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# IN2002 Data Structures and Algorithms

## Lecture 4 – Pointers and Singly Linked Lists

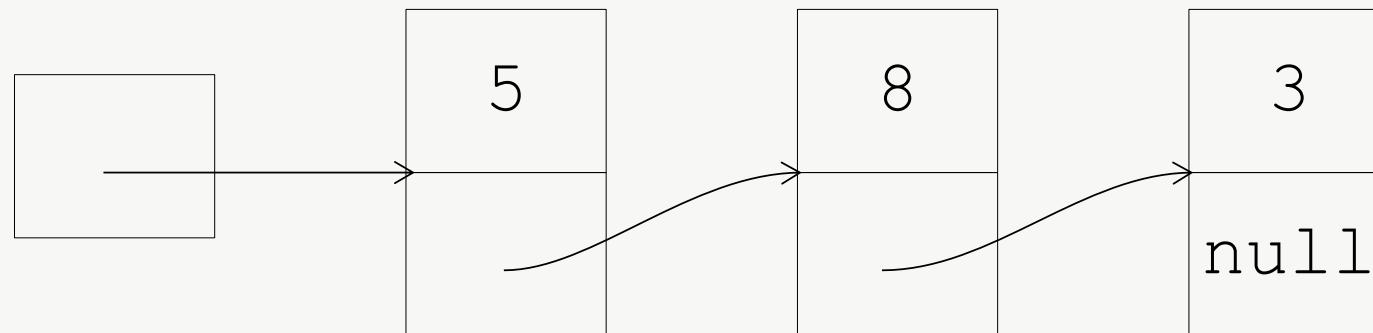
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Semester 1, 2018/19

# Learning Objectives

- Understand and be able to use the data structure singly linked lists
- Be able to understand, apply and develop algorithms to handle singly linked lists. Including:
  - Adding elements
  - Deleting elements
  - Traversing a list

# **Introduction**

# What is this?



# Database Records

- A record is a fixed collection of things of different types.
- E.g. records at a mobile phone provider:

Client ID	1234
Phone number	+44 7887 444444
Owner	Rupert Taylor
Start date	04/10/2001
Credit	£24.53
Type of contract	Ultra
International opt.	No
Free minutes left	42

# Records as Objects

```
class Mobile {  
    public int clientID;  
    public String number;  
    public String name;  
    public Date startdate;  
    public int credit;  
    public char contract;  
    public boolean intl;  
    public int seconds;  
    ...  
    public Mobile(int id, String nu, String na, ...) {  
        clientID = c; number = nu; name = na; ....  
    }  
}
```

# Pointers to Records

- In Java, variables cannot contain objects (record instances, in this case), but pointers to them
- Every reference type can have value null, which refers to no object at all

```
Mobile x = null;
```

- Several variables may refer to the same object:

```
Mobile y = new Mobile(1234,  
                     "+44 7887 4444444", "Rupert", ...);
```

```
Mobile z = y;
```

# Are two records equal?

Consider:

```
Mobile x = new Mobile(1234,  
                      "+44 7887 4444444", "Rupert", ...);  
Mobile y = new Mobile(1234,  
                      "+44 7887 4444444", "Rupert", ...);  
Mobile z = x;
```

x == z is

x == y is

x.compareTo(y), however, can be defined to return 0 to indicate that both objects contain the same values

# What would happen here?

Consider:

```
Mobile x = new Mobile(1234,  
                      "+44 7887 4444444", "Rupert", ...);  
  
Mobile y = new Mobile(1234,  
                      "+44 7887 4444444", "Rupert", ...);  
  
Mobile z = x;  
z.name = "Anne";  
y.clientID = 555;
```

What are the values of:

