Gormos:

Sharded Plasma for scalable DEX

Alex Xiong



Agenda

- Preliminary
- Problem Statement
- □ Challenges → Gormos: Layer 2 Sharding → Open Problem

→ ACK: Joint work with Dr. Loi Luu

Design Space → **Design Patterns/Lessons**







Preliminary

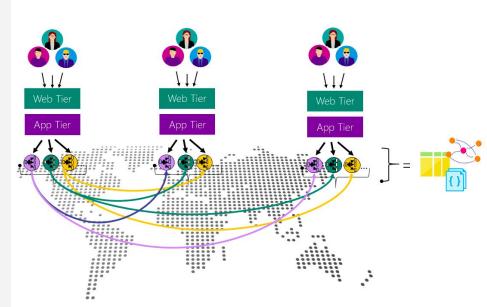


Stateful distributed database, with standardized execution engine, without central control.



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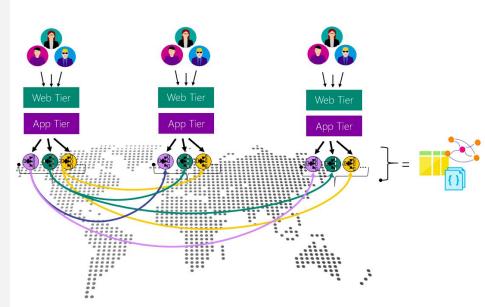
without central control.





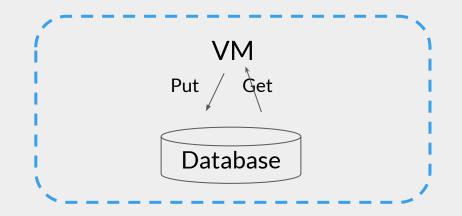
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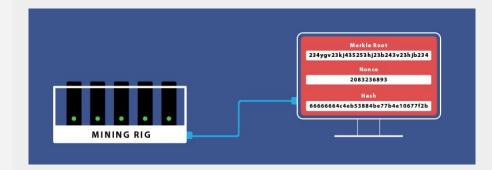


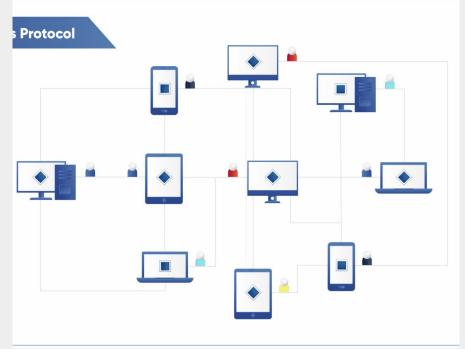
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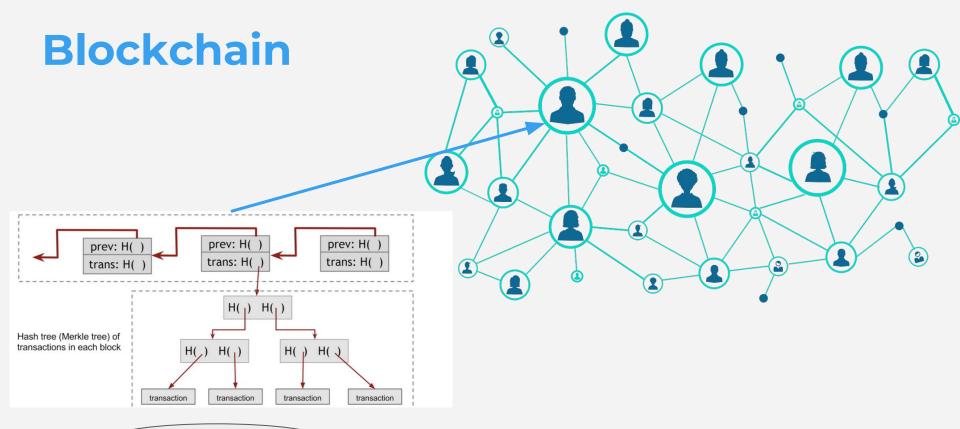




Stateful distributed database, with standardized execution engine, without central control.











