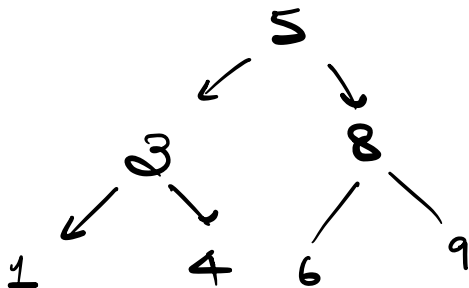
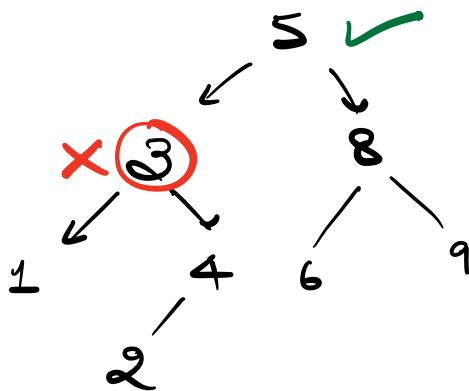
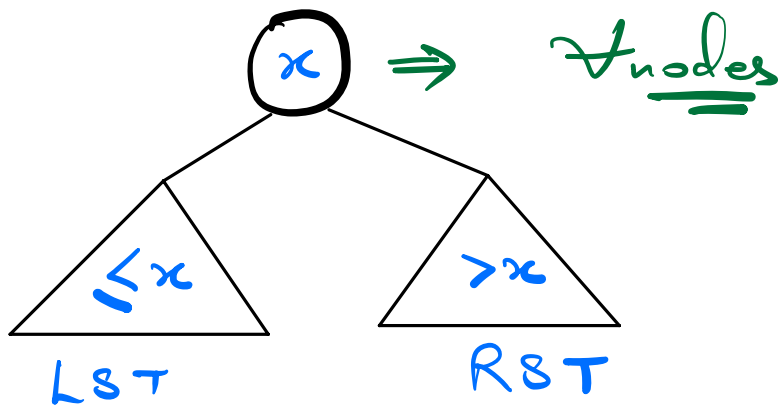
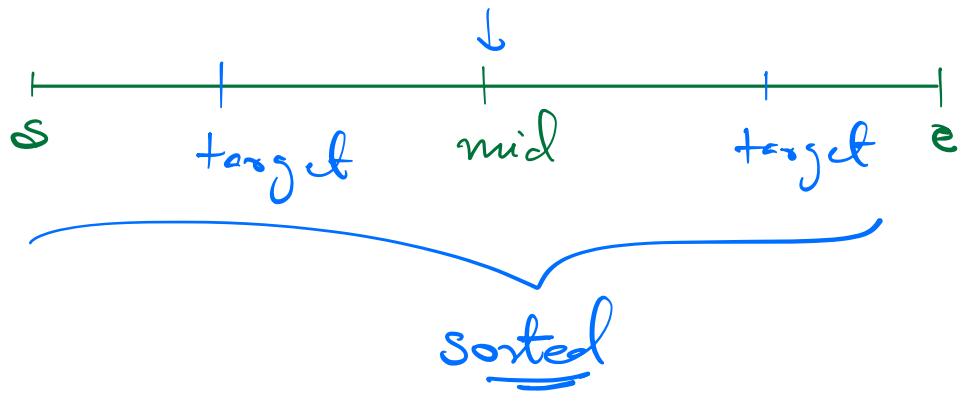
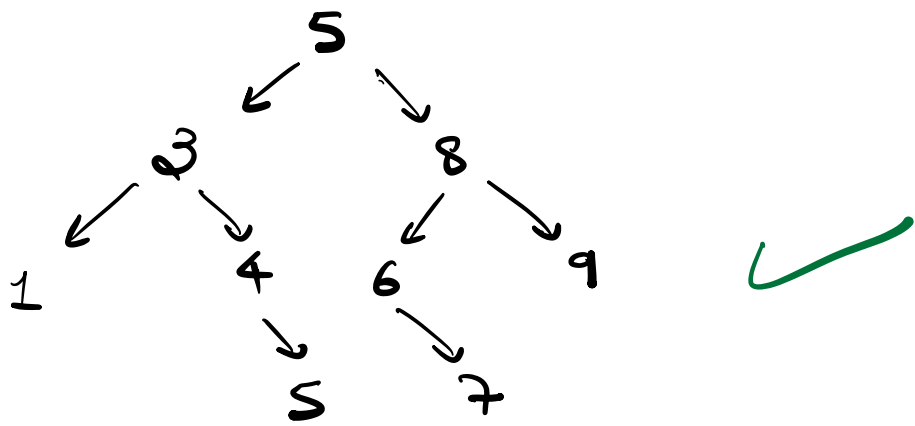


