

# A Tactical Communication Protocol Link32 and Reference Implementation Skynet

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## **Abstract**

Link32 is a tactical communication protocol designed for low-latency, secure, and scalable data exchange in contested environments, enabling swarm drone coordination, real-time position location information (PLI), command and control (C2), and tactical chat over UDP-based multicast networks. Skynet, its reference implementation, is a lightweight C99 framework inspired by LTE's QoS Class Identifier (QCI) framework, tailored for military applications and resource-constrained devices such as drones and embedded systems. With minimal dependencies and a focus on portability, Skynet supports dynamic TDMA slot management, AES-256-GCM encryption, and lock-free concurrency, ensuring robust performance in mission-critical scenarios.

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# 1 Introduction

Link32 is a tactical communication protocol inspired by standards such as VMF, Link16, TSM, SRW, and MQTT, tailored for military applications requiring robust, low-latency, and secure data exchange in contested environments. It facilitates swarm drone coordination, real-time position location information (PLI), command and control (C2), and tactical chat using UDP-based multicast networks. Skynet is the reference server implementation of Link32, developed in C99 with minimal dependencies to ensure portability and performance on resource-constrained devices. Its convergence layer draws inspiration from LTE's QoS Class Identifier (QCI) framework, enabling granular QoS control and dynamic resource allocation for mission-critical communications.

## 1.1 Properties

Link32 and Skynet are designed with the following properties:

- **Implementation:** Written in C99 for portability and performance.
- **Message Size:** Minimum 32-byte header (48 bytes with AES-256-GCM tag) to optimize bandwidth.
- **Security:** ECDH key exchange over secp384r1, AES-256-GCM encryption for all messages.
- **Latency:** Microsecond-precision timing using monotonic clocks and non-blocking I/O.
- **Concurrency:** Lock-free atomic operations (CMPXCHG) for thread-safe queue management.
- **Networking:** UDP multicast with topic-based subscriptions.
- **Swarm Scalability:** Supports large-scale drone swarms via dynamic slot allocation.
- **Footprint:** ~64KB L1 cache usage, ~2000 lines of code (LOC).
- **Dependencies:** Single dependency on OpenSSL for cryptography.
- **Threat Model:** Prioritizes, confidentiality and integrity, no non-repudiation (HMAC may be included in payload).

## 1.2 Principles

Link32 adheres to the following design principles:

- **Key Provisioning:** Manual public key distribution for controlled setup.
- **Mandatory Encryption:** All messages encrypted with AES-256-GCM.
- **Node Identification:** Node names hashed to 32-bit using FNV-1a.
- **Lock-Free Design:** Uses atomic compare-and-swap for concurrency.
- **Topic Architecture:** Topics map to IP multicast groups.
- **Queue Management:** Global network queue, per-topic subscribes.
- **Key Storage:** Separate key stores per executable.

## 2 Link32 Protocol

### 2.1 S-Message Format

The Link32 message structure, `SkyNetMessage`, is compact for large-scale swarm communication:

```
typedef struct {
    uint8_t version : 4;    // Protocol version (current: 1)
    uint8_t type : 4;       // Message type (0-6)
    uint8_t qos : 4;        // Quality of Service (0-3)
    uint8_t hop_count : 4;  // Hop count for routing (0-15)
    uint32_t npg_id;        // Topic identifier (1-103)
    uint32_t node_id;       // Sender node ID (FNV-1a hash)
    uint32_t seq_no;        // Sequence number for deduplication
    uint8_t iv[16];         // AES-256-GCM initialization vector
    uint16_t payload_len;   // Payload length (0-32767)
    uint8_t payload[MAX_BUFFER]; // GCM-TAG
} SkyNetMessage;
```

- **Header Size:** 32 bytes.
- **Total Size:** 48 bytes minimum (32-byte header + 16-byte GCM tag).
- **Payload:** Up to 32720 bytes.

## 2.2 Message Types

The protocol defines seven message types, as shown in Table 1.

Table 1: Link32 Message Types		
ID	Type	Description
0	Key Exchange	Exchanges ECC public keys for ECDH session setup.
1	Slot Request	Requests a TDMA slot from the server.
2	Chat	Sends tactical chat messages.
3	Ack	Acknowledges slot assignments or control messages.
4	Waypoint	Specifies navigation waypoints for C2.
5	Status	Reports position, velocity, or sensor data (e.g., PLI).
6	Formation	Coordinates swarm formations.

