SLE 66R35R / SLE 66R35I

Intelligent 1kByte Memory chip with NRG

Extended datasheet

Devices

- SLE 66R35R
- SLE 66R35I

Features

- Intelligent 1 kByte Memory Chip with NRG (ISO/IEC 14443-3 type A with CRYPTO1)
- Physical Interface and Anticollision compliant to ISO/IEC14443-2 and -3 Type A
    - Operation frequency 13.56 MHz
    - Data rate 106 kbit/s
- 1 kByte EEPROM
    - Block organization of memory, 16 Sectors with fixed 4 blocks of 16 bytes each
    - User definable access conditions for each memory block
- Security Features
    - SLE 66R35R: 4-byte reused identification number (r-ID), 4-byte UID previously issued
    - SLE 66R35I: 4-byte fixed non-unique number (FNUID)
    - Mutual three-pass authentication between card and reader for basic security
    - Selective memory access control secured by authentication and access conditions
    - Data encryption for RF channel
    - Dedicated Value Counter

About this document

Scope and purpose

This document describes the features, functionality and operational characteristics of SLE 66R35R/I.

Intended audience

This document is primarily intended for system and application developers.

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Features

Intelligent 1 kByte Memory Chip with NRG

Contactless Interface

- Physical Interface and Anticollision compliant to ISO/IEC14443-2 and -3 Type A
  - Operation frequency 13.56 MHz; data rate 106 kbit/s
  - Contactless transmission of data and supply energy
  - Anticollision logic: several cards may be operated in the field simultaneously
- Read and Write Distance up to 10 cm and more (influenced by external circuitry i.e. reader and inlay design)
- Short transaction times: typical ticketing transaction < 100 ms; transaction possible when card is moving

1 kByte EEPROM

- Block organization of memory, 16 Sectors with fixed 4 blocks of 16 bytes each
- EEPROM updating time per block < 5 ms
- Endurance > 100.000 erase/write cycles[1]
- Data Retention > 10 years[1]
- User definable access conditions for each memory block

Security Features

- SLE 66R35R: 4-byte reused identification number (r-ID), 4-byte UID previously issued
- SLE 66R35I: 4-byte fixed non-unique number (FNUID)
- Mutual three-pass authentication between card and reader for basic security
  - 48-bit key length
  - 2 keys per sector enabling key management
  - Transport key at chip delivery
- Selective memory access control secured by authentication and access conditions
- Suited to multifunctional applications: Individual key sets are available for each EEPROM sector
- Data encryption for RF channel
- Dedicated Value Counter
- Data integrity supported by CRC, Parity Check, etc.

Electrical characteristics

- On-chip capacitance 18.3 pF + 10 %
- ESD protection typical 2 kV
- Ambient temperature -25 ... +70°C for the chip

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[1] Values are temperature dependant

1 Ordering and packaging information

Table 1   Ordering information

Type	Package	Remark
SLE 66R35I C	Die (on wafer)	sawn / unsawn
SLE 66R35I NB	Die (on wafer)	NiAu-bumps, sawn
SLE 66R35I MCC8	MCC8-2-6	
SLE 66R35I MCC2	MCC2-2-1	
SLE 66R35R C	Die (on wafer)	sawn / unsawn
SLE 66R35R NB	Die (on wafer)	NiAu-bumps, sawn
SLE 66R35R MCC8	MCC8-2-6	
SLE 66R35R MCC2	MCC2-2-1	

Note: *For further information on technology, delivery forms (wafer thickness or height of NiAu-bump) please contact your local Infineon Technologies sales representative (www.infineon.com).*

Figure 1   Pin configuration Module Contactless Card – MCC2 (top view)
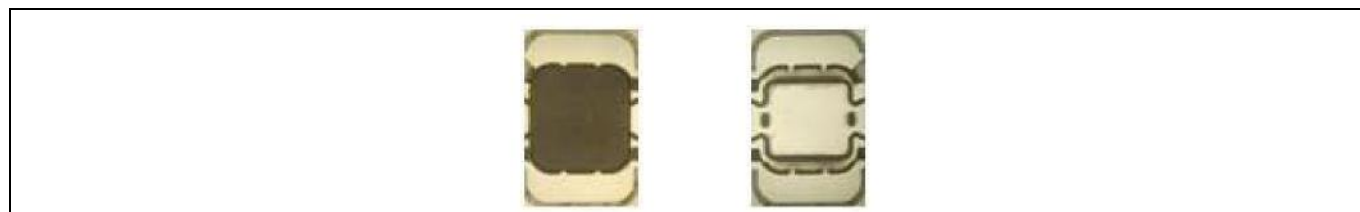
Figure 2   Pin configuration Module Contactless Card – MCC8-2-6 (top / bottom view)
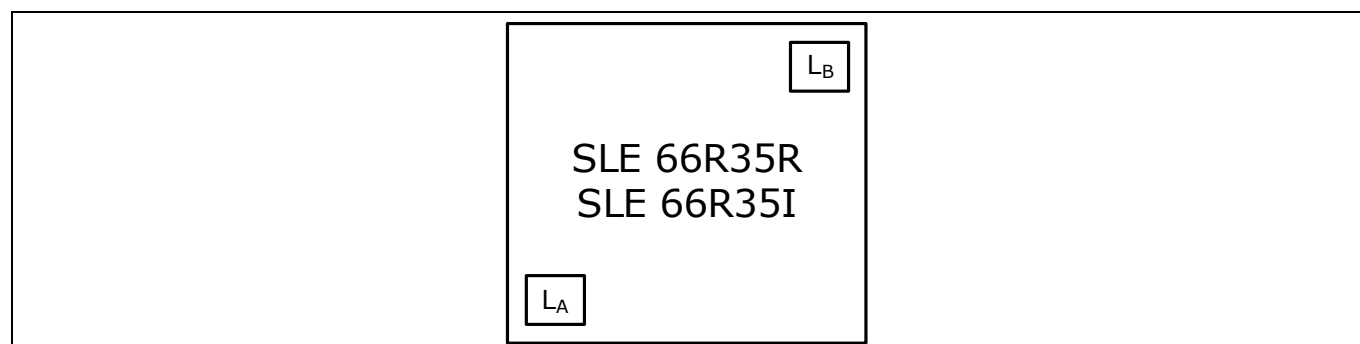
Figure 3   Pad configuration die

Table 2   Pin description and function

Symbol	Function
L_A	Antenna Connection
L_B	Antenna Connection

2 Overview of a NRG system

The SLE 66R35R/I is designed to operate in a NRG system. The system consists of a smart card and a card reader together with an antenna. The card's antenna consists of a simple coil with a few turns embedded in plastic.
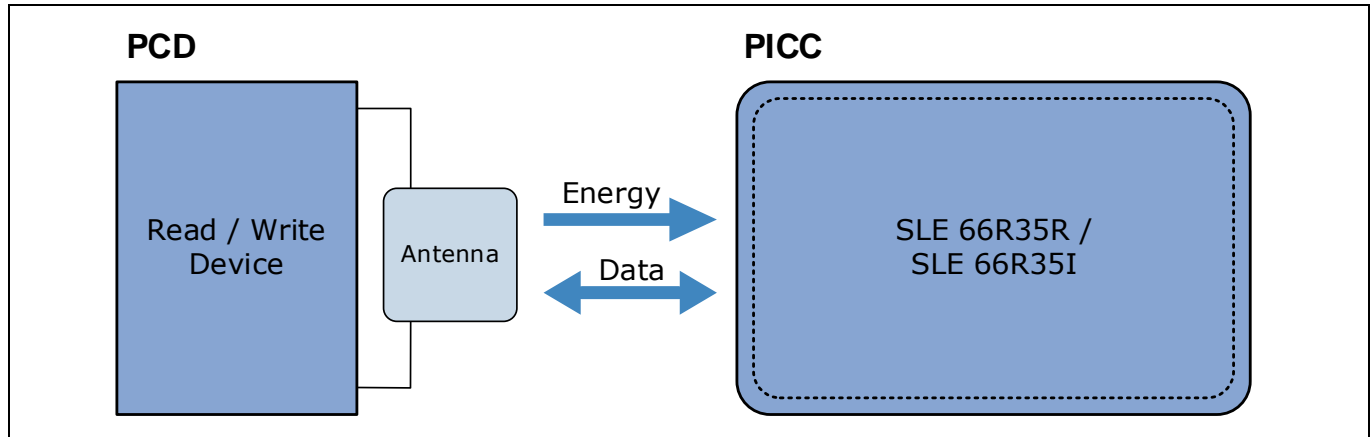
Figure 4 System overview

The operating distance between card and reader antenna is up to 10 cm and more (influenced by external circuitry i.e. reader-antenna configuration).

The RF communication interface transmits at 106 kbit/s resulting in short transaction times, the effect being that a card user can move freely through a reader gate with minimum disruption. A typical ticketing transaction can be handled in less than 100 ms. Robust contactless transmission means that the card with SLE 66R35R/I may also remain in the wallet of the user even if there are coins in it.

An intelligent anticollision function based on the chip's single size unique identifier (uid0-uid3) enables more than one card in the field to operate simultaneously. The anticollision algorithm selects each card individually and enables the execution of a transaction with a selected card is performed correctly without data corruption resulting from other cards in the field.

Access to SLE 66R35R/I is only allowed after a three-pass authentication. The serial number is unique for each card and cannot be changed. Each data transmission is enciphered. Protection from misuse is done by configurable access conditions that are protected by secret keys used for memory operations such as read or write.

Multi-Application Functionality

The SLE 66R35R/I is suited for the use in multi-application schemes, for example combining a transportation fare collection scheme and a ticketing system such as a stadium ticketing. Both applications can be performed with the same card, as hierarchical key management is supported. This means that two different keys for each memory sector can be assigned to enable authentication to that sector.

2.1 Supported standards

The SLE 66R35R/I supports the following standards:

- ISO/IEC 14443-1, -2 and -3 Type A [1]
- Tested according to ISO/IEC 10373-6 [2]

2.2 Command set

A set of standard ISO/IEC 14443-3 Type A commands is implemented to operate the chip.

Additionally the SLE 66R35R/I specific command set is implemented. This facilitates the access to the on-chip integrated memory, supports the execution of authentication, encryption and decryption data as well as an increment or a decrement of a dedicated value counter.

3 Circuit Description

SLE 66R35R/I consists of an EEPROM memory of 1 kByte organized in 16 sector with 4 blocks each containing 16 bytes, an analog interface for contactless energy and data transmission and a control unit.

The power supply and data are transferred to SLE 66R35R/I via an antenna, which consists of a coil with few turns directly connected to the module. No further external components are necessary. The circuit is designed to communicate with a card-reader at an operating distance of up to 10 cm (or more) depending on the reader-antenna configuration.

