

Miniboxing

Load-time Specialization on the JVM

DSLAB, EPFL, 19th of August 2013

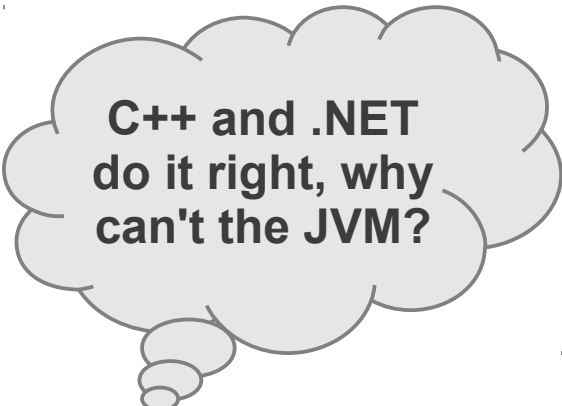


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Generic Code

```
def identity[T](t: T): T = t
```

- Reuse, type safety
- Uniform interface
- Low level code: **non-uniform**
 - Registers, stack
 - Sizes and semantics
 - Current solution on the JVM → **erasure**



C++ and .NET
do it right, why
can't the JVM?

```
identity(3) // = unbox_int(identity(box_int(3)))
```

Optimal Low Level Code

- C++
 - **Compile-time on-demand** template expansion
 - Undermined separate compilation
- .NET CLR
 - **Load-time on-demand** bytecode specialization
 - Type system burnt in the virtual machine
- Specialization in Scala
 - **Compile-time static** bytecode specialization

What's the difference?

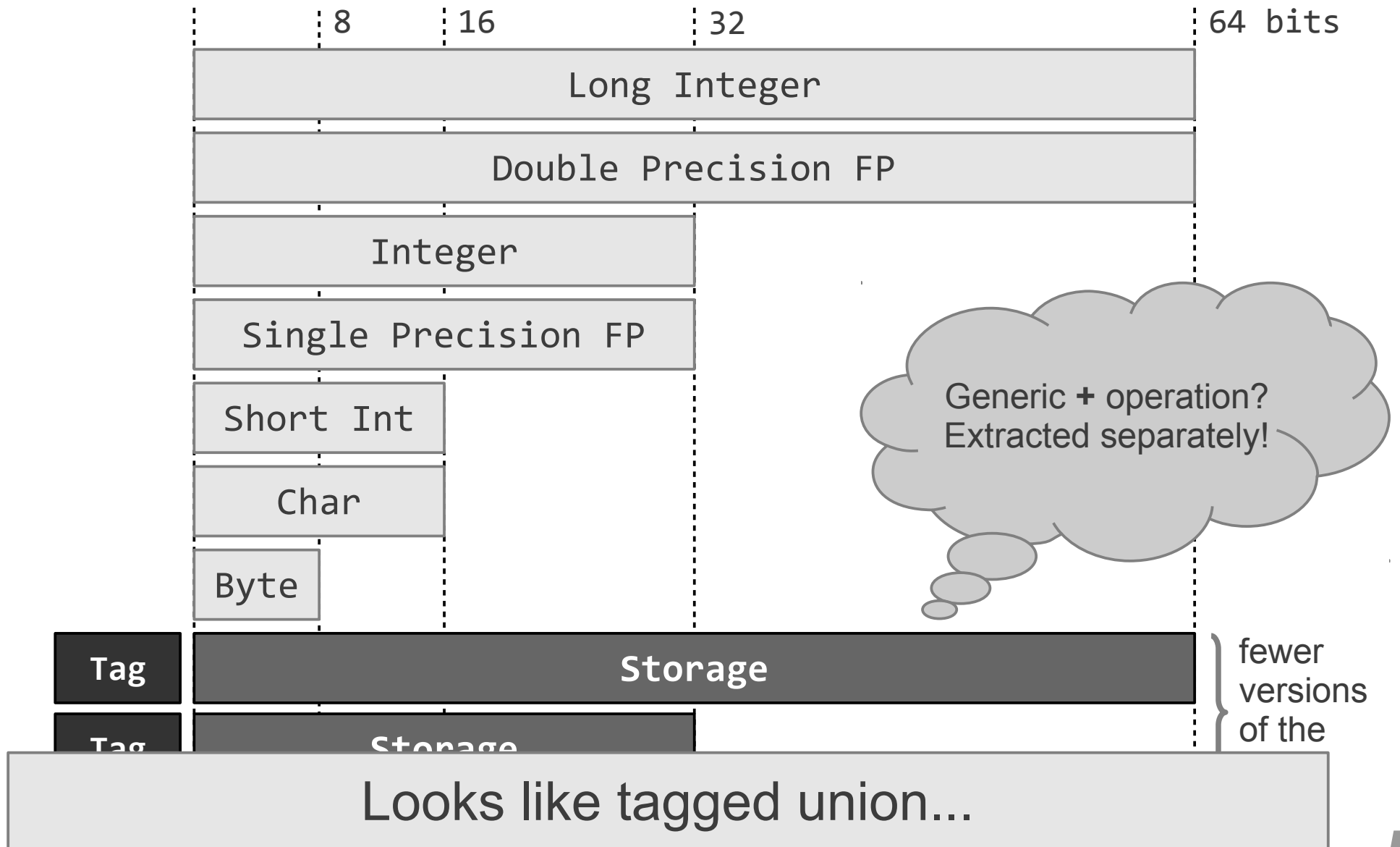
Specialization

```
def identity[T](t: T): T = t
```

```
def identity(t: Any): Any = t  
def identity_D(t: Double): Double = t  
def identity_J(t: Long): Long = t  
def identity_I(t: Int): Int = t  
def identity_F(t: Float): Float = t  
def identity_S(t: Short): Short = t  
def identity_C(t: Char): Char = t  
def identity_B(t: Byte): Byte = t  
def identity_Z(t: Bool): Bool = t  
def identity_U(t: Unit): Unit = t
```


```
def apply[T, R](arg: T): R // 102 versions
```

Miniboxing




But It's Not

Tagged Union

- a single version of the code, uniform
 - dynamic
 - attaches type tags to each value
- 
- dispatch overhead
 - no optimizations across dispatches

Miniboxing

- multiple versions of the code, as specialization
 - static
 - tags attached to classes and methods
- 
- eliminates dispatch overhead
 - optimizations across

What's the magic?

Dispatch Overhead

- `ArrayBuffer.reverse()`

```
def reverse(): Unit {  
  var index = 0  
  while (index * 2 < length) {  
    val opposite = length-index-1  
    val tmp1: T = array(index)  
    val tmp2: T = array(opposite)  
    array(index) = tmp2  
    array(opposite) = tmp1  
    index += 1  
  }  
}
```

`T$type match {
 case INT => ...
 ...
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Lift the `T$type match { ... }` above the loop?

Object-Oriented Dispatching

- Dispatch object
 - Encodes array interactions

```
class Dispatcher[T] {  
  def array_get(arr, idx): Storage  
  def array_set(arr, idx, value: Storage): Unit  
}
```

```
object IntDispatcher extends Dispatcher[Int] {  
  ...  
}
```

Passing a dispatcher = switch already done

Object-oriented Dispatching

- `ArrayBuffer.reverse()`

```
def reverse(): Unit {  
  var index = 0  
  while (index * 2 < length) {  
    val opposite = length-index-1  
    val tmp1: T = array(index)  
    val tmp2: T = array(other)  
    array(index) = tmp2  
    array(opposite) = tmp1  
    index += 1  
  }  
}
```

