

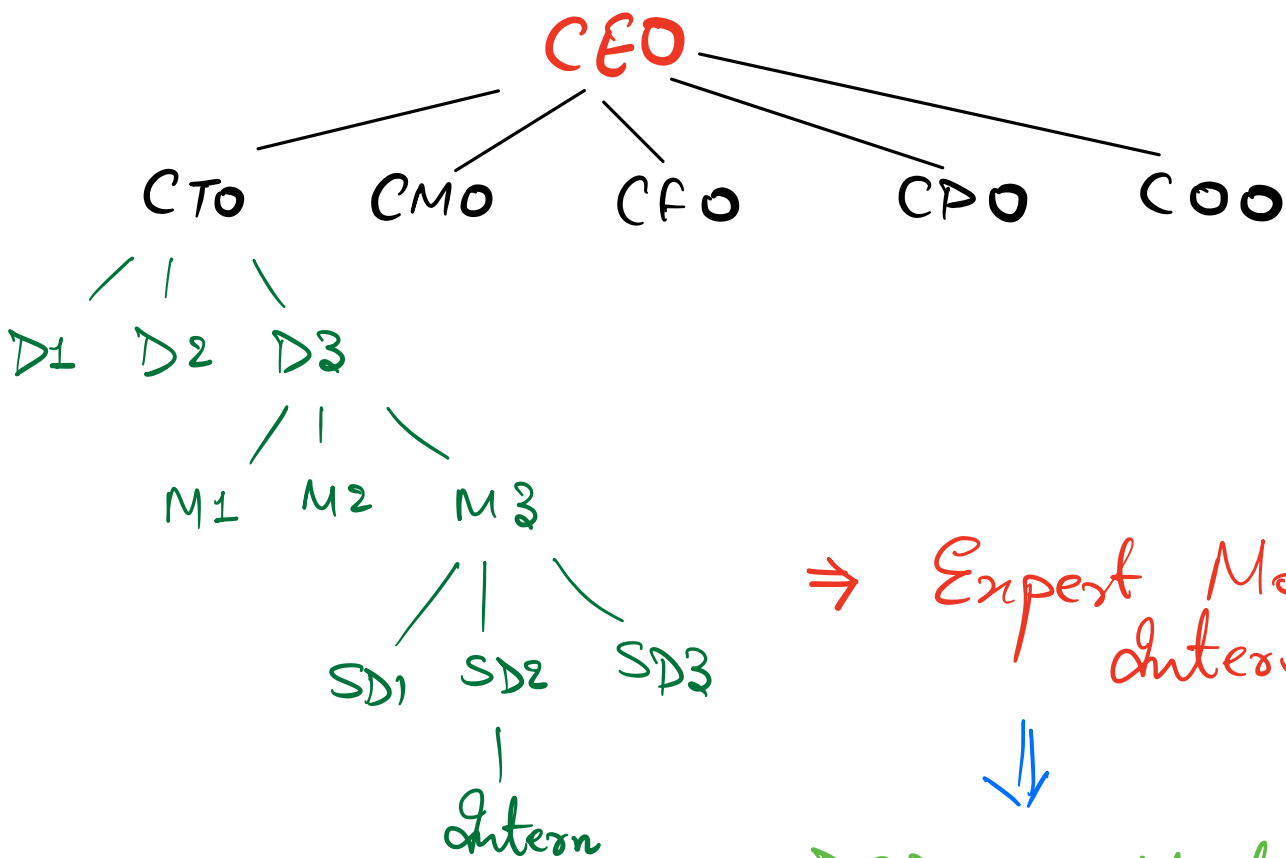
Ayush Sharma \Rightarrow Scaler \Rightarrow Lead DSA

\Downarrow
2019 \Rightarrow CSE, IIT Patna

\Downarrow
Strand life Sciences

\Downarrow
Scaler (2021)

