Week 5 - Notes

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2021-22

1 Harlow Shows that $RT \implies Complementary Recovery$

• The first thing to consider is a small variation in the entropy:

$$\delta S(\tilde{\rho}, M) = S(\tilde{\rho} + \delta \tilde{\rho}, M) - S(\tilde{\rho}, M)
= -\hat{\text{Tr}} \left(\left(\hat{\rho}_M + \delta \hat{\rho}_M \right) \log \left(\hat{\rho}_M + \delta \hat{\rho}_M \right) \right) + \hat{\text{Tr}} \left(\hat{\rho}_M \log \hat{\rho}_M \right)
= -\hat{\text{Tr}} \left(\hat{\rho}_M \left(\log \hat{\rho}_M + \hat{\rho}_M^{-1} \delta \hat{\rho}_M \right) \right) - \hat{\text{Tr}} \left(\delta \hat{\rho}_M \log \hat{\rho}_M \right) + \hat{\text{Tr}} \left(\hat{\rho}_M \log \hat{\rho}_M \right) + \mathcal{O} \left(\delta \hat{\rho}_M^2 \right)
= -\hat{\text{Tr}} \left(\delta \hat{\rho}_M \right) - \hat{\text{Tr}} \left(\delta \hat{\rho}_M \log \hat{\rho}_M \right) + \mathcal{O} \left(\delta \hat{\rho}_M^2 \right)
= -\hat{\text{Tr}} \left(\delta \tilde{\rho} \log \left(\tilde{\rho}_A \otimes I_{\overline{A}} \right) \right) + \mathcal{O} \left(\delta \tilde{\rho}^2 \right)
= -\hat{\text{Tr}} \left(\delta \tilde{\rho} \log \tilde{\rho}_A \right) + \mathcal{O} \left(\delta \tilde{\rho}^2 \right)$$

$$(1.1)$$

using that $Tr(\delta \tilde{\rho}) = 0$ and the standard log expansion

• We can play the same game more generally if we take the variation $\delta \tilde{\rho}$ about a state σ :

$$S(\tilde{\sigma} + \delta \tilde{\rho}, M) - S(\tilde{\sigma}, M) = -\text{Tr}\left(\delta \tilde{\rho} \log \tilde{\sigma}_A\right)$$
(1.2)

A similar calculation shows

$$\delta S(\tilde{\rho}) = -\text{Tr}(\delta \tilde{\rho} \log \tilde{\rho}) \implies S(\tilde{\sigma} + \delta \tilde{\rho}) - S(\tilde{\sigma}) = -\text{Tr}(\delta \tilde{\rho} \log \tilde{\sigma})$$
 (1.3)

• Recall the RT formula on A:

$$S(\tilde{\rho}_A) = \text{Tr}\tilde{\rho}\mathcal{L}_A + S(\tilde{\rho}, M) \tag{1.4}$$

• So, taking the variation $\delta \tilde{\rho}$ on $\tilde{\sigma}$, we get

$$S(\tilde{\sigma}_A + \delta \tilde{\rho}_A) - S(\tilde{\sigma}_A) = \text{Tr} \left((\tilde{\sigma} + \delta \tilde{\rho}) \mathcal{L}_A \right) - \text{Tr} \left(\tilde{\sigma} \mathcal{L}_A \right) + S(\tilde{\sigma} + \delta \tilde{\rho}, M) - S(\tilde{\sigma}, M)$$

$$\implies -\text{Tr} \left(\delta \tilde{\rho}_A \log \tilde{\sigma}_A \right) = \text{Tr} \left(\delta \tilde{\rho} \left(\mathcal{L}_A - \log \tilde{\sigma} \right) \right)$$
(1.5)

to linear order

• Since both sides are linear in $\delta \tilde{\rho}$, we can integrate over all such variations on both sides, giving us

$$-\operatorname{Tr}\left(\tilde{\rho}_{A}\log\tilde{\sigma}_{A}\right) = \operatorname{Tr}\left(\tilde{\rho}\left(\mathcal{L}_{A} - \log\tilde{\sigma}_{A}\right)\right) \tag{1.6}$$

