

3<sup>rd</sup> Year

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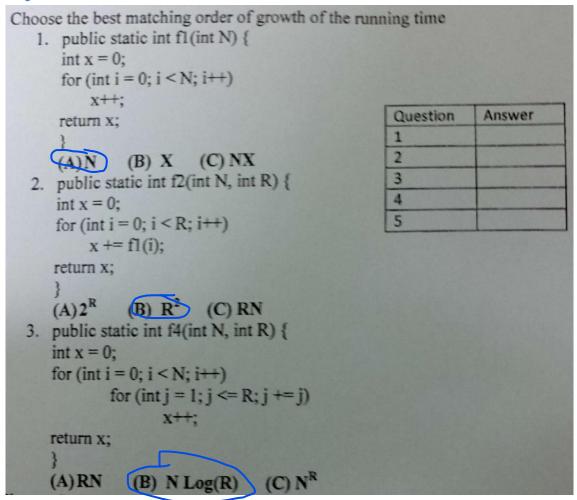
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#### PART I: Analysis of Normal Code

#### **QUESTION#1**



#### **QUESTION#2**

For each group of functions, sort the functions in increasing order of asymptotic (big-O) complexity:

Group1	Group2 Group3	
$   f_1(n) = n^{0.999999} \log n    f_2(n) = 10000000n    f_3(n) = 1.000001^n    f_4(n) = n^2 $	$f_1(n) = 2^{2^{1000000}}$ $f_2(n) = 2^{100000n}$ $f_3(n) = \binom{n}{2}$ $f_4(n) = n\sqrt{n}$	$   f_1(n) = n^{\sqrt{n}}    f_2(n) = 2^n    f_3(n) = n^{10} \cdot 2^{n/2}    f_4(n) = \sum_{i=1}^n (i+1) $

1 -> 2 -> 4 -> 3

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# Algorithms Design and Analysis ANALYSIS Sheet



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#### Group1

**Solution:** The correct order of these functions is  $f_1(n)$ ,  $f_2(n)$ ,  $f_4(n)$ ,  $f_3(n)$ . To see why  $f_1(n)$  grows asymptotically slower than  $f_2(n)$ , recall that for any c > 0,  $\log n$  is  $O(n^c)$ . Therefore we have:

$$f_1(n) = n^{0.999999} \log n = O(n^{0.999999} \cdot n^{0.000001}) = O(n) = O(f_2(n))$$

The function  $f_2(n)$  is linear, while the function  $f_4(n)$  is quadratic, so  $f_2(n)$  is  $O(f_4(n))$ . Finally, we know that  $f_3(n)$  is exponential, which grows much faster than quadratic, so  $f_4(n)$  is  $O(f_3(n))$ .

#### Group2

**Solution:** The correct order of these functions is  $f_1(n)$ ,  $f_4(n)$ ,  $f_3(n)$ ,  $f_2(n)$ . The variable n never appears in the formula for  $f_1(n)$ , so despite the multiple exponentials,  $f_1(n)$  is constant. Hence, it is asymptotically smaller than  $f_4(n)$ , which does grow with n. We may rewrite the formula for  $f_4(n)$  to be  $f_4(n) = n\sqrt{n} = n^{1.5}$ . The value of  $f_3(n) = \binom{n}{2}$  is given by the formula n(n-1)/2, which is  $\Theta(n^2)$ . Hence,  $f_4(n) = n^{1.5} = O(n^2) = O(f_3(n))$ . Finally,  $f_2(n)$  is exponential, while  $f_3(n)$  is quadratic, meaning that  $f_3(n)$  is  $O(f_2(n))$ .



