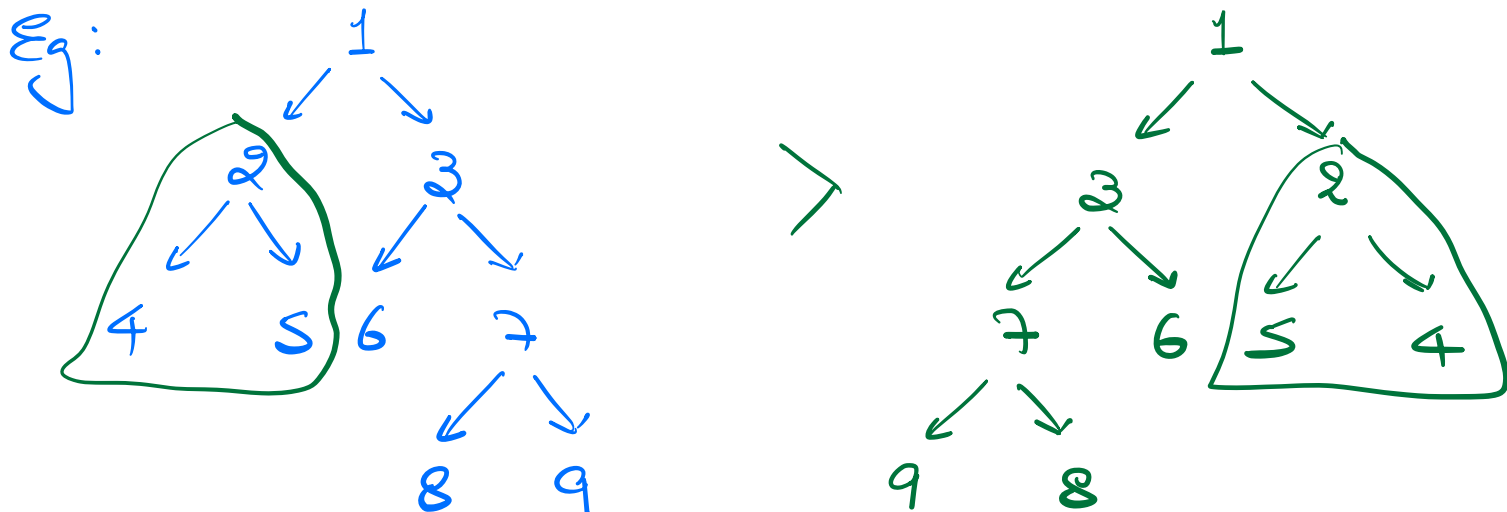


Agenda

- 1) Invert Binary Tree
- 2) Equal Tree Partition
- 3) Next pointer in Binary Tree
- 4) Root to leaf path sum = K
- 5) Diameter of Binary Tree.

Q Given the root node of a Binary Tree, write a function to invert the tree



Inversion

Steps

- 1) Invert the LST
- 2) Invert the RST
- 3) Make left as right & vice versa.