• We can now calculate relative entropies as:

$$S(\tilde{\rho}_{A}|\tilde{\sigma}_{A}) = \operatorname{Tr}(\tilde{\rho}_{A}\log\tilde{\rho}_{A}) - \operatorname{Tr}(\tilde{\rho}_{A}\log\tilde{\sigma}_{A})$$

$$= -S(\tilde{\rho}_{A}) + \operatorname{Tr}(\tilde{\rho}(\mathcal{L}_{A} - \log\tilde{\sigma}_{A}))$$

$$= -S(\tilde{\rho}_{A}) + S(\tilde{\rho}_{A}) - S(\tilde{\rho}, M) - \operatorname{Tr}(\tilde{\rho}\log\tilde{\sigma}_{A})$$

$$= -S(\tilde{\rho}, M) - \operatorname{Tr}(\tilde{\rho}_{M}\log\tilde{\sigma}_{M})$$

$$= S(\tilde{\rho}|\tilde{\sigma}, M)$$
(1.7)

• A similar argument on \overline{A} shows that

$$S(\tilde{\rho}_{\overline{A}}|\tilde{\sigma}_{\overline{A}}) = S(\tilde{\rho}|\tilde{\sigma}, M') \tag{1.8}$$

- These two between them imply the second condition of Harlow 5.1
- Consider a state $|\tilde{\psi}\rangle \in \mathcal{H}_{\text{code}}$, an operator $X_{\overline{A}}$ on $\mathcal{H}_{\overline{A}}$, and an operator $\tilde{O} \in M$
- Since M is spanned by its Hermitian elements, we take \tilde{O} to be Hermitian
- We now consider

$$\langle \tilde{\psi} | e^{-i\lambda \tilde{O}} X_{\overline{A}} e^{i\lambda \tilde{O}} | \tilde{\psi} \rangle = \langle \tilde{\psi} | e^{-i\lambda \tilde{O}} P_{\text{code}} X_{\overline{A}} P_{\text{code}} e^{i\lambda \tilde{O}} | \tilde{\psi} \rangle$$
(1.9)

since $\tilde{O} |\tilde{\psi}\rangle = P_{\text{code}} \tilde{O} P_{\text{code}} P_{\text{code}} |\tilde{\psi}\rangle = P_{\text{code}} \tilde{O} |\tilde{\psi}\rangle$

• Consider the states

$$|\tilde{\psi}(\lambda)\rangle \equiv e^{i\lambda\tilde{O}}|\tilde{\psi}\rangle \tag{1.10}$$

and note that for any $\tilde{O}' \in M'$, the expectation $\langle \tilde{\psi}(\lambda) | \tilde{O}' | \tilde{\psi}(\lambda) \rangle$ is independent of λ

- Recalling that for any state ρ , there is a unique state $\rho_{M'} \in M'$ such that $\mathbb{E}_{\rho}(x') = \mathbb{E}_{\rho_{M'}}(x')$ for any $x' \in M'$, this means that the states $\tilde{\psi}(\lambda)_{M'}$ corresponding to $\tilde{\psi}(\lambda) = |\tilde{\psi}(\lambda)\rangle \langle \tilde{\psi}(\lambda)|$ on M' are independent of λ too
- So, for any two $\lambda, \lambda', \tilde{\psi}(\lambda)_{M'} = \tilde{\psi}(\lambda')_{M'}$, which means

$$S(\tilde{\psi}(\lambda)|\tilde{\psi}(\lambda'), M') = 0 \tag{1.11}$$

• So from (1.8):

$$0 = S(\tilde{\psi}(\lambda)|\tilde{\psi}(\lambda'), M') = S(\tilde{\psi}(\lambda)_{\overline{A}}|\tilde{\psi}(\lambda')_{\overline{A}})$$
(1.12)

and so $\operatorname{Tr}_A\left(|\tilde{\psi}(\lambda)\rangle\langle\tilde{\psi}(\lambda)|\right)$ is independent of λ too