`T$Dispatcher.array_get`

`T$Dispatcher.array_get`

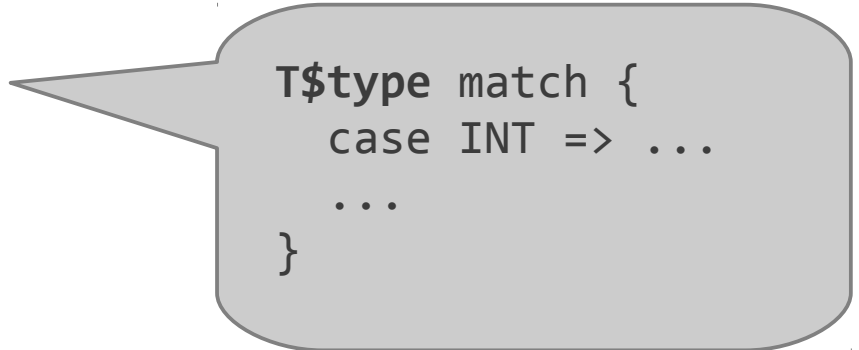
`T$Dispatcher.array_set`

`T$Dispatcher.array_set`

Call site becomes megamorphic

Load-time Specialization

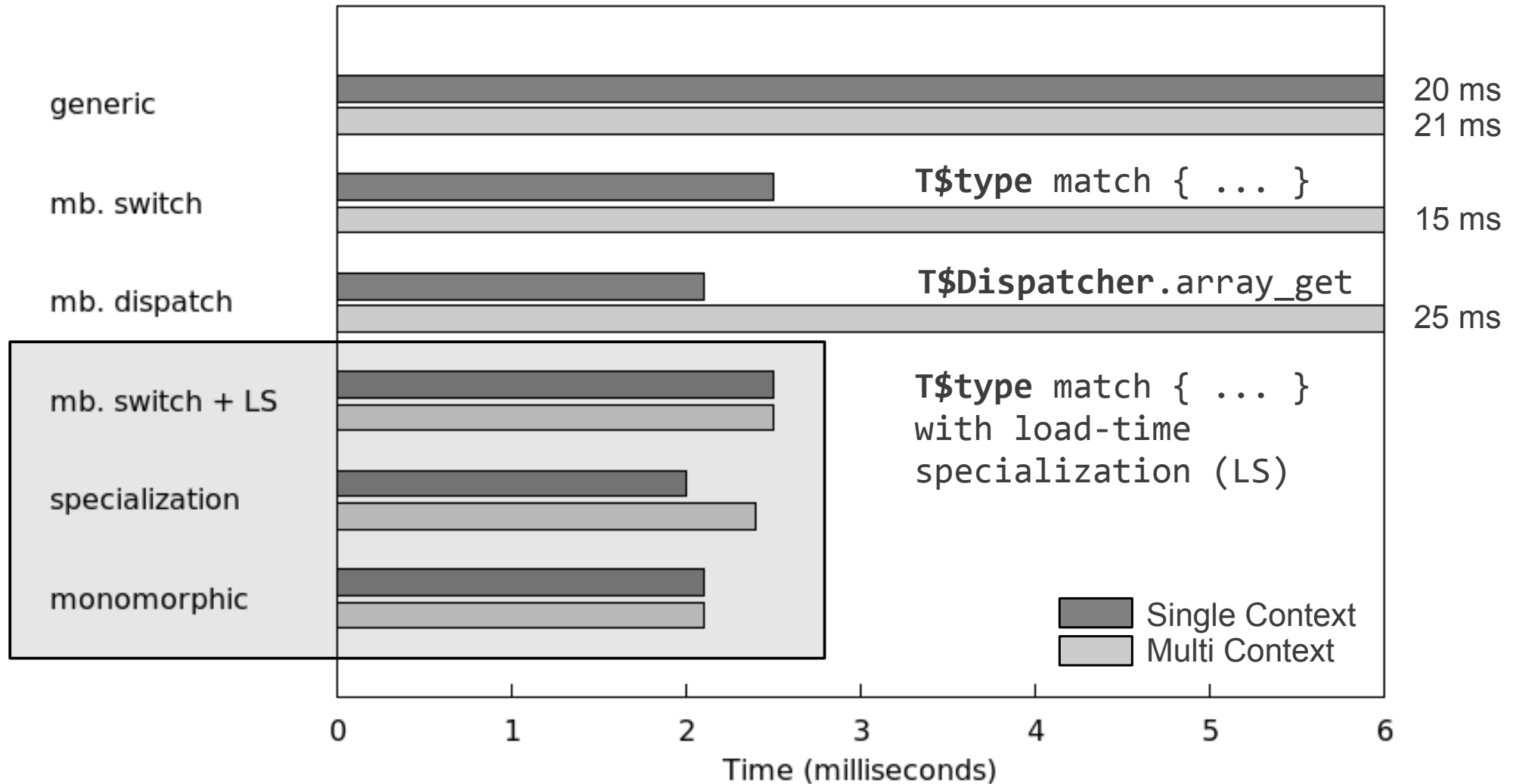
- Switch-based dispatching
- At load-time
 - set **T\$type** statically
 - perform constant propagation
 - perform dead code elimination
- We get the specialized class
 - no dispatch overhead
 - optimizations can kick in across dispatches



```
T$type match {  
  case INT => ...  
  ...  
}
```

Does it work in practice?