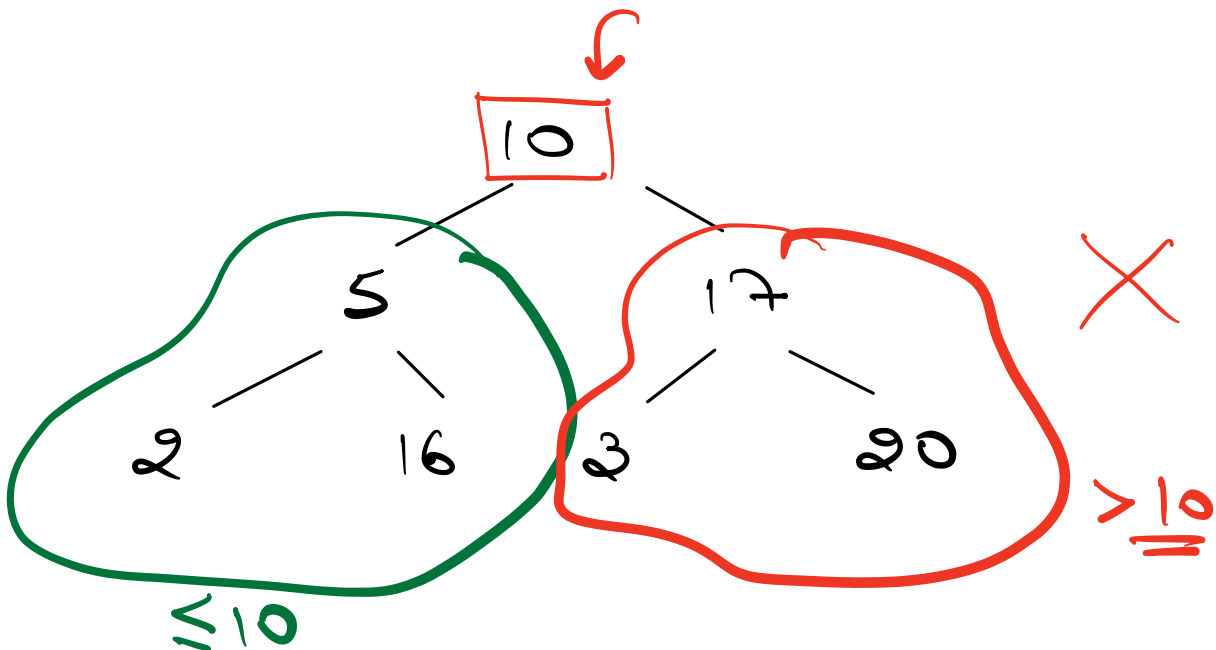
# Binary Search Tree



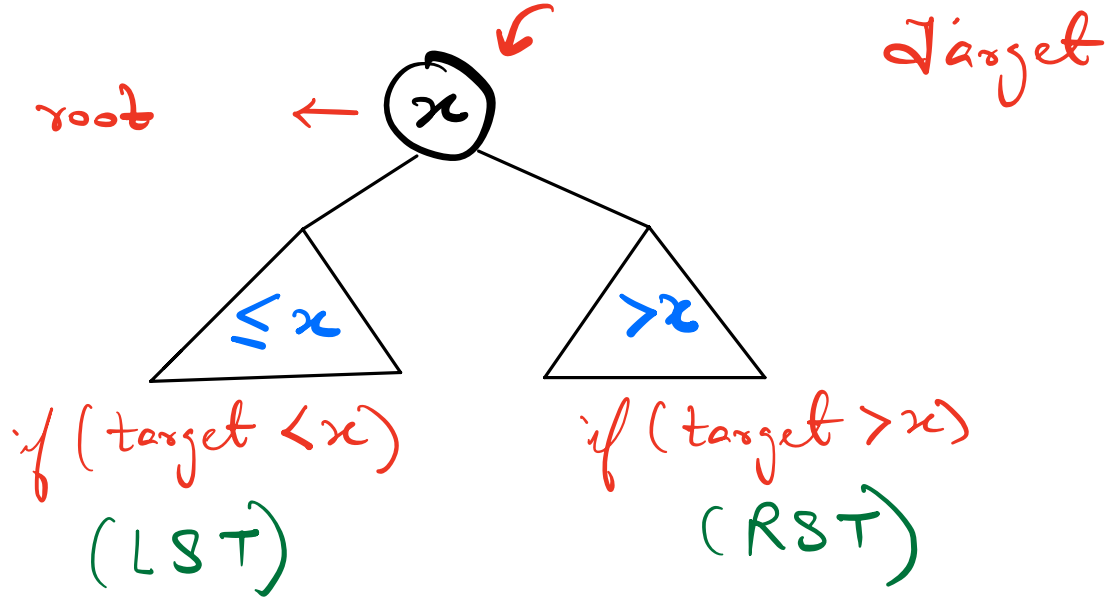


$\forall \text{ nodes} \Rightarrow$ 

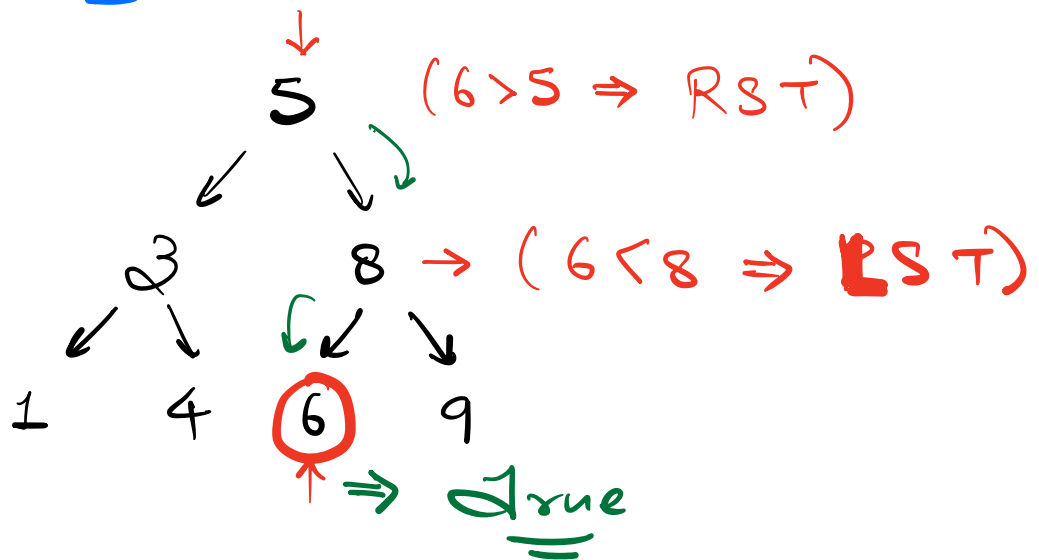
$$\begin{aligned} \text{node.left.data} &\leq \text{node.data} \\ \& \\ \text{node.right.data} &> \text{node.data} \end{aligned}$$



