

8<sup>th</sup> December : Recursion, Maths & OOPS

↳ 2<sup>nd</sup> Dec → PS session  
7<sup>th</sup> Dec → PS session

Contest 1 Reattempt 3

↳ 5<sup>th</sup> December.

Revision : ~~Re Watch Lecture Recording~~

~~Go through lecture notes.~~

1) Solve unseen problems.

↳ Pending assign.  
↳ Pending add. problem  
↳ LeetCode, InterviewBit, GfG,  
HR .....

Able to  
solve

Not able to  
solve (Revise)

- 1) Solution
- 2) Lecture Notes
- 3) Recording (2x speed)

Medium

# Hashing : Introduction

↳ Int: Hashing.

2) Problems

3) Internal Implementation

HM/H\$ ⇒ Data Structure

impl ⇒ Hashing ⇒ Concept

## Hash Map

Data Structure which implements a Hash Table.

Eg

| Key                    | Value. |
|------------------------|--------|
| Student_email          | PSP    |
| ayush.sharma@scdn.com  | 100    |
| parth.thewol@gmail.com | 100    |

awesome guy berrhit @ ju  
bipin deool 22 @ gmail . in  
hackerzy @ gmail

98

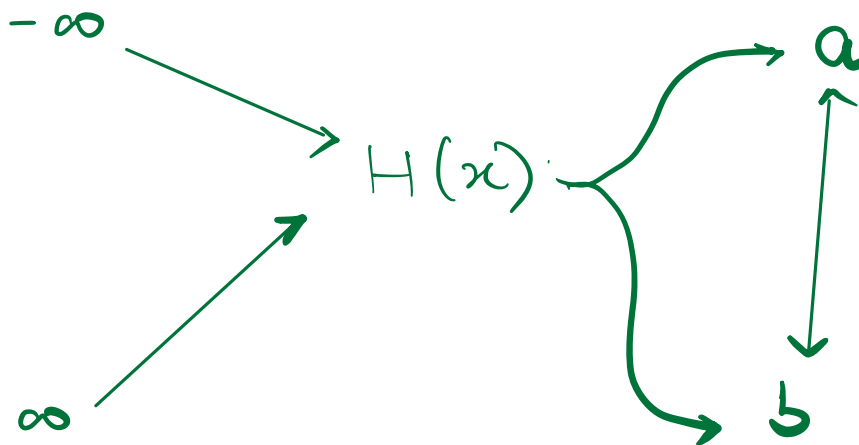
35

56

⋮

list of  $\langle \text{Key}, \text{Value} \rangle$  pair.

Key  $\rightarrow$  (Hash Function)  $\rightarrow$  where to store  
value for the  
Key



Modulo  $\Rightarrow x \% M = [0, M-1]$

$\downarrow$

$[-\infty, \infty]$   
Range

T.C. of Search, insert, deletion  
 $= O(1)$

Cond:

1) Key must be unique.

2) Values can be anything.

No  $\Rightarrow$  Ass/H.W.

Advanced batch  $\Rightarrow$  Add. problems

---

HashMap  $\langle \frac{\text{Key}}{\text{datatype}}, \frac{\text{Value}}{\text{datatype}} \rangle$  name;

! Value  
Population of every Key  
country

HashMap  $\langle \text{String}, \text{Long} \rangle$

! Value.  
No. of states for every Key  
country

HashMap  $\langle \text{String}, \text{int} \rangle$

! Value.  
Name of all states for every Country. <sup>Key</sup>

~~< Key, Value >~~  
~~India, UP~~  
~~India, Rajasthan~~  
~~India, MP~~

< Key, Value >  
India, [UP, Rejas., MP....]

HashMap < String, List <String> >

! Value  
Population of Each State for every Country <sup>Key</sup>

HashMap < String, HashMap < String, Long > >

Hash Map

|         |               |            |     |            |
|---------|---------------|------------|-----|------------|
| Java    | C++           | Python     | JS  | C#         |
| HashMap | unordered_map | dictionary | map | dictionary |

Hash Set

|         |               |        |     |          |
|---------|---------------|--------|-----|----------|
| Java    | C++           | Python | JS  | C#       |
| HashSet | unordered_set | set    | set | HashSet. |

Q1 Given  $N$  elements &  $Q$  no. of queries.

