Weisfeiler and Leman Go Neural: Higher-order Graph Neural Networks

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Abstract

In recent years, graph neural networks (GNNs) have emerged as a powerful neural architecture to learn vector representations of nodes and graphs in a supervised, end-to-end fashion. Up to now, GNNs have only been evaluated empirically—showing promising results. The following work investigates GNNs from a theoretical point of view and relates them to the 1-dimensional Weisfeiler-Leman graph isomorphism heuristic (1-WL). We show that GNNs have the same expressiveness as the 1-WL in terms of distinguishing non-isomorphic (sub-)graphs. Hence, both algorithms also have the same shortcomings. Based on this, we propose a generalization of GNNs, so-called k-dimensional GNNs (k-GNNs), which can take higher-order graph structures at multiple scales into account. These higher-order structures play an essential role in the characterization of social networks and molecule graphs. Our experimental evaluation confirms our theoretical findings as well as confirms that higher-order information is useful in the task of graph classification and regression.

Introduction

Graph-structured data is ubiquitous across application domains ranging from chemo- and bioinformatics to image and social network analysis. To develop successful machine learning models in these domains, we need techniques that can exploit the rich information inherent in graph structure, as well as the feature information contained within a graph's nodes and edges. In recent years, numerous approaches have been proposed for machine learning graphs—most notably, approaches based on graph kernels (Vishwanathan et al. 2010) or, alternatively, using graph neural network algorithms (Hamilton, Ying, and Leskovec 2017b).

Kernel approaches typically fix a set of features in advance—e.g., indicator features over subgraph structures or features of local node neighborhoods. For example, one of the most successful kernel approaches, the *Weisfeiler-Lehman subtree kernel* (Shervashidze et al. 2011), which is based on the 1-dimensional Weisfeiler-Leman graph isomorphism heuristic (Grohe 2017, pp. 79 ff.), generates node features through an iterative relabeling, or *coloring*, scheme: First, all nodes are assigned a common initial color; the algorithm then iteratively recolors a node by aggregating over the multiset of colors

in its neighborhood, and the final feature representation of a graph is the histogram of the resulting node colors. By iteratively aggregating over local node neighborhoods in this way, the WL subtree kernel is able to effectively summarize the neighborhood substructures present in a graph. However, while powerful, the WL subtree kernel—like other kernel methods—is limited because this feature construction scheme is fixed (i.e., it does not adapt to the given data distribution). Moreover, this approach—like the majority of kernel methods—focuses only on the graph structure and cannot interpret continuous node and edge labels, such as real-valued vectors which play an important role in applications such as bio- and chemoinformatics.

Graph neural networks (GNNs) have emerged as a machine learning framework addressing the above challenges. Standard GNNs can be viewed as a neural version of the 1-WL algorithm, where colors are replaced by continuous feature vectors and neural networks are used to aggregate over node neighborhoods (Hamilton, Ying, and Leskovec 2017a; Kipf and Welling 2017). In effect, the GNN framework can be viewed as implementing a continuous form of graph-based "message passing", where local neighborhood information is aggregated and passed on to the neighbors (Gilmer et al. 2017). By deploying a trainable neural network to aggregate information in local node neighborhoods, GNNs can be trained in an end-to-end fashion together with the parameters of the classification or regression algorithm, possibly allowing for greater adaptability and better generalization compared to the kernel counterpart of the classical 1-WL algorithm.

Up to now, the evaluation and analysis of GNNs has been largely empirical, showing promising results compared to kernel approaches, see, e.g., (Ying et al. 2018b). However, it remains unclear how GNNs are actually encoding graph structure information into their vector representations, and whether there are theoretical advantages of GNNs compared to kernel based approaches.

Present Work. We offer a theoretical exploration of the relationship between GNNs and kernels that are based on the 1-WL algorithm. We show that GNNs cannot be more powerful than the 1-WL in terms of distinguishing non-isomorphic (sub-)graphs, e.g., the properties of subgraphs around each node. This result holds for a broad class of GNN architectures and all possible choices of parameters for them. On the positive side, we show that given the right parameter initialization GNNs have the same expressiveness as the 1-WL algorithm,

completing the equivalence. Since the power of the 1-WL has been completely characterized, see, e.g., (Arvind et al. 2015; Kiefer, Schweitzer, and Selman 2015), we can transfer these results to the case of GNNs, showing that both approaches have the same shortcomings.

Going further, we leverage these theoretical relationships to propose a generalization of GNNs, called k-GNNs, which are neural architectures based on the k-dimensional WL algorithm (k-WL), which are strictly more powerful than GNNs. The key insight in these higher-dimensional variants is that they perform message passing directly between subgraph structures, rather than individual nodes. This higher-order form of message passing can capture structural information that is not visible at the node-level.

Graph kernels based on the k-WL have been proposed in the past (Morris, Kersting, and Mutzel 2017). However, a key advantage of implementing higher-order message passing in GNNs—which we demonstrate here—is that we can design hierarchical variants of k-GNNs, which combine graph representations learned at different granularities in an end-to-end trainable framework. Concretely, in the presented hierarchical approach the initial messages in a k-GNN are based on the output of lower-dimensional k'-GNN (with $k' < k$), which allows the model to effectively capture graph structures of varying granularity. Many real-world graphs inherit a hierarchical structure—e.g., in a social network we must model both the ego-networks around individual nodes, as well as the coarse-grained relationships between entire communities, see, e.g., (Newman 2003)—and our experimental results demonstrate that these hierarchical k-GNNs are able to consistently outperform traditional GNNs on a variety of graph classification and regression tasks. Across twelve graph regression tasks from the QM9 benchmark, we find that our hierarchical model reduces the mean absolute error by 54.45% on average. For graph classification, we find that our hierarchical models leads to slight performance gains.

Key Contributions. Our key contributions are summarized as follows:

1. We show that GNNs are not more powerful than the 1-WL in terms of distinguishing non-isomorphic (sub-)graphs. Moreover, we show that, assuming a suitable parameter initialization, GNNs have the same power as the 1-WL.
2. We propose k-GNNs, which are strictly more powerful than GNNs. Moreover, we propose a hierarchical version of k-GNNs, so-called 1-k-GNNs, which are able to work with the fine- and coarse-grained structures of a given graph, and relationships between those.
3. Our theoretical findings are backed-up by an experimental study, showing that higher-order graph properties are important for successful graph classification and regression.

Related Work

Our study builds upon a wealth of work at the intersection of supervised learning on graphs, kernel methods, and graph neural networks.

Historically, kernel methods—which implicitly or explicitly map graphs to elements of a Hilbert space—have been the dominant approach for supervised learning on graphs.

Important early work in this area includes random-walk based kernels (Gärtner, Flach, and Wrobel 2003; Kashima, Tsuda, and Inokuchi 2003)) and kernels based on shortest paths (Borgwardt and Kriegel 2005). More recently, developments in graph kernels have emphasized scalability, focusing on techniques that bypass expensive Gram matrix computations by using explicit feature maps. Prominent examples of this trend include kernels based on graphlet counting (Shervashidze et al. 2009), and, most notably, the Weisfeiler-Lehman subtree kernel (Shervashidze et al. 2011) as well as its higher-order variants (Morris, Kersting, and Mutzel 2017). Graphlet and Weisfeiler-Leman kernels have been successfully employed within frameworks for smoothed and deep graph kernels (Yanardag and Vishwanathan 2015a; Yanardag and Vishwanathan 2015b). Recent works focus on assignment-based approaches (Kriege, Giscard, and Wilson 2016; Nikolentzos, Meladianos, and Vazirgiannis 2017; Johansson and Dubhashi 2015), spectral approaches (Kondor and Pan 2016), and graph decomposition approaches (Nikolentzos et al. 2018). Graph kernels were dominant in graph classification for several years, leading to new state-of-the-art results on many classification tasks. However, they are limited by the fact that they cannot effectively adapt their feature representations to a given data distribution, since they generally rely on a fixed set of features. More recently, a number of approaches to graph classification based upon neural networks have been proposed. Most of the neural approaches fit into the graph neural network framework proposed by (Gilmer et al. 2017). Notable instances of this model include *Neural Fingerprints* (Duvenaud et al. 2015), *Gated Graph Neural Networks* (Li et al. 2016), *GraphSAGE* (Hamilton, Ying, and Leskovec 2017a), *SplineCNN* (Fey et al. 2018), and the spectral approaches proposed in (Bruna et al. 2014; Defferrard, X., and Vandergheynst 2016; Kipf and Welling 2017)—all of which descend from early work in (Merkwirth and Lengauer 2005) and (Scarselli et al. 2009b). Recent extensions and improvements to the GNN framework include approaches to incorporate different local structures around subgraphs (Xu et al. 2018) and novel techniques for pooling node representations in order perform graph classification (Zhang et al. 2018; Ying et al. 2018b). GNNs have achieved state-of-the-art performance on several graph classification benchmarks in recent years, see, e.g., (Ying et al. 2018b)—as well as applications such as protein-protein interaction prediction (Fout et al. 2017), recommender systems (Ying et al. 2018a), and the analysis of quantum interactions in molecules (Schütt et al. 2017). A survey of recent advancements in GNN techniques can be found in (Hamilton, Ying, and Leskovec 2017b).