The chip is designed to meet the cost-optimized requirements of a basic security level. The targeted applications are transport, corporate access, events and loyalty cards with basic security requirements.
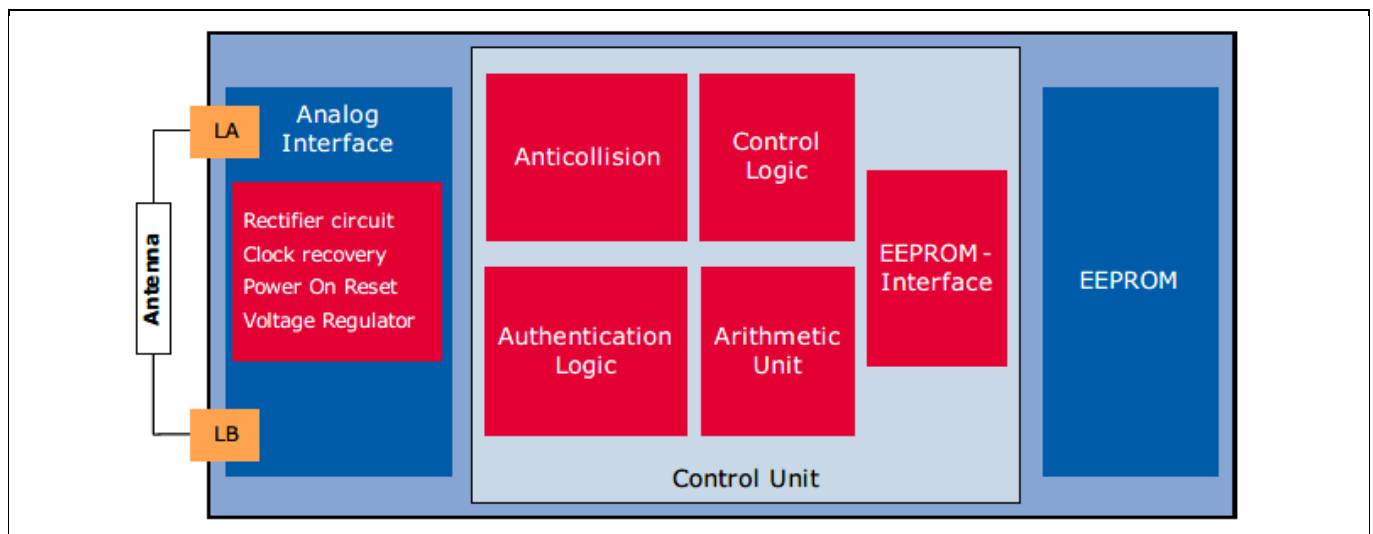
Figure 5 **Block diagram**

- **Analog Contactless Interface:**
  - The Analog Contactless Interface comprises the voltage rectifier, voltage regulator and system clock to supply the IC with appropriate power. Additionally the data stream is modulated and demodulated.
- **Anticollision**
  - Internal logic of SLE 66R35R/I ensures the recognition of several cards in the field which may be selected and operated in sequence.
- **Authentication Logic**
  - Correct execution of any memory operation can only occur after the authentication procedure with a specific key.
- **Control Logic**
  - Access to a block is defined by the associated access conditions for that block. These are programmed individually for each block in a sector.
- **Arithmetic Unit**
  - Arithmetic Capability: increment and decrement of values stored in a special redundant format.
- **EEPROM:**
  - 1 kByte organized in 16 sectors with 4 blocks by 16 bytes each. The last block of each sector is called "Sector Trailer" and is used to store for a pair of secret keys and programmable access conditions for each block.

4 SLE 66R35x options

SLE 66R35R / SLE66R35I support systems based on single size UIDs:

- SLE 66R35R: 4-byte reused ID (r-ID), 4-byte UID previously issued
- SLE 66R35I: 4-byte fixed non-unique number (FNUID) according to ISO/IEC 14443-3 Type A
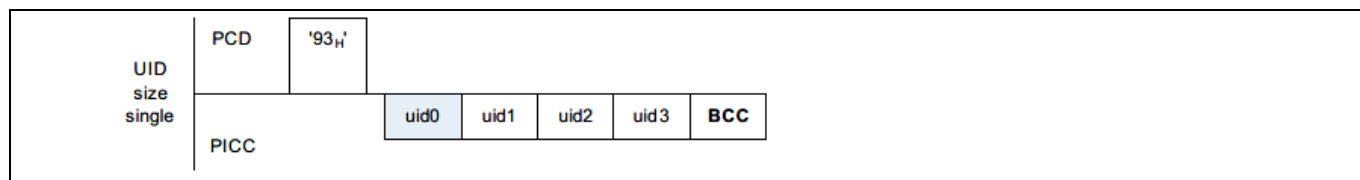
Figure 6   UIDs according to ISO/IEC 14443-3 Type A

4.1 Content of Block 00$_H$

The SLE 66R35R/I is delivered with a single size UID.

Block 00$_H$ is configured as shown in Figure 7.
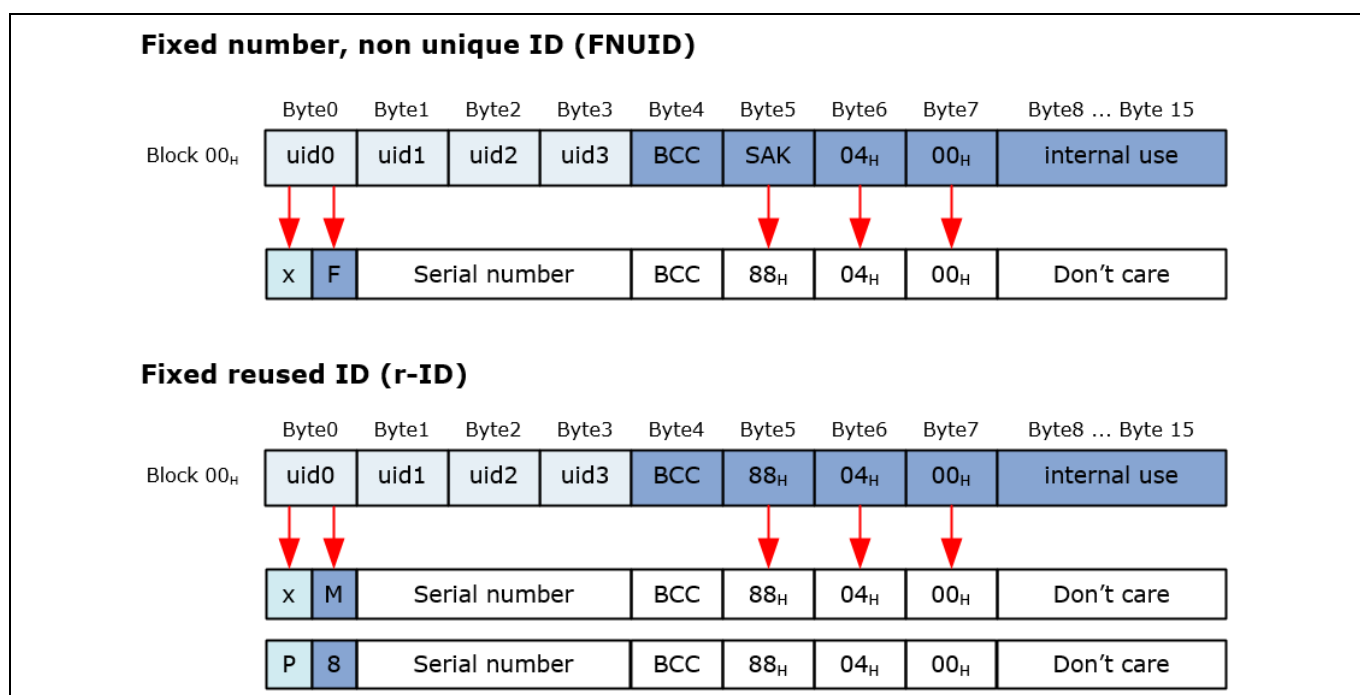
Figure 7   Memory structure block 00$_H$

Serial Number Check, BCC

According to the ISO/IEC14443-3 Type A the BCC is the UID CLn checkbyte, calculated as exclusive-or over the four previous bytes (as described in ISO/IEC 14443-3 Type A).

SLE 66R35x options

4.2 Anticollision

The SLE 66R35R/I supports single cascade anticollision scheme.
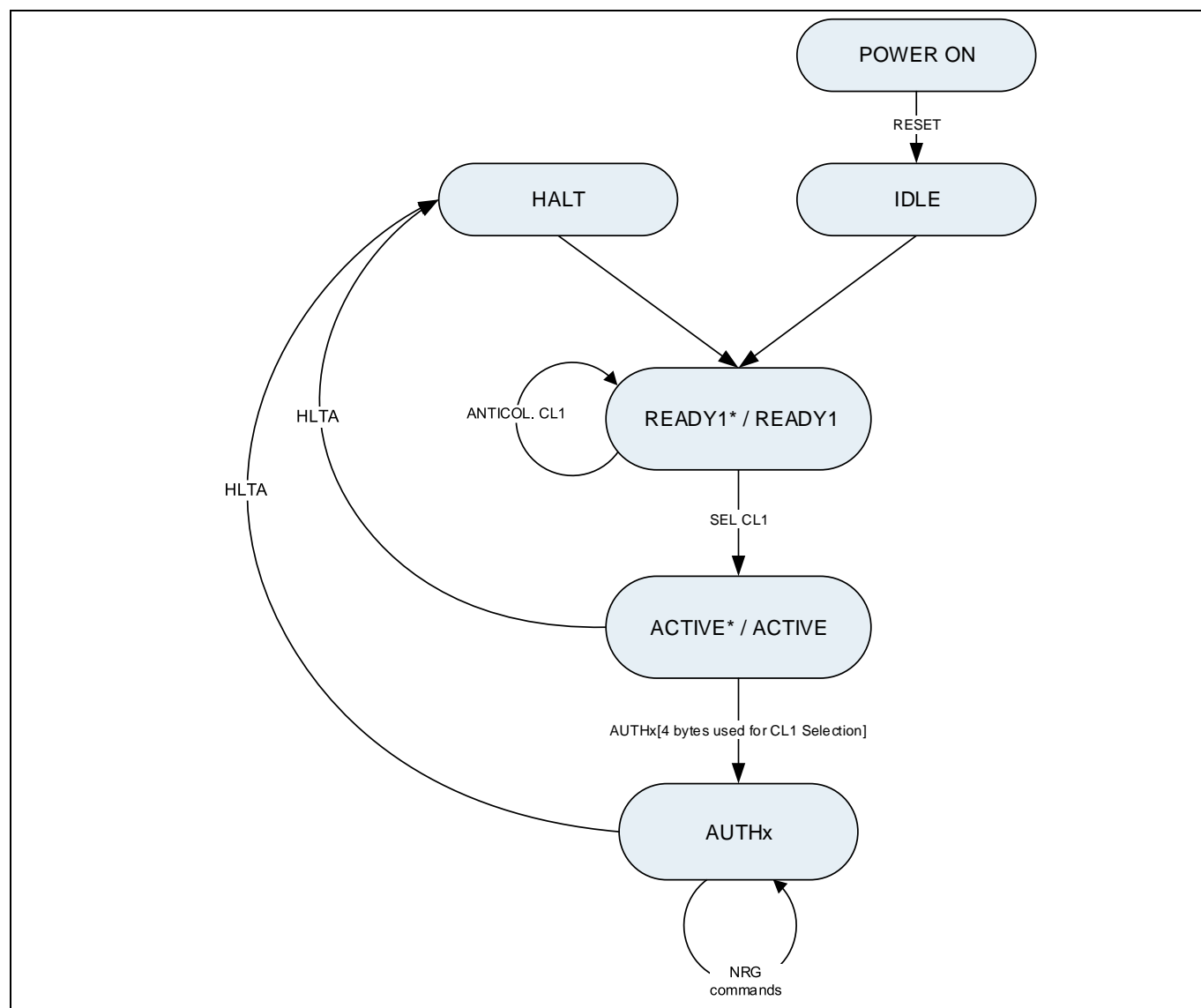
Figure 8   **Single cascade anticollision**

