

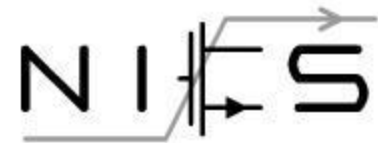


NXgraph: An Efficient Graph Processing System on a Single Machine

Yuze Chi¹, Guohao Dai¹, Yu Wang¹, Guangyu Sun²,
Guoliang Li¹ and Huazhong Yang¹

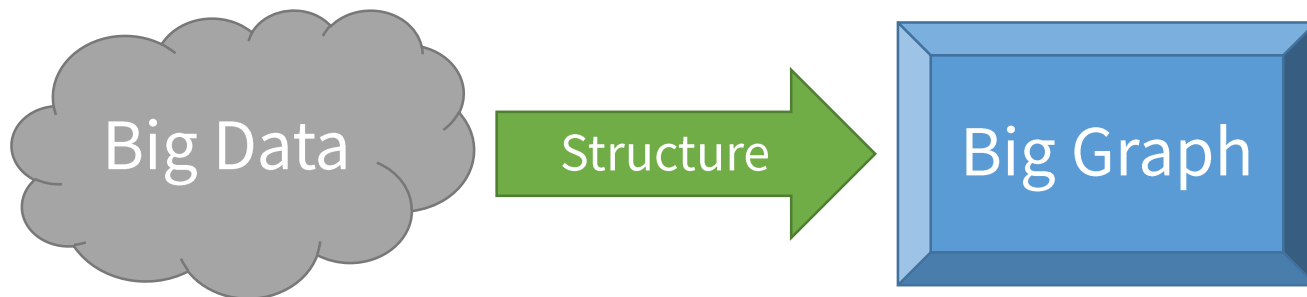
Tsinghua National Laboratory for Information Science and Technology¹,
Tsinghua University

Center for Energy Efficient Computing and Applications²,
Peking University





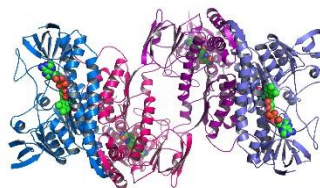
Motivation



social media



science



advertising



web





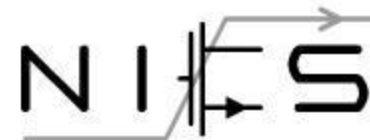
Motivation

Can we
exploit it well?

- Computation capacity of a single CPU
 - Intel i7-5820K: 12 hyper-threads, 3.3GHz
 - Assume 50 clock cycles/edge: **792MEPS** (Million Edges Per Second)

System	Throughput/MEPS (PageRank on Twitter)	Notes
Spark ^{HotCloud2010}	15 in total/0.15 each	100 CPUs in 50 nodes
PowerGraph ^{OSDI2012}	408 in total/6.4 each	64 CPUs in 64 nodes
GraphChi ^{OSDI2012}	50	
TurboGraph ^{SIGKDD2013}	108	
X-stream ^{SOSP2013}	20	no pre-processing
VENUS ^{ICDE2015}	15.4	on HDD
GridGraph ^{ATC2015}	61	

- Large gap!
- Our objective: **higher MEPS/CPU**





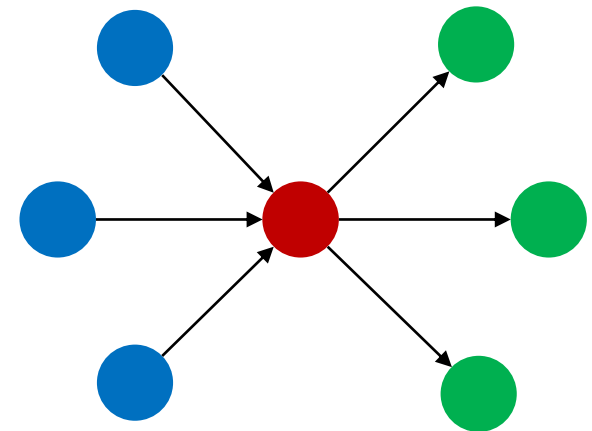
Vertex-centric Model SIGMOD2010

- Graph $G = (V, E)$
 - vertex $v = (\text{id}, \text{attr})$
 - edge $e = (\text{src vid}, \text{dst vid}, \text{optional attribute})$
- "Think like a vertex"

Update my neighbors...

- Example: Breadth-First Search (BFS)

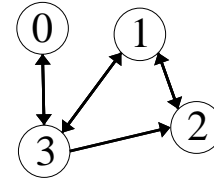
```
for each dst in my.out_edges
  if dst.depth > my.depth+1
  then
    dst.depth = my.depth+1
```





Interval-shard Partitioning^{OSDI2012}

- Vertices \rightarrow Intervals
- Edges \rightarrow Shards
 - Edges in each shard \rightarrow Sub-Shards
- Objective
 - Limit memory access to a small region
 - Improve locality

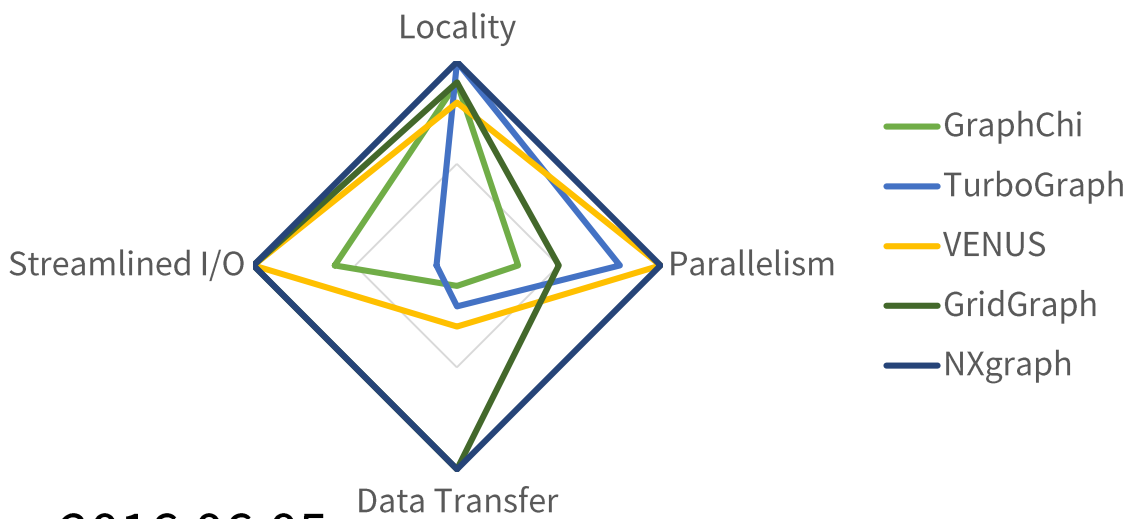


I ₁	I ₂
0, 1	2, 3
S ₁	S ₂
SS _{1.1}	SS _{1.2}
	1 \rightarrow 2 0, 1 \rightarrow 3
SS _{2.1}	SS _{2.2}
3 \rightarrow 0 2, 3 \rightarrow 1	3 \rightarrow 2