Up to this point (and despite their empirical success) there has been very little theoretical work on GNNs—with the notable exceptions of Li et al.'s (Li, Han, and Wu 2018) work connecting GNNs to a special form Laplacian smoothing and Lei et al.'s (Lei et al. 2017) work showing that the feature maps generated by GNNs lie in the same Hilbert space as some popular graph kernels. Moreover, Scarselli et al. (Scarselli et al. 2009a) investigates the approximation capabilities of GNNs.

Preliminaries

We start by fixing notation, and then outline the Weisfeiler-Leman algorithm and the standard graph neural network framework.

Notation and Background

A *graph* G is a pair (V, E) with a finite set of *nodes* V and a set of *edges* $E \subseteq \{\{u, v\} \subseteq V \mid u \neq v\}$. We denote the set of nodes and the set of edges of G by $V(G)$ and $E(G)$, respectively. For ease of notation we denote the edge $\{u, v\}$ in $E(G)$ by (u, v) or (v, u). Moreover, $N(v)$ denotes the *neighborhood* of v in $V(G)$, i.e., $N(v) = \{u \in V(G) \mid (v, u) \in E(G)\}$. We say that two graphs G and H are *isomorphic* if there exists an edge preserving bijection $\varphi : V(G) \to V(H)$, i.e., (u, v) is in $E(G)$ if and only if $(\varphi(u), \varphi(v))$ is in $E(H)$. We write $G \simeq H$ and call the equivalence classes induced by $\simeq$ *isomorphism types*. Let $S \subseteq V(G)$ then $G[S] = (S, E_S)$ is the *subgraph induced* by S with $E_S = \{(u, v) \in E(G) \mid u, v \in S\}$. A *node coloring* is a function $V(G) \to \Sigma$ with arbitrary codomain Σ. Then a *node colored* or *labeled graph* (G, l) is a graph G endowed with a node coloring $l : V(G) \to \Sigma$. We say that $l(v)$ is a *label* or *color* of $v \in V(G)$. We say that a node coloring c *refines* a node coloring d, written $c \sqsubseteq d$, if $c(v) = c(w)$ implies $d(v) = d(w)$ for every v, w in $V(G)$. Two colorings are *equivalent* if $c \sqsubseteq d$ and $d \sqsubseteq c$, and we write $c \equiv d$. A *color class* $Q \subseteq V(G)$ of a node coloring c is a maximal set of nodes with $c(v) = c(w)$ for every v, w in Q. Moreover, let $[1 : n] = \{1, \ldots, n\} \subset \mathbb{N}$ for $n > 1$, let S be a set then the set of k-*sets* $[S]^k = \{U \subseteq S \mid |U| = k\}$ for $k \geq 2$, which is the set of all subsets with cardinality k, and let $\{\!\{\ldots\}\!\}$ denote a multiset.

Weisfeiler-Leman Algorithm

We now describe the 1-WL algorithm for labeled graphs. Let (G, l) be a labeled graph. In each iteration, $t \geq 0$, the 1-WL computes a node coloring $c_l^{(t)} : V(G) \to \Sigma$, which depends on the coloring from the previous iteration. In iteration 0, we set $c_l^{(0)} = l$. Now in iteration $t > 0$, we set

$$c_l^{(t)}(v) = \mathrm{H\,A\,S\,H}\Big(\big(c_l^{(t-1)}(v), \{\!\{c_l^{(t-1)}(u) \mid u \in N(v)\}\!\}\big)\Big) \quad (1)$$

where $\mathrm{H\,A\,S\,H}$ bijectively maps the above pair to a unique value in Σ, which has not been used in previous iterations. To test two graph G and H for isomorphism, we run the above algorithm in "parallel" on both graphs. Now if the two graphs have a different number of nodes colored σ in Σ, the 1-WL concludes that the graphs are not isomorphic. Moreover, if the number of colors between two iterations does not change, i.e., the cardinalities of the images of $c_l^{(t-1)}$ and $c_l^{(t)}$ are equal, the algorithm terminates. Termination is guaranteed after at most $\max\{|V(G)|, |V(H)|\}$ iterations. It is easy to see that the algorithm is not able to distinguish all non-isomorphic graphs, e.g., see (Cai, Fürer, and Immerman 1992). Nonetheless, it is a powerful heuristic, which can successfully test isomorphism for a broad class of graphs (Babai and Kucera 1979).

The k-dimensional Weisfeiler-Leman algorithm (k-WL), for $k \geq 2$, is a generalization of the 1-WL which colors tuples from $V(G)^k$ instead of nodes. That is, the algorithm computes a coloring $c_{l,k}^{(t)} : V(G)^k \to \Sigma$. In order to describe the algorithm, we define the j-th neighborhood

$$N_j(s) = \{(s_1, \ldots, s_{j-1}, r, s_{j+1}, \ldots, s_k) \mid r \in V(G)\} \quad (2)$$

of a k-tuple $s = (s_1, \ldots, s_k)$ in $V(G)^k$. That is, the j-th neighborhood $N_j(t)$ of s is obtained by replacing the j-th component of s by every node from $V(G)$. In iteration 0, the algorithm labels each k-tuple with its *atomic type*, i.e., two k-tuples s and s' in $V(G)^k$ get the same color if the map $s_i \mapsto s_i'$ induces a (labeled) isomorphism between the subgraphs induced from the nodes from s and s', respectively. For iteration $t > 0$, we define

$$C_j^{(t)}(s) = \mathrm{H\,A\,S\,H}\big(\{\!\{c_{l,k}^{(t-1)}(s') \mid s' \in N_j(s)\}\!\}\big), \quad (3)$$

and set

$$c_{k,l}^{(t)}(s) = \mathrm{H\,A\,S\,H}\Big(\big(c_{k,l}^{(t-1)}(s), (C_1^{(t)}(s), \ldots, C_k^{(t)}(s))\big)\Big). \quad (4)$$

Hence, two tuples s and s' with $c_{k,l}^{(t-1)}(s) = c_{k,l}^{(t-1)}(s')$ get different colors in iteration t if there exists j in $[1 : k]$ such that the number of j-neighbors of s and s', respectively, colored with a certain color is different. The algorithm then proceeds analogously to the 1-WL. By increasing k, the algorithm gets more powerful in terms of distinguishing non-isomorphic graphs, i.e., for each $k \geq 2$, there are non-isomorphic graphs which can be distinguished by the $(k+1)$-WL but not by the k-WL (Cai, Fürer, and Immerman 1992). We note here that the above variant is not equal to the *folklore* variant of k-WL described in (Cai, Fürer, and Immerman 1992), which differs slightly in its update rule. However, it holds that the k-WL using Equation (4) is as powerful as the folklore $(k-1)$-WL (Grohe and Otto 2015).

WL Kernels. After running the WL algorithm, the concatenation of the histogram of colors in each iteration can be used as a feature vector in a kernel computation. Specifically, in the histogram for every color σ in Σ there is an entry containing the number of nodes or k-tuples that are colored with σ.

Graph Neural Networks

Let (G, l) be a labeled graph with an initial node coloring $f^{(0)} : V(G) \to \mathbb{R}^{1 \times d}$ that is *consistent* with l. This means that each node v is annotated with a feature $f^{(0)}(v)$ in $\mathbb{R}^{1 \times d}$ such that $f^{(0)}(u) = f^{(0)}(v)$ if and only if $l(u) = l(v)$. Alternatively, $f^{(0)}(v)$ can be an arbitrary real-valued feature vector associated with v. Examples include continuous atomic properties in chemoinformatic applications where nodes correspond to atoms, or vector representations of text in social network applications. A GNN model consists of a stack of neural network layers, where each layer aggregates local neighborhood information, i.e., features of neighbors, around each node and then passes this aggregated information on to the next layer.

