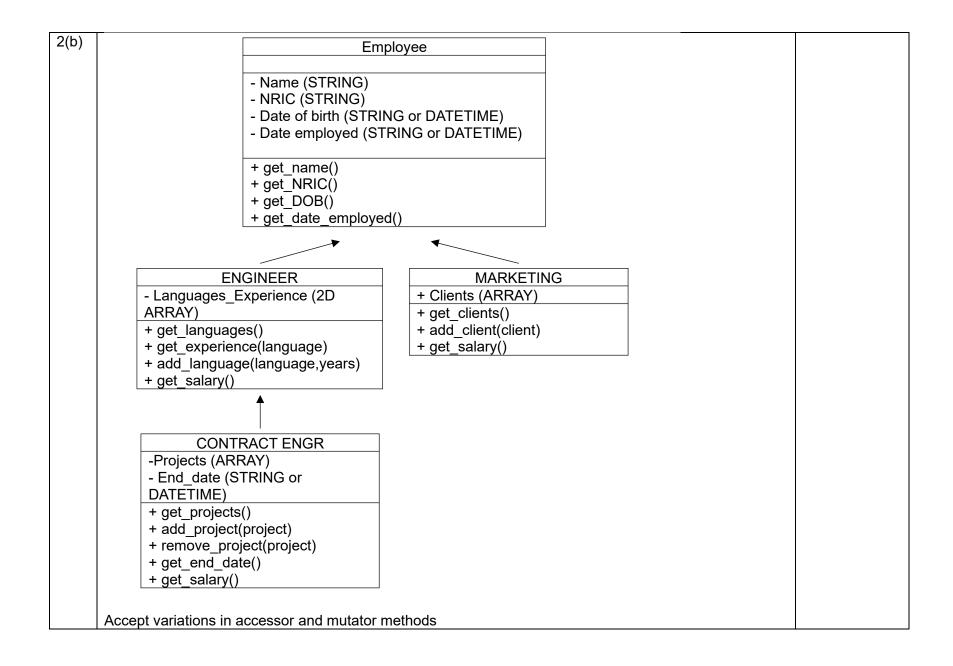
1(i)		110000 ₂									
	5/ ₁₀ =	111001 ₂									
		t 4 digits interpreted as a bir	nary r	numbe	er con	vert to	the b	ase 10) digit	they r	represent (0000 ₂ = 0 ₁₀ ,
	1001 ₂ =										
(ii)		here are 128 characters, 7 b									
(iii)		k digit is an addition digit sud used). This allows single-di						ts is e	ven (c	or odd,	depending on the
(d)	There a	are more than 128 character e languages, there would ne aracters.	s bei	ng use	ed by t	the lar	nguag				
2(a)	Full tab										_
		Software engineer	Υ	Y	Υ	Y	N	N	N	N	
	Conditions	Contract based	Y	Υ	N	N	Υ	Y	N	N	-
	Con	Been with company more than 15 years	Υ	N	Υ	N	Υ	N	Υ	N	
		35 days of leave			X						-
		28 days of leave	Х			Х					-
	Actions	25 days of leave					X		X		
	A	21 days of leave		Х							
		18 days of leave						X		Х	
	of engi	ot necessary to create a sepneer then they are in marketing instead.									

	Software engineer	Υ	Υ	Υ	Υ	N	N
Conditions	Contract based	Υ	Υ	N	N	-	-
Cond	Been with company more than 15 years	Υ	N	Υ	N	Υ	N
	35 days of leave			Х			
	28 days of leave	Х			Х		
Actions	25 days of leave					X	
Ac	21 days of leave		Х				
	18 days of leave						X