Four Optimizing Rules

	GraphChi OSDI2012	TurboGraph SIGKDD2013	VENUS ICDE2015	GridGraph ATC2015	NXgraph ICDE2016
1. Exploit the locality of graph data	😊	😊	😊	😊	😊
2. Utilize the parallelism of multi-thread CPU		😊	😊		😊
3. Reduce the amount of disk data transfer				😊	😊
4. Streamline the disk I/O	😊		😊	😊	😊



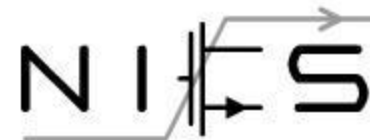
No previous system considers them all





Our Effort

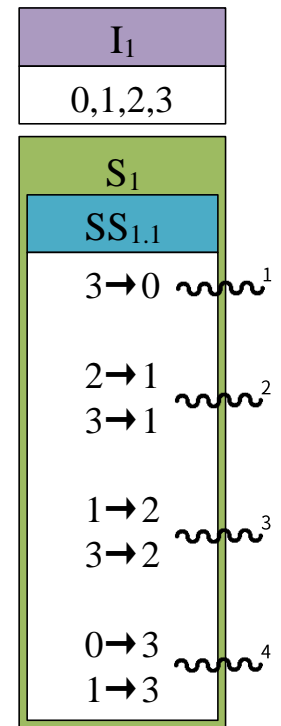
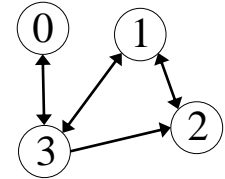
- **Follows all four optimizing rules**
- According to optimizing rule 1 & 2, we design
 - Destination-Sorted Sub-Shard (DSSS) structure
- According to optimizing rule 3 & 4, we design
 - Adaptive updating strategies
 - Single-Phase Update (SPU)
 - Double-Phase Update (DPU)
 - Mixed-Phase Update (MPU)

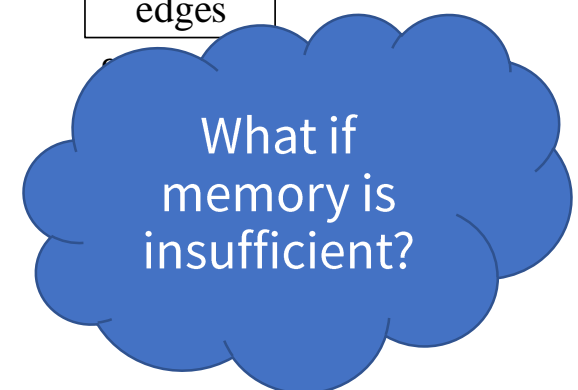




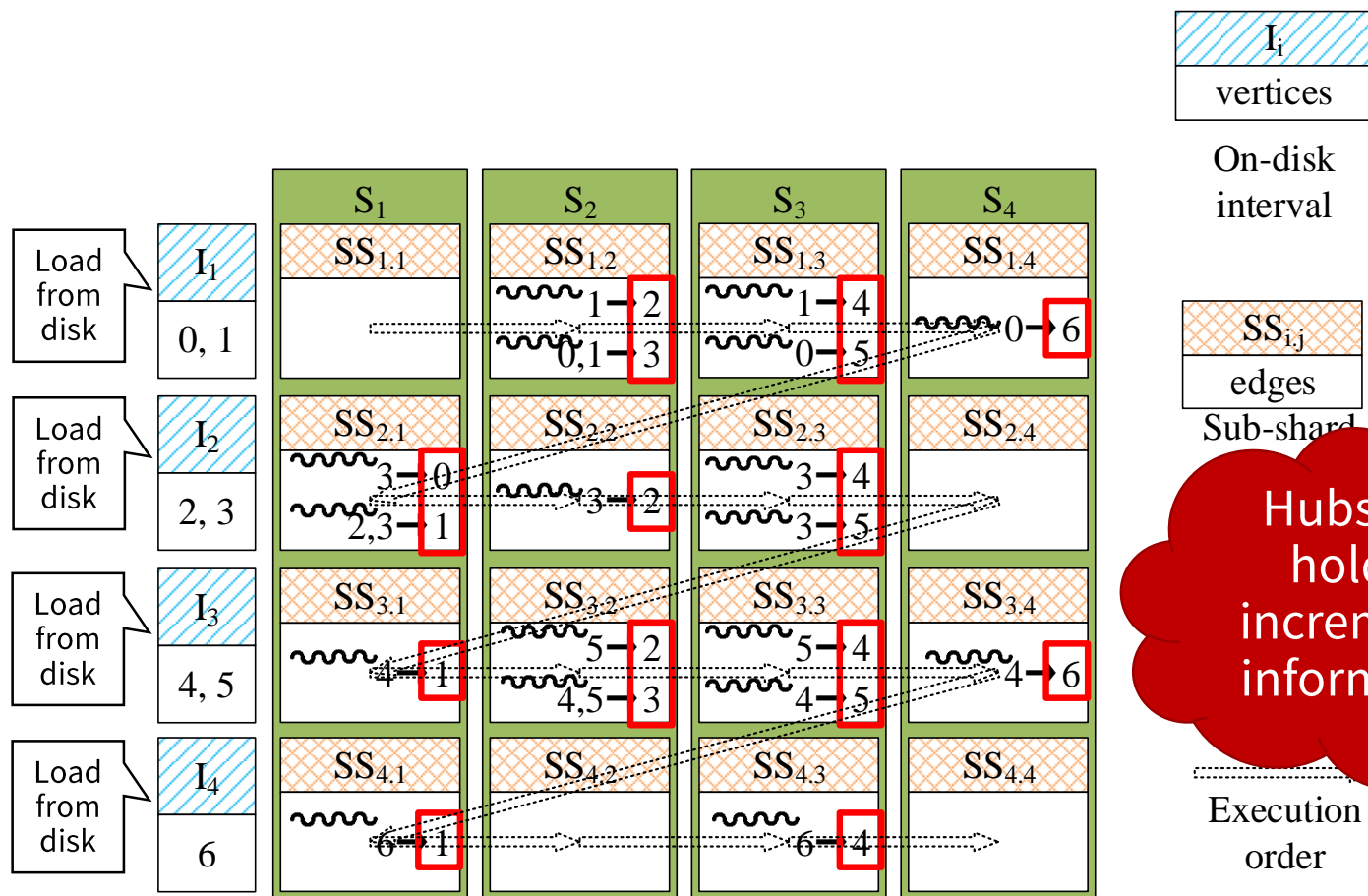
Destination-Sorted Sub-Shards

- Sort edges in each sub-shard
- Sequential reads from source interval (Rule 1)
 - Potentially improves cache hit rate
- Parallel writes to destination vertices (Rule 2)
 - No write conflict among threads
- Defines behavior **inside** each sub-shard