- This can only hold if $|\tilde{\psi}(\lambda)\rangle$ itself is independent of λ
- Now return to (1.9); since this is therefore independent of λ , its first variation with respect to λ must vanish
- Therefore:

$$\delta \langle \tilde{\psi} | e^{-i\lambda \tilde{O}} P_{\text{code}} X_{\overline{A}} P_{\text{code}} e^{i\lambda \tilde{O}} | \tilde{\psi} \rangle \propto \langle \tilde{\psi} | [P_{\text{code}} X_{\overline{A}} P_{\text{code}}, \tilde{O}] | \tilde{\psi} \rangle = 0$$
 (1.13)

which just implies $P_{\text{code}}X_{\overline{A}}P_{\text{code}}=X'P_{\text{code}}$ for some $X'\in M'$, which is just the second condition as claimed

• Since we can repeat this argument with $M \leftrightarrow M'$ and $A \leftrightarrow \overline{A}$, we get that a two sided RT formula implies complementary recovery

2 Pollack 4.8

• We jump straight in with Pollack's theorem 4.8

Theorem 2.1. Consider an encoding isometry V, a subregion A, and a von Neumann algebra M so that (V, A, M) have complementary recovery. Then (V, A, M) and (V, \overline{A}, M') both have an RT formula with the same area operator L, and L is in the center Z_M .

Proof. Suppose M induces decomposition $\mathcal{H}_L = \bigoplus_{\alpha} (\mathcal{H}_{L_{\alpha}} \otimes \mathcal{H}_{\overline{L}_{\alpha}})$. Then, we can decompose M and M' together as

$$M = \bigoplus_{\alpha} (\mathcal{L}(\mathcal{H}_{L_{\alpha}}) \otimes I_{\overline{L}_{\alpha}}), \quad M' = \bigoplus_{\alpha} (I_{L_{\alpha}} \otimes \mathcal{L}(\mathcal{H}_{\overline{L}_{\alpha}}))$$
 (2.1)

Define $|\alpha, i, j\rangle = |i_{\alpha}\rangle_{L_{\alpha}} \otimes |j_{\alpha}\rangle_{\overline{L}_{\alpha}}$ as the compatible basis for \mathcal{H}_{L} . By earlier lemmas, we know there exist factorisations of the form

$$\mathcal{H}_{A} = \bigoplus_{\alpha} (\mathcal{H}_{A_{1}^{\alpha}} \otimes \mathcal{H}_{A_{2}^{\alpha}}) \oplus \mathcal{H}_{A_{3}}, \quad \mathcal{H}_{\overline{A}} = \bigoplus_{\alpha} (\mathcal{H}_{\overline{A}_{1}^{\alpha}} \otimes \mathcal{H}_{\overline{A}_{2}^{\alpha}}) \oplus \mathcal{H}_{\overline{A}_{3}}$$
(2.2)

and unitaries U_A and $U_{\overline{A}}$ such that

$$(U_{A} \otimes I_{\overline{A}})V |\alpha, i, j\rangle = |\psi_{\alpha, i}\rangle_{A_{1}^{\alpha}} \otimes |\chi_{\alpha, j}\rangle_{A_{2}^{\alpha}\overline{A}}$$

$$(I_{A} \otimes U_{\overline{A}})V |\alpha, i, j\rangle = |\overline{\chi}_{\alpha, i}\rangle_{A_{A}^{\alpha}} \otimes |\overline{\psi}_{\alpha, j}\rangle_{\overline{A}_{1}^{\alpha}}$$

$$(2.3)$$

If we apply $(U_A \otimes I_{\overline{A}})$ and $(I_A \otimes U_{\overline{A}})$ sequentially, we see that