Query: Given  $x \Rightarrow$  Return the frequency of  $x$  in the array.

Eg:  $A = [2, 6, 3, 8, 2, 8, 3, 8, 10, 6]$

$Q = 3$

$x$

2

8

5

freq of  $x$ .

2

3

0

Sol<sup>n</sup>  $\Rightarrow$  Brute force

$\Rightarrow \underbrace{\forall \text{ query}}_{Q} \Rightarrow \underbrace{\text{Iterate the array}}_{N} \text{ \& find freq.}$

T.C. =  $O(Q \times N)$

S.C. =  $O(1)$

Q How can we improve T.C. ??

Key

Value

$\langle \text{Element}, \text{Frequency} \rangle$

$A = [2, 6, 3, 8, 2, 8, 3, 8, 10, 6]$

$i$

$\langle \text{Key}$

$\text{Value} \rangle$

2  
6  
3  
8  
10

~~1~~ 2  
~~1~~ 2  
~~1~~ 2  
~~1~~ ~~2~~ 3  
1

$O(N)$

$Q = 3$

Elements

2  
8  
5

Frequency

2  $\Rightarrow O(1)$   
3  $\Rightarrow O(1)$   
0  $\Rightarrow O(1)$

$\} Q$

$T.C. = O(N + Q)$

---

Hash Map

- 1) insert (Key, Value) : insert a new entry
- 2) containsKey (Key) : Returns True if the given Key is present in HM.
- 3) update (Key, updated value) : Update the value for given Key
- 4) size() : Return the size of the HM.  
↓  
No. of entries  
↓,  
No. of Keys.
- 5) delete (Key) : Delete the entry for given Key.

## Hash Set

- 1) insert (Key)
  - 2) containsKey (Key) : Returns True if the given Key is present in HS
  - 3) size() : Total no. of Keys
  - 4) delete (Key) : Removes the given Key from HS
-



# Code

1) Create a frequency map (HM)

```
HashMap < int, int > hm;
```

```
for (i=0; i < N; i++) {
```

```
    if (! hm.containsKey (A[i])) {
```

```
        hm.insert (A[i], 1);
```

```
    }
```

```
    else {
```

```
        hm.update (A[i], hm.get(A[i]) + 1);
```

```
    }
```

```
}
```

2) Answer all queries // queries[N]

```
for (i=0; i < Q; i++) {
```

```
    if (hm.containsKey (queries[i])) {
```

```
        print (hm.get (queries[i]));
```

```
    }
```

```
    else {
```

```
        print (0);
```

```
    }
```

```
}
```

Q Given an integer array of size  $N$ .  
Return the first non repeating element

Eg  $N = 6$   
 $A = [1, 2, \underline{3}, 1, 2, 5]$   
 $Ans = 3$

$N = 8$

$A = [4, 3, 3, 2, 5, 6, 8, 5]$

$Ans = 4$

Sol<sup>n</sup>  $\Rightarrow$  Brute Force

$\forall i \Rightarrow$  Iterate & find freq. of  $A[i]$  } HM

$\Downarrow$   
The first  $i$  for which freq of  $A[i]$  is  
1 becomes my answer.

2) Opt.

- 1) Create a frequency map (HM)
- 2) Iterate over the array & check the freq of each ele.
- 3) The first ele with freq 1 is ans.

! In point 2, can we iterate over HM  
NO

The elements of HM are not stored in the order of insertion.

! Given an integer array of size  $N$ .  
Return the count of distinct elements in the array.

$N = 5$ ,

$A = [3, 5, 6, 5, 4]$

Ans = 4

Sol<sup>n</sup> Use HashSet

Sol<sup>n</sup> Use HashSet

- 1) Insert all elements in a HashSet
- 2) Return its size.

Code

```

HashSet <int> hs;

for (i=0; i<N; i++) {
    hs.insert (A[i]);
}

return hs.size();

```

T.P. =  $O(N)$

S.P. =  $O(N)$

$$A = \left[ \begin{array}{c} \vdots \\ \text{Range } [0, 100] \end{array} \right]$$
~~Range =  $[0, 10^9]$~~ 
$$\text{Size} = [10^4]$$