

1 Linear Models

1.1 Types of errors

Errors come from: imprecise data, mistakes in the model, computational precision,.. We know two types of errors:

- **Absolute** error = approximate value - correct value

$$\Delta x = \bar{x} - x$$

- **Relative** error = $\frac{\text{absolute err}}{\text{correct value}}$

$$\delta x = \frac{\Delta x}{x}$$

1.2 Mathematical model is **linear**, when the function F is a linear function of the parameters:

$$F(x, a_1, \dots, a_p) = a_1\phi_1(x) + \dots + a_p\phi_p(x)$$

where ϕ_1, \dots, ϕ_p are functions of a specific type.

1.3 Least squares method Given points

$$\{(x_1, y_1), \dots, (x_m, y_m)\}, x_i \in R^n, y_i \in R$$

the task is to find a function $F(x, a_1, \dots, a_p)$ that is good fit for the data. The values of the parameters a_1, \dots, a_p should be chosen so that the equations

$$y_i = F(x, a_1, \dots, a_p), i = 1, \dots, m$$

are satisfied or, if this is not possible, that the error is as small as possible.

We use **Least squares method** to determine that the sum of squared errors is as small as possible.

$$\sum_{i=1}^m (F(x_i, a_1, \dots, a_p) - y_i)^2$$

1.4 Systems of linear equations

A system of linear equations in the matrix form is given by $A\vec{x} = \vec{b}$, where:

- A is the matrix of coefficients of order $m \times n$ where m is the number of equations and n is the number of unknowns,
- \vec{x} is the vector of unknowns and
- \vec{b} is the right side vector

$$\begin{bmatrix} \phi_1(x_1) & \phi_2(x_1) & \dots & \phi_p(x_1) \\ \phi_1(x_2) & \phi_2(x_2) & \dots & \phi_p(x_2) \\ \vdots & \vdots & \ddots & \vdots \\ \phi_1(x_n) & \phi_2(x_n) & \dots & \phi_p(x_n) \end{bmatrix} \begin{bmatrix} a_1 \\ a_2 \\ \vdots \\ a_p \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_p \end{bmatrix}$$

1.5 Existence of solutions in linear equations

Let $A = [\vec{a}_1, \dots, \vec{a}_n]$, where \vec{a}_i are vector representing the columns of A. For any vector

$$\vec{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_p \end{bmatrix} \text{ the product } A\vec{x} \text{ is a linear combination } A\vec{x} = \sum_i x_i a_i.$$

The system is **solvable** iff the vector \vec{b} can be expressed as a linear combination of the columns of A, that is it is in the column space of A, $\vec{b} \in C(A)$.

By adding \vec{b} to the columns of A we obtain the extended matrix of the system:

$$[A|\vec{b}] = [\vec{a}_1, \dots, \vec{a}_n | \vec{b}]$$

The system $A\vec{x} = \vec{b}$ is solvable iff the rank of A equals the rank of the extended matrix $[A|\vec{b}]$, i.e.:

$$\text{rank} A = \text{rank}[A|\vec{b}] =: r$$

The solution is unique if the rank of the two matrices equals num of unknowns ($r = n$).

1.6 Properties of squared matrices

Let $A \in R^{n \times n}$ be a square matrix. The following conditions are equivalent and characterize when a matrix A is **invertible** or **non-singular**:

- The matrix A has an inverse
- $\text{rank } A = n$
- $\det(A) \neq 0$
- The null space $N(A) = \{\vec{x} : A\vec{x} = 0\}$ is trivial
- All eigenvalues of A are nonzero
- For each \vec{b} the system of equations $A\vec{x} = \vec{b}$ has precisely one solution

1.7 Generalized inverse of a matrix $A \in R^{n \times m}$ is a matrix $G \in R^{m \times n}$ such that

$$AGA = A$$

Let G be a generalized inverse of A. Multiplying $AGA = A$ with A^{-1} from the left and the right side we obtain:

$$\text{LHS: } A^{-1}GAA^{-1} = IGI = G$$

$$\text{RHS: } A^{-1}AA^{-1} = IA^{-1} = A^{-1}$$

where I is the identity matrix. The equality LHS=RHS implies that $G = A^{-1}$.