$$(I_{A} \otimes U_{\overline{A}})(U_{A} \otimes I_{A})V |\alpha, i, j\rangle = |\psi_{\alpha, i}\rangle_{A_{1}^{\alpha}} \otimes (I_{A_{2}^{\alpha}} \otimes U_{\overline{A}}) |\chi_{\alpha, j}\rangle_{A_{2}^{\alpha}\overline{A}} (U_{A} \otimes I_{A})(I_{A} \otimes U_{\overline{A}})V |\alpha, i, j\rangle = (U_{A} \otimes I_{\overline{A}_{2}^{\alpha}}) |\overline{\chi}_{\alpha, i}\rangle_{A_{A}^{\alpha}} \otimes |\overline{\psi}_{\alpha, j}\rangle_{\overline{A}_{1}^{\alpha}}$$
(2.4)

In order for both of these to be true simultaneously, there therefore must exist states $|\chi_{\alpha}\rangle_{A_{2}^{\alpha}\overline{A}_{2}^{\alpha}}$ such that $(I_{A_{2}^{\alpha}}\otimes U_{\overline{A}})|\chi_{\alpha,j}\rangle_{A_{2}^{\alpha}\overline{A}}=|\chi_{\alpha}\rangle_{A_{2}^{\alpha}\overline{A}_{2}^{\alpha}}\otimes|\overline{\psi}_{\alpha,j}\rangle_{\overline{A}_{1}^{\alpha}}$, which implies

$$(U_A \otimes U_{\overline{A}})V |\alpha, i, j\rangle = |\psi_{\alpha, i}\rangle_{A_1^{\alpha}} \otimes |\chi_{\alpha}\rangle_{A_2^{\alpha}\overline{A}_2^{\alpha}} \otimes |\overline{\psi}_{\alpha, j}\rangle_{\overline{A}_1^{\alpha}}$$
(2.5)

This result will imply the RT formula when we consider a logical operator in this basis. Suppose ρ is a state with support on \mathcal{H}_L ; we will compute $S(M, \rho)$ and $S(\operatorname{Tr}_{\overline{A}}(V\rho V^{\dagger}))$ and take the difference.

Recall that instead of considering $S(\operatorname{Tr}_{\overline{A}}(V\rho V^{\dagger}))$ we might as well consider $S(\operatorname{Tr}_{\overline{A}}(V\rho_M V^{\dagger}))$: to see why, set $O_A \in \mathcal{L}(\mathcal{H}_A)$, and compute

$$\operatorname{Tr}(O_A \cdot \operatorname{Tr}_{\overline{A}}(V \rho V^{\dagger})) = \operatorname{Tr}((O_A \otimes I_{\overline{A}}) \cdot V \rho V^{\dagger}) = \operatorname{Tr}(V^{\dagger}(O_A \otimes I_{\overline{A}}) V \cdot \rho)$$
 (2.6)

by cyclicity of the trace. But $V^{\dagger}(O_A \otimes I_{\overline{A}})V \in M$, and for any $O \in M$ we have $\operatorname{Tr}(O\rho) = \operatorname{Tr}(O\rho_M)$, we can just consider ρ_M in the above. Since the states $\operatorname{Tr}_{\overline{A}}(V\rho V^{\dagger})$ and $\operatorname{Tr}_{\overline{A}}(V\rho_M V^{\dagger})$ for all observables in M, they are the same state and so have the same entropy. Acting with a unitary on \mathcal{H}_A and $\mathcal{H}_{\overline{A}}$ separately does not change the unitary, we have:

$$S(\operatorname{Tr}_{\overline{A}}(V\rho_M V^{\dagger})) = S(\operatorname{Tr}_{\overline{A}}((U_A \otimes U_{\overline{A}})V\rho_M V^{\dagger}(U_A \otimes U_{\overline{A}})^{\dagger}))$$
(2.7)

