

Chapter 2

C Pointer Arithmetic and Corresponding Assembly Code

Variable, Constant, and Variable Name

- Variable Declarations
- `int x;`
 - define integer **variable** “x” (allocate memory area to store an integer value and name the area “x” – associate label “_x” (“x”) with the address of the area)
 - x (in LHS): denotes that the target variable of the assignment is “x”
 - `x = 3;` // 3 is stored in variable “x”
 - x (in RHS): denotes the value of variable “x” (the contents of memory area named “x”)
 - `&x`: denotes the address of variable “x” (the address of the memory area named “x”)



Variable, Constant, and Variable Name (cont)

- `int *px;`
 - define pointer to integer **variable** “px” (allocate memory area to store a pointer to integer and name the area “px”)
 - px (in LHS): denotes that the target variable of the assignment is “px”
 - `px = &x;` // the address of variable “x” is stored in variable “px”
 - px (in RHS): denotes the value of variable “px” (the contents of the memory area named “px” --- the address of variable “x”)
 - `&px`: denotes the address of variable “px” (the address of the memory area named “px”)
 - `*px` : **dereference** of px
 - `*px` (LHS)
 - In this case, it assigns the RHS value in the memory area addressed by the value of “px”. However, it is not always so
 - `*px` (RHS)
 - In this case, it denotes the contents of the memory area addressed by the value of “px” . However, it is not always so



Variable, Constant, and Variable Name (cont)

- 3
 - denotes constant 3
 - no memory area is allocated; in common implementations, value 3 is embedded in an machine instruction (immediate addressing))
 - therefore, both “3 = x” and “&3” are illegal



Pointers and arrays

- `int a[10];`
 - reserve a memory area to store 10 integers and name the starting address of the area “a”
 - “a” is a **constant** and there is no memory area named “a” (therefore, `&a` is illegal)
 - try `printf(“%d, %d\n”, a, x);`
 - the type of “a” is “pointer to integer”
 - “a[i]” is a variable (the contents of the ith integer area in the allocated area)
 - “&a[i]” is the address of variable “a[i]”
- C/C++ do not have an array (arithmetic)
 - exception : declaration of an array, such as “`int a[10];`”, that allocates a memory area
- “a[i]” is shorthand for “*(a+i)” (“a” is a pointer to [some object], ” i” is an integer)
 - when the compiler encounters “a[i],” it executes “*(a+i)”
 - you can write “i[a]” in place of “a[i]”
 - `(a[i] == *(a + i) == *(i+a) == i[a])`
 - after declaring “`int *pa=a;`”, you can write “pa[i]” to access the same area as “a[i]”



Pointer declaration: the right spiral rule

- Declaration of pointers
 - types that appear in the right side of the variable name
 - `[] [] ... []` : array `[] [] ... []` of (read “`[] [] ... []`” at once)
 - `(...)` : function that takes ... and returns
 - types that appear in the left side of the variable name
 - `*` : pointer to
- interpretation of a pointer declaration
 1. find the variable name and start from there
 2. draw a half right arc to the first C symbol and read it
 3. draw a half right arc to find the next C symbol and read it
 - repeat Step 3 until C data type (int, float, etc) is found



Pointer declaration: the right spiral rule (cont)

- `char *x[3];`
 - x is of type “array [3] of pointer to char”
- `char (* x) [2];`
 - x is of type “pointer to array [2] of char”
- `char * (* (x [])) ();` `[== char * (* x []) ();]`
 - x is of type “array [] of pointer to function that returns pointer to char”
- `int * (* (* x [2][3]) [4]) ();`
 - x is of type “array[2][3] of pointer to array[4] of pointer to function that returns pointer to int”



Pointer arithmetic

- Consider the following type declaration

```
TYPE *ptr;      // “ptr” is a pointer to an instance of TYPE (“ptr” is a variable)
TYPE a[10];     // “a” is a pointer to an instance of TYPE (“a” a constant)
```

- Only two types of operations may be applied to pointer variables/constants
 - pointer + integer (similarly, pointer – pointer)
 - *pointer // dereferencing
- Type “array[i][j][k]...[z]” is equivalent to “pointer to array[j][k]....[z]”
 - since there is no “array arithmetic in C”, when the type under analysis is [i][j][k]...[z], it must be converted to “pointer to array[j][k]...[z]” (**RULE1**)



Pointer arithmetic (cont)

Apply the following derivation rules using “T” (type) and “V” (value) on an expression containing a pointer

1. pointer+integer

ptr: T suppose “pointer to TYPE”
 V suppose “VAL”

then

ptr + i (i is an integer) is

 T pointer to TYPE

 the type does not change when an integer is added

 V $VAL + i * \text{sizeof}(\text{TYPE})$

 when added 1, ptr addresses the next element in “the array”



Multi-dimensional arrays

Consider declaration “int a[3][4];”

