1) Pair with given Sum
2) Pair with given difference
3) Subarray with given sum
4) Container w/t most water

Given a **sorted** integer array A of size N.
find any pair (i, j) s.t. A[i] + A[j] = K ⇒ Given (i != j)

0	1	2	3	4	5	6
-5	-2	1	8	10	12	15

A = [-5, -2, 1, 8, 10, 12, 15]

K = 11 ⇒ (2, 4)

Soln 1) Brute Force

⇒ ∀ pairs (i, j) check if (A[i] + A[j]) == K.

```
for (i = 0 to N-1) {
    for (j = i+1 to N-1) {
        if (A[i] + A[j] == K) {
            return (i, j);
        }
    }
}
```

Use Binary Search to find (K - A[i]) since A is sorted

T.C. = $O(N^2)$

2) Binary Search

```
for (i=0; i<N; i++) {                    N
    s = i+1;
    e = N-1;
    target = K - A[i];                   } log N
    ( B.S. )
}
```

T.C. = $O(N \log N)$

S.C. = $O(1)$

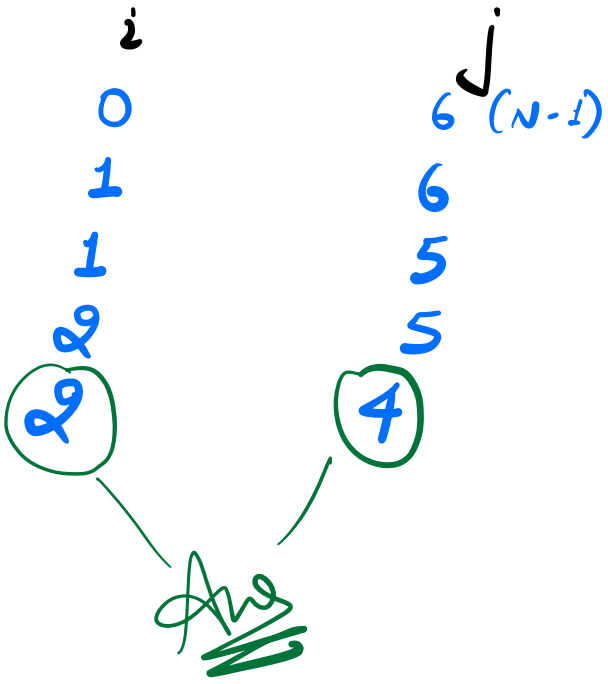
3) Two Pointer

↳ 2 variables pointing to 2 indexes of the array.

Q1 Where to initialise the 2 pointers ??

Q2 How to move the 2 pointers.

	0	1	2	3	4	5	6	
A =	-5,	-2,	1,	8,	10,	12,	15]	K = 11
	i	i	i		j	j	j	



Sum

10 < 11 (Discard i)