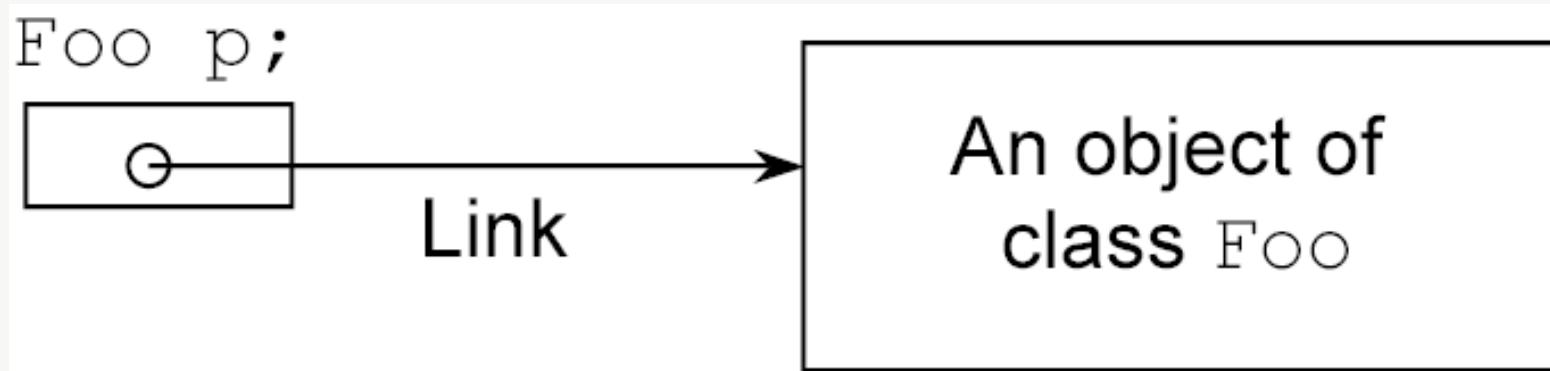
x.name	x.clientID
y.name	y.clientID
z.name	z.clientID

# What are links?

All object variables in Java are actually pointers

A pointer is a link to some object, i.e. a piece of memory:

- Links can be thought of as arrows, addresses or references
- They point at an object that is somewhere in the system's memory



# Pointers to an address

```
Foo p;
```

```
null
```

When you create and assign objects, the system allocates memory and assigns addresses

```
Foo p = new Foo();      memory address E1E42FB:
```

```
E1E42FB
```

New object of  
class Foo

```
Foo q = p;
```

```
E1E42FB
```

Creating a pointer to the same object  
at address **E1E42FB**

# Pointers are links

```
Foo p;  
  null
```

In Java you don't need to worry about the address, just what you are pointing at.

```
Foo p = new Foo();      memory address E1E42FB:
```



New object of  
class Foo

```
Foo q = p;
```



In official Java parlance, they are called  
*references*.

# Links in Java

**All** variables whose type is a class (or interface) of objects are actually pointers (“references”) to an instance of such class (or to null if empty).

This means that the **same** object can be assigned to different variables

... and, importantly, the same object’s methods and parameters can be accessed through different variables

# **Singly Linked Lists**

# Singly linked lists

How do you remove a value from the middle of an array?

How would you do it in real life, e.g. in a long written list?

In real life, to add something at the middle of a long list what do you do?

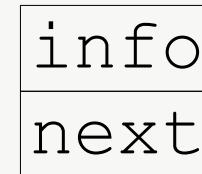
How would you do it with arrays?

Singly linked lists use pointers as links to the next element in the list.

# Nodes in singly linked lists

Each node contains a value and a pointer to the next node:

```
public class Node {  
    public int info;  
    public Node next;  
    ...  
}
```



```
Node h;
```

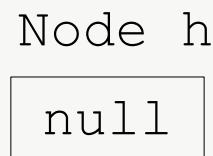


# A node implementation in a singly linked list

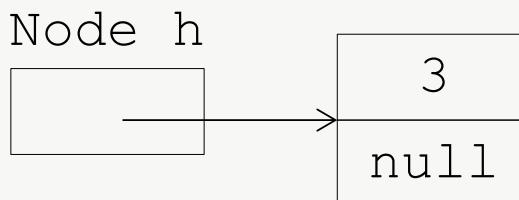
```
public class Node {  
    public int info;  
    public Node next;  
  
    public Node(int i, Node n) {  
        info = i; next = n;  
    }  
  
    public Node(int i) {  
        this(i, null);  
    }  
}
```