- a memory area for 3*4 integers is allocated
- “a” is a constant and its value is the starting address of the allocated area
- a[0][0] is a variable and denotes the value of the element in the 0th (the first) row and the 0th column of the array
- then, what is “a[0]” ? Execute the following program:

```
int a[3][4]={1,2,3,4,5,6,7,8,9,10,11,12};
int main( ) {
    printf("%x\t%x\t%x\n", a, a[0], a[0][0]);
    printf("%x\t%x\t%x\n", a+1, a[0]+1, a[0][0]+1);
    printf("%x\t%x\t%x\n", a+1, *a+1, **a+1);
    return 0;
}
```

output:

```
4227136 4227136 1
4227152 4227140 2
4227152 4227140 2
```



Multi-dimensional arrays (cont)

- from the execution, what do you observe ? what is the difference between “a” and “a[0]” ?
- Clearly understand that “ *pointer” is not necessarily the contents of the location addressed by “pointer”



Pointer arithmetic (important)

2. *pointer (dereference: remove the leftmost “pointer to” in the type)

1. Case 1:

- ptr: T “pointer to a simple (non-array) object type”
V suppose “VAL”
- *ptr T the simple object type
V the contents of the memory area addressed by “VAL”

2. Case 2:

- ptr: T “pointer to [i][j][k]...[z] of object type”
(an array object type)
V suppose “VAL”
- *ptr T [i][j][k]...[z] of object type
= pointer to [j][k]...[z] of object type
(C has “pointer arithmetic” but does not have “array arithmetic”; when the array notation appears in the derivation, transform it to the equivalent pointer notation)
V VAL (the same value as ptr)
if the pointer points to an array object, the value does not change after dereferencing



Pointer representation and array representation

- $a[i] = *(a + i)$, where a is a pointer (to some object) and i is an integer
- $a[i][j] = (a[i])[j] = *(a[i] + j) = (*(a + i) + j)$
- $a[i][j][k] = ((a[i])[j])[k] = *((a[i])[j] + k) = (*(a[i] + j) + k) = (*(a + i) + j) + k)$
- Thus, to analyze $a[i][j][k]$, apply the “+int” and “* (dereference)” rules in the following order:
 1. a
 2. $a+i$
 3. $*(a+i)$ ($= a[i]$)
 4. $a[i] + j$
 5. $*(a[i] + j)$ ($= a[i][j]$)
 6. $a[i][j] + k$
 7. $*(a[i][j] + k)$ ($= a[i][j][k]$)



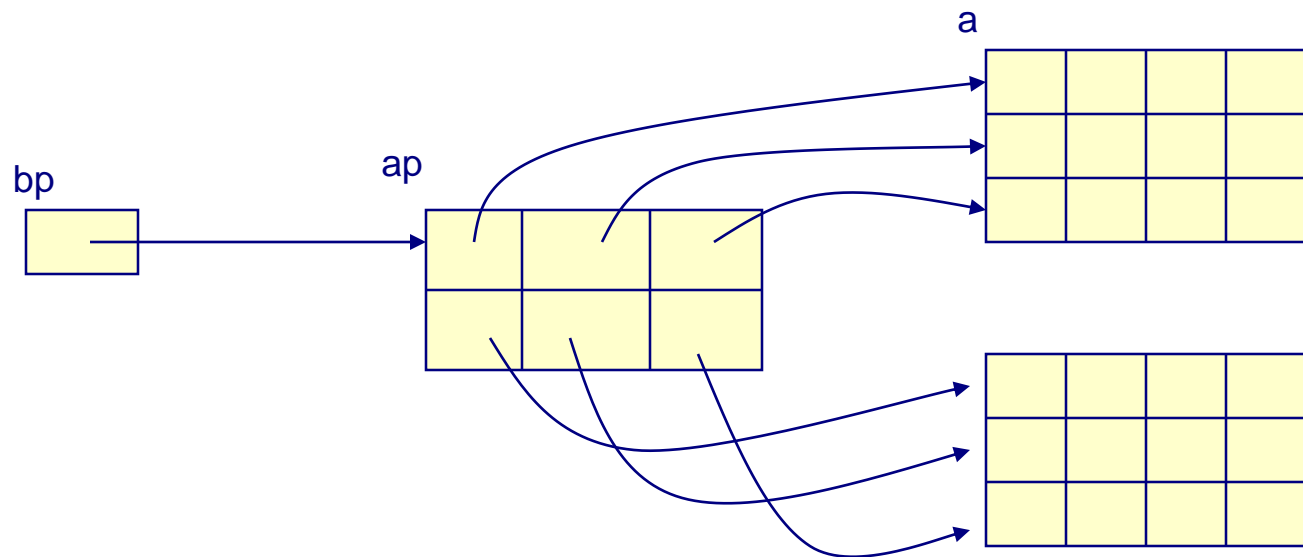
Exercise on pointer arithmetic

- Consider the following declaration

```
int a[2][3][4] = {1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17,18,19,20,21,22,23,24};
```

```
int *ap[2][3] = {{a[0][0], a[0][1], a[0][2]}, {a[1][0], a[1][1], a[1][2]}};
```

```
int *(*bp)[3] = ap;
```