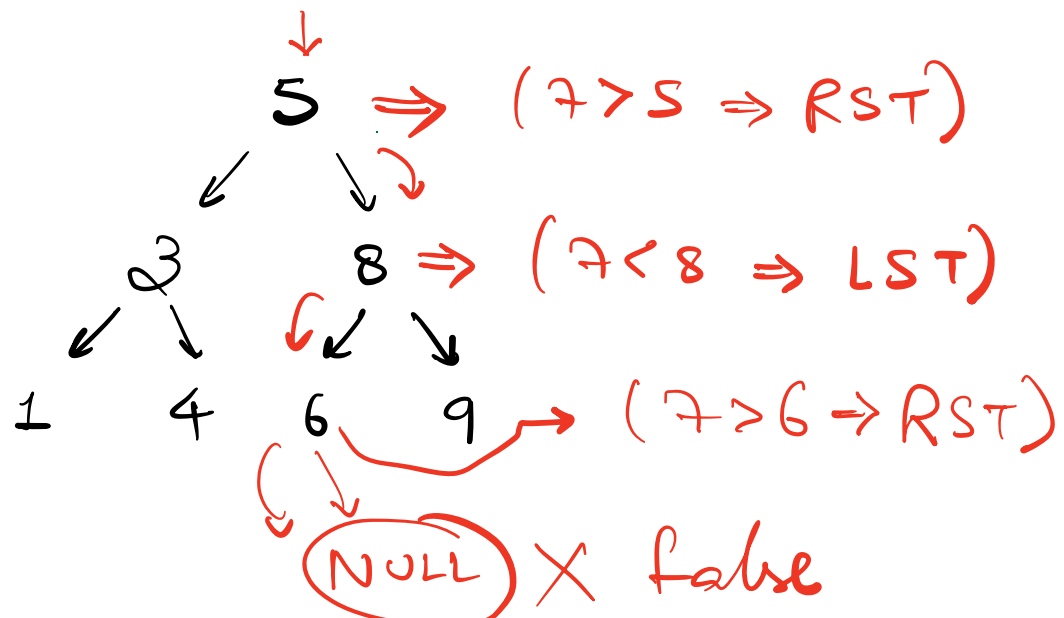
! Searching in a BST



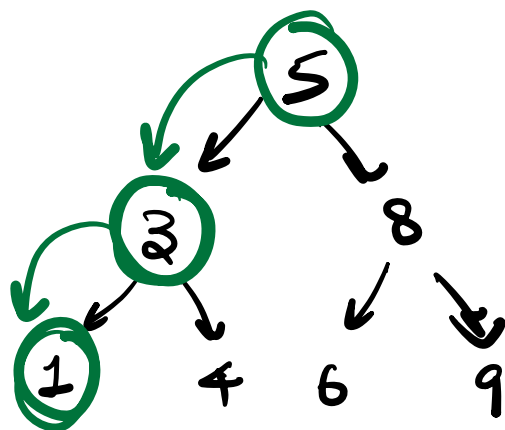
Target = 6



Target = 7  $\Rightarrow$  false



Quiz



Ans = 3

## Code

```
boolean search (Node root, int target) {
```

```
    if (root == NULL) {  
        return false;  
    }
```

```
    if (root.data == target) {  
        return true;  
    }
```

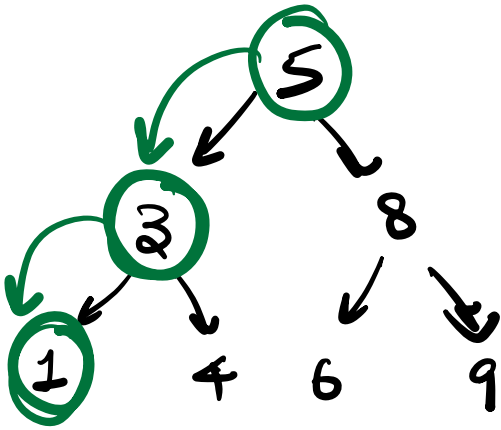
```
    else if (root.data > target) {  
        return search (root.left, target);  
    }
```

```
    else {  
        return search (root.right, target);  
    }
```

