# CMPL Coliop/Coin Mathematical Programming Language



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# **Manual**

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# 1 About CMPL

CMPL (Coliop / Coin Mathematical Programming Language) is a mathematical programming language and a system for modelling, solving and analysing linear programming (LP) problems and mixed integer programming (MIP) problems.

The CMPL syntax is similar in formulation to the original mathematical model but also includes syntactic elements from modern programming languages. CMPL is intended to combine the clarity of mathematical models with the flexibility of programming languages. CMPL contains the COIN-OR OSSolverService including the COIN-OR solvers. Since it is also possible to transform the mathematical problem into MPS, Free-MPS or OSiL files, alternative solvers can be used.

CMPL is an open source project licensed under GPL. It is written in C++ and is available for most of the relevant operating systems. CMPL is a COIN-OR project initiated by the Technical University of Applied Sciences Wildau and the Institute for Operations Research and Business Management at the Martin Luther University Halle-Wittenberg.

For further information please visit the CMPL website (www.coliop.org/cmpl).

# 2 Syntactic elements

# 2.1 General structure of a CMPL program

The structure of a CMPL program follows the standard model of linear programming (LP), which is defined by a linear objective function and linear constraints. Apart from the variable decision vector x all other components are constant.

```
c^{T} \cdot x \rightarrow max / min
s.t.
A \cdot x \leq b
x \geq 0
```

A CMPL program consists of four sections, the parameters section, the variables section, the objectives section and the constraints section, which can be inserted several times and mixed in a different order. Each sector can contain one or more lines with user-defined expressions.

```
parameters:
    # definition of the parameters
variables:
    # definition of the variables
objectives:
    # definition of the objective(s)
constraints:
    # definition of the constraints
```

A typical LP problem is optimal production planning. The aim is to find an optimal quantity for the products, depending on given capacities. The objective function is defined by the profit contribution per unit c and the

variable quantity of the products x. The constraints consist of the use of the capacities and the ranges for the decision variables. The use of the capacities is given by the product of the coefficient matrix A and the vector of the decision variables x and restricted by the vector of the available capacities b.

For example,

```
\begin{array}{l} 1 \cdot x_1 + 2 \cdot x_2 + 3 \cdot x_3 \rightarrow max \: ! \\ s.t. \\ 5.6 \cdot x_1 + 7.7 \cdot x_2 + 10.5 \cdot x_3 \leq 15 \\ 9.8 \cdot x_1 + 4.2 \cdot x_2 + 11.1 \cdot x_3 \leq 20 \\ 0 \leq x_n \  \  ; n = 1(1)3 \end{array}
```

#### can be formulated in CMPL as follows:

# 2.2 Keywords and other syntactic elements

**Keywords** 

parameters, variables, objectives,	section markers
constraints	
real, integer, binary	types of variable
real, integer, binary, string, set	types of parameter expression
	also used for type casts
max, min	objective senses
set, in, element, len, defset	key words for sets
max, min, dim, def, format, type	functions for parameter expressions
sqrt, exp, ln, lg, ld, srand, rand,	mathematical functions that can be used for para-
sin, cos, tan, acos, asin, atan, sinh,	meter expressions
cosh, tanh, abs, ceil, floor, round	

include	include of CMPL file
readcsv, readstdin	data import from a CSV file or from user input

error, echo	error and user message
sum	summation
continue, break, default, repeat	key words for control structures

**Arithmetic operators** 

+ -	signs for parameters or addition/subtraction
^	to the power of
* /	multiplication and division
div mod	integer division and remainder on division
:=	assignment operator

**Condition operators** 

= <= >=	conditions for constraints, while-loops and if-then
	clauses
== < > != <>	additional conditions in while-loops and if-then
	clauses
&&    !	logical operations (and, or, not)

**Other syntactic elements** 

- arithmetical bracketing in constant expressions - lists for initialising vectors of constants - parameters for constant functions - increment in an algorithmic set - indexing of vectors - range specification in variable definitions - control structures - algorithmic set (e.g. range for indices or loop counters) - range specification in variable definitions - element separation in an initialisation list for constant vectors and enumeration sets - separation of function parameters - separation of indices - separation of loop heads in a loop - separation of variables in a variable definition - mark indicating beginning of sections - definition of variables - definition of parameter type - separation of loop header from loop body - separation of line names - separation of alternative blocks in a control structure	other syntactic elements	
- parameters for constant functions - increment in an algorithmic set  - indexing of vectors - range specification in variable definitions  - control structures - algorithmic set (e.g. range for indices or loop counters) - range specification in variable definitions  - element separation in an initialisation list for constant vectors and enumeration sets - separation of function parameters - separation of indices - separation of loop heads in a loop - separation of variables in a variable definition  - mark indicating beginning of sections - definition of variables - definition of parameter type - separation of loop header from loop body - separation of line names  - separation of alternative blocks in a control struc-	( )	- arithmetical bracketing in constant expressions
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- separation of line names - separation of alternative blocks in a control struc-		- definition of parameter type
- separation of alternative blocks in a control struc-		- separation of loop header from loop body
		- separation of line names
ture	1	- separation of alternative blocks in a control struc-
		ture

;	- mark indicating end of a statement - every state-
	ment is to be closed by a semicolon

#		- comment (up to end of line)
/*	*/	- comment (between /* and */)

# 2.3 Objects

#### 2.3.1 **Parameters**

A parameters section consists of parameter definitions and assignments to parameters. A parameter can only defined within the parameters section using a assignment.

Note that a parameter can be used as a constant in a linear optimization model as coefficients in objectives and constraints. Otherwise parameters can be used like variables in programming languages. Parameters are usable in expressions, for instance in the calculation and definition of other parameters. A user can assign a value to a parameter and can then subsequently change the value with a new assignment.

A parameter is identified by name and, if necessary, by one or more indices. A type can be specified but is not necessary. Possible types are real, integer, binary, string and set. A parameter can be a scalar or an array of parameter values (e.g. vector, matrix or another multidimensional construct). A parameter is defined by an assignment with the assignment operator :=.

#### Usage:

```
name [: type] := scalarExpression;
name[index] [: type] := scalarExpression;
name[[set]] [: type] := non-scalarExpression;
```

name name of the parameter

type optional specification of the type of parameter

> Possible types are real, integer, binary, string and set. If the type is not defined, the type of the parameter is given by the expres-

sion on the right hand side.

index a position in an array of parameters

> The index can be an integer or a string expression. For multidimensional arrays it is necessary to set the index for every dimension separated by

commas.

scalarExpression A scalar parameter or a single part of an array of parameters is assigned a

single integer or real number, a single string, the scalar result of a math-

ematical function.

The elements of an array can also be sets. But it is not possible to mix set

and non-set expressions in one array.

set

an optional set expression for the definition of the dimension of the array

A set is a collection of distinct objects. Distinction can be made between enumeration sets, algorithmic sets and sets which are based on set operations like unions or intersections.

non-scalarExpression

A non-scalar expression consists of a list of <code>scalarExpressions</code>. The elements of the list are separated by commas and imbedded in brackets.

The elements of the list can also be sets. But it is not possible to mix set and non-set expressions.

#### Examples:

k := 10;	parameter k with value 10
k[15] := (0.5, 1, 2, 3.3, 5.5);	vector of parameters with 5 elements
k[] := (0.5, 1, 2, 3.3, 5.5);	
A[]: integer := (16, 45.4);	Definition of a matrix with two integer values a [1] =16
	and a[2]=45. Since A[] is defined as an integer mat-
	rix the real value 45.4 is transformed to the rounded
	value 45.
a[,] := ((5.6, 7.7, 10.5),	matrix with 2 rows and 3 columns
( 9.8, 4.2, 11.1 ));	
<pre>products := set( "bike1", "bike2");</pre>	defines a vector for machine hours based on the set
machineHours[products]:= (5.4, 10);	products
<pre>myString := "this is a string";</pre>	string parameter
q := 3;	parameter $q$ with value 3
g[1q] := (1, 2, 3);	usage of ${\bf q}$ for the definition of the parameter ${\bf g}$

If a name is used for a defined parameter, different usages of this name with indices can only refer to parameter, but not to model variables (e.g. if a[1] is a parameter, then a[2], a or a[1,1] can only be defined as parameter and <u>not</u> as model variables. a can also not be used as a local parameter like a loop counter).

A special kind of parameter are local parameters, which can only be defined within the head of a control structure. A local parameter is only valid in the body of the control structure and can be used like any other parameter. Only scalar parameters are permitted as local parameters. The main application of local parameters are loop counters iterated over a set.

# 2.3.2 Variables

The variables section is intended to declare the variables of a decision model, which are necessary for the definition of objectives and constraints in the decision model. A model variable is identified by name and, if necessary, by an index. A type must be specified. A model variable can be a scalar or a part of a vector, a matrix or another array of variables. A variable cannot be assigned a value.

#### Usage:

```
variables:
  name : type [[[lowerBound]..[upperBound]]];
  name[index] : type [[[lowerBound]..[upperBound]]];
  name[set] : type [[[lowerBound]..[upperBound]]];
```

name of model variable

type type of model variable.

Possible types are real, integer, binary.

[lowerBound..upperBound] optional parameter for limits of model variable

lowerBound and upperBound must be a real or integer expression. For the type binary it is not possible to specify bounds.

## Examples:

x: real;	$_{ imes}$ is a real model variable with no ranges
x: real[0100];	$_{\rm x}$ is a real model variable, $0 \le x \le 100$
x[15]: integer[1020];	vector with 5 elements, $10 \le x_n \le 20$ ; $n=1(1)5$
x[15,15,15]: real[0];	a three-dimensional array of real model variables with
	125 elements identified by indices,
	$x_{i,j,k} \ge 0$ ; $i, j, k = 1(1)5$
parameters:	
<pre>prod := set( "bike1", "bike2");</pre>	
variables:	defines a vector of non-negative real model variables
x[prod]: real[0];	based on the set prod
y: binary;	$\mathbf{x}$ is a binary model variable $\mathbf{y} \in \{0,1\}$

Different indices may cause model variables to have different types. (e.g. the following is permissible: variables: x[1]: real; x[2]: integer; )

If a name is used for a model variable definition, different usages of this name with indices can only refer to model variables and <u>not</u> to parameters (e.g. if x[1] is a model variable, then x[2], x or x[1,1] can only be defined as model variables).

## 2.3.3 Indices and sets

Sets are used for the definitions of arrays of parameters or model variables and for the iterations in loops. Indices are necessary to identify an element of an array like a vector or matrix of parameters or model variables.

A set is a collection of distinct integer and string elements. Sets can be defined by an enumeration of elements or by algorithms within the parameters section. It is also possible to build sets using set operations

like condition sets or unions or intersections of defined sets. A set can be stored in a scalar parameter or in an element of an array of parameters.

#### Usage set definitions:

startNumber(in/decrementor)endNumber	#algorithmic set
[startNumber][endNumber]	#algorithmic set
.integerstring.	<pre>#algorithmic set #algorithmic set</pre>
set(listOfIntAndStrings)	#enumeration set

startNumber(in/decrementor) endNumber set of integers based on an algorithm

The set starts at the startNumber, is changed by an incrementer or decrementer at every iteration and ends at the endNumber.

set of integers based on an algorithm

The set starts at the startNumber, is changed by the number one at every iteration and ends at the endNumber.

startNumber and endNumber are optional elements.

infinite set with all integers greater than or equal to startNumber

infinite set with all integers less than or equal to endNumber

infinite set with all integers and strings

infinite set with all integers infinite set with all strings

elements of an enumeration set

An enumeration set consists of one or more integer expressions or string expressions separated by commas and imbedded in brackets, and is described by the key word set.

It is possible to define an empty set using an empty array within the statement set().

#### Examples:

s:=;	s is assigned an infinite set of all integers and strings
s:=6;	s is assigned $s \in (, 4, 5, 6)$
s:=6;	s is assigned $s \in (6,7,8,)$

startNumber..endNumber

startNumber..

..endNumber

.integer.