Code

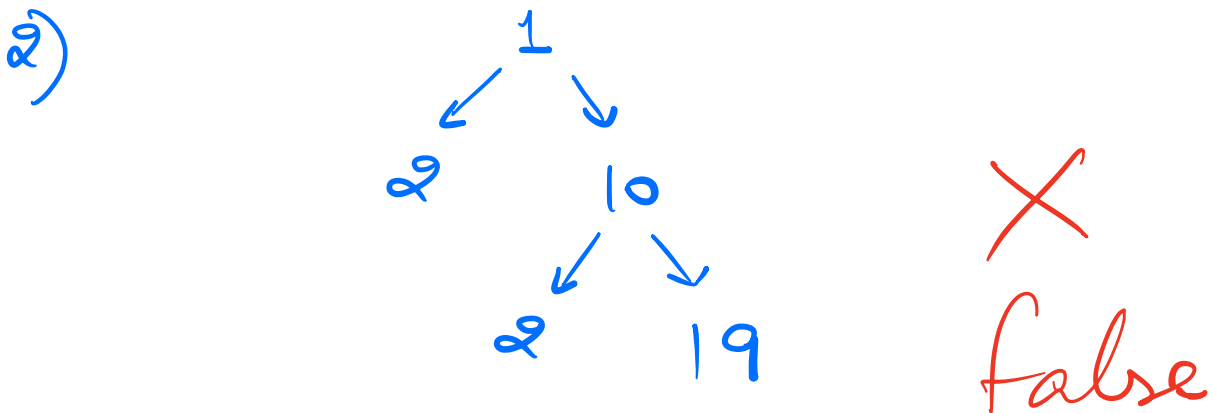
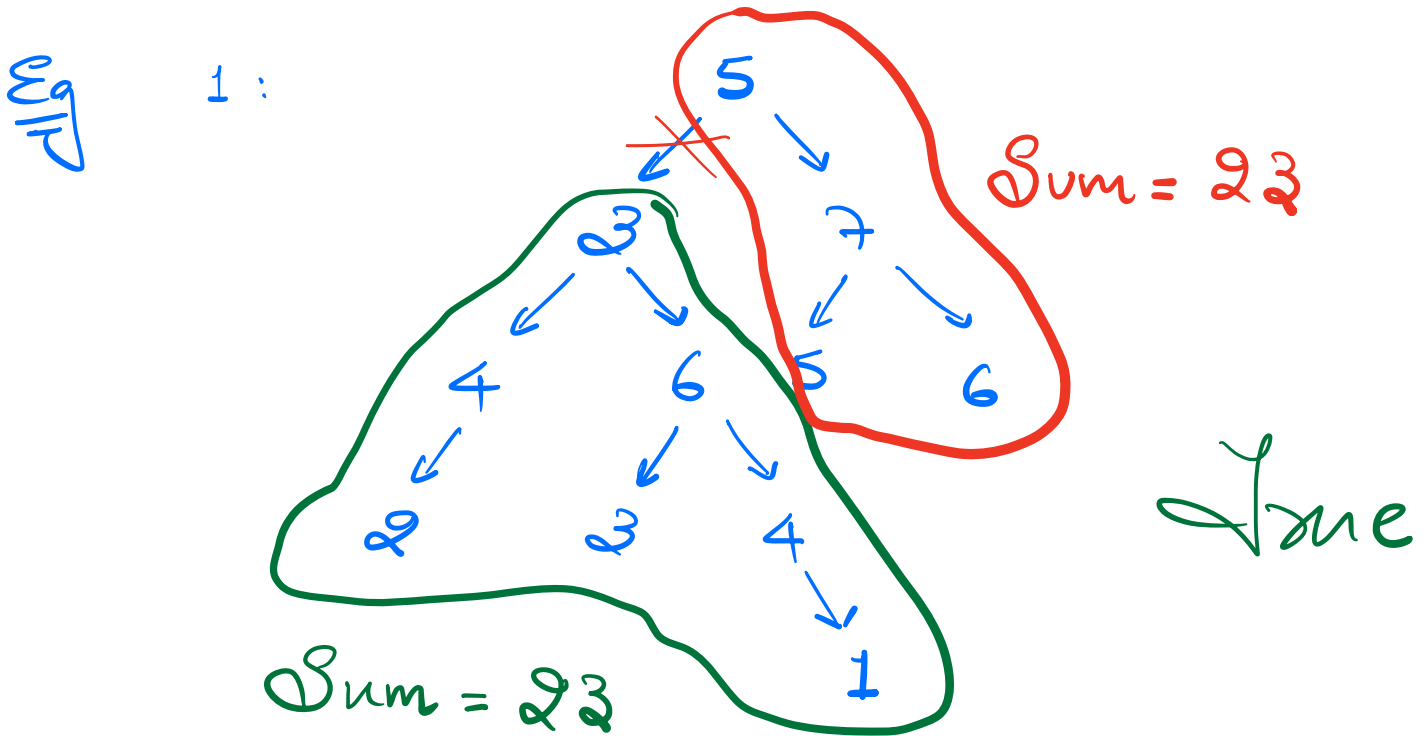
```
void invertTree( Node root) {  
    if (root == NULL) { return; }
```

```
    Node temp = root.left;  
    root.left = root.right;  
    root.right = temp;
```