Trees \Rightarrow Hierarchical Data



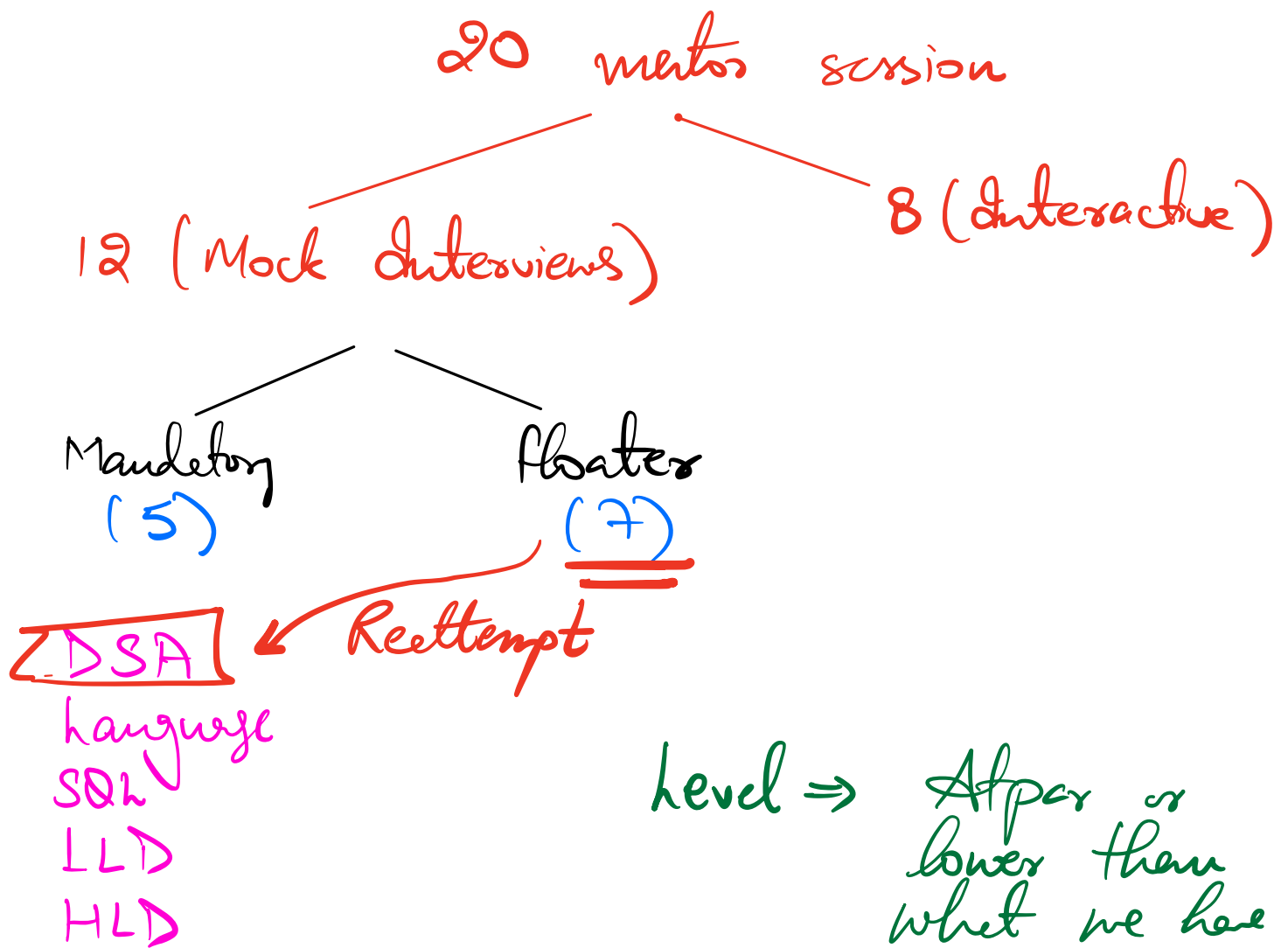
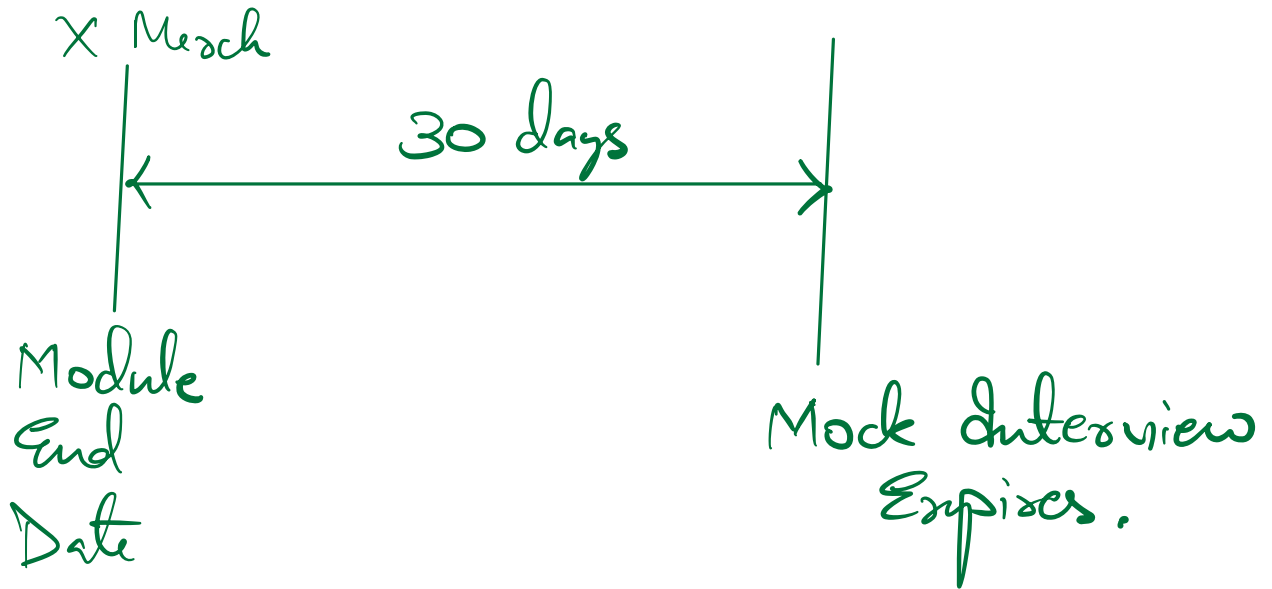
\Rightarrow Expert Mock interview.



DSA \Rightarrow Mock interview



100% Companies.