.string.

listOfIntAndStrings

s:=06;	s is assigned $s \in (0,1,\ldots,6)$
s:=0(1)6;	
s:=10(-2)4;	s is assigned $s \in (10, 8, 6, 4)$
<pre>prod := set( "bike1", "bike2");</pre>	enumeration set of strings
a:= set(1, "a", 3, "b", 5, "c");	enumeration set of strings and integers
x[a]:=(10,20,30,40,50,60);	vector $\ensuremath{\mathbf{x}}$ identified by the set $\ensuremath{\mathbf{a}}$ is assigned an integer vector
	The following user messages are displayed:
echo x[1];	10
echo x["a"];	20
{i in a: echo x[i];}	10 20 30 40 50 60

#### Usage set operations and set construction:

```
set{ setIteration , condition: localParameter }; #condition set

set1 + set2; #union set

set1 * set2; #intersection set
```

# Usage set operations and set construction:

```
set1 + set2
set1 * set2
set3 * set1teration, condition: param } The local parameter param is to be used for the definition
of an iteration over a set (defined by setIteration)
and is to be evaluated in the condition condition. The
result is a set of all elements which are in the iterated set
and fulfil the condition.
```

#### Examples:

```
s1 := set( "a", "b", "c", "d");
s2 := set( "a", "e", "c", "f");
s3 := s1 + s2;
s4 := s1 * s2;
s5 := set{i in 1..10, i mod 2 = 0: i};
s5 := set{i in s1, !(i element s2): i}; s6 is assigned ( "b", "d" )
```

## 2.3.4 Line names

Line names are useful in huge models to provide a better overview of the model. In CMPL a line name can be defined by characters, numbers and the underscore character \_ followed by a colon. Names that are used for cmpl variables or model variables cannot be used for a line name. Within a control structure a line

name can include the current value of local parameters. This is especially useful for local parameters which are used as a loop counter. It is also possible to include the current matrix line number using a substitution expression imbedded in \$ \$.

#### Usage:

```
lineName:
lineName$1$:
lineName$2$:
lineName$k$:
loopName { controlStructure }
```

\$k\$ is replaced by the value of the local parameter k
\$1\$ is replaced by the number of the current line of the matrix.

\$1\$ is replaced by the number of the current line of the matrix.

In an implicit loop \$2\$ is replaced by the specific value of the free index.

\$100pName{controlStructure} defines line name subject to the following control structure. The values of loop counters in the control structure are appended automatically.

#### Examples:

```
parameters:
A[1..2,1..3] := ((1,2,3),(4,5,6));
b[1..2] := (100,100);
c[1..3] := (20,10,10);
variables:
x[1...3]: real[0...];
objectives:
profit: c[]T *x[] ->max;
                                         generates a line named profit
constraints:
restriction $1$: A[,] * x[] <=b[];
                                          generates 3 lines named restriction 2
                                                              restriction 3
                                                              restriction 4
restriction $2$: A[,] * x[] <=b[];
                                         generates 3 lines named restriction 1
                                                              restriction 2
                                                              restriction 3
```

{ i:=1(1)3:	<pre>generates 3 lines named restriction_1</pre>
<pre>restriction_\$i\$: A[,]*x[]&lt;=b[];</pre>	restriction_2
}	restriction_3
restriction { i:=1(1)3:	<pre>generates 3 lines named restriction_1</pre>
A[,]*x[]<=b[];	restriction_2
}	restriction_3

# 3 Expressions

#### 3.1 Overview

Expressions are rules for computing a value during the run-time of a CMPL program. Therefore an expression generally cannot include a model variable. Exceptions to this include special functions whose value depends solely on the definition of a certain model variable. Expressions are a part of an assignment to a parameter or are usable within the echo function. Assignments to a parameter are only permitted within the parameters section or within a control structure. An expression can be a single number or string, a function or a set. Therefore only real, integer, binary, string or set expressions are possible in CMPL. An expression can contain the normal arithmetic operations.

# 3.2 Array functions

With the following functions a user may identify specific characteristics of an array or a single parameter or model variable.

#### Usage:

```
can be a list of numerical expressions separated by commas or can be a multi-
dimensional array of parameters

vector one-dimensional array of parameters or model variables

a scalar parameter or model variable or a multidimensional array of parameters or model variables

array[ [,[, ..]] ] array of parameters or variables with at least one free index
```

#### Examples:

a[]:= (1,2,5);	
echo max(a[]);	returns user message 5
echo min( a[] );	returns user message 1
echo dim( a[] );	returns user message 3
echo def(a[1]);	returns user message 1
echo def(a[5]);	returns user message 0
echo def( a[] );	returns user message 3

#### 3.3 Mathematical functions

In CMPL there are the following mathematical functions which can be used in expressions. With the exception of div and mod all these functions return a real value.

# Usage:

```
#integer division
p div q
p \mod q
              #remainder on division
              #sqrt function
sqrt( x )
exp(x)
              #exp function
ln(x)
              #natural logarithm
lg( x )
              #common logarithm
ld( x )
              #logarithm to the basis 2
srand( x )
              #Initialisation of a pseudo-random number generator using the
              argument x. Returns the value of the argument x.
rand( x )
              \#returns an integer random number in the range 0 \le r and r
sin(x)
              #sine measured in radians
cos(x)
              #cosine measured in radians
tan(x)
              #tangent measured in radians
acos(x)
              #arc cosine measured in radians
              #arc sine measured in radians
asin(x)
atan(x)
              #arc tangent measured in radians
sinh(x)
              #hyperbolic sine
cosh(x)
              #hyperbolic cosine
tanh(x)
              #hyperbolic tangent
abs(x)
              #absolute value
ceil(x)
              #smallest integer value greater than or equal to a given value
floor( x )
              #largest integer value less than or equal to a given value
round( x )
              #simple round
```

```
p, q integer expression
x real or integer expression
```

Examples:

```
value is:
                                         6.000000
c[1] := sqrt(36);
                                         22026.465795
c[2] := exp(10);
                                         2.302585
c[3] := ln(10);
                                         4.000000
c[4] := lg(10000);
                                         3.000000
c[5] := ld(8);
                                         3.000000
                                                   (random number)
c[6] := rand(10);
c[7] := sin(2.5);
                                         0.598472
                                         0.153374
c[8] := cos(7.7);
                                         0.800789
c[9] := tan(10.1);
c[10] := acos(0.1);
                                         1.470629
                                         0.411517
c[11] := asin(0.4);
c[12] := atan(1.1);
                                         0.832981
c[13] := sinh(10);
                                         11013.232875
c[14] := cosh(3);
                                         10.067662
                                         1.000000
c[15] := tanh(15);
c[16] := abs(12.55);
                                         12.000000
c[17] := ceil(12.55);
                                         13.000000
c[18] := floor(-12.55);
                                         -13.000000
c[19] := round(12.4);
                                         12.000000
c[20] := 35 \text{ div } 4;
                                         8
                                         3
c[21] := 35 \mod 4;
```

# 3.4 Type casts

It is useful in some situations to change the type of an expression into another type. A set expression can only be converted to a string. A string can only be converted to a numerical type if it contains a valid numerical string. Every expression can be converted to a string.

#### Usage:

```
type(expression) #type cast
```

#### Examples:

	returns the user messages:
a: real:= 6.666;	
echo integer(a);	7
echo binary(a);	1
a:=0;	
echo binary(a);	0

a := 6.6666;	
echo string(a);	6.666600
b: integer := 100;	
echo real(b);	100.000000
echo binary(b);	1
b := 0;	
echo binary(b);	0
b:= 100;	
echo string(b);	100
c: binary :=1;	
echo real(c);	1.000000
echo integer(c);	1
echo string(c);	1
e: string := "1.888";	
echo real(e);	1.888000
echo integer(e);	1
echo binary(e);	1
e := "";	
echo binary(e);	0

# 3.5 String operations

Especially for displaying strings or numbers with the echo function there are string operations to concatenate and format strings.

#### Usage:

expression which is converted to string

Cannot be a set expression. Such an expression must be converted to a strin expression by a type cast

expression by a type cast

a string expression containing format parameters

 $\ensuremath{\mathsf{CMPL}}$  uses the format parameters of the programming language C. For further

information please consult a C manual.

formatString

# Usage format parameters:

%<flags><width><.precision>specifier

specifier	
ld	integer
lf	real
s	string
flags	
_	left-justify
+	Forces the result to be preceded by a plus or minus sign (+ or -) even for positive numbers.
	By default only negative numbers are preceded with a - sign.
width	
(number)	Minimum number of characters to be printed. If the value to be printed is shorter than this
	number, the result is padded with blank spaces. The value is not truncated even if the result is
	larger.
*	The width is not specified in the format string, but as an additional integer value argument
	preceding the argument that has to be formatted.
.precisio	on
.number	For integer specifiers 1d: precision specifies the minimum number of digits to be written. If
	the value to be written is shorter than this number, the result is padded with leading zeros.
	The value is not truncated even if the result is longer. A precision of $^{\circ}$ means that no character
	is written for the value 0.
	For $lf:$ This is the number of digits to be printed after the decimal point.
	For ${\rm s}\colon$ this is the maximum number of characters to be printed. By default all characters are
	printed until the ending null character is encountered.
	When no precision is specified, the default is 1. If the period is specified without an explicit
	value for precision, 0 is assumed.
.*	The <i>precision</i> is not specified in the <i>format</i> string, but as an additional integer value argument preceding the argument that has to be formatted.

# Examples:

```
a:=66.77777;
echo type(a)+ " " + a + " to string returns the user message
" + format("%10.21f", a); real 66.777770 to string 66.78
```

If you would like to display an entire set concatenating with a string, then you have to use a string cast of your set.

# Example:

s:= set( 7, "qwe", 6, "fe", 5, 8 );	returns the user message
echo "set is " + string(s);	set is set( 7, "qwe", 6, "fe", 5, 8 )

#### 3.6 Set functions

With the following functions a user can identify the specific characteristics of a set.

#### Usage:

```
len(set) #count of the elements of the set - returns an integer

defset(array) #returns the set of the first free index of the array

element element set #returns 1 - if the element is an element of the set
    #returns 0 - otherwise
```

array of parameters or model variables with at least one free index.

set set expression

element an integer or string that is to be checked

#### Examples:

```
a:= set(1, "a", 3, "b", 5, "c");
                                        returns the user message length of the set: 6
echo "length of the set: "+ len(a);
A[,] := ( (1,2,3,4,5),
             (1,2,3,4,5,6,7));
                                        row is assigned the set 1..2
row := defset(A[,]);
col := defset(A[1,]);
                                        col is assigned the set 1..5
a:= set(1, "a", 3, "b", 5, "c");
                                        returns the user message 1
echo "a" element a;
                                        returns the user message 1
echo 5 element a;
echo "bb" element a;
                                        returns the user message 0
```

# 4 Input and output operations

The CMPL input and output operations can be separated into message function, a function that reads the external data and the include statement that reads external CMPL code.

# 4.1 Error and user messages

Both kinds of message functions display a string as a message. In contrast to the echo function an error message terminates the CMPL program after displaying the message.

#### Usage:

```
error expression; #error message - terminates the CMPL program

echo expression; #user message
```

expression

A message that is to be displayed. If the expression is not a string it will be automatically converted to string.

#### Examples:

{a<0: error "negative value"; }	If a is negative an error message is displayed and
	the CMPL program will be terminated.
echo "constant definitions finished";	A user message is displayed.
{ i:=1(1)3: echo "value:" + i;}	The following user messages are displayed:
	value: 1
	value: 2
	value: 3

## 4.2 Readcsv and readstdin

CMPL has two functions that enable a user to read external data. The function <code>readstdin</code> is designed to read a user's numerical input and assign it to a parameter. The function <code>readcsv</code> reads numerical data from a CSV file and assigns it to a vector or matrix of parameters.

#### Usage:

message

string expression for the message that is to be displayed

fileName

string expression for the file name of the CSV file (relative to the directory in

which the current CMPL file resides)

In CMPL CSV files that use a comma or semicolon to separate values are permit-

ted.

# Example:

a := readstdin("give me a number");	reads a value from stdin to be used as value for a.
	Only recommended when using CMPL as a command
	line interpreter.

