

SORTING WITH STL

Introduction

- STL comes with good algorithms that can satisfy all our sorting needs
- Basic function is sort() with complexity O(n logn) that comes in two versions:
 - o sort(begin, end)
 - Sorts range [begin, end) in ascending order using operator<
 - sort(begin, end, comparator)
 - Sorts range [begin, end) using comparator function
 - Comparator function takes two objects and returns whether first object should come before the second
 - Similar, but not the same as with the priority queue (there the comparator was a structure)

Strana • 3



An example of ascending sorting

2

Example of descending and object sorting

```
struct Rectangle {
    int a;
    int b;
    Rectangle(int a, int b) {
        this->a = a;
        this->b = b;
    }
    int area() {
        return this->a * this->b;
    }
};

bool descending(int a, int b) {
    return a > b;
}

bool asc_rectangles(Rectangle a, Rectangle b) {
    return a.area() < b.area();
}</pre>
```

Example of descending and object sorting

```
int main() {
    int numbers [] = { 3, 1, 5, 2, 4};
    sort(numbers, numbers + 5, descending);
    for (int i = 0; i < 5; i++) {
        cout << numbers[i] << ' ';
    }
    cout << endl;

    vector<Rectangle> p({ Rectangle(10, 10), Rectangle(2, 2), Rectangle(4, 4) });
    sort(p.begin(), p.end(), asc_rectangles);
    for (auto it = p.begin(); it != p.end(); ++it) {
        cout << it->area() << ' ';
    }
    cout << endl;
    return 0;
}
</pre>
```

Problem

- Prepare a vector with shuffled numbers from 1 to 2 million and sort it ascending. Let's display how long the sorting takes.
- Solution:

Algorithms under the hood

- Visual Studio implements the sort() function using Intro sort and Insertion sort
- It works like this:
 - o If < 32 items are sorted, use Insertion sort
 - If sorting >= 32 elements, use Intro sort
 - Start with Quick sort
 - As a pivot, it uses the median of the first, middle and last elements
 - If Quick sort breaks the defined recursion depth limit, it gets cancelled and Heap sort is used instead
 - Recursion depth limit is usually 2*logn

Strana ■ 8



Stable sorting

• If we have the initial elements and sort them by key:

12	49	12	3
Mireau Miric	Yuro Yurich	Anna Anicci	Schtepf Stefi

•Unstable sorting can produce a result:

3	12	12	49
Schtepf Stefi	Anna Anicci	Mireau Miric	Yuro Yurich

■ While stable sorting is guaranteed to yield this result:

	3	12	12	49
	Schtepf Stefi	Mireau Miric	Anna Anicci	Yuro Yurich
Strana •	9			



Example of stable sorting usage

- Example:
 - o We have students sorted by last name and then first name
 - o We want to also sort them by the number of points on the exam
 - o In addition, we want students who have the same number of points to keep their alphabetical order

Student	Score	Student	Score	Student	Score
Ana Anić	51	Željko Željkić	93	Branka Brankić	93
Branka Brankić	93	Branka Brankić	93	Željko Željkić	93
Bruno Brunić	43	Iva Ivić	55	Iva Ivić	55
Iva Ivić	55	Ana Anić	51	Ana Anić	51
Željko Željkić	93	Bruno Brunić	43	Bruno Brunić	43
Strana • 10 Sorted by names		Unstable		Stable	ALGE

5

Stable sorting

- In stable sorting, the elements retain their relative position after sorting
- •Algorithm sort() is not stable
- Function stable_sort() is a stable version that can be used instead of sort()
 - Same usage sort()
 - Has complexity O(n log²n)
 - If there is enough memory, it can complete in O(n logn)
 - o Uses Merge sort in combination with Insertion sort

Strana = 11



Sorting a list

- Some containers do not make sense to sort (queue, stack, map, etc.)
- Function sort() can sort all other containers instead listsList does not have a random access via [] nor at()
- •list<T>::sort() is also a stable sort based on the Merge
 sort
 - $\circ\,$ Has the same complexity



Example of sorting a list

```
bool desc_rectangles(Rectangle a, Rectangle b) {
    return a.area() > b.area();
}
int main() {
    list<Rectangle> l({
        Rectangle(10, 10),
        Rectangle(2, 2),
        Rectangle(4, 4) });

    l.sort(desc_rectangles);

    for (auto it = l.begin(); it != l.end(); ++it) {
            cout << it->area() << ' ';
        }
        cout << endl;

Strana-13 return 0;
}</pre>
```