A basic GNN model can be implemented as follows (Hamilton, Ying, and Leskovec 2017b). In each layer $t > 0$, we compute a new feature

$$f^{(t)}(v) = \sigma\Big(f^{(t-1)}(v) \cdot W_1^{(t)} + \sum_{w \in N(v)} f^{(t-1)}(w) \cdot W_2^{(t)}\Big) \quad (5)$$

in $\mathbb{R}^{1 \times e}$ for v, where $W_1^{(t)}$ and $W_2^{(t)}$ are parameter matrices from $\mathbb{R}^{d \times e}$, and σ denotes a component-wise non-linear function, e.g., a sigmoid or a ReLU.[1]

Following (Gilmer et al. 2017), one may also replace the sum defined over the neighborhood in the above equation by a permutation-invariant, differentiable function, and one may substitute the outer sum, e.g., by a column-wise vector concatenation or LSTM-style update step. Thus, in full generality a new feature $f^{(t)}(v)$ is computed as

$$f_{\text{merge}}^{W_1}\left(f^{(t-1)}(v), f_{\text{aggr}}^{W_2}\left(\{\!\!\{ f^{(t-1)}(w) \mid w \in N(v) \}\!\!\}\right)\right), \quad (6)$$

where $f_{\text{aggr}}^{W_1}$ aggregates over the set of neighborhood features and $f_{\text{merge}}^{W_2}$ merges the node's representations from step $(t-1)$ with the computed neighborhood features. Both $f_{\text{aggr}}^{W_1}$ and $f_{\text{merge}}^{W_2}$ may be arbitrary differentiable, permutation-invariant functions (e.g., neural networks), and, by analogy to Equation 5, we denote their parameters as W_1 and W_2, respectively. In the rest of this paper, we refer to neural architectures implementing Equation (6) as 1-*dimensional GNN architectures* (1-GNNs).

A vector representation f_{GNN} over the whole graph can be computed by summing over the vector representations computed for all nodes, i.e.,

$$f_{\text{GNN}}(G) = \sum_{v \in V(G)} f^{(T)}(v),$$

where $T > 0$ denotes the last layer. More refined approaches use differential pooling operators based on sorting (Zhang et al. 2018) and soft assignments (Ying et al. 2018b).

In order to adapt the parameters W_1 and W_2 of Equations (5) and (6), to a given data distribution, they are optimized in an end-to-end fashion (usually via stochastic gradient descent) together with the parameters of a neural network used for classification or regression.

Relationship Between 1-WL and 1-GNNs

In the following we explore the relationship between the 1-WL and 1-GNNs. Let (G, l) be a labeled graph, and let $\mathbf{W}^{(t)} = \left(W_1^{(t')}, W_2^{(t')}\right)_{t' \leq t}$ denote the GNN parameters given by Equation (5) or Equation (6) up to iteration t. We encode the initial labels $l(v)$ by vectors $f^{(0)}(v) \in \mathbb{R}^{1 \times d}$, e.g., using a 1-hot encoding.

Our first theoretical result shows that the 1-GNN architectures do not have more power in terms of distinguishing between non-isomorphic (sub-)graphs than the 1-WL algorithm. More formally, let $f_{\text{merge}}^{W_1}$ and $f_{\text{aggr}}^{W_2}$ be any two functions chosen in (6). For every encoding of the labels $l(v)$ as vectors $f^{(0)}(v)$, and for every choice of $\mathbf{W}^{(t)}$, we have that the coloring $c_l^{(t)}$ of 1-WL always refines the coloring $f^{(t)}$ induced by a 1-GNN parameterized by $\mathbf{W}^{(t)}$.

Theorem 1. Let (G, l) be a labeled graph. Then for all $t \geq 0$ and for all choices of initial colorings $f^{(0)}$ consistent with l, and weights $\mathbf{W}^{(t)}$,

$$c_l^{(t)} \sqsubseteq f^{(t)}.$$

Our second result states that there exist a sequence of parameter matrices $\mathbf{W}^{(t)}$ such that 1-GNNs have exactly the same power in terms of distinguishing non-isomorphic (sub-)graphs as the 1-WL algorithm. This even holds for the simple architecture (5), provided we choose the encoding of the initial labeling l in such a way that different labels are encoded by linearly independent vectors.

Theorem 2. Let (G, l) be a labeled graph. Then for all $t \geq 0$ there exists a sequence of weights $\mathbf{W}^{(t)}$, and a 1-GNN architecture such that

$$c_l^{(t)} \equiv f^{(t)}.$$

Hence, in the light of the above results, 1-GNNs may viewed as an extension of the 1-WL which in principle have the same power but are more flexible in their ability to adapt to the learning task at hand and are able to handle continuous node features.

Shortcomings of Both Approaches

The power of 1-WL has been completely characterized, see, e.g., (Arvind et al. 2015). Hence, by using Theorems 1 and 2, this characterization is also applicable to 1-GNNs. On the other hand, 1-GNNs have the same shortcomings as the 1-WL. For example, both methods will give the same color to every node in a graph consisting of a triangle and a 4-cycle, although vertices from the triangle and the vertices from the 4-cycle are clearly different. Moreover, they are not capable of capturing simple graph theoretic properties, e.g., triangle counts, which are an important measure in social network analysis (Milo et al. 2002; Newman 2003).

k-dimensional Graph Neural Networks

In the following, we propose a generalization of 1-GNNs, so-called k-GNNs, which are based on the k-WL. Due to scalability and limited GPU memory, we consider a set-based version of the k-WL. For a given k, we consider all k-element subsets $[V(G)]^k$ over $V(G)$. Let $s = \{s_1, \ldots, s_k\}$ be a k-set in $[V(G)]^k$, then we define the *neighborhood* of s as

$$N(s) = \{t \in [V(G)]^k \mid |s \cap t| = k - 1\}.$$

The *local neighborhood* $N_L(s)$ consists of all $t \in N(s)$ such that $(v, w) \in E(G)$ for the unique $v \in s \setminus t$ and the unique $w \in t \setminus s$. The *global neighborhood* $N_G(s)$ then is defined as $N(s) \setminus N_L(s)$.[2]

The set based k-WL works analogously to the k-WL, i.e., it computes a coloring $c_{s,k,l}^{(t)} : [V(G)]^k \to \Sigma$ as in **??** based on the above neighborhood. Initially, $c_{s,k,l}^{(0)}$ colors each element s in $[V(G)]^k$ with the isomorphism type of $G[s]$.

Let (G, l) be a labeled graph. In each k-GNN layer $t \geq 0$, we compute a feature vector $f_k^{(t)}(s)$ for each k-set s in $[V(G)]^k$. For $t = 0$, we set $f_k^{(0)}(s)$ to $f^{\text{iso}}(s)$, a one-hot

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[1]For clarity of presentation we omit biases.

[2]Note that the definition of the local neighborhood is different from the the one defined in (Morris, Kersting, and Mutzel 2017) which is a superset of our definition. Our computations therefore involve sparser graphs.
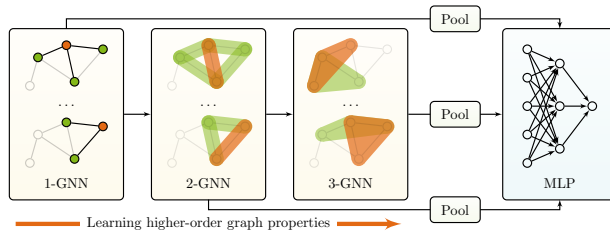
(a) Hierarchical 1-2-3-GNN network architecture
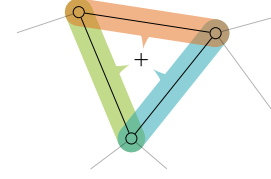
(b) Pooling from 2- to 3-GNN.