Blockchain: Permission(less/ed)

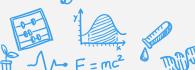
- Permission required to join & participate
- ☐ Potential problems of public blockchain:
 - Sybil attack
 Sybi
 - Communication overhead

 - → Byzantine failure











Public Blockchain: Properties

- Secure & fault tolerant
 - ∪ Under "honest majority" assumption
- ☐ Data availability & (somewhat) censorship resistant
 - Through replication & decentralized control
- Immutability w.h.p & enforced honest computing
 - Through authenticated, append-only data structure and public smart contract









Problems



Fact Check ...





56,000 tps only 1/3 is used



4,000 tps



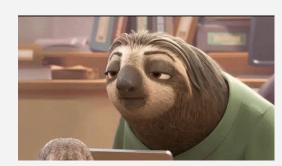
10 tps





7 tps









Blockchain: inherent limitation

- Each node: process all TXs, store all TXs.
 - Network throughput = O (capacity of a single node)
 - □ Does not scale with more computation
- ☐ Throughput (tps) = tx per block / block interval
 - ****









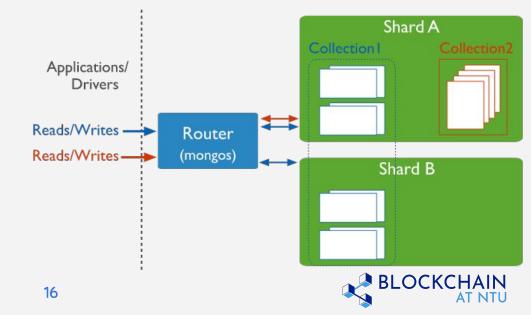
Key Idea:

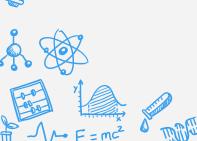
Scale network throughput by partitioning data processing & storage & parallelizing their execution within a smaller committee.



Sharding:: a good old friend

- Huge data volume
- ☐ Frequent data TX
- ☐ Higher throughput & faster response







Challenges



Traditional Sharding

- ☐ Trusted Infrastructure
 - Master-slave, coordinated sharding
 - → Non-byzantine behavior











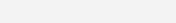


Non-trivial for public blockchain

- Stronger adversarial assumption
 - Single shard takeover?
- ☐ Generic sharding
- Data availability
 - Shard fraud inspection/detection

 - **Expensive cross-shard communication**
 - \rightarrow Asynchronous, non-global \rightarrow race condition, atomicity











Global enforcement on non-global data

-- what's so hard about Blockchain Layer 2 approaches





Gormos

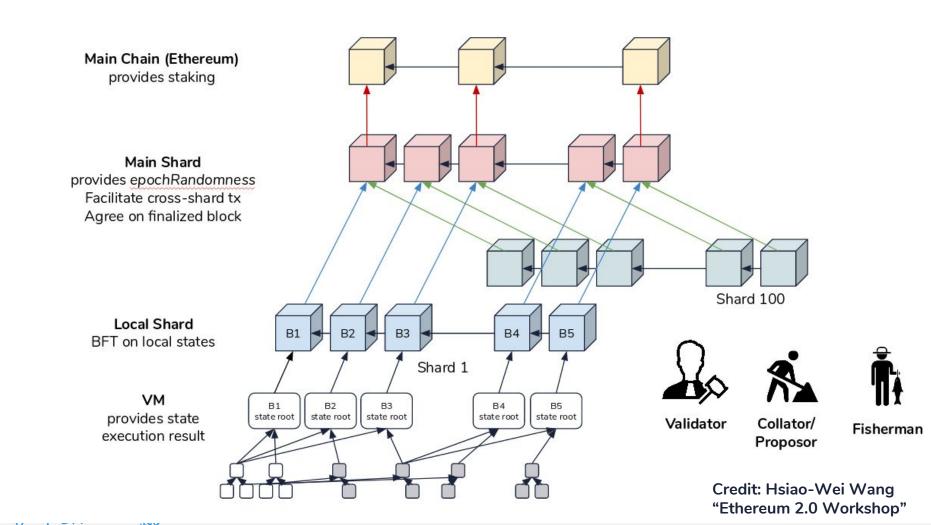


Key Observation & Insight

Sharded Plasma chain (layer 2 / off-chain)

- with cryptononmically bounded validators randomly assigned to each shard to notarize new block.
- 2. **Shard key is token pair** (e.g. {ETH, OMG}) to minimize cross-shard TX.
- 3. **Fisherman** guarantee data availability through probabilistic challenge.