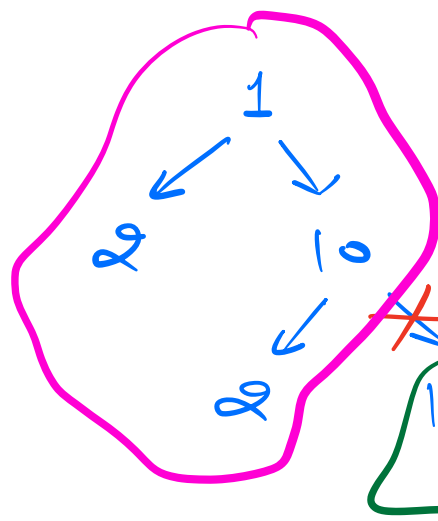
```
    invertTree( root.left);  
    invertTree( root.right);  
}
```

T.C. = $O(N)$
S.C. = $O(\text{Height})$;

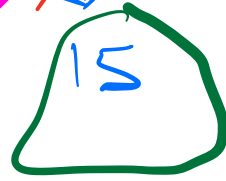
Q2 Given the root node of a Binary Tree,
Return True if the tree can be split into two non-empty subtrees with equal sums. False otherwise.



3)



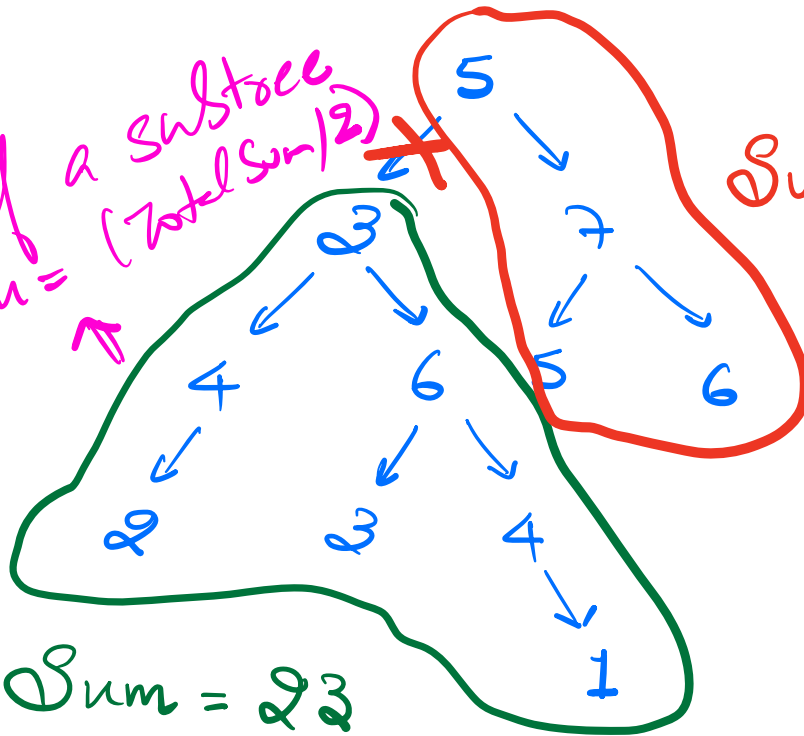
Sum = 15



Sum = 15

Soln

find if a subtree with sum = (Total Sum / 2) exists.



Sum = 23

Sum = 23

True

Total Sum = 46

Steps

1) Calculate Total Sum. (Return false if Odd)

2) \forall subtree \Rightarrow Calculate Sum & check if it is equal to half of Total Sum.