# Adding elements at the head of a list

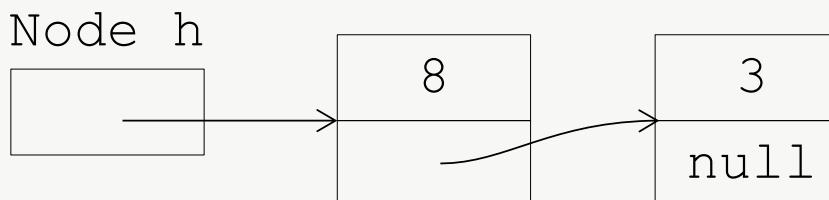
```
Node h = null;
```



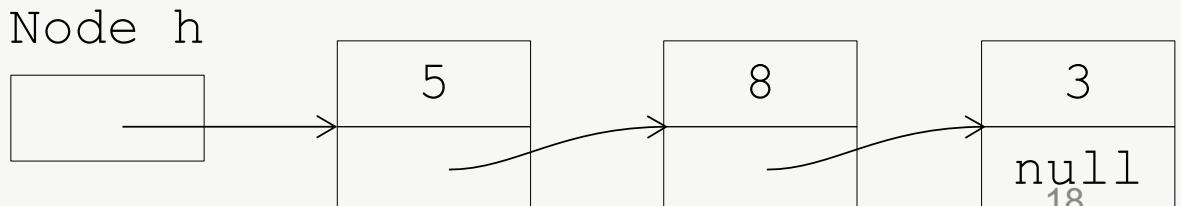
```
h = new Node(3);
```



```
h = new Node(8,h);
```

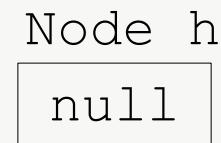


```
h = new Node(5,h);
```

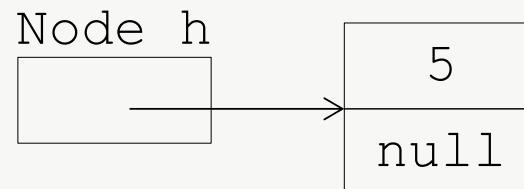


# Adding elements at the tail

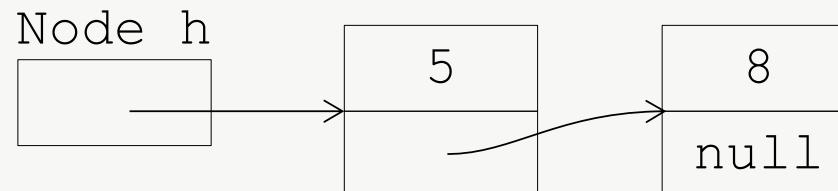
```
Node h = null;
```



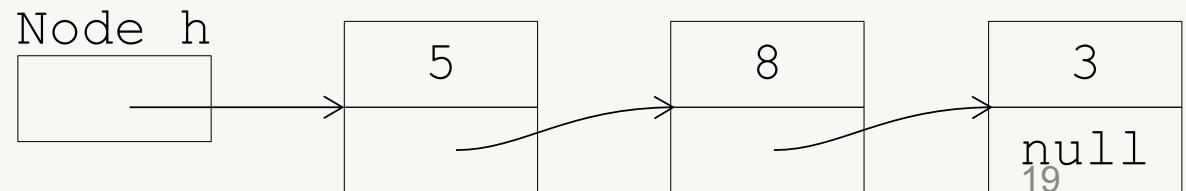
```
h = new Node(5);
```



```
h.next = new Node(8);
```



```
h.next.next = new Node(3);
```



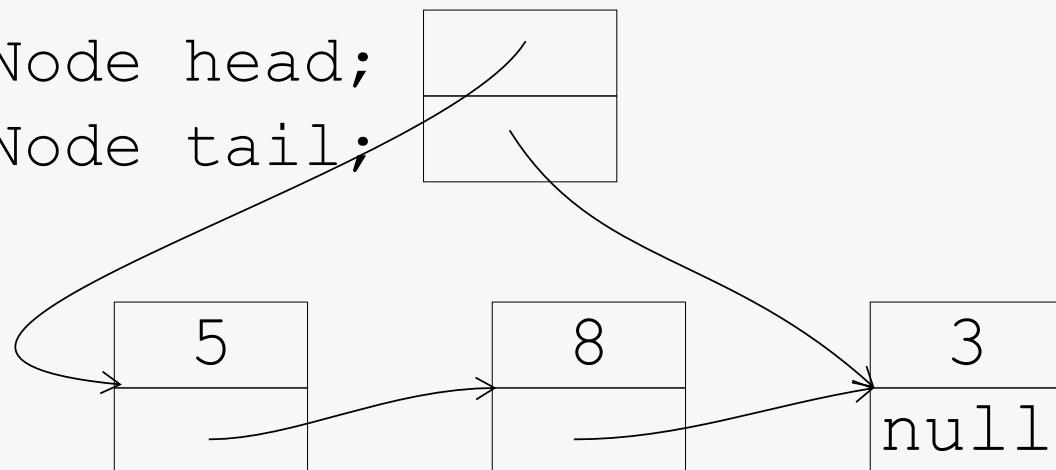
# Pointers to the head and tail

Adding elements at the tail is faster<sup>1</sup> if we keep a pointer to the last element (if any). We define a new class to hold the head and tail pointers:

SLList

Node head;

Node tail;



1. Faster than adding elements at the tail while having only a pointer to the head.

# The singly linked list class

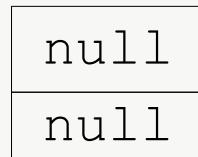
```
public class SLList {  
    private Node head = null;  
    private Node tail = null;  
  
    public boolean isEmpty() {  
        return head == null;  
    }  
  
    public void addToHead(int el)      {...}  
    public void addToTail(int el) { ... }  
    public int deleteFromHead() { ... }  
}
```

# Special cases

An empty list:

SLList

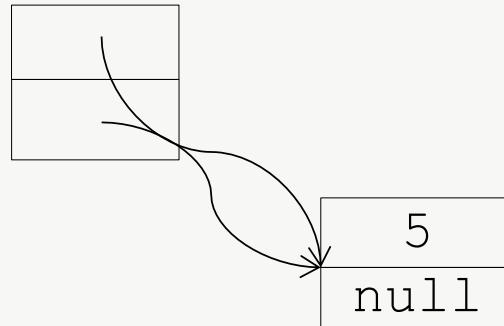
```
Node head;  
Node tail;
```



A list with one element:

SLList

```
Node head;  
Node tail;
```



# What shouldn't happen in singly linked lists

- Loops
- Errors with pointers, e.g.
  - tail does not point to head.next. ... .next
- Tail node with non-null pointer
- Non-tail node with null pointer

# Operations in singly linked lists

- Adding an element at the head
- Adding an element at the tail
- Deleting the head element
- Deleting the tail element

# Adding an element at the head

There are two cases:

- The list contains at least one element  
(so the tail stays the same)
- The list is empty (`head = tail = null`)  
(we need to set the tail)

```
public void addToHead(int el) {  
    head = new Node(el, head);  
    if(tail == null)  
        tail = head;  
}
```

# Adding an element at the tail

Two cases again:

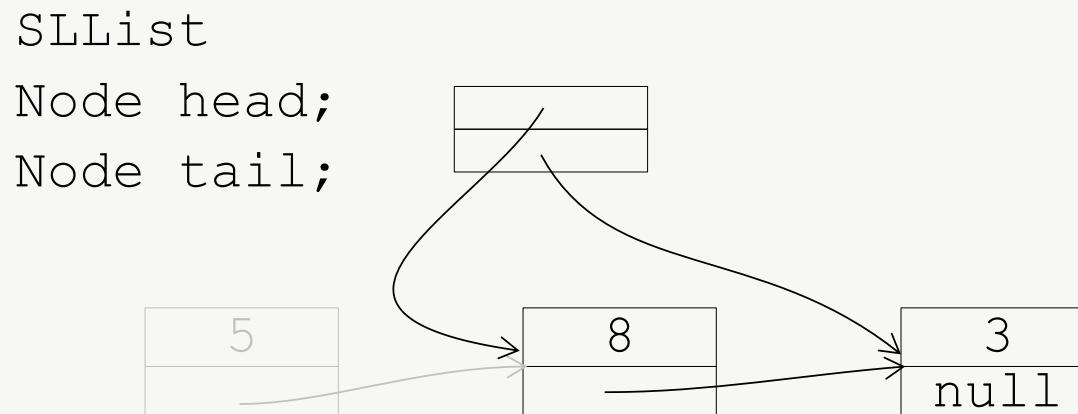
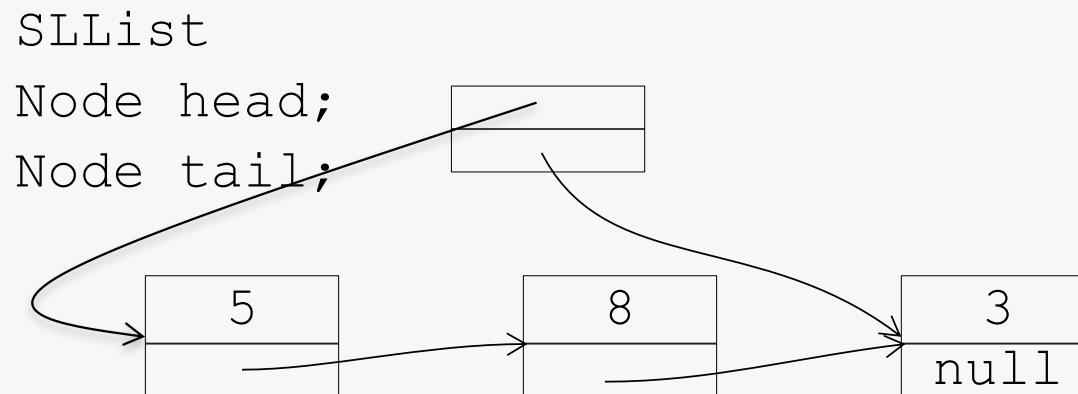
- If the list is empty, we need to set the head
- Otherwise, we need to set the pointer of the old tail

```
public void addToTail(int el) {  
    if (! isEmpty()) {  
        tail.next = new Node(el);  
        tail = tail.next;  
    }  
    else  
        head = tail = new Node(el);  
}
```

