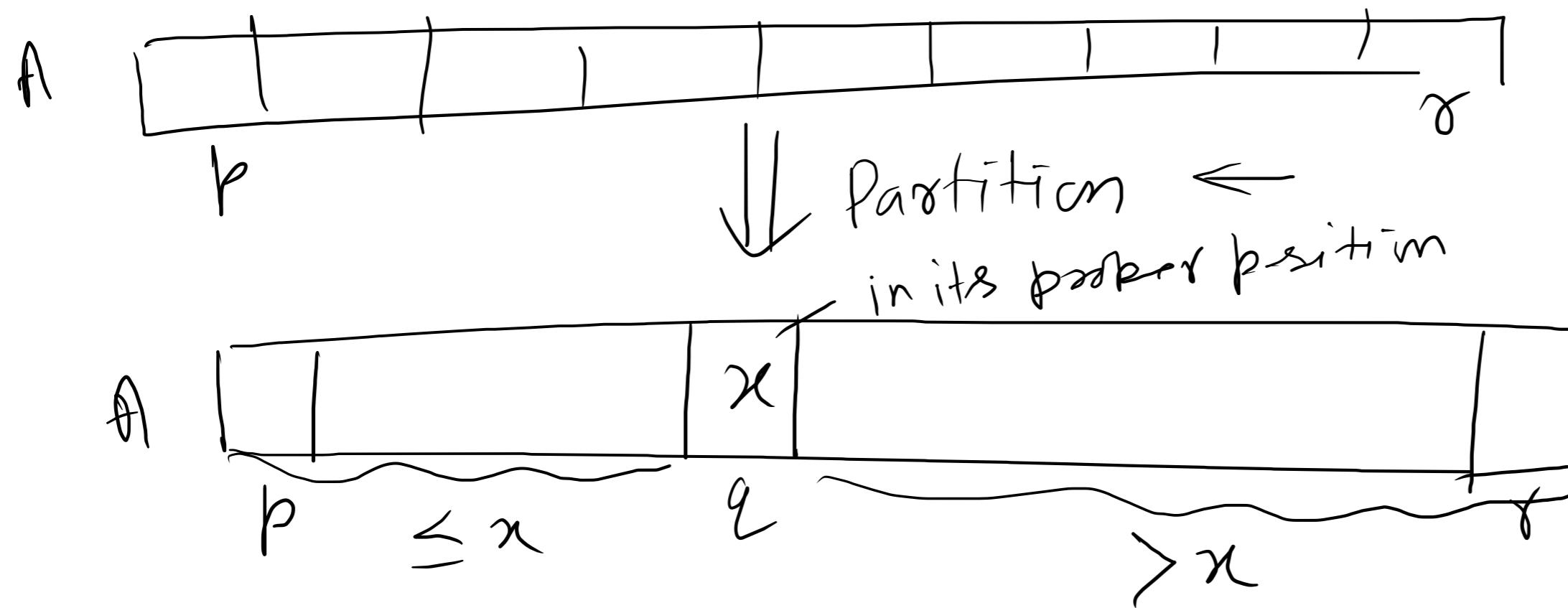


Quick sort

Helper function Partition.



Partition (A, p, r)

