
MODULE *ZabWithQTest2*

EXTENDS *Zab*

constants that uniquely used for constraining state space in model checking
 CONSTANTS *MaxElectionNum*, *MaxTotalRestartNum*, *MaxTransactionNum*

variables that uniquely used for constraining state space in model checking
 VARIABLES *electionNum*, the round of leader election, not equal to $Maximum\{currentEpoch[i] : i \in Server\}$,
 because *currentEpoch* will increase only when follower receives *NEWEPOCH*,
 and it is common that some round of election ends without leader broadcasting *NEWEPOCH*
 or follower receiving *NEWEPOCH*.
totalRestartNum the number of restart from all servers, also as a global variable.

testVars $\triangleq \langle electionNum, totalRestartNum \rangle$

varsT $\triangleq \langle vars, testVars \rangle$

InitT $\triangleq \wedge Init$
 $\wedge electionNum = 0$
 $\wedge totalRestartNum = 0$

ElectionT(*i*, *Q*) \triangleq test restrictions
 $\wedge electionNum < MaxElectionNum$
 $\wedge Election(i, Q)$
 $\wedge electionNum' = electionNum + 1$

InitialElectionT(*i*, *Q*) \triangleq
 $\wedge \forall s \in Server : state[s] = Follower \wedge leaderOracle[s] = NullPoint$
 $\wedge ElectionT(i, Q)$
 $\wedge UNCHANGED \langle currentEpoch, history, commitIndex, currentCounter,$
 $sendCounter, recoveryVars, proposalMsgsLog, totalRestartNum \rangle$

LeaderTimeoutT(*i*, *j*) \triangleq
 $\wedge state[i] \neq Follower$
 $\wedge j \neq i$
 $\wedge j \in cluster[i]$
 $\wedge LET newCluster \triangleq cluster[i] \setminus \{j\}$
 IN $\wedge \vee \wedge newCluster \in Quorums$
 $\wedge cluster' = [cluster \text{ EXCEPT } ![i] = newCluster]$
 $\wedge clean(i, j)$
 $\wedge UNCHANGED \langle state, cepochRecv, ackeRecv, ackldRecv, ackIndex,$
 $committedIndex, initialHistory, tempMaxEpoch, tempMaxLastEpoch,$
 $tempInitialHistory, leaderOracle, leaderEpoch, cepochSent, electionNum \rangle$