Code

total Sum (Calculate it first)
ans = false;

```
int sum (Node root) {  
    if (root == NULL) return 0; }
```

```
    int sum_left = sum (root.left);  
    int sum_right = sum (root.right);
```

```
    if (sum_left == totalSum/2 ||  
        sum_right == totalSum/2) {
```

```
        ans = true;
```

```
    }
```

```
    return (sum_left + sum_right +  
            root.data);
```

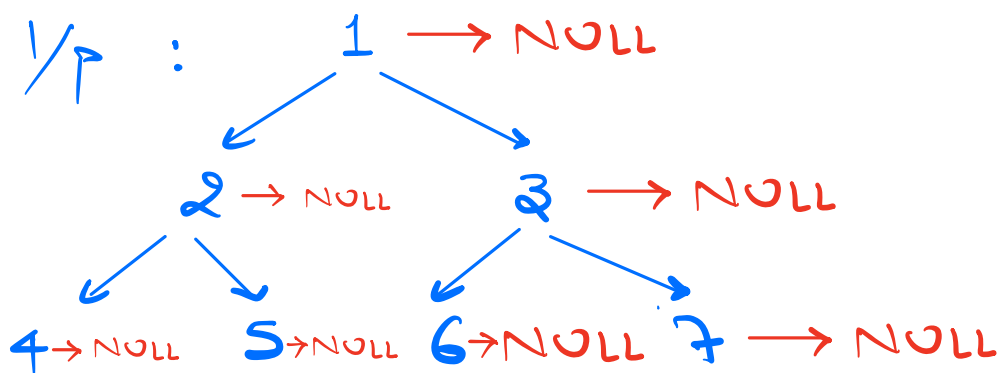
```
}
```

T.C. = $O(N)$

S.C. = $O(\text{Height})$

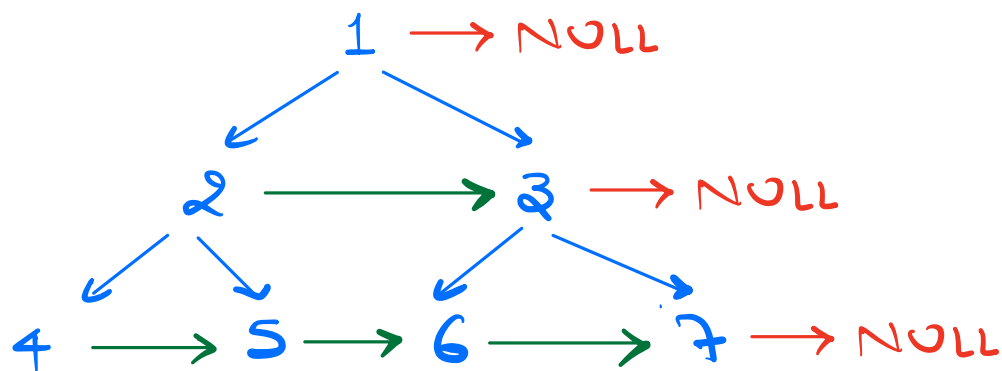
Q3 Given a perfect binary tree with all next pointers set to NULL,

Modify the tree in-place (S.C. = O(1)) to the next node in the same level from left to right.



Node &
int data
Node left
Node right
Node next;
&

O/p :



Solⁿ

> Brute force \Rightarrow level Order Traversal.

```
Node connect (Node root) &  
if (root == NULL) &  
    return NULL;
```

```
Queue <Node> ,
```

```
q.enqueue (root);
```

```
while (! q.isEmpty()) &
```

```
    int levelSize = q.size();
```

```
    for (i = 0; i < levelSize; i++) &
```

```
        Node temp = q.dequeue();
```

```
        if (i < (levelSize-1)) &  
            temp.next = q.front;
```

```
    }
```

```
    if (temp.left != NULL) &  
        q.enqueue (temp.left);
```

```
    }
```

```
    if (temp.right != NULL) &  
        q.enqueue (temp.right);
```