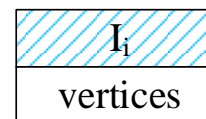
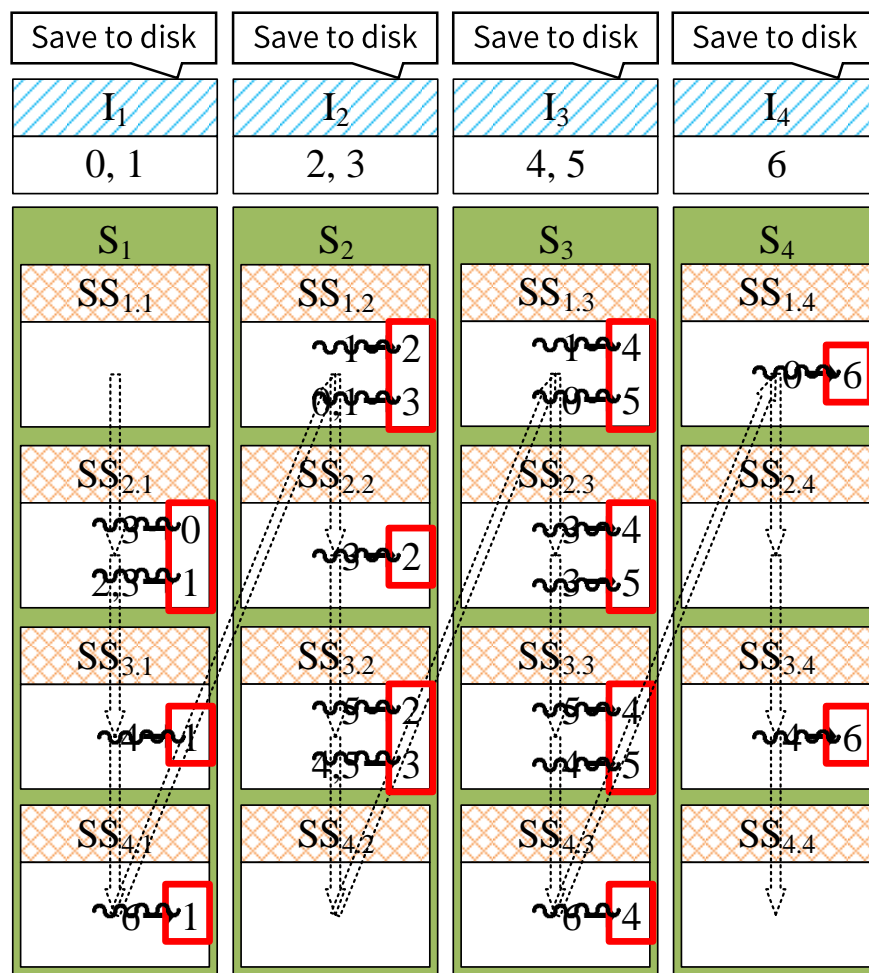
Double-Phase Update (To-Hub Phase)



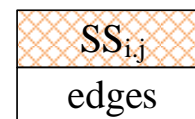
Hubs now hold all incremental information



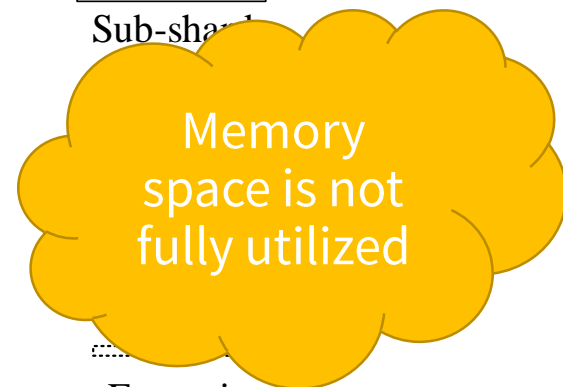
Double-Phase Update (From-Hub Phase)