Exercise on pointer arithmetic (cont)

- Derivation of `a[1][2][3]`

a:	T	[2][3][4] of int=ptr to [3][4] of int
	V	&a[0][0][0] (since “a” is an array name, it denotes the starting address of the array area)
a+1:	T	ptr to [3][4] of int
	V	&a[0][0][0]+1*sizeof([3][4] of int)=&a[1][0][0]
*(a+1)=a[1]	T	[3][4] of int = ptr to [4] of int
	V	&a[1][0][0] (Case 2)
a[1]+2	T	ptr to [4] of int
	V	&a[1][0][0]+2*sizeof([4] of int)=&a[1][2][0]
*(a[1]+2)=a[1][2]	T	[4] of int = ptr to int
	V	&a[1][2][0] (Case 2)
a[1][2]+3	T	ptr to int
	V	&a[1][2][0]+3*sizeof(int)=&a[1][2][3]
a[1][2][3]	T	int
	V	the contents of &a[1][2][3] =24 (Case 1)



Exercise on pointer arithmetic (cont)

- derivation of `ap[1][2][3]`

<code>ap</code>	T	<code>[2][3]</code> of ptr to int = ptr to <code>[3]</code> of ptr to int
	V	<code>&ap[0][0]</code>
<code>ap+1</code>	T	ptr to <code>[3]</code> of ptr to int
	V	<code>&ap[0][0]+1*sizeof([3] of ptr to int) = &ap[1][0]</code>
<code>ap[1]</code>	T	<code>[3]</code> of ptr to int = ptr to ptr to int
<code>= *(ap+1)</code>	V	<code>&ap[1][0]</code> (Case 2)
<code>ap[1]+2</code>	T	ptr to ptr to int
	V	<code>&ap[1][0]+2*sizeof(ptr to int) = &ap[1][2]</code>
<code>ap[1][2]</code>	T	ptr to int
<code>= *(ap[1]+2)</code>	V	the contents of <code>&ap[1][2]=&a[1][2][0]</code> (Case 1)
<code>ap[1][2]+3</code>	T	ptr to int
	V	<code>&a[1][2][0]+3*sizeof(int) = &a[1][2][3]</code>
<code>ap[1][2][3]</code>	T	int
<code>= *(ap[1][2]+3)</code>	V	the contents of <code>&a[1][2][3]=24</code> (Case 1)



Exercise on pointer arithmetic (cont)

- Derivation of `bp[1][2][3]`

<code>bp</code>	T	ptr to [3] of ptr to int
	V	<code>&ap[0][0]</code> (since <code>bp</code> is a simple variable, its value is the contents of the variable)
<code>bp+1</code>	T	ptr to [3] of ptr to int
	V	<code>&ap[0][0]+1*sizeof([3] of ptr to int)=&a[1][0]</code>
<code>bp[1]</code>	T	[3] of ptr to int = ptr to ptr to int
<code>= *(bp+1)</code>	V	<code>&ap[1][0]</code> (Case 2)
<code>bp[1]+2</code>	T	ptr to ptr to int
	V	<code>&ap[1][0]+2*sizeof(ptr to int) = &ap[1][2]</code>
<code>bp[1][2]</code>	T	ptr to int
<code>= *(bp[1]+2)</code>	V	the contents of <code>&ap[1][2] = &a[1][2][0]</code> (Case 1)
<code>bp[1][2]+3</code>	T	ptr to int
	V	<code>&a[1][2][0]+3*sizeof(int) = &a[1][2][3]</code>
<code>bp[1][2][3]</code>	T	int
<code>= *(bp[1][2]+3)</code>	V	the contents of <code>&a[1][2][3] = 24</code> (Case 1)



Assignments to pointer variables

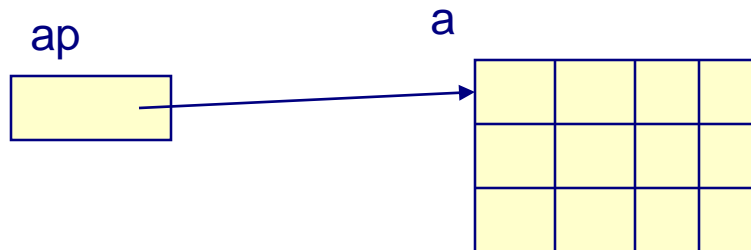
- In an assignment statement to a pointer variable, the type of the right hand side must match the type of the left hand side
- Recall the previous declaration of a, ap, and bp
- `int *ap[2][3] = {{a[0][0], a[0][1], a[0][2]}, {a[1][0], a[1][1], a[1][2]}}`;
 - each element of ap is of type “pointer to int”
 - from the derivation of `a[1][2][3]`, the type of `a[1][2]` is also “pointer to int”
- `int *(*bp)[3] = ap;`
 - the type of bp is “pointer to [3] of pointer to int”
 - the type of ap is “[2][3] of pointer to int”, which is equivalent to “pointer to [3] of pointer to int”