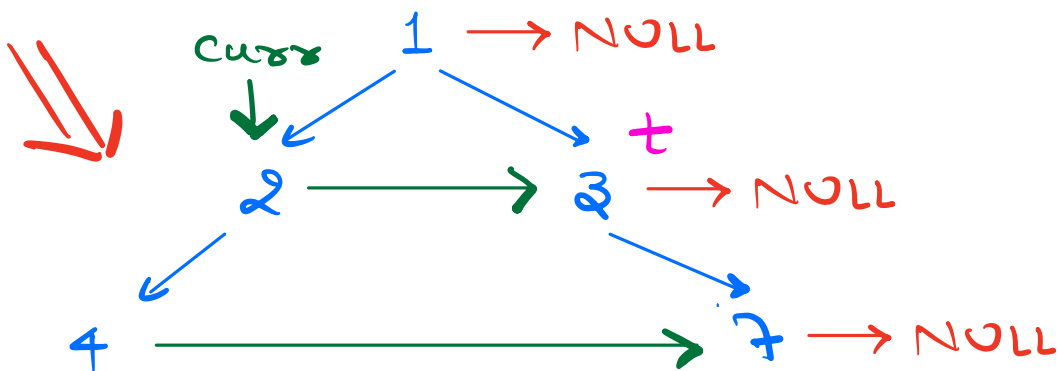
```
    }
```

return root;

$$T.C. = O(N)$$

$$S.C. = \underline{O(N)}$$

H.W.



Code

curr = root;

while (curr != NULL) {

temp = curr;

while (temp != NULL) {

temp.left.next = temp.right;

if (temp.next != NULL) {

temp.right.next = temp.next.left;

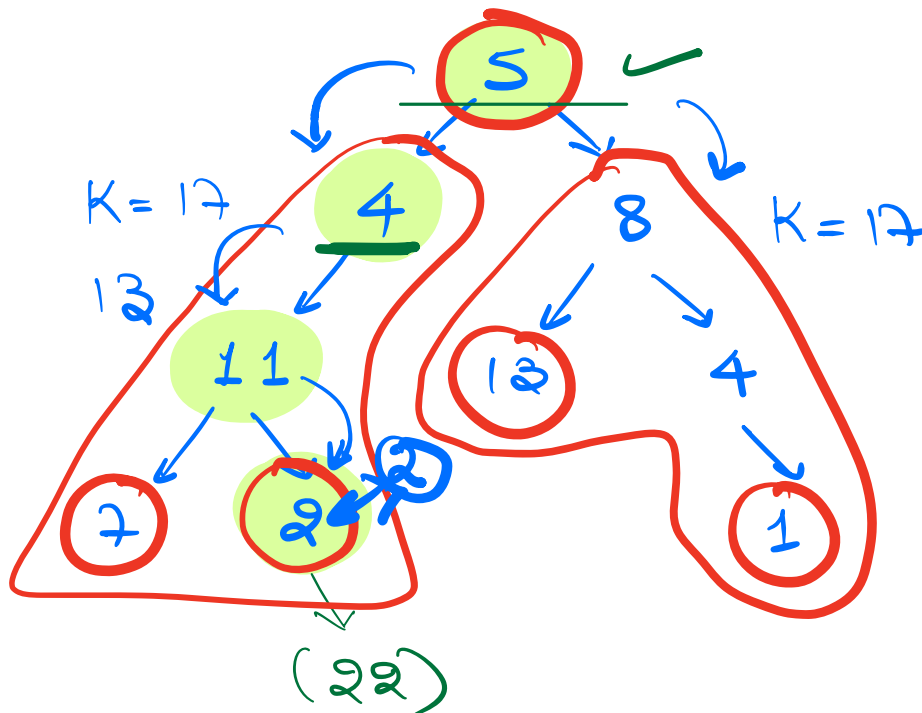
temp = temp.next;

curr = curr->left;

$$T.C. = O(N)$$
$$S.C. = O(1)$$

Q4 Given a Binary Tree. Check if there exists a root to leaf path in the tree such that adding all the node values along the path has sum = K. (Given).