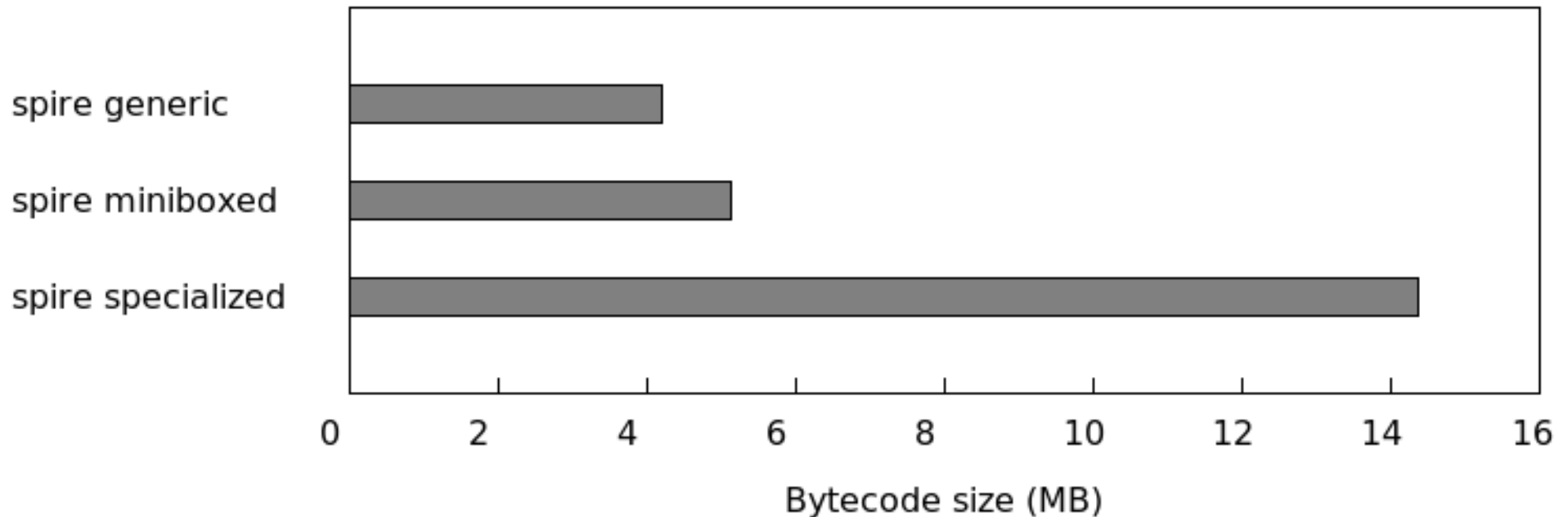
Evaluation – reverse on 3M elements



Similar results on other benchmarks

Evaluation – Code Size

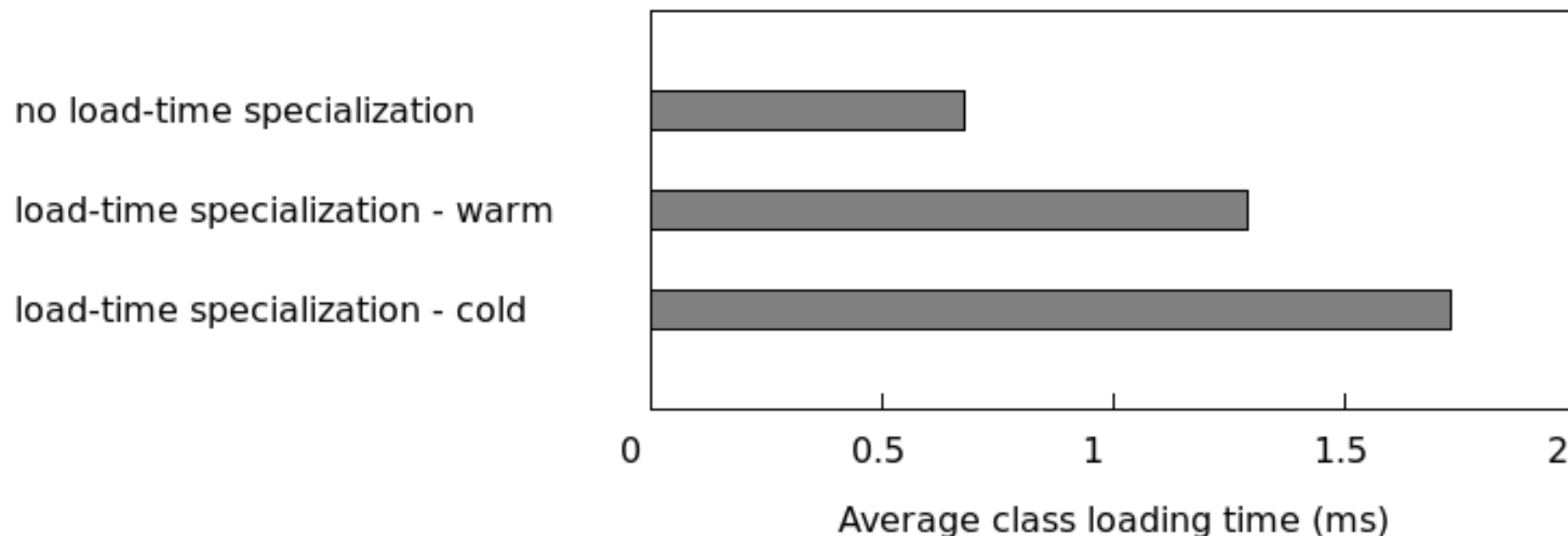
- Spire – numeric abstractions library (12KLOC)



2.8x bytecode reduction (4.7x for Vector in std. lib)

Evaluation – LS Overhead

- `scala.collection.immutable.Vector`



Load-time specialization overhead – at most 2.5x

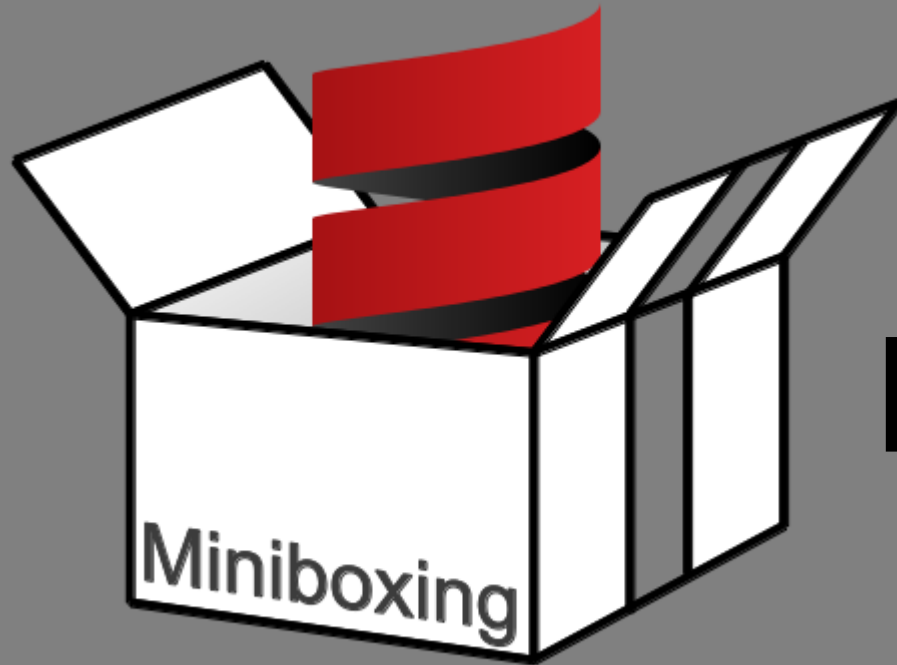
Conclusions

- Miniboxing – same performance, less bytecode
 - Encoding
 - Transformation
 - Runtime Support
 - Load-time Specialization
- Virtual machines and performance
 - **High-level self-modifying code**



Thank you!

github.com/miniboxing
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WE ARE HERE

Backup slides

Generic Code

```
def identity[T](t: T): T = t  
identity(3) // = unbox_int(identity(box_int(3)))
```

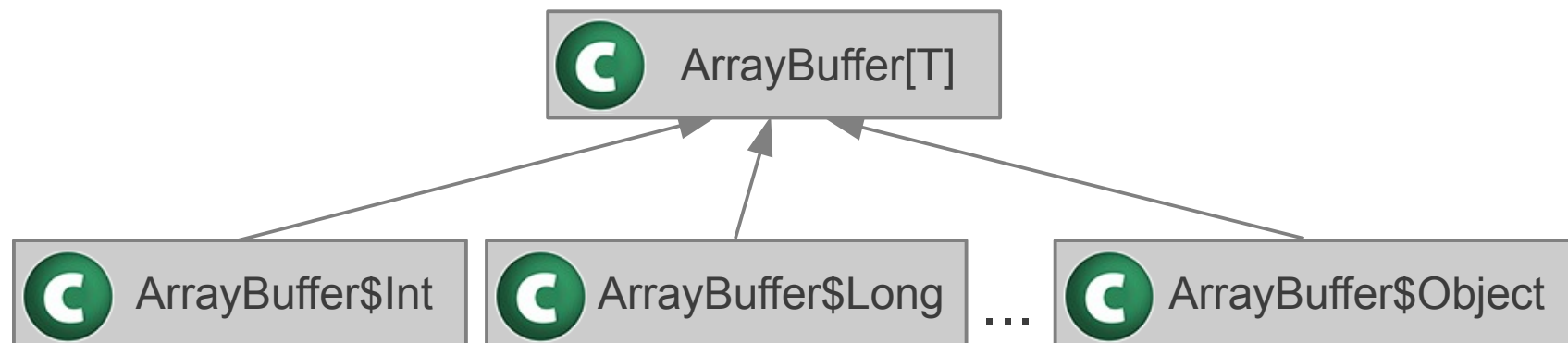
- common representation – boxing?
 - **allocation** and **garbage collection**
 - **wasteful** use of heap memory
 - affects **cache locality**

Specialization

- duplicate and adapt methods

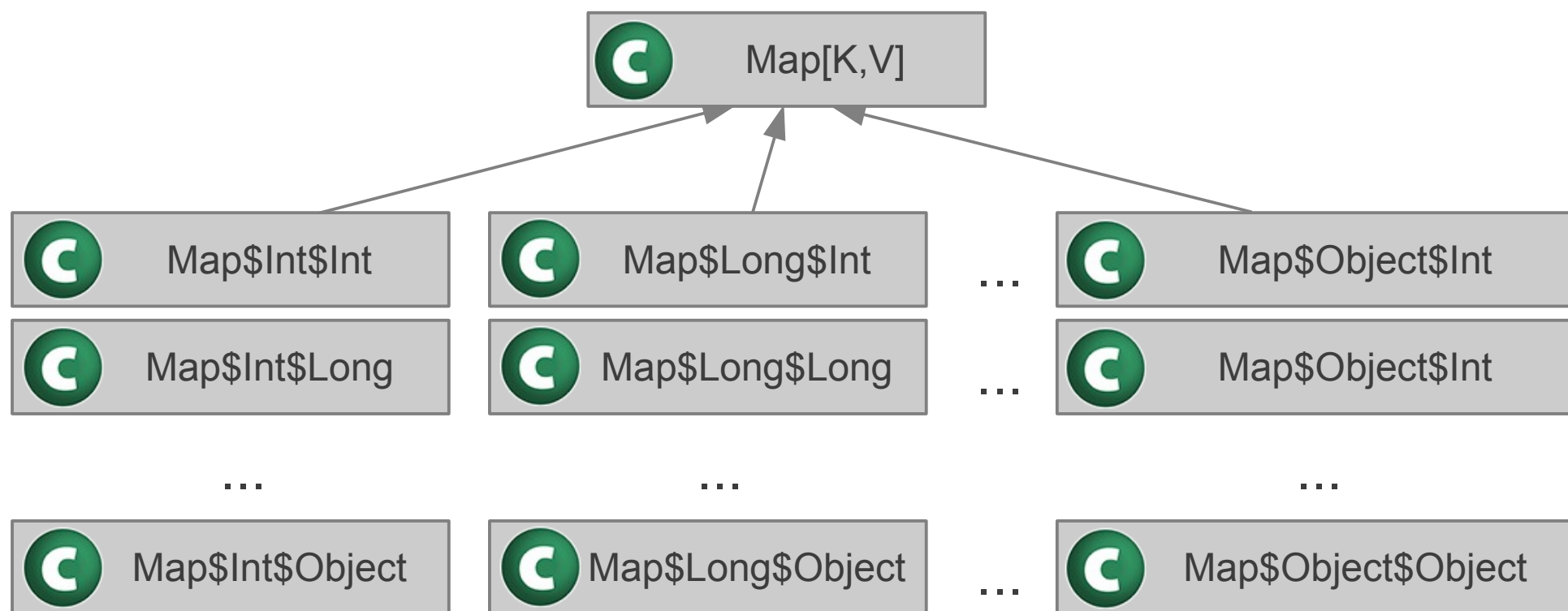
```
def identity[T](t: T): T = t
def identity$Int(t: Int): Int = t
...
```