Figure 1: Illustration of the proposed hierarchical variant of the k-GNN layer. For each subgraph S on k nodes a feature f is learned, which is initialized with the learned features of all $(k-1)$-element subgraphs of S. Hence, a hierarchical representation of the input graph is learned.

encoding of the isomorphism type of $G[s]$ labeled by l. In each layer $t > 0$, we compute new features by

$$f_k^{(t)}(s) = \sigma\left(f_k^{(t-1)}(s) \cdot W_1^{(t)} + \sum_{u \in N_L(s) \cup N_G(s)} f_k^{(t-1)}(u) \cdot W_2^{(t)}\right).$$

Moreover, one could split the sum into two sums ranging over $N_L(s)$ and $N_G(s)$ respectively, using distinct parameter matrices to enable the model to learn the importance of local and global neighborhoods. To scale k-GNNs to larger datasets and to prevent overfitting, we propose *local k-GNNs*, where we omit the global neighborhood of s, i.e.,

$$f_{k,\mathrm{L}}^{(t)}(s) = \sigma\left(f_{k,\mathrm{L}}^{(t-1)}(s) \cdot W_1^{(t)} + \sum_{u \in N_L(s)} f_{k,\mathrm{L}}^{(t-1)}(u) \cdot W_2^{(t)}\right).$$

The running time for evaluation of the above depends on $|V|$, k and the sparsity of the graph (each iteration can be bounded by the number of subsets of size k times the maximum degree). Note that we can scale our method to larger datasets by using sampling strategies introduced in, e.g., (Morris, Kersting, and Mutzel 2017; Hamilton, Ying, and Leskovec 2017a). We can now lift the results of the previous section to the k-dimensional case.

Proposition 3. Let (G, l) be a labeled graph and let $k \geq 2$. Then for all $t \geq 0$, for all choices of initial colorings $f_k^{(0)}$ consistent with l and for all weights $\mathbf{W}^{(t)}$,

$$c_{s,k,l}^{(t)} \sqsubseteq f_k^{(t)}.$$

Again the second result states that there exists a suitable initialization of the parameter matrices $\mathbf{W}^{(t)}$ such that k-GNNs have exactly the same power in terms of distinguishing non-isomorphic (sub-)graphs as the set-based k-WL.

Proposition 4. Let (G, l) be a labeled graph and let $k \geq 2$. Then for all $t \geq 0$ there exists a sequence of weights $\mathbf{W}^{(t)}$, and a k-GNN architecture such that

$$c_{s,k,l}^{(t)} \equiv f_k^{(t)}.$$

Hierarchical Variant

One key benefit of the end-to-end trainable k-GNN framework—compared to the discrete k-WL algorithm—is that we can hierarchically combine representations learned at different

granularities. Concretely, rather than simply using one-hot indicator vectors as initial feature inputs in a k-GNN, we propose a *hierarchical* variant of k-GNN that uses the features learned by a $(k-1)$-dimensional GNN, in addition to the (labeled) isomorphism type, as the initial features, i.e.,

$$f_k^{(0)}(s) = \sigma\left(\left[f^{\mathrm{iso}}(s), \sum_{u \subset s} f_{k-1}^{(T_{k-1})}(u)\right] \cdot W_{k-1}\right),$$

for some $T_{k-1} > 0$, where W_{k-1} is a matrix of appropriate size, and square brackets denote matrix concatenation.

Hence, the features are recursively learned from dimensions 1 to k in an end-to-end fashion. This hierarchical model also satisfies Propositions 3 and 4, so its representational capacity is theoretically equivalent to a standard k-GNN (in terms of its relationship to k-WL). Nonetheless, hierarchy is a natural inductive bias for graph modeling, since many real-world graphs incorporate hierarchical structure, so we expect this hierarchical formulation to offer empirical utility.

Experimental Study

In the following, we want to investigate potential benefits of GNNs over graph kernels as well as the benefits of our proposed k-GNN architectures over 1-GNN architectures. More precisely, we address the following questions:

Q1 How do the (hierarchical) k-GNNs perform in comparison to state-of-the-art graph kernels?

Q2 How do the (hierarchical) k-GNNs perform in comparison to the 1-GNN in graph classification and regression tasks?

Q3 How much (if any) improvement is provided by optimizing the parameters of the GNN aggregation function, compared to just using random GNN parameters while optimizing the parameters of the downstream classification/regression algorithm?

Datasets

To compare our k-GNN architectures to kernel approaches we use well-established benchmark datasets from the graph kernel literature (Kersting et al. 2016). The nodes of each graph in these dataset is annotated with (discrete) labels or no labels.

To demonstrate that our architectures scale to larger datasets and offer benefits on real-world applications, we conduct experiments on the QM9 dataset (Ramakrishnan et al. 2014; Ruddigkeit et al. 2012; Wu et al. 2018), which consists of

Table 1: Classification accuracies in percent on various graph benchmark datasets.

	Method	Dataset						
		PRO	IMDB-BIN	IMDB-MUL	PTC-FM	NCI1	MUTAG	PTC-MR
Kernel	GRAPHLET	72.9	59.4	40.8	58.3	72.1	87.7	54.7
	SHORTEST-PATH	**76.4**	59.2	40.5	62.1	74.5	81.7	58.9
	1-WL	73.8	72.5	**51.5**	62.9	83.1	78.3	61.3
	2-WL	75.2	72.6	50.6	**64.7**	77.0	77.0	61.9
	3-WL	74.7	73.5	49.7	61.5	83.1	83.2	62.5
	WL-OA	75.3	73.1	50.4	62.7	**86.1**	84.5	**63.6**
GNN	DCNN	61.3	49.1	33.5	—	62.6	67.0	56.6
	PATCHYSAN	75.9	71.0	45.2	—	78.6	**92.6**	60.0
	DGCNN	75.5	70.0	47.8	—	74.4	85.8	58.6
	1-GNN NO TUNING	70.7	69.4	47.3	59.0	58.6	82.7	51.2
	1-GNN	72.2	71.2	47.7	59.3	74.3	82.2	59.0
	1-2-3-GNN NO TUNING	75.9	70.3	48.8	60.0	67.4	84.4	59.3
	1-2-3-GNN	75.5	**74.2**	49.5	62.8	76.2	86.1	60.9

Table 2: Mean absolute errors on the QM9 dataset. The far-right column shows the improvement of the best k-GNN model in comparison to the 1-GNN baseline.

Target	Method						
	DTNN (Wu et al. 2018)	MPNN (Wu et al. 2018)	1-GNN	1-2-GNN	1-3-GNN	1-2-3-GNN	Gain
μ	**0.244**	0.358	0.493	0.493	**0.473**	0.476	4.0%
α	0.95	0.89	0.78	**0.27**	0.46	**0.27**	65.3%
$\varepsilon_{\text{HOMO}}$	0.00388	0.00541	**0.00321**	0.00331	0.00328	0.00337	–
$\varepsilon_{\text{LUMO}}$	0.00512	0.00623	0.00355	**0.00350**	0.00354	0.00351	1.4%
$\Delta\varepsilon$	0.0112	0.0066	0.0049	0.0047	**0.0046**	0.0048	6.1%
$\langle R^2 \rangle$	**17.0**	28.5	34.1	21.5	25.8	22.9	37.0%
ZPVE	0.00172	0.00216	0.00124	**0.00018**	0.00064	0.00019	85.5%
U_0	2.43	2.05	2.32	**0.0357**	0.6855	0.0427	98.5%
U	2.43	2.00	2.08	**0.107**	0.686	0.111	94.9%
H	2.43	2.02	2.23	0.070	0.794	**0.0419**	98.1%
G	2.43	2.02	1.94	0.140	0.587	**0.0469**	97.6%
C_{v}	0.27	0.42	0.27	0.0989	0.158	**0.0944**	65.0%

133 385 small molecules. The aim here is to perform regression on twelve targets representing energetic, electronic, geometric, and thermodynamic properties, which were computed using density functional theory.

Baselines

We use the following kernel and GNN methods as baselines for our experiments.

Kernel Baselines. We use the Graphlet kernel (Shervashidze et al. 2009), the shortest-path kernel (Borgwardt and Kriegel 2005), the Weisfeiler-Lehman subtree kernel (WL) (Shervashidze et al. 2011), the Weisfeiler-Lehman Optimal Assignment kernel (WL-OA) (Kriege, Giscard, and Wilson 2016), and the global-local k-WL (Morris, Kersting, and Mutzel 2017) with k in $\{2, 3\}$ as kernel baselines. For each kernel, we computed the normalized Gram matrix. We used the C-SVM implementation of LIBSVM (Chang and Lin 2011) to compute the classification accuracies using 10-fold cross validation. The parameter C was selected from $\{10^{-3}, 10^{-2}, \ldots, 10^2, 10^3\}$ by 10-fold cross validation on the training folds.