$$x = A[r]$$

$$i = p - 1$$

$$\text{for } j = p \text{ to } r - 1$$

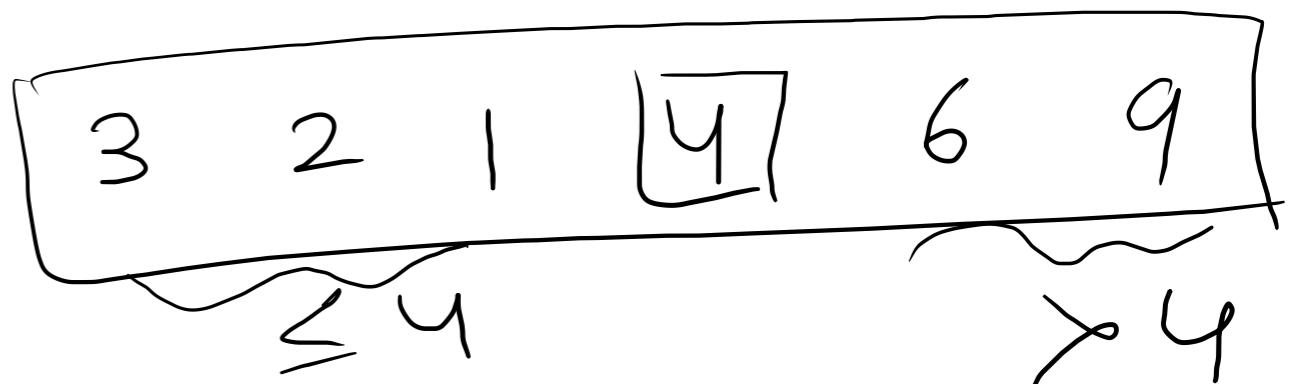
