

Lab:07 Infix to Postfix Conversion

Theory:

Infix expression: The expression of the form “a operator b” ($a + b$) i.e., when an operator is in-between every pair of operands.

Postfix expression: The expression of the form “a b operator” ($ab+$) i.e., When every pair of operands is followed by an operator.

Why postfix representation of the expression?

The compiler scans the expression either from left to right or from right to left.

Consider the expression: $a + b * c + d$

- The compiler first scans the expression to evaluate the expression $b * c$, then again scans the expression to add a to it.
- The result is then added to d after another scan.

The repeated scanning makes it very inefficient. Infix expressions are easily readable and solvable by humans whereas the computer cannot differentiate the operators and parenthesis easily so, it is better to convert the expression to postfix (or prefix) form before evaluation.

The corresponding expression in postfix form is $abc*+d+$. The postfix expressions can be evaluated easily using a stack.

Algorithm for infix to postfix conversion

1. Scan the infix expression **from left to right**.
2. Initialize an empty stack
3. If the scanned character is an operand, put it in the postfix expression.
4. Otherwise, do the following
 - If the precedence of the current scanned operator is higher than the precedence of the operator on top of the stack, or if

the stack is empty, or if the stack contains a '(', then push the current operator onto the stack.

- Else, pop all operators from the stack that have precedence higher than or equal to that of the current operator. After that push the current operator onto the stack.
5. If the scanned character is a '(', push it to the stack.
 6. If the scanned character is a ')', pop the stack and output it until a '(' is encountered, and discard both the parenthesis.
 7. Repeat steps **3-6** until the infix expression is scanned.
 8. Once the scanning is over, Pop the stack and add the operators in the postfix expression until it is not empty.
 9. Finally, print the postfix expression.

Source code:

```
#include <stdio.h>
#include <stdlib.h>
#define MAX 50

char stk[50];
int top = -1;

int isEmpty()
{
    return top == -1;
}

int isFull()
{
    return top == MAX - 1;
}

char peek()
{
    if (isEmpty())
        return -1;
    return stk[top];
}

char pop()
{
    if (isEmpty())
```

```

        return -1;

    char ch = stk[top];
    top--;
    return(ch);
}

void push(char oper)
{
    if (isFull())
        printf("Stack Full!!!!");
    else {
        top++;
        stk[top] = oper;
    }
}

int checkIfOperand(char ch)
{
    return (ch >= 'a' && ch <= 'z') || (ch >= 'A' && ch <= 'Z');
}

int precedence(char ch)
{
    switch (ch)
    {
        case '+':
        case '-':
            return 1;

        case '*':
        case '/':
            return 2;

        case '^':
            return 3;
    }
    return -1;
}

void convertInfixToPostfix(char* expression)
{
    int i, j = 0;
    char postfix[MAX];

```

```

for (i = 0; expression[i]; ++i)
{
    if (checkIfOperand(expression[i]))
        postfix[j++] = expression[i];

    else if (expression[i] == '(')
        push(expression[i]);

    else if (expression[i] == ')')
    {
        while (!isEmpty() && peek() != '(')
            postfix[j++] = pop();
        if (!isEmpty() && peek() != '(')
            return; // Invalid expression
        else
            pop();
    }
    else
    {
        while (!isEmpty() && precedence(expression[i]) <=
precedence(peek()))
            postfix[j++] = pop();
        push(expression[i]);
    }
}

while (!isEmpty())
    postfix[j++] = pop();

postfix[j] = '\0'; // Null-terminate the postfix expression

printf("Postfix Expression: %s\n", postfix);
}

int main()
{
    char expression[50];
    printf("Enter the infix expression:");
    scanf("%s", expression);
    convertInfixToPostfix(expression);
    return 0;
}

```

Process feedback link: https://app.processfeedback.org/coding/c_2025-01-27_11-18_45d8c6da-723d-441e-b144-33560f61720d/?token=eca46b09-07ba-496a-8da0-df81b2d0f60e

Discussion and Conclusion

In this lab, we did the conversion of infix expressions to postfix notation by using a stack. Infix expressions are those that follow the standard mathematical notations, while postfix gets rid of the parentheses and precedence of operators by having the operands and operators in a certain order. By utilizing a stack, one can easily manage operators and operands during the process of conversion. In summary, the conversion from infix to postfix simplifies the evaluation of expressions, especially by computers, since it eliminates cumbersome precedence rules and parentheses, thus making it easier to implement in a programming language or calculator.

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