

Q: Regard to the recurrence equation below:

$$T(n) = 4T(n/2) + n$$

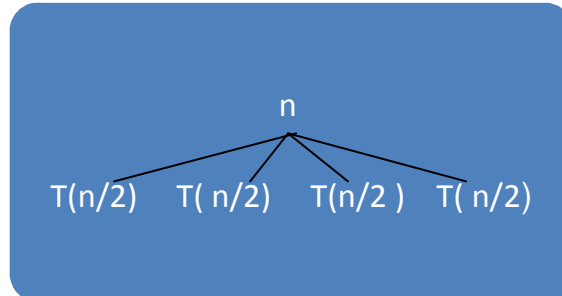
- a. Draw the recursion tree.**
- b. Determine the number of nodes and cost at each level.**
- c. Determine the total cost for all levels (i.e. in summation form).**

Steps:

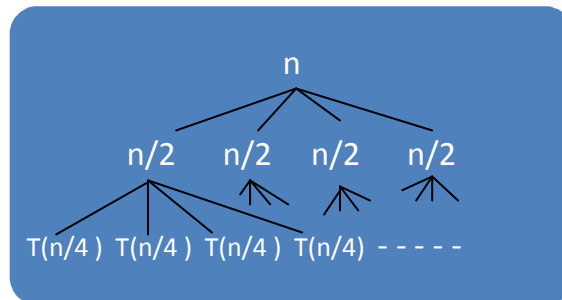
- 1. Draw the tree based on the recurrence**
- 2. From the tree determine:**
 - a. # of levels in the tree**
 - b. cost per level**
 - c. # of nodes in the last level**
 - d. cost of the last level (which is based on the number found in 2c)**
- 3. Write down the summation using Σ notation – this summation sums up the cost of all the levels in the recursion tree**
- 4. Recognize the sum or look for a closed form solution for the summation created in 3). Use Appendix A.**
- 5. Apply that closed form solution to your summation coming up with your “guess” in terms of Big-O, or Θ , or Ω (depending on which type of asymptotic bound is being sought).**

Solution :

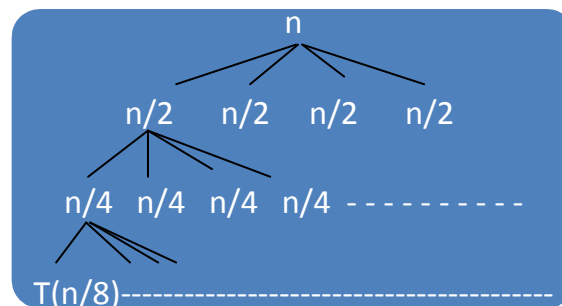
$$T(n) = 4T(n/2) + n$$

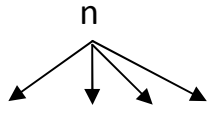
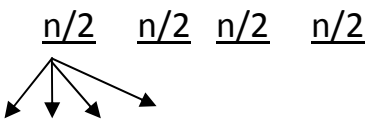
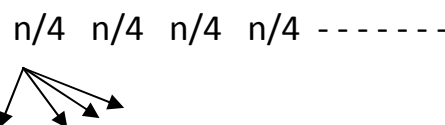
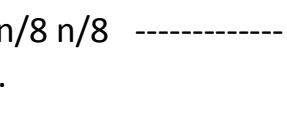


$$T(n/2) = 4T(n/4) + n/2$$



$$T(n/4) = 4T(n/8) + (n/4)$$



Level #	#of node	Recursion tree	Level Sum
0	1		n
1	4		$4(n/2) = 2n$
2	4^2		$4^2(n/4) = 4n$
3	4^3		$4^3(n/8) = 8n$
.	.	.	.
.	.	.	.
.	.	.	.
i	4^i	$n/2^i$	$4^i(n/2^i) = 2^i n$
.	.	.	$4^{h-1}(n/2^{h-1}) = 2^{h-1} n$
.	.	.	.
h	4^h	T(1) T(1) T(1) T(1)	4^h

$$T(n) = 4^h T(1) + \sum_{i=0}^{h-1} 2^i n$$

$$\frac{n}{2^h} = 1 \rightarrow h = \log_2 n \leftarrow \text{equation \#1}$$

from equation #1

$$T(n) = 4^{\log_2 n} T(1) + \sum_{i=0}^{h-1} 2^i n$$

$$T(n) = n^{\log_2 4} T(1) + \sum_{i=0}^{h-1} 2^i n$$

$$n^{\log_2 4} = n^2 \text{ since } \log_2 4 = 2 \leftarrow \text{equation \#2}$$

from equation #2

$$T(n) = n^2 T(1) + n \sum_{i=0}^{h-1} 2^i$$

$$T(n) = n^2 T(1) + n [1 + 2 + 2^2 + \dots + 2^{h-1}]$$

Increasing Geometric Series $r = 2 \quad a = 1$

$$S = a + ar + ar^2 + \dots + ar^{k-1} = \frac{a(r^k - 1)}{r - 1} \quad r > 1$$

Since $r = 2 \quad a = 1$

$$n [1 + 2 + 2^2 + \dots + 2^{h-1}] = \frac{1(2^h - 1)}{2 - 1} = 2^h - 1$$

from Equation #1

$$= 2^{\log_2 n} - 1$$

$$= n^{\log_2 2} - 1 = n - 1 \quad \text{Since } \log_2 2 = 1$$

$$T(n) = n^2 T(1) + n(n - 1)$$

$$T(n) = \theta(n^2)$$