On-disk
interval

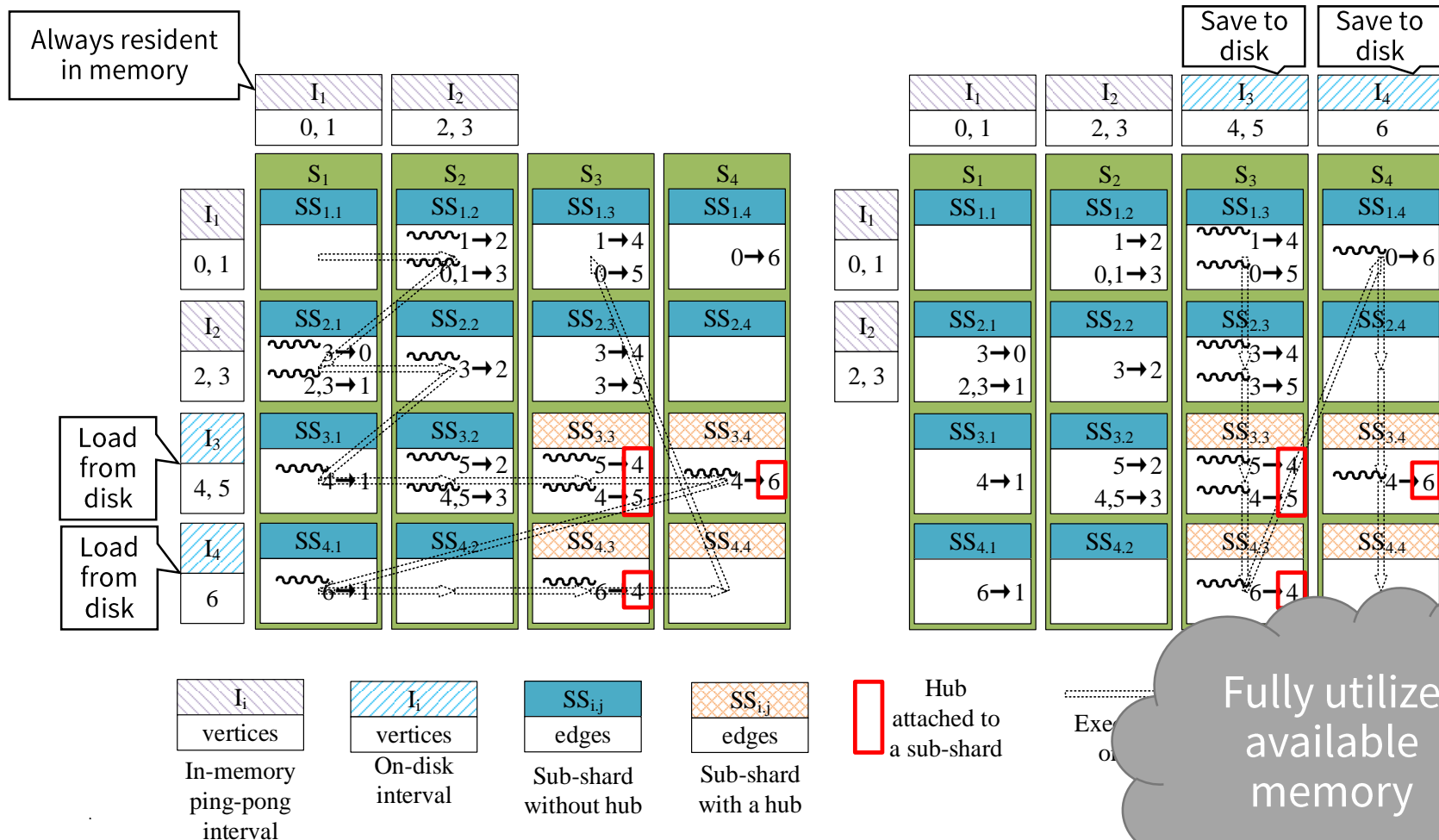


Sub-shard



Execution
order

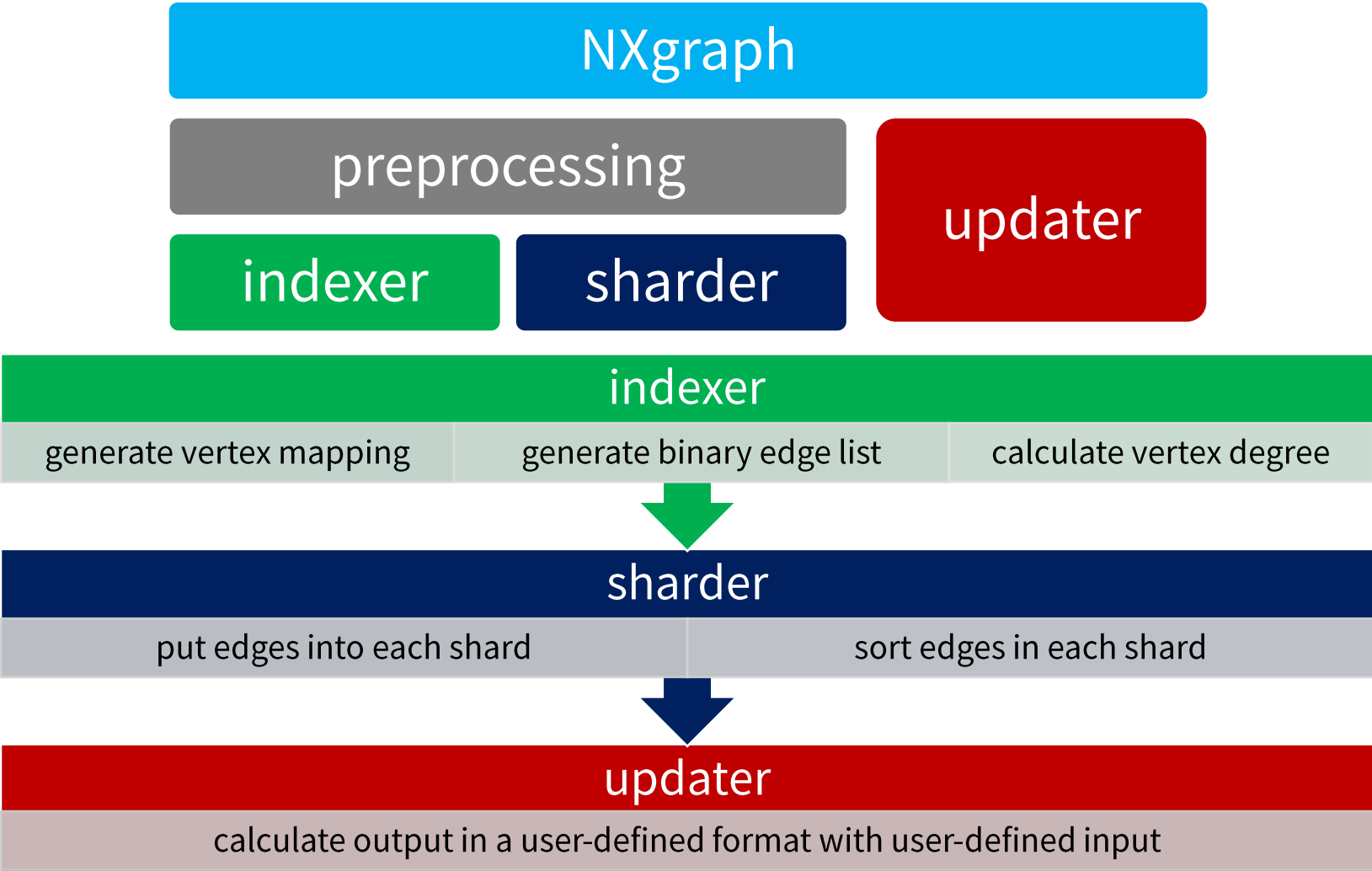
Mixed-Phase Update



Fully utilize available memory



System Architecture





Results

- Evaluation platform
 - Hex-core hyper-threading Intel i7-5820K CPU @ 3.3GHz
 - 8x8G DDR4 RAM, 2x128G RAID0 SSD, 1T HDD
 - Ubuntu 14.04 LTS 64bit/Windows 10 Edu 64bit