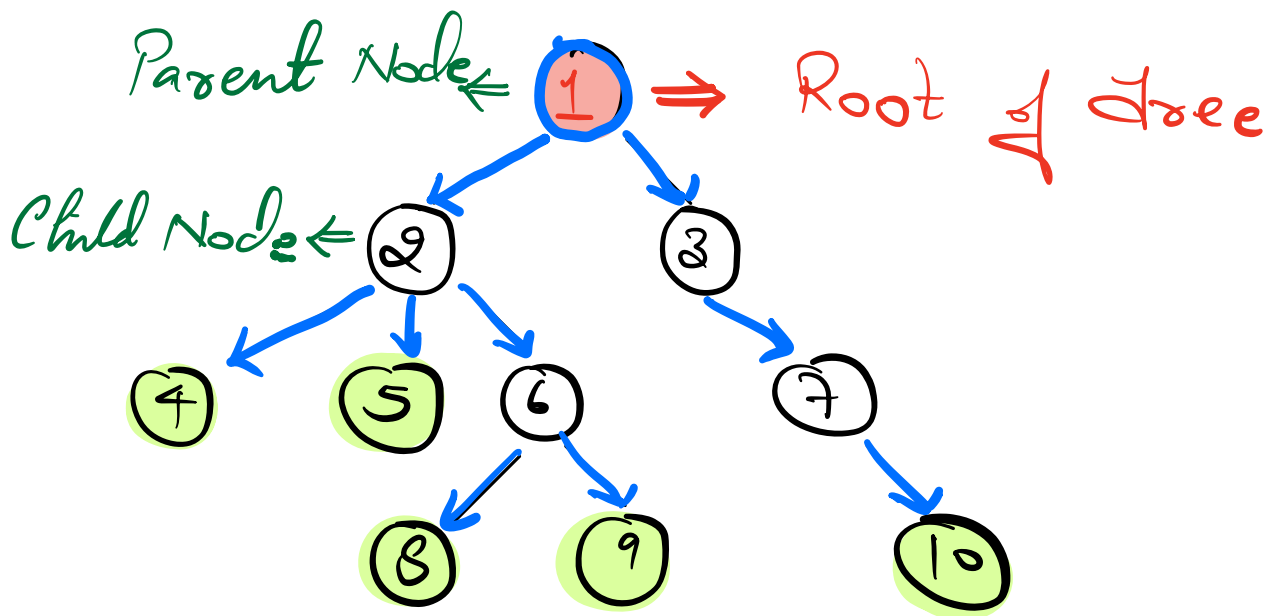
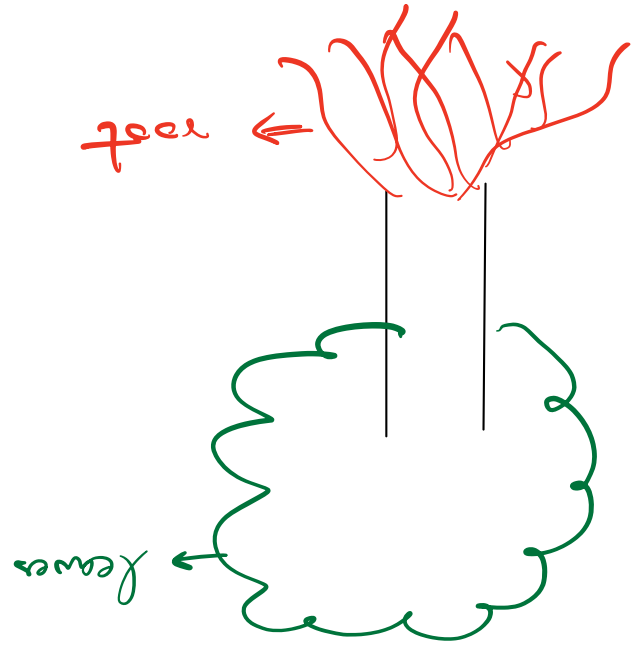
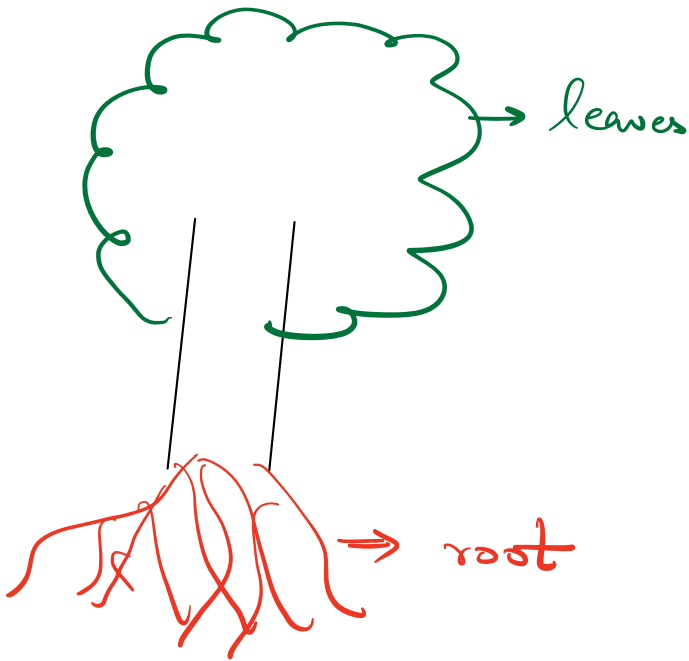


level \Rightarrow After or lower than what we have covered

\Downarrow
PSP $\geq 70\%$

90% chance of
clearing the mob

Trees



An element in a tree \Rightarrow Node.

Q Child nodes of Node 2
④, ⑤, ⑥

Q Parent Node of ⑦
③

Leaf Nodes

\hookrightarrow Nodes with no children

4, 5, 8, 9, 10.

Siblings

\hookrightarrow Nodes with same parent

(2, 3)