# The following example uses three CSV files:

1;2;3	c.csv
5.6;7.7;10.5	a.csv
9.8;4.2;11.1	
15;20	b.csv
parameters:	Using readcsv CMPL generates the
<pre>c[] := readcsv("c.csv");</pre>	following model:
<pre>b[] := readcsv("b.csv");</pre>	$1 \cdot x_1 + 2 \cdot x_2 + 3 \cdot x_3 \rightarrow max!$
A[,] := readcsv("a.csv");	s.t.
variables:	$\begin{array}{c c} 5.t. \\ 5.6 \cdot x_1 + 7.7 \cdot x_2 + 10.5 \cdot x_3 \le 15 \end{array}$
x[1dim(c[])]: real[0];	1 2
objectives:	$9.8 \cdot x_1 + 4.2 \cdot x_2 + 11.1 \cdot x_3 \le 20$
c[]T * x[]->max;	$x_{j} \ge 0$ ; $j = 1(1)3$
constraints:	
A[,] * x[] <= b[];	

#### 4.3 Include

Using the include directive it is possible to read external CMPL code in a CMPL program. The CMPL code in the external CMPL file can be used by several CMPL programs. This makes sense for sharing basic data in a couple of CMPL programs or for the multiple use of specific CMPL statements in several CMPL programs. The include directive can stand in any position in a CMPL file. The content of the included file is inserted at this position before parsing the CMPL code. Because include is not a statement it is not closed with a semi-colon.

#### Usage:

```
include "fileName" #include external CMPL code
```

fileName

file name of the CMPL file (relative to the directory in which the current CMPL file resides)

Note that fileName can only be a literal string value. It cannot be a string expression or a string parameter.

The following CMPL file "const-def.gen" is used for the definition of a couple of parameters:

# 5 Statements

As mentioned earlier, every CMPL program consists of at least one of the following sections: parameters:, variables:, objectives: and constraints:. Each section can be inserted several times and mixed in a different order. Every section can contain special statements.

Every statement finishes with a semicolon.

# 5.1 parameters and variables section

Statements in the parameters section are assignments to parameters. These assignments define parameters or reassign a new value to already defined parameters. Statements in the variables sections are definitions of model variables.

All the syntactic and semantic requirements are described in the chapters above.

# 5.2 objectives and constraints section

In the objectives and constraints sections a user has to define the content of the decision model in linear terms. In general, an objective function of a linear optimization model has the form:

$$c_1 \cdot x_1 + c_2 \cdot x_2 + \dots + c_n \cdot x_n \rightarrow max!$$
 (or min!)

with the objective function coefficient  $c_j$  and model variables  $x_j$ . Constraints in general have the form:

$$\begin{array}{ll} k_{11} \cdot x_1 + k_{12} \cdot x_2 + \ldots + k_{1n} \cdot x_n & \leq b_1 \\ k_{21} \cdot x_1 + k_{22} \cdot x_2 + \ldots + k_{2n} \cdot x_n & \leq b_2 \\ \vdots \\ k_{ml} \cdot x_1 + k_{m2} \cdot x_2 + \ldots + k_{mn} \cdot x_n & \leq b_m \end{array}$$

with constraint coefficients  $k_{ii}$  and model variables  $x_i$ .

An objective or constraint definition in CMPL must use exactly this form or a sum loop that expresses this form. A coefficient can be an arbitrary numerical expression, but the model variables cannot stand in expres-

sions that are different from the general form formulated. The rule that model variables cannot stand in bracketed expressions serves to enforce this.

#### Please note, it is not permissible to put model variables in brackets!

The example (a and b are parameters, x and y model variables) a\*x + a\*y + b\*x + b\*ycan be written alternatively (with parameters in brackets) as:

(a + b) \*x + (a + b) \*y

but not (with model variables in brackets) as:

a\*(x + y) + b\*(x + y)

For the definition of the objective sense in the objectives section the syntactic elements ->max or ->min are used. A line name is permitted and the definition of the objective function has to have a linear form.

Usage of an objective function:

objectives: [lineName:] linearTerm ->max|->min;

lineName optional element

description of objective

linearTerm definition of linear objective function

The definition of a constraint has to consist of a linear definition of the use of the constraint and one or two relative comparisons. Line names are permitted.

#### Usage of a constraint:

constraints: [lineName:] linearTerm <=|>=|= linearTerm [<=|>=|= linearTerm];

lineName optional element

description of objective

linearTerm linear definition of use of constraint

# 6 Control structure

#### 6.1 Overview

A control structure is imbedded in { } and defined by a header followed by a body separated off by :.

General usage of a control structure:

```
[controlName]|[sum|set] { controlHeader : controlBody }
```

A control structure can be started with an optional name for the control structure. In the objectives and in the constraints section this name is also used as the line name.

It is possible to define different kinds of control structures based on different headers, control statements and special syntactical elements. Thus the control structure can used for for loops, while loops, if-then-else clauses and switch clauses. Control structures can be used in all sections.

A control structure can be used for the definition of statements. In this case the control body contains one or more statements which are permissible in this section.

It is also possible to use control structures for sum and set as expressions. Then the body contains a single expression. A control structure as an expression cannot have a name because this place is taken by the keyword sum or set. Moreover a control structure as an expression cannot use control statements because the body is an expression and not a statement.

#### 6.2 Control header

A control header consists of one or more control headers. Where there is more than one header, the headers must be separated by commas. Control headers can be divided into iteration headers, condition headers, local assignments and empty headers.

#### 6.2.1 Iteration headers

Iteration headers define how many repeats are to be executed in the control body. Iteration headers are based on sets.

Usage:

```
localParam :=|in set # iteration over a set
```

localParam

name of the local parameter

set

The defined local parameter iterates over the elements of the set and the body is executed for every element in the set.

#### Examples:

s1 := set( "a", "b", "c", "d");	
{k in s1: }	${\tt k}$ is iterated over all elements of the set ${\tt s1}$
s2 := 1(1)10;	
{k in s2: }	$k$ is iterated in the sequence $k \in \{1, 2,, 10\}$
s3 := 26;	
{k := s3:}	$k$ is iterated in the sequence $k \in \{2,3,,6\}$

# **6.2.2 Condition headers**

A condition returns 1 (True) or 0 (False) subject to the result of a comparison or the properties of a parameter or a set. If the condition returns 1 (True) the body is executed once or else the body is skipped.

#### Comparison operators for parameters:

=, ==	equality
<>, !=	inequality
<	less than
>	greater than
<=	equal to or less than
>=	equal to or greater than

## Comparison operators for sets:

=	equality
==	tests whether the iteration order of two sets is equal
<>	inequality
!=	tests whether the iteration order of two sets is not equal
<	subset or not equal
>	greater than
<=	subset or equal
>=	equal to or greater than

## Logical operators:

& &	AND
H	OR
!	NOT

If a real or integer parameter is assigned 0, the condition returns 0 (false). Alternatively if the parameter is assigned 1 the condition returns 1 (true).

#### Examples:

```
\begin{array}{c} \texttt{i}:=1;\\ \texttt{j}:=2;\\ \texttt{\{i>j:=3;}\\ \texttt{\{i>j:=1;}\\ \texttt{\{i>j:=2;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=2;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=2;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=2;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=2;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=2;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=2;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=1;}\\ \texttt{(i>j:=2;}\\ \texttt{(i>j:=1;}\\ \texttt{(i=j:=1;}\\ \texttt{(i
```

# 6.2.3 Local assignments

A local assignment as control header is useful if a user wishes to make several calculations in a local environment. Assigning expression to a parameter within the constraints section is generally not allowed with the exception of a local assignment within a control structure. The body will be executed once.

#### Usage:

```
localParam := expression  # assignment to a local parameter
```

10calParam Defines a local parameter with this name.

expression which is assigned to the local parameter.

#### Examples:

constraints:	${\bf k}$ is assigned ${\bf 1}$ and used as local parameter within the
{ k:=1 : }	control structure.

# 6.3 Alternative bodies

If a control header consists of at least one condition, it is possible to define alternative bodies. Structures like that make sense if a user wishes to combine a for loop with an if-then clause.

The first defined body after the headers is the main body of the control structure. Subsequent bodies must be separated by the syntactic element |. Alternative bodies are only executed if the main body is skipped.

#### Usage:

```
{ controlHeader: mainBody [ | condition1: alternativeBody1 ]
        [ | ... ] [ | default: alternativeDefaultBody ] }
```

controlHeader header of the control structure including at least one condition

The alternative bodies belong to last header of control header. This header cannot be an assignment of a local parameter, because in this case the

main body is never skipped.

main body of control structure

condition1 will be evaluated if alternative body is executed

alternativeBody1 The first alternative body with a condition that evaluates to true is ex-

ecuted. The remaining alternative bodies are skipped without checking the

conditions.

alternativeDefaultBody If no condition evaluates to true then the alternative default body is ex-

ecuted. If the control structure has no alternative default body, then no

body is executed.

#### 6.4 Control statements

It is possible to change or interrupt the execution of a control structure using the keywords continue, break and repeat. A continue stops the execution of the specified loop, jumps to the loop header and executes the next iteration. A break only interrupts the execution of the specified loop. The keyword repeat starts the execution again with the referenced header.

Every control statement references one control header. If no reference is given, it references the innermost header. Possible references are the name of the local parameter which is defined in this head, or the name of the control structure. The name of the control structure belongs to the first head in this control structure.

#### Usage:

```
continue [reference];
break [reference];
repeat [reference];
```

reference

a reference to a control header specified by a name or a local parameter

break [reference]

The execution of the body of the referenced head is cancelled. Remaining

statements are skipped.

If the referenced header contains iteration over a set, the execution for the

remaining elements of the set is skipped.

continue [reference]

The execution of the body of the referenced head is cancelled. Remaining

statements are skipped.

If the referenced header contains iteration over a set, the execution is continued with the next element of the set. For other kinds of headers continue is

equivalent to break.

repeat [reference]

The execution of the body of the referenced header is cancelled. Remaining

statements are skipped.

The execution starts again with the referenced header. The expression in this header is to be evaluated again. If the header contains iteration over a set, the execution starts with the first element. If this header is an assignment to a local parameter, the assignment is executed again. If the header is a condition, the expression is to be checked prior to execution or skipping the body.

# 6.5 Specific control structures

## **6.5.1** For loop

A for loop is imbedded in { } and defined by at least one iteration header followed by a loop body separated off by :. The loop body contains user-defined instructions which are repeatedly carried out. The number of repeats are based on the iteration header definition.

#### Usage:

```
{ iterationHeader [, iterationHeader1] [, ...] : controlBody }
```

iterationHeader
iterationHeader1

controlBody

defined iteration headers

CMPL statements that are executed in every iteration

# Examples:

{ i := 1(1)3 : }	loop counter $\mathtt{i}$ with a start value of 1, an increment of 1 and an end condition of 3
{ i in 13 : }	alternative definition of a loop counter; loop counter i with a start value of 1 and an end condition of 3. (The increment is automatically defined as 1)
<pre>products:= set( "p1", "p2", "p3"); hours[products]:=(20,55,10); {i in products:     echo "hours of product " +         i + " : "+ hours[i]; }</pre>	<pre>for loop using the set products returns user messages hours of product: p1 : 20     hours of product: p2 : 55     hours of product: p3 : 10</pre>
{i := 1(1)2: {j := 2(2)4: A[i,j] := i + j; } }	defines $A[1,2] = 3$ , $A[1,4] = 5$ , $A[2,2] = 4$ and $A[2,4] = 6$

# Several loop heads can be combined. The above example can thus be abbreviated to:

#### 6.5.2 If-then clause

An if-then consists of one condition as control header and user-defined expressions which are executed if the if condition or conditions are fulfilled. Using an alternative default body the if-then clause can be extended to an if-then-else clause.

#### Usage:

```
{ condition: thenBody [| default: elseBody ]}
```

condition If the evaluated condition is true, the code within the body is executed.

thenBody This body is executed if the condition is true.

elseBody This body is executed if the condition is false.

