{\}]]
148
           \land cur = [r \in Replica \mapsto \{\}]
149
           \land state = [r \in Replica \mapsto InitState]
150
          For communication between the server and the clients:
           \land comm!Init
154
          For model checking:
           \wedge chins = Char
158
159 ⊦
     Locate the node in rcss which matches the context ctx of cop.
     rcss: the css at replica r \in Replica
165 Locate(cop, rcss) \stackrel{\Delta}{=} CHOOSE \ n \in (rcss.node) : n = cop.ctx
     xForm: iteratively transform cop with a path through the css at replica r \in Replica, following
     the first edges.
     xForm(cop, r) \triangleq
171
           LET rcss \stackrel{\triangle}{=} css[r]
172
                 u \triangleq Locate(cop, rcss)
173
                 v \triangleq u \cup \{cop.oid\}
174
                 RECURSIVE xFormHelper(\_, \_, \_, \_)
175
                  'h' stands for "helper"; xcss: eXtra css created during transformation
176
                 xFormHelper(uh, vh, coph, xcss) \stackrel{\Delta}{=}
177
                      IF uh = cur[r]
178
                       THEN xcss
179
                       ELSE LET fedge \stackrel{\Delta}{=} \text{CHOOSE } e \in rcss.edge :
180
                                                     \wedge e.from = uh
181
                                                     \land \forall uhe \in rcss.edge:
182
                                                         (uhe.from = uh \land uhe \neq e) \Rightarrow tb(e.cop, uhe.cop, r)
183
                                     uprime \triangleq fedge.to
184
                                     fcop \triangleq fedge.cop
185
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\begin{array}{ccc} coph2fcop & \triangleq & COT(coph, fcop) \\ fcop2coph & \triangleq & COT(fcop, coph) \end{array}
186
187
                                          vprime \stackrel{\triangle}{=} vh \cup \{fcop.oid\}
188
                                          xFormHelper(uprime, vprime, coph2fcop,
                                   IN
189
                                               [xcss \ EXCEPT \ !.node = @ \circ \langle vprime \rangle,
190
                                                                    the order of recording edges here is important
191
                                                                   !.edge = @ \circ \langle [from \mapsto vh, to \mapsto vprime, cop \mapsto fcop2coph],
192
                                                                                       [from \mapsto uprime, to \mapsto vprime, cop \mapsto coph2fcop]\rangle])
193
                  xFormHelper(u, v, cop, [node \mapsto \langle v \rangle,
194
                                                      edge \mapsto \langle [from \mapsto u, to \mapsto v, cop \mapsto cop] \rangle])
195
      Perform cop at replica r \in Replica.
      Perform(cop, r) \triangleq
200
            LET xcss \stackrel{\triangle}{=} xForm(cop, r)
201
                   xn \stackrel{\Delta}{=} xcss.node
202
                   xe \stackrel{\triangle}{=} xcss.edge
203
                  xcur \triangleq Last(xn)
204
                   xcop \stackrel{\triangle}{=} Last(xe).cop
205
                  \wedge css' = [css \ EXCEPT \ ![r].node = @ \cup Range(xn),
206
                                                    ![r].edge = @ \cup Range(xe)]
207
                   \wedge cur' = [cur \ EXCEPT \ ![r] = xcur]
208
                   \land state' = [state \ EXCEPT \ ![r] = Apply(xcop.op, @)]
209
210
      Client c \in Client issues an operation op.
      DoOp(c, op) \stackrel{\Delta}{=} op: the raw operation generated by the client c \in Client
214
                \land cseq' = [cseq \ EXCEPT \ ![c] = @ + 1]
215
                \land LET cop \stackrel{\triangle}{=} [op \mapsto op, oid \mapsto [c \mapsto c, seq \mapsto cseq'[c]], ctx \mapsto cur[c], sctx \mapsto \{\}]
216
                          \land Perform(cop, c)
217
                           \land comm! CSend(cop)
218
      DoIns(c) \triangleq
220
            \exists ins \in Ins :
221
                \land ins.pos \in 1 .. (Len(state[c]) + 1)
222
                \land ins.ch \in chins
223
                \wedge ins.pr = Priority[c]
224
                \wedge chins' = chins \setminus \{ins.ch\} We assume that all inserted elements are unique.
225
                \wedge DoOp(c, ins)
226
                \land UNCHANGED sVars
227
       DoDel(c) \triangleq
229
            \exists del \in Del:
230
                \land del.pos \in 1 .. Len(state[c])
231
                \wedge DoOp(c, del)
232
                \land UNCHANGED \langle sVars, eVars \rangle
233
235 Do(c) \triangleq
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\vee DoIns(c)
236
            \vee DoDel(c)
237
     Client c \in Client receives a message from the Server.
     Rev(c) \triangleq
241
            \land comm! CRev(c)
242
            \wedge LET cop \stackrel{\triangle}{=} Head(cincoming[c]) the received original operation
243
                    Perform(cop, c)
244
            \wedge PrintT(css)
245
            \land UNCHANGED \langle eVars, eVars, sVars \rangle
246
247
     The Server receives a message.
      SRev \triangleq
251
           \land comm! SRev
252
           \land LET cop \triangleq [Head(sincoming) \ EXCEPT !.sctx = soids]
                                                                                     set its sctx field
253
                    \land soids' = soids \cup \{cop.oid\}
254
                     \land Perform(cop, Server)
255
                     \land comm! SSendSame(cop.oid.c, cop) broadcast the original operation
256
           \land UNCHANGED \langle eVars, cVars \rangle
257
258
     The next-state relation.
     Next \triangleq
^{262}
           \vee \exists c \in Client : Do(c) \vee Rev(c)
263
           \vee SRev
264
     The Spec. There is no requirement that the clients ever generate operations.
     Spec \stackrel{\Delta}{=} Init \wedge \Box [Next]_{vars} \wedge WF_{vars}(SRev \vee \exists c \in Client : Rev(c))
269
270 ⊦
     The compactness of CJupiter: the css at all replicas are essentially the same.
     IgnoreSctx(rcss) \triangleq
275
          [rcss except !.edge = {[e except !.cop.sctx = {}] : e \in @}]
276
      Compactness \triangleq
278
          comm!EmptyChannel \Rightarrow Cardinality(\{IgnoreSctx(css[r]) : r \in Replica\}) = 1
279
     Theorem Spec \Rightarrow Compactness
282 L
      \* Modification History
      \* Last modified Thu Nov 01 14:15:01 CST 2018 by hengxin
      \* Created Sat Sep 01 11:08:00 CST 2018 by hengxin
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