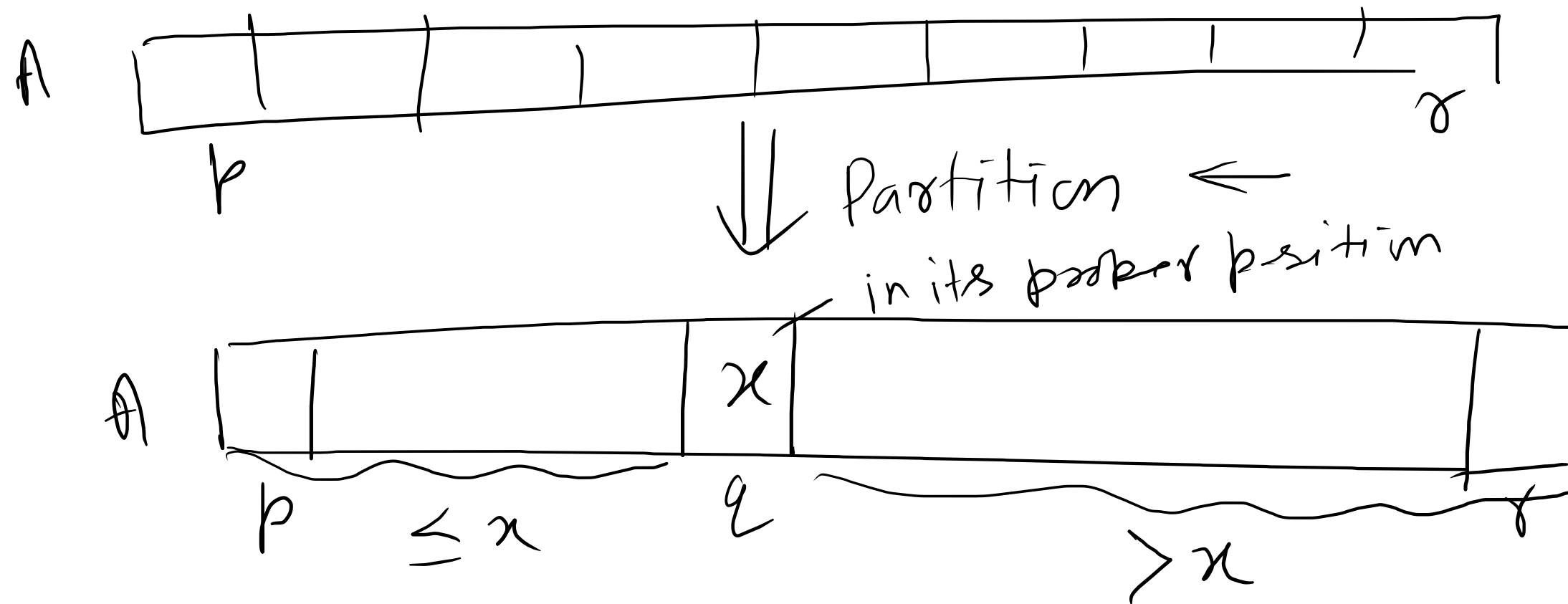


# Quick sort

Helper function Partition.



Partition (A, p, r)