SLE 66R35x options

4.3 Product overview

Following memory chips with NRG are available (see Table 3):

Table 3 Overview on chip types using NRG technology

Type	UID size	UID type	uid0	Description
SLE 66R35[1][2]	4-byte	UID	xM_H $P8_H$	Fixed unique number programmed by manufacturer ($M = 1_H, 5_H, 7_H, 9_H$) and ($P = 1_H, 2_H, 3_H, 4_H, 5_H$)
SLE 66R35I	4-byte	FNUID	xF_H	Fixed number, non-unique programmed by manufacturer
SLE 66R35R	4-byte	r-ID	xM_H $P8_H$	Fixed reused identity number programmed by manufacturer ($M = 1_H, 5_H, 7_H, 9_H$) and ($P = 1_H, 2_H, 3_H, 4_H, 5_H$)
SLE 66R35E7(H)	7-byte	UID	05_H	Fixed unique number programmed by manufacturer (delivery default)
	4-byte	FNUID	xF_H	Fixed number, non-unique derived from 7-byte UID (personalization option). The FNUID is not stored in Block 00_H, it is derived from the 7-byte UID stored in Block 00_H. The derived value for the uid0 byte is logically OR-ed with $1F_H$; due to that x may have following values: $1_H, 3_H, 5_H, 7_H, 9_H, B_H, D_H, F_H$
	4-byte	RND-ID	08_H	uid1 to uid3 is a random number (RND1 - RND3) (personalization option). The RND-ID is not stored in Block 00_H; a new RND-ID is generated with every power-up.

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[1] The available numbers are already exhausted.

[2] Discontinued. Consider to use successor products SLE 66R35I, SLE 66R35R, SLE 66R36E7, SLE 66R35E7H.

5 Memory organization

The 1024 Byte EEPROM is organized in 16 sectors with 4 blocks of 16 bytes each.

Figure 9 Memory organization

Each sector contains four 16 bytes blocks.

- 3 blocks of data are configurable
  - as Read / Write blocks for storing general data
  - as Value blocks for e.g. electronic purse applications
  - Sector 0, Block 00_H contains Manufacturer Data (serial number, SAK, ATQA, etc.), read only
- the Sector Trailer contains the individual secret authentication Key A and optional Key B as well as the block accesses conditions of the respective sector
  - authentication Key A
  - (optional) authentication Key B
  - access bits to define the access conditions for the specified blocks for every sector individually.
    The Sector Trailer and the data blocks are controlled independently.

A successful authentication procedure for the desired sector has to be carried out to allow access to the memory by the appropriate commands.

Furthermore, the access to the EEPROM is controlled by the access conditions (set by the access bits) depending on the application. Applications for contactless access control (e.g. identification) only require read and write operations whereas in revenue control systems (e.g. public transport applications) additional commands like increment, decrement for direct control of the value stored are provided.

All sectors can be assigned to different applications by use of different keys. The authentication procedure is performed between the Reader and the contactless card automatically. Access to stored data is only permitted after successful authentication to that sector.

In erased state the EEPROM cells are as a 1_B, the written state is represented by a 0_B.

Memory organization

5.1 Manufacturer block SLE 66R35R/I

The information within the manufacturer block (Block 00_H, Sector 0) is programmed and locked during the manufacturing process. It is reserved to store:

- 4-byte r-ID or 4-byte FNUID
  - uid1 – uid3 = part of the serial number
- SAK: Select Acknowledge
- ATQA: Answer to Request A
- Manufacturer specific data

Figure 10 and Figure 11 give some details on the content of Block 00_H.
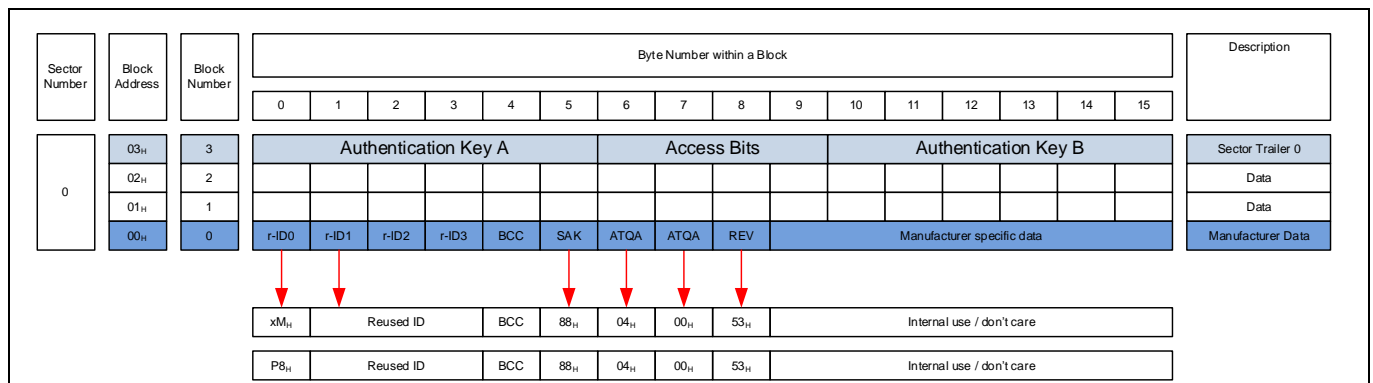
Figure 10   Manufacturer block SLE 66R35R (r-ID)
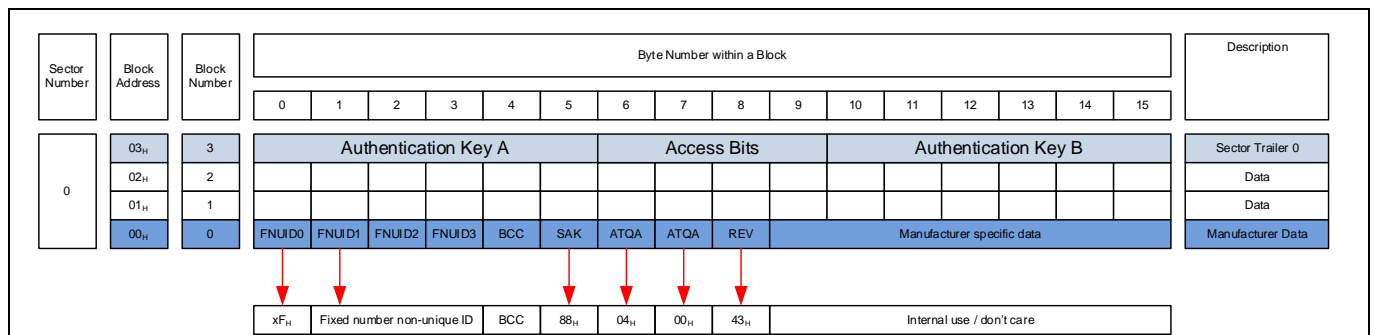
Figure 11   Manufacturer block SLE 66R35I (FNUID)

5.2 Answer to Request (ATQA) and Select Acknowledge (SAK)

Following valid responses are returned to valid …

- ATQA to REQA or WUPA commands and
- SAK to the SELECT command

Table 4   ATQA and SAK responses

Product	ATQA response	SAK Cascade Level 1	Description
SLE 66R35R	00_H 04_H	88_H	4-byte reused number (r-ID) and
SLE 66R35I	00_H 04_H	88_H	4-byte fixed, non-unique number (FNUID)

Memory organization

5.3 Sector Structure

Each sector has three data blocks (block numbers 0, 1 and 2) and the Sector Trailer block (block number 3).

5.3.1 Data blocks / Value blocks

Each data block can be defined as Read / Write block or as Value block by setting the specific access conditions. According to these conditions data can be read, written, incremented, decremented, restored or transferred to the card after a successful authentication with either authentication Key A or authentication Key B.

Read / Write blocks

A Read / Write block is used to store general application data (valid commands: READ, WRITE).

Value block

The Value blocks allows electronic purse functions to be performed. Valid commands are READ, WRITE, INCREMENT, DECREMENT, RESTORE and TRANSFER. The Value blocks have a fixed data format that permits error detection, error correction and a backup management.
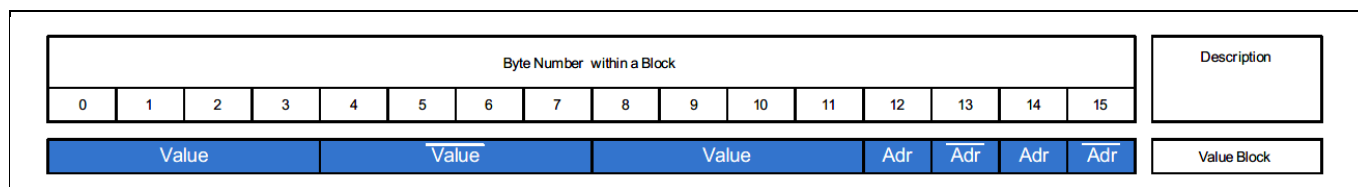
Figure 12   Data structure of a Value block

- Value is a signed four byte value. For the security and integrity purposes it is stored three times; twice non-inverted and once bit-inverted. Values are stored in 2^s complement binary format with the most significant bit as a sign bit.

  The value is stored in the big endian order, hence the most significant byte is stored at the highest address of the reserved memory field.

- The address (Adr) signifies a 1-byte block address. It is used to keep the storage address of a block that is particularly useful for implementing a powerful backup management. The address byte is stored four times, twice inverted and twice non-inverted and it can only be changed by a WRITE command. During INCREMENT, DECREMENT, RESTORE and TRANSFER operations the address is not altered.

The WRITE command must also be used to initialize a Value block. The value and the address must be written in the appropriate format. The arithmetic commands INCREMENT, DECREMENT are available to modify the value content, but Value blocks may still be accessed with READ and WRITE commands. The block management commands RESTORE and TRANSFER are available to manage the backup management and programming of the Value blocks.

Memory organization

5.3.2 Sector Trailer

Every sector has a Sector Trailer (block 3 within every sector) containing the authentication keys (Key A and Key B) and the access condition information of the associated sector, i.e. access to the data blocks (0, 1 and 2) and the Sector Trailer itself.