Example of Incorrect Pointer Declaration

```
int a[3][4] = {1,2,3,4,5,6,7,8,9,10,11,12};  
int **ap = a;    // warning "type mismatch"
```

access ap[1][2]



Example of Incorrect Pointer Declaration (cont)

ap T: ptr to ptr to int
 V: &a[0][0]

ap+1 T: ptr to ptr to int
 V: &a[0][0]+1*sizeof(ptr to int) // assume sizeof(int) == sizeof(ptr) == 4
 = &a[0][1]

ap[1] T: ptr to int
 V: the contents of &a[0][1]=2

ap[1]+2 T: ptr to int
 V: 2 + 2*sizeof(int) = 10

ap[1][2] T: int
 V: the contents of address 10 (segmentation fault)

Output:
4202496
4202500
2
10
Segmentation fault (core dumped)

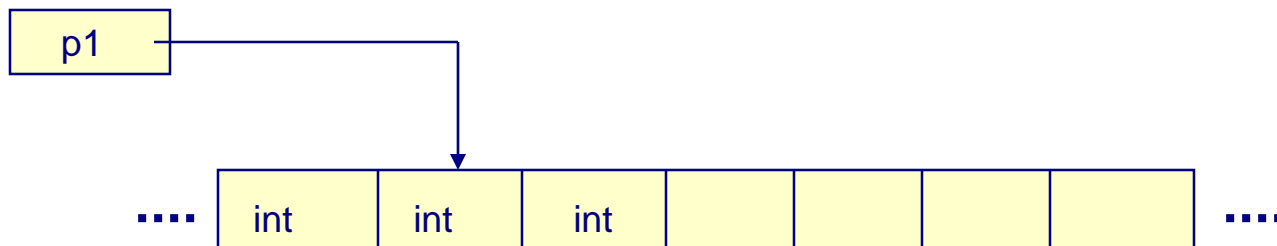


Declaration of pointer variables

- Consider a declaration “int *p1”
- By applying the pointer arithmetic on “p1”, you can access any element in the virtual one-dimensional integer array starting at the address pointed to by “p1”
- Thus, the image of the object pointed to by “p1” is not