Neural Baselines. To compare GNNs to kernels we used the basic 1-GNN layer of Equation (5), DCNN (Wang et al. 2018), PatchySan (Niepert, Ahmed, and Kutzkov 2016), DGCNN (Zhang et al. 2018). For the QM9 dataset we used a 1-GNN layer similar to (Gilmer et al. 2017), where we replaced the inner sum of Equation (5) with a 2-layer MLP in order incorporate edge features (bond type and distance information). Moreover, we compare against the numbers provided in (Wu et al. 2018).

Model Configuration

We always used three layers for 1-GNN, and two layers for (local) 2-GNN and 3-GNN, all with a hidden-dimension size of 64. For the hierarchical variant we used architectures that use features computed by 1-GNN as initial features for the 2-GNN (1-2-GNN) and 3-GNN (1-3-GNN), respectively. Moreover, using the combination of the former we componentwise concatenated the computed features of the 1-2-GNN and the 1-3-GNN (1-2-3-GNN). For the final classification and regression steps, we used a three layer MLP, with binary cross entropy and mean

squared error for the optimization, respectively. For classification we used a dropout layer with $p = 0.5$ after the first layer of the MLP. We applied global average pooling to generate a vector representation of the graph from the computed node features for each k. The resulting vectors are concatenated column-wise before feeding them into the MLP. Moreover, we used the Adam optimizer with an initial learning rate of 10^{-2} and applied an adaptive learning rate decay based on validation results to a minimum of 10^{-5}. We trained the classification networks for 100 epochs and the regression networks for 200 epochs.

Experimental Protocol

For the smaller datasets, which we use for comparison against the kernel methods, we performed a 10-fold cross validation where we randomly sampled 10% of each training fold to act as a validation set. For the QM9 dataset, we follow the dataset splits described in (Wu et al. 2018). We randomly sampled 10% of the examples for validation, another 10% for testing, and used the remaining for training. We used the same initial node features as described in (Gilmer et al. 2017). Moreover, in order to illustrate the benefits of our hierarchical k-GNN architecture, we did not use a complete graph, where edges are annotated with pairwise distances, as input. Instead, we only used pairwise Euclidean distances for connected nodes, computed from the provided node coordinates. The code was built upon the work of (Fey et al. 2018) and is provided at `https://github.com/chrsmrrs/k-gnn`.

Results and Discussion

In the following we answer questions **Q1** to **Q3**. Table 1 shows the results for comparison with the kernel methods on the graph classification benchmark datasets. Here, the hierarchical k-GNN is on par with the kernels despite the small dataset sizes (answering question **Q1**). We also find that the 1-2-3-GNN significantly outperforms the 1-GNN on all seven datasets (answering **Q2**), with the 1-GNN being the overall weakest method across all tasks.[3] We can further see that optimizing the parameters of the aggregation function only leads to slight performance gains on two out of three datasets, and that no optimization even achieves better results on the PROTEINS benchmark dataset (answering **Q3**). We contribute this effect to the one-hot encoded node labels, which allow the GNN to gather enough information out of the neighborhood of a node, even when this aggregation is not learned.

Table 2 shows the results for the QM9 dataset. On eleven out of twelve targets all of our hierarchical variants beat the 1-GNN baseline, providing further evidence for **Q2**. For example, on the target H we achieve a large improvement of 98.1% in MAE compared to the baseline. Moreover, on ten out of twelve datasets, the hierarchical k-GNNs beat the baselines from (Wu et al. 2018). However, the additional structural information

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[3]Note that in very recent work, GNNs have shown superior results over kernels when using advanced pooling techniques (Ying et al. 2018b). Note that our layers can be combined with these pooling layers. However, we opted to use standard global pooling in order to compare a typical GNN implementation with standard off-the-shelf kernels.

extracted by the k-GNN layers does not serve all tasks equally, leading to huge differences in gains across the targets.

It should be noted that our k-GNN models have more parameters than the 1-GNN model, since we stack two additional GNN layers for each k. However, extending the 1-GNN model by additional layers to match the number of parameters of the k-GNN did not lead to better results in any experiment.

Conclusion

We presented a theoretical investigation of GNNs, showing that a wide class of GNN architectures cannot be stronger than the 1-WL. On the positive side, we showed that, in principle, GNNs possess the same power in terms of distinguishing between non-isomorphic (sub-)graphs, while having the added benefit of adapting to the given data distribution. Based on this insight, we proposed k-GNNs which are a generalization of GNNs based on the k-WL. This new model is strictly stronger then GNNs in terms of distinguishing non-isomorphic (sub-)graphs and is capable of distinguishing more graph properties. Moreover, we devised a hierarchical variant of k-GNNs, which can exploit the hierarchical organization of most real-world graphs. Our experimental study shows that k-GNNs consistently outperform 1-GNNs and beat state-of-the-art neural architectures on large-scale molecule learning tasks. Future work includes designing task-specific k-GNNs, e.g., devising k-GNNs layers that exploit expert-knowledge in bio- and chemoinformatic settings.

Acknowledgments

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Appendix

In the following we provide proofs for Theorem 1, Theorem 2, Proposition 3, and Proposition 4.

Proof of Theorem 1

Theorem 5 (Theorem 1 in the main paper). Let (G, l) be a labeled graph. Then for all $t \geq 0$ and for all choices of initial colorings $f^{(0)}$ consistent with l, and weights $\mathbf{W}^{(t)}$,

$$c_l^{(t)} \sqsubseteq f^{(t)}.$$

For the theorem we consider a single iteration of the 1-WL algorithm and the GNN on a single graph.

Proof of Theorem 1. We show for an arbitrary iteration t and nodes $u, v \in V(G)$, that $c_l^{(t+1)}(u) = c_l^{(t+1)}(v)$ implies $f^{(t+1)}(u) = f^{(t+1)}(v)$. In iteration 0 we have $c_l^{(0)}(u) = c_l^{(0)}(v) \iff f^{(0)}(u) = f^{(0)}(v)$ as the initial node coloring $f^{(0)}$ is chosen consistent with l.

Let $u, v \in V(G)$ and $t \in \mathbb{N}$ such that $c_l^{(t+1)}(u) = c_l^{(t+1)}(v)$. Assume for the induction that $c_l^{(t)}(u) = c_l^{(t)}(v) \implies f^{(t)}(u) = f^{(t)}(v)$ holds. As $c_l^{(t+1)}(u) = c_l^{(t+1)}(v)$ we know from the refinement step of the 1-WL that the old colors $c_l^{(t)}(u) = c_l^{(t)}(v)$ of u and v as well as the multisets $\{\!\{c_l^{(t)}(w) \mid w \in N(u)\}\!\}$ and $\{\!\{c_l^{(t)}(w) \mid w \in N(v)\}\!\}$ of colors of the neighbors of u and v are identical.

Let $M_u = \{\!\{f^{(t)}(w) \mid w \in N(u)\}\!\}$ and $M_v = \{\!\{f^{(t)}(v) \mid w \in N(v)\}\!\}$ be the multisets of feature vectors of the neighbors of u and v respectively. By the induction hypothesis, we know that $M_u = M_v$ and $f^{(t)}(u) = f^{(t)}(v)$ such that independent of the choice of f_{merge} and f_{aggr} we get $f^{(t+1)}(u) = f^{(t+1)}(v)$. This holds as the input to both functions f_{merge} and f_{aggr} is identical. This proves $c_l^{(t+1)}(u) = c_l^{(t+1)}(v) \implies f^{(t+1)}(u) = f^{(t+1)}(v)$ and thereby the theorem. $\square$

Proof of Theorem 2

Theorem 6 (Theorem 2 in the main paper). Let (G, l) be a labeled graph. Then for all $t \geq 0$ there exists a sequence of weights $\mathbf{W}^{(t)}$, and a 1-GNN architecture such that

$$c_l^{(t)} \equiv f^{(t)}.$$

For the proof we start by giving the proof for graphs where all nodes have the same initial color and then extend it to colored graphs. In order to do that we use a slightly adapted but equivalent version of the 1-WL. Note that the extension to colored graphs is mostly technical, while the important idea is already contained in the first case.