2(c)	Subclasses inherit attributes and methods from their parent classes. For example, contract engineers inherit the list of programming languages and the years of experience from the parent engineer class. The	
0(-1)	code is reused, so that changes can be easily made and debugging is easily done.	
2(d)	The salary calculation is different for each type of employee. However, using the same method name for	
2(-)	all of them makes it easy to run code to compute the salary regardless of what type of employee they are.	
2(e)	The NRIC number and date of birth should be private attributes as those are personal data.	
2(f)	A 2-dimensional array where each row consists of two elements, the programming language and the number of years' experience with it.	
3(a)	Local DNS checks if URL is inside its jurisdiction, sends back IP address to the browser if it is.	
0(4)	2. If URL is not in its domain, it looks up URL in its cache to see if IP address is there. If it is, the IP	
	address is sent back to the browser	
	3. If not, the DNS sends out a request to a root server, which provides an address for a DNS server	
	with jurisdiction over the top-level domain, which can provide an address for a DNS server for the	
	next level domain, and so on, until a server which can provide the IP address is found	
	4. The DNS adds this URL and IP address to its cache, and sends the IP address to the browser	
3(b)(2 ³² (about 4.2 billion)	
i)` ^`		
3(b)(There are more than 4.2 billion unique devices.	
ii)		
3(b)(There are $16^{32} = 2^{128}$ possible addresses, which is far more than the present number of devices, and will	
iii)	be for a long time to come.	
3(c)(1. A message from one computer to another is broken up into individual (numbered) packets and each	
i)	packet is sent separately.	
	2. The route for each packet is not determined in advance. It is sent from one computer to another in	
	the same network, which then decides to send it to another computer which may be nearer to the	
	destination, and so on, until it reaches the destination. Packets thus travel by different routes.	
	The destination computer receives the packets and assembles them in the correct order.	
3(c)(If computer is down, a connection between two computers is damaged, the packet can be rerouted to the	
ii)	destination by another route. Since the route is not predetermined, the path can be modified to work around	
	changes (damage) in the network.	
3(d)	 A malicious user sends many repeated requests to the server. 	
	2. This overwhelms the server by using up available resources (memory and bandwidth) on the server,	
	causing it to overload.	
	Legitimate users are not able to use the server, and their normal operations are disrupted.	
4(a)	Any two of: Worm, virus, trojan, ransomware, spyware	
4(b)	Downloading from FTP or webpage	

	2. Email att	achn	nent							
4(c)	1. Firewall									
, ,	2. Antivirus									
4(d)				ility to k	keep softw	are up to	date against latest kind	d of vi	ruses, and push	
	updates o									
							lvertised, and not make fa			
			ould tes	t softwa	are thorou	ghly befo	re release, and resolve	bugs	that they have	
5 (-)	discovere		4: 0	-1 / А А Г		NITEOED	\			
5(a)	1 PROCEDURE			τ(Arr: Ar	RRAY OF I	INTEGER	.)			
	2 N ← LENGT 3	ПП(А	u1 <i>)</i>							
	4 FORi←21	O N	I DO							
	5 key ← Ari		100							
	6 j ← i - 1	1.1								
	7									
	9 WHILE j >	> 0 A	ND Arr[j]	> key						
	10 Arr[j +	1] ←	- A[j]							
	11 j←j-									
	12 ENDWH	ILE								
	13									
	13 Arr[j + 1]	← K	ey							
	15 ENDFOR	חור	_							
	16 ENDPROCEI	JUK	E							
5(b)	O(N ²)									
5(c)	` '	alrea	dv sorte	d in asce	endina ord	er. The co	ondition for the while loop	from li	ne 9 to line 12	
,							ted. The time complexity			
5(d)					•					
		i	key	left	right	mid	Α	j		
		2	1	1	1	1	3,1,6,5,4,2,7	1		
							1,3,6,5,4,2,7			
		3	6	1	2	1	1,3,6,5,4,2,7	2		
							1,3,6,5,4,2,7			