- duplicate and adapt classes



Specialization - Scaling

- static expansion, in the library



Specialization - Compile-time

- Link-time
 - Could be seen as JAR assembly
 - Not very common, people prefer mix-n-matching
- Load-time
 - Storing ASTs or source code
 - Compilation is a heavyweight process
 - especially for scalac
 - **Need a lightweight solution**

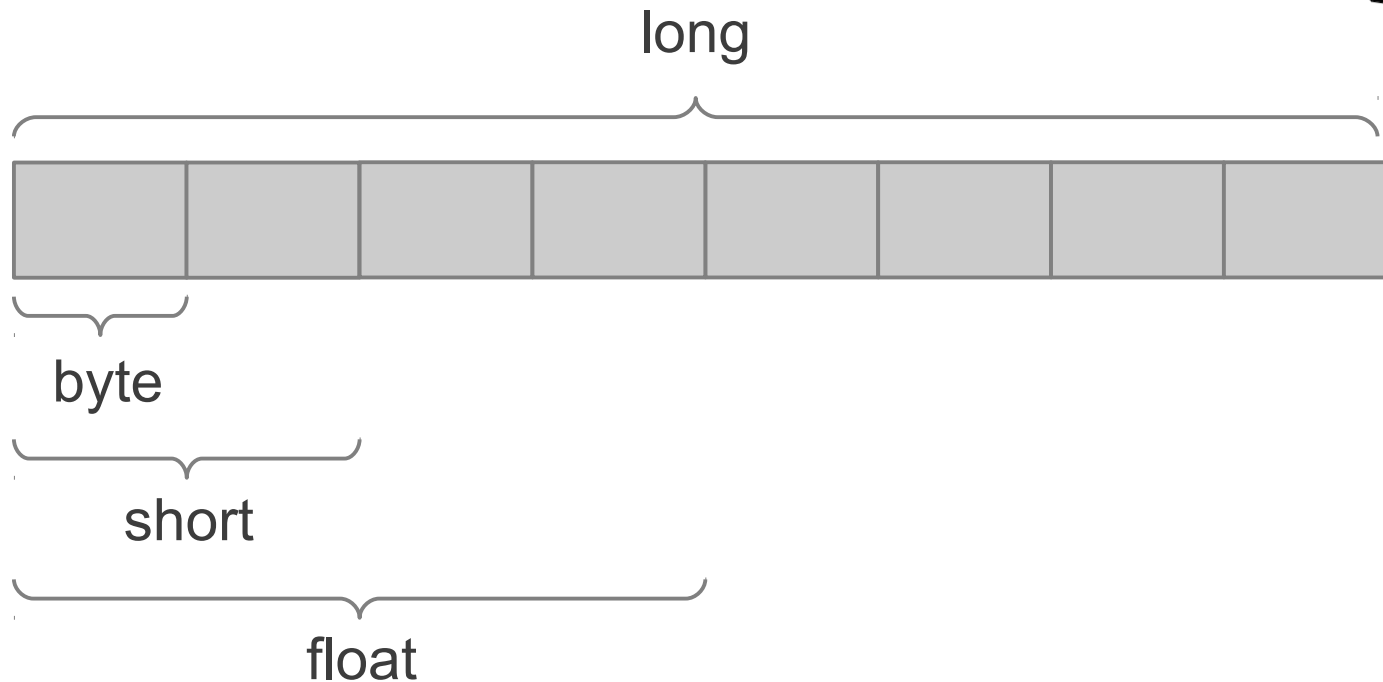
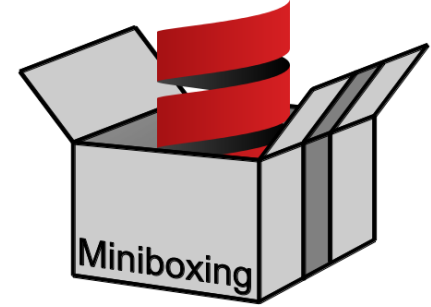
Specialization - Static

- On-demand
 - Suppose we attached the AST or the source code
 - And expanded it at compile-time
 - Multiple versions of common classes:
 - `Function[Int, Int]`
 - in the standard library
 - in Akka and other middleware
 - in your application
 - We can't perform de-duplication like C++ linkers do
 - since few people do JAR assemblies

Specialization – On source code

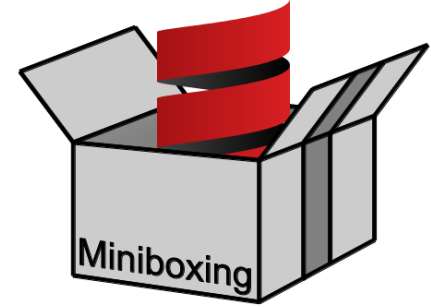
- On bytecode
 - Complex machinery, would prefer a higher level
 - Full Specialization would be easy
 - But it requires reified types
 - And those are slow*
 - Opportunistic specialization
 - Compatibility is very complex to maintain