Figure 13   Data structure of a Sector Trailer

- Authentication Key A (mandatory): these field contains the cryptographic key A which can never be read (a READ command returns the value 00_H 00_H 00_H 00_H 00_H 00_H.
- Authentication Key B (optional): depending on the access conditions this portion of the Sector Trailer can be read after an authentication or never be read (when used as key; in this case a READ command returns the value 00_H 00_H 00_H 00_H 00_H 00_H).
- If the Key B is readable then it can be used for the authentication however any subsequent access to the memory (Read, Write...) is denied and will result in NACK response.
- Access Bits define the access conditions for each block of the sector and the Sector Trailer.
- Byte 9 of the Sector Trailer is reserved for future use and shall not be checked by the application. It does not contain access condition information and might be used for other application data.

Each sector should have different values for authentication Key A and (the optional) authentication Key B as all sectors can be assigned to different applications offered by different system providers. The authentication procedure is performed between the reader and the SLE 66R35R/I card. Only after successful authentication to a sector the access to data blocks is enabled.

5.4 Memory access

The SLE 66R35R/I memory can be accessed using a defined memory commands. Before execution of any memory command, the desired card has to be selected via appropriate operation and the three-pass authentication procedure has to be completed. All possible memory operations for an addressed block in the memory depend on key which is used for the authentication as well as on the assigned access conditions stored in the associated sector trailer, see also 5.4.1.

Table 5   Memory access command set

Command	Op-code	Description	Valid for blocks
READ	30_H	Reads the data from an addressed block if access conditions permit an access.	Data block Value block Sector Trailer
WRITE	$A0_H$	Writes the data to an addressed block if access conditions permit an access.	Data block Value block Sector Trailer
DECREMENT	$C0_H$	Decrement the addressed value by the received value and stores the result into the internal buffer.	Value block

Memory organization

Command	Op-code	Description	Valid for blocks
INCREMENT	C1$_H$	Increment the addressed value by the received value and stores the result into the internal buffer.	Value block
RESTORE	C2$_H$	Loads the content of an addressed block into the internal buffer. Only allowed if the addressed block in a Value block format.	Value block
TRANSFER	B0$_H$	Writes the content of the internal buffer to the addressed block. The addressed block will be programed in Value block format.	Value block

5.4.1 Access conditions

The access conditions (AC) for every data block and Sector Trailer are stored in byte 6, 7 and 8 of the Sector Trailer of each block. These access bits control the access rights of memory access for different commands and keys. The byte 9 is accessible under the same access condition as for Sector Trailer and can be used for the additional data storage.

Bits $C1_Y$, $C2_Y$ and $C3_Y$ (Y = block number, 0..3), are stored twice (non-inverted and inverted) for data integrity reasons. They define the access conditions for every block inside each sector.

If the format of access condition bits in one specific Sector Trailer is incorrect, an authentication to this sector is still possible, but any subsequent memory access will be rejected.
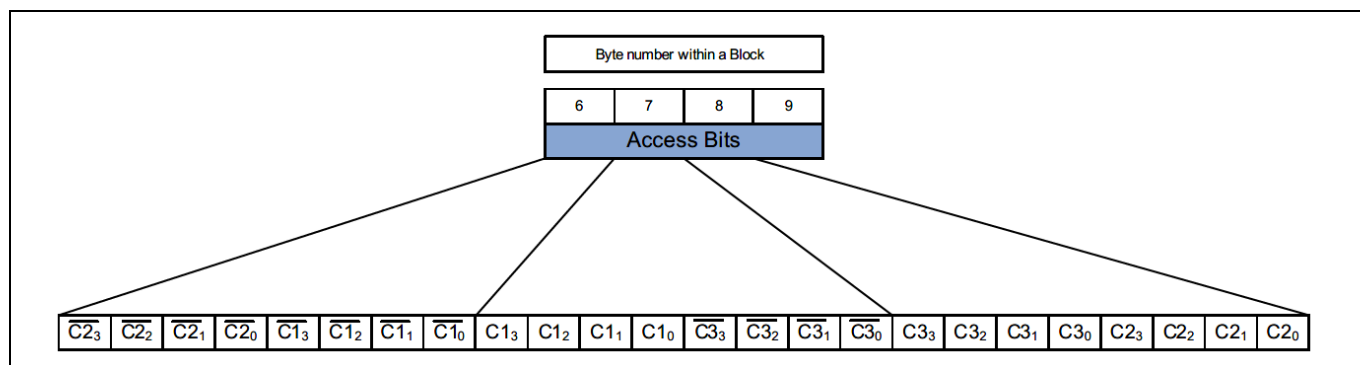
Figure 14   Access conditions

Table 6   Access conditions

Access Bits	Valid for block	Description	Valid commands
$C3_3$ $C2_3$ $C1_3$	3	Sector Trailer	READ, WRITE
$C3_2$ $C2_2$ $C1_2$	2	Data block	READ, WRITE, INCREMENT, DECREMENT, TRANSFER, RESTORE
$C3_1$ $C2_1$ $C1_1$	1	Data block	READ, WRITE, INCREMENT, DECREMENT, TRANSFER, RESTORE
$C3_0$ $C2_0$ $C1_0$	0	Data block	READ, WRITE, INCREMENT, DECREMENT, TRANSFER, RESTORE

5.4.2 Access conditions for Sector Trailer

Depending on the access conditions for the Sector Trailer read / write access to either authentication Key A or Key B or to the access bits is specified as 'Never', 'Key A' or 'Key B'. 'Key A / Key B' means that the access is possible only after an authentication to the sector using authentication Key A or authentication Key B.

Memory organization

Table 7   Access condition for Sector Trailer

Access Bits			Access condition for						Remark
			Authentication Key A		Access Bits[1]		Authentication Key B		
$C3_3$	$C2_3$	$C1_3$	READ	WRITE	READ	WRITE	READ	WRITE	
0	0	0	Never	Key A	Key A	Never	Key A	Key A	Key B may be read[2]
0	0	1	Never	Key B	Key A / Key B	Never	Never	Key B	
0	1	0	Never	Never	Key A	Never	Key A	Never	Key B may be read
0	1	1	Never	Never	Key A / Key B	Never	Never	Never	
1	0	0	Never	Key A	Key A	Key A	Key A	Key A	Transport Configuration
1	0	1	Never	Never	Key A / Key B	Key B	Never	Never	
1	1	0	Never	Key B	Key A / Key B	Key B	Never	Key B	
1	1	1	Never	Never	Key A / Key B	Never	Never	Never	

5.4.3   Access conditions for data blocks

Depending on the access bits for the data blocks (block number Y = 0…2) the read / write access is specified as 'Never', 'Key A', 'Key B' or 'Key A / Key B' (Key A or Key B). The setting of the relevant access bits defines the application and the corresponding applicable commands.

Table 8   Access condition for data blocks

Access Bits (Y = 0 … 2)			Access condition Data block (block number Y = 0 … 2)				Application
$C3_Y$	$C2_Y$	$C1_Y$	READ	WRITE[3]	INCREMENT	DECREMENT TRANSFER RESTORE	
0	0	0	Key A / Key B	Key A / Key B	Key A / Key B	Key A / Key B	Transport Configuration Read / Write / Value block
0	0	1	Key A / Key B	Key B	Never	Never	Read / Write block
0	1	0	Key A / Key B	Never	Never	Never	Read / Write block
0	1	1	Key A / Key B	Key B	Key B	Key A / Key B	Value block
1	0	0	Key A / Key B	Never	Never	Key A / Key B	Value block
1	0	1	Key B	Never	Never	Never	Read / Write block

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[1]   Access Bits can also be locked which prevents any further changes of the access conditions

[2]   The bytes reserved to store for Key B may be used to store data. Access to data blocks is not possible.

[3]   Write access conditions do not apply for Block 00_H of Sector 0.

Memory organization

Access Bits (Y = 0 … 2)			Access condition Data block (block number Y = 0 … 2)				Application
C3$_Y$	C2$_Y$	C1$_Y$	READ	WRITE[3]	INCREMENT	DECREMENT TRANSFER RESTORE	
1	1	0	Key B	Key B	Never	Never	Read / Write block
1	1	1	Never	Never	Never	Never	Read / Write block

5.5   Transport configuration

At delivery the memory is predefined. After a successful authentication with either Key A or Key B of a sector the respective data can be accessed.

- Data area is set to 00$_H$ as default

- Key A and Key B are set to FF$_H$ FF$_H$ FF$_H$ FF$_H$ FF$_H$ FF$_H$ as default

Reading of Key A is never allowed, thus the value 00$_H$ 00$_H$ 00$_H$ 00$_H$ 00$_H$ 00$_H$ is read. Sector Trailer byte 9 is set arbitrary.

Figure 15   Memory map – Transport Configuration SLE 66R35R/I

6 Communication Principle

6.1 State diagram

The SLE 66R35R/I is fully compliant to the ISO/IEC 14443-3 Type A specification.

All operations are initiated by an appropriate reader and controlled by the internal logic of the SLE 66R35R/I. Prior to any memory access the card has to be selected according to the ISO/IEC 14443-3 Type A anticollision and selection scheme.

After the anticollision and selection the reader may enter the protected state by performing the authentication procedure to any sector by sending the AUTHENTICATE command (either with Key A or with Key B and a valid block address). Any other command will cause an error and the SLE 66R35R/I will return either to IDLE or to HALT state.