#### Examples:

```
\{i := 1(1)5, j := 1(1)5:
       \{i = j: A[i,j] := 1; \}
                                               definition of the identity matrix with combined loops
       {i != j: A[i,j] := 0; }
                                               and two if-then clauses
\{i := 1(1)5, j := 1(1)5:
       \{i = j: A[i,j] := 1; |
                                               same example, but with one if-then-else clause
        default: A[i,j] := 0; }
i:=10;
                                               example of an if-then-else clause
{ i<10: echo "i less than 10";
                                               {\it returns user message} \; {\it i} \; {\it greater than} \; {\it 9}
  | default: echo "i greater than 9";
sum{ i = j : 1 | default: 2 }
                                               conditional expression, evaluates to 1 if i = j, oth-
                                               erwise to 2
```

## 6.5.3 Switch clause

Using more than one alternative body the if-then clause can be extended to a switch clause.

#### Usage:

```
{ condition1: body1 [| condition2: body2>] [| ... ] [| default: defaultBody ]}
```

If the first condition returns TRUE, only body1 will be executed. Otherwise the next condition condi-tion1 will be verified. body2 is executed if all of the previous conditions are not fulfilled. If no condition returns true, then the defau1tBody is executed.

#### Example:

```
i:=2;
{  i=1: echo "i equals 1";
  | i=2: echo "i equals 2";
  | i=3: echo "i equals 3";
  | default: echo "any other value";
}
example of a switch clause
returns user message i equals 2
```

# 6.5.4 While loop

A while loop is imbedded in { } and defined by a condition header followed by a loop body separated off by : and finished by the keyword repeat. The loop body contains user-defined instructions which are repeatedly carried out until the condition in the loop header is false.

#### Usage:

```
{ condition : statements repeat; }
```

condition

If the evaluated condition is true, the code within the body is executed. This repeats until the condition becomes false.

statements

one or more user-defined CMPL instructions

To prevent an infinite loop the statements in the control body must have an impact on the <code>condition</code>.

#### Examples:

Litarriples.	
i:=2;	while loop with a global parameter
<pre>{i&lt;=4:     A[i] := i;     i := i+1;     repeat; }  {a := 1, a &lt; 5:     echo a;     a := a + 1;     repeat; }</pre>	Can only be used in the parameters section, because the assignment to a global parameter is not permitted in other sections.  defines A[2] = 2, A[3] = 3 and A[4] = 4  while loop using a local parameter Can be used in all sections.  returns user messages 1
repeat;	2 3 4
<pre>{a:=1:     xx {:         echo a;         a := a + 1;         {a&gt;=4: break xx;}         repeat;     } }</pre>	Alternative formulation: The outer control structure defines the local parameter a. This control structure is used as a loop with a defined name and an empty header. The name is necessary, because it is needed as reference for the break statement in the inner control structure. (Without this reference the break statement would refer to the condition a>=4)

# 6.6 Set and sum control structure as expression

Starting with the keyword sum or the keyword set a control structure returns an expression. Only expressions are permitted in the body of the control structure. Control statements are not allowed, because the body cannot contain a statement. It is possible to define alternative bodies.

#### Usage:

```
sum { controlHeader : bodyExpressions }
set { controlHeader : bodyExpressions }
```

controlHeader header of the control structure

The header of a sum or a set control structure is usually an iteration header,

but all kinds of control header can be used.

bodyExpressions user-defined expressions

A sum expression repeatedly summarises the user-defined expressions in the <code>bodyExpressions</code>. If the body is never executed, it evaluates to 0. A set expression returns a set subject to the <code>controlHeader</code> and the <code>bodyExpressions</code>. The element type included in <code>bodyExpressions</code> must be integer or string.

#### Examples:

```
x[1..3] := (2, 4, 6);
                                                           a is assigned 12
a := sum\{i := 1(1)3 : x[i] \};
products:= set( "p1", "p2", "p3");
hours [products] := (20, 55, 10);
                                                           totalHours is assigned 85
totalHours:= sum{i in products: hours[i] };
                                                           using sum with more then one con-
                                                           trol header
x[1..3,1..2] := ((1,2),(3,4),(5,6));
                                                           b is assigned 21.
b:= sum\{i := 1(1)3, j := 1(1)2: x[i,j] \};
s:=set();
                                                           sums up all elements in the set s.
d:= sum{i in s: i |default: -1 };
                                                           Since s is an empty set, d is as-
                                                           signed to the alternative default
                                                           value -1.
e:= set{i:= 1..10: i^2 };
                                                           e is assigned the set
                                                           (1, 4, 9, 16, 25, 36, 49,
                                                           64, 81, 100)
f := set\{i := 1..100, round(sqrt(i))^2 = i: i \};
                                                           f is assigned the set
                                                           (1, 4, 9, 16, 25, 36, 49,
                                                           64, 81, 100)
```

The sum expression can also be used in linear terms for the definition of objectives and constraints. In this case the body of the control structure can contain model variables.

# Examples:

```
parameters: a[1..2,1..3] := ((1,2,3),(4,5,6)); b[1..2] := (100,100); c[1..3] := (20,10,10);
```

# 6.7 Implicit loops

As mentioned above it is possible to define objectives and constraints using control structures as loops. The syntax of these control structures is easy to understand and to use, but it follows the idea of programming languages. For a formulation of objectives and constraints in a more mathematical way it is simpler to use implicit loops. Implicit loops allow users to define objectives and constraints in a mathematical notation (e.g. matrix vector multiplication). All mathematical requirements are applied for implicit loops. Implicit loops are only possible in the objectives section and the constraints section.

Implicit loops are formed by matrices and vectors, which are defined by the use of free indices.

A free index is an index which is not specified by a position in an array. It can be specified by an entire set or without any specification. But the separating commas between indices must in any case be specified.

A multidimensional array with one free index is always treated as a column vector, regardless of where the free index stands. A column vector can be transposed to a row vector with  $\mathbb{T}$ . A multidimensional array with two free indices is always treated as a matrix. The first free index is the row, the second the column.

#### Usage:

```
vector[[set]] #column vector
vector[[set]]T #transpose of column vector - row vector

matrix[index, [set]] #column vector

matrix[[set], index] #also column vector

matrix[index, [set]]T #transpose of column vector - row vector

matrix[[set], index]T #transpose of column vector - row vector

matrix[[set], index]T #transpose of column vector - row vector

matrix[[set1], [set2]] #matrix
```

#### Examples:

x[]	vector with free index across the entire defined area
x[25]	vector with free index in the range 2 – 5
A[,]	matrix with two free indices
A[1,]	matrix with one fixed and one free index; this is a
	column vector.
A[,1]	matrix with one fixed and one free index; this is also a
	column vector.

The most important ways to define objectives and constraints with implicit loops are vector-vector multiplication and matrix-vector multiplication. A vector-vector multiplication defines a row of the model (e.g. an objective or one constraint). A matrix-vector multiplication can be used for the formulation of more than one row of the model.

#### Usage of multiplication using implicit loops:

paramVectorname of a vector of parametersvarVectorname of a vector of model variablesparamMatrixname of a matrix of parametersTsyntactic element for transposing a vector

#### Examples:

```
parameters:
        a[1..2,1..3] := ((1,2,3),
                                 (4,5,6));
        b[1..2] := (100,100);
        c[1..3] := (20,10,10);
variables:
        x[1...3]: real[0...];
objectives:
                                                      objective definition using implicit loops
        c[]T * x[] ->max;
                                                         20 \cdot x_1 + 10 \cdot x_2 + 10 \cdot x_3 \rightarrow max!
constraints:
                                                      constraint definition using implicit loops
        a[, ] * x[] <=b[];
                                                         1 \cdot x_1 + 2 \cdot x_2 + 3 \cdot x_3 \le 100
                                                        4 \cdot x_1 + 5 \cdot x_2 + 6 \cdot x_3 \le 100
```

Aside from vector-vector multiplication and matrix-vector multiplication vector subtractions or additions are also useful for the definition of constraints. The addition or subtraction of a variable vector adds new columns to the constraints. The addition or subtraction of a constant vector changes the right side of the constraints.

Usage of additions or subtractions using implicit loops:

```
linearTerms + varVector[[set]]  #variable vector addition
linearTerms - varVector[[set]]  #variable vector subtraction

linearTerms + paramVector[[set]]  #parameter vector addition
linearTerms - paramVector[[set]]  #parameter vector subtraction
```

linearTerms

other linear terms in an objective or constraint

#### Examples:

```
parameters:
      a[1..2,1..3] := ((1,2,3),
                            (4,5,6));
      b[1..2] := (100,100);
      d[1..2] := (10,10);
      c[1..3] := (20,10,10);
variables:
      x[1..3]: real[0..];
objectives:
      c[]T * x[] ->max;
constraints:
                                             constraints definition using implicit loops
      a[, ] * x[] + d[] <=b[];
                                               1 \cdot x_1 + 2 \cdot x_2 + 3 \cdot x_3 \le 90
                                               4 \cdot x_1 + 5 \cdot x_2 + 6 \cdot x_3 \le 90
                                             equivalent to
                                             a[, ] * x[] <=b[] - d[];
0 \le x[1..3] + y[1..3] + z[2] \le b[1..3];
                                             implicit loops for a column vector.
0 \le x[1] + y[1] + z[2] \le b[1];
                                             equivalent formulation
0 \le x[2] + y[2] + z[2] \le b[2];
0 \le x[3] + y[3] + z[2] \le b[3];
parameters:
      a[1..2,1..3] := ((1,2,3),
                            (4,5,6));
      b[1..2] := (100,100);
      d[1..2] := (10,10);
      c[1..3] := (20,10,10);
variables:
      x[1...3]: real[0..];
      z[1..2]: real[0..];
```

objectives:	
c[]T * x[] ->max;	
constraints:	constraints definition using implicit loops
a[,] * x[] + z[] <=b[];	$1 \cdot x_1 + 2 \cdot x_2 + 3 \cdot x_3 + z_1 \leq 90$
	$4 \cdot x_1 + 5 \cdot x_2 + 6 \cdot x_3 + z_2 \le 90$

# 7 Automatic code generating

#### 7.1 Overview

CMPL includes two types of automatic code generation which release the user from additional modelling work. CMPL automatically optimizes the generated model by means of matrix reductions. The second type of automatic code generation is the equivalent transformation of variable products. If a CMPL program includes a variable product with at least one integer factor, CMPL will transform this non-linear form equivalent in a set of linear inequations.

## 7.2 Matrix reductions

Matrix reductions are subject to constraints of a specific form.

a) If a constraint contains only one variable or only one of the variables with a coefficient not equal to 0, then the constraint is taken as a lower or upper bound.

```
For the following summation (x[] is a variable vector) sum\{i:=1 (1) 2: (i-1) * x[i]\} \le 10; no matrix line is generated; rather x[2] has an upper bound of 10.
```

b) If there is a constraint in the coefficients of all variables proportional to another constraint, only the more strongly limiting constraint is retained.

```
Only the second of the two constraints (x[] is a variable vector)
2*x[1] + 3*x[2] <= 20;
10*x[1] + 15*x[2] <= 50;
is used in generating a model line.
```

# **7.3** Equivalent transformations of Variable Products

In general a product of variables like  $x \cdot y$  cannot be a part of an LP or MIP model, because such a variable product is a non-linear term. But it is possible to formulate an equivalent transformation using a set of

specific inequations. The automatic generation of an equivalent transformation of a variable product is a special capability characteristic of CMPL.