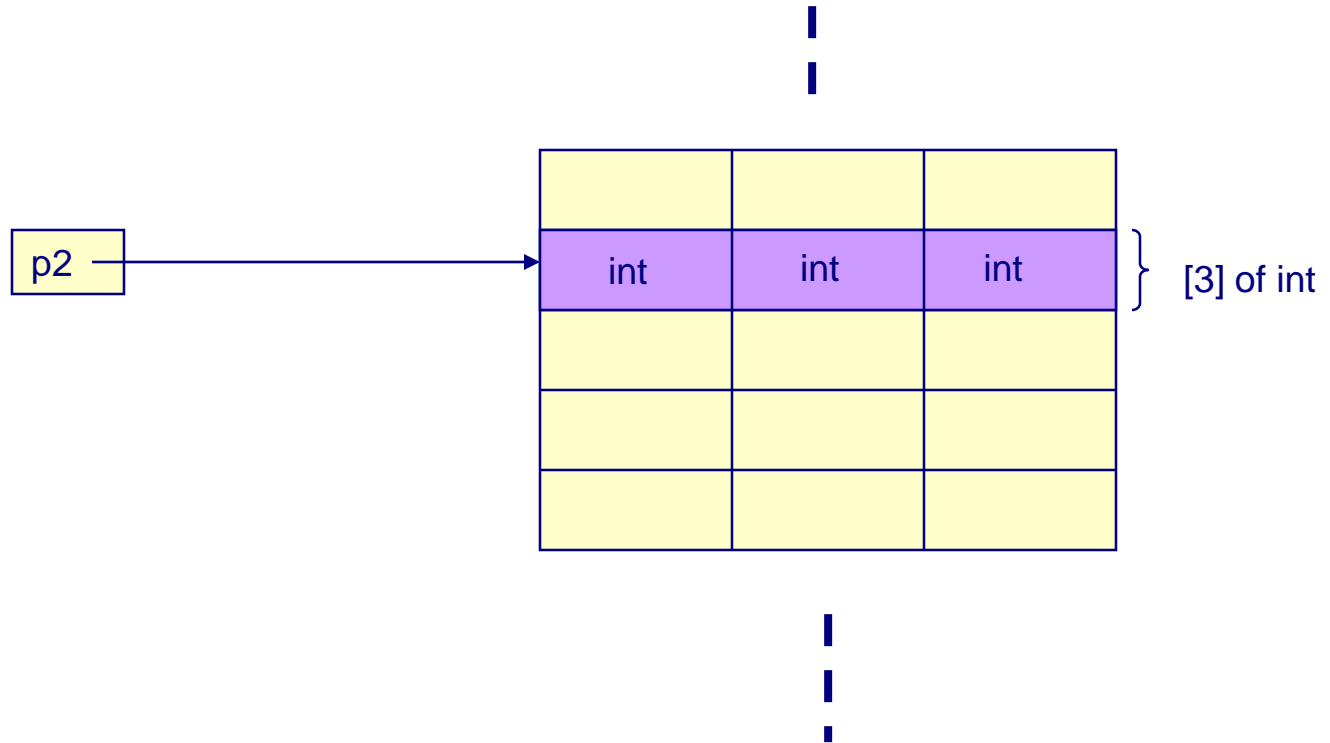


- but the following infinite length of one-dimensional integer array

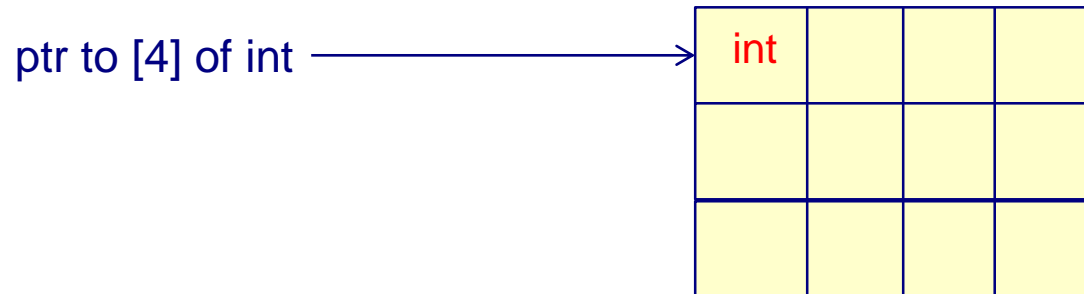


Declaration of pointer variables (cont)

- Similarly, “`int (*p2)[3]`” declares pointer variable “`p2`” that points to the following infinite length of one-dimensional array [3] of int

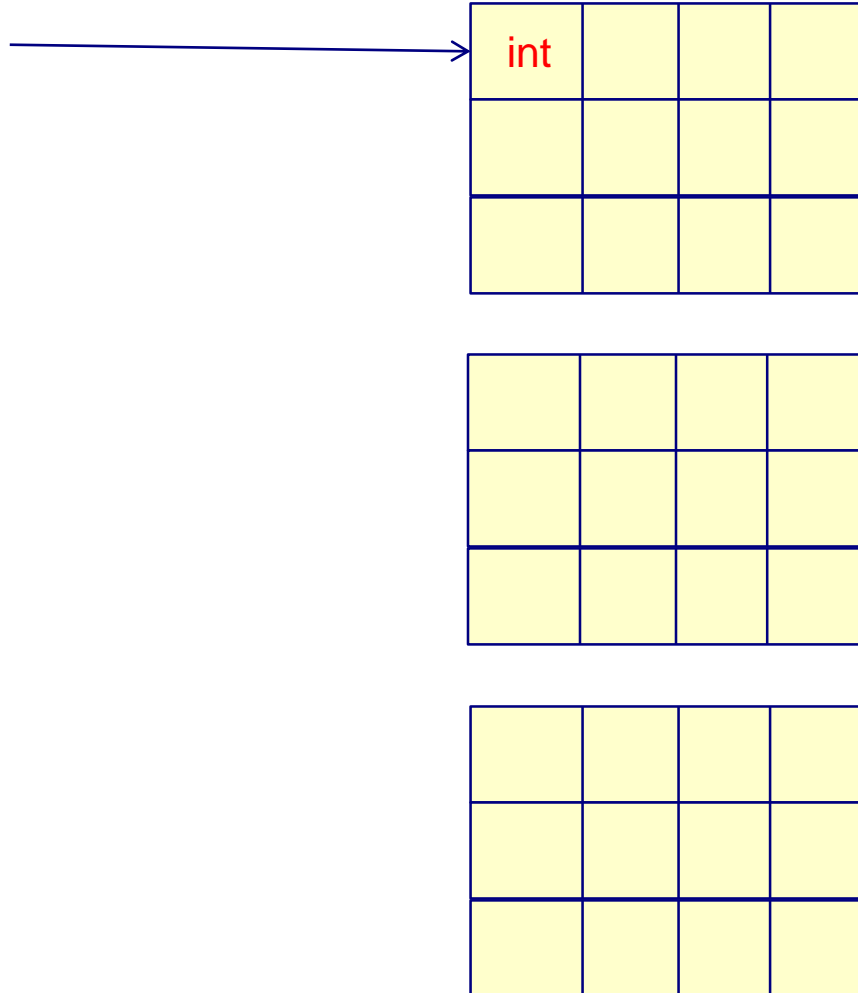


Declaration of pointer variables (cont)