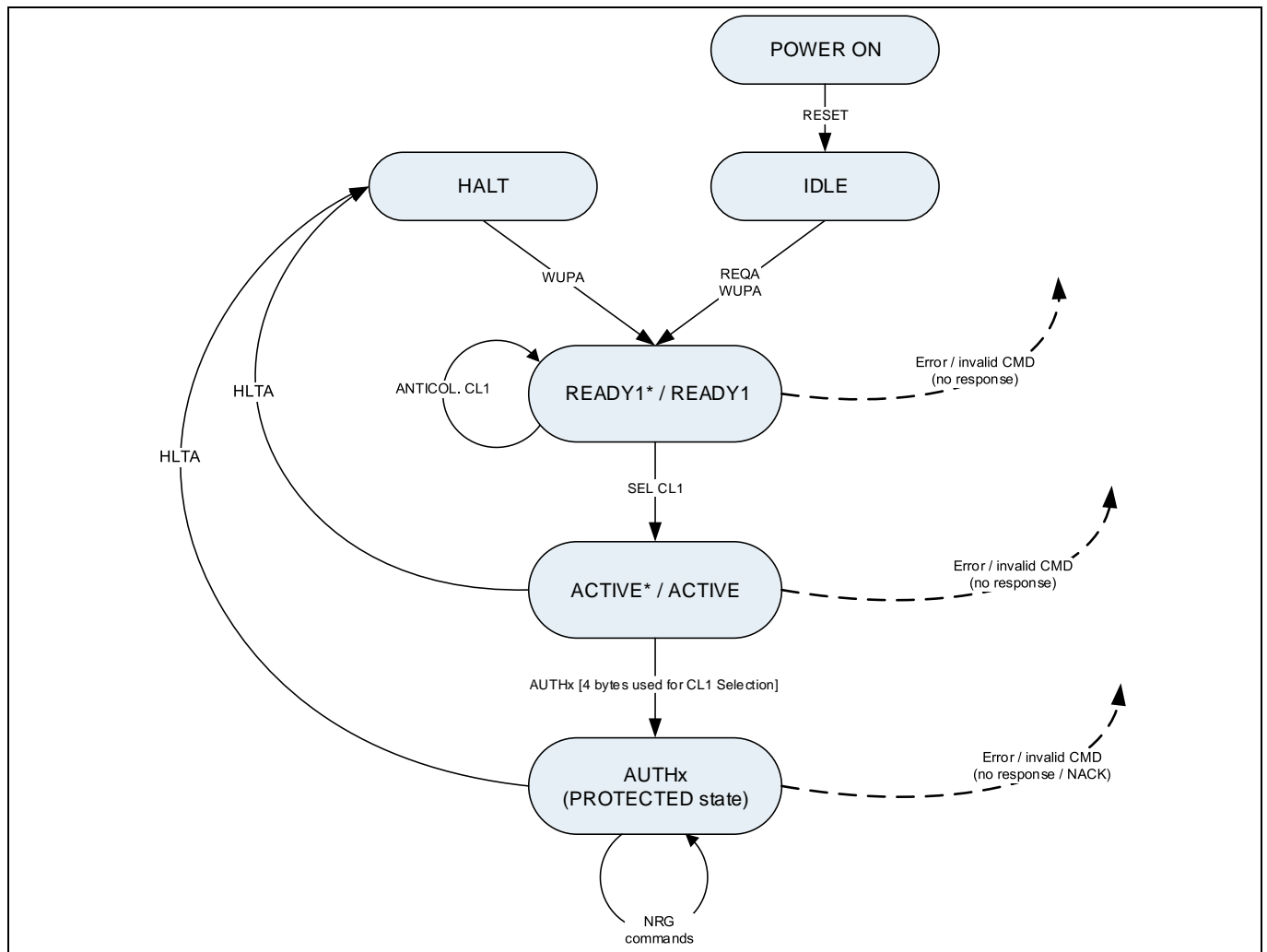Figure 16 shows the state diagram of the SLE 66R35R/I.

Figure 16   SLE 66R35R/I state diagram

Communication Principle

6.1.1 IDLE state

After POWER ON, the SLE 66R35R/I is in IDLE state.

If REQA or WUPA command is executed in this state, the SLE 66R35R/I transits to READY1 state. Any other command is interpreted as an error and the SLE 66R35R/I stays in IDLE state without any response.

6.1.2 READY1 / READY1* state

In READY1 / READY1* state:

- the whole UID can be retrieved by using ISO/IEC 14443-3 Type A Anticollision and Select commands. After the Select command is executed properly the IC transits to ACTIVE / ACTIVE* state. The answer to a Select command in READY1 / READY1* state is Select Acknowledge (SAK), which indicates that the UID is complete (see also ISO/IEC 14443-3 Type A).
  Any other command or any other interruption is interpreted as an error and the SLE 66R35R/I returns back to IDLE or HALT state without any response, depending from which state it came from.

6.1.3 ACTIVE / ACTIVE* state

In ACTIVE / ACTIVE* state during an execution of a valid AUTHENTICATE command 4 bytes of the UID are used to initialize the CRYPTO1.

The ACTIVE / ACTIVE* state is left with a HLTA command. The SLE 66R35R/I transits to HALT state and waits until a WUPA command is received.

If any error is detected the SLE 66R35R/I sends "No Response" (NR) or "Not Acknowledge" (NACK) and transits to IDLE or HALT state depending on the previous state.

6.1.4 PROTECTED state (AUTHx)

In the PROTECTED State SLE 66R35R/I memory access commands can be executed. These commands can be applied to the currently authenticated sector in any order, including the authentication of another sector with the following exceptions and recommendations:

- arithmetic commands such as INCREMENT and DECREMENT shall only be applied to blocks in Value block format
- successful execution of a RESTORE, INCREMENT or DECREMENT command is recommended before executing the TRANSFER command
- SLE 66R35R/I exits the PROTECTED state upon reception of the HLTA command or in a case of an error

6.1.5 HALT state

The HLTA command sets the SLE 66R35R/I in the HALT state. The SLE 66R35R/I sends no response to the HLTA command. The HALT state is left by a WAKE-UP A (WUPA) request. Any other data received is interpreted as an error, the SLE 66R35R/I sends no response and remains in HALT state.

6.2 Start up

Latest 250 µs after entering the powering field the SLE 66R35R/I is ready to receive a command.

6.3 Frame Delay Time

For detailed timings see ISO/IEC 14443-3 Type A Standard.

Communication Principle

Note: The response timing of a particular command in specified in the command description.

6.4   Error handling

The SLE 66R35R/I responds to valid frames only. In case of any error the SLE 66R35R/I returns to its initial state, either IDLE or HALT state. Depending on the error type, the SLE 66R35R/I responds either with a Not Acknowledge (NACK) or does not respond (NR).

Table 9 lists ACK and NACK responses.

Table 9   ACK and NACK responses

Response	Code (4 bits)[1]
ACK	1010_B
NACK0	0000_B
NACK1	0001_B
NACK4	0100_B
NACK5	0101_B
NR[2]	n.a.

Table 10 describes the behavior of the SLE 66R35R/I in different error cases.

Table 10   Behavior in case of error

Current States	Command or Error	Response	Next State
IDLE / HALT READY1 / READY1*	Invalid op-code	NR	IDLE / HALT
	Parity, Miller or CRC error	NR	IDLE / HALT
	Command too short or too long	NR	IDLE / HALT
	Invalid address	NR	IDLE / HALT
	Other errors	NR	IDLE / HALT
ACTIVE / ACTIVE*	Invalid op-code	NR	IDLE / HALT
	Parity, Miller, CRC error	NACK1 or NACK5	IDLE / HALT
	Command too short or too long	NR	IDLE / HALT
	Invalid address	NACK0 or NACK4	IDLE / HALT
	Other errors	NACK0 or NACK4	IDLE / HALT

6.5   Data integrity

Reliable data transmission in the contactless communication link is supported by following mechanisms:

- 16 bit CRC (Cyclic Redundancy Check) for each data transmission
- Parity bits for each byte
- Monitoring of protocol sequence
    - bit coding to distinguish between "1", "0" and "no information"
    - bit stream analysis
    - bit count checking

---

[1]   No integrity mechanism, the response is encrypted.

[2]   NR = no response. Depending on the state the does not send a response on errors.

6.6 Three-pass authentication

A basic security level is provided using a three-pass authentication between the SLE 66R35R/I and the reader. The three-pass authentication sequence is split into following steps:

– The reader selects the sector to be accessed by addressing a block within this sector. The reader transmits the AUTH command using either authentication Key A or Key B.
– SLE 66R35R/I reads the secret key and the associated access conditions from the Sector Trailer. Then SLE 66R35R/I transmits a random number as a challenge to the reader (step 1).
– The reader calculates the response using the secret key and the challenge from SLE 66R35R/I. The reader transmits its response together with its own challenge (random number) to SLE 66R35R/I (step 2).
– SLE 66R35R/I verifies the reader response. Then SLE 66R35R/I calculates the response to the challenge from the reader and returns it to the reader (step 3).
– The reader verifies the received response to its own challenge.

After the first random challenge (step 1) further data transmission is encrypted.

6.7 Memory access / operations

After a mutual authentication any of the following operations may be performed:

– READ block
– WRITE block
– DECREMENT: decrements the content of a Value block and stores the result in an internal data register
– INCREMENT: increments the content of a Value block and stores the result in an internal data register
– RESTORE: moves the content of a Value block into an internal data register
– TRANSFER: writes the content of the temporary internal data register to a block

7 Command set

7.1 Supported ISO/IEC 14443-3 Type A command set

Table 11 describes the ISO14443-3 Type A command set which is supported by SLE 66R35R/I.

Table 11   ISO/IEC 14443-3 Type A command set

Command	Abbreviation	Command Op-Code	Description
Request A	REQA	26_H	Short Frame Command Type A request to all ISO/IEC 14443-3, Type A compatible chips in IDLE State
Wake Up A	WUPA	52_H	Short Frame Command Type A Wake Up request to all ISO/IEC 14443-3 Type A compatible chips
Anticollision	AC	93_H NVB_H	Cascade level 1 with the Number of Valid Bits (NVB)
Select	SELA	93_H 70_H	Select the UID of Cascade level 1
HaltA	HLTA	$5x_H$[1][2]	Sets a chip to a HALT state

For a detailed command description please refer to ISO/IEC 14443-3 Type A standard.

7.2 Memory access command set

There are two command types implemented:

- One step commands:
  - the PCD sends a command, the PICC sends a response
- Two step commands:
  - the PCD sends the first command, the PICC sends the first response
  - if the PCD does not detect an error the PCD sends the second command and the PICC responds with the answer to the second command

The command set of the SLE 66R35R/I is listed in Table 12.

Table 12   SLE 66R35R/I command set

Command	Abbreviation	Command Op-code	Description
Authenticate with Key A	AUTHA	60_H	Authentication with Key A to the sector in which the address block is located.
Authenticate with Key B	AUTHB	61_H	Authentication with Key B to the sector in which the address block is located.
READ	RD	30_H	Reads the data from an addressed block if access conditions permits access.
WRITE	WR	$A0_H$	Writes the data to an addressed block if access conditions permits the access.
DECREMENT	DCR	$C0_H$	Decrement the addressed value by the received value and writes the result into the internal transfer buffer.

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[1]   SLE 66R35R/I accepts also command op-code values $5x_H$.

[2]   Within the parameter field of the HLTA command values from 00_H to $3F_H$ are accepted.

Command set

Command	Abbreviation	Command Op-code	Description
INCREMENT	INC	C1$_H$	Increment the addressed value by the received value and writes the result into the internal transfer buffer.
RESTORE	RSTR	C2$_H$	Loads the content of an addressed block into the transfer buffer. Only allowed if the addressed block in a value block format. Data is lost in case of power loss.
TRANSFER	TRFR	B0$_H$	Writes the content of the transfer buffer to the addressed block. The addressed block will be programed in value block format.

The data in the command are always sent LSByte first. Within a byte the LSBit is sent first

7.2.1   AUTHENTICATE (AUTHA and AUTHB)

AUTHA or AUTHB command performs the authentication using either Key A or Key B to a sector. The specified address indicates the block number. The valid address range is 00$_H$ to 3F$_H$. If any other address is specified the SLE 66R35R/I replies with an error message. This command shall be executed after the SLE 66R35R/I has been selected.

The AUTHENTICATE command is performed as a three-pass authentication:

− Step 1:
The PCD sends the AUTHA or AUTHB command to the PICC. The paramenter, a valid block address, indicates the sector for the authentication. The PICC decodes this command and replies with a challenge, TokenRB, which is a 32 bit random number[1].