# **7.3.1** Variable Products with at least one binary variable

For the following given variables

```
variables: x: binary;
   y: real[YU..YO];
```

each occurrence of the term x\*y in the CMPL model description is replaced by an implicit newly-defined variable x y, and the following additional statements are generated automatically:

```
constraints:
    min(YU, 0) <= x_y <= max(YO, 0);
    {YU < 0: x_y - YU*x >= 0; }
    {YO > 0: x_y - YO*x <= 0; }
    y - x_y + YU*x >= YU;
    y - x_y + YO*x <= YO;</pre>
```

# 7.3.2 Variable Product with at least one integer variable

For the following given variables

```
variables: x: integer[XU..XO];
    y: real[YU..YO];
```

each occurrence of the term x\*y in the CMPL model description is replaced by an implicit newly-defined variable  $x_y$ , and the following additional statements are generated automatically (here n stands for the number of binary positions needed for xo-xu+1):

```
variables:
    _x[1..n]: binary;
    _x_y[1..n]: real;

constraints:

min(XU*YU, XU*YO, XO*YU, XO*YO) <= x_y <= max(XU*YU, XU*YO, XO*YU, XO*YO);

x = XU + sum{i=1(1)n: (2^(i-1))*_x[i]};

x_y = XU*y + sum{i=1(1)n: (2^(i-1))*_x_y[i]};</pre>
```

```
{i = 1(1)n:
    min(YU, 0) <= _x_y[i] <= max(YO, 0);
    {YU < 0: _x_y[i] - YU*_x[i] >= 0; }
    {YO > 0: _x_y[i] - YO*_x[i] <= 0; }
    y - _x_y[i] + YU*_x[i] >= YU;
    y - _x_y[i] + YO*_x[i] <= YO;
}</pre>
```

# 8 CMPL as command line tool

# 8.1 Usage

The CMPL command line tool can be used in two modes. Using the solver mode, an LP or MIP can be formulated, solved and analysed. In this mode, the OSSolverservice is invoked. In the model mode it is possible to transform the mathematical problem into MPS, Free-MPS or OSiL files that can be used by certain alternative LP or MIP solvers.

```
cmpl [<options>] <input file>
```

#### Model mode:

-i <cmplFile> : input file

-m [<File>] : export model in MPS format in a file or stdout

-x [<File>] : export model in OSiL XML format in a file or stdout

-noOutput : no generating of a MPS or OSiL file

# **Solver mode:**

-solver <solver> : name of the solver you want to use

Possible values are clp (COIN-OR Clp), cbc (COIN-OR Cbc), symphony

(COIN Symphony) and glpk (glpk)

-solverUrl <url> : URL of the solver service

w/o a defined remote solver service, a local solver is used

-osol <file> : name of the file that contains the solver options

-solutionCsv : optimization results in CSV format

A file <cmplFileName>.csv will be created.

-solutionAscii : optimization results in ASCII format

A file <cmplFileName>.sol will be created.

-silent : optimization results are not displayed

-obj <objName> : name of the objective function

-objSense <max/min> : objective sense

### **General options:**

-e [<File>] : output for error messages and warnings

-e simple output to stderr

-e<File> ouput in MprL XML format to file

-l [<File>] : output for replacements for products of variables

-s [<File>] : output for short statistic info

-p [<File>] : output for protocol

-gn : generation option: do not make reductions

-gf : generation option: constraints for products of variables follow the product

-cd : no warning at multiple parameter definition

-ca : no warning at deprecated '=' assignment

-ci<X> : mode for integer expressions (0 - 3), defaults to 1

-fc<X> : format option: maximal length of comment, defaults to 60

-ff : format option: generate free MPS

-f%<format> : format option: format spezifier for float number output, defaults to %f

-h : get this help

-v : version

## Examples - solver mode:

cmpl test.cmpl	solves the problem test.cmpl locally with the de-
	fault solver and displays a standard solution report
cmpl -solver glpk test.cmpl	solves the problem test.cmpl locally using GLPK
	and displays a standard solution report
cmpl -solverUrl http://gsbkip.chica-	solves the problem test.cmpl remotely with the
gogsb.edu/os/OSSolverService.jws	defined web service and displays a standard solution
test.cmpl	report
cmpl -solutionCsv test.cmpl	solves the problem test.cmpl locally with the de-
	fault solver writes the solution in the CSV-file
	test.csv and displays a standard solution report

## Examples - model mode:

cmpl -i test.cmpl -m test.mps	reads the file test.cmpl and generates the MPS-
	file test.mps.
cmpl -ff -i test.cmpl -m test.mps	reads the file test.cmpl and generates the Free-
	MPS-file test.mps.
cmpl -i test.cmpl -x test.osil	reads the file test.cmpl and generates the OS-
	iL-file test.osil.

# 8.2 Input and output file formats

## 8.2.1 Overview

CMPL uses several ASCII files for the communication with the user and other programs such as solvers.

CMPL	input file for CMPL - syntax as described above
MPS	output file for the generated model in MPS format
	Can be used with most solvers.
	This format is very restrictive and therefore not recommended.
Free-MPS	output file for the generated model in Free-MPS format
	Can be used with most solvers.

OSiL	output file for the generated model in OSiL format
	The OSiL XML schema is developed by the COIN-OR community (COmputational IN-
	frastructure for Operations Research - open source for the operations research com-
	munity).
	Can be used with solvers which are supported by the COIN-OR Optimization Services
	(OS) Framework.
OSoL	OSoL is an XML-based language for representing options that get passed to an op-
	timization solver or a hosted optimization solver Web service.
MprL	output file for the status of the results or errors of a CMPL model
	XML file in accordance with the MprL schema

## 8.2.2 CMPL

A CMPL file is an ASCII file that includes the user-defined CMPL code with a syntax as described in this manual.

The example

$$\begin{array}{l} 1 \cdot x_1 + 2 \cdot x_2 + 3 \cdot x_3 \rightarrow max \ ! \\ s.t. \\ 5.6 \cdot x_1 + 7.7 \cdot x_2 + 10.5 \cdot x_3 \leq 15 \\ 9.8 \cdot x_1 + 4.2 \cdot x_2 + 11.1 \cdot x_3 \leq 20 \\ 0 \leq x_n \ \ ; n = 1(1)3 \end{array}$$

#### can be formulated in CMPL as follows:

#### 8.2.3 MPS

An MPS (Mathematical Programming System) file is a ASCII file for presenting linear programming (LP) and mixed integer programming problems.

MPS is an old format and was the de facto standard for most LP solvers. MPS is column-oriented and is set up for punch cards with defined positions for fields. Owing to these requirements the length of column or row names and the length of a data field are restricted. MPS is very restrictive and therefore not recommended. For more information please see http://en.wikipedia.org/wiki/MPS\_(format).

The MPS file for the CMP example given in the section above is generated as follows:

```
* CMPL - MPS - Export
NAME
           test.mps
* OBJNAME profit
* OBJSENSE max
ROWS
N profit
L machine1
L machine2
COLUMNS
   x1
          profit
                            1 machine1 5.600000
          machine2
                     9.800000
   x1
          profit
                               machine1 7.700000
   x2
   x2
          machine2
                      4.200000
           profit
                                machine1 10.500000
   xЗ
           machine2 11.100000
   xЗ
RHS
                               machine2
   RHS
           machine1
                     15
                                                  20
RANGES
BOUNDS
ENDATA
```

## 8.2.4 Free - MPS

The Free-MPS format is an improved version of the MPS format. There is no standard for this format but it is widely accepted. The structure of a Free-MPS file is the same as an MPS file. But most of the restricted MPS format requirements are eliminated, e.g. there are no requirements for the position or length of a field. For more information please visit the project website of the lp\_solve project. (http://lpsolve.sourceforge.net)

The Free-MPS file for the given CMP example is generated as follows:

```
* CMPL - Free-MPS - Export
NAME
             test.mps
* OBJNAME profit
* OBJSENSE max
ROWS
N profit
L machine1
L machine2
COLUMNS
x1 profit 1 machine1 5.600000
x1 machine2 9.800000
x2 profit 2 machine1 7.700000
x2 machine2 4.200000
 x3 profit 3 machine1 10.500000
x3 machine2 11.100000
RHS
RHS machine1 15 machine2 20
RANGES
BOUNDS
ENDATA
```

## 8.2.5 OSiL

OSiL is an XML-based format which can be used for presenting linear programming (LP) and mixed integer programming problems. The OSiL XML schema was developed by the COIN-OR community (COmputational INfrastructure for Operations Research - open source for the operations research community). The format makes it very easy to save and present a model and so is particularly suitable for defining an interface to several solvers. An OSiL file can be used with solvers which are supported by the COIN-OR Optimization Services (OS) Framework.

For more information please visit the project website of COIN-OR OS project. (https://projects.coin-or.org/OS or http://www.optimizationservices.org)

The OSiL file for the given CMP example is generated as follows:

```
<name>test.gen</name>
           <description>generated by CMPL v1.4.2</description>
     </instanceHeader>
     <instanceData>
           <variables numberOfVariables="3">
                 <var name="x1" type="C" lb="0"/>
                 <var name="x2" type="C" lb="0"/>
                 <var name="x3" type="C" lb="0"/>
           </variables>
           <objectives numberOfObjectives="1">
                 <obj name="profit" maxOrMin="max" numberOfObjCoef="3">
                       <coef idx="0">1</coef>
                       <coef idx="1">2</coef>
                       <coef idx="2">3</coef>
                 </obj>
           </objectives>
           <constraints numberOfConstraints="2">
                 <con name="machine1" ub="15"/>
                 <con name="machine2" ub="20"/>
           </constraints>
           <linearConstraintCoefficients numberOfValues="6">
                 <start>
                       <el>0</el>
                       <el>2</el>
                       <el>4</el>
                       <el>6</el>
                 </start>
                 <rowIdx>
                       <el>0</el>
                       <el>1</el>
                       <el>0</el>
                       <el>1</el>
                       <el>0</el>
                       <el>1</el>
                 </rowIdx>
                 <value>
                       <el>5.600000</el>
                       <el>9.800000</el>
                       <el>7.700000</el>
                       <el>4.200000</el>
                       <el>10.500000</el>
                       <el>11.100000</el>
                 </value>
           </linearConstraintCoefficients>
     </instanceData>
</osil>
```

## 8.2.6 OsoL

OSoL is an XML-based language for representing options that get passed to an optimization solver or a hosted optimization solver Web service. For more information please visit the project website of COIN-OR OS project. (https://projects.coin-or.org/OS or http://www.optimizationservices.org)

The following OSoL-file describes a couple of parameters for the CBC solver. (See the Optimization Services User's Manual by Horand Gassmann, Jun Ma, Kipp Martin, and Wayne Sheng)

```
<?xml version="1.0" encoding="UTF-8"?>
<osol xmlns="os.optimizationservices.org"</pre>
    xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance"
    xsi:schemaLocation="os.optimizationservices.org
    http://www.optimizationservices.org/schemas/2.0/OSoL.xsd">
    <general>
      <solverToInvoke>cbc</solverToInvoke>
    </general>
    <job>
    </job>
    <optimization>
      <solverOptions numberOfSolverOptions="4">
         <solverOption name="reslim" solver="cbc" type="numeric" value="1"/>
         <solverOption name="maxN" solver="cbc" value="5" />
         <solverOption name="cuts" solver="cbc" value="off" />
         <solverOption name="max active nodes" solver="symphony" value="2" />
      </solverOptions>
    </optimization>
</osol>
```

## 8.2.7 MprL

MprL is an XML-based format for representing the general status and/or errors of the transformation of a CMPL model in one of the described output files. MprL is intended for communication with other software that uses CMPL for modelling linear optimization problems.

An MprL file consists of two major sections. The <qeneral> section describes the general status and the name of the model and a general message after the transformation. The <mplResult> section consists of one or more messages about specific lines in the CMPL model.

After the transformation of the given CMPL model, CMPL will finish without errors. The general status is represented in the following MprL file.