Miniboxing - Insight

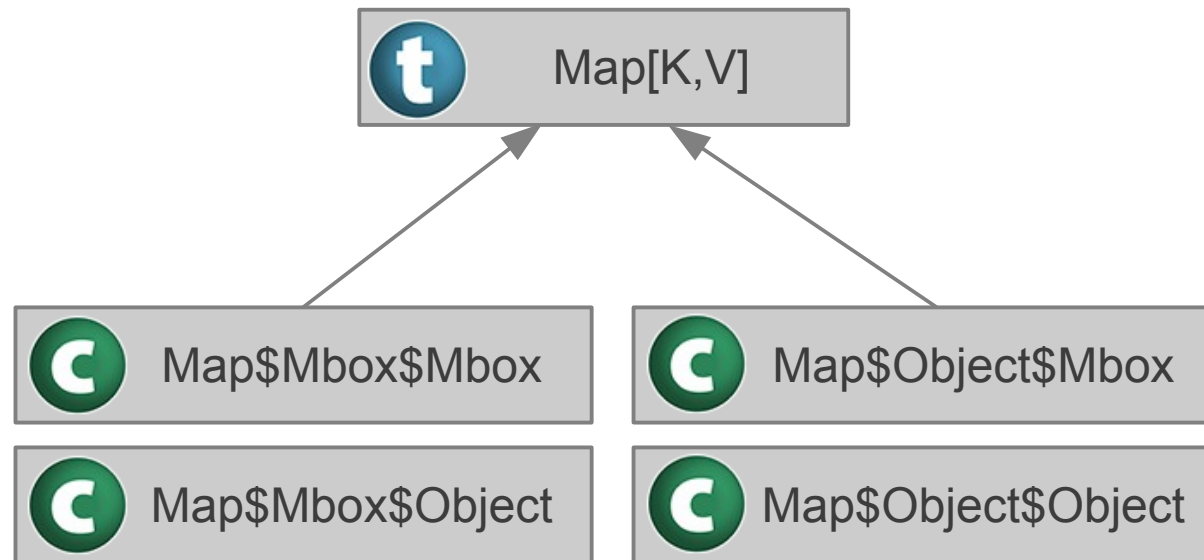


- Only applied to stack values
- We'll look at arrays separately

Miniboxing - Scaling

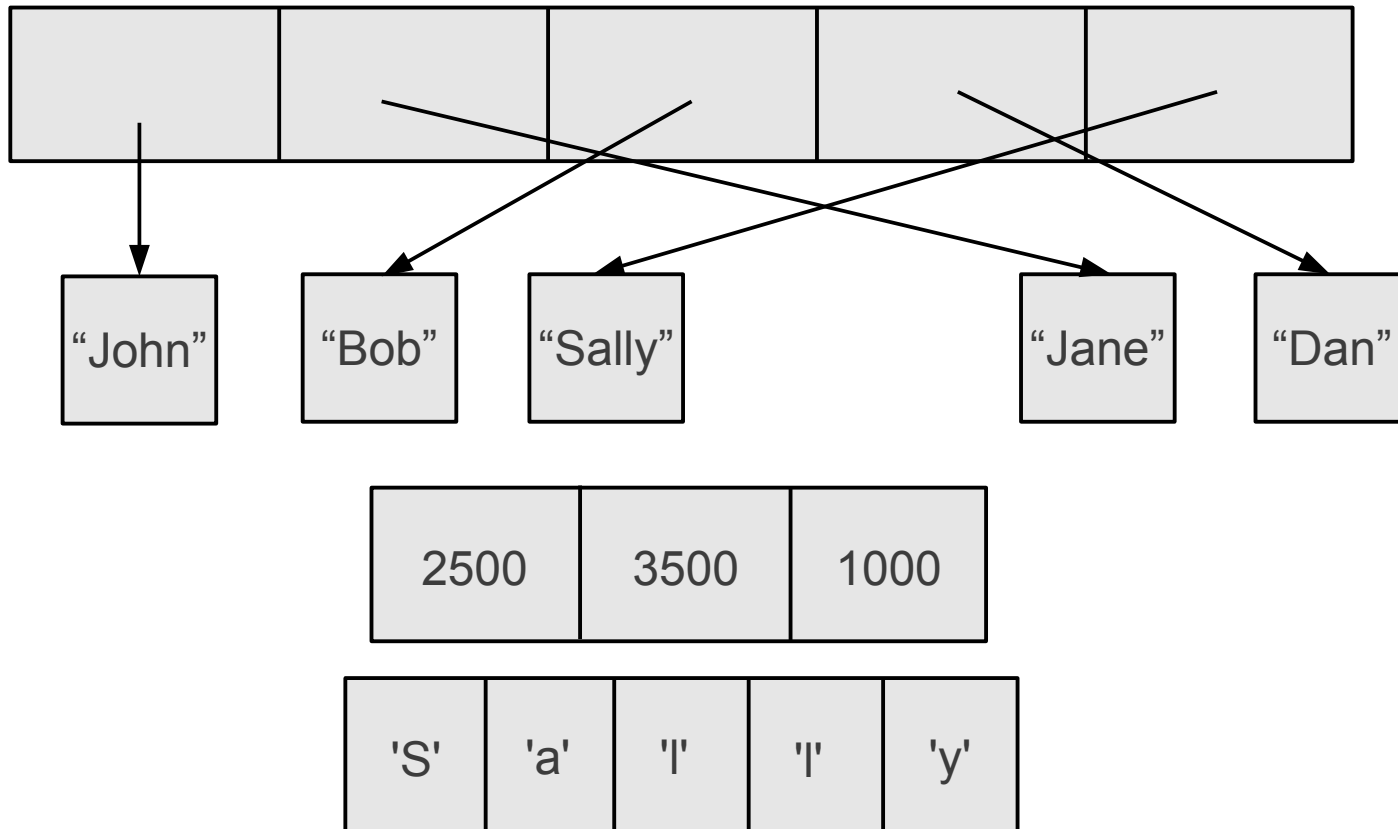


- static expansion, in the library



Arrays

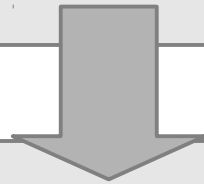
- Monomorphic
 - `Array[String]`, `Array[Int]`, `Array[Float]`



Library Support

- Executing code on value types (.hashCode...)
- Interfacing with arrays

`array(i)`



```
T$type match {  
  case INT => array.asInstanceOf[Array[Int]](i).toLong  
  ...  
}
```

Casting and converting to long is fast, noop-fast

Library Support

- `ArrayBuffer.reverse()`

```
def reverse(): Unit {  
  var index = 0  
  while (index * 2 < length) {  
    val opposite = length-index-1  
    val tmp1: T = array(index)  
    val tmp2: T = array(opposite)  
    array(index) = tmp2  
    array(opposite) = tmp1  
    index += 1  
  }  
}
```

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Can the JVM optimize such code?

Java Virtual Machine Support

- We confuse the JVM heuristics by inlining manually
- The JVM will do its best
 - Using a single value of T\$Type (one type)
 - hoist the switch outside the loop (loop unswitching)
 - get rid of multiple casts (common subexpression elim)
 - get rid of multiple bounds checks (check elimination)
 - **good speed**
 - Using more types...
 - can't hoist the switch outside → bad performance
 - **we need to optimize this code too**

Can we lift the switch ourselves?

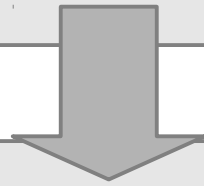
Dispatching

- Systems approach: T\$type
 - byte value
 - encodes the type of a parameter (INT, ...)
 - switch to do the right operation
- Object-oriented approach
 - common operations object (Dispatcher)
 - specialized instances for all types
 - a type class for Haskell folks

Object-oriented Dispatching

- T\$Dispatcher: Dispatcher
 - IntDispatcher, LongDispatcher, ...
- Common operators

array(i)



T\$Dispatcher.array_get(array, i)

So far so good

Object-oriented Dispatching

- `ArrayBuffer.reverse()`

```
def reverse(): Unit {  
  var index = 0  
  while (index * 2 < length) {  
    val other = length-index-1  
    val tmp1: T = array(index)  
    val tmp2: T = array(other)  
    array(index) = tmp2  
    array(other) = tmp1  
    index += 1  
  }  
}
```