- 1. Validator pool formation
- 2. Random assignment of validators to shards
- 3. Intra-committee BFT consensus
- 4. Final-committee/main shard consensus and root-chain commit
- 5. Generate randomness for next epoch







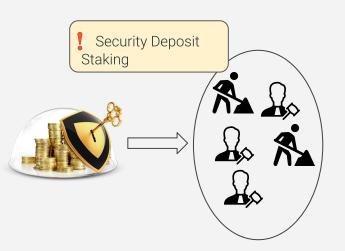
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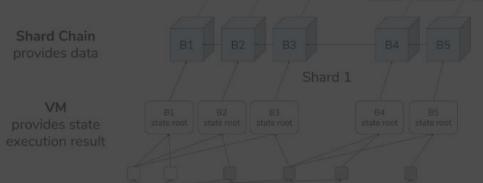




Main Chain provides staking

Staking - Sharding Manager Contract (SMC)

- ♦ Security deposit
- ♦ Validator & Collator pool



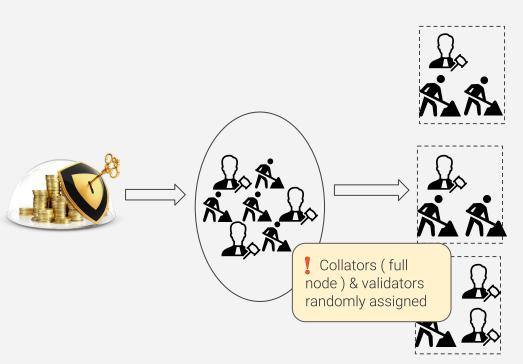


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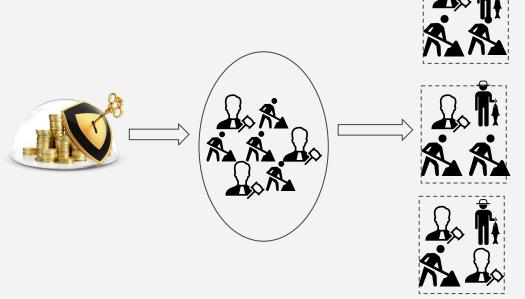








Fisherman (small bond) challenge collator, ensure data availability





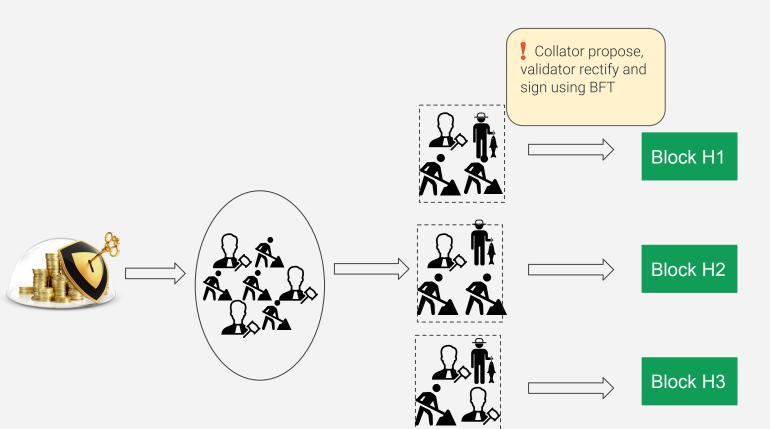
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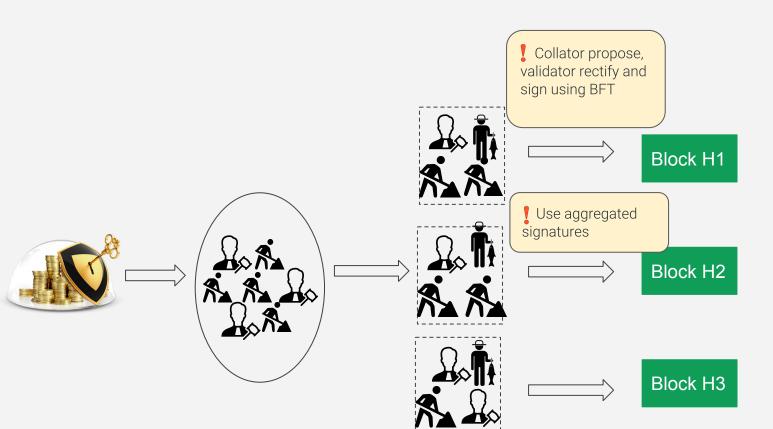








