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1  |----- MODULE XJupiter -----|
   | Specification of the Jupiter protocol described in CSCW'2014 by Yi Xu, Chengzheng Sun, and |
   | Mo Li. We call it XJupiter, with 'X' for "Xu". |
7  | EXTENDS StateSpace |
8  |-----|
9  | VARIABLES |
   | The 2D state spaces (2ss, for short). Each client maintains one 2D state space. The server |
   | maintains n 2D state spaces, one for each client. |
15 |   c2ss,   c2ss[c]: the 2D state space at client c ∈ Client |
16 |   s2ss    s2ss[c]: the 2D state space maintained by the Server for client c ∈ Client |
18 |   vars ≜ ⟨intVars, ctxVars, c2ss, s2ss⟩ |
19 |-----|
20 |   TypeOK ≜ |
21 |     ∧ TypeOKInt |
22 |     ∧ TypeOKCtx |
23 |     ∧ Comm(Cop)! TypeOK |
24 |     ∧ ∀ c ∈ Client : IsSS(c2ss[c]) ∧ IsSS(s2ss[c]) |
25 |-----|
26 |   Init ≜ |
27 |     ∧ InitInt |
28 |     ∧ InitCtx |
29 |     ∧ Comm(Cop)! Init |
30 |     ∧ c2ss = [c ∈ Client ↦ EmptySS] |
31 |     ∧ s2ss = [c ∈ Client ↦ EmptySS] |
32 |-----|
   | xForm: iteratively transform cop with a path through the 2D state space ss at some client. |
37 | xForm(cop, ss, current) ≜ |
38 |   LET u ≜ Locate(cop, ss) |
39 |   v ≜ u ∪ {cop.oid} |
40 |   RECURSIVE xFormHelper(-, -, -, -, -) |
41 |   'h' stands for "helper"; xss: eXtra ss created during transformation |
42 |   xFormHelper(uh, vh, coph, xss, xcoph) ≜ |
43 |     IF uh = current |
44 |       THEN ⟨xss, xcoph⟩ |
45 |     ELSE LET e ≜ CHOOSE e ∈ ss.edge : e.from = uh ∧ ClientOf(e.cop) ≠ ClientOf(cop) |
46 |       uprime ≜ e.to |
47 |       copprime ≜ e.cop |
48 |       coph2copprime ≜ COT(coph, copprime) |
49 |       copprime2coph ≜ COT(copprime, coph) |
50 |       vprime ≜ vh ∪ {copprime.oid} |
51 |     IN xFormHelper(uprime, vprime, coph2copprime, |
52 |       [node ↦ xss.node ∪ {vprime}, |
53 |       edge ↦ xss.edge ∪ {[from ↦ vh, to ↦ vprime, cop ↦ copprime2coph], |
54 |       [from ↦ uprime, to ↦ vprime, cop ↦ coph2copprime]}, |

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coph2copprime
56 IN $x\text{FormHelper}(u, v, \text{cop}, [\text{node} \mapsto \{v\}, \text{edge} \mapsto \{\text{from} \mapsto u, \text{to} \mapsto v, \text{cop} \mapsto \text{cop}\}], \text{cop})$
57 \vdash
Client $c \in \text{Client}$ perform operation cop .
61 $\text{ClientPerform}(\text{cop}, c) \triangleq$
62 LET $x\text{form} \triangleq x\text{Form}(\text{cop}, c2ss[c], ds[c])$ $x\text{form}: \langle xss, xcop \rangle$
63 $xss \triangleq x\text{form}[1]$
64 $xcop \triangleq x\text{form}[2]$
65 IN $\wedge c2ss' = [c2ss \text{ EXCEPT } ![c] = @ \oplus xss]$
66 $\wedge \text{state}' = [\text{state} \text{ EXCEPT } ![c] = \text{Apply}(xcop.op, @)]$
Client $c \in \text{Client}$ generates an operation op .