level, Height & Depth

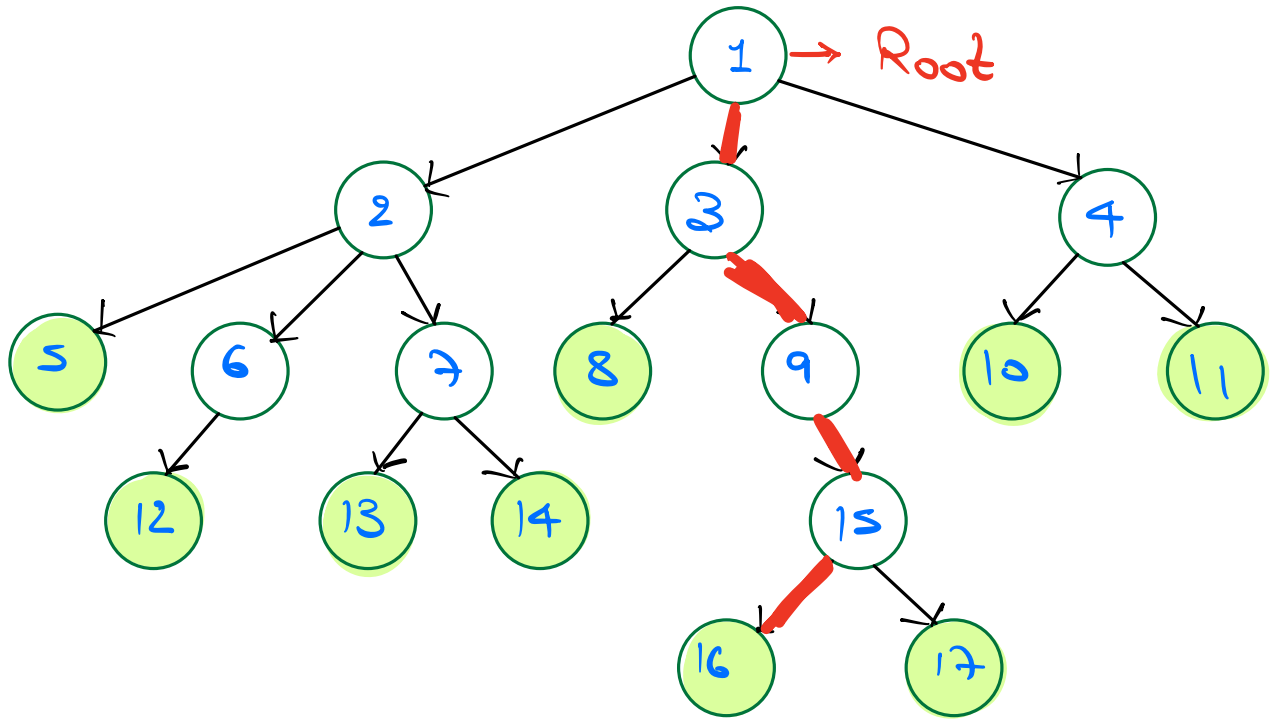
level 0

level 1

level 2

level 3

level 4



Depth of a Node

⇒ level at which a node is present
(Distance of the node from root)

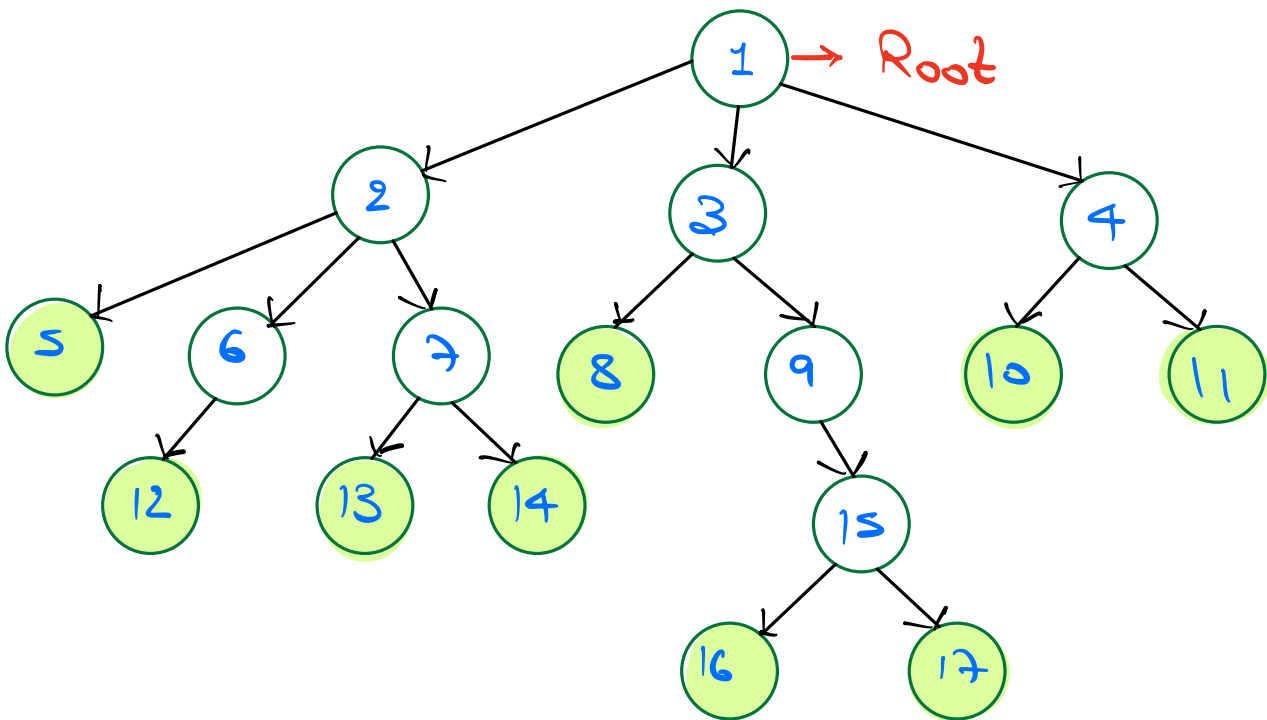
Depth of 14 = 3

Height of a Node

length of the longest path from a node to the furthest leaf.

Height of Node ② \Rightarrow 2

Height of root node = 4 = Height of the tree.



Height of leaf Node = 0.

! For a tree with N nodes.
find the total no. of subtrees possible ??

N unique
root nodes possible $\Rightarrow N$ unique subtrees possible

Nomenclature of a tree

\Rightarrow A tree is named on the basis of max. no. children any node can have.

Max no. of children

Name

2

Binary tree

3

Ternary tree

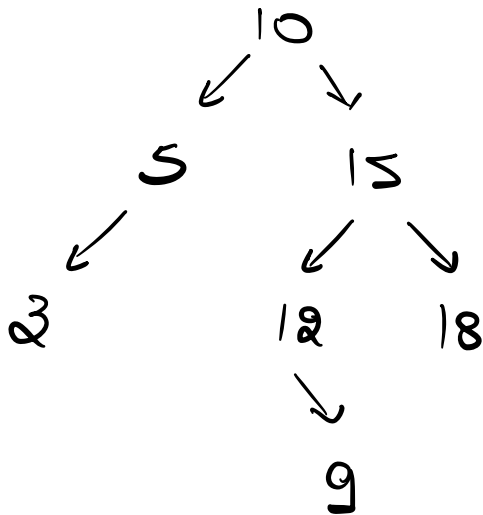
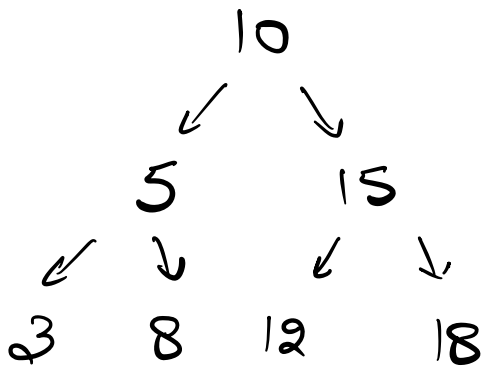
\vdots