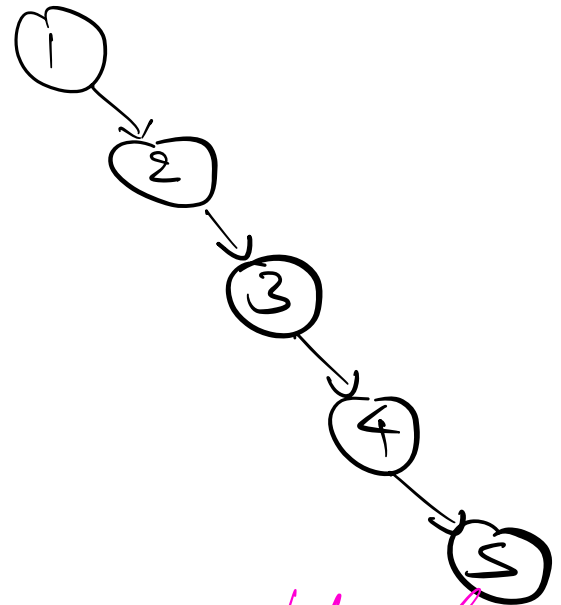
```
}
```

$$T.C. = O(\text{Height})$$

$\Downarrow$   
 $O(N) \Rightarrow \text{Skewed Tree}$

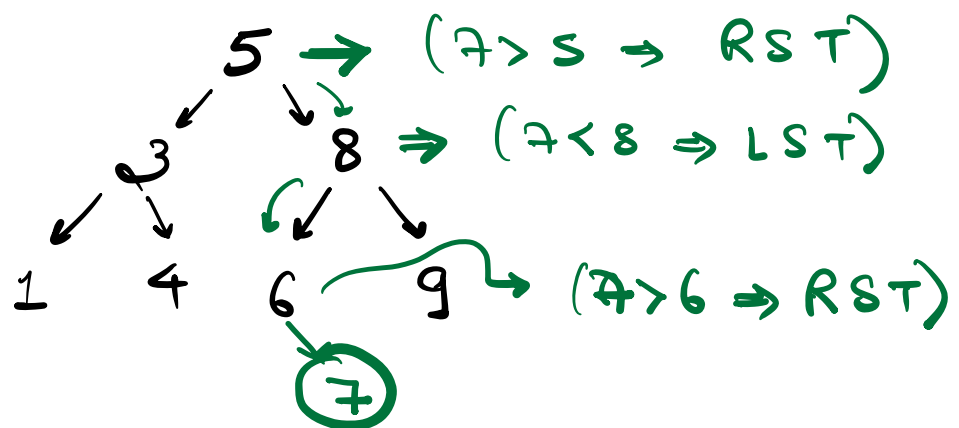


Height =  $\log N$



Height =  $N$

## Insertion in a BST



Insert a Node with a value 7

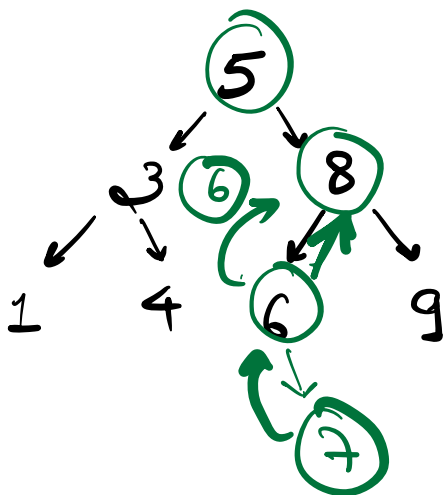
# Code

```
Node insert ( root, value ) {
```

```
    if ( root == NULL ) {  
        return new Node (value);  
    }
```

```
    if ( value <= root->data ) {  
        root->left = insert ( root->left, value );  
    } else {  
        root->right = insert ( root->right, value );  
    }
```

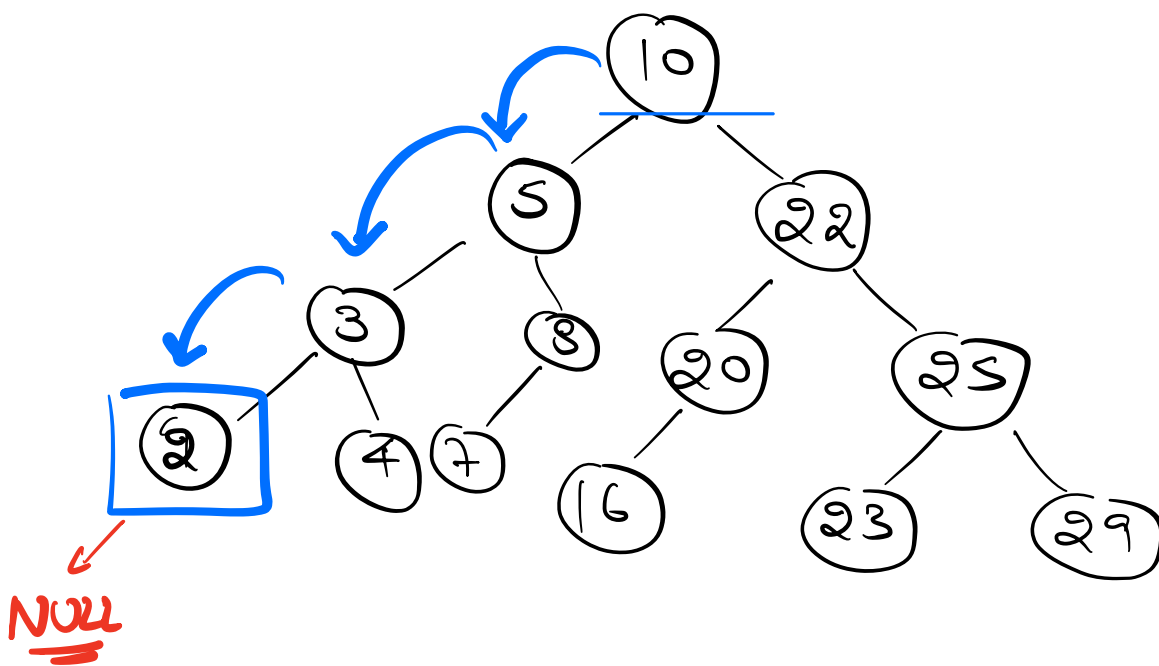
```
    return root;  
}
```



6. right = insert ( NULL, 7 )  
8. left ← insert ( 6, 7 )  
5. right = insert ( 8, 7 )  
insert ( 5, 7 )

$$T.C. = O(\text{Height}) = O(n)$$

$$S.C. = O(\text{Height}) = O(\underline{\underline{n}})$$



Code

Node smallest (root) &

if (root == NULL) & return NULL;

temp = root;

while (temp->left != NULL) &

temp = temp->left;