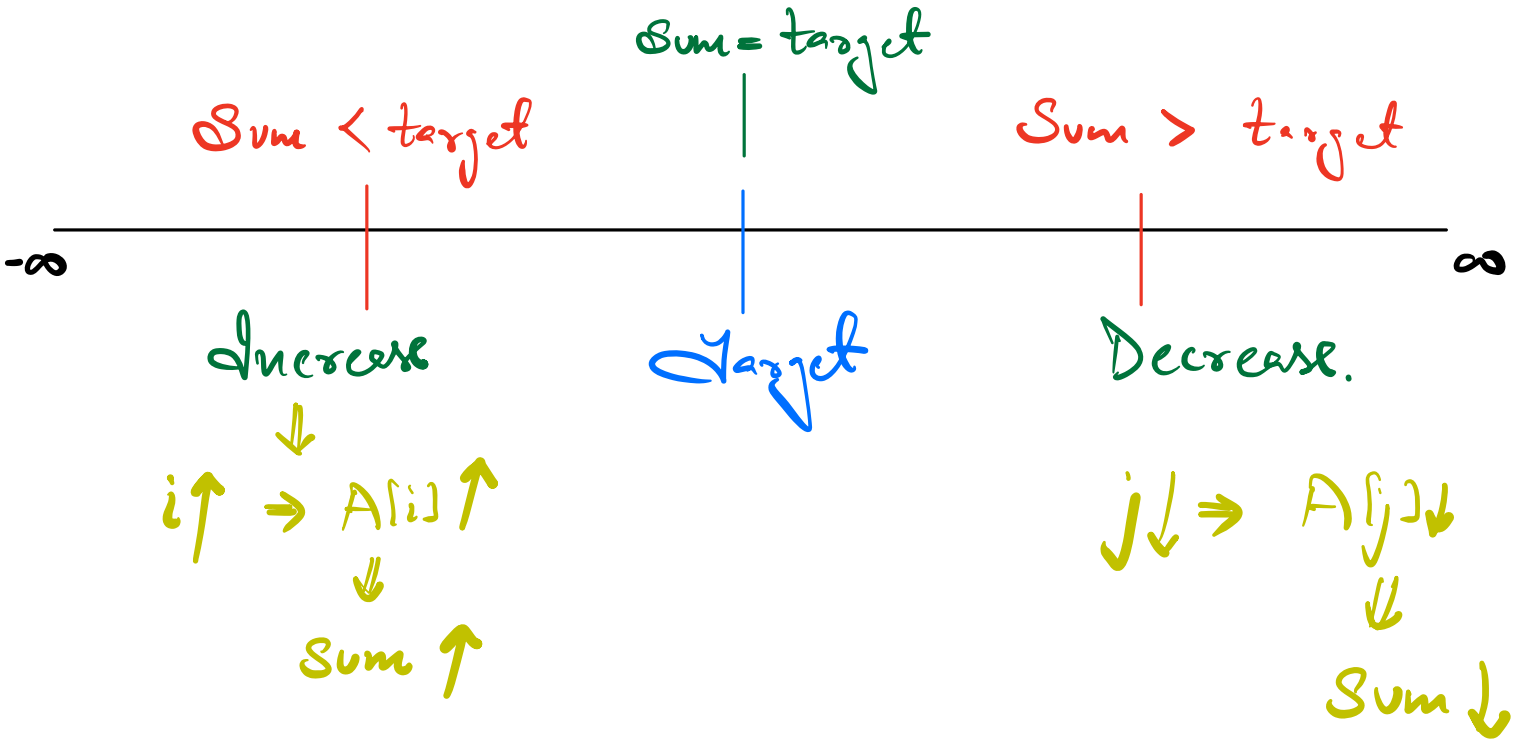
13 > 11 (Discard j)

10 < 11 (Discard i)

12 > 11 (Discard j)

(11 = 11)

Sum = A[i] + A[j]



Code

```
i = 0, j = N-1

while (i < j) {

    Sum = A[i] + A[j];
```

```
if (sum == K) {
        return (i, j);
}
else if (sum > target) {
        j--;
} else {
        i++;
}
}
return <-1, -1>
```

T.C. = O(N)
S.C. = O(1)

Given a sorted integer array A of size N. Find count of pair (i, j) s.t. A[i] + A[j] = K ⇒ Given (i != j)

$A = [1, 2, 3, 4, 5, 6, 8]$ (indices 0 1 2 3 4 5 6)

K = 10

Ans = 2 ((1,6), (3,5))

Case-1 Distinct Elements

Code

```
i=0, j= N-1 , count = 0;
while (i<j) {
    sum = A[i]+A[j];
    if (sum == K) {
        count++,   i++,   j--;
    }
    else if (sum > target) {
        j--;
    } else {
        i++;
    }
}
return count;
```

T.C. = O(N)
S.C. = O(1)

Case - II Duplicates

K = 10

freq i

freq j

A = [1, 1, 1, 2, 3, 3, 3, 3, 4, 5, 5, 5, 5, 7, 7, 7]
indices: 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
i: 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
j: 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15

A[4] + A[13] = 10(K)

(4, 11) (5, 11) (6, 11) (7, 11)
(4, 12) (5, 12) (6, 12) (7, 12) 12
(4, 13) (5, 13) (6, 13) (7, 13)

$$^4C_2 = \frac{4 \times 3}{2} = 6$$

$$+$$

$$18$$

Q) Given a sorted array of size N & an integer ($K > 0$)
find any pair (i, j) s.t. $(A[j] - A[i]) == K$
$(i \neq j)$

$$A = \underset{0}{-5}, \underset{1}{-2}, \underset{2}{1}, \underset{3}{8}, \underset{4}{10}, \underset{5}{12}, \underset{6}{15}$$

$K = 11 \Rightarrow (\underset{i}{2}, \underset{j}{5}) \Rightarrow (12 - 1) = 11$

Ans

1) Brute force → Check all pairs.