N

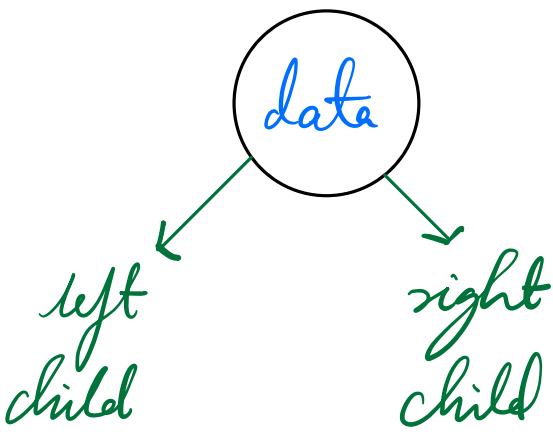
N -ary tree

Binary Tree

⇒ A tree where any node can have at max 2 children (0, 1 or 2)



Structure of a Node in B.T.



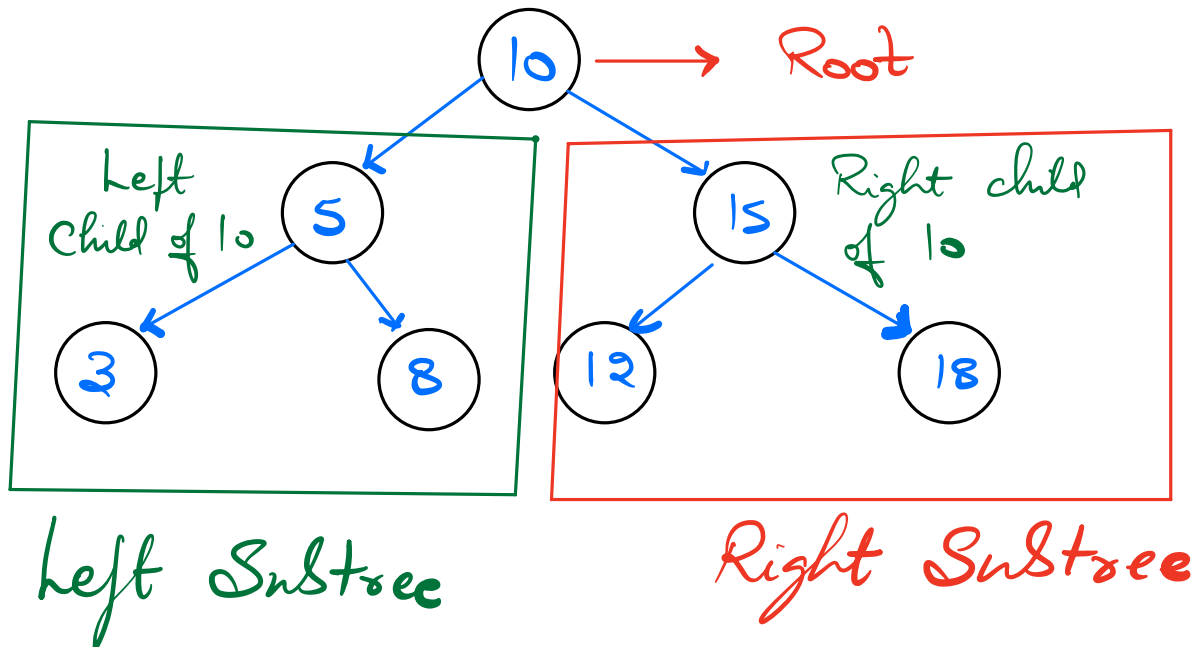
```
class Node {  
    int data;  
    Node left;  
    Node right;  
}
```

```
Node (int x) {  
    data = x;  
    left = right = NULL;  
}
```