}

return temp;

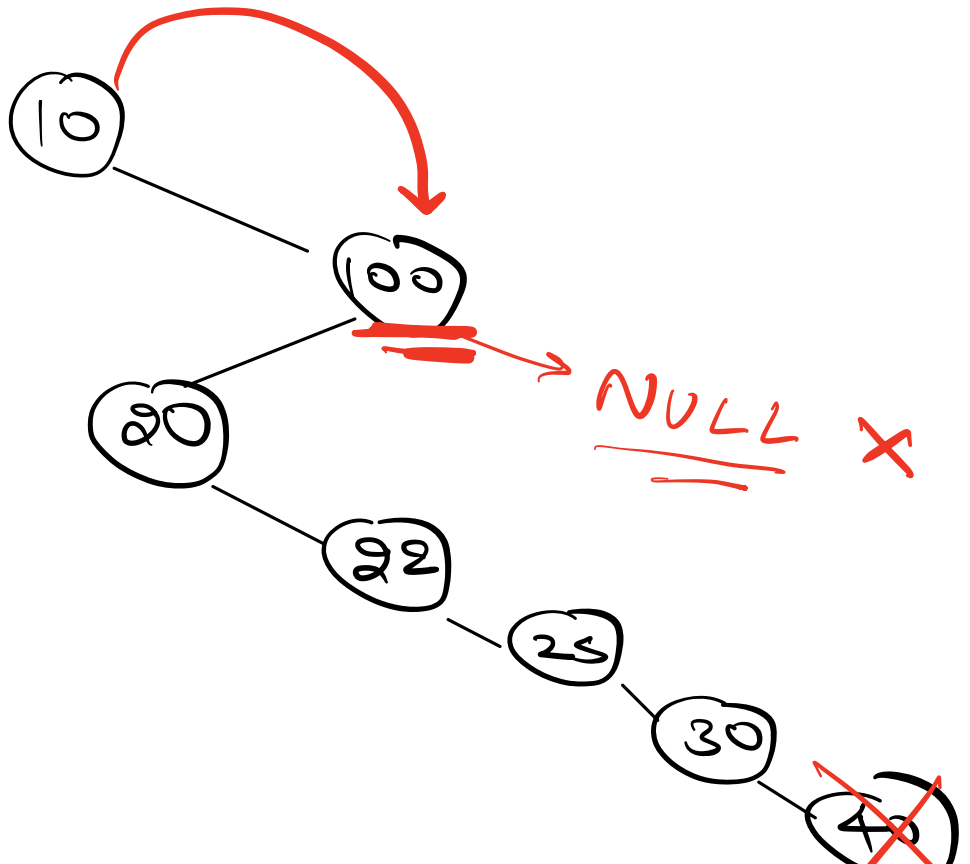
}

T.C. =  $O(\text{Height}) = O(n)$

S.C. =  $O(1)$

find the largest Node of

BST







Node largest (root) &

if (root == NULL) & return NULL;

temp = root;

while (temp->right != NULL) &

temp = temp->right;

&

return temp;

