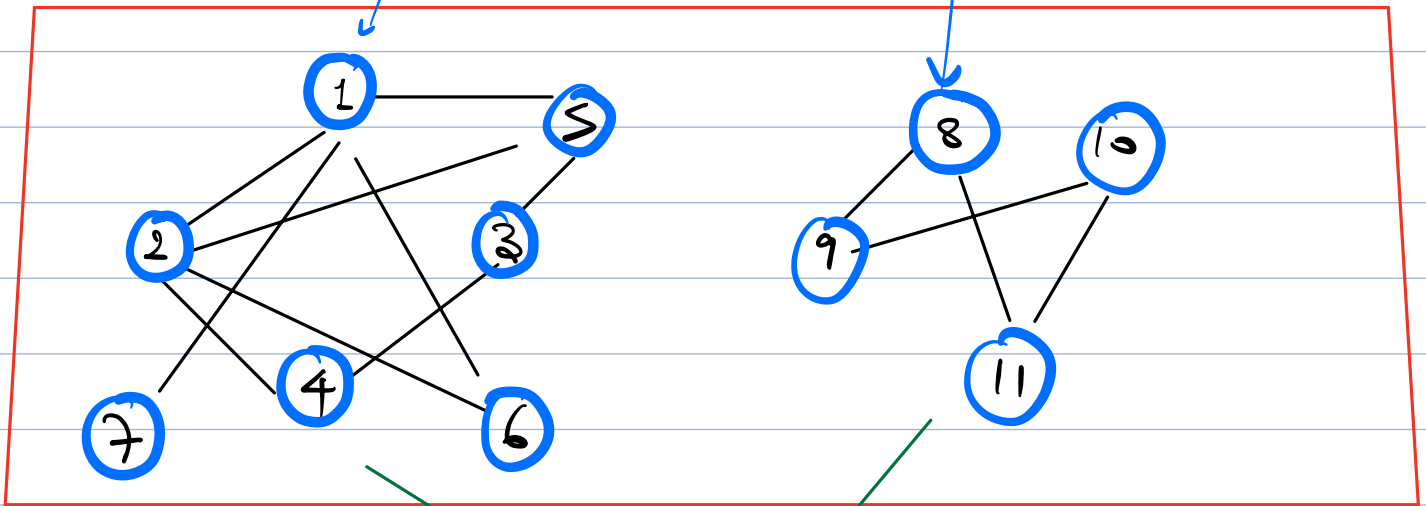


Connected Components

PS Session \rightarrow DP
Graphs
Tomorrow 4PM

Facebook

dfs



Connected Components (cc)

Q Find the no. of CC present in a graph with \checkmark nodes numbered from 1 to n & E edges. (Undirected)

Code

1) Create adjacency list.

visited [$\checkmark+1$] = {false}

count = 0;

for (i = 1; i ≤ ~~N~~; i++) {

if (visited[i] == false) {

count++;

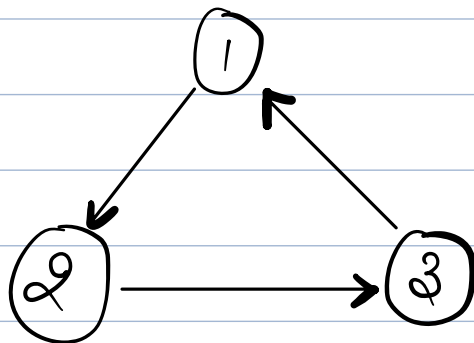
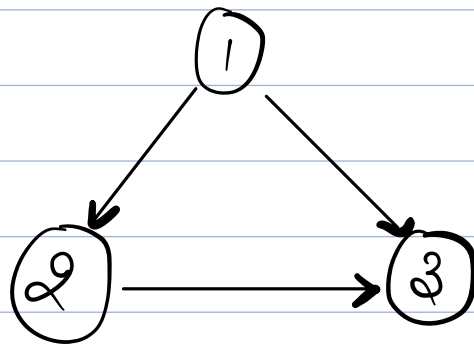
dfs(i);

}

}

T.C. = $O(N + E)$

Strongly Connected Components.