# If a semicolon is not set in line 26, CMPL will finish with errors that are represented in the following MprL file.

#### The MprL schema is defined as follows:

```
<?xml version="1.0" encoding="UTF-8"?>
<xsd:schema xmlns:xsd="http://www.w3.org/2001/XMLSchema" >
    <xsd:element name="mprl">
      <xsd:complexType>
         <xsd:sequence>
           <xsd:element ref="general" minOccurs="1" maxOccurs="1"/>
            <xsd:element ref="mplResult" minOccurs="0" maxOccurs="unbounded"/>
         </xsd:sequence>
      </xsd:complexType>
    </xsd:element>
    <xsd:element name="general">
      <xsd:complexType>
         <xsd:sequence>
           <xsd:element ref="generalStatus" minOccurs="1" maxOccurs="1"/>
           <xsd:element name="instanceName" type="xsd:string" minOccurs="1"</pre>
                  maxOccurs="1"/>
           <xsd:element name="mplName" type="xsd:string" minOccurs="1"</pre>
                  maxOccurs="1"/>
           <xsd:element name="message" type="xsd:string" minOccurs="0"</pre>
                  maxOccurs="unbounded"/>
         </xsd:sequence>
      </xsd:complexType>
    </xsd:element>
    <xsd:element name="generalStatus">
```

```
<xsd:simpleType>
         <xsd:restriction base="xsd:string">
           <xsd:enumeration value="error"/>
           <xsd:enumeration value="warning"/>
           <xsd:enumeration value="normal"/>
         </xsd:restriction>
      </xsd:simpleType>
    </xsd:element>
    <xsd:element name="mplResult">
      <xsd:complexType>
         <xsd:sequence>
           <xsd:element ref="mplMessage" minOccurs="0" maxOccurs="unbounded"/>
         </xsd:sequence>
         <xsd:attribute name="numberOfMessages" type="xsd:nonNegativeInteger"</pre>
           use="required"/>
      </xsd:complexType>
    </xsd:element>
    <xsd:element name="mplMessage">
      <xsd:complexType>
         <xsd:attribute name="type" use="required">
           <xsd:simpleType>
             <xsd:restriction base="xsd:string">
                  <xsd:enumeration value="error"/>
                  <xsd:enumeration value="warning"/>
             </xsd:restriction>
           </xsd:simpleType>
         </xsd:attribute>
         <xsd:attribute name="file" type="xsd:string" use="required"/>
         <xsd:attribute name="line" type="xsd:nonNegativeInteger"</pre>
             use="required"/>
         <xsd:attribute name="description" type="xsd:string" use="required"/>
      </xsd:complexType>
    </xsd:element>
</xsd:schema>
```

## 8.3 Using CMPL with several solvers

Since CMPL transforms a CMPL model into an MPS, a Free-MPS or an OSiL file, the model can be solved using most free or commercial solvers.

## 8.3.1 Coliop3

CMPL is recommended for use as an integral part of Coliop3. Coliop3 is an IDE (Integrated Development Environment) intended to solve linear programming (LP) problems and mixed integer programming (MIP) problems. This project contains CMPL and two solvers for LP and MIP problems (glpk and lpsolve). CMPL and Coliop3 are projects of the Technical University of Applied Sciences Wildau and the Institute for Operations Research and Business Management at the Martin Luther University Halle-Wittenberg. Coliop3 and CMPL are

open source projects licensed under GPL and available for all relevant operating systems. For more information please visit the Coliop3 project website: http://www.coliop.org.

### 8.3.2 **GLPK**

The GLPK (GNU Linear Programming Kit) package is intended for solving large-scale linear programming (LP), mixed integer programming (MIP), and other related problems. It is a set of routines written in ANSI C and organized in the form of a callable library. For more information please visit the GLPK project website: http://www.gnu.org/software/glpk/. The CMPL package contains GLPK. But it is also possible to use CMPL and GLPK separately.

Using CMPL with glpk:

```
cmpl -ff -i cmplFilename -m mpsFileName
glpsol --freemps --max|min --output solutionName mpsFileName
```

cmp1Filename name of the CMPL file - CMPL input
mpsFileName name of the MPS file - CMPL output
solutionName name of the solution file - GLPK output

### 8.3.3 LPSolve

Mixed Integer Linear Programming (MILP) solver lp\_solve solves pure linear, (mixed) integer/binary, semi-continuous and special ordered set (SOS) models. lp\_solve is written in ANSI C and can be compiled on many different platforms including Linux and WINDOWS.

For more information please visit the GLPK project website: http://sourceforge.net/projects/lpsolve/.

Using CMPL with LPSolve:

```
cmpl -ff -i cmplFilename -m mpsFileName
lp_solve -max|min -fmps mpsFileName -S4 solutionName
```

cmplFilename name of the CMPL file - CMPL input
mpsFileName name of the MPS file - CMPL output
solutionName name of the solution file - LPSolve output

## 9 Examples

## 9.1 Selected decision problems

## 9.1.1 The diet problem

The goal of the diet problem is to find the cheapest combination of foods that will satisfy all the daily nutritional requirements of a person for a week.

The following data is given (example based on Fourer/Gay/Kernigham: AMPL, 2nd ed., Thomson 2003, p. 27ff.):

food	cost per package	provision of daily	vitamin requireme	nts in percentages	;
		Α	B1	B2	С
BEEF	3.19	60	20	10	15
СНК	2.59	8	2	20	520
FISH	2.29	8	10	15	10
НАМ	2.89	40	40	35	10
MCH	1.89	15	35	15	15
MTL	1.99	70	30	15	15
SPG	1.99	25	50	25	15
TUR	2.49	60	20	15	10

The decision is to be made for one week. Therefore the combination of foods has to provide at least 700% of daily vitamin requirements. To promote variety, the weekly food plan must contain between 2 and 10 packages of each food.

The mathematical model can be formulated as follows:

$$3.19 \cdot x_{\textit{BEEF}} + 2.59 \cdot x_{\textit{CHK}} + 2.29 \cdot x_{\textit{FISH}} + 2.89 \cdot x_{\textit{HAM}} + 1.89 \cdot x_{\textit{MCH}} + 1.99 \cdot x_{\textit{MTL}} + 1.99 x_{\textit{SPG}} + 2.49 \cdot x_{\textit{TUR}} \rightarrow \textit{min} \,! \\ s.t. \\ 60 \cdot x_{\textit{BEEF}} + 8 \cdot x_{\textit{CHK}} + 8 \cdot x_{\textit{FISH}} + 40 \cdot x_{\textit{HAM}} + 15 \cdot x_{\textit{MCH}} + 70 \cdot x_{\textit{MTL}} + 25 x_{\textit{SPG}} + 60 \cdot x_{\textit{TUR}} \leq 700 \\ 20 \cdot x_{\textit{BEEF}} + 0 \cdot x_{\textit{CHK}} + 10 \cdot x_{\textit{FISH}} + 40 \cdot x_{\textit{HAM}} + 35 \cdot x_{\textit{MCH}} + 30 \cdot x_{\textit{MTL}} + 50 x_{\textit{SPG}} + 20 \cdot x_{\textit{TUR}} \leq 700 \\ 10 \cdot x_{\textit{BEEF}} + 20 \cdot x_{\textit{CHK}} + 15 \cdot x_{\textit{FISH}} + 35 \cdot x_{\textit{HAM}} + 15 \cdot x_{\textit{MCH}} + 15 \cdot x_{\textit{MTL}} + 25 x_{\textit{SPG}} + 15 \cdot x_{\textit{TUR}} \leq 700 \\ 15 \cdot x_{\textit{BEEF}} + 20 \cdot x_{\textit{CHK}} + 10 \cdot x_{\textit{FISH}} + 10 \cdot x_{\textit{HAM}} + 15 \cdot x_{\textit{MCH}} + 15 \cdot x_{\textit{MTL}} + 15 x_{\textit{SPG}} + 10 \cdot x_{\textit{TUR}} \leq 700 \\ x_{i} \in \{2, 3, \dots, 10\} \quad ; \ j \in \{\textit{BEEF}, \textit{CHK}, \textit{DISH}, \textit{HAM}, \textit{MCH}, \textit{MTL}, \textit{SPG}, \textit{TUR}\} \\ \end{cases}$$

The CMPL model diet.cmpl is formulated as follows:

```
parameters:
    NUTR := set("A","B1", "B2", "C");
    FOOD := set("BEEF", "CHK", "FISH", "HAM", "MCH", "MTL", "SPG", "TUR");
```

```
#cost per package
     costs[FOOD] := ( 3.19, 2.59, 2.29, 2.89, 1.89, 1.99, 1.99, 2.49 );
     #provision of the daily requirements for vitamins in percentages
     vitamin[NUTR, FOOD] := ( (60, 8, 8, 40, 15, 70, 25, 60) ,
                                (20, 0, 10, 40, 35, 30, 50, 20),
                                (10, 20, 15, 35, 15, 15, 25, 15),
                                (15, 20, 10, 10, 15, 15, 15, 10)
                             );
     #weekly vitamin requirements
     vitMin[NUTR] := (700,700,700,700);
variables:
     x[FOOD]: integer[2..10];
objectives:
     cost: costs[]T * x[]->min;
constraints:
     # capacity restriction
     $2$: vitamin[,] * x[] >= vitMin[];
```

#### CMPL command:

```
cmpl diet.cmpl
```

#### Solution:

Problem : SolverName Nr. of Solutions:	1				
Solution nr.:	1				
Objective name : Objective value :					
Variables					
Nr. of variables:	8				
Jame	Туре	Activity	Lower Bound	Upper Bound	Marginal
BEEF	I	2.00	2.00	10.00	
_ K_CHK	I	8.00	2.00	10.00	-
k FISH	I	2.00	2.00	10.00	-
	I	2.00	2.00	10.00	-
_				10.00	_
_ ×_HAM	I	10.00	2.00	10.00	
_ x_HAM x_MCH	I	10.00 10.00	2.00 2.00	10.00	-
x_HAM x_MCH x_MTL x_SPG	_				-

Constraints					
Nr. of constr	aints: 4				
Name	Туре	Activity	Lower Bound	Upper Bound	Marginal
A	G	1500.00	700.00	Infinity	-
B1	G	1330.00	700.00	Infinity	-
В2	G	860.00	700.00	Infinity	-
С	G	700.00	700.00	Infinity	_

## 9.1.2 Production mix

This model calculates the production mix that maximizes profit subject to available resources. It will identify the mix (number) of each product to produce and any remaining resource.

The example involves three products which are to be produced with two machines. The following data is given:

		P1	P2	P3	upper bounds [h]
upper bound of a product	[units]	250	240	250	
selling price per unit	[€/unit]	500	600	450	
direct costs per unit	[€/unit]	425	520	400	
profit contribution per unit	[€/unit]	75	80	50	
machine hours required per unit					
machine 1	[h/unit]	8	15	12	1,000
machine 2	[h/unit]	15	10	8	1,000

The mathematical model can be formulated as follows:

```
75 \cdot x_1 + 80 \cdot x_2 + 50 \cdot x_3 \rightarrow max ?
s.t.
8 \cdot x_1 + 15 \cdot x_2 + 12 \cdot x_3 \le 1,000
15 \cdot x_1 + 10 \cdot x_2 + 8 \cdot x_3 \le 1,000
x_1 \in \{0, 1, \dots, 250\}
x_2 \in \{0, 1, \dots, 240\}
x_3 \in \{0, 1, \dots, 250\}
```

The CMPL model production-mix.cmpl is formulated as follows:

```
parameters:
    price[] := (500, 600, 450 );
    costs[] := (425, 520, 400);

#machine hours required per unit
    a[,] := ((8, 15, 12), (15, 10, 8));

#upper bounds of the machines
    b[] := (1000, 1000);
```

```
#profit contribution per unit
    {j:=1(1)dim(price[]): c[j] := price[j]-costs[j]; }

#upper bound of the products
    xMax[] := (250, 240, 250);

variables:
    x[1..dim(price[])]: integer;

objectives:
    profit: c[]T * x[] ->max;

constraints:
    res_$2$: a[,] * x[] <= b[];
    0<=x[]<=xMax[];</pre>
```

#### CMPL command:

```
cmpl production-mix.cmpl -solver glpk
```

#### Solution:

```
General status: normal
Problem: production-m
SolverName COIN-OR glpk
            production-mix.cmpl
Nr. of Solutions: 1
-----
Solution nr.:
Objective name : profit
Objective value : 6395 (max!)
Variables
Nr. of variables: 3
Name Type Activity Lower Bound Upper Bound Marginal
              I 33.00 0.00 250.00
x1
                        49.00
              I
                                     0.00
                                                 240.00
            I
                   0.00
                               0.00 250.00
Nr. of constraints: 2
Name Type Activity Lower Bound Upper Bound
                                                           Marginal
             L 999.00 -Infinity
res_1
                                                 1000.00
                                                 1000.00
res 2
              L
                        985.00
                                   -Infinity
```

## 9.1.3 Production mix including thresholds and step-wise fixed costs

This model calculates the production mix that maximizes profit subject to available resources. When a product is produced, there are fixed set-up costs. There is also a threshold for each product. The quantity of a product is zero or greater than the threshold.