&

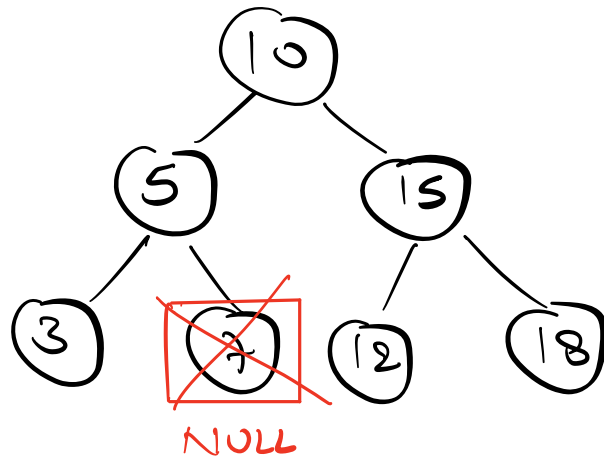
T.C. =  $O(\text{Height}) = O(N)$

S.C. =  $O(1)$

---

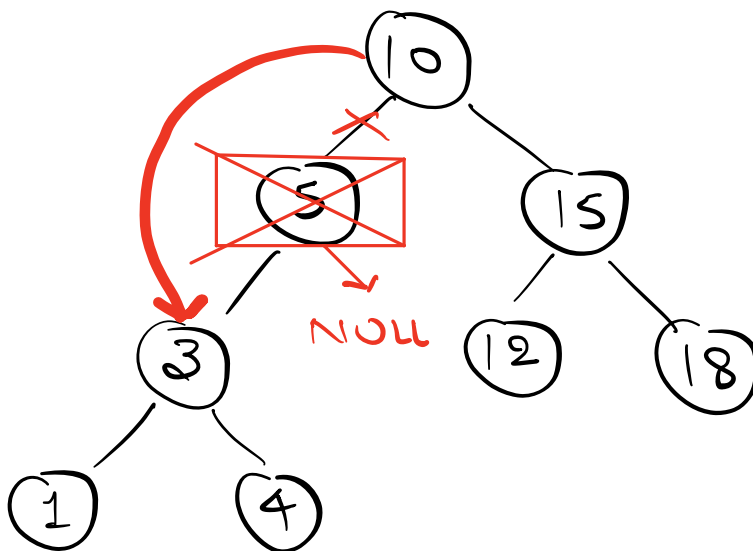
Delete in a Binary Search  
Tree

Case I: Node with No Children (leaf Node)  
Delete the Node

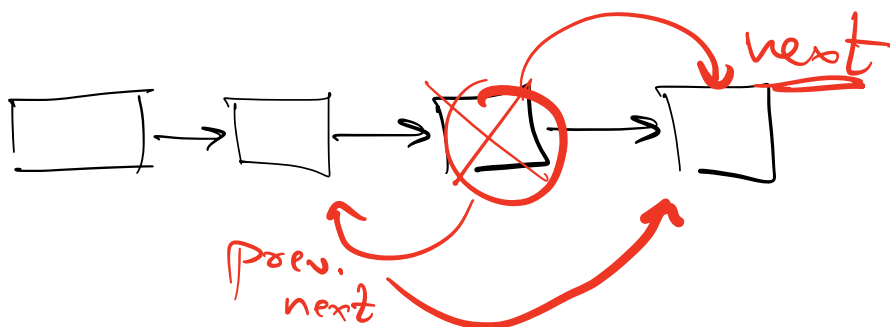


Target = 7

Case II Node with One Child.  
Replace with the non NULL child.