Uncolored Graphs

Let Γ_G be the refinement operator for the 1-WL, mapping the old coloring $c_l^{(t-1)}$ to the updated one $c_l^{(t)}$:

$$c_l^{(t)}(v) = \left(\Gamma_G\big(c_l^{(t-1)}\big)\right)(v) = \text{HASH}\left(\big(c_l^{(t-1)}(v), \{\!\{c_l^{(t-1)}(u) \mid u \in N(v)\}\!\}\big)\right).$$

We first show that for uncolored graphs this is equivalent to the update rule $\tilde{\Gamma}_G$:

$$\tilde{c}_l^{(t)}(v) = \left(\tilde{\Gamma}_G\big(\tilde{c}_l^{(t-1)}\big)\right)(v) = \text{HASH}\left(\big(\{\!\{c_l^{(t-1)}(u) \mid u \in N(v)\}\!\}\big)\right).$$

We denote J as the all-1 matrix where the size will always be clear from the context.

Lemma 7. Let G be a graph, $v, w \in V(G)$, and $t \in \mathbb{N}$ such that $\tilde{c}^{(t)}(u) \neq \tilde{c}^{(t)}(v)$. Then $\tilde{c}^{(t')}(u) \neq \tilde{c}^{(t')}(v)$ for all $t' \geq t$.

Proof. Let $t \in \mathbb{N}$ be minimal such that there are v, w with

$$\tilde{c}^{(t)}(u) \neq \tilde{c}^{(t)}(v), \tag{7}$$

and

$$\tilde{c}^{(t+1)}(u) = \tilde{c}^{(t+1)}(v). \tag{8}$$

Then $t \geq 1$, because $\tilde{c}^{(0)} = J$ as there are no initial colors. Let $P_1, \ldots, P_p$ be the color classes of $\tilde{c}^{(t-1)}$. That is, for all $x, y \in V(G)$ we have $\tilde{c}^{(t-1)}(x) = \tilde{c}^{(t-1)}(y)$ if any only if there is an $i \in [1 : p]$ such that $x, y \in P_i$. Similarly, let $Q_1, \ldots, Q_q$ be the color classes of $\tilde{c}^{(t)}$. Observe that the partition $\{Q_1, \ldots, Q_q\}$ of $V(G)$ refines the partition $\{P_1, \ldots, P_p\}$. Indeed, if there were $i \neq i' \in [1{:}p]$, $k \in [1{:}q]$ such that $P_i \cap Q_k \neq \emptyset$ and $P_{i'} \cap Q_k \neq \emptyset$, then all $x \in P_i \cap Q_k, y \in P_{i'} \cap Q_k$ would satisfy $\tilde{c}^{(t-1)}(x) \neq \tilde{c}^{(t-1)}(y)$ and $\tilde{c}^{(t)}(x) = \tilde{c}^{(t)}(y)$, contradicting the minimality of q.

Choose $v, w \in V(G)$ satisfying (7) and (8). By (7), there is an $i \in [1{:}p]$ such that $|N_G(v) \cap P_i| \neq |N_G(w) \cap P_i|$. Let $j_1, \ldots, j_\ell \in [1{:}q]$ such that $P_i = Q_{j_1} \cup \ldots \cup Q_{j_\ell}$. By (8), for all $k \in [1 : \ell]$ we have $|N_G(v) \cap Q_{j_k}| = |N_G(w) \cap Q_{j_k}|$. As the Q_j are disjoint, this implies $|N_G(v) \cap P_i| = |N_G(w) \cap P_i|$, which is a contradiction. $\square$

Hence, the two update rules are equivalent.

Corollary 8. For all G and all $t \in \mathbb{N}$ we have $\tilde{c}^{(t)} \equiv c^{(t)}$.

Thus we can use the update rule $\tilde{\Gamma}$ for the proof on unlabeled graphs. For the proof, it will be convenient to assume that $V(G) = [1{:}n]$ (although we still work with the notation $V(G)$). It follows that $n = |V|$. A node coloring $f^{(t)}$ defines a matrix $F^{(t)} \in \mathbb{R}^{n \times d}$ where the ith row of $F^{(t)}$ is defined by $f^{(t)}(i) \in \mathbb{R}^{1 \times d}$. Here we interpret i as a node from $V(G)$. As colorings and matrices can be interpreted as one another, given a matrix $F \in \mathbb{R}^{V(G) \times d}$ we write $\Gamma_G(F)$ (or $\tilde{\Gamma}_G(F)$) for a Weisfeiler-Leman iteration on the coloring c_F induced by the matrix F. For the GNN computation we provide a matrix based notation. Using the adjacency matrix $A \in \mathbb{R}^{n \times n}$ of G and a coloring $F \in \mathbb{R}^{n \times d}$, we can write the update rule of the GNN layer as

$$F^{t+1} = \Lambda_{A,W,b}(F^{(t)}) = \sigma(AF^{(t)}W^{(t)} + bJ),$$

where $\Lambda_{A,W,b}$ is the refinement operator of GNNs corresponding to a single iteration of the 1-WL. For simplicity of the proof, we choose

$$\sigma = \text{sgn} = \begin{cases} 1 & \text{if } x > 0, \\ -1 & \text{otherwise,} \end{cases}$$

and the bias as

$$b = -1.$$

Note that we later provide a way to simulate the sign-function using ReLu operations to indicate that choosing the sign function is not really a hard restriction.

Lemma 9. Let $B \in \mathbb{Z}^{s \times t}$ be a matrix such that $0 \leq B_{ij} \leq n - 1$ for all i, j and the rows of B are pairwise distinct. Then there is a matrix $X \in \mathbb{R}^{t \times s}$ such that the matrix $\text{sgn}(BX - J) \in \{-1, 1\}^{s \times s}$ is non-singular.

Proof. Let $z = (1, n, n^2, \ldots, n^{t-1})^T \in \mathbb{R}^t$ where n is the upper bound on the matrix entries of B and $b = Bz \in \mathbb{R}^s$. Then the entries of b are nonnegative and pairwise distinct. Without loss of generality, we assume that $\mathbf{b} = (b_1, \ldots, b_s)^T$ such that $b_1 > b_2 > \cdots > b_s \geq 0$. Now we choose numbers $x_1, \ldots, x_s \in \mathbb{R}$ such that

$$\begin{cases} b_i \cdot x_j < 1 & \text{if } i \geq j, \\ b_i \cdot x_j > 1 & \text{if } i < j \end{cases} \tag{9}$$

for all $i, j \in [s]$ as the b_i are ordered. Let $\mathbf{x} = (x_1, \ldots, x_s) \in \mathbb{R}^{1 \times s}$ and $C = b \cdot x \in \mathbb{R}^{s \times s}$ and $\hat{C} = \text{sgn}(C - J)$.

Then C has entries $C_{ij} = b_i \cdot x_j$, and thus by (9),

$$\hat{C} = \begin{pmatrix} -1 & 1 & 1 & 1 & \cdots & 1 & 1 \\ -1 & -1 & 1 & & \cdots & & 1 \\ \vdots & & & \ddots & \ddots & & \vdots \\ -1 & & \cdots & & -1 & 1 \\ -1 & & \cdots & & -1 & -1 \end{pmatrix} \tag{10}$$

Thus $\hat{C}$ is non-singular. Now we simply let $X = z \cdot x$. Then $BX = C$. $\qquad \square$

Let us call a matrix *row-independent modulo equality* if the set of all rows appearing in the matrix is linearly independent.

Example 10. The matrix

$$\begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$

is row-independent modulo equality.

Note that the all-1 matrix J is row-independent modulo equality in all dimensions.

Lemma 11. Let $d \in \mathbb{N}$, and let $F \in \mathbb{R}^{n \times d}$ be row independent modulo equality. Then there is a $W \in \mathbb{R}^{d \times n}$ such that the matrix $\Lambda_{A,W,-1}(F)$ is row independent modulo equality and

$$\Lambda_{A,W,-1}(F) \equiv \tilde{\Gamma}_G(F).$$

Proof. Let $Q_1, \ldots, Q_r$ be the color classes of F (that is, for all $v, v' \in V(G)$ it holds that $F_v = F_{v'} \iff \exists j \in [r] : v, v' \in Q_j$). Let $\tilde{F} \in \mathbb{R}^{r \times d}$ be the matrix with rows $\tilde{F}_j = F_v$ for all $j \in [r], v \in Q_j$. Then the rows of $\tilde{F}$ are linearly independent, and thus there is a matrix $M \in \mathbb{R}^{d \times r}$ such that $\tilde{F}M$ is the $(r \times r)$ identity matrix. It follows that $FM \in \mathbb{R}^{n \times r}$ is the matrix with entries

$$(FM)_{vj} = \begin{cases} 1 & \text{if } v \in Q_j, \\ 0 & \text{otherwise.} \end{cases} \tag{11}$$

Let $D \in \mathbb{Z}^{n \times r}$ be the matrix with entries $D_{vj} := |N_G(v) \cap Q_j|$. Note that

$$AFM = D, \tag{12}$$

because for all $v \in V$ and $j \in [t]$ we have

$$(AFM)_{vj} = \sum_{v' \in V(G)} A_{vv'}(FM)_{v'j} = \sum_{v' \in Q_j} A_{vv'} = D_{vj},$$

where the second equality follows from Equation (11). By the definition of Γ_G as the 1-WL operator on uncolored graphs, we have

$$\Gamma_G(F) \equiv D \tag{13}$$

if we view D as a coloring of V.