## 2.3 Multicast Topics

Link32 uses multicast topics mapped to IP multicast groups, as in Table 2.

Table 2: Multicast Topics			
NPG	Name	Multicast	Purpose
1	npg.control	239.255.0.1	Kkey exchange, Slot requests.
6	npg.pli	239.255.0.6	Position information (status).
7	npg.surveillance	239.255.0.7	Sensor data forwarding.
29	npg.chat	239.255.0.29	Tactical chat and acks.
100	npg.c2	239.255.0.100	C2 (waypoints, formations).
101	npg.alerts	239.255.0.101	Network alerts and self-healing.
102	npg.logistics	239.255.0.102	Logistical coordination.
103	npg.coord	239.255.0.103	Swarm Coordination.

## 2.4 Slot Management

Link32 employs a Time Division Multiple Access (TDMA)-like slot manager to minimize collisions:

- **Slot Array:** Fixed-size array (`slots[SLOT_COUNT=256]`) in `ServerState`.
- **Dynamic Topics:** Each slot creates a temporary multicast group (`239.255.1.slot_id % 256`).
- **Allocation:** First-come, first-serve with no timeouts.
- **Timing:** Slots cycle every `TIME_SLOT_INTERVAL_US=1000 $\mu$ s`.

Clients send `SKYNET_MSG_SLOT_REQUEST` to NPG 1, and the server assigns slots via `SKYNET_MSG_ACK`.

## 2.5 Deduplication

A fixed-size circular buffer (`seq_cache`) prevents message loops:

- **Structure:** Stores `node_id`, `seq.no`.
- **Memory:**  $\sim 16\text{KB}$  ( $1024 \times 16$  bytes).
- **Complexity:**  $O(1)$  lookup using FNV-1a hashing.
- **Threshold:** Discards duplicates within 1 second.

## 2.6 Security

Security mechanisms include:

- **Key Exchange:** ECDH over `secp384r1` for 256-bit AES keys.
- **Encryption:** AES-256-GCM with 16-byte IV and 16-byte tag.
- **Key Storage:** Per Process in `ec_priv` and `ec_pub`.
- **Key Derivation:** HKDF-SHA256 for AES keys.
- **Self-messages:** Skips (drops) messages.

## 2.7 Subscriptions

Nodes subscribe to topics based on roles, as shown in Table 3:

Table 3: Role-Based Subscriptions		
Role	NPGs	Purpose
Infantry	1, 29	Network control and tactical chat.
Drone	1, 6, 7, 100, 101	C2, PLI, surv., alerts.
Air	1, 6, 7, 100, 101, 103	C2, PLI, surv., alerts, coord.
Sea	1, 7, 29, 102, 103	C2, surveillance, chat, logistics, coord.
Ground	1, 7, 29, 102	C2, surveillance, chat, logistics.
Relay	1, 6, 101	C2, PLI, alerts for relaying.
Controller	1, 6, 100, 101	C2, PLI, alerts for command posts.

## 3 Skynet Implementation

### 3.1 Dependencies

- **OpenSSL:** For ECC, ECDH, and AES-256-GCM.
- **C99 Compiler:** GCC or equivalent.
- **POSIX Environment:** For threading, epoll, and timerfd.

### 3.2 Build

To build Skynet:

```
$ git clone git@github.com:BitEdits/skynet
$ cd skynet
$ gcc -o skynet_client skynet_client.c skynet_proto.c -lcrypto
$ gcc -o skynet skynet.c skynet_proto.c skynet_conv.c -lcrypto
```

### 3.3 Installation

The provisioning script `skynet.sh` generates ECC key pairs:

```
# ./skynet.sh
Generated keys for node npg_control (hash: 06c5bc52) in /secp/
Generated keys for node npg_pli (hash: c9aef284) in /secp/
Generated keys for node npg_surveillance (hash: 4d128cdc) in /secp/
Generated keys for node npg_chat (hash: 9c69a767) in /secp/
Generated keys for node npg_c2 (hash: 89f28794) in /secp/
Generated keys for node npg_alerts (hash: 9f456bca) in /secp/
Generated keys for node npg_logistics (hash: 542105cc) in /secp/
Generated keys for node npg_coord (hash: e46c0c22) in /secp/
Generated keys for node server (hash: 40ac3dd2) in /secp/
Generated keys for node client (hash: 8f929c1e) in /client/secp/
# cp /secp/*.ec_pub /client/secp/
```