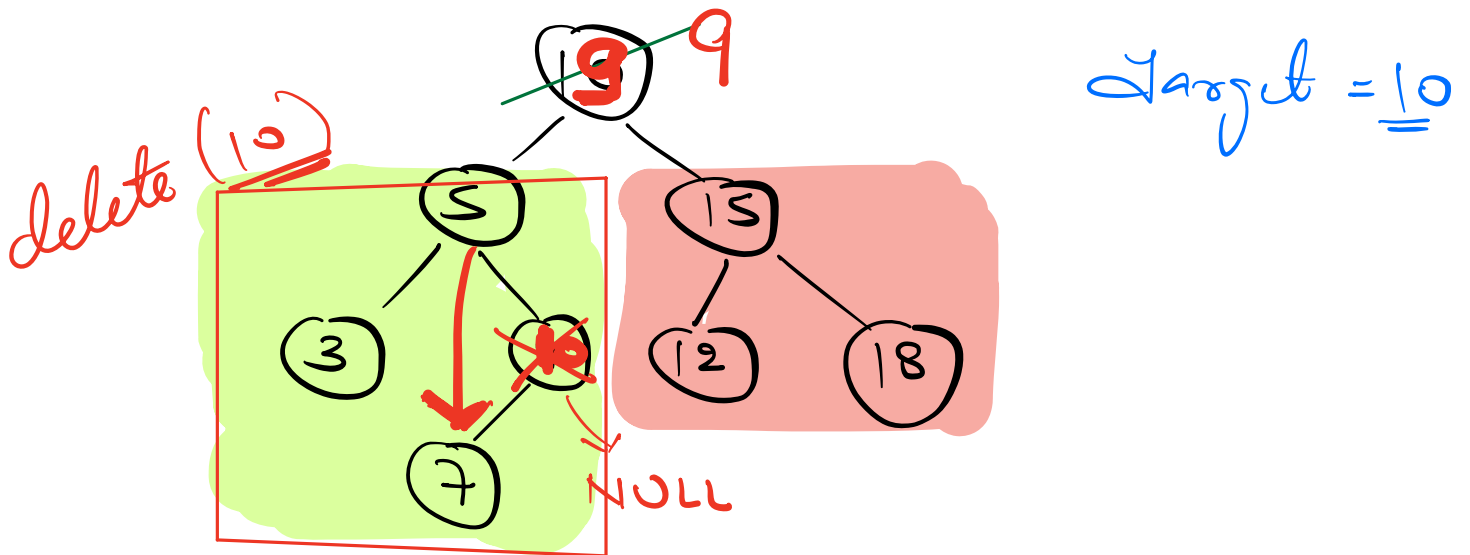


Target = 5



# Case III

Node with Two Children



In Order  
Predecessor

largest Node of  
LST

In Order  
Successor

Smallest Node  
of RST

In order  
traversal

=

Root

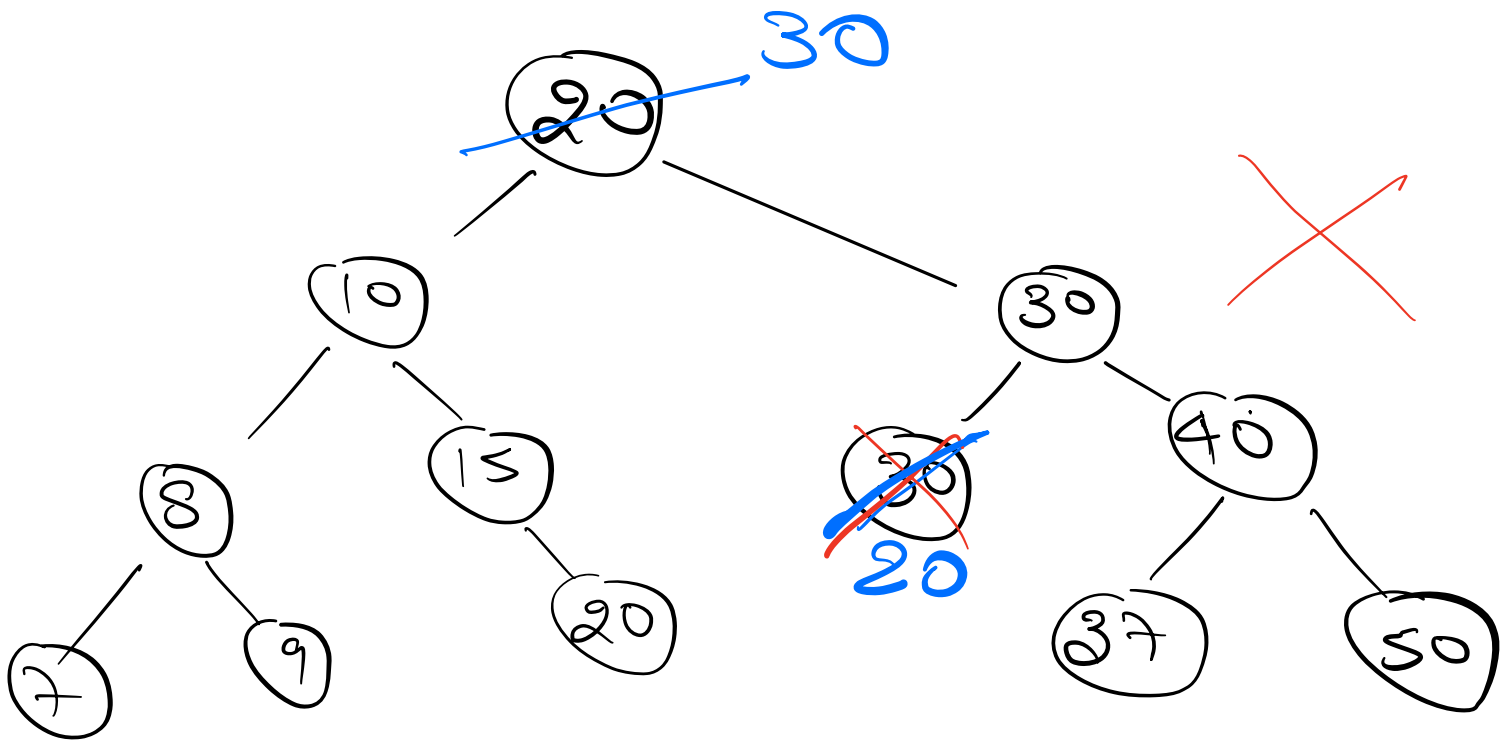
IP 3, 5, 7, 9, 10, IS 12, 15, 18

Sorted

L, N, R

$L \leq N < R$

H.W. Think if anything can go wrong in case of replacing with inorder successor if equality is in the left.



Code

Node delete (root, value) {

if (root == NULL) { return NULL; }

if (root.data == value) {

if (root.left == NULL && root.right == NULL) {  
return NULL;  
}

if (root.left == NULL) {  
return root.right;  
}

if (root.right == NULL) {  
return root.left;  
}

Node temp = root.left;

while (temp.right != NULL) {  
temp = temp.right;  
}

swap (root, temp); // swap values.

root.left = delete (root.left, value);

else if

```
if (value < root.data) {
    root.left = delete(root.left, value);
} else if {
    root.right = delete(root.right, value);
}
```

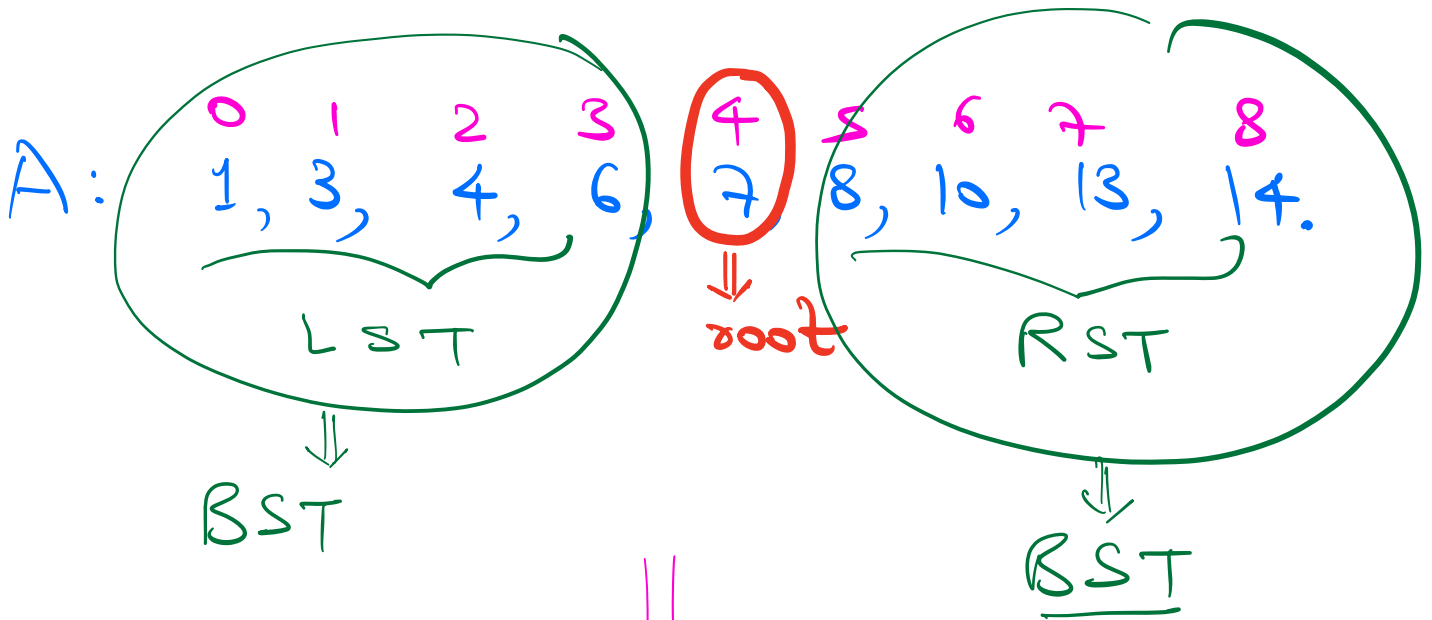
return root

}

Create a BST

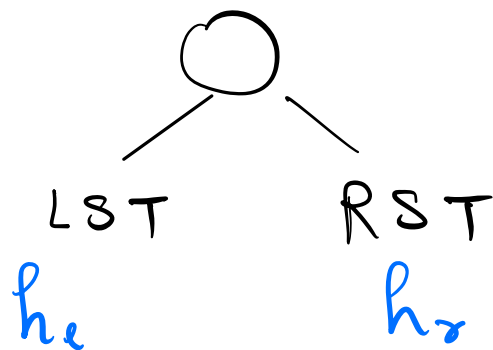
Given a sorted array

↳ Create a BST.