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#### PART II: Analysis of Recursive Code

#### **QUESTION#1**

What's the exact <b>complexity</b> of the following code?	
Fun2(A, S, E)	Θ(N)
<pre>IF (S == E) return A[S]</pre>	<b>b.</b> Θ(N <sup>3</sup> )
S1 = (S + E) / 3	<b>c.</b> $\Theta((\log(N))^3)$
$S2 = 2 \times (S + E) / 3$	<b>d.</b> Θ(log(N))
R1 = Fun2(A, S, S1) R2 = Fun2(A, S1+1, S2)	<b>e.</b> $\Theta(3 \times \log(N))$
R3 = Fun2(A, S1+1, S2) R3 = Fun2(A, S2+1, E)	f. $\Theta(1)$
Return Max(R1, R2, R3)	

#### **QUESTION#2**

#### For the following "Main" function:

```
Main(A, N)
                                               Fun(\mathbb{Z}, \mathbb{L})
       I = N
                                                      If (Z == 2)
       Sum = 0
                                                             Return 0
       For (J = 0; J < I; J += N/2)
                                                      For I = 1 to L
                                                              If (random(1,1000) % 2 == 0)
              While (I > 1)
                     Sum = Sum + 1
                                                                     X += I \times \mathbf{Fun}(\frac{Z}{L}, L)
                      I = I - 1
              EndWhile
                                                                    X += (I + 1) \times \mathbf{Fun}(\frac{Z}{2 \times I}, L)
              I = N
       EndFor
                                                      EndFor
       Return 5 \times Fun(N, 4)
                                               End Fun
End Main
                                               //random(1,1000): generates a random integer
                                               between 1 & 1000 in O(1)
```

- 1. What's the complexity of the non-recursive part of code? theta(n)
- 2. What's the recurrence equation (T(N)) that represents the LOWER BOUND of the entire code? T(N) = 4T(N/8) + theta(1); --> theta(n^0.6667)
- 3. What's the LOWER BOUND complexity of entire code? theta(n)
- **4.** What's the recurrence equation (T(N)) that represents the UPPER BOUND of the entire code? T(N) = 4T(N/4) + theta(1); --> theta(n)
- **5.** What's the UPPER BOUND complexity of entire code? theta(n)

so, as lower bound complexity == upper bound complexity then, theta of entire code = n;



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## PART III: **Asymptotic Notations**

## QUESTION#1

202011111	_				
Chose the notat	tion that <b>BEST</b>	REPRESEN	TS the		
<b>WORST</b> case?					
Fun2 (N)				<b>a.</b> O(1)	
Sorte	d = true			<b>b.</b> Θ(1)	
I = 1				<b>c.</b> O(N)	
While (I < N AND Sorted ==		d ==	<b>d.</b> Θ(N) e. Ω(1)		
true)					
If (A[I] > A[I + 1])					
	Sorte	d = fal	se	f. Ω(N)	
I	End If			g. Θ(N-1)	
I = I + 1			<b>h.</b> None of the choices		
End wh					
Print	Sorted				
QUESTION#2		NI) – NI vyh	at's the nor	ssible value(s) of N <sub>0</sub> that make	c f(NI) -
$\Omega(g(N))$ for cons		IN) – IN, WII	at s the pos	ssible value(s) of N <sub>0</sub> that make:	S I(IV) -
<b>a.</b> 1≤ N0≤32	<b>b.</b> 0< N0 ≤ 2	<b>c.</b> N0 ≥ 2	<b>d.</b> N0 ≥ 4	e. No values exist. The relation	n is no
				valid	
0.1.1.0.0.1.1.1.0		•	•		
QUESTION#3					
Given the follow	ving orders ( <i>lo</i>	<b>g</b> is base 2	):		
$1)\sqrt{e^{\log(N)}+N}$	$\overline{I^3}$ , 2) $1/\sqrt{\log}$	$\overline{(N)}$ ,3) $N$	$(4) N^3, 5)$	$4^{\log(N)}, 6) N \log(N)$	
Arrange them in	n increasing or	der of gro	wth rate (v	vith g(n) following f(n) in your l	ist if
and only if f(n)=	O(g(n))). Chos	e the corre	ct order?		
226544	224654	2.63	4.5.4.	f. N	one of
a. 2,3,6,5,1,4 k	<b>2</b> ,3,1,6,5,4	<b>c.</b> 2,6,3,	1,5,4 <b>a.</b> 2	2,3,6,1,5,4 <b>e.</b> 4,5,1,3,6,2	hoice



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### **PART IV: Recurrence Equations**