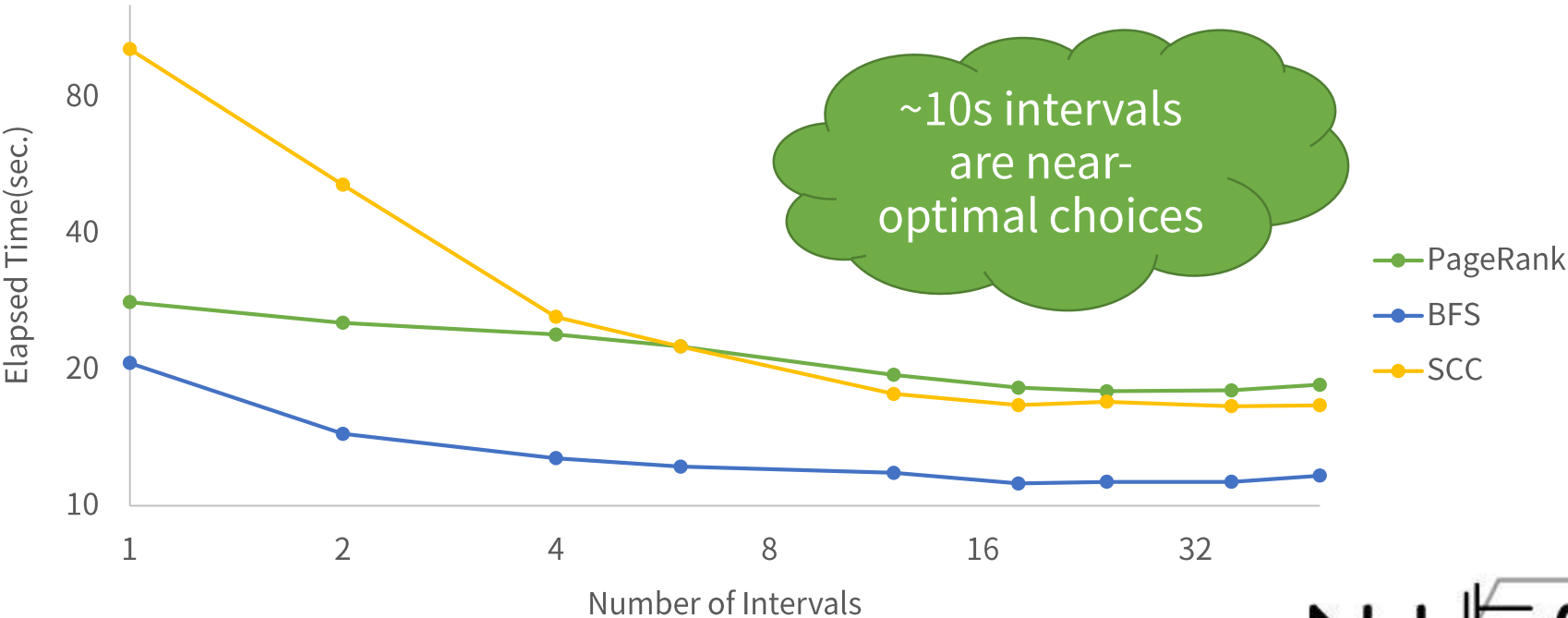
Results: Design Decisions

Performance with different sub-shard model

Model	Elapsed Time	
	Live-journal	Twitter
src-sorted, coarse-grained	1.44s	72.5s
dst-sorted, fine-grained	1.00s	20.50s

Destination-sorted is the right choice!

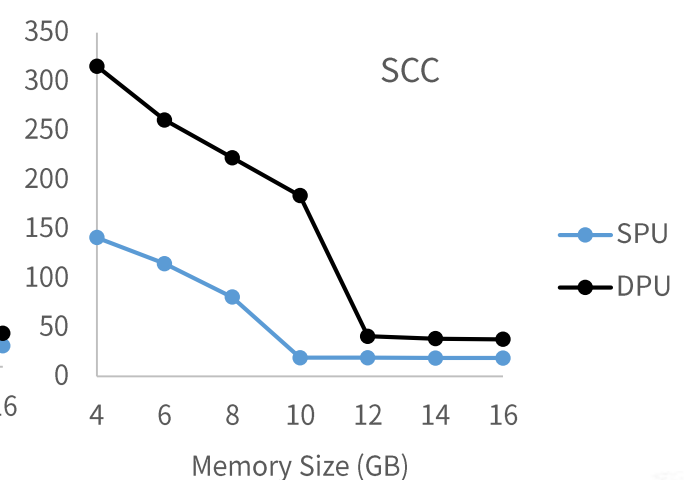
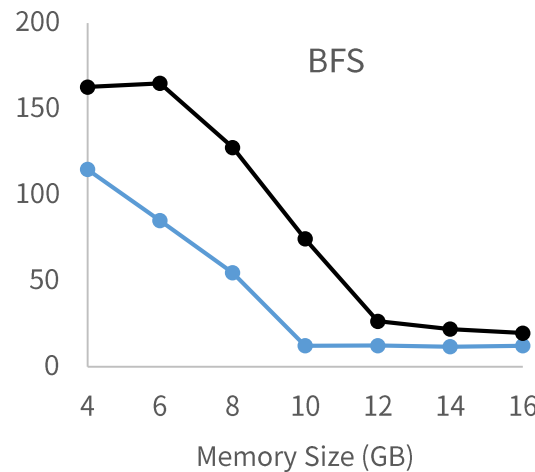
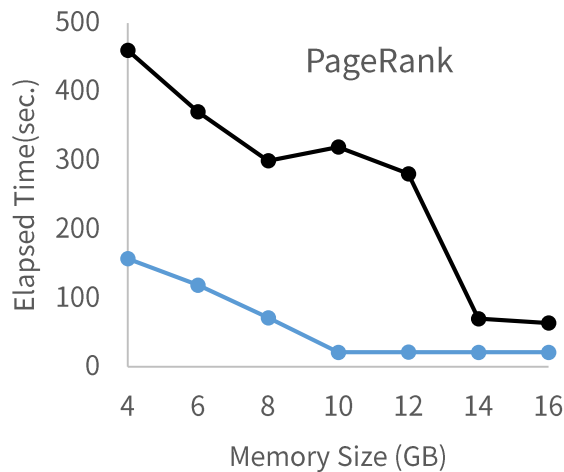
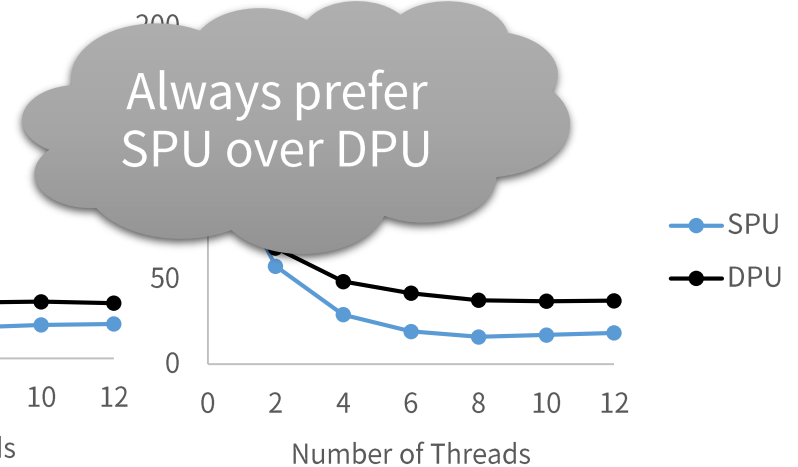
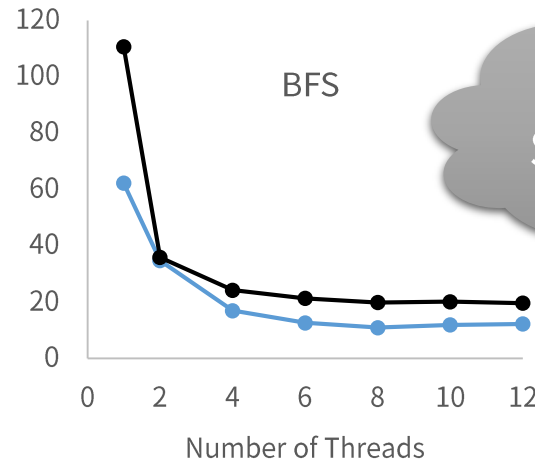
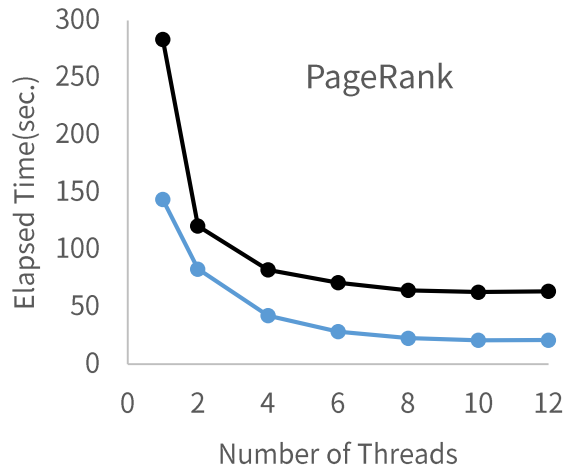
Performance with different numbers of interval on Twitter





