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1  |----- MODULE CJupiter -----|
   | Model of our own CJupiter protocol. |
5  | EXTENDS JupiterInterface |
6  |-----|
   | Cop: operation of type Op with context |
10 |  $Oid \triangleq [c : Client, seq : Nat]$  operation identifier
11 |  $Cop \triangleq [op : Op \cup \{Nop\}, oid : Oid, ctx : SUBSET\ Oid]$ 

   | tb: Is cop1 totally ordered before cop2? |
   | This can be determined according to the serial view (sv) of any replica. |
18 |  $tb(cop1, cop2, sv) \triangleq$ 
19 |   LET  $pos1 \triangleq FirstIndexOfElementSafe(sv, cop1.oid)$ 
20 |        $pos2 \triangleq FirstIndexOfElementSafe(sv, cop2.oid)$ 
21 |   IN IF  $pos1 \neq 0 \wedge pos2 \neq 0$  at the server or both are remote operations
22 |       THEN  $pos1 < pos2$  at a client: one is a remote operation and the other is a local operation
23 |       ELSE  $pos1 \neq 0$ 
   | OT of two operations of type Cop. |
27 |  $COT(lcop, rcop) \triangleq [lcop\ EXCEPT\ !.op = Xform(lcop.op, rcop.op), !.ctx = @ \cup \{rcop.oid\}]$ 
28 |-----|
29 | VARIABLES |
   | For the client replicas: |
33 |  $cseq,$   $cseq[c]$ : local sequence number at client  $c \in Client$ 
   | For all replicas: the n-ary ordered state space |
37 |  $css,$   $css[r]$ : the n-ary ordered state space at replica  $r \in Replica$ 
38 |  $cur,$   $cur[r]$ : the current node of  $css$  at replica  $r \in Replica$ 
39 |  $state,$   $state[r]$ : state (the list content) of replica  $r \in Replica$ 
   | For edge ordering in CSS |
43 |  $serial,$   $serial[r]$ : the serial view of replica  $r \in Replica$  about the server
44 |  $cincomingSerial,$ 
45 |  $sincomingSerial,$ 
   | For communication between the Server and the Clients: |
49 |  $cincoming,$   $cincoming[c]$ : incoming channel at the client  $c \in Client$ 
50 |  $sincoming,$  incoming channel at the Server
   | For model checking: |
54 |  $chins$  a set of chars to insert
55 |-----|
56 |  $serialVars \triangleq \langle serial, cincomingSerial, sincomingSerial \rangle$ 
57 |  $vars \triangleq \langle chins, cseq, css, cur, state, cincoming, sincoming, serialVars \rangle$ 
58 |-----|
59 |  $comm \triangleq$  INSTANCE CComm WITH  $Msg \leftarrow Cop$ 
60 |  $commSerial \triangleq$  INSTANCE CComm WITH  $Msg \leftarrow Seq(Oid),$ 
61 |  $cincoming \leftarrow cincomingSerial, sincoming \leftarrow sincomingSerial$ 
62 |-----|

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A *css* is a directed graph with labeled edges, represented by a record with node field and edge field.
Each node is characterized by its context, a set of oids. Each edge is labeled with an operation.

69 $IsCSS(G) \triangleq$
70 $\wedge G = [node \mapsto G.node, edge \mapsto G.edge]$
71 $\wedge G.node \subseteq (SUBSET\ Oid)$
72 $\wedge G.edge \subseteq [from : G.node, to : G.node, cop : Cop]$

74 $EmptySS \triangleq [node \mapsto \{\{\}\}, edge \mapsto \{\}]$

76 $TypeOK \triangleq$
For the client replicas:
80 $\wedge cseq \in [Client \rightarrow Nat]$
For edge ordering in *CSS*:
84 $\wedge serial \in [Replica \rightarrow Seq(Oid)]$
85 $\wedge commSerial! TypeOK$
For all replicas: the *n*-ary ordered state space
89 $\wedge \forall r \in Replica : IsCSS(css[r])$
90 $\wedge cur \in [Replica \rightarrow SUBSET\ Oid]$
91 $\wedge state \in [Replica \rightarrow List]$
For communication between the server and the clients:
95 $\wedge comm! TypeOK$
For model checking:
99 $\wedge chins \subseteq Char$

101 $Init \triangleq$
For the client replicas:
105 $\wedge cseq = [c \in Client \mapsto 0]$
For the server replica:
109 $\wedge serial = [r \in Replica \mapsto \langle \rangle]$
110 $\wedge commSerial! Init$
For all replicas: the *n*-ary ordered state space
114 $\wedge css = [r \in Replica \mapsto EmptySS]$
115 $\wedge cur = [r \in Replica \mapsto \{\}]$
116 $\wedge state = [r \in Replica \mapsto InitState]$
For communication between the server and the clients:
120 $\wedge comm! Init$
For model checking:
124 $\wedge chins = Char$

125
Locate the node in *rcss* (the *css* at replica $r \in Replica$) that matches the context *ctx* of cop.
129 $Locate(cop, rcss) \triangleq \text{CHOOSE } n \in rcss.node : n = cop.ctx$
Take union of two state spaces *ss1* and *ss2*.