The next step is to define isometries $\tilde{V}_{\alpha}: (\mathcal{H}_{L_{\alpha}} \otimes \mathcal{H}_{\overline{L}_{\alpha}}) \to (\mathcal{H}_{A_{1}^{\alpha}} \otimes \mathcal{H}_{\overline{A}_{1}^{\alpha}})$ using the states $|\psi_{\alpha,i}\rangle_{A_{1}^{\alpha}}$ and $|\overline{\psi}_{\alpha,j}\rangle_{\overline{A}_{1}^{\alpha}}$:

$$\tilde{V}_{\alpha} |\alpha, i, j\rangle \equiv |\psi_{\alpha, i}\rangle_{A_{1}^{\alpha}} \otimes |\psi_{\alpha, j}\rangle_{\overline{A}_{1}^{\alpha}}$$
(2.8)

which is certainly an isometry since $|\psi_{\alpha,i}\rangle_{A_1^{\alpha}}$ and the other one are bases for $\mathcal{H}_{A_1^{\alpha}}$ and $\mathcal{H}_{\overline{A}_1^{\alpha}}$ respectively.

 \tilde{V}_{α} then lets us rewrite (2.5) as

$$(U_A \otimes U_{\overline{A}})V |\alpha, i, j\rangle = \tilde{V}_{\alpha} |\alpha, i, j\rangle \otimes |\chi_{\alpha}\rangle_{A_{\alpha}^{\alpha} \overline{A}_{\alpha}^{\alpha}}$$
(2.9)

which lets us further simplify $(U_A \otimes U_{\overline{A}})V \rho_M V^{\dagger} (U_A \otimes U_{\overline{A}})^{\dagger}$:

$$(U_{A} \otimes U_{\overline{A}})V \rho_{M} V^{\dagger} (U_{A} \otimes U_{\overline{A}})^{\dagger})$$

$$= \sum_{\alpha} p_{\alpha} \cdot (U_{A} \otimes U_{\overline{A}})V \rho_{A_{\alpha}} V^{\dagger} (U_{A} \otimes U_{\overline{A}})^{\dagger}$$

$$= \sum_{\alpha} p_{\alpha} \cdot \frac{1}{p_{\alpha}} \sum_{i,j} \sum_{i',j'} \rho[\alpha]_{i,j,i',j'} (U_{A} \otimes U_{\overline{A}})V |\alpha,i,j\rangle \langle \alpha,i',j'| V^{\dagger} (U_{A} \otimes U_{\overline{A}})^{\dagger}$$

$$= \sum_{\alpha} p_{\alpha} \cdot \frac{1}{p_{\alpha}} \sum_{i,j} \sum_{i',j'} \rho[\alpha]_{i,j,i',j'} \tilde{V}_{\alpha} |\alpha,i,j\rangle \langle \alpha,i',j'| \tilde{V}_{\alpha}^{\dagger} \otimes |\chi_{\alpha}\rangle \langle \chi_{\alpha}|$$

$$= \sum_{\alpha} \cdot \tilde{V}_{\alpha} \rho_{\alpha} \tilde{V}_{\alpha}^{\dagger} \otimes |\chi_{\alpha}\rangle \langle \chi_{\alpha}|$$

$$= \sum_{\alpha} \cdot \tilde{V}_{\alpha} \rho_{\alpha} \tilde{V}_{\alpha}^{\dagger} \otimes |\chi_{\alpha}\rangle \langle \chi_{\alpha}|$$

$$= \sum_{\alpha} \cdot \tilde{V}_{\alpha} \rho_{\alpha} \tilde{V}_{\alpha}^{\dagger} \otimes |\chi_{\alpha}\rangle \langle \chi_{\alpha}|$$