Strongly
Connected
Components

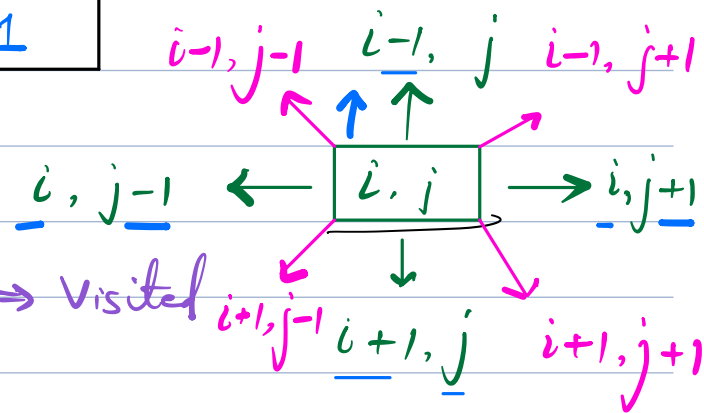
~~Ans~~ No. of Islands.

1	1	0	0	1
0	1	0	1	0
1	0	0	1	1
1	1	0	0	0
1	0	1	1	1

$N \times M$

1 \rightarrow Islands
0 \rightarrow Water

$i, j = 0, 1 \Rightarrow$ Visited
 $\hookrightarrow 2$



Q Count the no. of islands given.

```
count = 0;
for (i = 0; i < N; i++) {
    for (j = 0; j < M; j++) {
        if (M[i][j] == 1) {
            count++;
            dfs(i, j);
        }
    }
}
```

Row & Col Modifiers.

row: $[-1, \underline{0}, \underline{1}, \underline{0}, 1, -1, 1, -1]$
col: $[0, \underline{1}, \underline{0}, \underline{-1}, 1, -1, -1, 1]$

void dfs (i, j) {

$M[i][j] = 2;$ // Mark as Visited.

for ($k=0$; $k < \text{row.size}()$; $k++$) {

$\text{row_index} = i + \text{row}[k];$
 $\text{col_index} = j + \text{col}[k];$

if (check(row_index, col_index) == True) {

if ($M[\text{row_index}][\text{col_index}] == 1$

dfs(row_index, col_index);

}

}

}

}

bool check (i, j) {

if (i < 0 || j < 0 || i >= N || j >= N) {

return false;

}

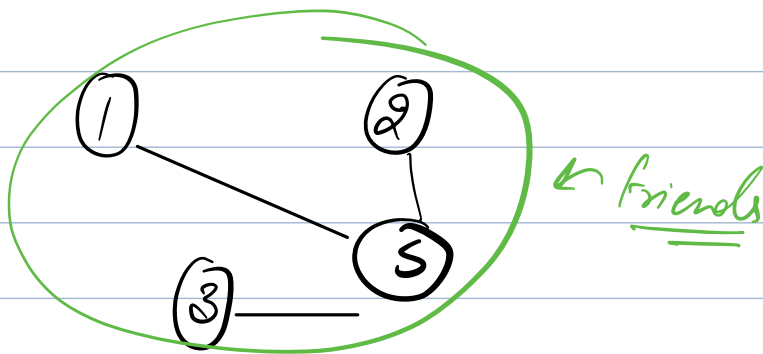
return True;

}

T.C = $O(N^2)$

Audible → Prime Users (Force Audible Subscription)