# Deleting the head element

Just move the head pointer:

(Special case:  
`head == tail`  
->  
`head = null;`  
`tail = null;`)



## Deleting the head element (2)

If the list has more than one element, change the head  
Otherwise, make the list empty

```
public int deleteFromHead() {  
    int el = head.info;  
    if(head == tail) // one or no elements  
        head = tail = null;  
    else  
        head = head.next;  
    return el;  
}
```

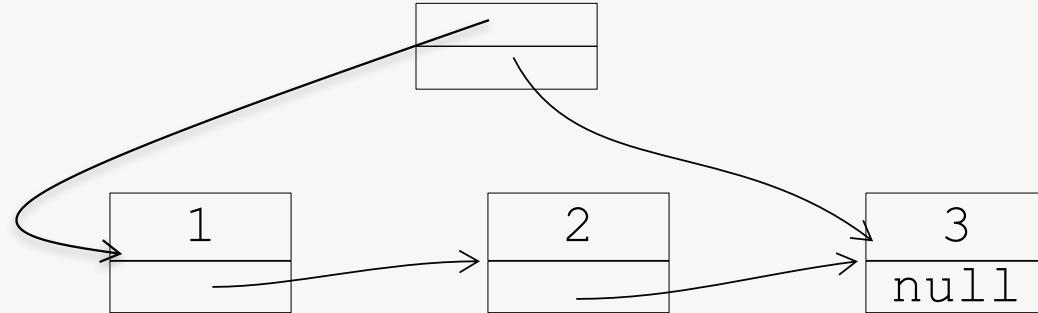
# Deleting the last element

We need to set tail to the second last node  
... but it's more easily said than done:

```
if(head == tail) // one or no elements
    head = tail = null;
else
    aux = sll.head
    while (aux.next!=tail)
        aux = aux.next
    aux.next = null
    sll.tail = aux
```

# Traversing a list

```
Function printList()
aux = sll.head
While (aux! = null)
    print aux.info
    aux = aux.next
```



... so to traverse the list, iteratively (or recursively!) evaluate the rest of the list (i.e., next)

# Applications of singly linked lists

- A list requires  $O(n)$  space
- The operations we have defined each take  $O(1)$  time
- Lists can represent the following ADTs:
  - Stack uses
    - for push
    - for pop
  - Queue uses
    - for enqueue
    - for dequeue
- Lists may also be used directly in many algorithms
- Important high-level programming languages strongly rely on linked lists (e.g., Lisp)

## Other uses?

It is possible to implement priority queues, stacks, queues and heaps using linked lists. However, is this efficient?

Let us have a closer look:

- priority queues: inserting a new element in a sorted singly linked list takes  $O(?)$
- heaps: searching a parent or child in a singly linked list is  $O(?)$