}

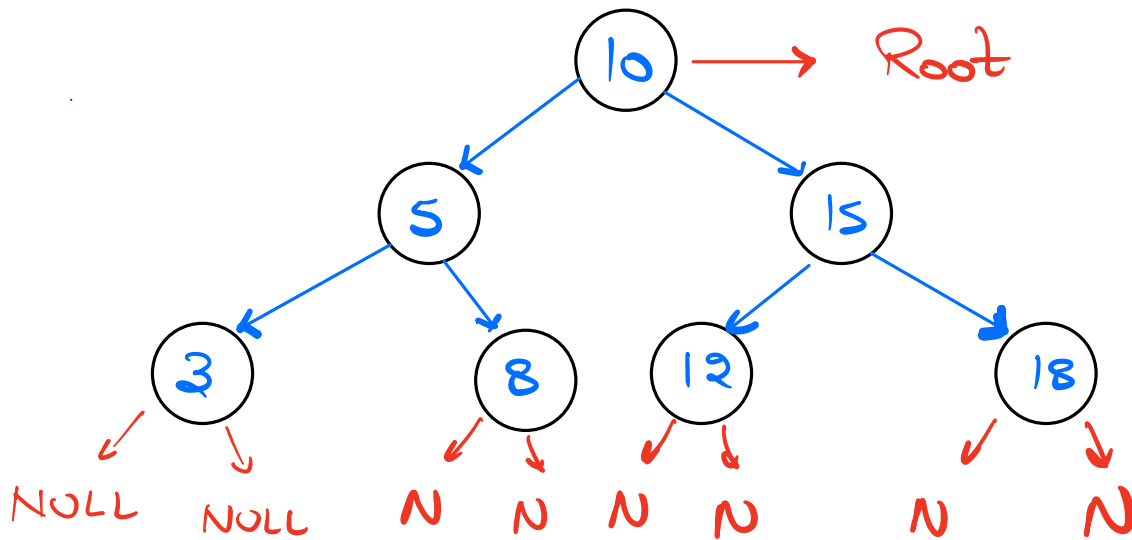
}

Traversals of Tree



Subtree rooted at 10 \Rightarrow Entire Tree

left subtree of 10 \Rightarrow Subtree rooted at left child of 10



1) Pre Order Traversal \Rightarrow Node, Left, Right

2) In Order Traversal \Rightarrow Left, Node, Right

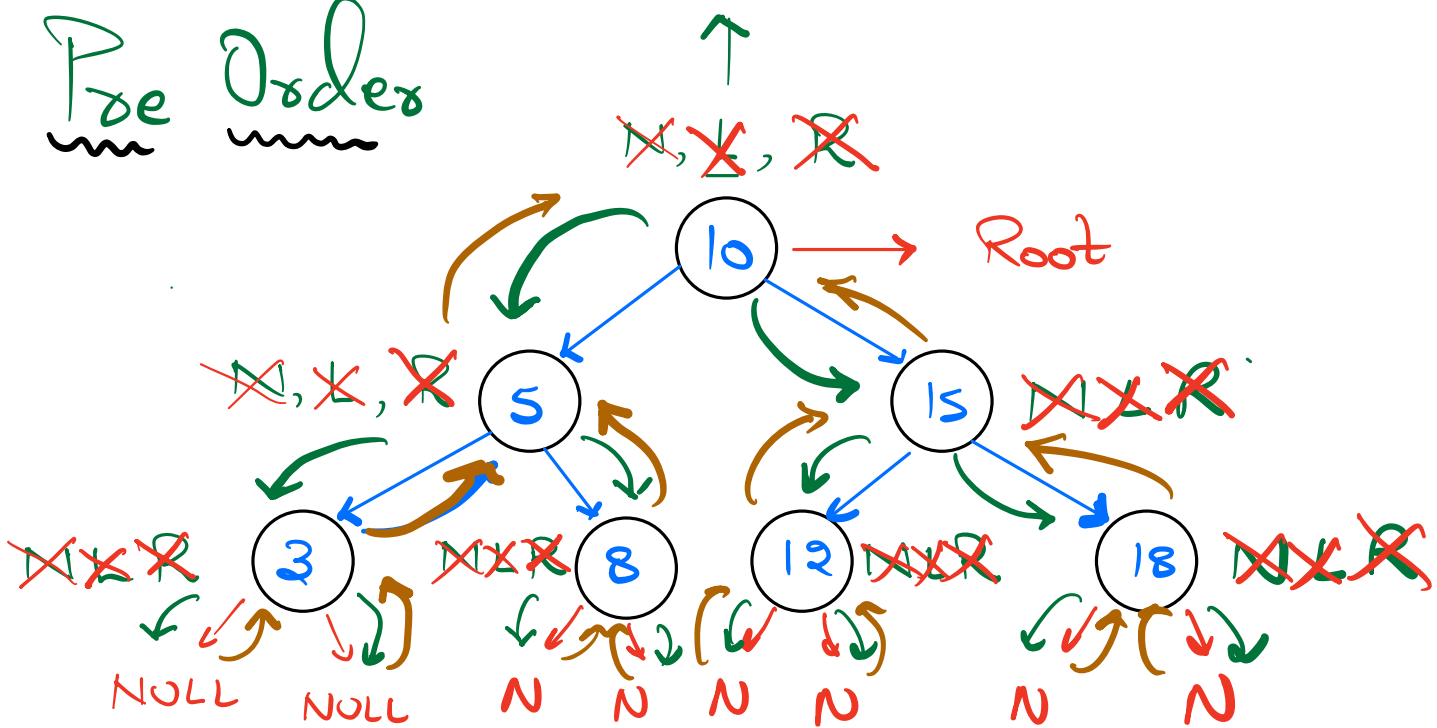
3) Post Order Traversal \Rightarrow Left, Right, Node

4) Level Order Traversal

5) Vertical Order Traversal

} Next Class

Pre Order



Ans: 10, 5, 3, 8, 15, 12, 18

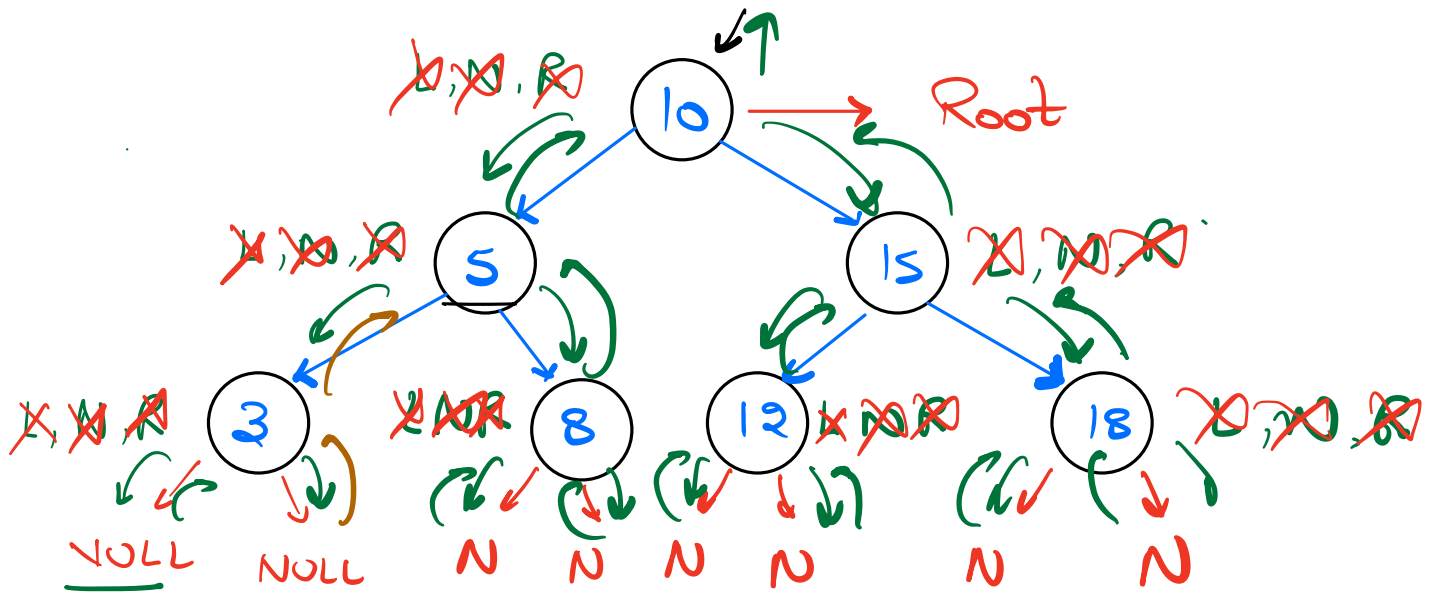
Code

```
void preOrder (Node root) {  
    if (root == NULL) return; }
```