Declaration of pointer variables (cont)

ptr to [3][4] of int



Declaration of pointer variables (cont)

Consider the previous declaration of a, ap, and bp

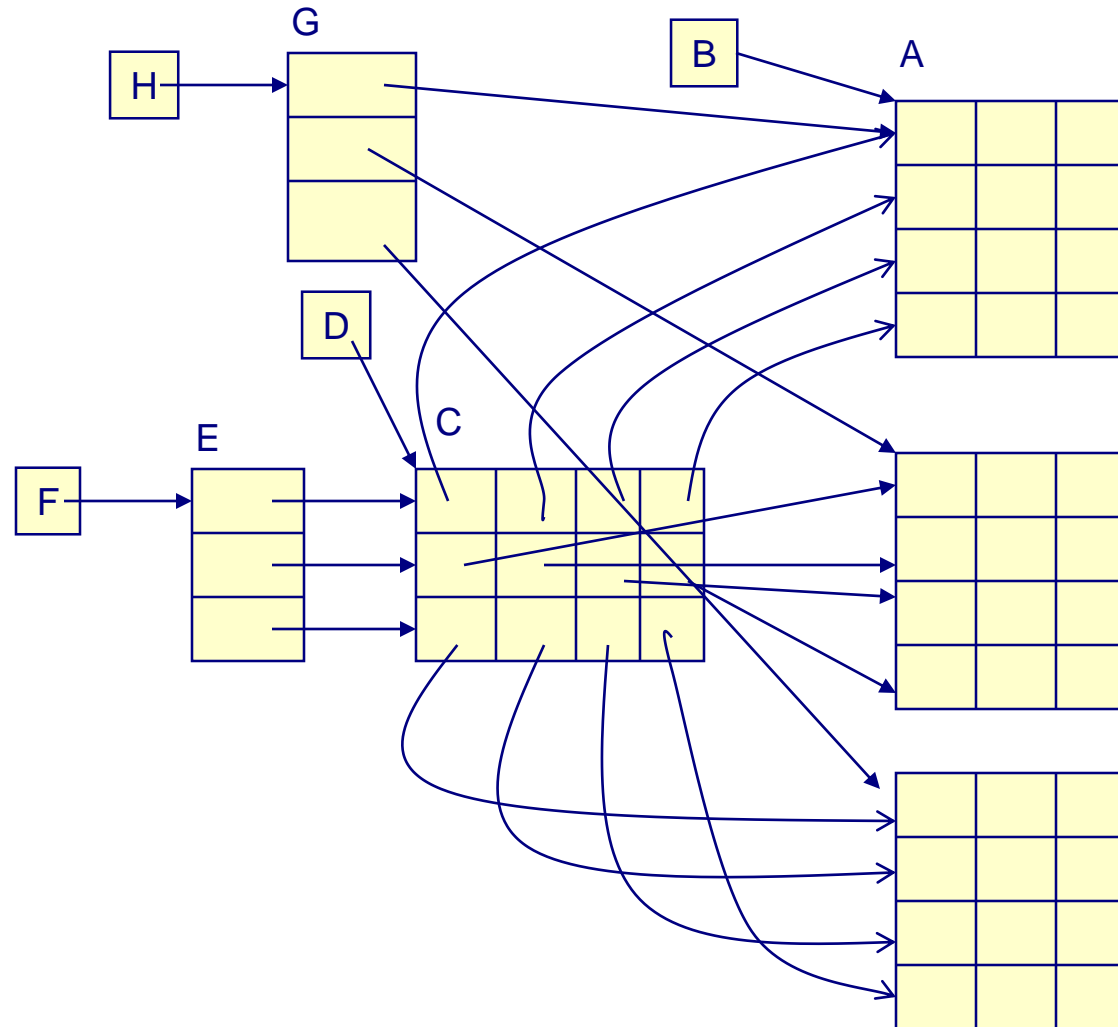
- The type of ap is “array[2][3] of pointer to int”
 - ap is an array with [2][3] elements, and each element is of type “pointer to int”
 - each element points to one row of array “a” (each row consists of 4 elements)
- The type of bp is “pointer to [3] of pointer to int”
 - bp itself is a simple (non-array) pointer variable that points to “[3] of pointer to int”
 - bp points to ap, that is, “[2][3] of pointer to int” -- a two-dimensional array of pointer to integer with column length 3
 - each element of the two-dimensional array (of type “pointer to integer”) points to one row of “a”