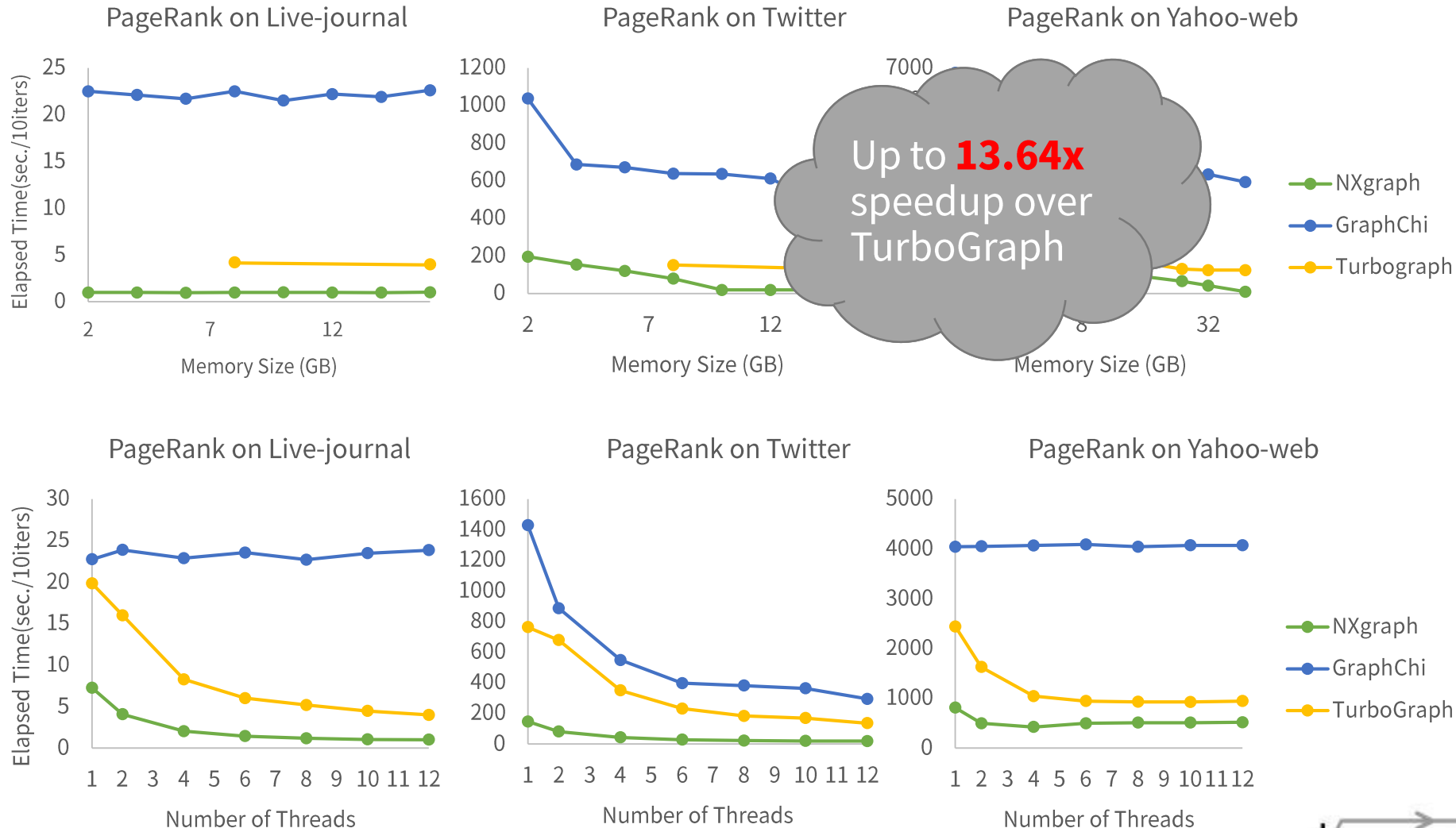
Results: Design Decisions

SPU vs DPU on performance





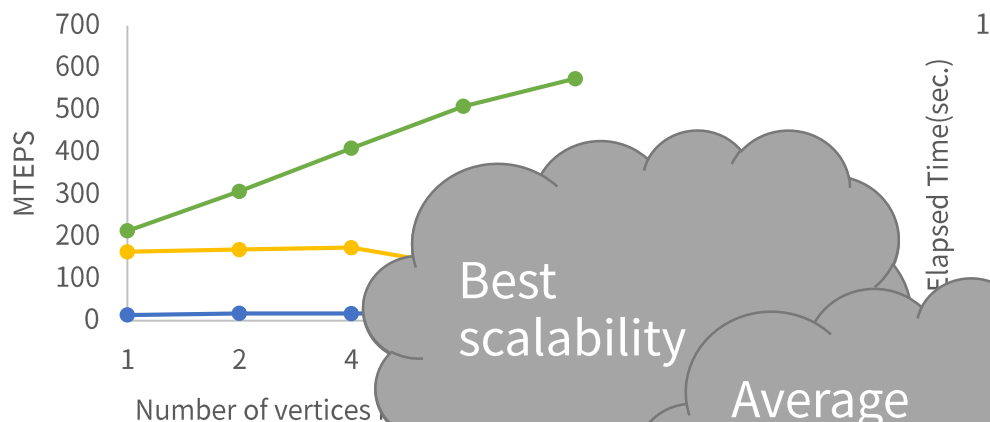
Results: Different Environments



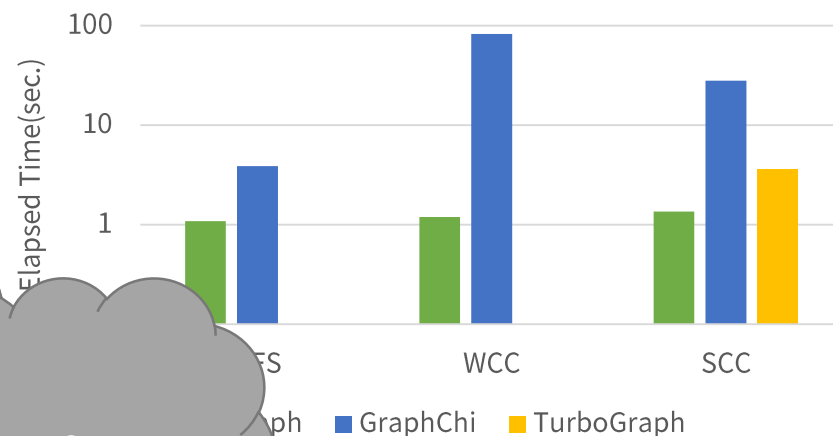


Results: Scalability and More Tasks

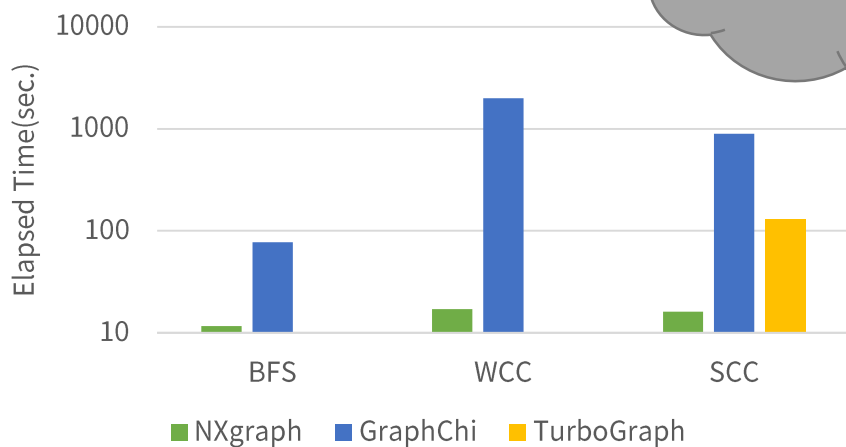
Scalability



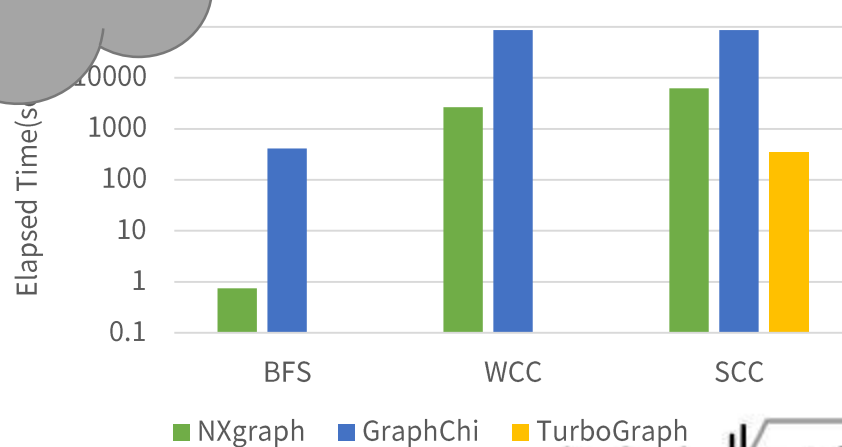
More tasks on Live-journal



More tasks on Twitter



More tasks on Yahoo-web





Results: More Systems (Limited Resources)

System performance with limited resources

System	Time (s)	Speedup	Evaluation environment
NXgraph	7.13	1.00	Intel i7 3.3GHz, 8t, 8G, SSD
GridGraph	26.91	3.77	AWS EC2 8t, 8G/30.5G, SSD
X-stream	88.95	12.48	
NXgraph	12.55	1.00	
VENUS	95.48	7.60	
GridGraph	24.11	1.92	
X-stream	81.70	6.51	AWS EC2, 8t, 8G/30.5G, HDD

Average speedup of
6.6x over various
state-of-the-art single-
machine systems with
limited resources

Task: 1 iteration of PageRank on Twitter graph



Results: More Systems (Best Case)

System performance in the best case

System	Time (s)	Speedup	Evaluation environment
NXgraph	2.05	1.00	Intel i7 3.3GHz, 8t, 16G, SSD
X-stream	23.25	11.57	Apple M1, 8t, 16G, SSD
GridGraph	24.11	11.99	Apple M1, 8t, 16G, SSD
MMap	13.10	6.52	Apple M1, 8t, 16G, SSD
PowerGraph	3.60	1.79	Intel i7 3.3GHz, 8t, 23G)