Each of the states $\tilde{V}_{\alpha}\rho_{\alpha}\tilde{V}_{\alpha}^{\dagger}\otimes|\chi_{\alpha}\rangle\langle\chi_{\alpha}|$ are normalised and act on different blocks, we can compute the entropy as

$$S(\operatorname{Tr}_{\overline{A}}(V\rho V^{\dagger})) = -\sum_{\alpha} p_{\alpha} \log p_{\alpha} + \sum_{\alpha} p_{\alpha} S(\operatorname{Tr}_{\overline{A}}(\tilde{V}_{\alpha}\rho_{\alpha}\tilde{V}_{\alpha}^{\dagger} \otimes |\chi_{\alpha}\rangle \langle \chi_{\alpha}|))$$

$$= \sum_{\alpha} p_{\alpha} \log(p_{\alpha}^{-1}) + \sum_{\alpha} p_{\alpha} S(\operatorname{Tr}_{\overline{A}}(\tilde{V}_{\alpha}\rho_{\alpha}\tilde{V}_{\alpha}^{\dagger})) + \sum_{\alpha} p_{\alpha} S(\operatorname{Tr}_{\overline{A}}(|\chi_{\alpha}\rangle \langle \chi_{\alpha}|))$$
(2.11)

Now, since $|\psi_{\alpha,j}\rangle$ is independent of i, we further have that

$$S(\operatorname{Tr}_{\overline{A}}(\tilde{V}_{\alpha}\rho_{\alpha}\tilde{V}_{\alpha}^{\dagger})) = S(\operatorname{Tr}_{\overline{A}_{1}^{\alpha}}(\tilde{V}_{\alpha}\rho_{\alpha}\tilde{V}_{\alpha}^{\dagger})) = S(\operatorname{Tr}_{\overline{L}_{\alpha}}(\rho_{\alpha}))$$
(2.12)

Now, the first two terms of (2.11) are just $S(\rho, M)$, so we have

$$S(\operatorname{Tr}_{\overline{A}}(V\rho V^{\dagger})) - S(\rho, M) = \sum_{\alpha} p_{\alpha} S(\operatorname{Tr}_{\overline{A}}(|\chi_{\alpha}\rangle \langle \chi_{\alpha}|))$$
 (2.13)

The RHS is linear in p_{α} , so must be linear in ρ also. This means there exists an area operator \mathcal{L} such that the RHS can be written as $\text{Tr}(\rho \mathcal{L})$. Explicitly, we define

$$I_{\alpha} \equiv \sum_{i,j} |\alpha, i, j\rangle \langle \alpha, i, j|$$

$$\mathcal{L} \equiv \sum_{\alpha} S(\operatorname{Tr}_{\overline{A}}(|\chi_{\alpha}\rangle \langle \chi_{\alpha}|)) \cdot I_{\alpha}$$
(2.14)

which gives $\mathcal{L} \in M$, and so $\text{Tr}(\rho \mathcal{L}) = \text{Tr}(\rho_M \mathcal{L})$, and we can write

$$\operatorname{Tr}(\rho_{M}\mathcal{L}) = \operatorname{Tr}\left(\sum_{\alpha} p_{\alpha}\rho_{\alpha} \cdot \sum_{\alpha} S(\operatorname{Tr}_{\overline{A}}(|\chi_{\alpha}\rangle\langle\chi_{\alpha}|)) \cdot I_{\alpha}\right)$$

$$= \sum_{\alpha} p_{\alpha}S(\operatorname{Tr}_{\overline{A}}(|\chi_{\alpha}\rangle\langle\chi_{\alpha}|)) \cdot \operatorname{Tr}(\rho_{\alpha}I_{\alpha})$$

$$= S(\operatorname{Tr}_{\overline{A}}(V\rho V^{\dagger})) - S(\rho, M)$$
(2.15)

so (V, A, M) satisfy an RT formula with area operator \mathcal{L} , and $\mathcal{L} \in M$. We can do the same derivation for (V, \overline{A}, M') with $i \leftrightarrow j$, and since $\mathcal{L} \in M'$, $\mathcal{L} \in Z_M$ as claimed.