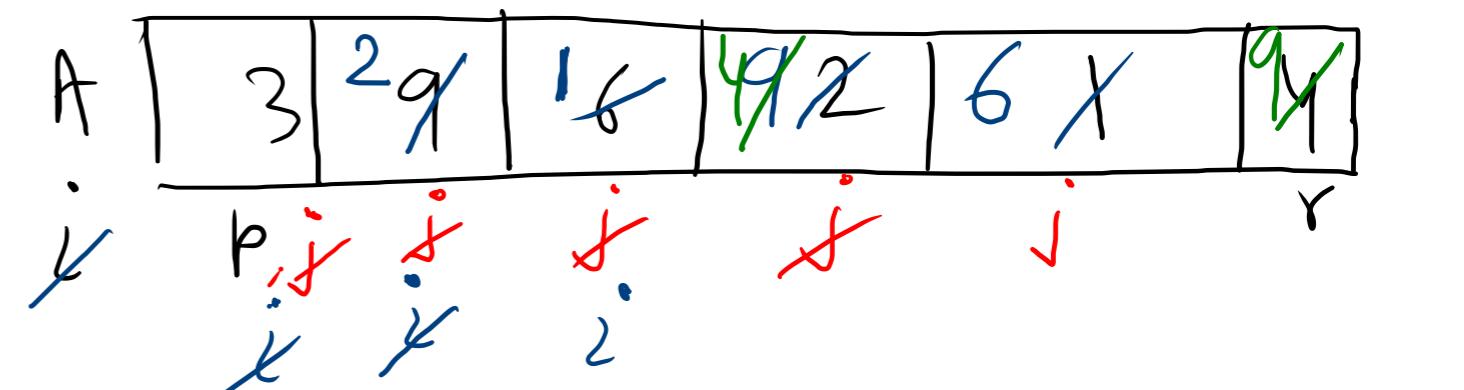
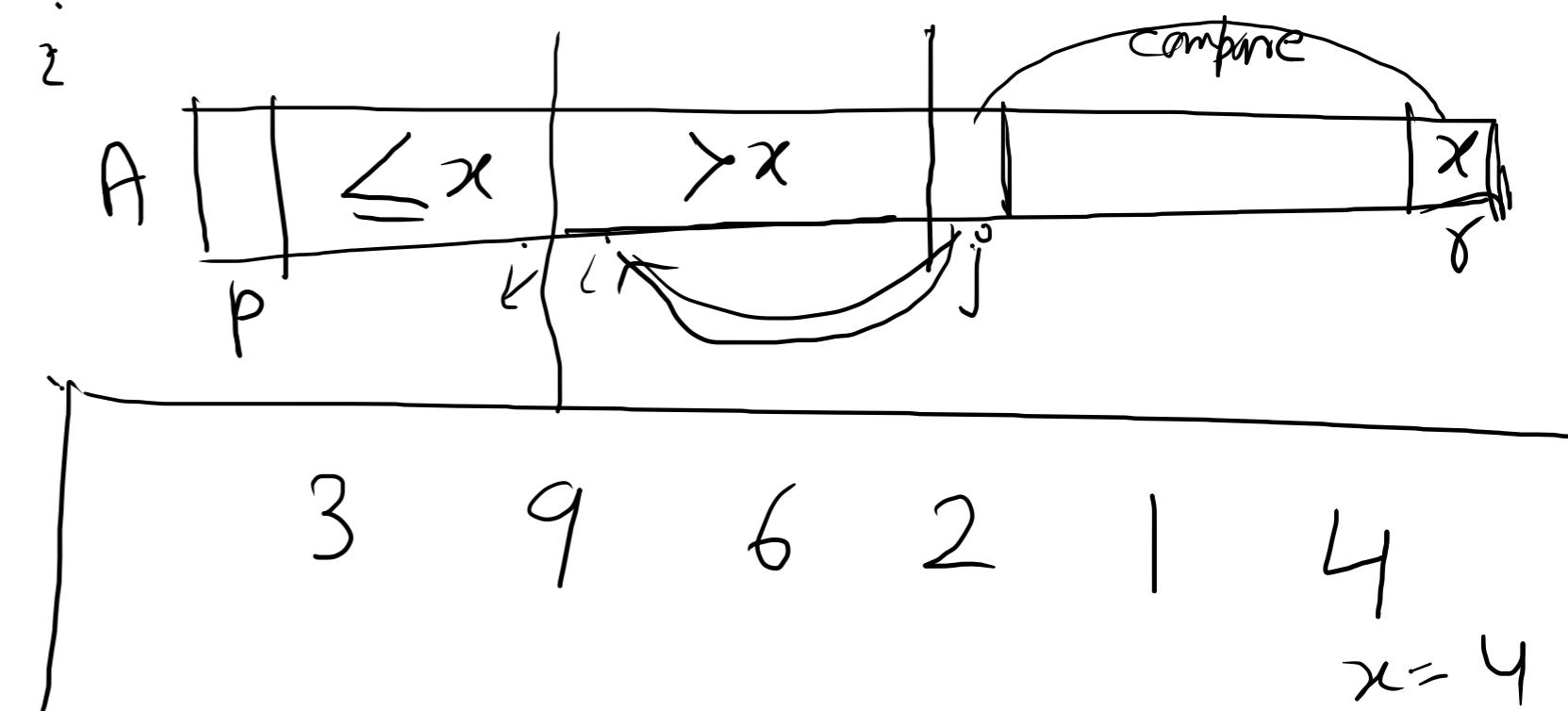
if $A[j] \leq x$
 $i = i + 1$
swap $A[i]$ and $A[j]$