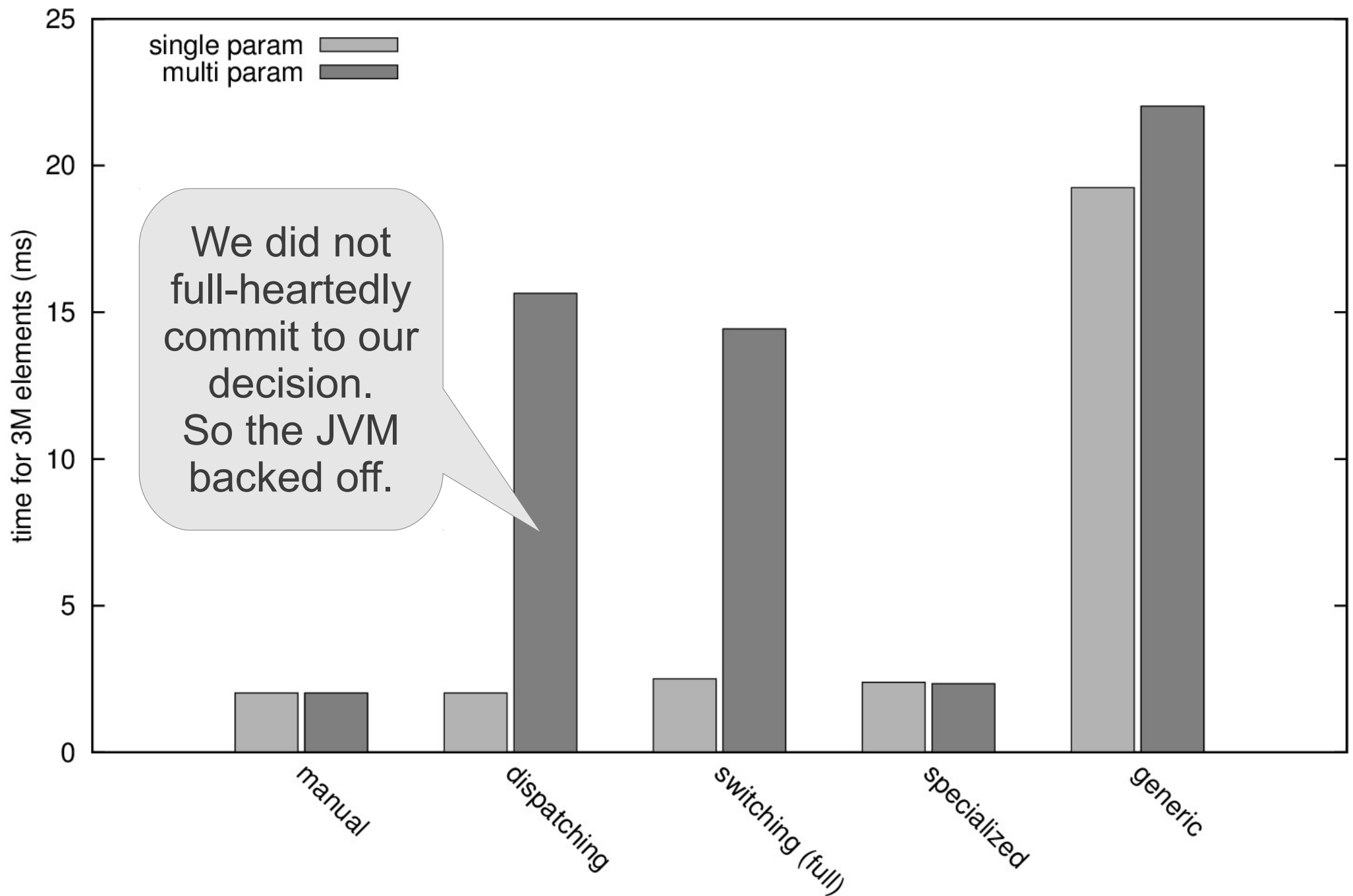
`T$Dispatcher.array_get`

`T$Dispatcher.array_get`

`T$Dispatcher.array_update`

`T$Dispatcher.array_update`

We committed to the data representation before



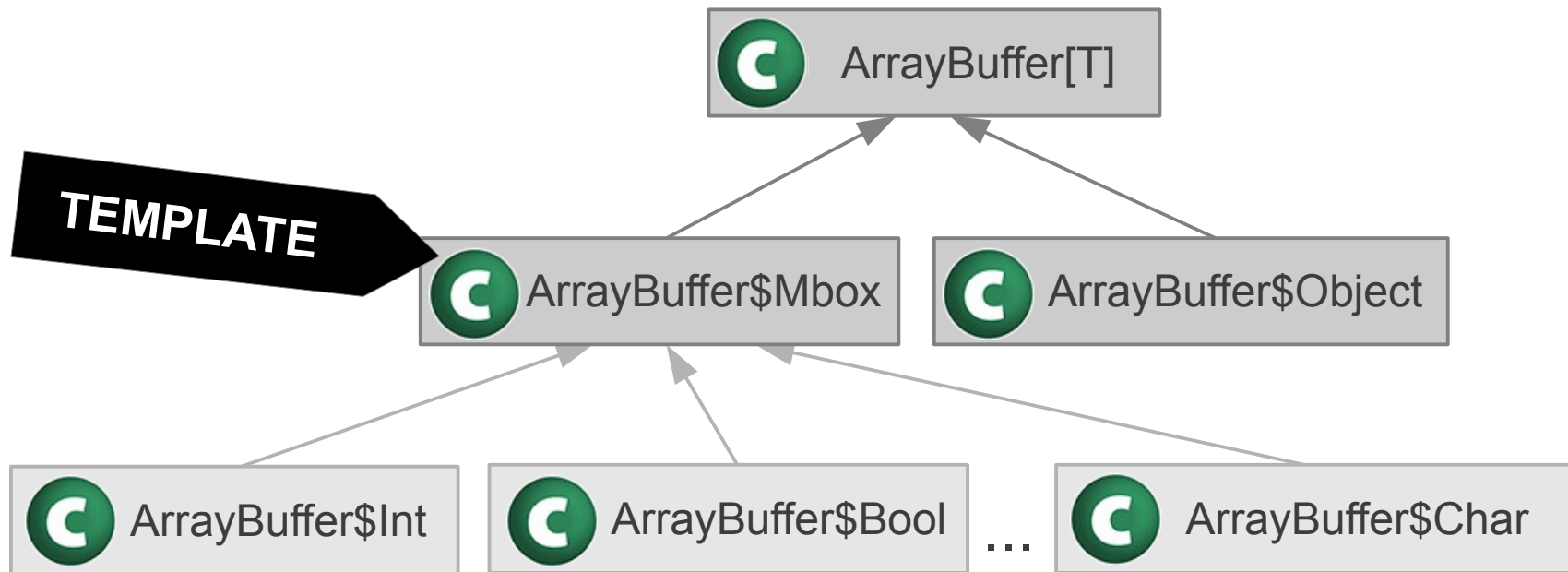
`ArrayBuffer.reverse()`

OO Dispatching Performance

- One or two dispatchers in a callsite
 - as fast as manually specialized code
- Three+ dispatchers
 - call to `array_get` becomes **megamorphic**
 - it is not inlined anymore
 - no optimization in the loop
 - awful results, in some cases slower than generic

Runtime Specialization

- In practice, dispatcher set per INSTANCE
- Semantically - dispatcher per SPECIALIZATION
- Remember the ArrayBuffer diagram?



Runtime Specialization

- `ArrayBuffer$Int`
 - has `IntDispatcher` set statically
- Operations can be inlined
 - since `IntDispatcher` is final
 - calls never become megamorphic
 - optimizations can be done inside the loop

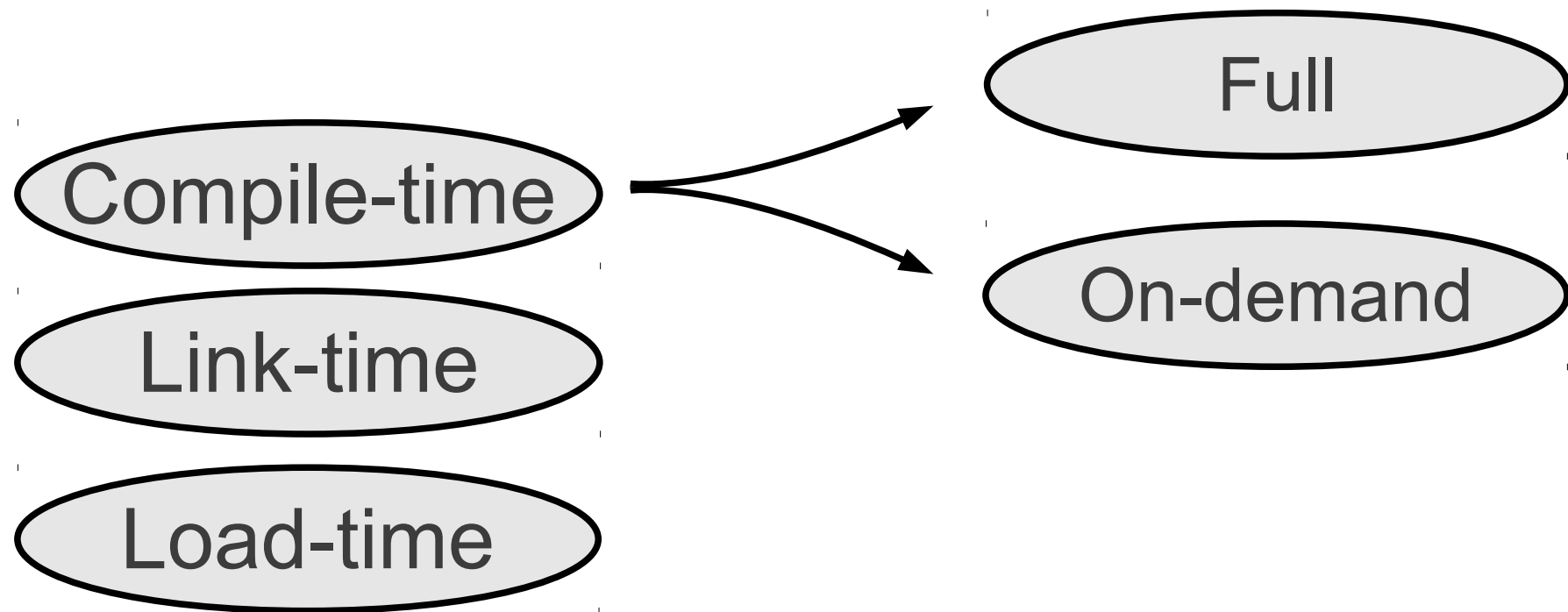
Class modification and loading cost?