### 3.4 Server Operation

The server binds to UDP port 6566, joins multicast groups, and processes messages using a global queue. Example output:

```
# skynet server
Node 40ac3dd2 bound to 0.0.0.0:6566.
Joined multicast group 239.255.0.1 (NPG 1: control).
Joined multicast group 239.255.0.6 (NPG 6: PLI).
Message received, from=8f929c1e, to=1, size=231.
Decryption successful, from=8f929c1e, to=1, size=215.
Saved public key for client 8f929c1e.
Assigned slot 0 to node 8f929c1e.
Message received, from=8f929c1e, to=6, size=40.
Decryption successful, from=8f929c1e, to=6, size=24.
Message sent from=8f929c1e, to=6, seq=3, multicast=239.255.1.0.
```

### 3.5 Client Operation

The client joins topic-specific multicast groups and sends key exchange, slot requests, and status messages. Example output:

```
# skynet_client client
Node 8f929c1e connecting to port 6566.
Joined multicast group 239.255.0.1 (NPG 1).
Joined multicast group 239.255.0.6 (NPG 6).
Sent key exchange message to server.
Sent slot request message to server.
Received slot assignment: slot=0.
Joined slot group 2399540.1.
Sent status message: multicast=239.255.1.0,
                        pos=[0.1, 0.1, 0.1],
                        vel=[0.0, 0.0, 0.0],
                        seq=2.
```



### 3.6 Usage

Skynet includes five utilities:

#### Keys Provisioning

Generates ECC key pairs.

```
skynet_keygen <node> [--server|--client]
```

#### Message Encryption

Encrypts a test message to <npg\_id>.sky.

```
skynet_encrypt <sender> <recipient> <file>
```

#### Message Decryption

Decrypts <file.sky>.

```
skynet_decoder <sender> <recipient> <file.sky>
```

#### Skynet Server

Runs the server with FNV-1a hashed <node>.

```
skynet <node>
```

#### Skynet Client

Runs the client with FNV-1a hashed <node>.

```
skynet_client <node>
```

### 3.7 Limitations

- **Slot Scalability:** Fixed SLOT\_COUNT=256 limits nodes to 256.
- **No Retransmission:** Dropped messages are not retransmitted.
- **Key Management:** Manual public key copying required.
- **Deduplication:** SEQ\_CACHE\_SIZE=1024 may cause collisions.

## 4 Convergence Architecture

The Skynet system’s convergence layer employs a 3-level structural hierarchy comprising Quality of Service (QoS), Bearer, and Entity components. This design is driven by the need for granular QoS control, per-node resource management, reliable packet delivery, dynamic slot allocation, and scalability in a Time Division Multiple Access (TDMA)-based tactical network. Inspired by the Long-Term Evolution (LTE) QoS Class Identifier (QCI) framework, the hierarchy ensures military-grade performance, including latency below 50ms for command-and-control (C2) traffic and reliable delivery for critical communications, such as swarm drone coordination. Compared to alternative designs, such as flat or centralized structures, this approach excels in dynamic, resilient scenarios, providing robust and scalable QoS management.

### 4.1 Flat QoS Structure

A flat QoS structure assigns slots directly to Network Protocol Groups (NPGs) based on their QoS profiles, without intermediate bearer or entity layers. For example, an NPG like `SKYNET_NPG_C2` might be statically allocated three slots with QoS level 3. While simpler and requiring less memory, this approach lacks per-node isolation, making it unsuitable for dynamic networks with multiple nodes sharing the same NPG. It also complicates reliable delivery, as there is no mechanism for per-flow reordering or sequence tracking. Additionally, static slot assignments cannot adapt to node arrivals or departures, leading to inefficient resource utilization. The 3-level hierarchy overcomes these limitations by introducing bearers for flow isolation and entities for node-level coordination, enabling dynamic and scalable QoS management.

```
typedef struct {
    uint32_t npg_id;
    uint8_t qos;
    uint32_t slot_count;
    uint32_t slot_ids[MAX_QOS_SLOTS];
    uint8_t priority;
} QoSSlotAssignment;
```