T.C. = $O(N^2)$

$(A[j] - A[i]) > 0 \rightarrow A[j] > A[i]$

$K > 0$

$(j \geq i)$

2) Binary Search

```
for (i=0; i<N; i++) {          N
    s = i+1;
    e = N-1;
    target = K + A[i];          } log N
    (B.S.)
}
```

T.C. = $O(N \log N)$

S.C. = $O(1)$

3) Two Pointers

Q1 Where to initialise two pointers

1) $i = 0, \quad j = N-1$ ✗ ⇒ Sum

2) $i = mid, \quad j = mid+1$ ✗

3) $i = 0, \quad j = 1$ ⇒ Diff.

Q2 How to move the pointers.

$$A = [-5, -2, 1, 8, 10, 19, 15]$$

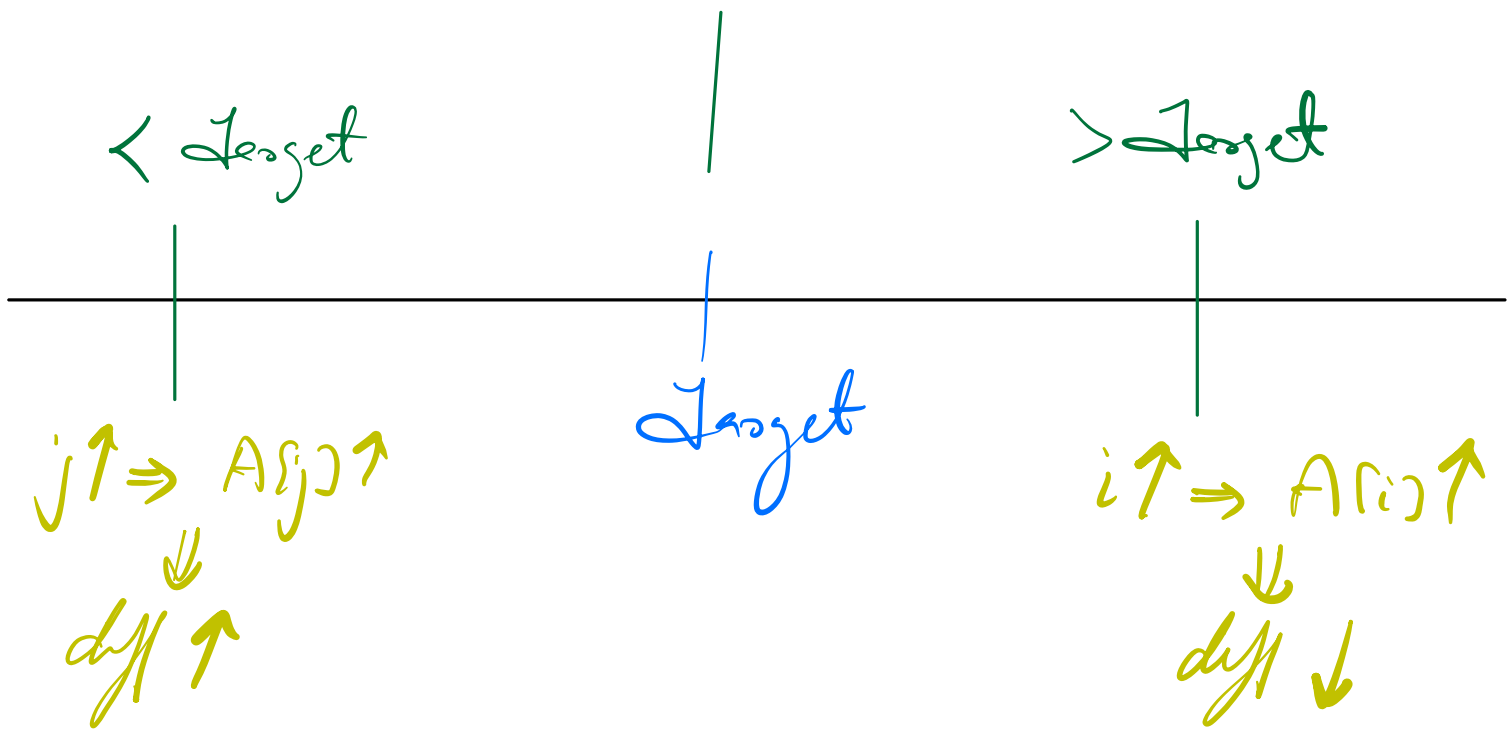
indices: 0 1 2 3 4 5 6

i j

$K = 11$

i	j	diff		
0	6 (N-1)	20	> 11	✗
0	5	17	j--	diff ↓
1	6	17	i++	diff ↓

3 (mid)	4 (mid+1)	2	$<$ target	✗
2	4	9	i--	diff ↑
3	5	4	j++	diff ↑



$$A = [-5, -2, 1, 8, 10, 12, 15]$$

indices: 0 1 2 3 4 5 6

$K = 11$

$diff = A[j] - A[i]$

i	j	diff
0	1	$-2-(-5) = 3 < 11 \Rightarrow j{+}{+}$
0	2	$1-(-5) = 6 < 11 \Rightarrow j{+}{+}$
0	3	$8-(-5) = 13 > 11 \Rightarrow i{+}{+}$
1	3	$8-(-2) = 10 < 11 \Rightarrow j{+}{+}$
1	4	$10-(-2) = 12 > 11 \Rightarrow i{+}{+}$
