

CH-232-A

Answers to ICS 2020 Problem Sheet #6

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1. The two elementary Boolean functions \rightarrow (implication) and \neg (negation) combined are sufficient to express all the possible Boolean functions. But individually, they are not universal function and they can't be used to sufficiently express the rest of the Boolean functions like that of NOT-AND and NOT-OR.

i.e. NOT-AND and NOT-OR are universal functions since they can be used to derive all elementary Boolean functions.

But combined:

$$\text{Let } S = \{\rightarrow, \neg\},$$

For \neg , it is obvious why it is expressible – It is a member of S and it itself is an elemental Boolean function. The remaining \rightarrow can express the terms of \wedge and \vee with \neg .

Such that:

$$x \vee y = \neg x \rightarrow y$$

De Morgans law

$$x \wedge y = \neg(x \rightarrow \neg y)$$

...

Thus, $\{\rightarrow, \neg\}$ are sufficient to express all the possible Boolean functions but not individually.

2.

a) $\varphi(A, B) = (\neg A \vee \neg B) \wedge (\neg A \vee B) \wedge (A \vee \neg B)$

$$\varphi(A, B) = (A' + B')(A' + B)(A + B')$$

(writing it in a simplified form)

$$(A'A' + A'B + A'B' + B'B)(A + B')$$

(distribution)

$$(A' + A'(B + B') + 0)(A + B')$$

(idempotent law)

$$(A' + A'(1) + 0)(A + B')$$

(complement)

$$(A')(A + B')$$

(idempotent law)

$$(A'A)(A'B')$$

(distribution)

$$(0 + A'B')$$

(complement)

$$A'B'$$

$$\therefore \neg A \wedge \neg B$$

(re-writing)

b) $\varphi(A, B, C) = (A \wedge \neg B) \vee (A \wedge \neg B \wedge C)$

$$\varphi(A, B, C) = (AB') + (AB'C)$$

(writing it in a simplified form)

$$(AB')(1 + C)$$

(distribution)

$$(AB')(1)$$

(domination)

$$AB'$$

$$\therefore A \wedge \neg B$$

(re-writing)

c) $\varphi(A, B, C, D) = (A \vee \neg(B \wedge A)) \wedge (C \vee (D \vee C))$

$$\varphi(A, B, C, D) = (A + (BA)').(C + (D + C))$$

(writing it in a simplified form)

$$(A + (BA'))(C + D)$$

(idempotent law)

$$(A + B' + A')(C + D)$$

(De Morgan's law)

$$\begin{aligned}
&(1 + B')(C + D) && \text{(Complement law)} \\
&(1)(C + D) && \text{(Domination)} \\
&C + D && \text{(Identity law)} \\
&\therefore C \vee D && \text{(re-writing)}
\end{aligned}$$

$$\begin{aligned}
\text{d) } \varphi(A, B, C) &= (\neg(A \wedge B) \vee \neg C) \wedge (\neg A \vee B \vee \neg C) \\
\varphi(A, B, C) &= ((AB)' + C')(A' + B + C') && \text{(writing it in a simplified form)} \\
&= (A' + B' + C')(A' + B + C') && \text{(De Morgan's law)}
\end{aligned}$$

By using distributions, Let us first distribute A' with $(A' + B + C')$ and simplify it

$$\begin{aligned}
&A'A' + A'B + C'A' \\
&A' + A'B + C'A' && \text{(Idempotent law)} \\
&A'(1 + B) + A'C' && \text{(distribution)} \\
&A' + A'C' && \text{(dominance law)} \\
&A'(1 + C') && \text{(distribution law)} \\
&\therefore A'
\end{aligned}$$

Secondly distribute B' with $(A' + B + C')$ and simplify it

$$\begin{aligned}
&B'A' + B'B + B'C' \\
&\therefore B'A' + B'C' && \text{(complement law)}
\end{aligned}$$

Thirdly distribute B' with $(A' + B + C')$ and simplify it

$$\begin{aligned}
&C'A' + C'B + C'C' \\
&C'A' + C'B + C' && \text{(idempotent law)} \\
&C'A' + C'(B + 1) \\
&C'A' + C' && \text{(dominance law)} \\
&C'(A + 1) && \text{(distribution)} \\
&\therefore C'
\end{aligned}$$

