

1. (a)

T		m	c	e	t	y		e
h	a	a	h	s			a	r
e	r	n	o		o	o	n	.
r	e	y	i	b	n	n	s	■
e			c	u	l	e	w	■

- (b) Number of shaded boxes = $\lceil n/k \rceil \times k - n$
Number of shaded boxes = $\lceil 43/5 \rceil \times 5 - 43 = (9 \times 5) - 43 = 2$
- (c) Plaintext = **There are many choices but only one answer.**

2. (a) 00010 10100 11001
x+x+x ++x+x xx++x
\-\|\ |-/-\ //--/

(b) \-\|\ |-/-\ //--/
x+x+x ++x+x xx++x
+xx+x x++++ xxx++
??\|\ ?-?-? //?-?
+x\|\ x-++ //x-+

Bob needs Alice to tell him which of the filter selections he made were correct.

(c) 00010 10100 11001
\-\|\ |-/-\ //--/
x+x+x ++x+x xx++x
+xx+x x++++ xxx++
??\|\ ?-?-? //?-?
010 0 0 11 0

final key: 01000110

- (d) Alice and Bob can check for Eve's eavesdropping by setting aside cw a number of bits of the key for verification.
- (e) No, because photons cannot be observed without changing their polarity. The speed of computation has nothing to do with it.

3. (a) $IC = (4 * 3 * 2 + 3 * 2 * 1 + 2 * 1 * 5) / (30 * 29) = 40/870 = 0.046$

(b) 3 is the more likely keyword length. Evidence:

- more of the repeated trigram sequences have offsets divisible by 3
- the subsequences induced by key length 3 have higher avg IC

(c) BGH probably occurs by accident, since $91 = 7 \times 13$

None of the other repeated sequences have offsets divisible by 7, and only one, GHH, by 13.

(d) keyword: DNA

(e) **the novel feature of the**

4. (a) The cipher must be of a type in which the order of encryption and decryption operations does not matter. Such ciphers include Caesar, Vigenere, and one-time pad. It would not work with a substitution cipher which is “last on, first off”.
- (b) $C1 = VWDDSD$ $C1 = m + a$ $b = C2 - C1 \pmod{26}$
 $C2 = WAWWW$ $C2 = m + a + b$ $a = C2 - C3 \pmod{26}$
 $C3 = LMLTIL$ $C3 = m + b$ $m = C1 + C3 - C2 \pmod{26}$
- Alice's secret key $a = LOL$
Bob's secret key $b = BET$
Alice's message $m = KISSES$
- (c) Since Alice is using Vigenere, C1 can be cracked using techniques like Kasiski Examination and IC/IMC.
- (d) *This question was withdrawn.*

5. (a) Public key: $(n = 15; e)$
The possibilities are $d = e = 3$ or 5 or 7 .
- (b) Private key: $(n = 15; d)$
The possibilities are $d = e = 3$ or 5 or 7 .
- (c) B: $M = 1$ $C = 1^3 \pmod{15} = 1$
A: $M = 0$ $C = 0^3 \pmod{15} = 0$
C: $M = 2$ $C = 2^3 \pmod{15} = 8$
K: $M = 10$ $C = 10^3 \pmod{15} = 10$
 $[1, 0, 8, 10]$
- (d) $M = C^d \pmod{n} = 12^3 \pmod{15} = 1728 \pmod{15} = 3 \Rightarrow 'D'$
 $M = 12^3 \pmod{15} = 3 \Rightarrow 'D'$
 $M = 12^5 \pmod{15} = 12 \Rightarrow 'M'$
 $M = 12^7 \pmod{15} = 3 \Rightarrow 'D'$
- (e)
 - The numbers are too small.
 - The public and private keys are identical.
 - The block size of 1 makes it a glorified substitution cipher.

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6. def ngramScore(decipherment, englishText, n):
    score = 0
    num_dictionary_ngrams = len(englishText) - n + 1
    # i is the start index of the n-gram
    for i in range(len(decipherment)-n+1):
        # Get the current n-gram
        gram = decipherment[i:i+n]
        # Get the frequency of the n-gram in the englishText text
        gram_freq = englishText.count(gram) / num_dictionary_ngrams
        # Update the score
        score += gram_freq
    return score
```

7. (a) Organize the words that share inflectional patterns:

	Word A	Word B
Case 1	$\perp \Delta \sqcap$ (i-de-mu)	$\bowtie \uplus \exists \sqcap$ (do-bi-ye-mu)
Case 2	$\perp \Delta \mp$ (i-de-te)	$\bowtie \uplus \exists \mp$ (do-bi-ye-te)
Case 3	$\perp \bowtie$ (i-do)	$\bowtie \uplus \in$ (do-bi-yo)

(b) Which symbols represent the bridging syllables?

$$\Delta (de) \bowtie (do) \exists (ye) \in (yo)$$

(c) Organize the symbols into the grid below.

	Vowel 1	Vowel 2
Consonant 1	$\Delta (de)$	$\bowtie (do)$
Consonant 2	$\exists (ye)$	$\in (yo)$

(d) Suppose that the following words have been correctly deciphered:

$$\bowtie \exists = \text{doye}$$

$$\uplus \mp = \text{bite}$$

$$\odot \uplus \sqcap = \text{lubimu}$$

Reconstruct the case endings for all three cases in the table below.

Case 1	-emu
Case 2	-ete
Case 3	-o