Bonus ↓ Refer & Join



Count the no. of such groups

2+(-1), 3+0

2+0, 3+(-1) ← 2, 3 → 2+0, 3+1

	0	1	2	3	4
0	1	1	0	0	1
1	0	1	0	1	0
2	1	0	1	1	1

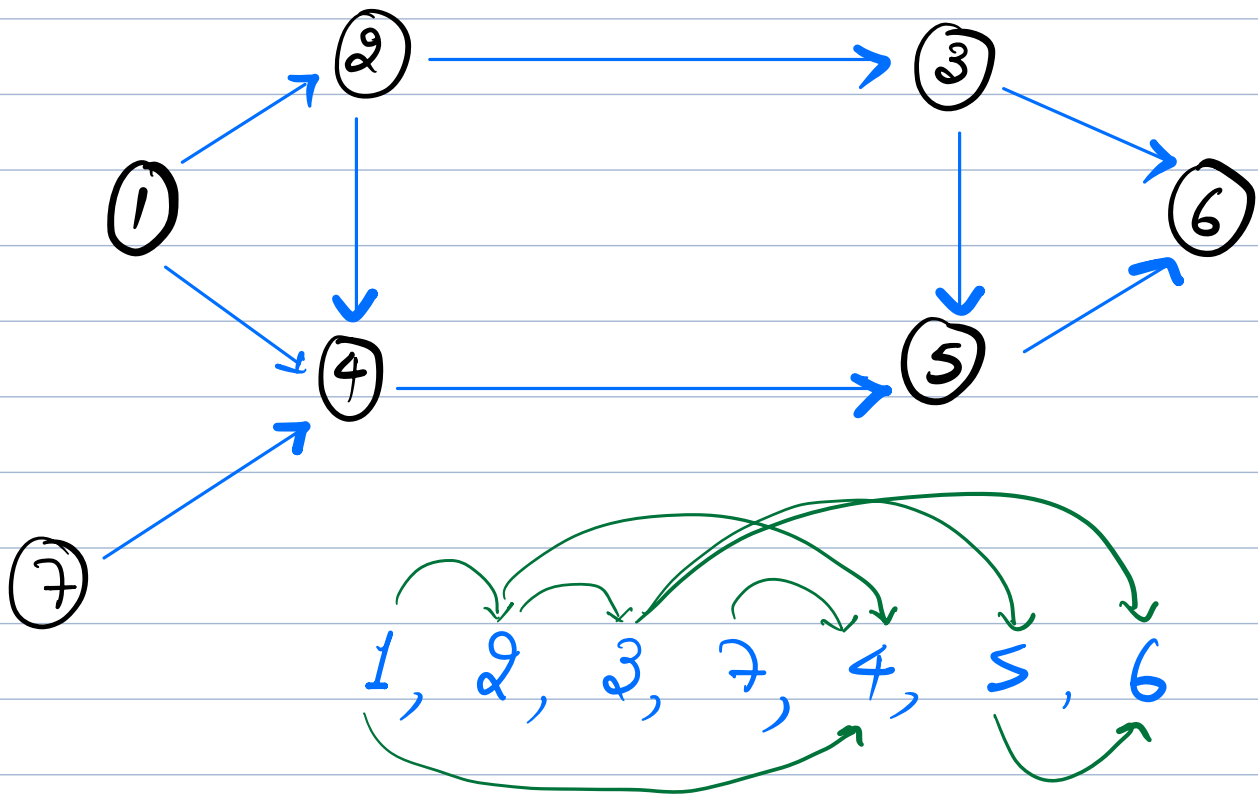
2+1, 3+0

3	1	1	0	0	0
4	1	0	1	1	1

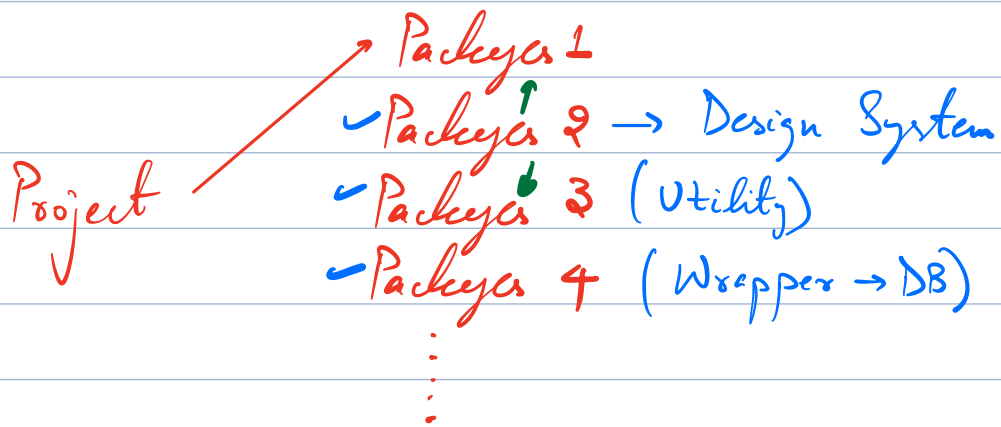
Topological Sort

Only for Directed Acyclic Graph.

linear ordering of vertices (Nodes) such that for every directed edge (u, v) , vertex u comes before vertex v .