The example involves three products which are to be produced with two machines. The following data is given:

		P1	P2	P3	upper bounds [h]
production minimum of a product	[units]	45	45	45	
upper bound of a product	[units]	250	240	250	
selling price per unit	[€/unit]	500	600	450	
direct costs per unit	[€/unit]	425	520	400	
profit contribution per unit	[€/unit]	75	80	50	
set-up costs	[€]	500	400	500	
machine hours required per unit					
machine 1	[h/unit]	8	15	12	1,000
machine 2	[h/unit]	15	10	8	1,000

The mathematical model can be formulated as follows:

$$75 \cdot x_1 + 80 \cdot x_2 + 50 \cdot x_3 - 500 \cdot y_1 - 400 \cdot y_2 - 500 \cdot y_3 \rightarrow \max!$$

$$s.t.$$

$$8 \cdot x_1 + 15 \cdot x_2 + 12 \cdot x_3 \le 1,000$$

$$15 \cdot x_1 + 10 \cdot x_2 + 8 \cdot x_3 \le 1,000$$

$$45 \cdot y_1 \le x_1 \le 250 \cdot y_1$$

$$45 \cdot y_2 \le x_2 \le 240 \cdot y_2$$

$$45 \cdot y_3 \le x_3 \le 250 \cdot y_3$$

$$x_1 \in \{0,1, \dots, 250\}$$

$$x_2 \in \{0,1, \dots, 240\}$$

$$x_3 \in \{0,1, \dots, 250\}$$

$$y_j \in \{0,1\} \quad ; j = 1(1)3$$

The CMPL model production-mix-fixed-costs.cmpl is formulated as follows:

```
parameters:
    price[] := (500, 600, 450 );
    costs[] := (425, 520, 400);

#machine hours required per unit
    a[,] := ((8, 15, 12), (15, 10, 8));
    #upper bounds of the machines
    b[] := (1000, 1000);

#profit contribution per unit
    {j:=1(1)dim(price[]): c[j] := price[j]-costs[j]; }
```

```
#upper bound of a product
    xMax[] := (250, 240, 250 );
    xMin[] := (45, 45, 45 );

#fixed setup costs
    FC[] := ( 500, 400, 500);

variables:
    {j:=1(1)dim(c[]): x[j]: integer[0..xMax[j]]; }
    y[1..dim(c[])]: binary;

objectives:
    profit: c[]T * x[] - FC[]T * y[] ->max;

constraints:
    res_$2$: a[,] * x[] <= b[];
    {j:=1(1)dim(c[]): xMin[j] * y[j] <= x[j] <= xMax[j] * y[j]; }</pre>
```

#### CMPL command:

```
cmpl production-mix-fixed-costs.cmpl
```

## Solution:

General status: Problem : SolverName Nr. of Solutions:	COIN-OR cbc	fixed-costs	.cmpl			
Solution nr.:	1					
Objective name : Objective value :	-					
Variables Nr. of variables:	6					
Name		Type	Activity	Lower Bound	Upper Bound	Marginal
x1		I	0.00	0.00	250.00	
x2		I	66.00	0.00	240.00	-
x3		I	0.00	0.00	250.00	=
y1		В	0.00	0.00	1.00	=
y2		В	1.00	0.00	1.00	=
у3 		В	0.00	0.00	1.00	
Constraints Nr. of constraints	· 8					
Name	. •	Туре	Activity	Lower Bound	Upper Bound	Marginal
res_1		L	990.00	-Infinity	1000.00	
res_2		L	660.00	-Infinity	1000.00	=
line_4		L	0.00	-Infinity	0.00	=
line_5		L	0.00	-Infinity	0.00	-
.ine_6		L	-21.00	-Infinity	0.00	=
.ine_7		L	-174.00	-Infinity	0.00	-
line_8		L	0.00	-Infinity	0.00	-
line 9		L	0.00	-Infinity	0.00	

## 9.1.4 The knapsack problem

Given a set of items with specified weights and values, the problem is to find a combination of items that fills a knapsack (container, room, ...) to maximize the value of the knapsack subject to its restricted capacity or to minimize the weight of items in the knapsack subject to a predefined minimum value.

In this example there are 10 boxes, which can be sold on the market at a defined price.

box number	weight	price
	[pounds]	[€/box]
1	100	10
2	80	5
3	50	8
4	150	11
5	55	12
6	20	4
7	40	6
8	50	9
9	200	10
10	100	11

- 1. What is the optimal combination of boxes if you are seeking to maximize the total sales and are able to carry a maximum of 60 pounds?
- 2. What is the optimal combination of boxes if you are seeking to minimize the weight of the transported boxes bearing in mind that the minimum total sales must be at least €600 ?

#### Model 1: maximize the total sales

The mathematical model can be formulated as follows:

```
100 \cdot x_1 + 80 \cdot x_2 + 50 \cdot x_3 + 150 \cdot x_4 + 55 \cdot x_5 + 20 \cdot x_6 + 40 \cdot x_7 + 50 \cdot x_8 + 200 \cdot x_9 + 100 \cdot x_{10} \rightarrow \max! s.t. 10 \cdot x_1 + 5 \cdot x_2 + 8 \cdot x_3 + 11 \cdot x_4 + 12 \cdot x_5 + 4 \cdot x_6 + 6 \cdot x_7 + 9 \cdot x_8 + 10 \cdot x_9 + 11 \cdot x_{10} \le 60 x_j \in \{0,1\} \quad ; j = 1(1)10
```

The basic data is saved in the CMPL file knapsack-data.cmpl:

```
parameters:
   boxes := 1(1)10;
   #weight of the boxes
   w[boxes] := (10,5,8,11,12,4,6,9,10,11);
   #price per box
   p[boxes] := (100,80,50,150,55,20,40,50,200,100);
   #max capacity
   maxWeight := 60;
```

```
#min sales
minSales := 600;
```

## A simple CMPL model knapsack-max-basic.cmpl can be formulated as follows:

```
include "knapsack-data.cmpl"
variables:
    x[boxes] : binary;
objectives:
    sales: p[]T * x[] ->max;
constraints:
    weight: w[]T * x[] <= maxWeight;</pre>
```

### CMPL command:

cmpl knapsack-max-basic.cmpl

#### Solution:

<pre>General status: Problem : SolverName</pre>	•				
Nr. of Solutions:	1				
Solution nr.:	1				
Objective name :	sales				
Objective value :	700 (max!)				
Variables					
Nr. of variables:	10				
Name	Type	Activity	Lower Bound	Upper Bound	Marginal
x1	В	1.00	0.00	1.00	
x2	В	1.00	0.00	1.00	=
x3	В	0.00	0.00	1.00	_
x4	В	1.00	0.00	1.00	=
х5	В	0.00	0.00	1.00	=
x6	В	1.00	0.00	1.00	-
x7	В	0.00	0.00	1.00	-
x8	В	1.00	0.00	1.00	=
x9	В	1.00	0.00	1.00	=
x10	В	1.00	0.00	1.00	-
Constraints					
Nr. of constraints					
Name	Type	Activity	Lower Bound	Upper Bound	Marginal
weight	L	60 00	-Infinity	60.00	

## Model 2: minimize the weight

The mathematical model can be formulated as follows:

$$\begin{array}{l} 10 \cdot x_1 + 5 \cdot x_2 + 8 \cdot x_3 + 11 \cdot x_4 + 12 \cdot x_5 + 4 \cdot x_6 + 6 \cdot x_7 + 9 \cdot x_8 + 10 \cdot x_9 + 11 \cdot x_{10} \rightarrow \min! \\ s.t. \\ 100 \cdot x_1 + 80 \cdot x_2 + 50 \cdot x_3 + 150 \cdot x_4 + 55 \cdot x_5 + 20 \cdot x_6 + 40 \cdot x_7 + 50 \cdot x_8 + 200 \cdot x_9 + 100 \cdot x_{10} \\ \geq 600 \\ x_j \in \{0,1\} \quad ; j = 1(1)10 \end{array}$$

## A simple CMPL model knapsack-min-basic.cmpl can be formulated as follows:

```
include "knapsack-data.cmpl"
variables:
    x[boxes] : binary;
objectives:
    weight: w[]T * x[] ->min;
constraints:
    sales: p[]T * x[] >= minSales;
```

#### CMPL command:

```
cmpl knapsack-min-basic.cmpl
```

#### Solution:

SolverName Nr. of Solutions:					
Solution nr.: Objective name: Objective value:	1 weight				
Variables					
Nr. of variables: Name	Туре	=		Upper Bound	-
x1	В	1.00		1.00	
x2	В	1.00	0.00	1.00	=
x3	В	0.00	0.00	1.00	=
x4	В	1.00	0.00	1.00	=
x5	В	0.00	0.00	1.00	=
хб	В	0.00	0.00	1.00	=
x7	В	0.00	0.00	1.00	=
x8	В	0.00	0.00	1.00	=
x9	В	1.00	0.00	1.00	-
x10	В	1.00	0.00	1.00	- 
Constraints					
Nr. of constraints					
Name	Type	Activity		Upper Bound	
sales	G	630.00	600.00	Infinity	

# 9.1.5 Quadratic assignment problem

Assignment problems are special types of linear programming problems which assign assignees to tasks or locations. The goal of this quadratic assignment problem is to find the cheapest assignments of n machines to n locations. The transport costs are influenced by

- the distance  $d_{jk}$  between location j and location k and
- the quantity  $t_{hi}$  between machine h and machine i, which is to be transported.

The assignment of a machine h to a location j can be formulated with the Boolean variables

$$x_{hj} = \begin{cases} 1 \text{ , if machine } h \text{ is assigned to location } j \\ 0 \text{ , if not} \end{cases}$$

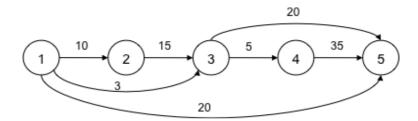
The general model can be formulated as follows:

$$\begin{split} &\sum_{h=1}^{n} \sum_{\substack{i=1\\i\neq h}}^{n} \sum_{j=1}^{n} \sum_{\substack{k=1\\i\neq j}}^{n} t_{hi} \cdot d_{jk} \cdot x_{hj} \cdot x_{ik} \to min! \\ &s.t. \\ &\sum_{j=1}^{n} x_{hj} = 1 \quad ; h = 1(1)n \\ &\sum_{h=1}^{n} x_{hj} = 1 \quad ; j = 1(1)n \\ &x_{hj} \in \{0,1\} \quad ; h = 1(1)n, j = 1(1)n \end{split}$$

Because of the product  $x_{ij} \cdot x_{ik}$  in the objective function the model is not a linear model. But it is possible to use a set of inequations to make an equivalent transformation of such multiplications of variables. This transformation is implemented in CMPL and the set of inequations will be generated automatically.

## Consider the following case:

There are 5 machines and 5 locations in the given factory. The quantities of goods which are to be transported between the machines are indicated in the figure below.