I/P :



$$K = 22$$

True

Solⁿ

Code

boolean hasPathSum (root, K) {

if (root == NULL) {
return false;

}

if (root.left == NULL && root.right == NULL) {

return (K == root.data);

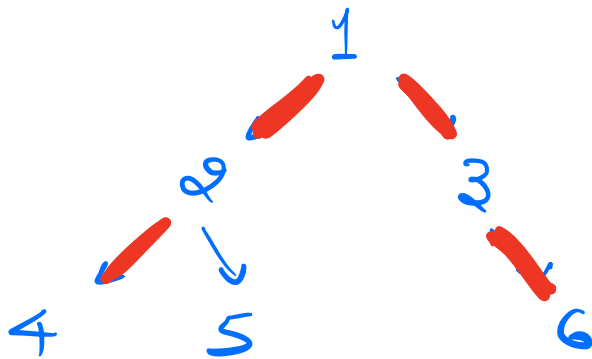
}

return hasPathSum(root.left, K - root.data) ||
hasPathSum(root.right, K - root.data);

}

Q5 Given a Binary Tree. Find the length of the longest path between any two nodes. (diameter)

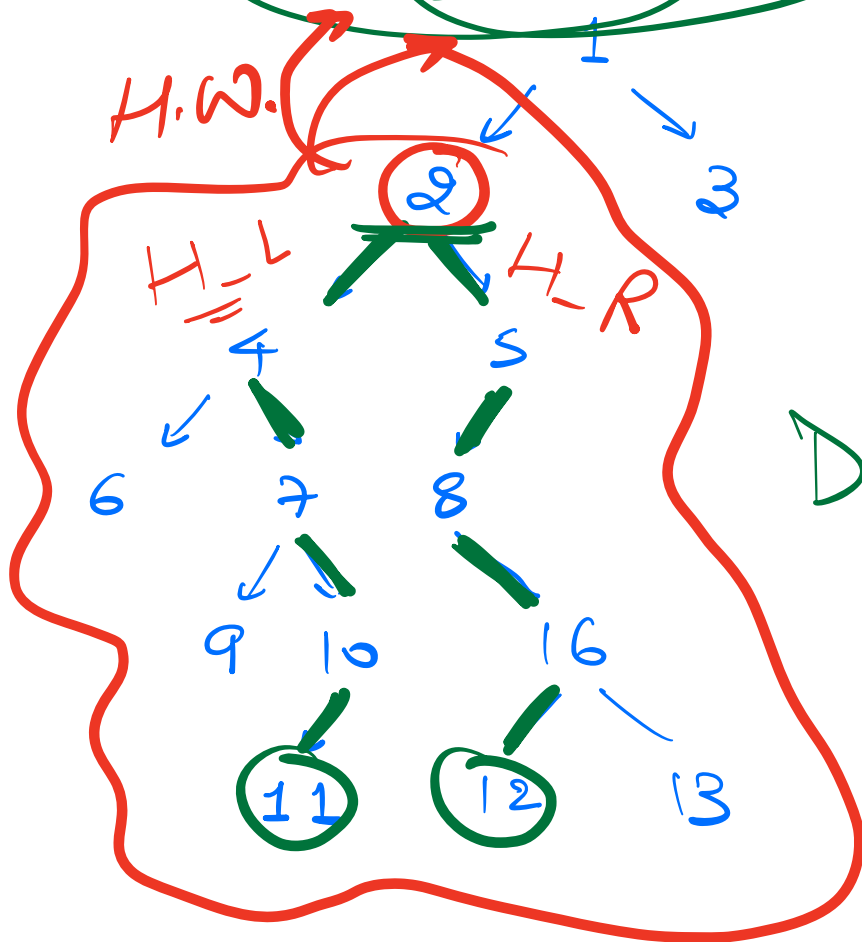
I/P :