1, 7, 2, 4, 3, 5, 6

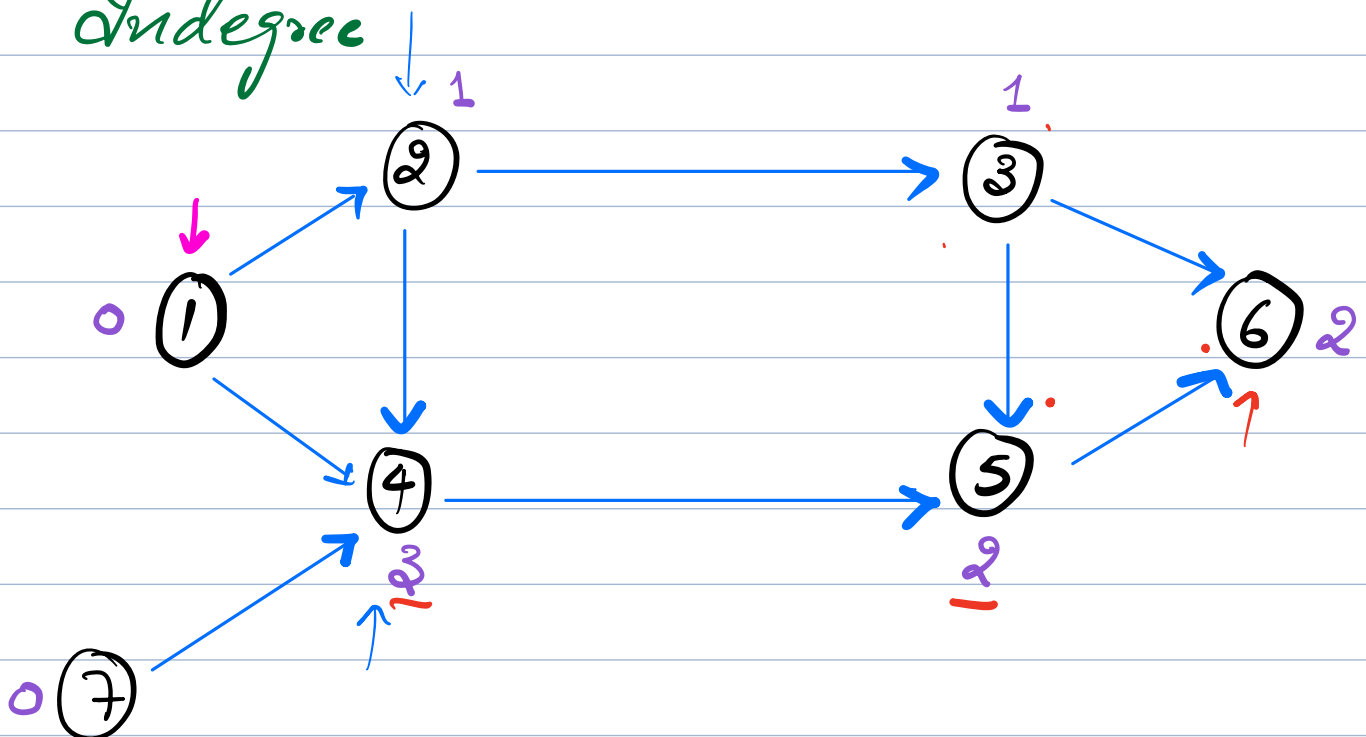


Error ⇒ Circular dependency X Not Cycle

Order in which packages needs to be built ⇒

! Find Topological Sorted Order

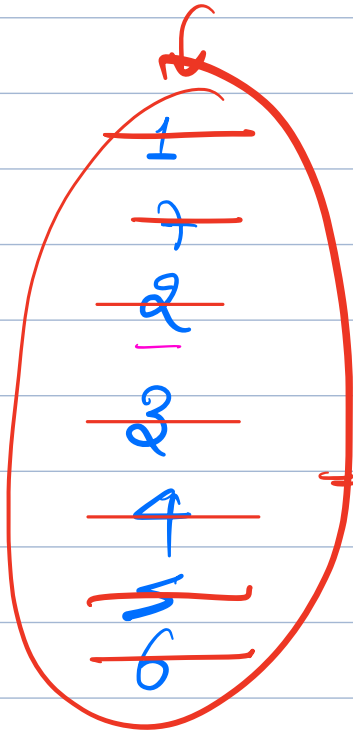
1) Indegree



0	1	2	3	4	5	6	7
X	0	1	1	2	2	2	0

0 0 2 1 1
~~10~~ 0 0

1, 2, 3, 7, 4, 5, 6



⇒ Hashset

$\text{indegree}[N+1] = \{0\}$

V+E for ($i=1; i \leq N; i++$) {

(Adjacency List)

for (all nodes u connected to i) {

$\text{indegree}[u]++;$

}

}

Hashset $\langle \text{int} \rangle$ zeroDegree;

for ($i=1; i \leq N; i++$) {


```
if (indegree[i] == 0) {  
    zerodegrees.add(i);  
}
```

```
}
```

```
while (zerodegrees.size() != 0) {
```

```
    int node = zerodegrees.remove();  
    print(node);
```

```
    for (all u connected to node) {
```

```
        indegree[u]--;
```

```
        if (indegree[u] == 0) {
```

```
            zerodegrees.add(u);  
        }
```

```
}
```

```
}
```

$V + E$

$$T.C = O(V + E)$$

H.W.