$x = A[r]$

$i = p - 1$

for  $j = p$  to  $r - 1$

if  $A[j] \leq x$

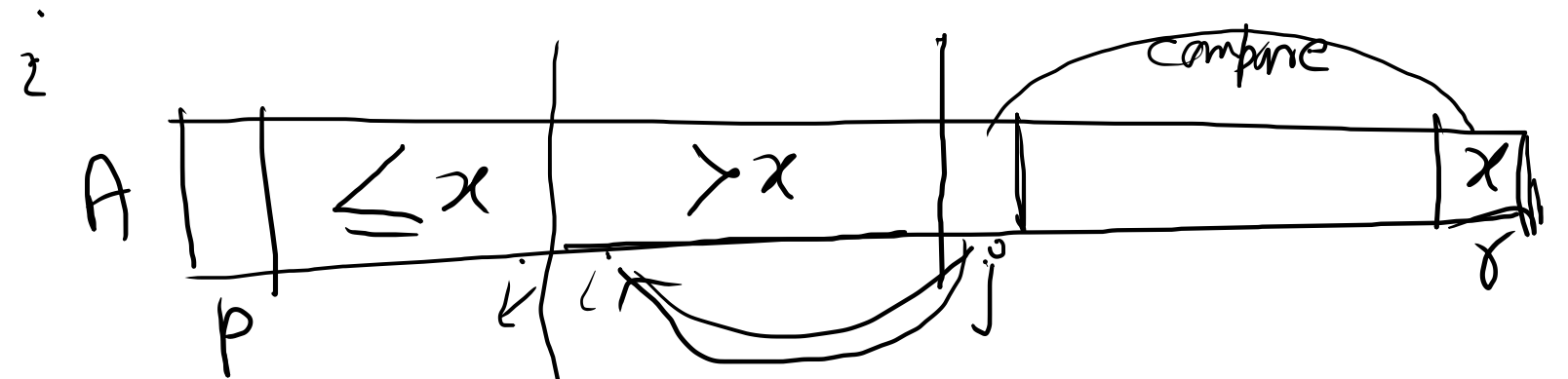
$i = i + 1$

swap  $A[i]$  and  $A[j]$

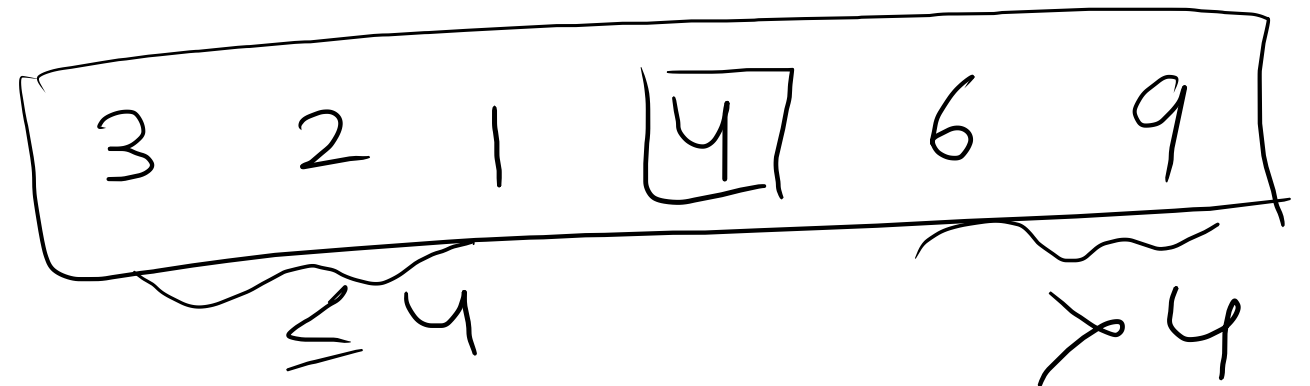
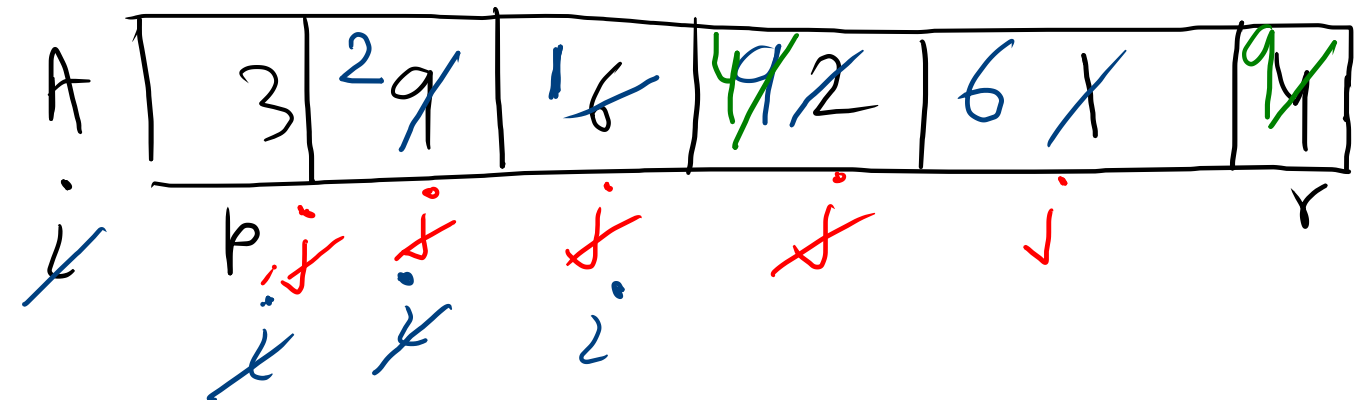
Swap  $A[i + 1]$  and  $A[r]$