Sorting a part of the range

- partial sort(begin, middle, end)
 - The elements in front of the middle are the smallest elements in the entire range and are sorted
 - o There is no guarantee how all other elements will be arranged
 - Uses Heap sort, complexity is O(n logm), where m is the distance from begin to middle



Problem

- Display three lowest elements from the vector.
- Solution:

```
vector<int> v({ 8, 3, 1, 5, 2, 4, 6, 7, 9, 10, 13, 11, 15,
12, 14, 16, 17, 19, 20, 18 });

partial_sort(v.begin(), v.begin() + 3, v.end());

for (auto it = v.begin(); it != v.begin() + 3; ++it) {
    cout << *it << ' ';
}
cout << endl;</pre>
```

Strana • 15



Sorting check

- ■is_sorted(begin, end)
 - o Return whether the range [begin, end) is sorted
- •is_sorted_until(begin, end)
 - Returns the iterator to the first element in the range [begin, end) that is not sorted
 - o If it returns the end(), the entire range is sorted

Strana = 16



Example

```
srand(unsigned(time(0)));
vector<int> v;
for (int i = 1; i <= 35; i++) { v.push_back(i); }

random_shuffle(v.begin(), v.end());
for (auto it = v.begin(); it != v.end(); ++it) {
        cout << *it << ' ';
}
cout << endl;

cout << "Sorted elements: ";
auto first_unsorted = is_sorted_until(v.begin(), v.end());
for (auto it = v.begin(); it != first_unsorted; ++it) {
        cout << *it << ' ';
}
cout << endl;</pre>
```

Sorting nth element

- •nth_element(begin, nth, end) in linear complexity
 modifies the range in a way:
 - At position nth is the value that would be there if the array was sorted (that value is now in place)
 - o The elements in front of *nth* are certainly not larger than the element on *nth*
 - The elements behind nth are certainly not smaller than the element on nth
- Task: calculate which number is the third smallest number in the range of shuffled numbers from 1 to 2 million. Let's compare the execution speed compared to sorting the entire range.

Strana = 18



Solution

Conclusion

- ■The main tool
 - o Function sort() sorts the range according to the criteria
- We want to sort while maintaining the previous relationship of equal elements
 - Function stable_sort(), but performance might be lower
- We want to find the value that would be in the nth place of the sorted elements
 - o Function nth_element()
- We want to find all the values from the beginning to the *nth* position of the sorted elements

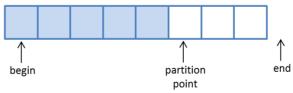
```
Function partial_sort()
```





Partitioning

•Sometimes we need to effectively divide the elements of a container into those that meet the condition and those that do not.



- In this example:
 - o Blue-marked elements satisfy the condition
 - o White-marked are those who do not satisfy the condition
 - The partitioning point is the first element that does not satisfy the condition



Partitioning functions

- The basic functions are:
 - o partition(begin, end, predicate)
 - Performs partitioning according to the predicate and returns the iterator to the partitioning point
 - A predicate is a function that receives a value and returns a bool
 - Complexity is O(n)
 - o is partitioned(begin, end, predicate)
 - Checks if the range is partitioned
 - o partition point(begin, end, predicate)
 - Returns the iterator to the partitioning point

Strana = 23



Problem

- Let's make a vector with shuffled numbers from 1 to 2 million and partition it so that first come even and then odd. Display the first odd number.
- Solution:

```
bool is_even(int n) {
    return n % 2 == 0;
}
int main() {
    srand(unsigned(time(0)));
    vector<int> v;
    for (int i = 1; i <= 2000000; i++) { v.push_back(i); }
    random_shuffle(v.begin(), v.end());
    auto start = chrono::high_resolution_clock::now();
    auto pp = partition(v.begin(), v.end(), is_even);
    auto end = chrono::high_resolution_clock::now();
    cout << "Partition point: " << *pp << endl;

Strana*24 return 0;</pre>
```



BINARY SEARCH

Strana = 25



Binary search

- Binary search is a procedure in which we look for a value in a sorted container in complexity O(logn)
- ■Two versions:
 - o binary search(begin, end, value)
 - Returns true if the requested value exists within the specified sorted range
 - o binary_search(begin, end, value, comparator)
 - Returns true if the requested value exists in the specified sorted range, but using the same comparator function as when sorting

Strana = 26



Problem

- Prepare a vector with shuffled values of 1 to 2 million. Look for the value of 1,234,456 in the following ways and discuss durations:
 - o Linear search in unsorted vector
 - o Linear search on a sorted vector
 - o Sort + binary search
 - o Binary search without sorting

