- 1. In the finite automaton with minimum state deterministic that accepts a given language  $L=\{w \mid w \in \{0,1\} *, \text{ the total number of 0s as well as 1s in w that would be divisible by 3 & 5, respectively} would have:$
- a. 9 states
- b. 10 states
- c. 11 states
- d. 15 states

Answer: (d) 15 states

 ${f 2.}$  If we consider an arbitrary NFA (non-deterministic finite automaton) with N states in total, the maximum number of states that are there in an equivalent DFA (minimized) is at

least:
a. N!
b. 2N
c. 2^N
d. N^2

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Answer: (c) 2^N

3. Which one of these given regular expressions isn't equivalent to this regular expression:

$$(m + n + 0) *$$

a. 
$$(m*n* + o*)*$$

b. 
$$((mn) * + o*) *$$

$$d. (m* + n* + o*)*$$

**Answer:** (b) ((mn) \* + o\*) \*

- **4.** Consider that we have a G ambiguous grammar along with its D disambiguated version. If the language that is recognized by these two grammars is denoted by L(G) and L(D), then which one of these would be true?
- a. L (D) = L (G)
- b. L (D)  $\subset$  L (G)
- c. L (D) is empty
- d. L (D)  $\supset$  L (G)

**Answer:** (a) L (D) = L (G)

- 5. If you consider a regular expression r, in which r = (11 + 111)\* over  $\Sigma$  = {0, 1}, then the number of states in minimal DFA and NFA respectively are:
- a. DFA 4, NFA 3
  b. DFA 3, NFA 3
  c. DFA 3, NFA 4
- d. DFA 4, NFA 4

Answer: (a) DFA - 4, NFA - 3

- **6.** The language that a Pushdown Automation accepts in which the stack stays limited to about 10 items is described best as:
- a. Recursive
- b. Deterministic Context Free
- c. Regular
- d. Context Free

Answer: (c) Regular

7. The C language is a:

- a. Regular language
- b. Context free language
- c. Language parsable fully by a Turing machine only
- d. Context sensitive language

Answer: (b) Context free language

8. Consider the language given below:

$$\{a^m b^n C^m + n \mid m, n \ge 1\}$$

It is a \_\_\_\_\_ language.

- a. regular
- b. not context free but context sensitive
- c. not context sensitive but type-0
- d. not regular but context-free

Answer: (d) not regular but context-free

9. Which of these is a regular set?

I. 
$$\{a^nb^{2m} \mid n \ge 0, m \ge 0\}$$

II. 
$$\{a^nb^m \mid n = 2m\}$$

III. 
$$\{a^nb^m \mid n \neq m\}$$

IV. 
$$\{xcy \mid x,y \in \{a,b\}^*\}$$

- a. I
- b. IV
- c. I and III
- d. I and IV

Answer: (d) I and IV

10. Consider the following languages:

 $L1 = \{0^i1^j \mid i != 2j\}$ 

 $L2 = \{0^i1^j \mid i = 2j+1\}$ 

 $L3 = \{0^i1^j \mid i = j\}$ 

 $L4 = \{0^i1^j | i != j\}$ 

Which of these is/are context free:

- a. Only L3
- b. Only L3 & L2
- c. Only L4 & L3
- d. All LA, L2, L3, and L4

Answer: (d) All LA, L2, L3, and L4

11. The L=  $\{0^i21^i \mid i \ge 0\}$  language over the  $\{0,1,2\}$  alphabet is:

a. a CFL but not a deterministic CFL

- b. a regular language
- c. is recursive as well as a deterministic CFL
- d. not recursive

Answer: (c) is recursive as well as a deterministic CFL

12. If G is the CFG, r is the total number of rightmost derivations, l is the total number of leftmost derivations, as well as P refers to the total number of parse trees, then assume that r, l, and P are computed for some given particular string. Here, for a given 'G' CFG and given 'w' string, what is the relation between all three of these?

a.  $r \le P \ge 1$ 

b. r = P = 1

- c.  $r \ge P \le 1$
- d. None

**Answer:** (b) r = P = 1

- 13. When L and L' happen to be recursively enumerable, here L is:
- a. context-free
- b. regular
- c. recursive
- d. context-sensitive

Answer: (c) recursive

- 14. Which of these problems given below is undecidable?
- a. The Ambiguity problem for the CFGs
- b. The Membership problem for the CFGs
- c. The Equivalence problem for the FSAs
- d. The Finiteness problem for the FSAs

Answer: (a) The Ambiguity problem for the CFGs

- **15.** If L refers to the language that is generated by S OSO/OO, then which of these is true?
- a. L is not O but regular
- b. L is regular but not context-free
- c. L is not context-free
- d. L = 0

Answer: (b) L is regular but not context-free