2	4	$10-1 = 9 < 11 \Rightarrow j{+}{+}$
2	5	$12-1 = 11 = 11$

Code

```
i = 0, j = 1;
while ( j<N  && i<N ){
    diff = A[j]- A[i];
    if (diff==K){
```

```
            return {i, j};
        } else if (diff > k) {
            i++;
        } else {
            j++;
        }
    }
    return {-1, -1};
}
```

T.C. = O(N)
S.C. = O(1)

Given an integer array of +ve elements. Check if a subarray with sum = k is present.

0	1	2	3	4	5	6
1,	3,	15,	10,	20,	3,	23

A = [1, 3, 15, 10, 20, 3, 23]

K = 33 ⇒ True

1) PS + Hashing

2) Optimal

K - 33 ⇒ True

$$A = [\underset{0}{1}, \underset{1}{3}, \underset{2}{15}, \underset{3}{10}, \underset{4}{20}, \underset{5}{3}, \underset{6}{23}]$$

$$PS: = [1, 4, 19, 29, 49, 52, 75] \Rightarrow \text{Sorted}$$

$$Sum[i,j] == K$$

$$\Downarrow$$

$$PS[j] - PS[i-1] == K \quad (\forall\ i != 0)$$

$$(PS[j] == K) \quad (\forall\ i == 0)$$ loop

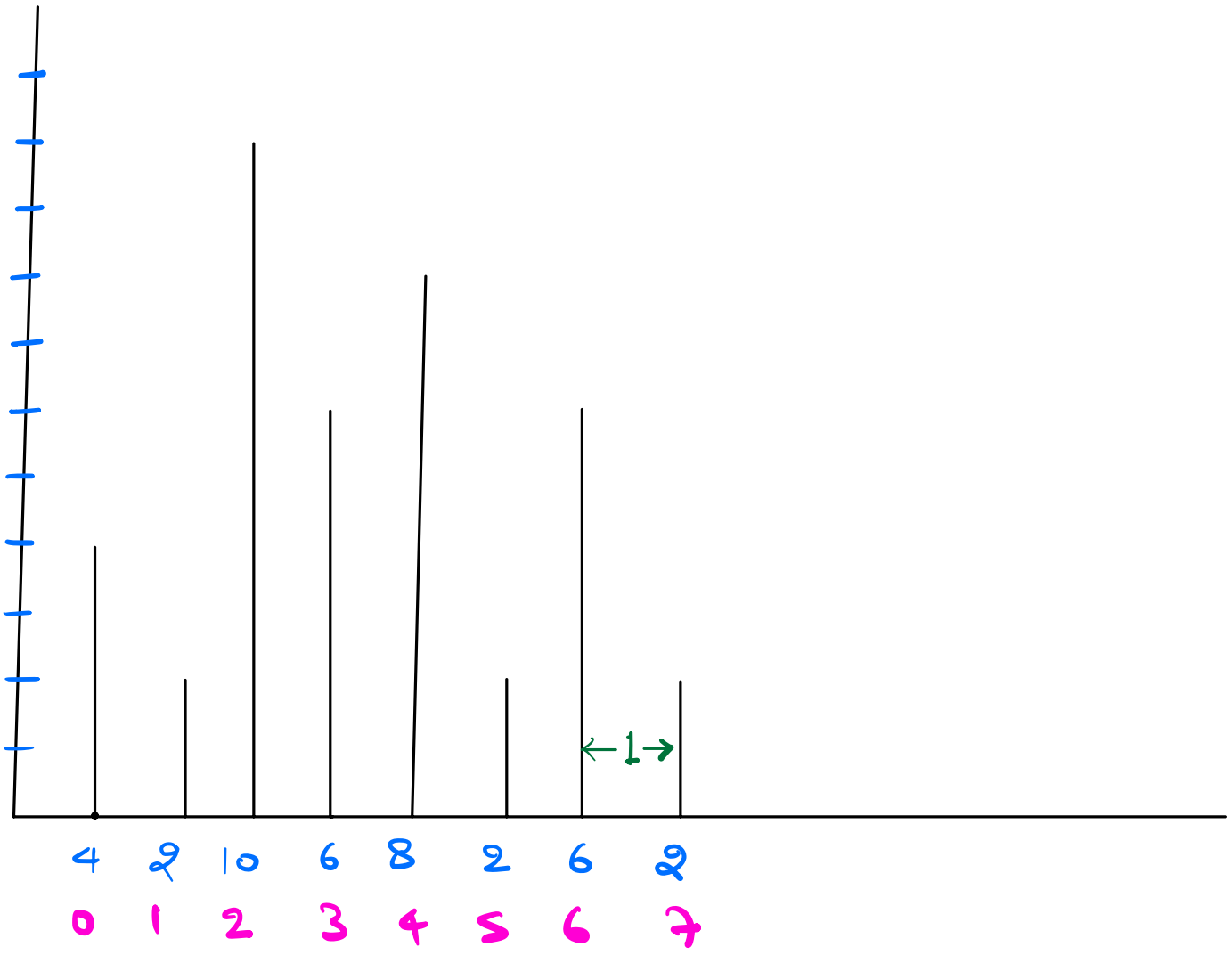
Q) Given an integer array of size N. representing the height of N walls.