Swap $A[i+1]$ and $A[r]$

return $i + 1$

Running time

$O(n)$



Quicksort (A, p, r) — T(n)

if $p < r$

$q = \text{Partition}(A, p, r) \leftarrow$

Quicksort (A, p, q-1)

Quicksort (A, q+1, r)

Running time $x \leftarrow$ is called the pivot element.

efficient when the pivot element partitions the array into roughly two equal arrays.

$$\frac{1}{100} : \frac{99}{100}$$

$$\begin{aligned} T(n) &= T\left(\frac{n}{100}\right) + T\left(\frac{99n}{100}\right) + O(n) \\ &= O(n \log n) \end{aligned}$$

H.W
How?

non-efficient

constant ; n-constant .

Recurrence

$$\begin{aligned} T(n) &= +(\text{constant}) + T(\text{n-constant}) + O(n) \\ &= O(n^2) \end{aligned}$$

H.W

How ?

Ex^m

worst case happens

- All elements are same (roughly same)
- Sorted Scanner .

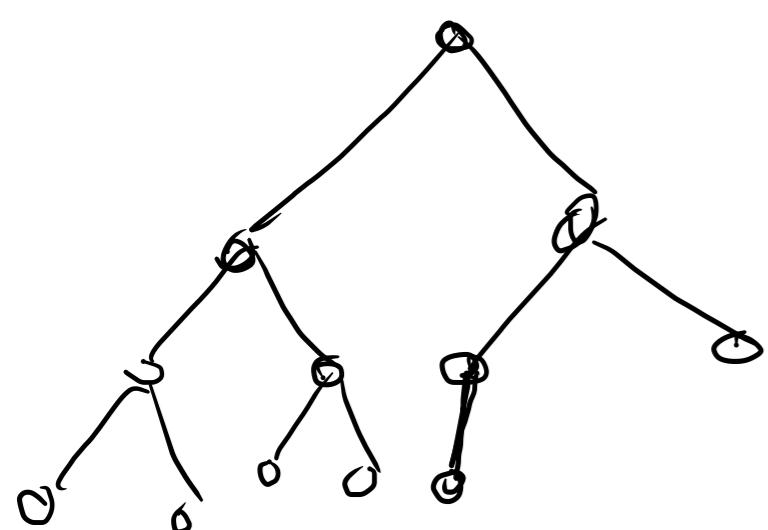
Randomized version

$O(n \lg n)$ - expected running time.