Outdegree ??



↓
No need to create outdegree array.
↓
How to use Outdegree

DSU (Disjoint Set Union)

1, 3, 5
2, 4, 8

Disjoint Sets

1
4
7

S1

2
5

S2

6
8

S3

Given N elements \rightarrow N different sets
no. from 1 to N .

1

2

3

4

5

Queries $\rightarrow (u, v) \Rightarrow$ If u, v belongs to different sets, merge (union) the 2 sets & return True. False otherwise

$\langle 1, 2 \rangle \Rightarrow$ True.

1

3

4

5

$\langle 4, 5 \rangle \Rightarrow$ True.

1

3

4

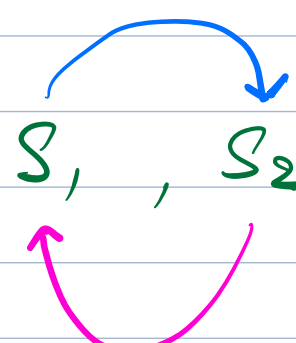
$\langle 2, 3 \rangle \Rightarrow$ True

1


4

$\langle 1, 3 \rangle \Rightarrow \underline{\underline{\text{false}}}$

Merge $\Rightarrow S_1, S_2$

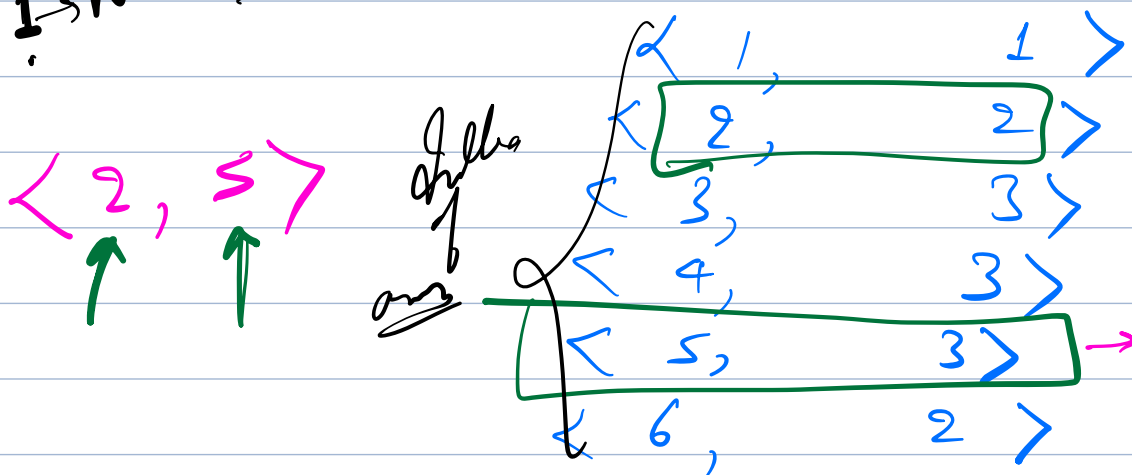


$S_1, S_2, S_3, \dots, S_n$



X \rightarrow

(Array $1 \rightarrow n$) Map $\Rightarrow \langle \text{Node}, \text{SetNo.} \rangle$



Set \Rightarrow Tree / Elements \Rightarrow Nodes

Identify \Rightarrow Root of
Tree

~~Node of
int data
Node present;~~

parent = [~~0~~ 4 1 1 4 4]