find any 2 walls that can form a container to store max amount of water

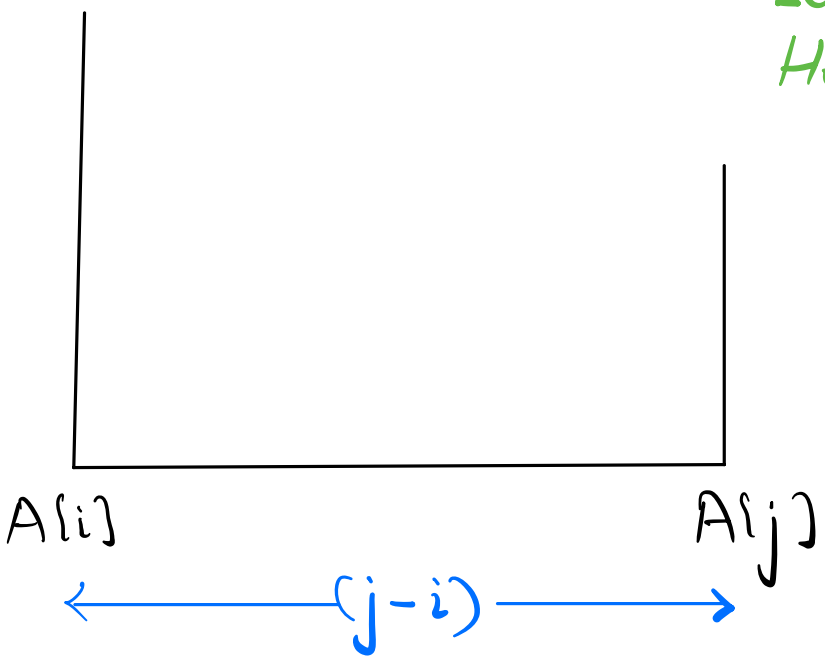
1) Volume = Area ←

2) The distance b/w every consecutive wall is 1 unit

A = [4, 2, 10, 6, 8, 2, 6, 2]



Length = (j - i)
Height = min(A[i], A[j])



$$\text{Amount of water} = (j-i) \times \min(A[i], A[j])$$

1) Brute Force

$\forall i, j$ where $i < j \Rightarrow$ Calculate the area & take max.