Numbers

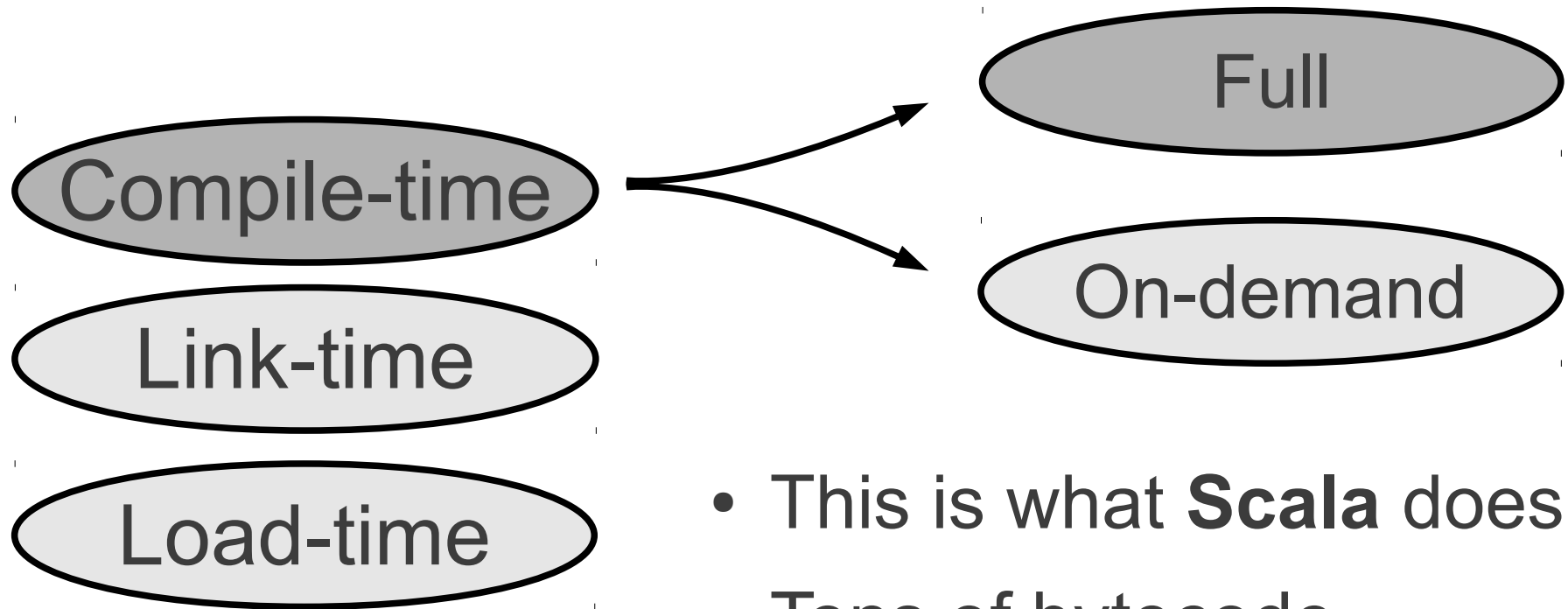
- `ArrayBuffer.reverse()`

ideal :	0.85883 ms +/-	0.05687
miniboxed switching mono :	0.62966 ms +/-	0.21116
miniboxed switching mega :	4.21423 ms +/-	1.15425
miniboxed switching w/cl :	0.87526 ms +/-	0.05992
miniboxed dispatch mono :	0.75230 ms +/-	0.21551
miniboxed dispatch mega :	4.90191 ms +/-	1.52366
miniboxed dispatch w/cl mono :	0.77984 ms +/-	0.14872
miniboxed dispatch w/cl mega :	0.80932 ms +/-	0.14671
specialized mono :	0.75813 ms +/-	0.16000
specialized mega :	0.61374 ms +/-	0.20144
generic :	9.76250 ms +/-	1.55297