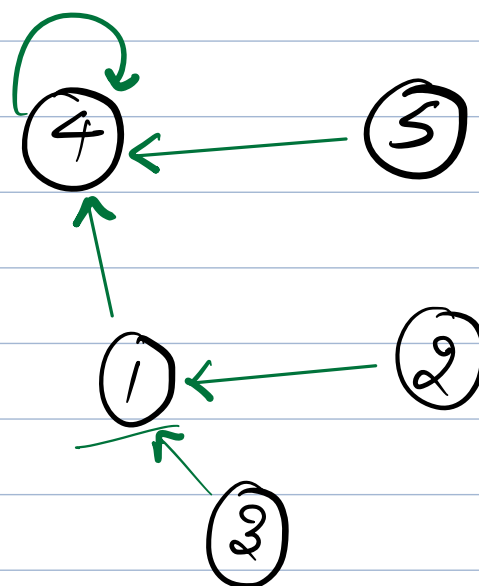
$\langle 1, 2 \rangle$

$\langle 2, 3 \rangle$

$\langle 4, 5 \rangle$

$\langle 3, 5 \rangle$

1) $\langle 1, 2 \rangle$



parent[1] = 2

parent[2] = 1

2) $\langle 2, 3 \rangle$

root(2) = 1

root(3) = 2

$$\text{parent}[1] = 3$$

$$\text{parent}[3] = 1 \quad \checkmark$$

3) $\langle 4, 5 \rangle$

$$\begin{aligned} \text{root}(4) &= 4 \\ \text{root}(5) &= 5 \end{aligned} \quad \}$$

$$\text{parent}[4] = 5$$

$$\text{parent}[5] = 4 \quad \checkmark$$

4) $\langle 3, 5 \rangle$

$$\begin{aligned} \text{root}(3) &= 1 \\ \text{root}(5) &= 4 \end{aligned}$$

$$\text{parent}[2] = 4 \quad \checkmark$$

$$\text{parent}[4] = 1$$

```
int root(u) {
```

```
    while (parent[u] != u) {
```

```
        u = parent[u];
```

```
    }
```

```
    return u;
```

```
}
```

T.P. = $O(\text{Height})$



$O(N)$

```
bool union(u, v) {
```

```
    x = root(u);
```

```
    y = root(v);
```

```
    if (x == y) {
```

```
        return false;
```

```
    }
```

```
    parent[x] = y;
```

```
    return true;
```

```
}
```

$O(N)$

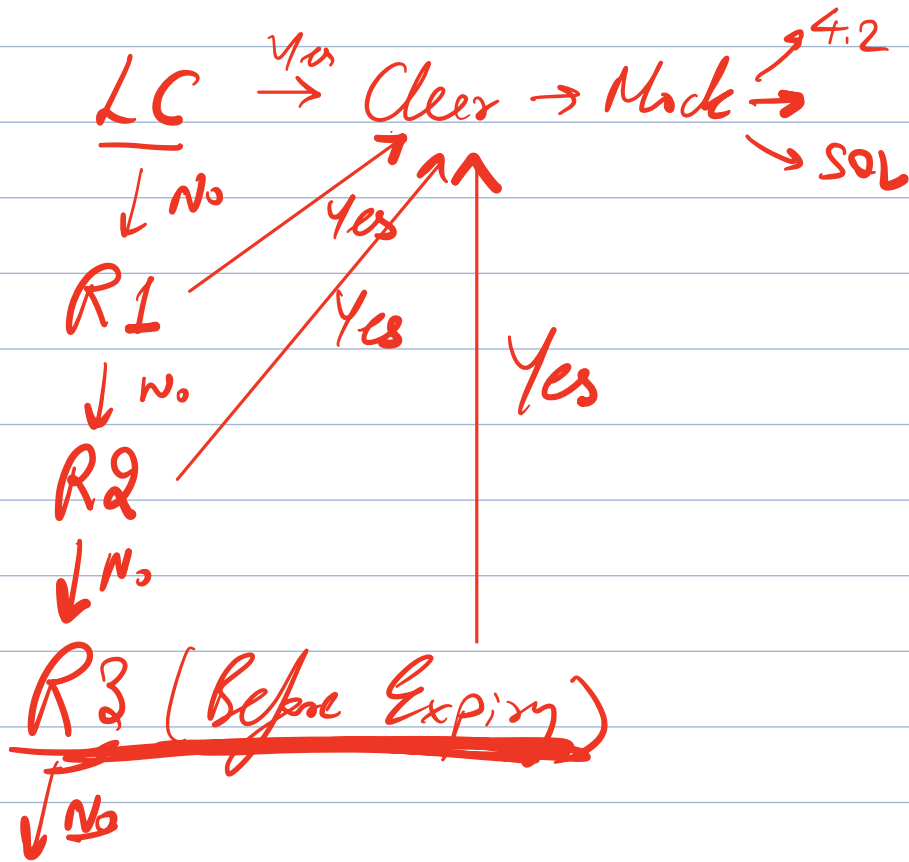


1) Optimise

2) Application in Graph



4PM PS session



↓
Jokes (DP)

↓
Doubts