Diameter = 4



I/P



Post Order

Diameter = 8

MAANG & Trac 1 \Rightarrow

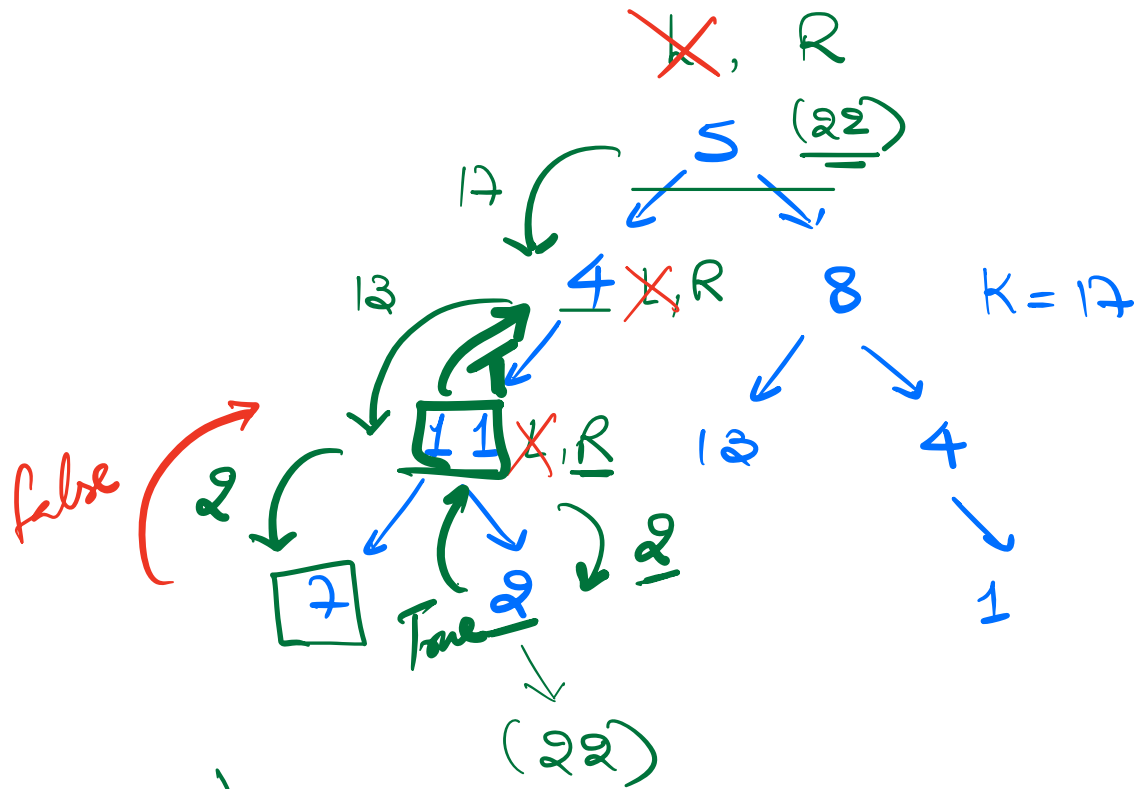
Difficult
Beatschky
methe
Graphs

DSA 4.2

LL
Stacks & Queues
Graphs
Binary Search
Trees
Heaps.
Greedy

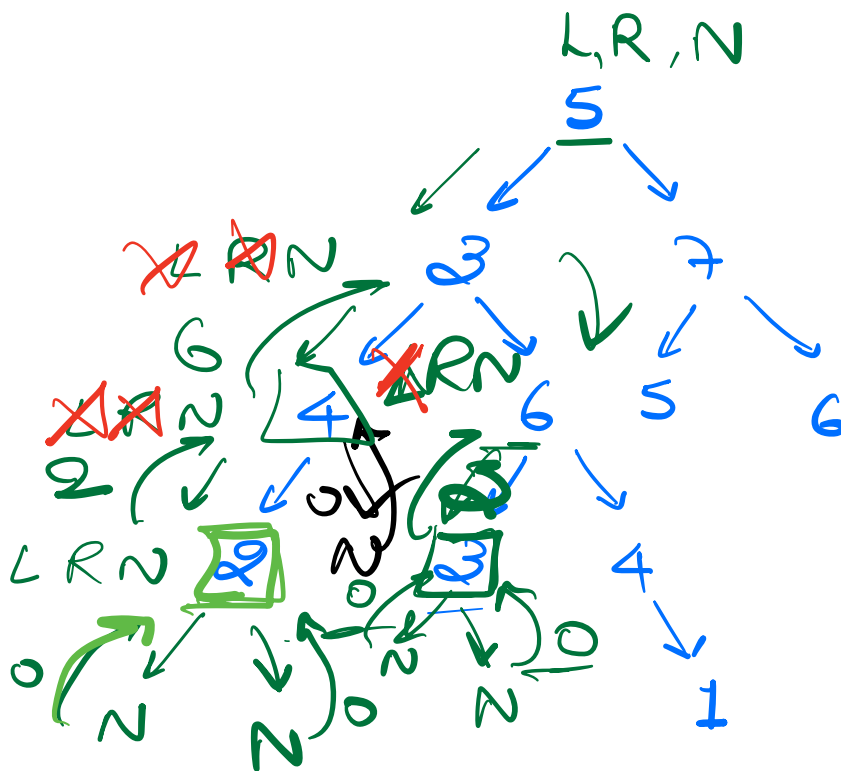
Practice

\Rightarrow Time Restriction