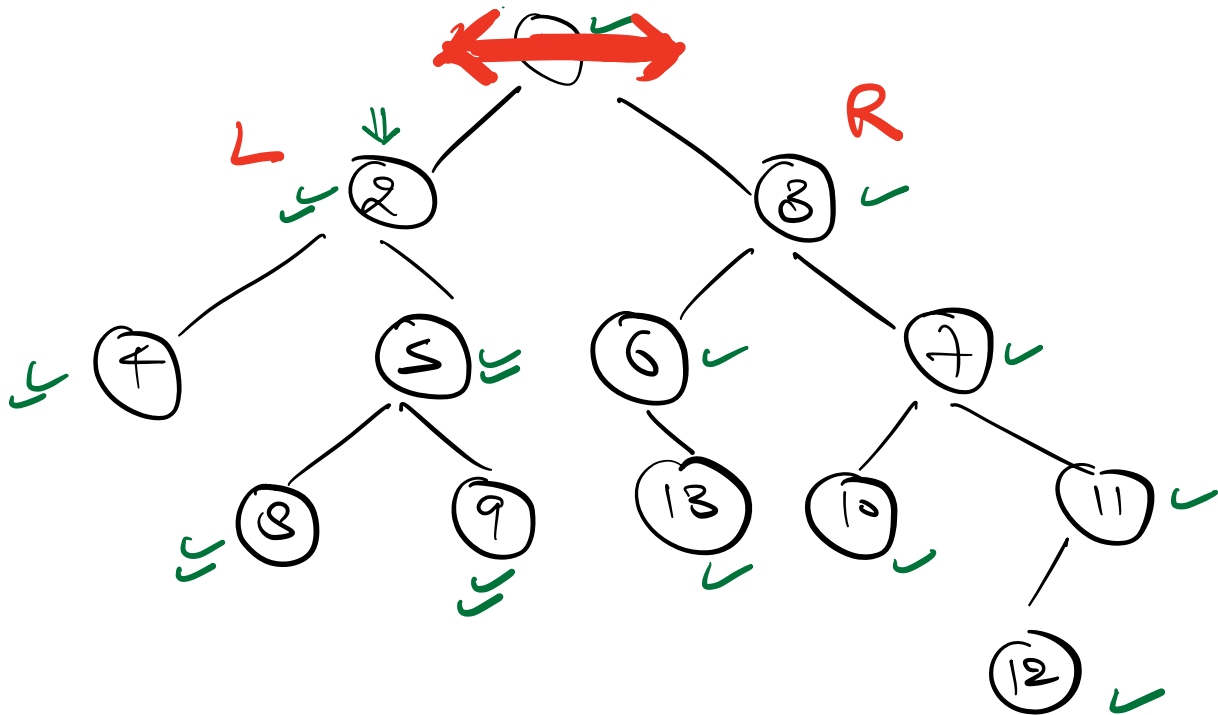


H.W.

Code



$$|h_l - h_r| \leq 1$$



$L, R, N \Rightarrow \underline{\underline{Post}}$

Given Inorder & Postorder

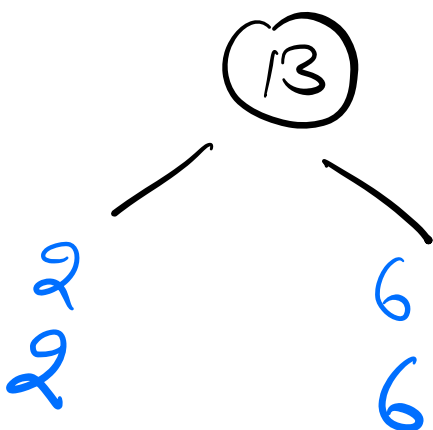


Construct the Binary Tree

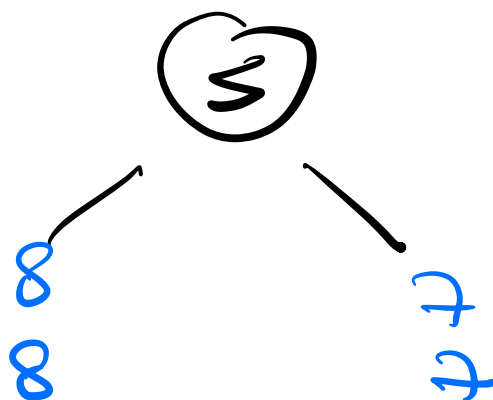
In: L, N, R : 2, 13, 6 (LST) (10) (RST) 8, 5, 7

Post: L, R, N : 2, 6, 13 (LST) 8, 7, 5 (RST) (10) (Root)

2ST RST  
In: 2, (13), 6  
Post: 2, 6, (13)  
L R



LST RST  
In: 8, (5), 7  
Post: 8, 7, (5)  
L R





②

⑥

⑧

⑦