Node
Left
Right
}

```
    print (root.data);  
    preOrder (root.left);  
    preOrder (root.right);
```

In-Order Traversal



3, 5, 8, 10, 12, 15, 18

void inOrder (Node root) {

if (root == NULL) return; {

left
Node
Right
{

inOrder (root.left);
print (root.data);
inOrder (root.right);
}

Post Order

H.W.

```
void postOrder (Node root) {  
    if (root == NULL) return ;
```

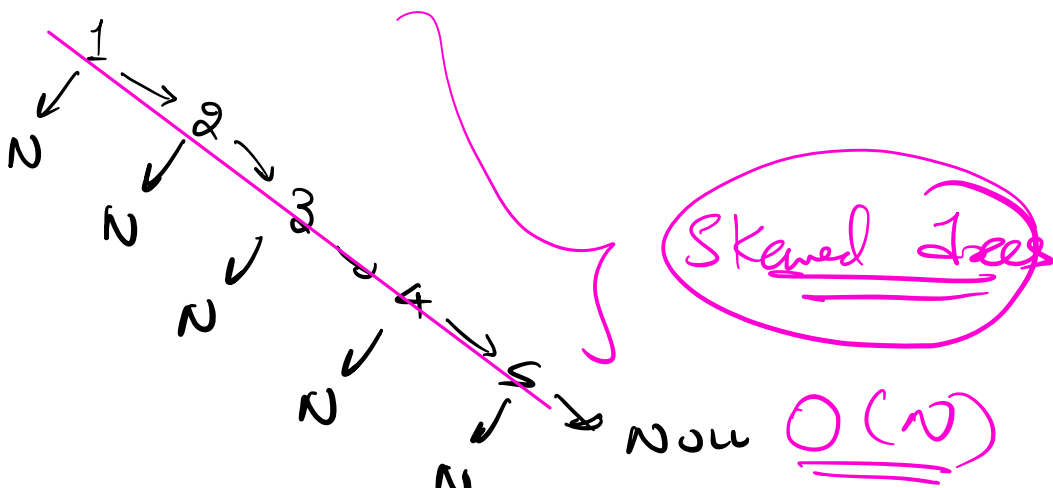
Left
Right
Node
{

```
    postOrder (root.left);  
    postOrder (root.right);  
    print (root.data);  
}
```

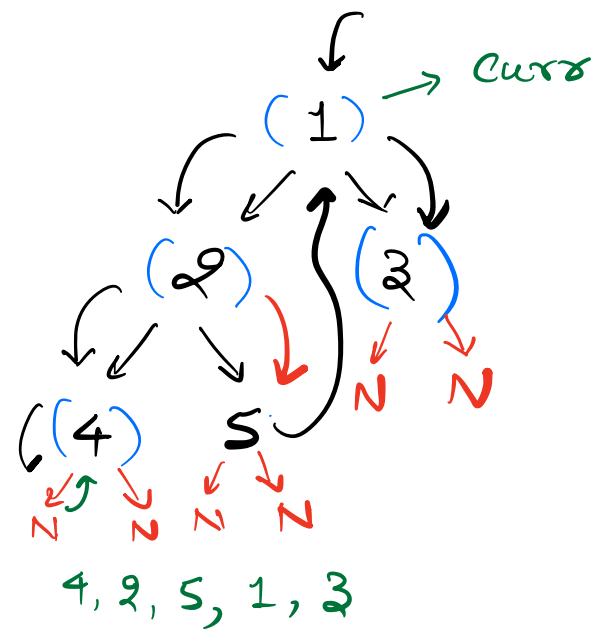
T.C. of Recursion = $\left(\# \text{ Recursive Cells} \right) \times \left(\text{Time per Recursive Cell} \right)$

T.C. = $O(N) \times O(1)$

T.C. = $O(N)$
S.C. = $O(\text{Height})$



Q Can we perform an iterative inOrder traversal of a Binary Tree



4 2 5 1 3

Code

```
Stack <Node> st  
Node curr = root;
```

```
while (curr != NULL || !st.isEmpty()) {
```

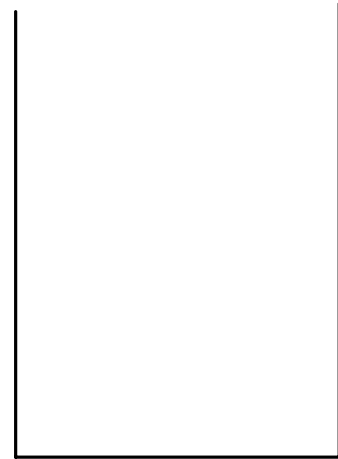
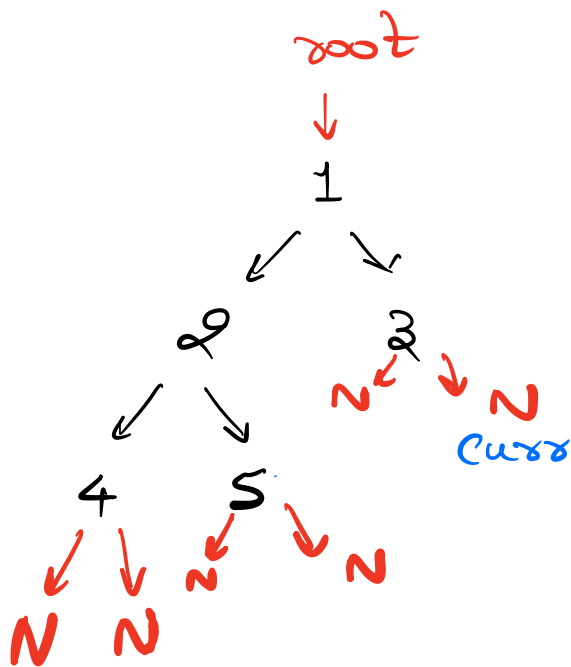
```
    if (curr != NULL) {  
        st.push(curr);
```