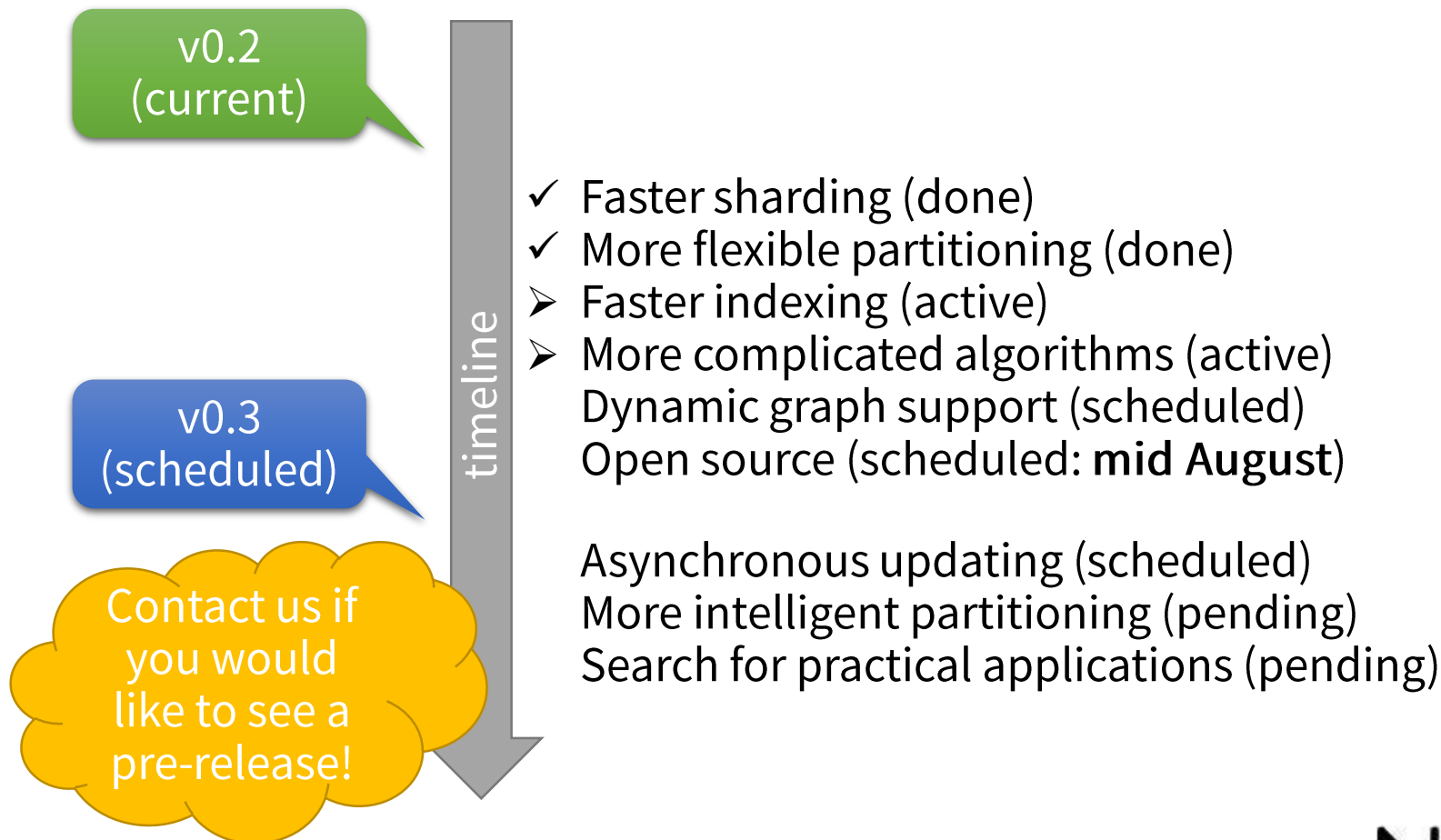
717MEPS actual
throughput vs
792MEPS
hypothetical limit

Task: 1 iteration of PageRank on Twitter graph



Future Work

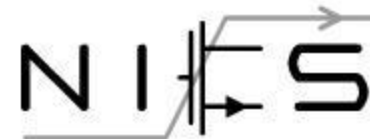
- NXgraph is still under development and subject to changes in data structures and APIs





Reference

- H. Kwak, C. Lee, H. Park, and S. Moon, “What is twitter, a social network or a news media?” in *WWW. ACM*, 2010, pp. 591–600.
- J. Gonzalez, Y. Low, and H. Gu, “Powergraph: Distributed graphparallel computation on natural graphs,” in *OSDI*, 2012, pp. 17–30.
- J. Cheng, Q. Liu, Z. Li, W. Fan, J. C. S. Lui, and C. He, “VENUS: Vertex-Centric Streamlined Graph Computation on a Single PC,” in *ICDE*, 2015, pp. 1131–1142.
- A. Kyrola, G. Blelloch, and C. Guestrin, “GraphChi: Large-Scale Graph Computation on Just a PC,” in *OSDI*, 2012, pp. 31–46.
- W.-S. Han, S. Lee, K. Park, J.-H. Lee, M.-S. Kim, J. Kim, and H. Yu, “TurboGraph: A Fast Parallel Graph Engine Handling Billion-Scale Graphs in a Single PC,” in *SIGKDD*, 2013, pp. 77–85.
- M. Zaharia, M. Chowdhury, M. J. Franklin, S. Shenker, and I. Stoica, “Spark: cluster computing with working sets,” in *HotCloud*, vol. 10, 2010, p. 10.
- L. Page, S. Brin, R. Motwani, and T. Winograd, “The pagerank citation ranking: bringing order to the web.” 1999.
- X. Zhu, W. Han, and W. Chen, “GridGraph : Large-Scale Graph Processing on a Single Machine Using 2-Level Hierarchical Partitioning,” in *ATC*, 2015, pp. 375–386.
- A. Roy, I. Mihailovic, and W. Zwaenepoel, “X-Stream: Edge-centric Graph Processing using Streaming Partitions,” in *SOSP*, 2013, pp. 472–488.





Reference

- Yahoo! altavisata web page hyperlink connectivity graph, circa 2002,”
- <http://webscope.sandbox.yahoo.com/>.
- “Livejournal social network,” <http://snap.stanford.edu/data/soc-LiveJournal1.html>.
- Z. Lin, M. Kahng, K. Sabrin, D. Horng, and P. Chau, “MMap : Fast Billion-Scale Graph Computation on a PC via Memory Mapping,” in *ICBD*. IEEE, 2014.
- G. Malewicz, M. H. Austern, A. J. C. Bik, J. C. Dehnert, I. Horn, N. Leiser, and G. Czajkowski, “Pregel : A System for Large-Scale Graph Processing,” in *SIGMOD*, 2010, pp. 135–145.



Thank you!

Q&A