$$\begin{aligned}
& \vee \wedge \text{newCluster} \notin \text{Quorums} \\
& \wedge \exists Q \in \text{Quorums} : \wedge i \in Q \\
& \wedge \exists v \in Q : \text{ElectionT}(v, Q) \\
& \wedge \text{UNCHANGED} \langle \text{currentEpoch}, \text{history}, \text{commitIndex}, \text{currentCounter}, \text{sendCounter}, \\
& \quad \text{recoveryVars}, \text{proposalMsgsLog}, \text{totalRestartNum} \rangle \\
\text{FollowerTimeoutT}(i) & \triangleq \\
& \wedge \text{state}[i] = \text{Follower} \\
& \wedge \text{leaderOracle}[i] \neq \text{NullPoint} \\
& \wedge \exists Q \in \text{Quorums} : \wedge i \in Q \\
& \wedge \exists v \in Q : \text{ElectionT}(v, Q) \\
& \wedge \text{UNCHANGED} \langle \text{currentEpoch}, \text{history}, \text{commitIndex}, \text{currentCounter}, \text{sendCounter}, \\
& \quad \text{recoveryVars}, \text{proposalMsgsLog}, \text{totalRestartNum} \rangle \\
\hline
\text{RestartT}(i) & \triangleq \text{test restrictions} \\
& \wedge \text{totalRestartNum} < \text{MaxTotalRestartNum} \\
& \wedge \text{totalRestartNum}' = \text{totalRestartNum} + 1 \\
& \wedge \text{Restart}(i) \\
& \wedge \text{UNCHANGED } \text{electionNum} \\
\text{RecoveryAfterRestartT}(i) & \triangleq \text{test restrictions} \\
& \wedge \text{totalRestartNum} < \text{MaxTotalRestartNum} \\
& \wedge \text{RecoveryAfterRestart}(i) \\
& \wedge \text{UNCHANGED } \text{testVars} \\
\text{HandleRecoveryRequestT}(i, j) & \triangleq \wedge \text{HandleRecoveryRequest}(i, j) \\
& \wedge \text{UNCHANGED } \text{testVars} \\
\text{HandleRecoveryResponseT}(i, j) & \triangleq \wedge \text{HandleRecoveryResponse}(i, j) \\
& \wedge \text{UNCHANGED } \text{testVars} \\
\text{FindClusterT}(i) & \triangleq \wedge \text{FindCluster}(i) \\
& \wedge \text{UNCHANGED } \text{testVars} \\
\hline
\text{FollowerDiscovery1T}(i) & \triangleq \wedge \text{FollowerDiscovery1}(i) \\
& \wedge \text{UNCHANGED } \text{testVars} \\
\text{LeaderHandleCEPOCHT}(i, j) & \triangleq \wedge \text{LeaderHandleCEPOCH}(i, j) \\
& \wedge \text{UNCHANGED } \text{testVars} \\
\text{LeaderDiscovery1T}(i) & \triangleq \wedge \text{LeaderDiscovery1}(i) \\
& \wedge \text{UNCHANGED } \text{testVars} \\
\text{FollowerDiscovery2T}(i, j) & \triangleq \wedge \text{FollowerDiscovery2}(i, j) \\
& \wedge \text{UNCHANGED } \text{testVars} \\
\text{LeaderHandleACKET}(i, j) & \triangleq \wedge \text{LeaderHandleACKE}(i, j)
\end{aligned}$$

$$\begin{array}{l}
\wedge \text{UNCHANGED } testVars \\
LeaderDiscovery2Sync1T(i) \triangleq \wedge LeaderDiscovery2Sync1(i) \\
\wedge \text{UNCHANGED } testVars \\
\hline
FollowerSync1T(i, j) \triangleq \wedge FollowerSync1(i, j) \\
\wedge \text{UNCHANGED } testVars \\
LeaderHandleACKLDT(i, j) \triangleq \wedge LeaderHandleACKLD(i, j) \\
\wedge \text{UNCHANGED } testVars \\
LeaderSync2T(i) \triangleq \wedge LeaderSync2(i) \\
\wedge \text{UNCHANGED } testVars \\
FollowerSync2T(i, j) \triangleq \wedge FollowerSync2(i, j) \\
\wedge \text{UNCHANGED } testVars \\
\hline
ClientRequestT(i, v) \triangleq \text{test restrictions} \\
\wedge Len(history[i]) < MaxTransactionNum \\
\wedge ClientRequest(i, v) \\
\wedge \text{UNCHANGED } testVars \\
LeaderBroadcast1T(i) \triangleq \wedge LeaderBroadcast1(i) \\
\wedge \text{UNCHANGED } testVars \\
FollowerBroadcast1T(i, j) \triangleq \wedge FollowerBroadcast1(i, j) \\
\wedge \text{UNCHANGED } testVars \\
LeaderHandleACKT(i, j) \triangleq \wedge LeaderHandleACK(i, j) \\
\wedge \text{UNCHANGED } testVars \\
LeaderAdvanceCommitT(i) \triangleq \wedge LeaderAdvanceCommit(i) \\
\wedge \text{UNCHANGED } testVars \\
LeaderBroadcast2T(i) \triangleq \wedge LeaderBroadcast2(i) \\
\wedge \text{UNCHANGED } testVars \\
FollowerBroadcast2T(i, j) \triangleq \wedge FollowerBroadcast2(i, j) \\
\wedge \text{UNCHANGED } testVars \\
\hline
LeaderHandleCEPOCHinPhase3T(i, j) \triangleq \wedge LeaderHandleCEPOCHinPhase3(i, j) \\
\wedge \text{UNCHANGED } testVars \\
LeaderHandleACKLDinPhase3T(i, j) \triangleq \wedge LeaderHandleACKLDinPhase3(i, j) \\
\wedge \text{UNCHANGED } testVars \\
\hline
BecomeFollowerT(i) \triangleq \wedge BecomeFollower(i) \\
\wedge \text{UNCHANGED } testVars
\end{array}$$