Exercise on pointer declaration

Declare the pointer variables so that the following relation holds

$A[i][j][k] ==$
 $B[i][j][k] ==$
 $C[i][j][k] ==$
 $D[i][j][k] ==$
 $E[i][j][k] ==$
 $F[i][j][k] ==$
 $G[i][j][k] ==$
 $H[i][j][k]$



Allocation of arrays in heap

- Library function “void *malloc(SIZE in bytes)” allocates “SIZE bytes” in the heap area and returns the base (starting) address of the allocated area (the return type is ptr to void)
- If an area for an imaginary array is allocated in heap and its starting address is set to a pointer that is appropriately declared, any element in the imaginary array is accessed via the pointer
 - Since malloc() returns the “pointer to void” type, the return value must be cast to an appropriate type



Allocation of arrays in heap (cont)

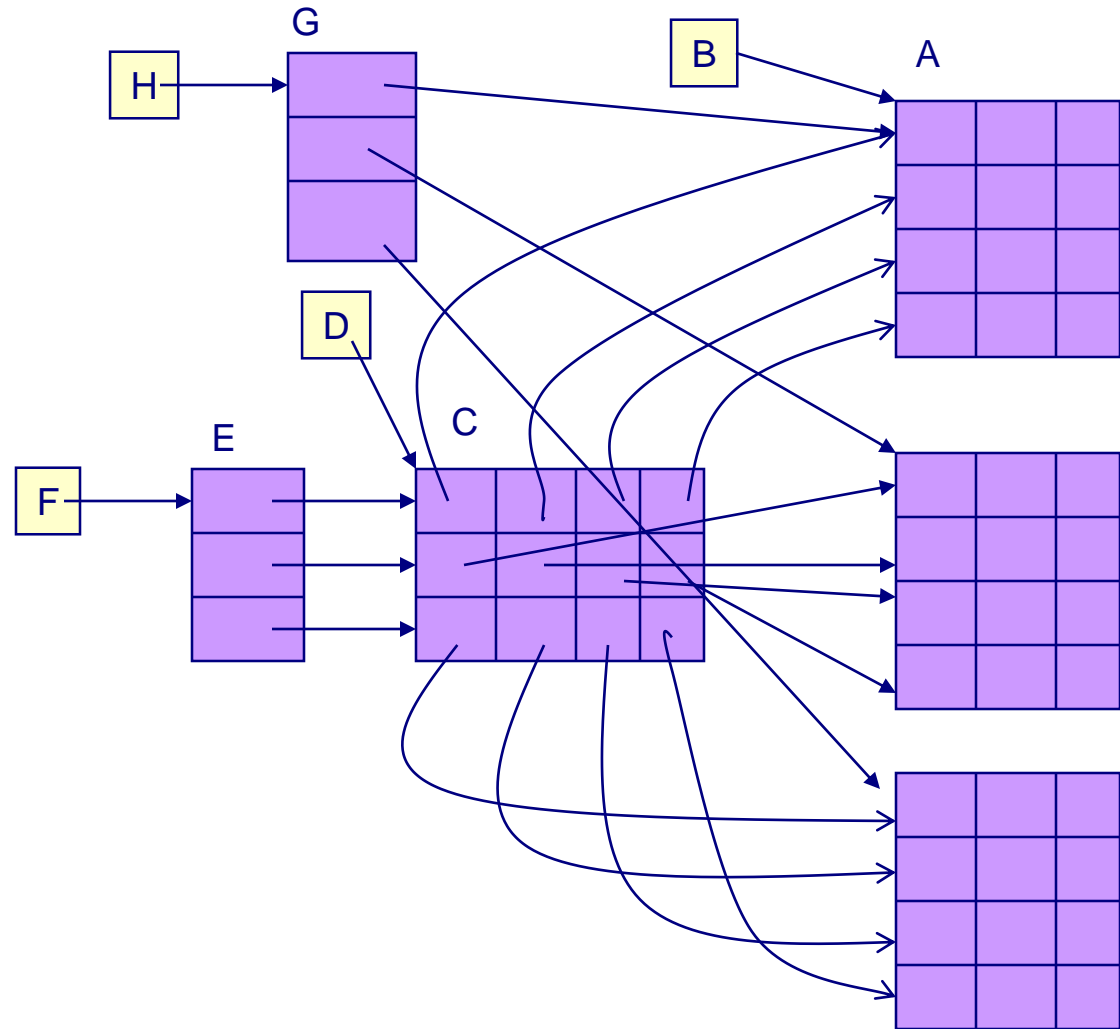
- To specify a type in cast and sizeof operations,
 - first, declare the type for variable “x” (any name is ok)
 - then, delete “x” in the declaration
- Example:
 - type “ptr to [4] of ptr to int”:
 1. declare the type for x: `int *(*x)[4]`
 2. delete x in the declaration: `(int *(*)[4])`
 - type “[3][4] of ptr to int”:
 1. declare the type for x: `int *x[3][4]`
 2. delete x in the declaration: `(int *[3][4])`
- Allocate an imaginary array of type “[3][4] of ptr to int” in heap and declare a pointer variable “d” to point to the array (the type of “d” is “ptr to [4] of ptr to int”)
 - `int *(*d)[4] = (int *(*)[4]) malloc(sizeof(int *[3][4]));`
- Note: Type declarations for the “sizeof” operation is not crucial. Only the size of instances of the type matters
 - You can declare `malloc(3*4*sizeof(int *))`, instead of `malloc(sizeof(int *[3][4]))`