$$T.C. = O(N^2)$$
$$S.C. = O(1)$$

2) Optimal Solution

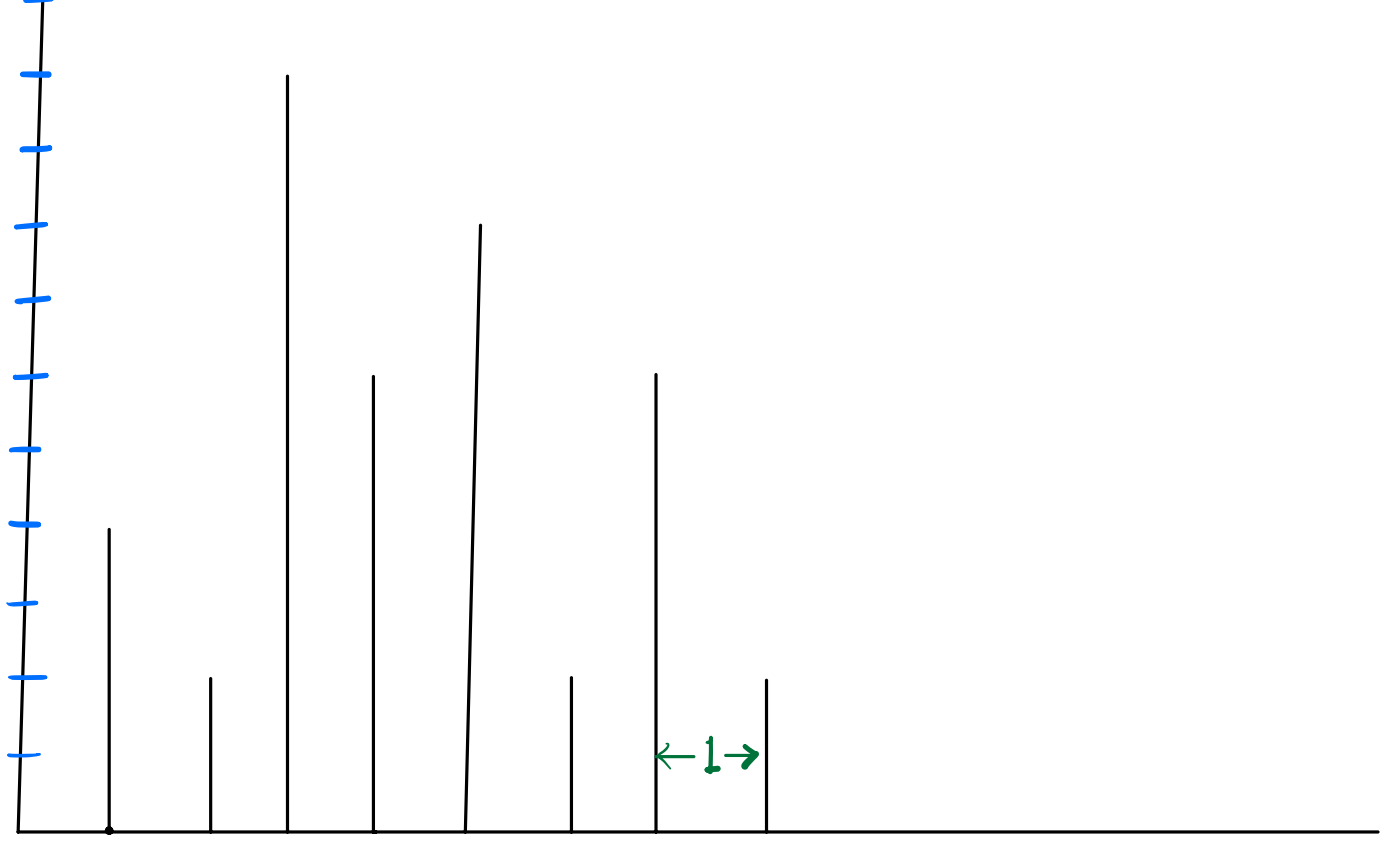
$$\text{Amount of water} = \underbrace{(j-i)}_{\text{Length}} \times \underbrace{\min(A[i], A[j])}_{\text{Height}}$$

Goal : Maximise amount of water.

Since we need to maximize area.
We start with max length.

$$A = [4, 2, 10, 6, 8, 2, 6, 2]$$

$\uparrow$



i	j	Length	Height	Amount
0	7	7	min(4,2) = 2(j)	14

move / discard smaller wall

Code

```
i = 0,   j = n-1;
ans = INT_MIN
while (i<j) {
```

```
area = min(A[i], A[j]) * (j - i);
ans = max(area, ans);

if (A[i] < A[j]) {
    i++;
} else if (A[j] < A[i]) {
    j++;
} else {
    i++;
    j++;
}
}
return ans;
```