Model Checking Op-based Counter Using TLA+

Hengfeng Wei

hfwei@nju.edu.cn

June 8, 2018



Our Goal: Model checking RDT using TLA+

(RDT: Replicated Data Types)

Our Goal: Model checking RDT using TLA+ (RDT: Replicated Data Types)

Counter, Register, Set, Graph, List, · · ·

Our Goal: Model checking RDT using TLA+ (RDT: Replicated Data Types)

Counter, Register, Set, Graph, List, · · ·

Various protocols + Various specifications

Modularity



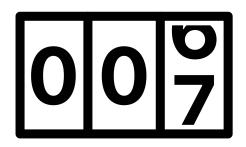
Modularity



Communication (FIFO, Causal, · · ·)

Composition (Set \Longrightarrow Graph)

Specification (Eventual Convergence)



$$\Sigma = \mathbb{N}_0 \times \mathbb{N}_0
M = \mathbb{N}_0
\sigma_0 = \langle 0, 0 \rangle$$

$$\Sigma = \mathbb{N}_0 \times \mathbb{N}_0
M = \mathbb{N}_0
\sigma_0 = \langle 0, 0 \rangle$$

 $\langle a,d \rangle$: $\langle \text{current value, } \# \text{ of incs since the last send} \rangle$

$$\Sigma = \mathbb{N}_0 \times \mathbb{N}_0
M = \mathbb{N}_0
\sigma_0 = \langle 0, 0 \rangle$$

$\langle a,d \rangle$: $\langle \text{current value}, \# \text{ of incs since the last send} \rangle$

$$\begin{split} \operatorname{rd}(\langle a, d \rangle) &= \langle a, d \rangle \\ \operatorname{inc}(\langle a, d \rangle) &= \langle a+1, d+1 \rangle \\ \operatorname{send}(\langle a, d \rangle) &= (\langle a, 0 \rangle, d) \\ \operatorname{rcv}(\langle a, d \rangle, \underline{d'}) &= \langle a+d', d \rangle \end{split}$$



EC: Eventual Consistency/Convergence

"if clients stop issuing ${\tt INCs}$, then the counters at all replicas will be eventually the same."

EC: Eventual Consistency/Convergence

"if clients stop issuing ${\tt INCs}$, then the counters at all replicas will be eventually the same."

$$\diamondsuit(\forall r_i, r_j \in \mathcal{R} : c_i@r_i = c_j@r_j \land c_i@r_i \neq 0)$$

QC: Quiescent Consistency

"if the system is at quiescent, then the counters at all replicas must be the same."

QC: Quiescent Consistency

"if the system is at quiescent, then the counters at all replicas must be the same."

$$\Box \Big((\forall r_i \in \mathcal{R} : d@r_i = 0 \land incoming@r = \emptyset)$$

$$\implies (\forall r_i, r_j \in \mathcal{R} : c_i@r_i = c_j@r_j) \Big)$$

SEC: Strong Eventual Consistency/Convergence

"if two replicas have processed the same set of INCs, then the counters at these two replicas must be the same."

SEC: Strong Eventual Consistency/Convergence

"if two replicas have processed the same set of INCs, then the counters at these two replicas must be the same."

$$\square\Big(\forall r_i, r_j \in \mathcal{R}: (\{C_i\} @ r_i = \{C_j\} @ r_j) \implies (c_i @ r_i = c_j @ r_j)\Big)$$



Cannot: loss, duplication

Can: reordering ($\{\}$ vs. $\langle\langle\rangle\rangle$)

Thank You!



Office 302

Mailbox: H016

hfwei@nju.edu.cn