

1 Section 1

2 Section 2

Focus:

INTERIOR-POINT METHODS

JEAN KOUAGOU
ARIEL KEMOGNE
BLESSING BASSEY
GEDEON MUHAWENAYO

AFRICAN MASTERS FOR MACHINE LEARNING

13 12 2019

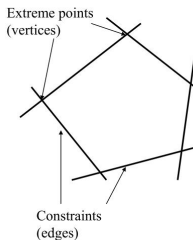


SECTION 1

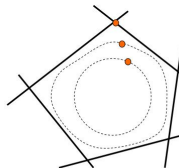
INTRODUCTION

An interior point method is a linear or nonlinear programming method that achieves optimization by going through the middle of the solid defined by the problem rather than around its surface.

Interior point methods



Simplex: search from vertex to vertex along the edges



Interior-point methods: go through the inside of the feasible space

© 2011 Daniel Kirschen and University of Washington

14

Figure: Simplex versus interior point method

INEQUALITY CONSTRAINED MINIMIZATION

Inequality constrained minimization

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b\end{array}$$

- f_i convex, twice continuously differentiable
- $A \in \mathbf{R}^{p \times n}$ with $\text{rank } A = p$
- we assume p^* is finite and attained
- we assume problem is strictly feasible: there exists \tilde{x} with $\tilde{x} \in \text{dom } f_0$, $f_i(\tilde{x}) < 0$, $i = 1, \dots, m$, $A\tilde{x} = b$ hence, strong duality holds and dual optimum is attained

LOGARITHMIC BARRIER

Approximation with the indicator function

$$\begin{array}{ll}\text{minimize} & f_0(x) + \sum_{i=1}^m I_{-}(f_i(x)) \\ \text{subject to} & Ax = b\end{array}$$

where $I_{-}(u) = 0$ if $u \leq 0$, $I_{-}(u) = \infty$ otherwise.

Approximation via logarithmic barrier

$$\begin{array}{ll}\text{minimize} & f_0(x) - (1/t) \sum_{i=1}^m \log(-f_i(x)) \\ \text{subject to} & Ax = b\end{array}$$

LOGARITHMIC BARRIER FUNCTION

$$\phi(x) = - \sum_{i=1}^m \log(-f_i(x)), \quad \text{dom } \phi = \{x | f_1(x) < 0, \dots, f_m(x) < 0\}$$

- convex (follows from composition rules)
- twice continuously differentiable, with derivatives

$$\nabla \phi(x) = \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla f_i(x)$$

$$\nabla^2 \phi(x) = \sum_{i=1}^m \frac{1}{f_i(x)^2} \nabla f_i(x) \nabla f_i(x)^T + \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla^2 f_i(x)$$

CENTRAL PATH

- for $t > 0$, define $x^*(t)$ as the solution of

$$\begin{array}{ll}\text{minimize} & tf_0(x) + \phi(x) \\ \text{subject to} & Ax = b\end{array}$$

(for now, assume $x^*(t)$ exists and is unique for each $t > 0$)

- central path is $\{x^*(t) | t > 0\}$

DUAL POINTS ON CENTRAL PATH

$x = x^*(t)$ if there exists a w such that

$$t \nabla f_0(x) + \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla f_i(x) + A^T w = 0, \quad Ax = b$$

- therefore, $x^*(t)$ minimizes the Lagrangian

$$L(x, \lambda^*(t), \nu^*(t)) = f_0(x) + \sum_{i=1}^m \lambda_i^*(t) f_i(x) + \nu^*(t)^T (Ax - b) \quad (1)$$

where we define $\lambda_i^*(t) = 1/(-t f_i(x^*(t)))$ and $\nu^*(t) = w/t$

DUAL POINTS ON CENTRAL PATH

$x = x^*(t), \lambda = \lambda^*(t), \nu = \nu^*(t)$ satisfy

1. primal constraints: $f_i(x) \leq 0, i = 1, \dots, m, Ax = b$
2. dual constraints: $\lambda \succeq 0$
3. approximate slackness: $-\lambda_i f_i(x) = 1/t, i = 1, \dots, m$
4. gradient of Lagrangian with respect to x vanishes:

$$\nabla f_0(x) + \sum_{i=1}^m \lambda_i \nabla f_i(x) + A^T \nu = 0$$

difference with KKT is that condition 3 replaces $\lambda_i f_i(x) = 0$

BARRIER METHOD

given strictly feasible x , $t := t^{(0)} > 0$, $\mu > 1$, tolerance $\epsilon > 0$
repeat