− Step 2:
The PCD generates a 64 bit TokenAB. The PCD encrypts its own 32 bit random number RA as well as the value RB' which is a modified RB value and sends it to the PICC. The PICC receives and encrypts the received value. Finally the PICC compares the received RB' it to its own RB'.

− Step3:
If the values are identical the PICC will repond the next challenge, TokenBA, which is encrypted value RB" confirming its authenticity. Otherwise an error (NR) will break the authentication. The PCD verifies the response and if no error is detected it allows further memory access i.e. both PCD and PICC are authenticated to each other.

Table 13   AUTHENTICATE command step 1

Command	Parameter	Integrity Mechanism		Response
AUTHA or AUTHB	**Block address**	**CRC**		**TokenRB**
60$_H$ or 61$_H$	00$_H$ - 3F$_H$	2 bytes CRC		32 bit random number or NR

Table 14   AUTHENTICATE command step 2

Command	Parameter	Integrity Mechanism		Response
AUTHA or AUTHB	**Block address**	**CRC**		**TokenBA or NR**
-	64 bit Token	-		32 bit random number or NR

Command set

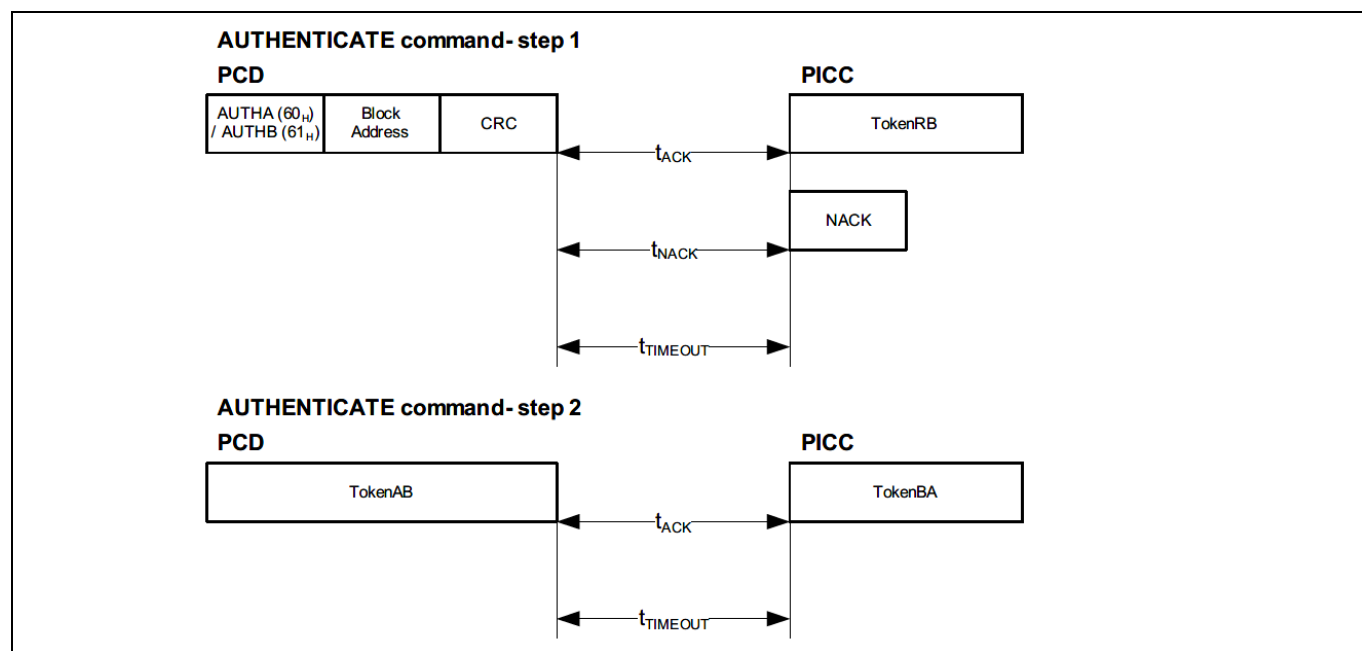**Figure 17   AUTHENTICATE command**

Table 15   Timing AUTHENTICATE command

AUTHENTICATE	$t_{ACK\ min}$	$t_{ACK\ max}$	$t_{NACK\ min}$	$t_{NACK\ max}$	$t_{TIMEOUT}$
Command step 1	147.82 µs	t_{nonNVM}	86.11 µs	t_{nonNVM}	$t_{nonNVM} \leq 950$ µs
Command step 2	86.11 µs	t_{nonNVM}	86.11 µs	t_{nonNVM}	$t_{nonNVM} \leq 950$ µs

Command set

7.2.2   READ (RD)

The READ command reads 16 bytes from the specified block address in the memory:

- − if the specified address is allocated in a previously authenticated sector
- − if the access conditions for the adressed block allow read access

Table 16   READ command

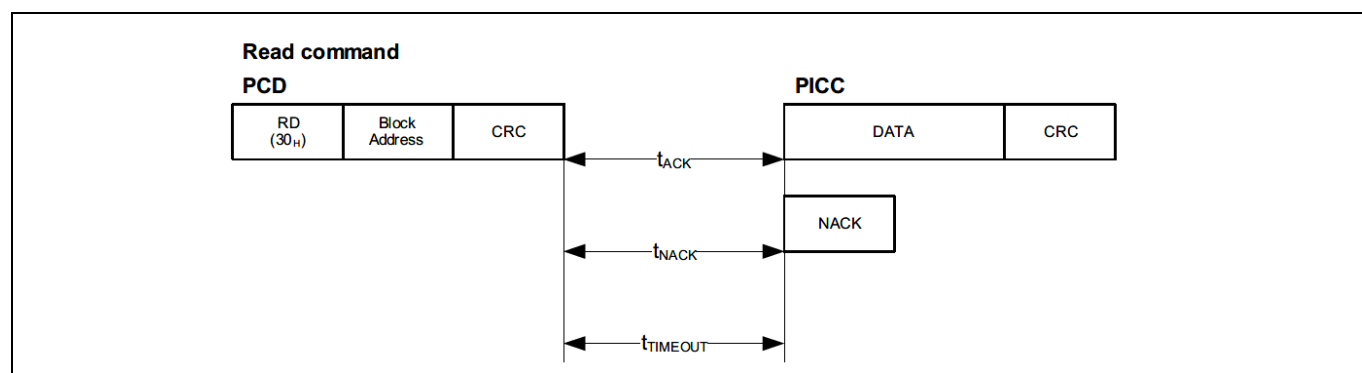| Command	Parameter	Integrity Mechanism		Response
READ	**Block address**	**CRC**		**DATA + CRC**
30_H	00_H - $3F_H$	2 bytes CRC		16 bytes data + 2 bytes CRC or NACK or NR

Figure 18   READ command

Table 17   Timing READ command

READ	$t_{ACK\ min}$	$t_{ACK\ max}$	$t_{NACK\ min}$	$t_{NACK\ max}$	$t_{TIMEOUT}$
	95.98 µs	t_{rdNVM}	86.11 µs	t_{rdNVM}	$t_{rdNVM} \leq 4750$ µs

Command set

7.2.3   WRITE (WR)

The WRITE command writes 16 bytes to the specified address in the memory:

- if the specified address is allocated in a previously authenticated sector
- if the access conditions for the adressed block allow write access

Note: The Block 00$_H$ is never writable independent on the access conditions.

Table 18   WRITE command step 1

Command	Parameter	Integrity Mechanism		Response
WRITE	**Block address**	**CRC**		
A0$_H$	00$_H$ – 3F$_H$	2 bytes CRC		ACK or NACK or NR

Table 19   WRITE command step 2

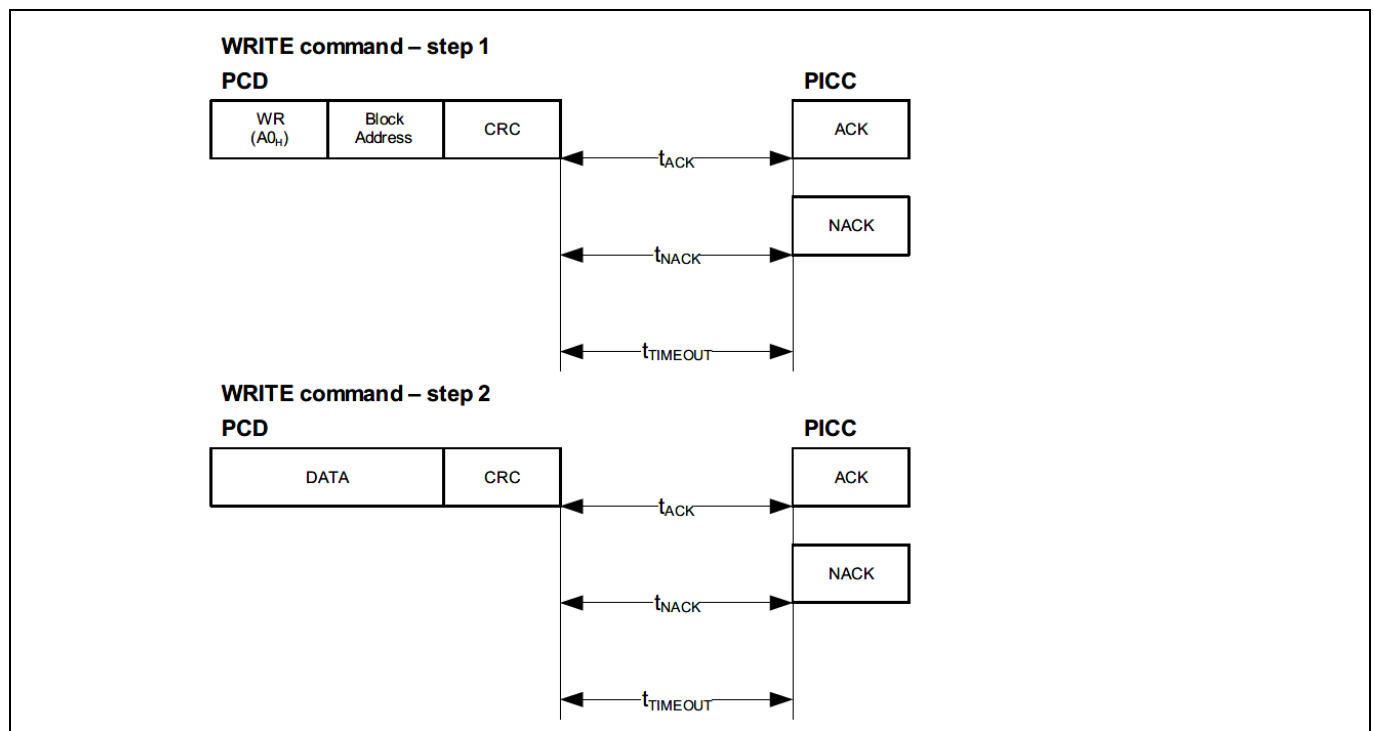| Command	Parameter	Integrity Mechanism		Response
WRITE	**DATA**	**CRC**		
-	16 bytes data	2 bytes CRC		ACK or NACK or NR

Figure 19   WRITE command

Table 20   Timing WRITE command

WRITE	$t_{ACK\ min}$	$t_{ACK\ max}$	$t_{NACK\ min}$	$t_{NACK\ max}$	$t_{TIMEOUT}$
Command step 1	86.38 μs	t_{rdNVM}	86.11 μs	t_{rdNVM}	t_{rdNVM} ≤ 4750 μs
Command step 2	3546.73 μs	t_{wrNVM}	86.11 μs	t_{wrNVM}	t_{wrNVM} ≤ 9500 μs

Command set

7.2.4 DECREMENT (DCR)

The DECREMENT command reads out the content of the addressed value block (block address) and decrements it by the value given in the command data field. The conditions are the following:

- − the specified block address is allocated in a previously authenticated sector
- − the data at specified block address is stored in a value block format
- − the access conditions for the adressed block allow decrement operation

The decremented data is stored internally and can be written to any authenticated data block using the TRANSFER command.