curr = curr.left;

else if

curr = st.pop();
print (curr.data);
curr = curr.right;

if



4, 2, 5, 1, 3

Wednesday \Rightarrow Contest

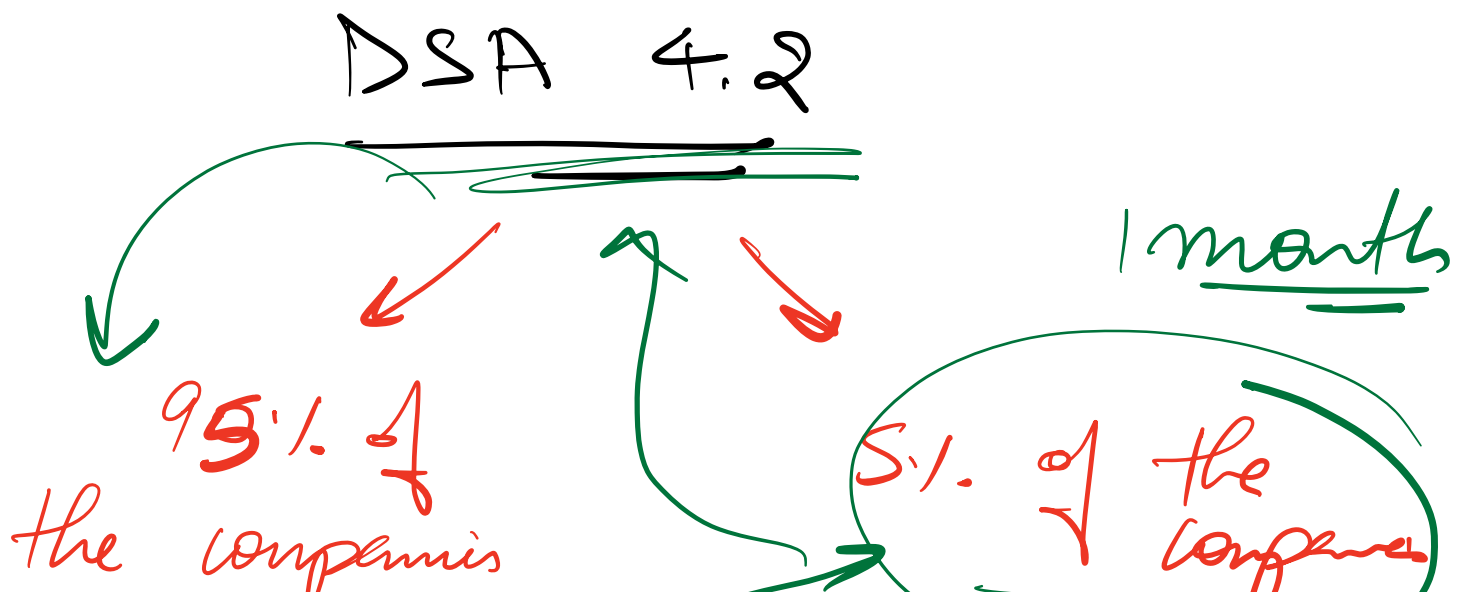
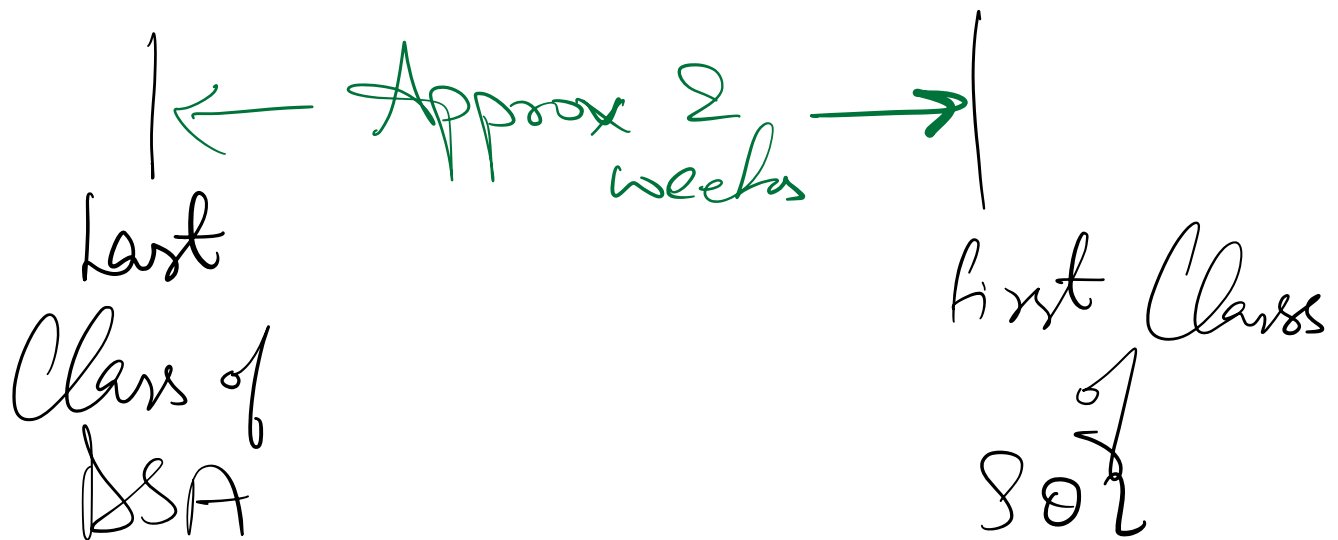


linked list, Stack
& Queue

Friday (26th January) \Rightarrow Off

\Downarrow Next Class

Monday \Rightarrow 29th January



Analysis

DSA 1, 2, 3, 4

→ DSA
→ SOL
→ End of
Course