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ICS 2020 Problem Sheet #3 Answers

Problem 3.1: Cartesian products

a) $(A \cap B) \times (C \cap D) = (A \times C) \cap (B \times D)$

$$\Rightarrow (x, y) \in (A \cap B) \text{ and } (C \cap D)$$

$$\Rightarrow x \in (A \cap B) \text{ and } y \in (C \cap D)$$

$$\Rightarrow \text{This would imply } x \in A \text{ and } x \in B, \text{ and also } y \in C \text{ and } y \in D$$

To show if $(A \times C) \cap (B \times D)$:

$$\Rightarrow (x \in A \text{ and } y \in C) \text{ and } (x \in B \text{ and } y \in D)$$

$$\Rightarrow (x, y) \in (A \times C) \cap (x, y) \in (B \times D)$$

$$\Rightarrow (x, y) \in (A \times C) \cap (B \times D)$$

Therefore $(A \cap B) \times (C \cap D) = (A \times C) \cap (B \times D)$

b) $(A \cup B) \times (C \cup D) = (A \times C) \cup (B \times D)$

$$(A \cup B) \times (C \cup D)$$

$$\Rightarrow \{(x, y) \mid x \in (A \cup B) \text{ and } y \in (C \cup D)\}$$

$$\Rightarrow \{(x, y) \mid x \in A \text{ or } x \in B \text{ and } y \in C \text{ or } y \in D\}$$

Four alternatives to what this means:

1st:

$$\Rightarrow \{(x, y) \mid (x, y) \in (A \times C) \text{ or } (x, y) \in (B \times D)\}$$

This would prove that the statement is actually correct. But we also have the other possibility since its and Or statement.

2nd:

$$\Rightarrow \{(x, y) \mid (x, y) \in (B \times C) \text{ or } (x, y) \in (A \times D)\}$$

3rd:

$$\Rightarrow \{(x, y) \mid (x, y) \in (B \times D) \text{ or } (x, y) \in (A \times C)\}$$

This would prove that the statement is actually correct.

4th:

$$\Rightarrow \{(x, y) \mid (x, y) \in (A \times D) \text{ or } (x, y) \in (B \times C)\}$$

This proves that it is not in fact the same.

Therefore: $(A \cup B) \times (C \cup D) \neq (A \times C) \cup (B \times D)$

Problem 3.2: Reflexive, symmetric, transitive

a) $R = \{(a, b) | a, b \in \mathbb{Z} \wedge |a - b| \leq 3\}$

i. **R is reflexive**

This is because there could exist an element in the relation that could have equal values. Since the condition for the relation is for a and b to be an integer and their absolute difference to be less than or equal to 3, then the condition will always be satisfied in the case of a and b having the same value.

ii. **R is Symmetric**

R is symmetric because since the difference is in absolute value, (a, b) could easily be (b, a) . i.e. the value would be the same.

iii. **R is not Transitive**

Transitive by definition is when $(a, b) \in R$ and $(b, c) \in R$, then $(a, c) \in R$.

This could not happen for this relation as the absolute difference of a and b if it exists in the relation and the absolute difference of b and c , if it exists in the relation, then this does not necessarily mean the absolute difference between a and c might exist in the relation.

b) $R = \{(a, b) | a, b \in \mathbb{Z} \wedge (a \bmod 10) = (b \bmod 10)\}$

i. **R is reflexive**

R is reflexive because there could exist an element in the relation that could have equal values.

ii. **R is Symmetric**

R is symmetric because for every a and b the relation could also have values with (b, a) that satisfy the condition.

iii. **R is Transitive**

R is transitive as $a \bmod 10 = b \bmod 10$ will mean that for whatever value $b \bmod 10$ is equal to, $a \bmod 10$ will also be equal to that number.

Problem 3.2: Proof by induction

Note: S and t are lists.

Proof of $\text{cnt } x (\text{con } s \ t) == (\text{cnt } x \ s) + (\text{cnt } x \ t)$

First lets just take the smallest possible value for s which is an empty list.

$S = []$

For the left side equation:

$\text{cnt } x (\text{con } [] \ t)$

$\text{Con } [] \ t = t$ --- from the Haskell code

Therefore:

$\text{cnt } x (\text{con } [] \ t) = \text{cnt } x \ t$

For the right-side equation:

$\text{cnt } x [] + \text{cnt } x \ t$

$\text{cnt } x []$ is 0 --- from the Haskell code

Therefore:

$(\text{cnt } x \ s) + (\text{cnt } x \ t) = \text{cnt } x \ t$

The proof holds for the smallest possible value.

The next smallest value I will take for s is a 1 item list and I will call it n.

For the left side equation:

$\text{cnt } x (\text{con } [n] \ t)$

$\text{con } [n] \ t = (n : t)$ --- from the Haskell code

Therefore:

$\text{cnt } x (\text{con } [n] \ t) = \text{cnt } x (n : t)$

if x happens to be equal to n, then:

$1 + \text{cnt } x \ t$

Else:

$0 + \text{cnt } x \ t$

For the right-side equation:

$\text{cnt } x [n] + \text{cnt } x \ t$

if x happens to be equal to n, then:

$1 + \text{cnt } x \ t$

Else:

$$0 + \text{cnt} \times t$$

This led me to conclude that $\text{cnt} \times (\text{con } s \ t) == (\text{cnt} \times s) + (\text{cnt} \times t)$ is true.