Every matrix $A \in R^{n \times m}$ has a generalized inverse. When computing a generalized inverse we come across two cases:

1. $\text{rank } A = \text{rank } A_{11}$ where

$$A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$$

And $A_{11} \in R^{r \times r}$, $A_{12} \in R^{r \times (m-r)}$, $A_{21} \in R^{(n-r) \times r}$, $A_{22} \in R^{(n-r) \times (m-r)}$. We claim that

$$G = \begin{bmatrix} A_{11}^{-1} & 0 \\ 0 & 0 \end{bmatrix}$$

where 0s denote zero matrices of appropriate sizes, is the generalized inverse of A.

2. The upper left $r \times r$ submatrix of A is **not** invertible.

One way to handle this case is to use permutation matrices P and Q , such that

$$PAQ = \begin{bmatrix} \tilde{A}_{11} & \tilde{A}_{12} \\ \tilde{A}_{21} & \tilde{A}_{22} \end{bmatrix},$$

$\tilde{A}_{11} \in R^{r \times r}$ and $\text{rank } \tilde{A}_{11} = r$. By case 1 generalized inverse of PAQ equals to

$$(PAQ)^g = \begin{bmatrix} \tilde{A}_{11}^{-1} & 0 \\ 0 & 0 \end{bmatrix}$$

Thus $(PAQ)(PAQ)^g(PAQ) = PAQ$ holds, Multiplying from the left by P^{-1} and from the right by Q^{-1} we get:

$$A(Q(PAQ)^g P)A = A$$

So

$$Q(PAQ)^g P$$

is a **generalized** inverse of A.

Algorithm for computing A^g :

1. Find any nonsingular submatrix B in A of order $r \times r$,
2. in A substitute
 - elements of submatrix B for corresponding elements of $(B^{-1})^T$,
 - all other elements with 0
3. the transpose of the obtained matrix is generalized inverse G

solutions:

Let $A \in R^{n \times m}$ and $\vec{b} \in R^m$. If the system $A\vec{x} = \vec{b}$ is solvable (that is, $\vec{b} \in C(A)$) and G is a generalized inverse of A, then $\vec{x} = G\vec{b}$ is a solution of the system. Moreover, all solutions of system are exactly vectors of the form

$$x_z = G\vec{b} + (GA - I)z$$

1.8 The Moore-Penrose generalized inverse

The MP inverse of $A \in R^{n \times m}$ is any matrix $A^+ \in R^{m \times n}$ satisfying the following four conditions:

1. A^+ is a generalized inverse of A: $AA^+A = A$
2. A is a generalized inverse of A^+ : $A^+AA^+ = A^+$
3. The square matrix $AA^+ \in R^{n \times n}$ is symmetric: $(AA^+)^T = AA^+$

4. The square matrix $A^+A \in R^{m \times m}$ is symmetric: $(A^+A)^T = A^+A$

Properties:

- If A is a square invertible matrix, then it $A^+ = A^{-1}$
- $((A^+))^+ = A$
- $(A^T)^+ = (A^+)^T$

Construction of the MP inverse (4 cases):

1. $A^T A \in R^{m \times m}$ is an invertible matrix ($m \leq n$)

$$A^+ = (A^T A)^{-1} A^T$$

2. AA^T is an invertible matrix ($n \leq m$)

$$A^+ = A^T (AA^T)^{-1}$$

3. $\Sigma \in R^{n \times m}$ is diagonal matrix of the form

$$\Sigma = \begin{bmatrix} \sigma_1 & & \\ & \ddots & \\ & & \sigma_n \end{bmatrix}$$

Then the MP inverse is:

$$\Sigma = \begin{bmatrix} \sigma_1^+ & & \\ & \ddots & \\ & & \sigma_n^+ \end{bmatrix}$$

$$\text{where } \sigma_i^+ = \begin{cases} \frac{1}{\sigma_i} & \sigma_i \neq 0 \\ 0 & \sigma_i = 0 \end{cases}$$

4. a general matrix A (using SVD)

$$A^+ = V \Sigma^+ U^T$$

1.9 SVD computation

1. Compute the eigenvalues and an orthonormal basis consisting of eigenvectors of the symmetric matrix $A^T A$ or AA^T (depending on which of them is of smaller size)
2. the singular values of the matrix $A \in R^{n \times m}$ are equal to $\sigma_i = \sqrt{\lambda_i}$
3. the left singular vectors are the corresponding orthonormal eigenvectors of AA^T
4. the right singular vectors are the corresponding orthonormal eigenvectors of $A^T A$
5. If u (resp. v) is a left (resp. right) singular vector corresponding to the singular value σ_i , then $v = Au$ (resp. $u = A^T v$) is a right (resp left) singular vector corresponding to the same singular value