Table 21   DECREMENT command step 1

Command	Parameter	Integrity Mechanism		Response
DECREMENT	**Block address**	**CRC**		
$C0_H$	00_H - $3F_H$	2 bytes CRC		ACK or NACK or NR

Table 22   DECREMENT command step 2

Command	Parameter	Integrity Mechanism		Reponse
DECREMENT	**Value**	**CRC**		
-	4 byte counter value[1]	2 bytes CRC		Nack or NR[2]
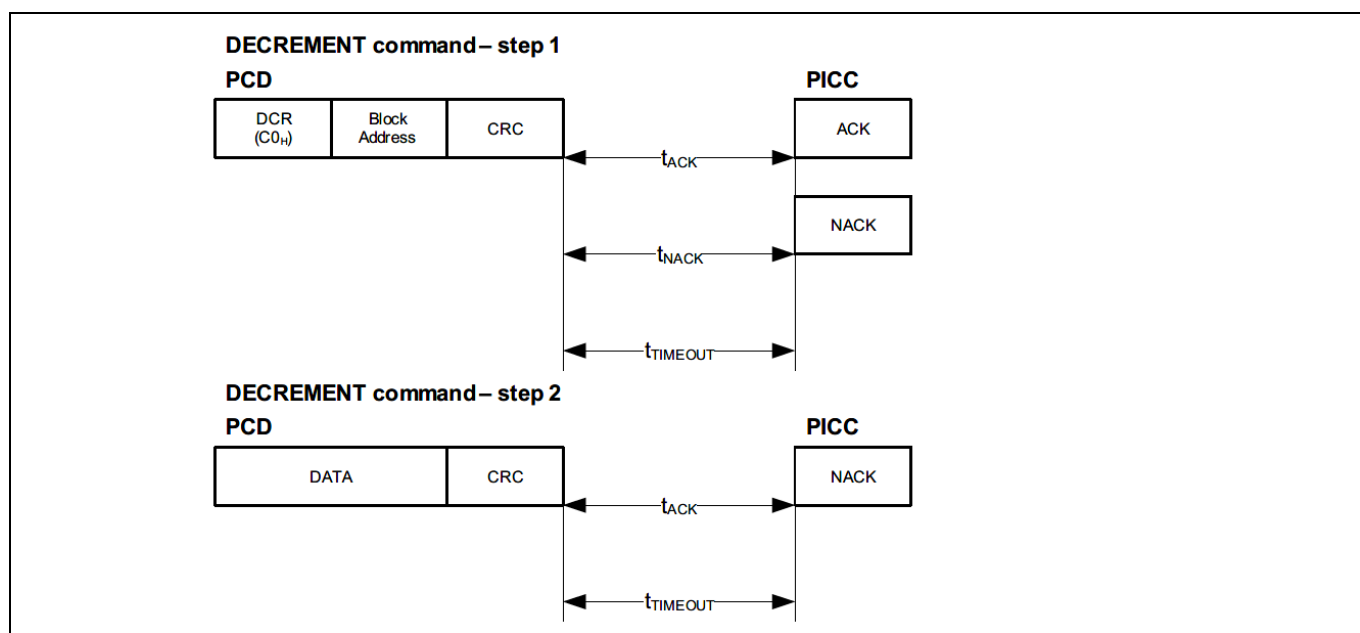
Figure 20   DECREMENT command

Table 23   Timing DECREMENT command

Decrement	$t_{ACK\ min}$	$t_{ACK\ max}$	$t_{NACK\ min}$	$t_{NACK\ max}$	$t_{TIMEOUT}$
Command step 1	256.53 µs	t_{rdNVM}	86.11 µs	t_{rdNVM}	$t_{rdNVM} \leq 4750$ µs
Command step 2	-	t_{rdNVM}	86.11 µs	t_{rdNVM}	$t_{rdNVM} \leq 4750$µs

---

[1]   The value is a 4 byte signed integer value and should be bigger than zero and positive. Please note that the sign (most significant bit of the value) will be ignored e.g. a decrement by $7F_H$ FF_H FF_H FF_H or FF_H FF_H FF_H FF_H returns the same result.

[2]   In the case of successful decrement there is NR (no response).

Command set

7.2.5   INCREMENT (INC)

The INCREMENT command internally reads the value block out of the specified block address. The value is incremented by the given value in command data field under the following conditions:

- the specified block address is allocated in a previously authenticated sector
- the data at the specified block address are stored in a value block format
- the access conditions for the adressed block allow decrement access

The incremented data is stored internally and can be written to any authenticated data block using the TRANSFER command.

Table 24   INCREMENT command step 1

Command	Parameter	Integrity Mechanism		Response
INCREMENT	**Block address**	**CRC**		
$C1_H$	00_H - $3F_H$	2 bytes CRC		ACK or NACK or NR

Table 25   INCREMENT command step 2

Command	Parameter	Integrity Mechanism		Response
INCREMENT	**Value**	**CRC**		
-	4 byte counter value[1]	2 bytes CRC		NACK or NR[2]
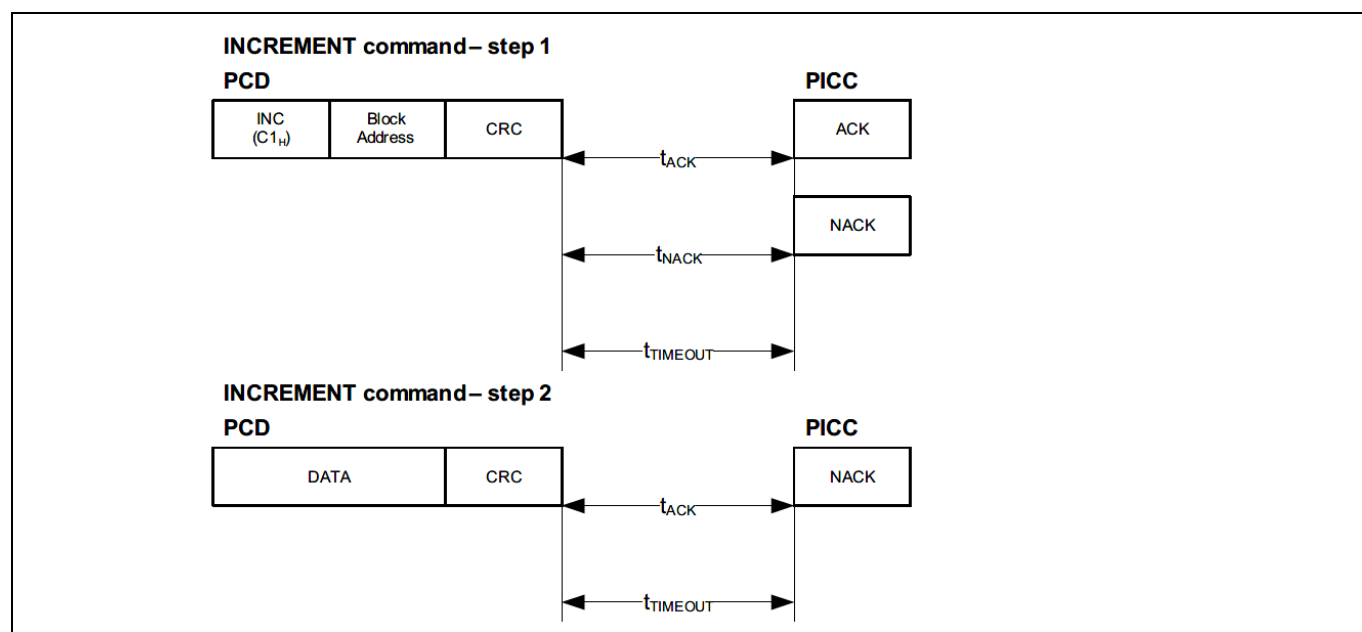
Figure 21   INCREMENT command

Table 26   Timing INCREMENT Command

Increment	$t_{ACK\ min}$	$t_{ACK\ max}$	$t_{NACK\ min}$	$t_{NACK\ max}$	$t_{TIMEOUT}$
Command step 1	256.53 µs	t_{rdNVM}	86.11 µs	t_{rdNVM}	t_{rdNVM} ≤ 4750 µs
Command step 2		t_{rdNVM}	86.11 µs	t_{rdNVM}	t_{rdNVM} ≤ 4750 µs

---

[1]   The value is a 4 byte signed integer value and should be bigger than zero and positive. Please note that the sign (most significant bit of the value) will be ignored e.g. a decrement by $7F_H$ FF_H FF_H FF_H or FF_H FF_H FF_H FF_H returns the same result.

[2]   In the case of successful increment there is NR (no response)

Command set

7.2.6 RESTORE (RSTR)

The RESTORE command reads the value (4 byte) and the address (1 byte) data from the specified block address in the memory and stores it into the Transfer Buffer under the following conditions:

  – the specified address is allocated in an authenticated sector

The restored data can be written to any authenticated block by the subsequent TRANSFER command.