Specialization in different languages

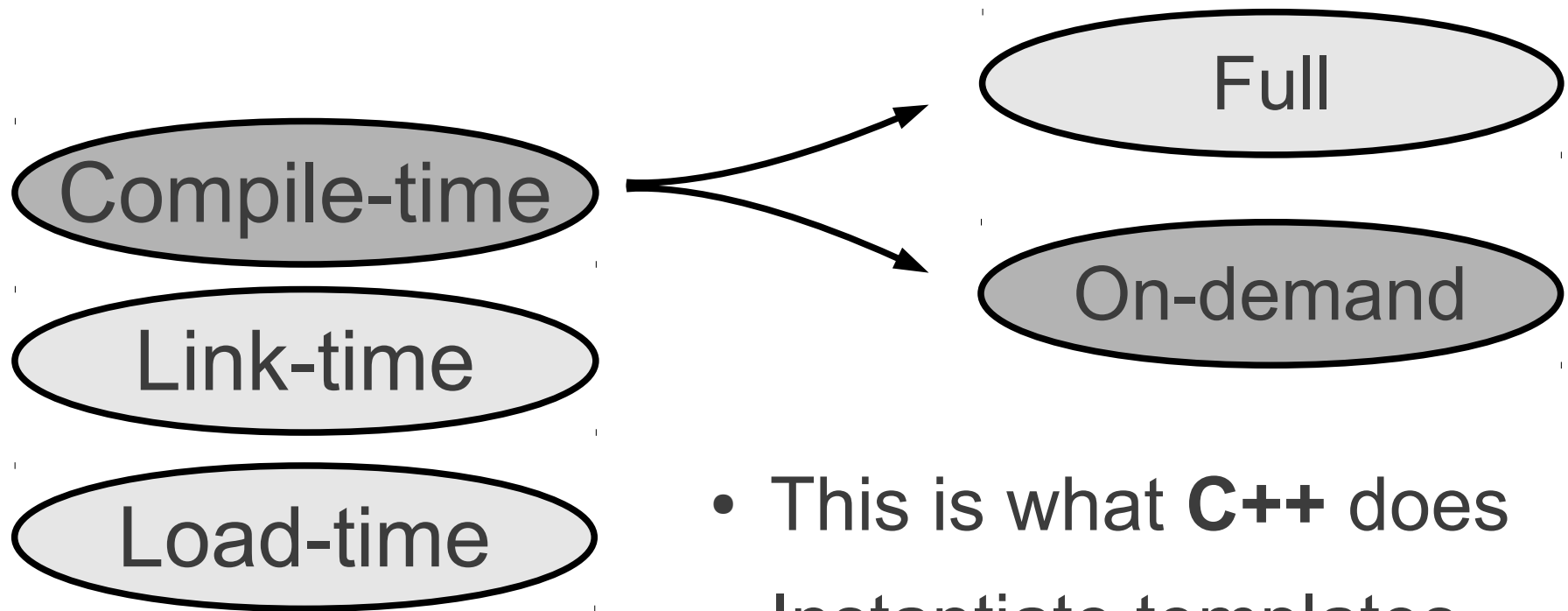


Specialization in different languages



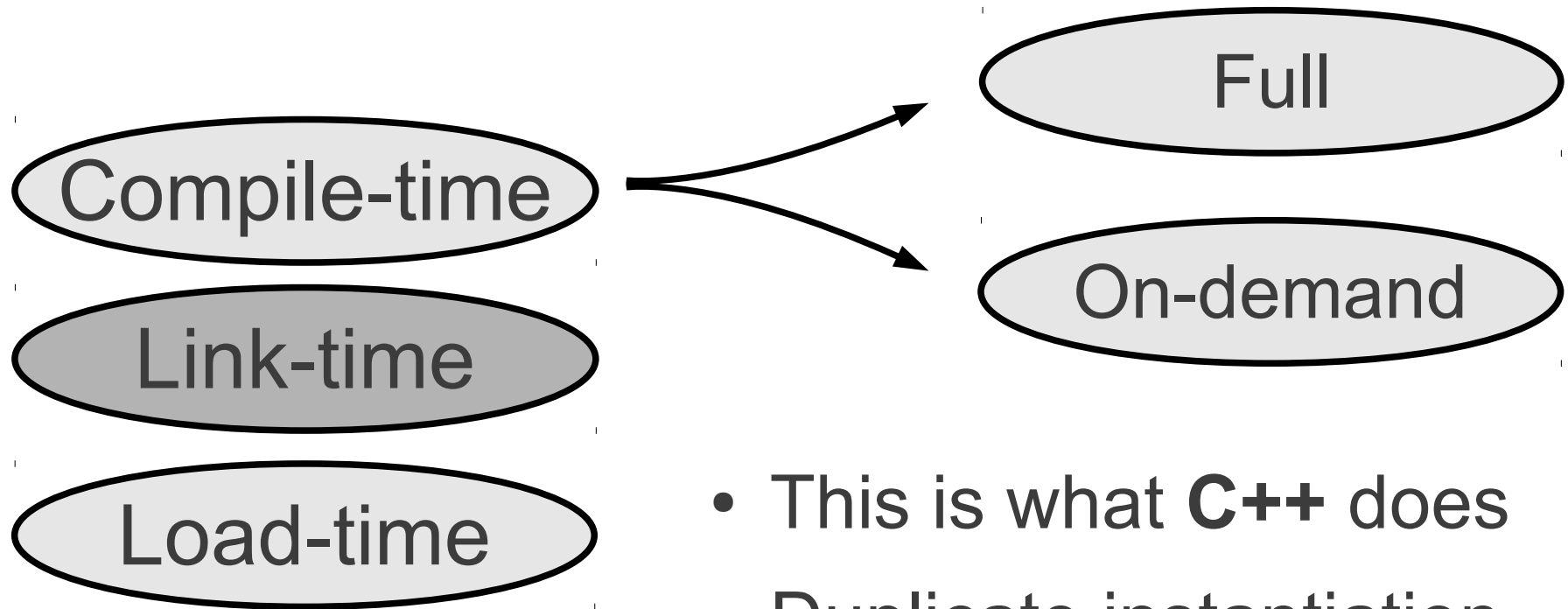
- This is what **Scala** does
- Tons of bytecode

Specialization in different languages



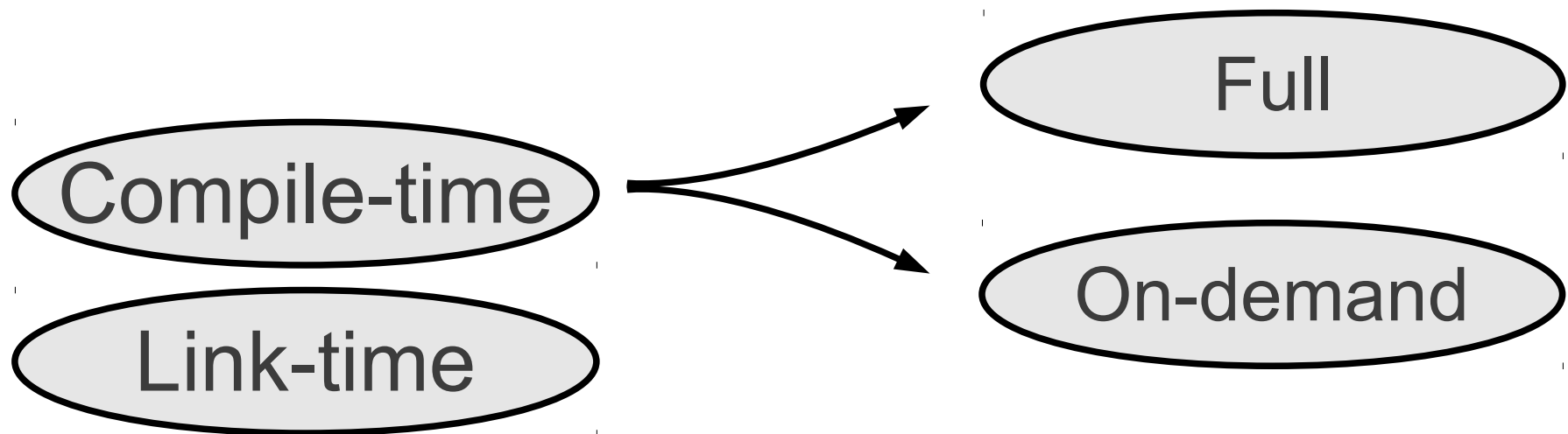
- This is what **C++** does
- Instantiate templates
 - Libraries are mostly headers
- Duplicate instantiation

Specialization in different languages



- This is what **C++** does
- Duplicate instantiation
- Cleaned up by the linker

Specialization in different languages

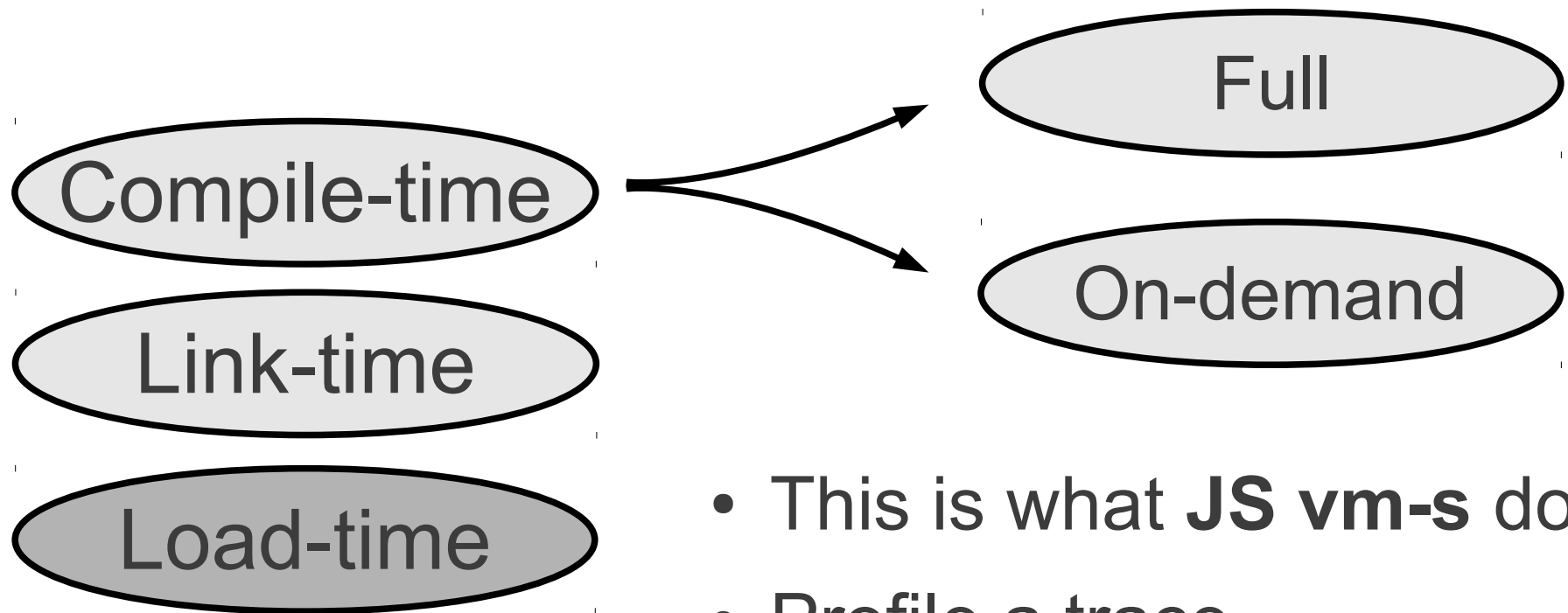


- This is what **.NET** does
- Runtime specialization
 - Based on templates
- Type system fixed
 - in the virtual machine

One more ACE up its sleeve:

Reified types

Specialization in different languages



- This is what **JS vm-s** do
- Profile a trace
- Specialize it
- Fallback to interpretation

In perspective