$\underline{K=22}$
 $\underline{\underline{22}}$

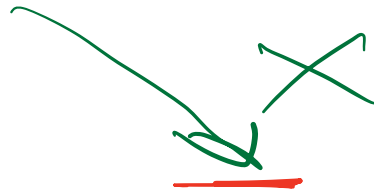
~~AN~~



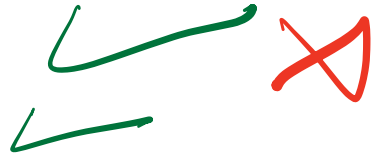
$\underline{46}$
 $\underline{\underline{46}}$

Revision →

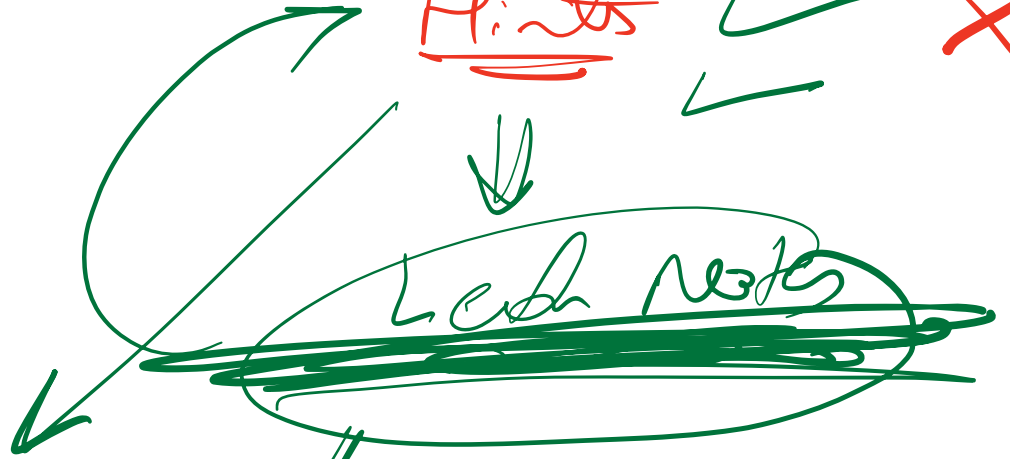
1) Solve new Questions X



Hints



Leads Notes



DSA

DSA 4.2



Reword