1. Compute $x^*(t)$ by minimizing $tf_0 + \phi$, subject to $Ax = b$
2. Update. $x := x^*(t)$
3. Stopping criterion. quit if $m/t < \epsilon$
4. Increase t . $t := \mu t$

- Terminates with $f_0(x) - p^* \leq \epsilon f_0(x^*(t)) - p^* \leq m/t$
- centering usually done using Newton's method, starting at current x
- choice of μ involves a trade-off: μ means fewer outer iterations
- more inner (Newton) iterations; typical values: $\mu = 10 - 20$
several heuristics for choice of $t^{(0)}$

BARRIER METHOD

given strictly feasible x , $t := t^{(0)} > 0$, $\mu > 1$, tolerance $\epsilon > 0$
repeat

1. Compute $x^*(t)$ by minimizing $tf_0 + \phi$, subject to $Ax = b$
2. Update. $x := x^*(t)$
3. Stopping criterion. quit if $m/t < \epsilon$
4. Increase t . $t := \mu t$

- Terminates with $f_0(x) - p^* \leq \epsilon f_0(x^*(t)) - p^* \leq m/t$
- centering usually done using Newton's method, starting at current x
- choice of μ involves a trade-off: μ means fewer outer iterations
- more inner (Newton) iterations; typical values: $\mu = 10 - 20$
several heuristics for choice of $t^{(0)}$

CONVERGENCE ANALYSIS

number of outer (centering) iterations: exactly

$$\left\lceil \frac{\log(m/(\epsilon t^{(0)}))}{\log \mu} \right\rceil$$

plus the initial centering step (to compute $x^*(t^{(0)})$)

CONVERGENCE ANALYSIS

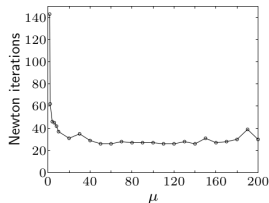
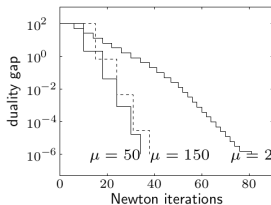
number of outer (centering) iterations: exactly

$$\left\lceil \frac{\log(m/(\epsilon t^{(0)}))}{\log \mu} \right\rceil$$

plus the initial centering step (to compute $x^*(t^{(0)})$)

Examples

inequality form LP ($m = 100$ inequalities, $n = 50$ variables)



Feasibility and phase I methods
feasibility problem: find x such that
$$f_i(x) \leq 0, \quad i = 1, \dots, m, \quad Ax = b$$

(2)

The packages `inputenc` and `FiraSans`^{1,2} are used to properly set the main fonts.

This theme provides styling commands to typeset *emphasized*, **alerted**, **bold**, *example text*, ...

FiraSans also provides support for mathematical symbols:

$$e^{j\pi} + 1 = 0.$$

¹<https://fonts.google.com/specimen/Fira+Sans>

²<http://mozilla.github.io/Fira/>

SECTION 2

These blocks are part of 1 slide, to be displayed consecutively.

Block

Text.

BLOCKS

These blocks are part of 1 slide, to be displayed consecutively.

Block

Text.

Alert block

Alert **text**.

BLOCKS

These blocks are part of 1 slide, to be displayed consecutively.

Block

Text.

Alert block

Alert **text**.

Example block

Example **text**.

This text appears in the left column and wraps neatly with a margin between columns.

Placeholder

Image

LISTS

Items:

- Item 1
 - ▶ Subitem 1.1
 - ▶ Subitem 1.2
- Item 2
- Item 3

Enumerations:

1. First
2. Second
 - 2.1 Sub-first
 - 2.2 Sub-second
3. Third

Descriptions:

First Yes.
Second No.

TABLE

Discipline	Avg. Salary
Engineering	\$66,521
Computer Sciences	\$60,005
Mathematics and Sciences	\$61,867
Business	\$56,720
Humanities & Social Sciences	\$56,669
Agriculture and Natural Resources	\$53,565
Communications	\$51,448
Average for All Disciplines	\$58,114

Table: Table caption

THANKS FOR USING **Focus!**

REFERENCES



DONALD E. KNUTH.

COMPUTER PROGRAMMING AS AN ART.

Commun. ACM, pages 667–673, 1974.



DONALD E. KNUTH.

TWO NOTES ON NOTATION.

Amer. Math. Monthly, 99:403–422, 1992.



LESLIE LAMPORT.

L^AT_EX: A DOCUMENT PREPARATION SYSTEM.

Pearson Education India, 1994.

BACKUP SLIDE

This is a backup slide, useful to include additional materials to answer questions from the audience.

The package `appendixnumberbeamer` is used to refrain from numbering appendix slides.