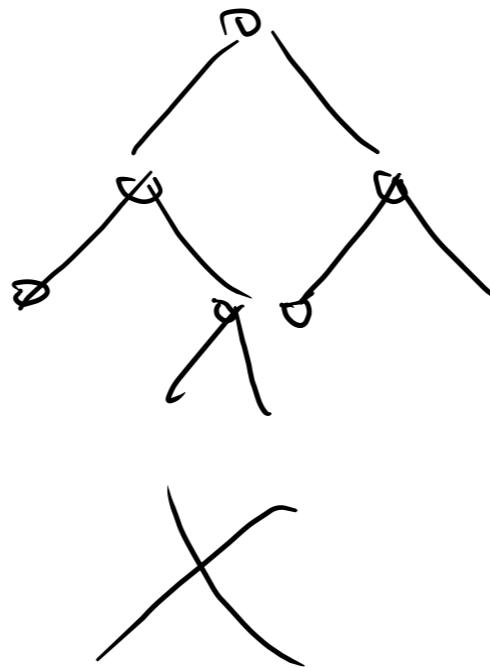
Heap sort

Heap data structure

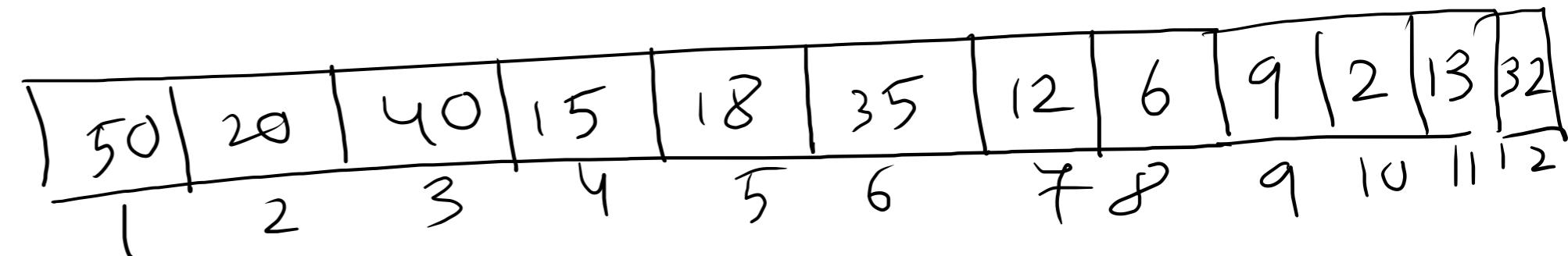
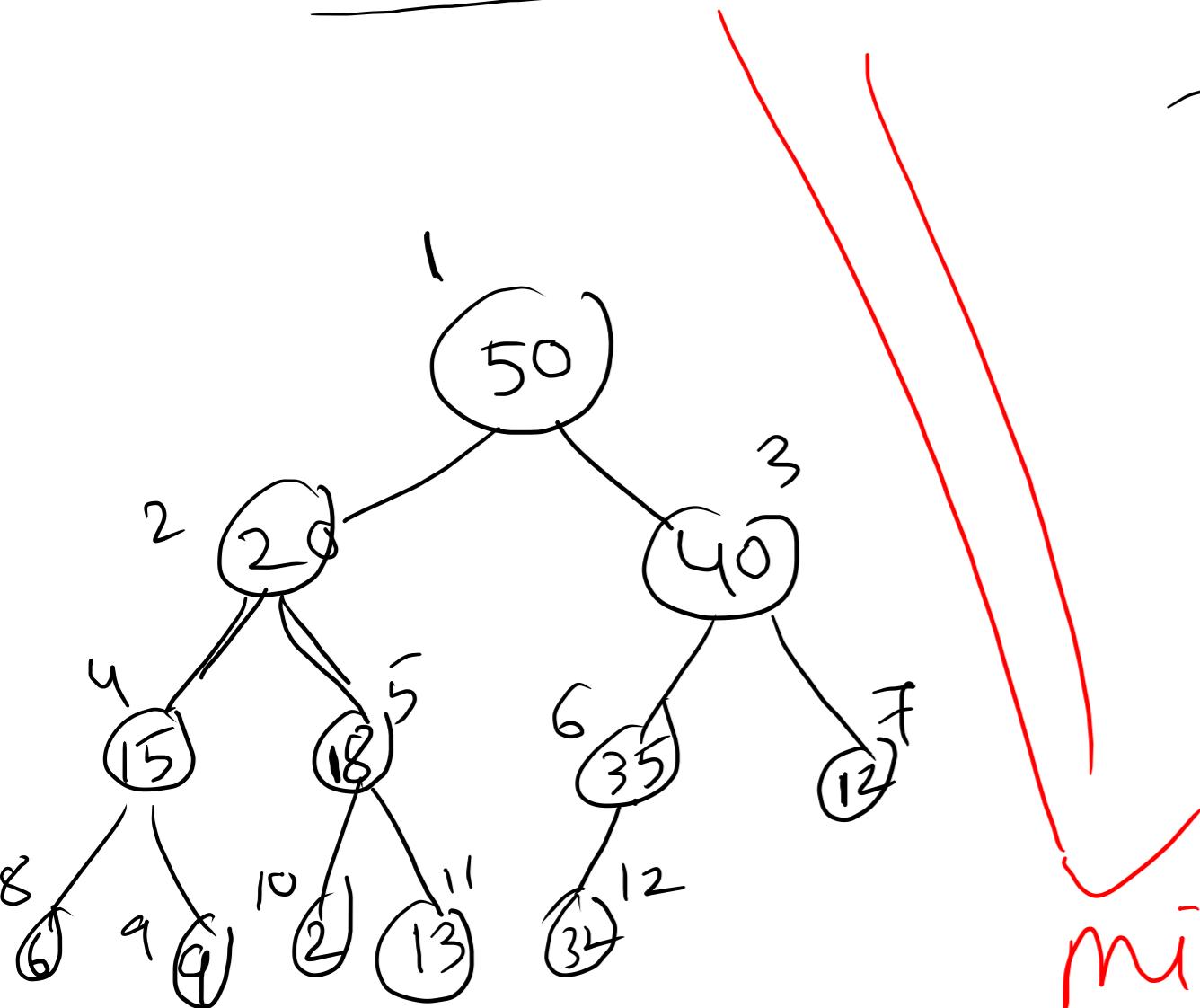
- It is an array
- can be viewed as a nearly complete binary tree where the lowest level leaves are filled from left to right.



✓



Max-heap property: The key of a node is $>$ the key of its children.



Min-heap property

The key of a node is $<$ the key of its children.