Let $P_1, \ldots, P_s$ be the color classes of D, and let $\tilde{D} \in \mathbb{Z}^{s \times r}$ be the matrix with rows $\tilde{D}_i = D_v$ for all $i \in [s]$ and $v \in P_i$. Then $0 \le \tilde{D}_{ij} \le n - 1$ for all i, j, and the rows of $\tilde{D}$ are pairwise distinct. By Lemma 9, there

is a matrix $X \in \mathbb{R}^{r \times s}$ such that the matrix $\mathrm{sgn}(\tilde{D}X - J) \in \mathbb{R}^{s \times s}$ is non singular. This implies that the matrix $\mathrm{sgn}(AFMX - J) = \mathrm{sgn}(DX - J)$ is row-independent modulo equality. Moreover, $\mathrm{sgn}(AFMX - J) \equiv D \equiv \Gamma_G(F)$ by (13). We let $W \in \mathbb{R}^{p \times n}$ be the matrix of obtained from $MX \in \mathbb{R}^{p \times s}$ by adding $n - s$ all-0 columns. Then

$$\Lambda_{A,W,-1}(F) = \mathrm{sgn}(AFW - J)$$

is row-independent modulo equality and $\Lambda_{A,W,-1}(F) \equiv \mathrm{sgn}(AFMX - J) \equiv \tilde{\Gamma}_G(F)$. $\qquad \square$

Corollary 12. There is a sequence $\mathbf{W} = (W^{(t)})_{t \in \mathbb{N}}$ with $W^{(t)} \in \mathbb{R}^{n \times n}$ such that for all $r \in \mathbb{N}$,

$$\tilde{c}^{(t)} \equiv \Lambda^{(t)}_{A,\mathbf{W},-1} \, .$$

where $\tilde{c}^{(t)}$ is given by the t-fold application of $\tilde{\Gamma}_G$ on the initial uniform coloring J.

Remark 13. The construction in Lemma 11 always outputs a matrix with as many columns as there are color classes in the resulting coloring. Thus we can choose d to be n and pad the matrix using additional 0-columns.

Colored Graphs

We now extend the computation to colored graphs. In order to do that, we again use an equivalent but slightly different variant of the Weisfeiler-Leman update rule leading to colorings $c^{(t)}_{l,0}$ instead of the usual $c^{(t)}_l$. We then start by showing that both update rules are equivalent.

We define $\Gamma_{G,l}$ to be the refinement operator for the 1-WL, mapping a coloring $c^{(t-1)}_{l,0}$ to the updated one $c^{(t)}_{l,0}$ as follows:

$$c^{(t)}_{l,0}(v) = \left(\Gamma_{G,l}\left(c^{(t-1)}_{l,0} \right) \right)(v) = \mathrm{HASH}\left(\left(c^{(0)}_{l,0}(v), \{\!\{ c^{(t-1)}_{l,0}(u) \mid u \in N(v) \}\!\} \right) \right).$$

Note that for $\Gamma_{G,l}$ we use the initial color $c^{(0)}_{l,0}(u)$ of a node u whereas Γ_G used the color $c^{(t-1)}_l(u)$ from the previous round. The idea of using those old colors is to make sure that any two nodes which got a different color in iteration t, get different colors in iteration $t' > t$. This is formalized by the following lemma.

Lemma 14. Let (G, l) be a colored graph, $v, w \in V(G)$, and $q \in \mathbb{N}$ such that $c^{(t)}_{l,0}(v) \neq c^{(t)}_{l,0}(w)$. Then $c^{(t')}_{l,0}(v) \neq c^{(t')}_{l,0}(w)$ for all $t' \geq t$.

Proof. Let $t \in \mathbb{N}$ be minimal such that there are v, w with

$$c^{(t)}_{l,0}(v) \neq c^{(t)}_{l,0}(w) \tag{14}$$

and

$$c^{(t+1)}_{l,0}(v) = c^{(t+1)}_{l,0}(w). \tag{15}$$

Then $t \geq 1$, because by (), $c^{(1)}_{l,0}(v) = c^{(1)}_{l,0}(w)$ implies $c^{(0)}_{l,0}(v) = c^{(0)}_{l,0}(w)$. Let $P_1, \ldots, P_p$ be the color classes of $c^{(t-1)}_{l,0}$, and let $Q_1, \ldots, Q_q$ be the color classes of $c^{(t)}_{l,0}$. Observe that the partition $\{Q_1, \ldots, Q_q\}$ of $V(G)$ refines the partition $\{P_1, \ldots, P_p\}$. The argument is the same as in the proof of Lemma 7.

Choose $v, w \in V(G)$ satisfying (14) and (15). By (14), either $c_{l,0}(v) \neq c_{l,0}(w)$ or there is an $i \in [1 : p]$ such that $|N_G(v) \cap P_i| \neq |N_G(w) \cap P_i|$. By (), $c_{l,0}(v) \neq c_{l,0}(w)$ contradicts (15). Thus $|N_G(v) \cap P_i| \neq |N_G(w) \cap P_i|$ for some $i \in [1:p]$. Let $j_1, \ldots, j_\ell \in [1:q]$ such that $P_i = Q_{j_1} \cup \ldots \cup Q_{j_\ell}$. By (14), for all $k \in [1:\ell]$ we have $|N_G(v) \cap Q_{j_k}| = |N_G(w) \cap Q_{j_k}|$. As the Q_j are disjoint, this implies $|N_G(v) \cap P_i| = |N_G(w) \cap P_i|$, which is a contradiction. $\qquad \square$

Corollary 15. For all graphs G and initial vertex colorings l of G we have $c^{(t)}_{l,0} \equiv c^{(t)}_l$ for all $t \in \mathbb{N}$.

We now consider the slightly modified update rule for 1-GNNs which takes the initial colors l into account. Let the matrix $F_{l,0}^{(0)} \in \mathbb{R}^{n \times d}$ be an encoding (such as the one-hot encoding) of l such that $F_{l,0}^{(0)}$ is linearly independent modulo equality. Then

$$F_{l,0}^{(t+1)} = \Lambda_{A,0,\mathbf{W},\mathbf{b}}(F_{l,0}^{(t)}) = \left(F_{l,0}^{(0)}, \ \Lambda_{A,W,b}(\Lambda_{A,0,\mathbf{W},\mathbf{b}}(F_{l,0}^{(t)}))\right).$$

We now show that this version of GNNs, which can be implemented by the simple variant (1) of GNNs provided in the main paper, is equivalent to 1-WL on colored graphs, proving the main theorem.

Proof of Theorem 2 (sketch). We prove the theorem by induction over the iteration t. The initial colorings are chosen consistent with l providing $c_{l,0}^{(0)} \equiv F_{l,0}^{(0)}$. For the induction step we assume $c_{l,0}^{(t)} \equiv F_{l,0}^{(t)}$ for iteration t. We know by Lemma 11 that the inner part $\Lambda_{A,W,b}(\Lambda_{A,0,\mathbf{W},\mathbf{b}}(F_{l,0}^{(t)}))$ of the update rule results in color classes which are identical to the ones that $\tilde{\Gamma}_G$ would produce. This implies that restricted to each color class Q from $Q_1^{(t)}, \ldots, Q_q^{(t)}$ of iteration t, the new color classes $Q_1^{(t+1)} \cap Q, \ldots, Q_{q'}^{(t+1)} \cap Q$ restricted to Q match the coloring $c_{l,0}^{(t+1)}|_Q$ that is $c_{l,0}^{(t+1)}$ restricted to Q. This holds as within one color class, the common color of the nodes contains no further information as shown in Lemma 7. Observe that colors are represented by linearly independent row vectors. This especially holds for $F^{(t+1)} = \Lambda_{A,W,b}(\Lambda_{A,0,\mathbf{W},\mathbf{b}}(F_{l,0}^{(t)}))$. In order to show that $F_{l,0}^{(t+1)}$ represents the coloring $c_{l,0}^{(t+1)}$ we have to prove two properties.