return  $i + 1$

Running time  
 $O(n)$



3 9 6 2 1 4  
 $x = 4$



Quicksort (A, p, r) —  $T(n)$

if  $p < r$

┌  $q = \text{Partition}(A, p, r) \leftarrow$   
├ Quicksort (A, p,  $q-1$ )  
└ Quicksort (A,  $q+1$ , r)

Running time

$x \leftarrow$  is called the pivot element.

efficient when the pivot element partitions the array into roughly two equal arrays.

$$\frac{1}{100} : \frac{99}{100}$$

$$\begin{aligned} T(n) &= T\left(\frac{n}{100}\right) + T\left(\frac{99n}{100}\right) + O(n) \\ &= O(n \log n) \end{aligned}$$

H.W  
How?

non efficient

constant :  $n$ -constant.

Recurrence

$$T(n) = T(\text{constant}) + T(n - \text{constant}) + O(n)$$

$$= O(n^2)$$

H.W

How?

Ex<sup>m</sup>

worst case happens

- All elements are same (roughly same)
- Sorted Sequence.

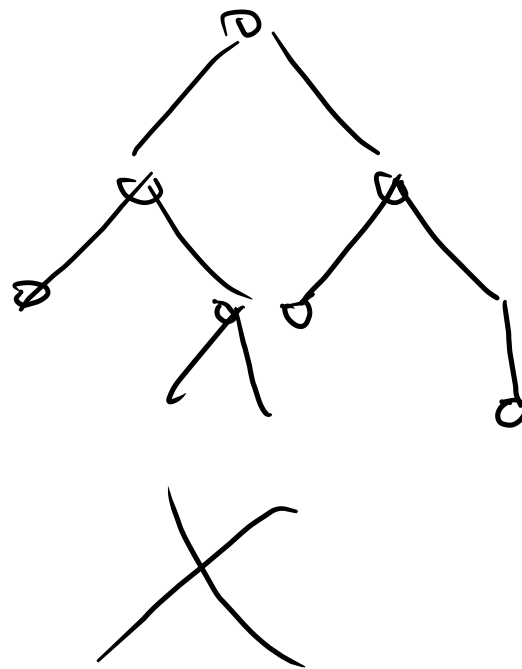
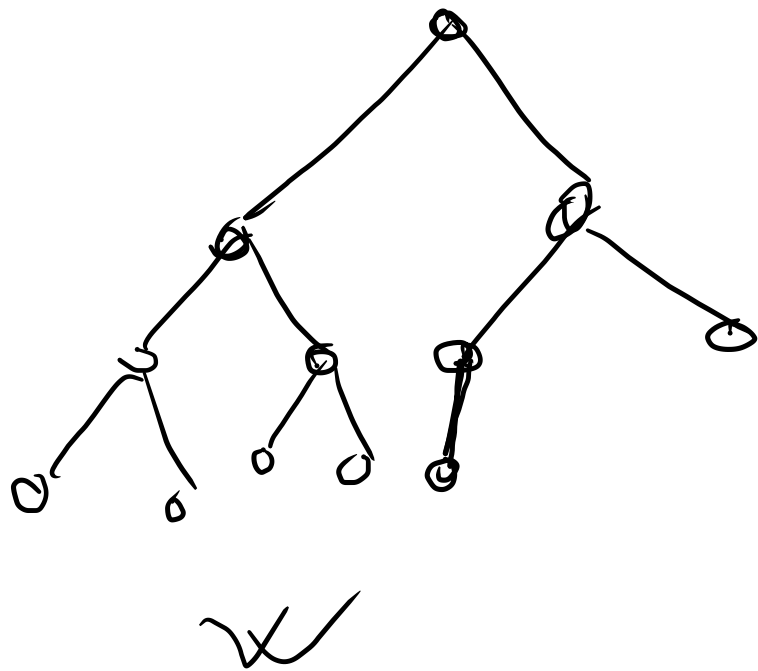
Randomized version