			7	•						
		4	5	1	3	2	1,3,6,5,4,2,7	3		
							1,3,5,6,4,2,7			
		5	4	1	4	2	1,3,5,6,4,2,7	4		
							1,3,4,5,6,2,7			
		6	2	1	5	3	1,3,4,5,6,2,7	5		
							1,2,3,4,5,6,7			
		7	7	1	6	3	1,2,3,4,5,6,7	6		
							1,2,3,4,5,6,7			
5(e)	Insertion sort us	es lir	near sear	ch to fin	d the corre	ect positio	n for the key to be inserte	d into.		
T/\$\	taken to find the Binary insertion Note: For insertion O(log i) steps. He take i steps in the We make a total This occurs whe	posi sort t ng th lower e wo l of N	tion. faster that ne i-th ele ver, to in- orst case I insertion e array is	ement in sert the (when v	ion sort. its correct element, we we have to o, the wors	position i ve need to insert at a t-case tin	t position. For a large list, in the sorted, finding the position all the elements from the starting position). The complexity of binary insignorder.	osition m pos	(pos) will take to i-1. This will	
5(f) 5(g)	Line 12: IF A[mid O(N ²)	d] < k	кеу							
5(h)	\ /						es data using only a sma	all, con	stant amount of	
	An algorithm is o	onsi	dered no	t in-plac	e if it requi	res additio	onal memory that scales v	vith the	size of the input	

5(i)	Advantage: Minimal additional memory usage, which can be important in memory-constrained environments.	
	Disadvantage: The original data is changed, which might not be desirable if the original data needs to be preserved.	
6(a)	Any 3 of the below points and a conclusion Hash table has constant search/insert/delete time O(1), while BST has O(lgn), so hash table search is faster	
	Collisions might occur in hash tables that increases the search time, and it could be O(n) in the worst case.	
	BST could have O(n) search time if the tree is unbalanced, while collisions are rare when the hash function is good.	
	BST can get list of sorted items by doing in-order traversal, but not for hash table	
	BST is more memory efficient as it does not require more memory than necessary but hash tables require a lot more memory than required to prevent collisions.	
6(b)	has a time complexity of O(1)	
	 returns a unique output when provided a unique input uses the entire input to determine the address 	
6(c)	10, 40, 30, 90, 80, 50	
6(d)	Function search(arr, Root, target) IF Root == None THEN // 1 RETURN False ELSE IF target == arr[Root][0] THEN // 1 RETURN True ELSE IF target < arr[Root][0] THEN // 1 RETURN search(arr, arr[Root][1], target) ELSE IF target > arr[Root][0] THEN // 1 RETURN search(arr, arr[Root][2], target)	

		_
6(e)	Search(arr, 8, 10) Return True Search(arr, 0, 10) Return True Return True Return True	
6(f)	Iterative algorithm has constant space requirement regardless of how long the list is. The amount of memory space required for recursive algorithm is depending on the number of recursive calls made. More recursive calls (which usually occurs when list is longer) leads to greater memory use.	
7(a)	Member(MemberID, MemberName, Email, phone, points)	
	MemberEvents(<u>MemberID*</u> , <u>EventName*</u>) Events(<u>EventName</u> , distance, location, number1*, number2*, number3*, organizerID*) Alternative answer Member(<u>MemberID</u> , MemberName, Email, phone, points) Events(<u>EventName</u> , distance, location, organizerID*)	
	MemberEvents(<u>MemberID*</u> , <u>EventName*</u> , points)	
7(b)	MemberID in MemberEvents ensures each event is joined by a valid member. EventName in MemberEvents ensures each member joins a valid event.	
7(c)	Member MemberEvents Events	

7(d)	Ensure all non-key attributes are fully functional dependent on the entire primary key in each table; hence the tables are in 2NF.	
	For the 2NF tables, ensure that there is no transitive dependency for all the non-primary key attributes in the tables.	
7(e)	SELECT Member.MemberName FROM Member INNER JOIN MemberEvents ON Member.MemberID = MemberEvents.MemberID WHERE EventName = "Charlestown Marathon 2002" AND Member.points < 20	