The distances between the locations are given in the following table:

from/to	1	2	3	4	5
1	М	1	2	3	4
2	1	М	1	2	3
3	2	1	М	1	2
4	3	2	1	М	1
5	4	3	2	1	М

The CMPL model quadratic-assignment.cmpl can be formulated as follows:

```
t[,] := ( (0, 10, 10, 0, 20),
             (0, 0, 15, 0, 0),
             (0, 0, 0, 5, 20),
             (0, 0, 0, 0, 35),
             (0, 0, 0, 0, 0);
variables:
     x[1..n,1..n]: binary;
     \#dummy variables to store the products x hj * x ik
     w[1..n,1..n,1..n]: real[0..1];
objectives:
     costs: sum\{ h:=1(1)n, i:=1(1)n, j:=1(1)n, k:=1(1)n :
                             t[h,i]*d[j,k]*w[h,j,i,k] } ->min;
constraints:
     { h:=1(1)n, i:=1(1)n, j:=1(1)n, k:=1(1)n:
                 { t[h,i] = 0: w[h,j,i,k] = 0; |
                   \# definition of the products x hj * x ik
                   default: w[h,j,i,k] = x[h,j] * x[i,k];
     { h:=1(1)n: sos1_{h}: sum{ <math>j:=1(1)n: x[h,j] } = 1; }
     { j:=1(1)n: sos2 $j$: sum{ h:=1(1)n: x[h,j] } = 1; }
```

#### CMPL command:

```
cmpl quadratic-assignment.cmpl -solver glpk
```

#### Solution:

```
General status: normal
Problem: quadratic-assignment.cmpl
SolverName
              COIN-OR glpk
Nr. of Solutions: 1
Solution nr.:
Objective name : costs
Objective value: 155 (min!)
Variables
Nr. of variables: 375
                                                          Upper Bound
                            Activity
                                          Lower Bound
Name
                Type
                                                                            Marginal
x1_1
                                0.00
                                                0.00
                                                                1.00
                 В
                                               0.00
x1_2
                               0.00
                                                                1.00
x1_3
                 В
                                0.00
                                                0.00
                                                                1.00
                                                0.00
x1_4
                  В
                                1.00
                                                                 1.00
                                               0.00
                                                                1.00
x1_5
                  В
                                0.00
x2 1
                               0.00
                                               0.00
                  В
                                                                1.00
x2_2
                               0.00
                                               0.00
                                                                1.00
x2_3
                  В
                                0.00
                                                0.00
                                                                1.00
x2_4
                   В
                                0.00
                                                 0.00
                                                                 1.00
```

x2 5					
x2_3	В	1.00	0.00	1.00	-
x3_1	В	0.00	0.00	1.00	-
k3_2	В	0.00	0.00	1.00	-
k3_3	В	1.00	0.00	1.00	-
k3_4	В	0.00	0.00	1.00	-
x3_5	В	0.00	0.00	1.00	-
<4_1	В	1.00	0.00	1.00	-
4_2	В	0.00	0.00	1.00	-
4_3	В	0.00	0.00	1.00	-
4_4	В	0.00	0.00	1.00	-
4_5	В	0.00	0.00	1.00	-
5_1	В	0.00	0.00	1.00	-
x5_2	В	1.00	0.00	1.00	-
F 2	В	0.00	0.00	1.00	-
32_3	D				
_	В	0.00	0.00	1.00	-
×5_4			0.00	1.00	-
x5_3 x5_4 x5_5 Constraints Nr. of constrai	В В	0.00			- - 
x5_4 x5_5 Constraints Nr. of constrai	В В	0.00	0.00		Marginal
x5_4 x5_5 Constraints Nr. of constrai	B B 	0.00	0.00	1.00	Marginal
x5_4 x5_5  Constraints Wr. of constrai	B B 	0.00	0.00	1.00	Marginal
x5_4 x5_5 Constraints Nr. of constrai Name sos1_1	B B nts: 710 Type	0.00 0.00	0.00	1.00 Upper Bound	Marginal
t5_4 t5_5 Constraints Wr. of constrai Name sos1_1 sos1_2	B B nts: 710 Type	0.00 0.00 Activity	0.00  Lower Bound  1.00	1.00 Upper Bound	Marginal 
x5_4 x5_5 	B B nts: 710 Type E E	0.00 0.00 Activity	0.00 Lower Bound 1.00 1.00	1.00 Upper Bound 1.00 1.00	Marginal 
x5_4 x5_4 x5_5  Constraints Nr. of constrai Name sos1_1 sos1_2 sos1_3	B B nts: 710 Type E E E	0.00 0.00 Activity 1.00 1.00	0.00 Lower Bound 1.00 1.00	1.00 Upper Bound 1.00 1.00 1.00	Marginal 
x5_4 x5_5 Constraints Nr. of constrai Name sos1_1 sos1_2 sos1_3 sos1_4	B B nts: 710 Type E E E E	0.00 0.00 Activity 1.00 1.00 1.00	1.00 1.00 1.00 1.00	1.00 Upper Bound 1.00 1.00 1.00	Marginal 
2.5_4 2.5_5 Constraints  Ir. of constrai  Iame 2.051_1 2.051_2 2.051_3 2.051_4 2.051_5 2.052_1	B B nts: 710 Type E E E E	0.00 0.00 Activity  1.00 1.00 1.00 1.00 1.00	1.00 1.00 1.00 1.00 1.00	1.00 Upper Bound 1.00 1.00 1.00 1.00	Marginal 
Constraints Ir. of constrai Iame cos1_1 cos1_2 cos1_3 cos1_4 cos1_5 cos2_1 cos2_2	B B nts: 710 Type E E E E E	0.00 0.00 1.00 1.00 1.00 1.00	1.00 1.00 1.00 1.00 1.00	1.00 Upper Bound 1.00 1.00 1.00 1.00 1.00 1.00 1.00	Marginal 
Constraints Wr. of constrai Name sos1_1 sos1_2 sos1_3 sos1_4 sos1_5	B B nts: 710 Type E E E E E	0.00 0.00 1.00 1.00 1.00 1.00 1.00	1.00 1.00 1.00 1.00 1.00 1.00	1.00  Upper Bound  1.00  1.00  1.00  1.00  1.00  1.00  1.00  1.00  1.00	Marginal 

The optimal assignments of machines to locations are given in the table below:

		location	IS			
		1	2	3	4	5
	1				Х	
es	2					X
machines	3			X		
ma	4	Х				
	5		X			

# 9.2 Using CMPL as a pre-solver

CMPL is not only intended to generate models in the MPS or OSIL format. CMPL can also be used as a presolver or simple solver. In this way it is possible to find a preliminary solution of a problem as a basis for the model which is to be generated.

## 9.2.1 Solving the knapsack problem

The knapsack problem is a very simple problem that does not necessarily have to be solved by an MIP solver. CMPL can be used as a simple solver for knapsack problems to approximate the optimal solution.

The idea of the following models is to evaluate each item using the relation between the value per item and weight per item. The knapsack will be filled with the items sorted in descending order until the capacity limit or the minimum value is reached.

Using the data from the examples in section 9.1.4 a CMPL model to maximize the total sales relative to capacity can be formulated as follows.

#### Model 1: maximize the total sales knapsack-max-presolved.cmpl

```
include "knapsack-data.cmpl"
#calculating the relative value of each box
{j in boxes: val[j]:= p[j]/w[j]; }
sumSales:=0;
sumWeight:=0;
#initial solution
x[] := (0,0,0,0,0,0,0,0,0,0);
{ i in boxes:
     maxVal:=max(val[]);
      {j in boxes:
            { maxVal=val[j] :
                        { sumWeight+w[j] <= maxWeight:
                                    x[j] := 1;
                                    sumSales:=sumSales + p[j];
                                    sumWeight:=sumWeight + w[j];
                        }
                        val[j]:=0;
                        break j;
            }
      }
echo "Solution found";
echo "Optimal total sales: "+ sumSales;
echo "Total weight: " + sumWeight;
{j in boxes: echo "x "+ j + ": " + x[j]; }
```

## CMPL command:

```
cmpl knapsack-max-presolved.cmpl -noOutput -cd
```

### Solution:

```
Solution found Optimal total sales: 690
```

```
Total weight: 57
x_1: 1
x_2: 1
x_3: 0
x_4: 1
x_5: 0
x_6: 1
x_7: 1
x_8: 0
x_9: 1
x_10: 1
```

This solution is not identical to the optimal solution on page 52 but good enough as an approximate solution.

## Model 2: minimize the total weight knapsack-min-presolved.cmpl

```
include "knapsack-data.cmpl"
#calculating the relative value of each box
{j in boxes: val[j]:= w[j]/p[j]; }
M:=10000;
sumSales:=0;
sumWeight:=0;
#initial solution
x[] := (0,0,0,0,0,0,0,0,0,0);
{sumSales < minSales:
     maxVal:=min(val[]);
      {j in boxes:
            { maxVal=val[j] :
                        { sumSales < minSales:
                                    x[j] := 1;
                                    sumSales:=sumSales + p[j];
                                    sumWeight:=sumWeight + w[j];
                        val[j]:=M;
                        break j;
            }
      repeat;
echo "Solution found";
echo "Optimal total weight: " + sumWeight;
echo "Total sales: "+ sumSales;
{j in boxes: echo "x "+ j + ": " + x[j]; }
```

## CMPL command:

```
cmpl knapsack-min-presolved.cmpl -noOutput -cd
```

#### Solution:

```
Optimal total weight: 47
Total sales: 630
x_1: 1
x_2: 1
x_3: 0
x_4: 1
x_5: 0
x_6: 0
x_7: 0
x_8: 0
x_9: 1
x_10: 1
```

This solution is identical to the optimal solution in section 9.1.4.

# 9.2.2 Finding the maximum of a negative convex function with the golden ratio method

One of the alternative methods for finding the maximum of a negative convex function is the golden ratio method. (see e.g. http://math.fullerton.edu/mathews/n2003/GoldenRatioSearchMod.html)

A CMPL program to find the maximum of  $f(x)=-0.5 \cdot x^2 - e^{-x}$  can be formulated as follows (max-negative-convex-function.cmpl):

```
parameters:
      #golden ratio delta
      d:=0.382;
      #distance epsilon
      e:=0.0001;
      #initial solution
      a:= 0;
      b := 1;
      1:= a + d*(b-a);
      m := a + (1-d) * (b-a);
      Fl:= -0.5 * 1^2 - exp(-1);
      Fm := -0.5 * m^2 - exp(-m);
      \{ (b-a) >= e :
             { Fl<Fm:</pre>
                   a:=1;
                   1:=m;
                   m:=a+(1-d)*(b-a);
                   F1:=Fm;
                   Fm:=-0.5 * m^2 - exp(-m);
             { F1>=Fm:
```

```
b:=m;
m:=1;
l:=a+d*(b-a);
Fm:=F1;
F1:= -0.5 * 1^2 - exp(-1);
}
repeat;
}
echo "Optimal solution found";
x:=round(b*1000)/1000;
echo "x: "+ format("%2.31f",x);
echo "function value: " + (-0.5 * x^2 - exp(-x));
```

#### CMPL command:

```
cmpl knapsack-min-presolved.cmpl -noOutput -cd
```

#### Solution:

```
Optimal solution found x: 0.562 function value: -0.727990
```

# 9.3 Several selected CMPL applications

## 9.3.1 Calculating the Fibonacci sequence

By definition, the first Fibonacci sequence starts with the numbers 0 and 1, and each remaining number is the sum of the previous two.

```
a_{n+1} = a_n + a_{n-1}; a_1 = 0, a_2 = 1, n \in \mathbb{N}
```

CMPL code to calculate the Fibonacci sequence (fibonacci.cmpl):

```
parameters:
    # initializing the first elements
    F[1..2] := (0, 1);
    # Calculating the Fibonacci sequence until the 10th element
    {i:=3(1)10: F[i] := F[i-2] + F[i-1]; }
    echo "The Fibonacci sequence for the first 10 elements";
    {i:=1(1)10: echo "element " + i + ": " + F[i]; }
```

#### CMPL command:

```
cmpl fibonacci.cmpl -noOutput
```

#### Calculated sequence:

```
The Fibonacci sequence for the first 10 elements
element 1: 0
element 2: 1
element 3: 1
element 4: 2
element 5: 3
element 6: 5
element 7: 8
element 8: 13
element 9: 21
element 10: 34
```

# 9.3.2 Calculating primes

A prime is defined as a natural number that has exactly two distinct natural number divisors: 1 and itself.

CMPL code to calculate the sequence of primes (primes.cmpl):

#### CMPL command:

```
cmpl primes.cmpl -noOutput
```

#### Calculated sequence:

```
The prime sequence in the range 3 until 10 element 1: 2 element 2: 3 element 3: 5 element 4: 7
```

# 10 Authors and Contact

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