Table 27   RESTORE command step 1

Command	Parameter	Integrity Mechanism	Response
RESTORE	**Block address**	**CRC**	
$C2_H$	00_H - $3F_H$	2 bytes CRC	ACK or NACK or NR

Table 28   RESTORE command step 2

Command	Parameter	Integrity Mechanism	Response
RESTORE	**DATA**	**CRC**	
-	4 byte data	2 bytes CRC	NACK or NR
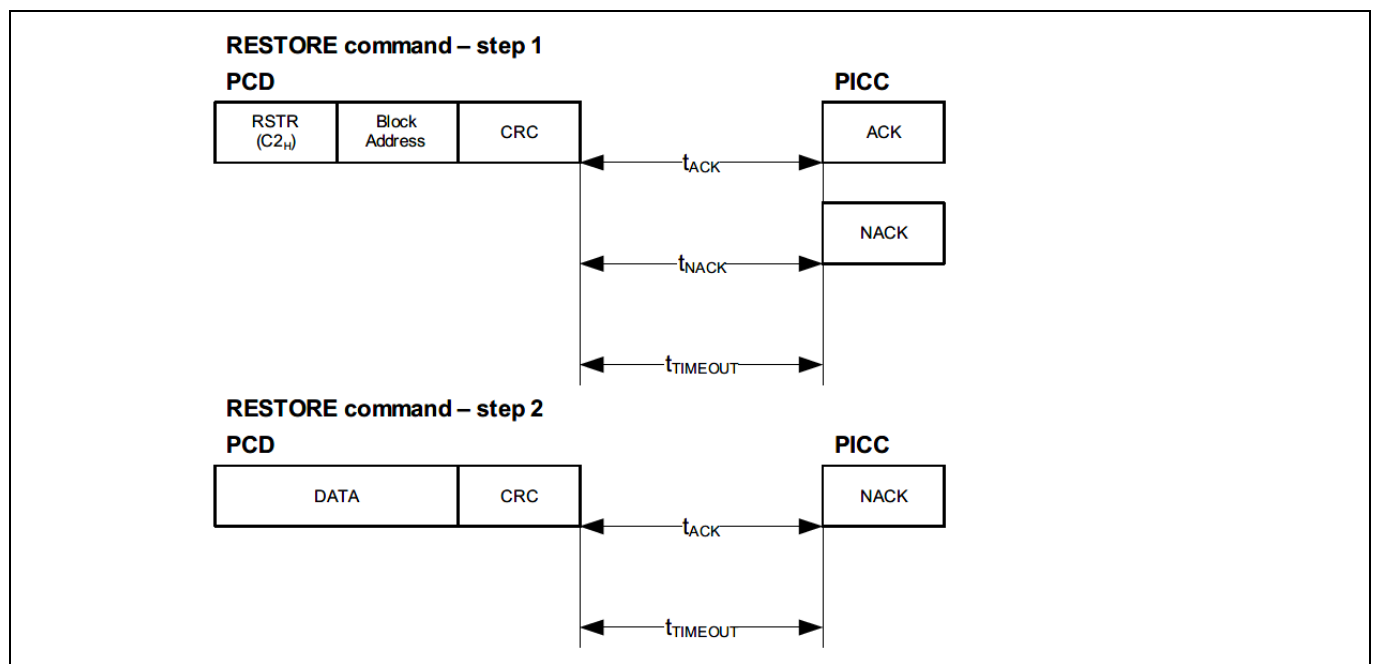
Figure 22   RESTORE command

Table 29   Timing RESTORE command

RESTORE	$t_{ACK\ min}$	$t_{ACK\ max}$	$t_{NACK\ min}$	$t_{NACK\ max}$	$t_{TIMEOUT}$
Command step 1	256.53 µs	t_{rdNVM}	86.11 µs	t_{rdNVM}	$t_{rdNVM} \leq 4750$ µs
Command step 2		t_{rdNVM}	86.11 µs	t_{rdNVM}	$t_{rdNVM} \leq 4750$µs

Command set

7.2.7 TRANSFER (TRFR)

The TRANSFER command programs the data to the specified block address using the RESTORE, INCREMENT or DECREMENT command under the following conditions:

- the specified block address is allocated in a authenticated sector
- pre-condition is a properly executed INCREMENT, DECREMENT or RESTORE command

The TRANSFER command completes a preceeding INCREMENT or DECREMENT or RESTORE operation. It is possible to transfer value blocks between different sectors.

Note: *Consecutive TRANSFER commands will cause a NACK response.*

Table 30   TRANSFER command

Command	Parameter	Integrity Mechanism	Response
RESTORE	**Block address**	**CRC**	
$B0_H$	00_H - $3F_H$	2 bytes CRC	ACK or NACK or NR
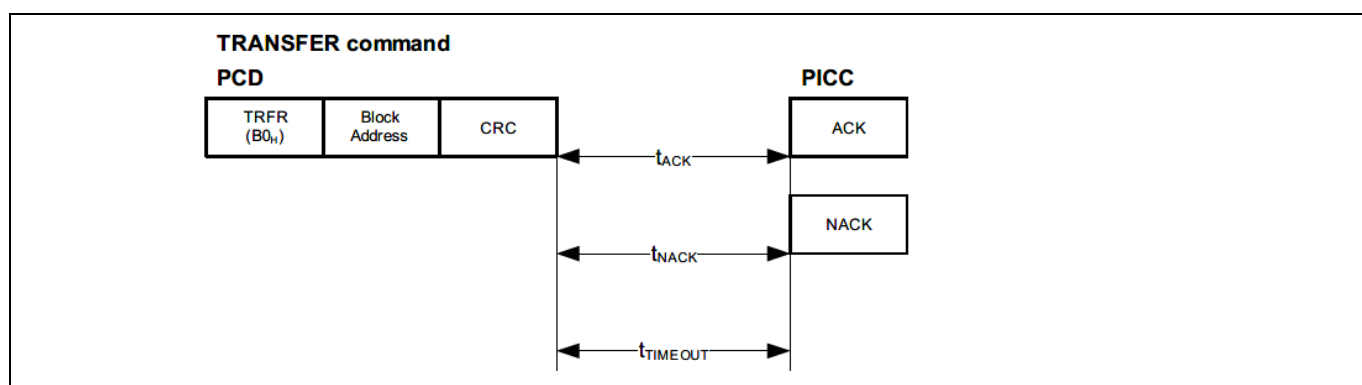
Figure 23   TRANSFER command

Table 31   Timing TRANSFER command

TRANSFER Command	$t_{ACK\,min}$	$t_{ACK\,max}$	$t_{NACK\,min}$	$t_{NACK\,max}$	$t_{TIMEOUT}$
	3546.73 µs	t_{wrNVM}	86.11 µs	t_{wrNVM}	$t_{wrNVM} \leq 9500$ µs

Command set

7.2.8 HLTA

The HLTA command sets the SLE 66R35R/I to the HALT state. The HALT State allows the user to separate already identified SLE 66R35R/I from each other.

Table 32   HLTA command

Command	Parameter	Integrity Mechanism		Response
HLTA	**Block address**	**CRC**		
$5x_H$	00_H - $3F_H$	2 bytes CRC		NACK or ACK

Note: *The HLTA command will be sent plain in ACTIVE state and encrypted in PROTECTED state. If the command is not received properly the SLE 66R35R/I replies a NACK.*
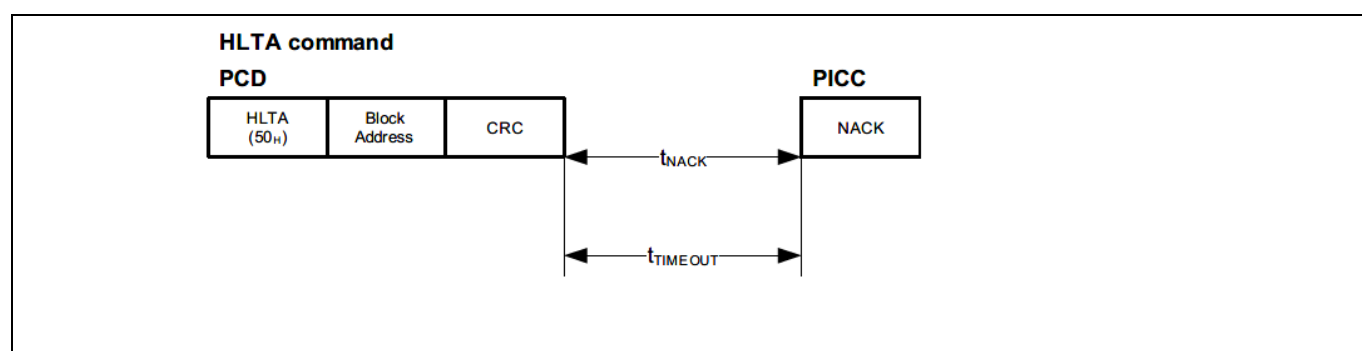
Figure 24   HLTA command

Table 33   Timing HLTA command

HLTA Command	$t_{ACK\ min}$	$t_{ACK\ max}$	$t_{NACK\ min}$	$t_{NACK\ max}$	$t_{TIMEOUT}$
			86.11 µs	t_{nonNVM}	t_{nonNVM} ≤ 950 µs

8 Performance and Operational Characteristics

The electrical characteristics ensure the operation of the SLE 66R35R/I over the listed range.

Typical characteristics specify mean values expected over the production spread. If not otherwise specified typical characteristics apply at $T_{ambient}$ = 25° C and the given supply voltage.

8.1 Electrical Characteristics

f_{CAR} = 13.56 MHz sinusoidal waveform, voltages refer to V_{SS}.

Table 34 Electrical Characteristics

Parameter	Symbol	Values			Unit	Note / Test Condition
		Min.	Typ.	Max.		
Chip input capacitance L_A – L_B	C_{IN}	16.5	18.3	20.1	pF	$V_{AB\,RMS}$ = 2.3 V, f_{CAR} = 13.56 MHz, $T_{ambient}$ = 25 °C Tolerance +/- 10%
Chip load resistance L_A – L_B	R_{IN}		17		kΩ	$V_{AB\,RMS}$ = 2.3 V, f_{CAR} = 13.56 MHz, $T_{ambient}$ = 25 °C Tolerance +/- 20%
Internally limited supply voltage	V_{DD}			3.3	V	Internal value
Supply current (idle mode)	I_{DD}		200	300	µA	Estimated value, not externally measured! f_{CAR} = 13.56 MHz, no modulation
Endurance (erase / write cycles)[1]		10^5				–
Data retention[1]		10			years	
EEPROM Erase and Write time	t_{prog}			5	ms	Combined erase + write; excluding time for command / response transfer between interrogator and chip, $T_{ambient}$ = 25 °C
ESD Protection voltage (L_A, L_B pins)	V_{ESD}	2			kV	EIA / JESD22 A114-B
Ambient temperature	$T_{ambient}$	-25		+70	°C	for chip
Junction temperature	$T_{junction}$	-25		+110	°C	for chip

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[1] Values are temperature dependent.

Performance and Operational Characteristics

8.2 Absolute Maximum Ratings

Stresses above the maximum values listed here may cause permanent damage to the device. Exposure to absolute maximum rating conditions for extended periods may affect device reliability, including EEPROM data retention and erase/write endurance. Maximum ratings are absolute ratings; exceeding only one of these values may cause irreversible damage to the integrated circuit. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this document is not implied.

Table 35   Absolute Maximum Ratings

Parameter	Symbol	Values			Unit	Note / Test Condition
		Min.	Typ.	Max.		
Input voltage between $L_A - L_B$	V_{INpeak}	-7		+7	V	On chip limitation by voltage regulator
Input current through $L_A - L_B$	I_{IN}			+50	mA	Maximal current
Storage temperature	$T_{storage}$	-40		+125	°C	

References

[1]   ISO/IEC 14443 Identification cards – Contactless integrated circuit(s) cards
        Proximity cards, Parts 1, 2 and 3

[2]   ISO/IEC 10373-6 Identification cards – Test methods – Proximity cards

Revision history

Reference	Description
Revision 2.0, 2021-05-28	
all	New document template
	Document classification changed
	Editorial changes
Revision 1.1, 2019-06-28	
All	Introduction of NRG terminology
Revision 1.0, 2010-01-11	
all	Initial version

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