$O(n \log n)$  - expected running time.

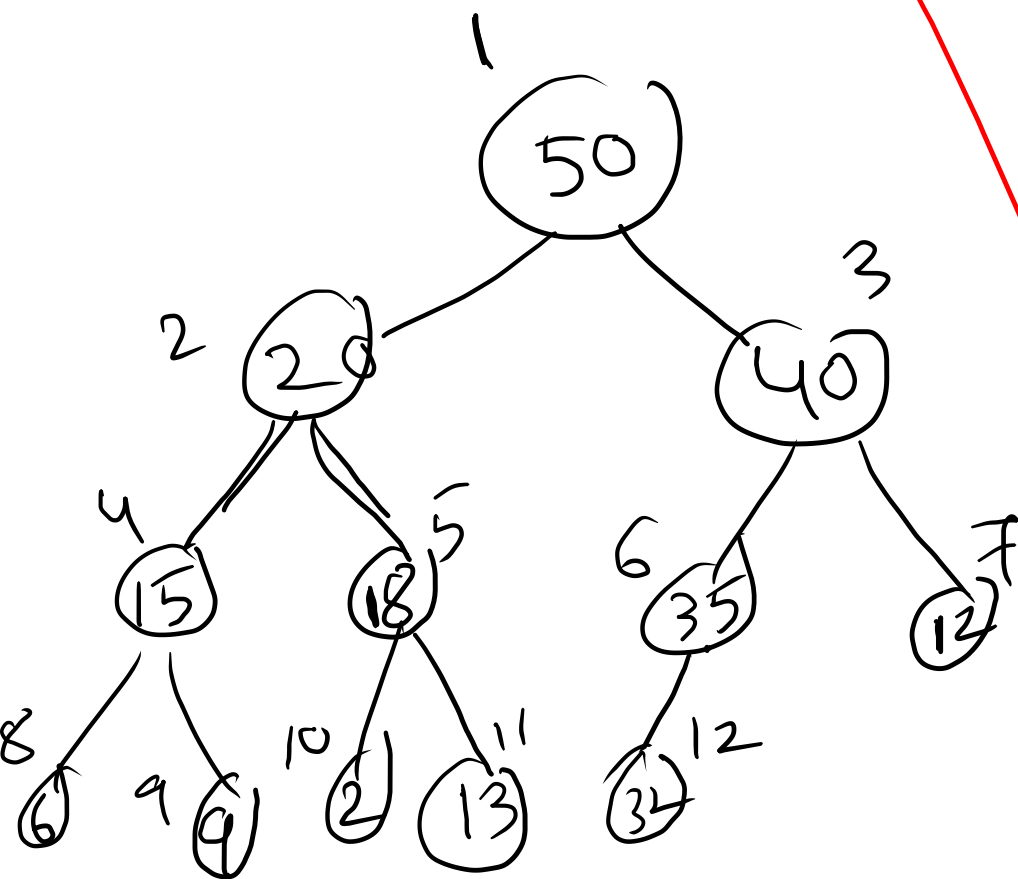
# Heap sort

## Heap data structure

- It is an array
- can be viewed as a nearly complete binary tree where the lowest level leaves are filled from left to right.



Max-heap property: The key of a node is  $>$   
the key of its children.



1	2	3	4	5	6	7	8	9	10	11	12
50	20	40	15	18	35	12	6	9	2	13	34

min-heap property

The key of a node is  $<$   
the key of its children.