

## Combinatorics

- Concerned with counting various objects, configurations, etc.
- Often solved by recurrences.
- Make sure the data type is large enough to store the result (e.g. use BigInteger).

## Fibonacci Numbers

- Fibonacci numbers grow very quickly.
- $f(0) = 0$ ,  $f(1) = 1$ ,  $f(n) = f(n-1) + f(n-2)$
- An alternate way to compute Fibonacci numbers (or other recurrences):

$$\begin{bmatrix} f(n) \\ f(n+1) \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} \cdot \begin{bmatrix} f(n-1) \\ f(n) \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}^n \cdot \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

- Use fast exponentiation to compute matrix power in  $O(\log n)$  steps.
- See `exp.cc` or `expmod.cc` for fast exponentiation of integers.
- Useful for computing  $f(n)$  for large  $n$ .

## Binomial Coefficients

- Compute with dynamic programming (Pascal's triangle).
- If answer does not overflow, no overflow can occur during the calculation using DP.

## Catalan Numbers

$$C(n) = \frac{1}{n+1} \binom{2n}{n}$$

- Number of binary trees with  $n$  nodes.
- Number of ways to bracket an expression of  $n$  operators.
- Number of ways to write  $n$  pairs of parenthesis which are correctly matched.
- Number of ways to triangulate a convex polygon with  $n + 2$  sides.
- Recurrence:

$$C(n+1) = \sum_{i=0}^n C(i)C(n-i)$$

## Prime Numbers

- A positive integer  $\geq 2$  is prime if its only divisors are 1 and itself.
- To determine if a number  $n$  is prime, use trial division up to  $\sqrt{n}$ . This is good enough for most 32-bit integers.

## Sieve of Eratosthenes

- If we want to test for primality in a range, it is inefficient to test each one by trial division.
- Given a prime  $p$ , it is easier to mark off the multiples of  $p$ .
- Set up a boolean array for the range, and initialize all of them to be true.
- For each relevant prime  $p$ , mark off the entries for  $p, 2p, 3p, \dots$
- Some optimization: consider primes in increasing order, and start from  $p \cdot p, (p + 1) \cdot p, \dots$
- Complexity to check the first  $n$  integers:  $O(n \log \log n)$ .

## GCD

- $\gcd(a, b)$  is the largest (positive) integer dividing both  $a$  and  $b$ .
- Computed by the Euclidean algorithm: `euclid.cc`. Very fast.
- To compute the GCD of three or more integers, use:

$$\gcd(a, b, c) = \gcd(\gcd(a, b), c)$$

## LCM

- LCM is related to the GCD (for two integers):

$$\text{lcm}(a, b) = \frac{a}{\text{gcd}(a, b)} \cdot b$$

It is important to divide by GCD first to avoid overflow.

- For three or more integers:

$$\text{lcm}(a, b, c) = \text{lcm}(\text{lcm}(a, b), c)$$

This is **NOT** the same as:

$$\frac{abc}{\text{gcd}(a, b, c)}$$



## Prime Factorization

- Many problems can be solved easier if we factored the integers into prime factors.
- Relatively easy for “small” (32-bit) integers.
- See `factor.cc`.
- For large integers this is very hard. See `factor_large.cc` if you want to use one possible advanced algorithm.
- See 5.5.5 and 5.5.6 for problems that can be solved easily once factorization is obtained.

## Linear Diophantine Equations

- Given  $a$ ,  $b$ , and  $c$ , you want to find integer solutions  $s$  and  $t$  such that

$$a \cdot s + b \cdot t = c.$$

- Solution exists iff  $\gcd(a, b)$  is a divisor of  $c$ .
- Use the extended Euclidean algorithm (`exteuclid.cc`) to find  $s'$  and  $t'$  such that

$$a \cdot s' + b \cdot t' = \gcd(a, b)$$

- Multiply solutions by  $c / \gcd(a, b)$  to get one solution.
- If  $(s, t)$  is a solution, the rest of the solutions can be obtained from

$$(s + k \cdot b / \gcd(a, b), t - k \cdot a / \gcd(a, b))$$

where  $k$  is any integer.

- Sometimes we are only interested in solutions in a certain range (e.g.

both are positive).