### 4.2 Hierarchical QoS Structure

The 3-level hierarchy separates concerns into QoS, Bearer, and Entity layers, each addressing specific aspects of network management. The `SkyNetBearerQoS` structure defines QoS parameters such as priority (1–15, where 1 is highest), delay budget (in milliseconds), reliability (best-effort or reliable), and minimum TDMA slots, allowing precise service differentiation. The `SkyNetBearer` represents a logical communication channel for a node-NPG pair, encapsulating QoS parameters, assigned slots, and state for reordering and reliability. The `SkyNetConvergenceEntity` aggregates bearers for a single node, managing slot requests and coordinating resource allocation. This modular design ensures

scalability, flexibility, and robustness, outperforming flat structures by providing flow isolation and dynamic adaptation.

```
typedef struct {
    uint8_t priority;           // 1-15 (1 = highest)
    uint32_t delay_budget_ms;   // Delay tolerance (ms)
    uint8_t reliability;        // 0 (best-effort), 1 (reliable)
    uint32_t min_slots;         // Minimum TDMA slots
} SkyNetBearerQoS;

typedef struct {
    uint32_t bearer_id;         // Unique bearer ID
    SkyNetBearerQoS qos;        // QoS parameters
    uint32_t node_id;           // Owning node
    uint32_t npg_id;            // Associated NPG
    uint32_t assigned_slots[SKYNET_MAX_SLOTS]; // Assigned slot IDs
    uint32_t slot_count;        // Number of assigned slots
    SkyNetMessage reorder_queue[SKYNET_REORDER_SIZE];
    uint32_t expected_seq_no;    // Next expected sequence number
    uint32_t last_delivered;     // Last delivered sequence number
    uint64_t last_reorder_time_us; // Last reorder check
} SkyNetBearer;

typedef struct {
    SkyNetBearer bearers[SKYNET_MAX_BEARERS]; // Active bearers
    uint32_t bearer_count; // Number of active bearers
    atomic_uint slot_requests_pending; // Pending slot requests
} SkyNetConvergenceEntity;
```

### 4.3 Granular QoS Control

Tactical networks handle diverse traffic types, such as C2, position location information (PLI), and chat, each with distinct latency, reliability, and bandwidth requirements. A uniform QoS approach fails to meet these needs. The `SkyNetBearerQoS` structure enables granular control by defining specific parameters for each bearer. For instance, `SKYNET_NPG_CONTROL` (NPG 1, QoS 3) is assigned three slots with a low delay budget to ensure timely C2 delivery, while `SKYNET_NPG_CHAT` (NPG 103, QoS 0) receives one slot for best-effort traffic. This differentiation, inspired by LTE QCI, guarantees that high-priority traffic meets stringent military requirements, such as sub-50ms latency for C2, while optimizing resource allocation for lower-priority flows.

### 4.4 Per-Node and Per-Flow Resource Management

Each node in the network may support multiple concurrent flows (e.g., C2, PLI, control) with varying QoS needs. Without isolation, these flows compete for resources, risking contention and degraded performance. The `SkyNetBearer` structure provides flow-level isolation by associating each bearer with a specific node-NPG pair, tracking its assigned slots and QoS parameters. The `SkyNetConvergenceEntity` groups all bearers for a node, enabling centralized

resource management and preventing one flow from starving others. This per-node and per-flow approach ensures efficient slot allocation, supports multiple simultaneous communications, and scales to accommodate dynamic network topologies.

#### 4.5 Reliability and Reordering

TDMA networks may deliver packets out of order due to slot scheduling or re-transmissions, particularly for reliable traffic (e.g., `reliability=1`). The bearer includes a `reorder_queue` and tracks `expected_seq_no` and `last_delivered` to reorder packets and ensure reliable delivery. The entity coordinates reorder checks across bearers using `last_reorder_time_us`, minimizing overhead. This mechanism is critical for applications requiring guaranteed delivery, such as C2 or control messages, and enhances robustness compared to flat structures, which lack per-flow reordering capabilities.

#### 4.6 Dynamic Slot Allocation

Tactical networks operate in dynamic environments where nodes join or leave, and traffic patterns shift. Static slot assignments are inefficient and inflexible. The entity tracks `slot_requests_pending` and manages bearer slot assignments via `assigned_slots` and `slot_count`. The bearer parameter `min_slots` ensures minimum resource guarantees, while the entity facilitates dynamic re-allocation through the `skynet_convergence_schedule_slots` function. Weighted Fair Queuing, driven by bearer priorities, optimizes slot distribution, ensuring high-QoS traffic receives preferential treatment. This adaptability is a key advantage over static or centralized designs.