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133  $ss1 \oplus ss2 \triangleq [node \mapsto ss1.node \cup ss2.node, edge \mapsto ss1.edge \cup ss2.edge]$ 
    xForm: Iteratively transform cop with a path through the css at replica  $r \in Replica$ , following
    the first edges.
138  $xForm(cop, r) \triangleq$ 
139   LET  $rcss \triangleq css[r]$ 
140    $u \triangleq Locate(cop, rcss)$ 
141    $v \triangleq u \cup \{cop.oid\}$ 
142   RECURSIVE  $xFormHelper(-, -, -, -, -, -)$ 
143   'h' stands for "helper"; xcss: eXtra css created during transformation
144    $xFormHelper(uh, vh, coph, xcss, xcoph, xcurh) \triangleq$ 
145     IF  $uh = cur[r]$ 
146     THEN  $\langle xcss, xcoph, xcurh \rangle$ 
147     ELSE LET  $fedge \triangleq$  CHOOSE  $e \in rcss.edge :$ 
148        $\wedge e.from = uh$ 
149        $\wedge \forall uhe \in rcss.edge :$ 
150          $(uhe.from = uh \wedge uhe \neq e) \Rightarrow tb(e.cop, uhe.cop, serial[r])$ 
151        $uprime \triangleq fedge.to$ 
152        $fcop \triangleq fedge.cop$ 
153        $coph2fcop \triangleq COT(coph, fcop)$ 
154        $fcop2coph \triangleq COT(fcop, coph)$ 
155        $vprime \triangleq vh \cup \{fcop.oid\}$ 
156     IN  $xFormHelper(uprime, vprime, coph2fcop,$ 
157        $[xcss \text{ EXCEPT } !.node = @ \cup \{vprime\},$ 
158        $!.edge = @ \cup \{[from \mapsto vh, to \mapsto vprime, cop \mapsto fcop2coph],$ 
159        $[from \mapsto uprime, to \mapsto vprime, cop \mapsto coph2fcop]\},$ 
160        $coph2fcop, vprime)$ 
161     IN  $xFormHelper(u, v, cop, [node \mapsto \{v\}, edge \mapsto \{[from \mapsto u, to \mapsto v, cop \mapsto cop]\}], cop, v)$ 
    Perform cop at replica  $r \in Replica$ .
165  $Perform(cop, r) \triangleq$ 
166   LET  $xform \triangleq xForm(cop, r)$  xform:  $\langle xcss, xcop, xcur \rangle$ 
167    $xcss \triangleq xform[1]$ 
168    $xcop \triangleq xform[2]$ 
169    $xcur \triangleq xform[3]$ 
170   IN  $\wedge css' = [css \text{ EXCEPT } ![r] = @ \oplus xcsc]$ 
171    $\wedge cur' = [cur \text{ EXCEPT } ![r] = xcur]$ 
172    $\wedge state' = [state \text{ EXCEPT } ![r] = Apply(xcop.op, @)]$ 
173 |-----|
    Client  $c \in Client$  issues an operation op.
177  $DoOp(c, op) \triangleq$  op: the raw operation generated by the client  $c \in Client$ 
178    $\wedge cseq' = [cseq \text{ EXCEPT } ![c] = @ + 1]$ 
179    $\wedge$  LET  $cop \triangleq [op \mapsto op, oid \mapsto [c \mapsto c, seq \mapsto cseq'[c]], ctx \mapsto cur[c]]$ 
180   IN  $\wedge Perform(cop, c)$ 
181    $\wedge comm! CSend(cop)$ 
183  $DoIns(c) \triangleq$ 

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184 $\exists ins \in \{op \in Ins : op.pos \in 1 \dots (Len(state[c]) + 1) \wedge op.ch \in chins \wedge op.pr = Priority[c]\} :$
 185 $\wedge DoOp(c, ins)$
 186 $\wedge chins' = chins \setminus \{ins.ch\}$ We assume that all inserted elements are unique.
 187 $\wedge UNCHANGED \langle serialVars \rangle$

189 $DoDel(c) \triangleq$
 190 $\exists del \in \{op \in Del : op.pos \in 1 \dots Len(state[c])\} :$
 191 $\wedge DoOp(c, del)$
 192 $\wedge UNCHANGED \langle chins, serialVars \rangle$

194 $Do(c) \triangleq$
 195 $\vee DoIns(c)$
 196 $\vee DoDel(c)$

Client $c \in Client$ receives a message from the *Server*.

200 $Rev(c) \triangleq$
 201 $\wedge comm!CRev(c)$
 202 $\wedge Perform(Head(cincoming[c]), c)$
 203 $\wedge commSerial!CRev(c)$
 204 $\wedge serial' = [serial \text{ EXCEPT } ![c] = Head(cincomingSerial[c])]$
 205 $\wedge UNCHANGED \langle chins, cseq \rangle$

The *Server* receives a message.

210 $SRev \triangleq$
 211 $\wedge comm!SRev$
 212 $\wedge LET \ cop \triangleq Head(sincoming)$
 213 $IN \ \wedge Perform(cop, Server)$
 214 $\wedge comm!SSendSame(cop.oid.c, cop)$ broadcast the original operation
 215 $\wedge serial' = [serial \text{ EXCEPT } ![Server] = Append(@, cop.oid)]$
 216 $\wedge commSerial!SSendSame(cop.oid.c, serial'[Server])$
 217 $\wedge UNCHANGED \langle chins, cseq, sincomingSerial \rangle$

219 $Next \triangleq$
 220 $\vee \exists c \in Client : Do(c) \vee Rev(c)$
 221 $\vee SRev$

Fairness: There is no requirement that the clients ever generate operations.

225 $Fairness \triangleq$
 226 $WF_{vars}(SRev \vee \exists c \in Client : Rev(c))$

228 $Spec \triangleq Init \wedge \Box[Next]_{vars} \wedge Fairness$ (We care more about safety.)

The compactness of *CJupiter*: the *CSSes* at all replicas are the same.

233 $Compactness \triangleq$
 234 $comm!EmptyChannel \Rightarrow Cardinality(Range(css)) = 1$

236 THEOREM $Spec \Rightarrow Compactness$

237 |
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