#### **QUESTION#1**

Using the RECURSION TREE method, Answer the following questions:

 $T(N) = T(N/5) + T(7N/10) + \Theta(N); T(1) = 1$ 

	Question	Answer
1.	What's the EXACT total number of levels in this tree?	lg10/7(n)
2.	What's the complexity of each level?	(9/10)^i * n
3.	What's the complexity of LAST level?	1
4.	What's the UPPER BOUND order of this tree?	O( n )

#### **QUESTION#2**

Given T(N) = 64 T(N / 16) + (22/7) NVN, using Master Method: what's the correct value of E

- 0 < ε ≤ 1.5
- 0 < ε ≤ 0.5
- ε = 1.6
- No possible ε

It's case2: no & is required

#### **QUESTION#3**

1. Select the correct asymptotic complexity of an algorithm with runtime T(n; n) where:

$$\begin{array}{lll} T(x,c) &=& \Theta(x) & \text{for } c \leq 2, \\ T(c,y) &=& \Theta(y) & \text{for } c \leq 2, \text{ and } \\ T(x,y) &=& \Theta(x+y) + T(x/2,y/2). \end{array}$$

Case 3



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- 1.  $\Theta(\log n)$ .
- $\underbrace{2.}_{3.}\Theta(n).$   $\Theta(n\log n).$ 

  - 4.  $\Theta(n \log^2 n)$ .
  - 5.  $\Theta(n^2)$ .
  - 6.  $\Theta(2^n)$ .

**Solution:** The correct answer is  $\Theta(n)$ . To see why, we rewrite the recurrence relation to avoid  $\Theta$  notation as follows:

$$T(x,y) = c(x + y) + T(x/2, y/2).$$

We may then begin to replace T(x/2, y/2) with the recursive formula containing it:

$$T(x,y) = c(x+y) + c\left(\frac{x+y}{2}\right) + c\left(\frac{x+y}{4}\right) + c\left(\frac{x+y}{8}\right) + \dots$$

This geometric sequence is bounded from above by 2c(x + y), and is obviously bounded from below by c(x+y). Therefore, T(x,y) is  $\Theta(x+y)$ , and so T(n,n)is  $\Theta(n)$ .

2. Select the correct asymptotic complexity of an algorithm with runtime T(n; n) where:

$$T(x,c) = \Theta(x)$$
 for  $c \le 2$ ,

$$T(c,y) = \Theta(y)$$
 for  $c \le 2$ , and

$$T(x,y) = \Theta(x) + T(x,y/2).$$

- 1.  $\Theta(\log n)$ .
- 2.  $\Theta(n)$ .

Case 3

- 3.  $\Theta(n \log n)$ .
- 4.  $\Theta(n \log^2 n)$ .
- 5.  $\Theta(n^2)$ .
- 6.  $\Theta(2^n)$ .



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**Solution:** The correct answer is  $\Theta(n \log n)$ . To see why, we rewrite the recurrence relation to avoid  $\Theta$  notation as follows:

$$T(x,y) = cx + T(x,y/2).$$

We may then begin to replace T(x, y/2) with the recursive formula containing it:

$$T(x,y) = \underbrace{cx + cx + cx + \ldots + cx}_{\Theta(\log y) \text{ times}}.$$

As a result, T(x,y) is  $\Theta(x \log y)$ . When we substitute n for x and y, we get that T(n,n) is  $\Theta(n \log n)$ .

**3.** Select the correct asymptotic complexity of an algorithm with runtime T(n; n) where:

$$\begin{array}{lll} T(x,c) &=& \Theta(x) & \text{for } c \leq 2, \\ T(x,y) &=& \Theta(x) + S(x,y/2), \\ S(c,y) &=& \Theta(y) & \text{for } c \leq 2, \text{and} \\ S(x,y) &=& \Theta(y) + T(x/2,y). \end{array}$$

- 1.  $\Theta(\log n)$ .
- $2. \Theta(n).$
- 3.  $\Theta(n \log n)$ .
- 4.  $\Theta(n \log^2 n)$ .
- 5.  $\Theta(n^2)$ .
- 6.  $\Theta(2^n)$ .