70 $\text{DoOp}(c, op) \triangleq$
71 LET $\text{cop} \triangleq [op \mapsto op, oid \mapsto [c \mapsto c, seq \mapsto cseq'[c]], ctx \mapsto ds[c]]$
72 IN $\wedge \text{ClientPerform}(\text{cop}, c)$
73 $\wedge \text{UpdateDS}(c, \text{cop})$
74 $\wedge \text{Comm}(\text{Cop})!C\text{Send}(\text{cop})$
76 $\text{DoIns}(c) \triangleq$
77 $\exists ins \in \{op \in \text{Ins} : op.pos \in 1 \dots (\text{Len}(\text{state}[c]) + 1) \wedge op.ch \in \text{chins} \wedge op.pr = \text{Priority}[c]\} :$
78 $\wedge \text{DoOp}(c, ins)$
79 $\wedge \text{chins}' = \text{chins} \setminus \{ins.ch\}$ We assume that all inserted elements are unique.
81 $\text{DoDel}(c) \triangleq$
82 $\exists del \in \{op \in \text{Del} : op.pos \in 1 \dots \text{Len}(\text{state}[c])\} :$
83 $\wedge \text{DoOp}(c, del)$
84 $\wedge \text{UNCHANGED } \text{chins}$
86 $\text{Do}(c) \triangleq$
87 $\wedge \text{DoCtx}(c)$
88 $\wedge \vee \text{DoIns}(c)$
89 $\vee \text{DoDel}(c)$
90 $\wedge \text{UNCHANGED } s2ss$
Client $c \in \text{Client}$ receives a message from the *Server*.
94 $\text{Rev}(c) \triangleq$
95 $\wedge \text{Comm}(\text{Cop})!C\text{Rev}(c)$
96 $\wedge \text{LET } \text{cop} \triangleq \text{Head}(\text{cincoming}[c])$ the received (transformed) operation
97 IN $\text{ClientPerform}(\text{cop}, c)$
98 $\wedge \text{RevCtx}(c)$
99 $\wedge \text{UNCHANGED } \langle \text{chins}, s2ss \rangle$
100 \vdash
The *Server* performs operation cop .
104 $\text{ServerPerform}(\text{cop}) \triangleq$
105 LET $c \triangleq \text{ClientOf}(\text{cop})$
106 $\text{scur} \triangleq ds[\text{Server}]$
107 $x\text{form} \triangleq x\text{Form}(\text{cop}, s2ss[c], \text{scur})$ $x\text{form}: \langle xss, xcop \rangle$

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108      $xss \triangleq xform[1]$ 
109      $xcop \triangleq xform[2]$ 
110      $xcur \triangleq scur \cup \{cop.oid\}$ 
111     IN  $\wedge s2ss' = [cl \in Client \mapsto$ 
112         IF  $cl = c$ 
113         THEN  $s2ss[cl] \oplus xss$ 
114         ELSE  $s2ss[cl] \oplus [node \mapsto \{xcur\},$ 
115              $edge \mapsto \{[from \mapsto scur, to \mapsto xcur, cop \mapsto xcop]\}]$ 
116         ]
117      $\wedge state' = [state \text{ EXCEPT } ![Server] = Apply(xcop.op, @)]$ 
118      $\wedge Comm(Cop)!SSendSame(c, xcop)$  broadcast the transformed operation
    The Server receives a message.
122  $SRev \triangleq$ 
123      $\wedge Comm(Cop)!SRev$ 
124      $\wedge LET cop \triangleq Head(sincoming)$ 
125     IN  $ServerPerform(cop)$ 
126      $\wedge SRevCtx$ 
127      $\wedge UNCHANGED \langle chins, c2ss \rangle$ 
128 |-----|
129  $Next \triangleq$ 
130      $\vee \exists c \in Client : Do(c) \vee Rev(c)$ 
131      $\vee SRev$ 
133  $Fairness \triangleq$ 
134      $WF_{vars}(SRev \vee \exists c \in Client : Rev(c))$ 
136  $Spec \triangleq Init \wedge \Box[Next]_{vars} \wedge Fairness$ 
137 |-----|
    In Jupiter (not limited to XJupiter), each client synchronizes with the server. In XJupiter, this
    is expressed as the following CSSync property.
142  $CSSync \triangleq$ 
143      $\forall c \in Client : (ds[c] = ds[Server]) \Rightarrow c2ss[c] = s2ss[c]$ 
144 |-----|
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