6. the remaining columns of U (resp. V) consist of an orthonormal basis of the kernel (i.e., the eigenspace of $\lambda = 0$) of AA^T (resp. $A^T A$)

1.10 General computation of A^+

1. For $A^T A$ compute its **nonzero** eigenvalues $\lambda_1 \geq \dots \geq \lambda_r > 0$, and the corresponding orthonormal eigenvectors v_1, \dots, v_r , and form the matrices:

$$S = \text{diag}(\sqrt{\lambda_1}, \dots, \sqrt{\lambda_r}) \in R^{r \times r}$$

$$V = [v_1 \dots v_r] \in R^{m \times r}$$

2. Put the vectors

$$u_1 = \frac{Av_1}{\sigma_1}, \dots, u_r = \frac{Av_r}{\sigma_r}$$

in the matrix $U = [u_1, \dots, u_r]$

3. Then $A^+ = V \Sigma^+ U^T$

1.11 Undetermined systems

Let $A \in R^{n \times m}$, where $m > n$. A system of equations that has more variables than constraints. Typically such system has infinitely many solutions, but it may happen that it has no solutions \rightarrow such system is undetermined.

1. An undetermined system of linear equations $Ax = b$ is solvable iff $AA^+b = b$.
2. If there are infinitely many solutions, the solution A^+b is the one with the smallest norm, i.e

$$\|A^+b\| = \min\{\|x\| : Ax = b\}$$

Moreover, it is the unique solution of smallest norm.

1.12 Overdetermined systems

Let $A \in R^{n \times m}$, where $n > m$. A system of equations that has more constraints than variables. Typically such system has no solutions, but it may happen that it may have one or even infinitely many solutions.

2 Nonlinear Models

2.1 Definition

Given is a sample of data points $(x_1, y_1), \dots, (x_m, y_m) \in R^n \times R$. We are searching for

$$F : R^n \times R^p \rightarrow R,$$

$$F(x_i, a_i, \dots, a_p) = y_i,$$

$$i = 1, \dots, m$$

which can be nonlinear in parameters a_i, \dots, a_p .

Examples: exponential decay/growth, Gaussian model, logistic model,...

We introduce a vector function

$$G : R^p \rightarrow R^m, G = (g_1, \dots, g_m)$$

$$g_1(a_1, \dots, a_p) = F(x_1, a_1, \dots, a_p) - y_1$$

$$\dots$$

$$g_m(a_1, \dots, a_p) = F(x_m, a_1, \dots, a_p) - y_m$$

Solving $F(x_i, a_i, \dots, a_p) = y_i$, is equivalent to solving $G(a_1, \dots, a_p) = (0, \dots, 0)$. Newton's method will solve the latter system.

Rate of convergence

Let us say that e_i is the i -th error of some method and α is a zero we are searching for. We can show that $|e_{i+1}| \leq C \cdot \frac{1}{2}|e_i|^r$, where $C \in R > 0$ and $r \in N$. r is the rate of convergence of the method.

2.2 Case 1

$p = m = 1$

$g : R \rightarrow R$. We have to find a zero of the function g (let's call it α).

Tangent method

1. Select an initial approximation x_0 .
2. Iterate $x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}, n = 0, 1, \dots$ until it converges to α

Problem: the convergence of this method really needs a good initial approximation.

Theorem: The sequence x_i from the tangent method converges with the rate $r = 2$ if:

- $f'(\alpha) \neq 0$
- x_0 was close enough to α

Bisection

1. Start with x_0, x_1 , such that $f(x_0)f(x_1) < 0$.
2. Compute $x_2 = \frac{x_0 + x_1}{2}$.
3. Choose $[x_0, x_2]$ if $f(x_0)f(x_2) < 0$, otherwise choose $[x_2, x_1]$.
4. Repeat.

Bisection always converges, but it is slow.

Rate of convergence: We can show that $|e_{i+1}| \approx \frac{1}{2}|e_i|$. From here we can see that bisection has a linear rate of convergence.

Stopping criteria Iterative method: x_0, \dots, x_n, x_{n+1} . How do we find n (how many steps of a method do we have to make)? This has to be true:

- $\frac{|x_{n+1} - x_n|}{|x_n|} < \text{tolerance}$
- $|f(x_{n+1})| < \text{tolerance}$

Tolerance is usually 10^{-10} .