# Some observations

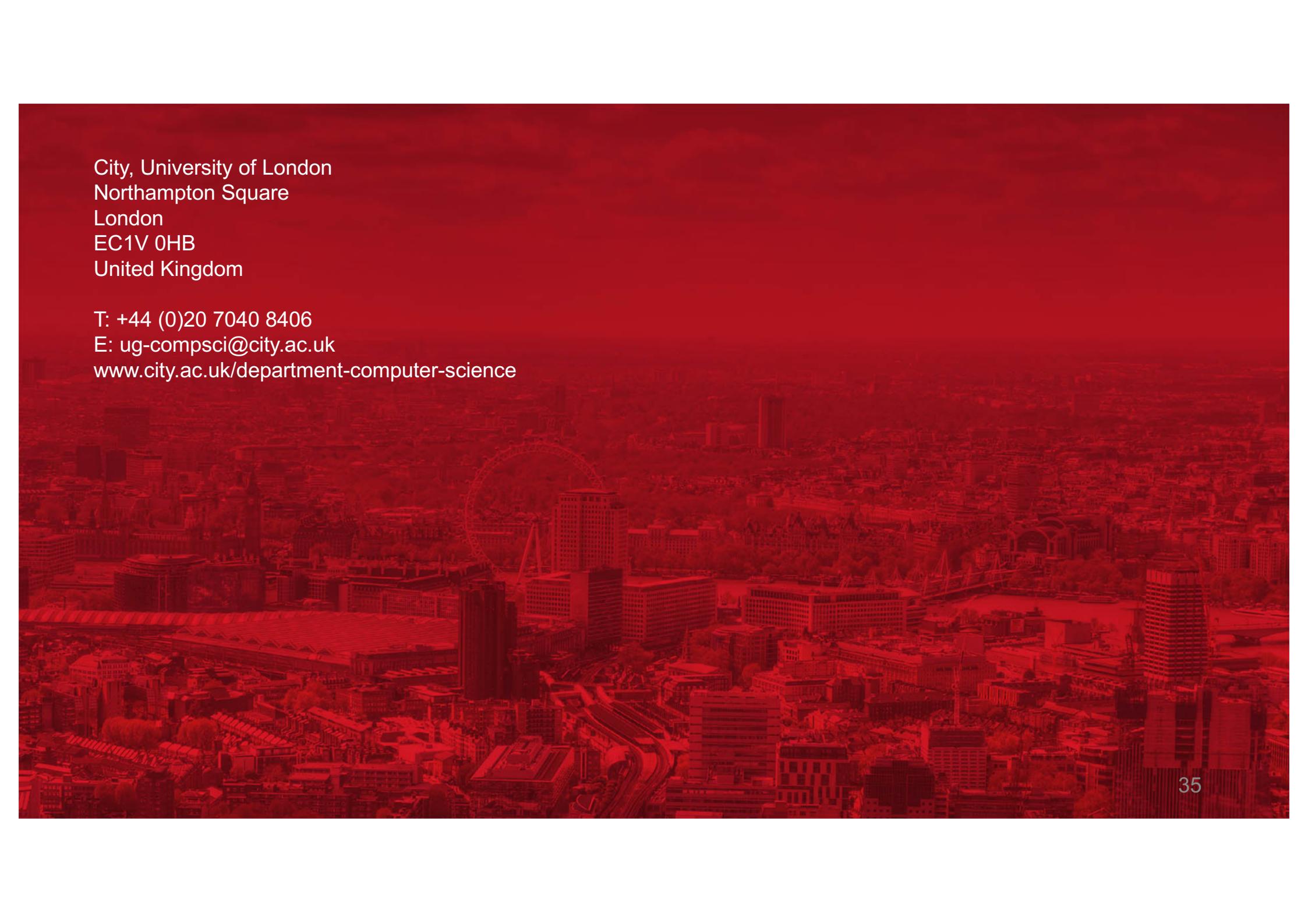
- Extensible arrays efficiently implement the Stack ADT, as well as the array operations, overhead for additional space
- Singly linked lists:
  - efficiently implement the Stack and Queue ADTs, and have many other uses
  - can only move in one direction, accessing elements in the middle or rear part is difficult
  - incur overhead for the pointers

# Reading

- Weiss: Chapter 15 (Stacks and Queues), and section 16.1 and 16.2 (Linked Lists)
- Drozdek: Sections 1.4 (Java and Pointers), 3.1 (Singly Linked Lists), 4.1 (Stacks) and 4.2 (Queues)

Next week: Circular lists and Doubly linked lists

Drozdek: sections 3.2 (doubly linked lists) and 3.3 (circular lists) OR Weiss: section 16.3 (doubly linked lists and circular lists)

The background of the slide is a high-angle aerial photograph of the London skyline. Key landmarks visible include the London Eye, the River Thames, and various modern skyscrapers and historical buildings. The sky is clear and blue.

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