#### 4.7 Scalability and Modularity

A flat QoS structure becomes unwieldy as the number of nodes and NPGs increases, complicating scheduling and state management. The 3-level hierarchy addresses this by separating concerns: QoS defines service requirements, bearers manage flow-specific state, and entities coordinate node-level convergence. This modular design scales to support large networks, simplifies debugging, and facilitates maintenance. The entity layer aggregates bearer state, reducing scheduling complexity from  $O(n)$  for  $n$  bearers to  $O(m)$  for  $m$  nodes. The hierarchy also supports future enhancements, such as preemption or adaptive QoS, without requiring a system overhaul.

#### 4.8 LTE QCI Compatibility

The 3-level hierarchy draws inspiration from LTE's QCI framework, which uses bearers with QoS profiles to manage diverse traffic types (e.g., VoIP, video, best-effort) per User Equipment (UE). By adapting this model to TDMA, Skynet replaces LTE's EPS bearers with TDMA slot-based bearers, leveraging telecom

best practices. The bearer parameters mirror QCI attributes, such as priority and delay budget, ensuring compatibility with established standards. This alignment reduces design risk, enhances interoperability with telecom systems, and provides a familiar framework for engineers, making it easier to develop and maintain the system.

## 4.9 MQTT Compatibility

Skynet’s convergence layer maps its QoS levels to MQTT’s QoS semantics (ISO/IEC 20922:2016) and LTE’s QCI framework (3GPP TS 23.203, Release 15), enabling compatibility with both standards for tactical communications. MQTT defines three QoS levels:

- **QoS 0 (At Most Once):** Best-effort delivery with no acknowledgment or retransmission. Suitable for non-critical data like tactical chat where packet loss is tolerable.
- **QoS 1 (At Least Once):** Guarantees at least one delivery with acknowledgment (ACK), allowing duplicates. Used for data requiring delivery, such as position location information (PLI), where duplicates are acceptable.
- **QoS 2 (Exactly Once):** Ensures exactly one delivery via a four-way handshake (PUBLISH, PUBREC, PUBREL, PUBCOMP). Critical for command and control (C2) messages requiring no loss or duplication.

Skynet implements these using its `SkyNetBearerQoS` and `SkyNetBearer` structures, configured via `skynet_convergence_init`. QoS 0 skips `reorder_queue` and ACKs, QoS 1 uses ACKs (`SKYNET_MSG_ACK`) with retransmission on timeout (100ms), and QoS 2 employs a handshake with strict ordering and deduplication via `seq_no`. All Network Protocol Groups (NPGs) and message types are mapped to QoS levels, as shown in Table 4.

Table 4: Skynet QoS Mapping to MQTT and LTE QCI

QoS	Priority	NPG	QCI	Budget	Slots	Type
0	15	29	QCI 9	300ms	1	CHAT
0	14	102	QCI 9	300ms	1	CHAT, STATUS
0	13	103	QCI 9	300ms	1	CHAT, WAYPOINT
1	7	6	QCI 7	100ms	2	STATUS
1	6	7	QCI 7	100ms	2	STATUS
2	3	1	QCI 5	50ms	3	EXCHANGE, ACK
2	2	100	QCI 3	50ms	3	FORMATION, ACK
2	1	101	QCI 5	50ms	3	ACK, STATUS

This mapping ensures military-grade performance: QoS 2 meets sub-50ms latency for C2 (`npg_c2`, NPG 100) and control (`npg_control`, NPG 1), QoS 1 supports reliable PLI (`npg_pli`, NPG 6) with 100ms latency, and QoS 0 optimizes

bandwidth for chat (`npg_chat`, NPG 29). The `SkyNetConvergenceEntity` dynamically allocates slots via Weighted Fair Queuing, prioritizing high-QoS bearers, while `reorder_queue` ensures reliability for QoS 1 and 2, aligning with LTE QCI’s packet error loss rates (e.g.,  $10^{-6}$  for QCI 5).

In conclusion, the 3-level hierarchy, inspired by LTE QCI, enables Skynet to replicate MQTT’s QoS semantics, ensuring scalable, low-latency, and reliable communication for swarm drone coordination and other tactical applications.

## 5 Conclusion

Link32 and Skynet provide a robust framework for tactical communication, combining low-latency, security, and scalability for military applications, including swarm drone coordination. Future improvements could address slot scalability beyond 256 nodes, automated key distribution to replace manual provisioning, retransmission mechanisms for dropped messages, and enhanced deduplication to support larger networks.

## References

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