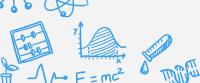




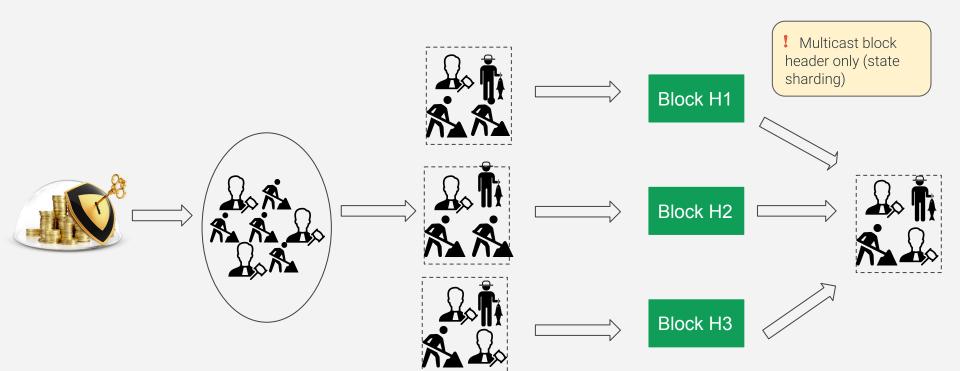


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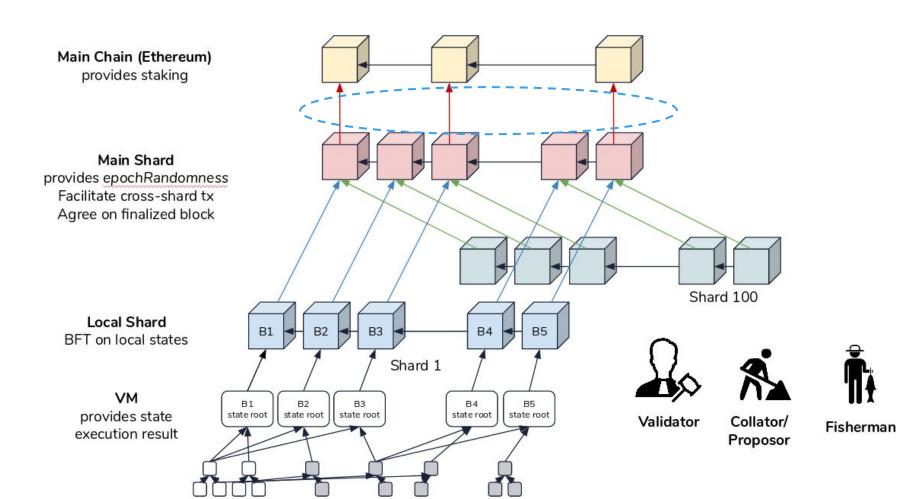












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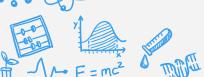
Epoch Randomness

- Unpredictable by adversaries
 - Avoid targeted attack & single shard takeover
- RNG is hard
 - ∪p to % faulty nodes
 - Excessive message complexity

Robust Random Number Generation for Peer-to-Peer Systems

Baruch Awerbuch^{1,⋆} and Christian Scheideler²



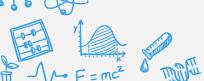




Epoch Randomness

- ☐ VRF [Micali, S. et al.]
- ☐ Or RANDAO









Open Problems



Open Problems

- Compatible incentive design for fisherman and collators
- Formalization of error bound on data availability guarantee
 - Probabilistic probing/challenging
- ☐ Fast sync for newly assigned validators
 - Recursive SNARK? Cryptonomic checkpointing?
- ☐ Efficient, secure cross-shard communication specs
- ✓□ Front-running attack







Thank you!