1. Given $Q_i^{(0)} \neq Q_j^{(0)}$ and $u \in Q_i^{(0)}, v \in Q_j^{(0)}$, we have $F_{l,0}^{(t+1)}(u) \neq F_{l,0}^{(t+1)}(v)$.

2. $F_{l,0}^{(t+1)}$ is linearly independent modulo equality.

In the first item we interpret $F_{l,0}^{(t+1)}$ as a coloring function. Both properties follow directly from the definition of $F_{l,0}^{(t+1)}$ as the concatenation of the matrices $F_{l,0}^{(0)}$ and $F^{(t+1)}$. The first property holds as the row vectors of $F_{l,0}^{(0)}$ are clearly different, as otherwise $Q_i = Q_j$. The second property follows from the fact that linear independence cannot be lost by extending a matrix. Thus within each old color class $Q_i^{(t)}$, all vectors are linearly independent modulo equality as $F^{(t+1)}$ (without subscript) is linearly independent modulo equality. This then extends to all combinations of old color classes and colors from $F^{(t+1)}$ as $F_{l,0}^{(0)}$ is also linearly independent modulo equality. By induction this chows that $c_{l,0}^{(t)} \equiv F_{l,0}^{(t)}$ for all t.

Note that the width of the matrices $F_{l,0}^{(t+1)}$ can be bounded by $2n$. The width of $F_{l,0}^{(0)}$ can trivially be bounded by n (in the worst case every node has a different initial color). The width of $F^{(t+1)}$ can also be bounded by n using Remark 13 as it is the output of a computation using Lemma 11. Using Corollary 15 for the step from $c_{l,0}^{(t+1)}$ to $c_l^{(t+1)}$ finishes the proof of 2. $\qquad\square$

Adaptation to the ReLu activation function.

The above framework can be easily adapted to use the ReLu function as the activation function of our choice, as follows. Recall that the sgn activation function was applied to the matrix $C - J$, where the matrix C is defined via Equation 9. We claim that a two-fold application of ReLu, interspersed by some elementary matrix operations, also yields the matrix $\text{sgn}(C - J)$. This ensures that we can use our GNN architecture to work with the ReLu activation function.

The proof is as follows. The first application of ReLu to the matrix $-(C - J)$ yields the matrix $C^{(1)} := \sigma_{\text{ReLu}}(J - C)$, which is a lower triangular matrix satisfying

$$\begin{cases} C_{ij}^{(1)} = 0 & \text{for } i < j \text{ and} \\ C_{ij}^{(1)} > 0 & \text{for } i \geq j \end{cases}$$

Let δ be the smallest positive value occurring in $C^{(1)}$. Note that δ is well defined as there are only a finite number of possible binary input matrices (with values $\{-1, 1\}$) for which $C^{(1)}$ can deterministically be computed. The matrix

$C^{(2)} := -\frac{2}{\delta} C^{(1)} + 2J$ then satisfies

$$\begin{cases} C_{ij}^{(2)} = 2 & \text{for } i < j \text{ and} \\ C_{ij}^{(2)} < 0 & \text{for } i \geq j. \end{cases}$$

Another application of ReLu yields an upper triangular matrix $C^{(3)} := \sigma_{\text{ReLu}}(C^{(2)})$, where every non-zero entry is equal to 2. Subtracting J from $C^{(3)}$ then yields $\text{sgn}\,(C)$. We now show that those operations can be represented by the GNN architecture. Recall that our GNN architecture was given by stacking a number of iterations of the form, in matrix notation

$$F^{(t)} = \sigma\left(F^{(t-1)} W_1^{(t)} + A F_l^{(t-1)} W_2^{(t)} + b W_3^{(t)} \right) \tag{16}$$

where the $W_3^{(t-1)}$ term is the bias that is added in each iteration and σ is a non-linear function such as $\tanh$ or ReLu written as σ_{ReLu}.

To compute $\sigma_{\text{ReLu}}(J - C)$ we can choose $W_1^{(1)}$ and $W_2^{(1)}$ appropriate to reach $-C$, $W_3^{(1)} = J$ as the all-1 matrix and $b = 1$. Correspondingly, for $\sigma_{\text{ReLu}}(-\frac{2}{\delta} C^{(1)} + 2J)$ we can choose $W_1^{(2)} = -\frac{2}{\delta} I$, $b = 2$, $W_3^{(2)} = J$ and $W_2^{(2)} = \mathbf{0}$. In the next round $W_3^{(3)}$ can be chosen according to the regular choices of $W_1^{(3)}$ and $W_2^{(3)}$ to simulate a computation on $\text{sgn}\,(C)$ instead of $\text{sgn}\,(C + J)$ which works by the linearity of the inner computation. Every layer described by Equation (16) is followed by an adjustment layer plus a few modifications in both the previous and upcoming step. As a result, we have a constant (two-fold) increase in the number of layers in our GNN set-up. The following corollary summarizes the above discussion.

Corollary 16 (ReLu activation). *Let (G, l) be a labeled graph. Then there exists a sequence $(\mathbf{W}^t)_{t>0}$ and a 1-GNN architecture based on the simple architecture described in (5) in the main paper using ReLu as activation function, such that $c_l^t = f^{(2t)}$ for all $t \geq 0$.*

Proofs of Proposition 3 and Proposition 4

Proposition 17 (Proposition 3 in the main paper). *Let (G, l) be a labeled graph and let $k \geq 2$. Then for all $t \geq 0$, for all choices of initial colorings $f_k^{(0)}$ consistent with l and for all weights $\mathbf{W}^{(t)}$,*

$$c_{s,k,l}^{(t)} \sqsubseteq f_k^{(t)} \,.$$

Proof. The proof follows the arguments in the proof of Theorem 1. We therefore only provide a brief proof by induction on the iteration t. For the base case, i.e. iteration $t = 0$, the statement holds because the initial coloring $f_k^{(0)}$ is chosen to be consistent with the isomorphism types $d_{k,l}^0$.

For the inductive step, assume that the statement holds until iteration $t - 1$. Consider two tuples u and v which (i) are not distinguished in the first $(t - 1)$ iterations and (ii) are not distinguished in the tth iteration. Therefore, u and v must have equal number of neighbors from every color class. This implies that the k-GNN update rule yields the same output for u and v. Hence, if two such tuples are distinguished by the k-GNN in the tth iteration, they must be distinguished by the k-WL as well. This finishes the induction and proves the proposition. $\square$

Proposition 18 (Proposition 4 in the main paper). *Let (G, l) be a labeled graph and let $k \geq 2$. Then for all $t \geq 0$ there exists a sequence of weights $\mathbf{W}^{(t)}$ and a k-GNN architecture such that*

$$c_{s,k,l}^{(t)} \equiv f_k^{(t)} \,.$$

Proof. Let us simulate the set-based k-WL on a n-vertex graph G via a 1-WL on a graph $G^{\otimes k}$ on $O(n^k)$ vertices, defined as follows. The vertex set of $G^{\otimes k}$ is the set $[V(G)]^k$ of all k-element subsets of $V(G)$. The edge set of $G^{\otimes k}$ is defined as follows: two sets s and t are connected by an edge in $G^{\otimes k}$ if and only if $|s \cap t| = k - 1$. Observe that the neighborhood of a vertex s in this graph is exactly the set $N(s)$ defined earlier. The initial labeling of the vertices of the graph $G^{\otimes k}$ is determined as follows: For $s \in V(G^{\otimes k})$ the initial label of s is its isomorphism type.

For the above construction, it immediately follows that performing the 1-WL on the graph $G^{\otimes k}$ yields the same coloring, as the one obtained by performing k-WL for the graph G. It remains to define the sequence $(\boldsymbol{W}^t)_{t>0}$ such that k-GNN simulates the set based k-WL on G. Applying Theorem 2 to the graph $G^{\otimes k}$ results in a sequence $\widetilde{\boldsymbol{W}}^t$ such that the 1-GNN can simulate 1-WL on $G^{\otimes k}$ using $\widetilde{\boldsymbol{W}}^t$. Hence, this sequence can be directly used in the k-GNN to simulate k-WL on G. $\qquad\square$