Allocation of arrays in heap (cont)

Allocate purple colored arrays in heap and declare pointer variables correctly so that the following relation holds

$B[i][j][k] ==$
 $D[i][j][k] ==$
 $F[i][j][k] ==$
 $H[i][j][k]$



Allocation of arrays in heap (cont)

Declarations:

```
int (*b)[4][3];
```

```
int *(*d)[4];
```

```
int ***f;
```

```
int (**h)[3];
```

```
// allocate array a and initialize it
```

```
b = (int (*)(4)[3]) malloc(sizeof(int [3][4][3]));
```

```
for(i = 0; i < 3; i++)
```

```
    for(j = 0; j < 4; j++)
```

```
        for(k = 0; k < 3; k++)
```

```
            b[i][j][k] = i*12 + j*3 + k + 1;
```

```
// allocate array c and initialize it
```

```
d = (int (*)(4)) malloc(sizeof(int *[3][4]));
```

```
for(i = 0; i < 3; i++)
```

```
    for(j = 0; j < 4; j++)
```

```
        d[i][j] = b[i][j];
```

```
// allocate array e and initialize it
```

```
f = (int ***) malloc(sizeof(int **[3]));
```

```
for(i = 0; i < 3; i++) f[i] = d[i];
```

```
// allocate array g and initialize it
```

```
h = (int (**)[3]) malloc(sizeof(int (*)(3)[3]));
```

```
for(i = 0; i < 3; i++) h[i] = b[i];
```



pointer to function

- Consider the following declaration
 - `int (*ptf) (int, int)`
 - `ptf` : a pointer to a function that takes (int, int) and returns int
 - that is, an area for variable `ptr` is allocated that stores a pointer to (address of) a function (in the text segment) that takes int and int as arguments and returns int
- In C, the name of a function denotes the starting address of the machine code of the function in the text segment
 - Suppose that function “`int intAdd(int x, int y) { return x + y;}`” is defined
 - By assignment “`ptf = intAdd;`”, the starting address of “`intAdd`” is set in `ptf`
 - To call “`intAdd(3, 59);`” through “`ptr`”, the following two ways are possible (both are end up with the identical assembly code):
 - `int z = ptf(3, 59);`
 - `int z = (*ptf)(3, 59);`

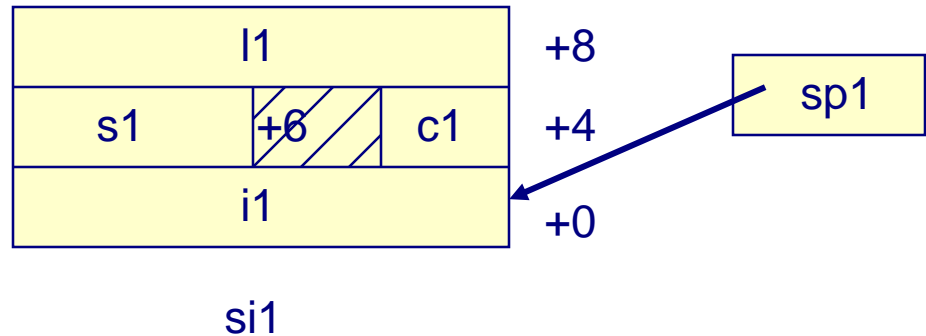


struct

- In struct, each field is converted to the offset relative to the starting address of the instance at compilation time
- The size of a struct instance is a multiple of the word size of the machine

In case of 4-byte, Little Endian machine

```
typedef struct {  
    int i1;  
    char c1;  
    short s1;  
    long l1;  
} struct1;
```



```
struct1 si1;  
struct1 *sp1 = &si1;
```

```
si1.l1 = 23;    // put 23 in the four-byte area starting &si1+8  
sp1->c1 = 'a'; // put 'a' in the byte at the value of sp1 +4  
sp1->s1 = 5;    // put 5 in the two-byte area starting the value  
                // of sp1 +6  
si1.i1 = 31;    // the four-byte area starting &si1+0
```