$$\text{DiscardStaleMessageT}(i) \triangleq \wedge \text{DiscardStaleMessage}(i) \\ \wedge \text{UNCHANGED } \text{testVars}$$

Defines how the variables may transition.

$$\text{NextT} \triangleq \\ \begin{aligned} & \vee \exists i \in \text{Server}, Q \in \text{Quorums} : \text{InitialElectionT}(i, Q) \\ & \vee \exists i \in \text{Server} : \text{RestartT}(i) \\ & \vee \exists i \in \text{Server} : \text{RecoveryAfterRestartT}(i) \\ & \vee \exists i, j \in \text{Server} : \text{HandleRecoveryRequestT}(i, j) \\ & \vee \exists i, j \in \text{Server} : \text{HandleRecoveryResponseT}(i, j) \\ & \vee \exists i, j \in \text{Server} : \text{FindClusterT}(i) \\ & \vee \exists i, j \in \text{Server} : \text{LeaderTimeoutT}(i, j) \\ & \vee \exists i \in \text{Server} : \text{FollowerTimeoutT}(i) \\ & \vee \exists i \in \text{Server} : \text{FollowerDiscovery1T}(i) \\ & \vee \exists i, j \in \text{Server} : \text{LeaderHandleCEPOCHT}(i, j) \\ & \vee \exists i \in \text{Server} : \text{LeaderDiscovery1T}(i) \\ & \vee \exists i, j \in \text{Server} : \text{FollowerDiscovery2T}(i, j) \\ & \vee \exists i, j \in \text{Server} : \text{LeaderHandleACKET}(i, j) \\ & \vee \exists i \in \text{Server} : \text{LeaderDiscovery2Sync1T}(i) \\ & \vee \exists i, j \in \text{Server} : \text{FollowerSync1T}(i, j) \\ & \vee \exists i, j \in \text{Server} : \text{LeaderHandleACKLDT}(i, j) \\ & \vee \exists i \in \text{Server} : \text{LeaderSync2T}(i) \\ & \vee \exists i, j \in \text{Server} : \text{FollowerSync2T}(i, j) \\ & \vee \exists i \in \text{Server}, v \in \text{Value} : \text{ClientRequestT}(i, v) \\ & \vee \exists i \in \text{Server} : \text{LeaderBroadcast1T}(i) \\ & \vee \exists i, j \in \text{Server} : \text{FollowerBroadcast1T}(i, j) \\ & \vee \exists i, j \in \text{Server} : \text{LeaderHandleACKT}(i, j) \\ & \vee \exists i \in \text{Server} : \text{LeaderAdvanceCommitT}(i) \\ & \vee \exists i \in \text{Server} : \text{LeaderBroadcast2T}(i) \\ & \vee \exists i, j \in \text{Server} : \text{FollowerBroadcast2T}(i, j) \\ & \vee \exists i, j \in \text{Server} : \text{LeaderHandleCEPOCHinPhase3T}(i, j) \\ & \vee \exists i, j \in \text{Server} : \text{LeaderHandleACKLDinPhase3T}(i, j) \\ & \vee \exists i \in \text{Server} : \text{DiscardStaleMessageT}(i) \\ & \vee \exists i \in \text{Server} : \text{BecomeFollowerT}(i) \end{aligned}$$

$$\text{SpecT} \triangleq \text{InitT} \wedge \Box[\text{NextT}]_{\text{varsT}}$$

\ * Modification History
\ * Last modified Mon May 17 22:00:38 CST 2021 by Dell
\ * Created Mon May 17 17:20:01 CST 2021 by Dell