Combinin the three distributions we get $A' + B'A' + B'C' + C'$

$$\begin{aligned}
&A'(1 + B') + C'(B' + 1) && \text{(dominance law)} \\
&A' + C' \\
&\therefore \neg A' \vee \neg C && \text{(re-writing)}
\end{aligned}$$

$$\begin{aligned}
\text{e) } \varphi(A, B) &= (A \vee B) \wedge (\neg A \vee B) \wedge (A \vee \neg B) \wedge (\neg A \vee \neg B) \\
\varphi(A, B) &= (A + B)(A' + B)(A + B')(A' + B') && \text{(writing it in a simplified form)} \\
&= (AA' + AB + BA' + BB)(AA' + AB' + B'A' + B'B') \\
&= (0 + AB + BA' + BB)(0 + B'A + A'B' + B'B') && \text{(complement law)} \\
&= (AB + BA' + B)(B'A + A'B' + B') && \text{(idempotent law)} \\
&= (B(A + A') + B)(B'(A + A') + B') \\
&= (B(1) + B)(B'(1) + B') && \text{(complement law)} \\
&= (B)(B') && \text{(complement law)} \\
&\therefore 0
\end{aligned}$$

$$3. \varphi(P, Q, R, S) = (\neg P \vee Q) \wedge (\neg Q \vee R) \wedge (\neg R \vee S) \wedge (\neg S \vee P)$$

a)

	P	Q	R	S	$\neg P$	$\neg Q$	$\neg R$	$\neg S$	$\neg P \vee Q$	$\neg Q \vee R$	$\neg R \vee S$	$\neg S \vee P$	$(\neg P \vee Q) \wedge (\neg Q \vee R) \wedge (\neg R \vee S) \wedge (\neg S \vee P)$
1	0	0	0	0	1	1	1	1	1	1	1	0	1
2	0	0	0	1	1	1	1	0	1	1	1	1	0
3	0	0	1	0	1	1	0	1	1	1	0	0	0
4	0	0	1	1	1	1	0	0	1	1	1	1	0
5	0	1	0	0	1	0	1	1	1	0	1	0	0
6	0	1	0	1	1	0	1	0	1	0	1	1	0
7	0	1	1	0	1	0	0	1	1	1	0	0	0
8	0	1	1	1	1	0	0	0	1	1	1	1	0
9	1	0	0	0	0	1	1	1	0	1	1	1	0
10	1	0	0	1	0	1	1	0	0	1	1	1	0
11	1	0	1	0	0	1	0	1	0	1	0	1	0
12	1	0	1	1	0	1	0	0	0	1	1	1	0
13	1	1	0	0	0	0	1	1	1	0	1	1	0
14	1	1	0	1	0	0	1	0	1	0	1	1	0
15	1	1	1	0	0	0	0	1	1	1	0	1	0
16	1	1	1	1	0	0	0	0	1	1	1	1	1

\therefore Only two conditions satisfy the φ .

i. e. when $\varphi(P, Q, R, S) = (0, 0, 0, 0)$

and when $\varphi(P, Q, R, S) = (1, 1, 1, 1)$

b) From the truth table we see that the conditions that satisfy φ are when $\varphi(P, Q, R, S) = (0, 0, 0, 0)$ and $\varphi(P, Q, R, S) = (1, 1, 1, 1)$.

This means, when $\varphi(P, Q, R, S)$ is 0 and $\varphi(P, Q, R, S)$ is negated.

P	Q	R	S	$\neg P$	$\neg Q$	$\neg R$	$\neg S$	φ
0	0	0	0	1	1	1	1	1
1	1	1	1	0	0	0	0	1

$$(\neg P \wedge \neg Q \wedge \neg R \wedge \neg S) \text{ OR } (P \wedge Q \wedge R \wedge S)$$

$$\therefore \text{DNF of } X = (\neg P \wedge \neg Q \wedge \neg R \wedge \neg S) \vee (P \wedge Q \wedge R \wedge S)$$