

# Security Audit Report for Coordinape Protocol

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### **Report Manifest**

Item	Description
Client	Coordinape
Target	Coordinape Protocol

### **Version History**

Version	Date	Description
1.1	April 11, 2022	Update the status for issue 5 described in Section 2.2.3
1.0	March 20, 2022	First Release

About BlockSec Team focuses on the security of the blockchain ecosystem, and collaborates with leading DeFi projects to secure their products. The team is founded by top-notch security researchers and experienced experts from both academia and industry. They have published multiple blockchain security papers in prestigious conferences, reported several zero-day attacks of DeFi applications, and released detailed analysis reports of high-impact security incidents. They can be reached at Email, Twitter and Medium.

# **Chapter 1 Introduction**

### 1.1 About Target Contracts

Information	Description
Туре	Smart Contract
Language	Solidity
Approach	Semi-automatic and manual verification

The repository that has been audited includes coordinape-protocol <sup>1</sup>. The auditing process is iterative. Specifically, we first audit the code in Version 1 and then audit the commits that fix the issues raised in Version 1. If there exist new issues, we will continue this process. The commit SHA values during the audit are shown in the following.

Project		Commit SHA	
Coordinape	Version 1	7a8e6173305696c72195fa4242126d284611270c	
Coordinape	Version 2	de2c00ad3421d7c72f2425c83a9eb950d03d81da	

#### 1.2 Disclaimer

This audit report does not constitute investment advice or a personal recommendation. It does not consider, and should not be interpreted as considering or having any bearing on, the potential economics of a token, token sale or any other product, service or other asset. Any entity should not rely on this report in any way, including for the purpose of making any decisions to buy or sell any token, product, service or other asset.

This audit report is not an endorsement of any particular project or team, and the report does not guarantee the security of any particular project. This audit does not give any warranties on discovering all security issues of the smart contracts, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit cannot be considered comprehensive, we always recommend proceeding with independent audits and a public bug bounty program to ensure the security of smart contracts.

The scope of this audit is limited to the code mentioned in Section 1.1. Unless explicitly specified, the security of the language itself (e.g., the solidity language), the underlying compiling toolchain and the computing infrastructure are out of the scope.

# 1.3 Procedure of Auditing

We perform the audit according to the following procedure.

• **Vulnerability Detection** We first scan smart contracts with automatic code analyzers, and then manually verify (reject or confirm) the issues reported by them.

1

<sup>&</sup>lt;sup>1</sup>https://github.com/coordinape/coordinape-protocol/



- Semantic Analysis We study the business logic of smart contracts and conduct further investigation on the possible vulnerabilities using an automatic fuzzing tool (developed by our research team).
   We also manually analyze possible attack scenarios with independent auditors to cross-check the result.
- Recommendation We provide some useful advice to developers from the perspective of good programming practice, including gas optimization, code style, and etc.

We show the main concrete checkpoints in the following.

#### 1.3.1 Software Security

- Reentrancy
- DoS
- Access control
- Data handling and data flow
- Exception handling
- Untrusted external call and control flow
- Initialization consistency
- Events operation
- Error-prone randomness
- Improper use of the proxy system

#### 1.3.2 DeFi Security

- Semantic consistency
- Functionality consistency
- Access control
- Business logic
- Token operation
- Emergency mechanism
- Oracle security
- Whitelist and blacklist
- Economic impact
- Batch transfer

#### 1.3.3 NFT Security

- Duplicated item
- Verification of the token receiver
- Off-chain metadata security

#### 1.3.4 Additional Recommendation

- Gas optimization
- Code quality and style

2



**Note** The previous checkpoints are the main ones. We may use more checkpoints during the auditing process according to the functionality of the project.

### 1.4 Security Model

To evaluate the risk, we follow the standards or suggestions that are widely adopted by both industry and academy, including OWASP Risk Rating Methodology <sup>2</sup> and Common Weakness Enumeration <sup>3</sup>. The overall *severity* of the risk is determined by *likelihood* and *impact*. Specifically, likelihood is used to estimate how likely a particular vulnerability can be uncovered and exploited by an attacker, while impact is used to measure the consequences of a successful exploit.

In this report, both likelihood and impact are categorized into two ratings, i.e., *high* and *low* respectively, and their combinations are shown in Table 1.1.

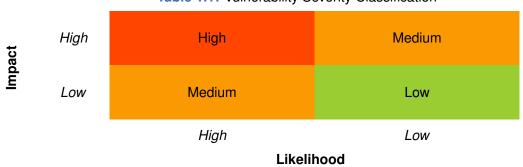


Table 1.1: Vulnerability Severity Classification

Accordingly, the severity measured in this report are classified into three categories: **High**, **Medium**, **Low**. For the sake of completeness, **Undetermined** is also used to cover circumstances when the risk cannot be well determined.

Furthermore, the status of a discovered issue will fall into one of the following four categories:

- Undetermined No response yet.
- Acknowledged The issue has been received by the client, but not confirmed yet.
- **Confirmed** The issue has been recognized by the client, but not fixed yet.
- **Fixed** The issue has been confirmed and fixed by the client.

<sup>&</sup>lt;sup>2</sup>https://owasp.org/www-community/OWASP\_Risk\_Rating\_Methodology

<sup>3</sup>https://cwe.mitre.org/

# **Chapter 2 Findings**

In total, we find **five** potential issues in the smart contract. We also have **five** recommendations, as follows:

High Risk: 2Medium Risk: 2Low Risk: 1

Recommendations: 5

ID	Severity	Description	Category	Status
1	Medium	Incorrect delay time in schedule	Software Security	Fixed
2	Medium	Potential unsafe check in TimeLock	Software Security	Fixed
3	Low	Insufficient check for _token and _simpleToken	DeFi Security	Fixed
4	High	Lack of authorization of the apeWithdrawSim- pleToken function	DeFi Security	Fixed
5	High	Insufficient sanitation in uploadEpochRoot	DeFi Security	Fixed
6	-	Do not use unlimited approval	Recommendation	Fixed
7	-	Use safeTransfer	Recommendation	Fixed
8	-	Ensure that the same storage layout of the proxy	Recommendation	Acknowledged
9	-	Ensure that state variables in ApeRegistry must be properly initialized	Recommendation	Fixed
10	-	Remove unused variables in ApeVault.sol	Recommendation	Fixed

The details are provided in the following sections.

# 2.1 Software Security

#### 2.1.1 Incorrect delay time in schedule

**Status** Fixed in Version 2 **Introduced by** Version 1

**Description** The schedule function in the TimeLock has an implementation error that the minDelay on line 58 should be \_delay. This can make the TimeLock mechanism work unexpectedly, since the task will be scheduled immediately, instead of with the expected delay (\_delay).

Listing 2.1: TimeLock.sol

**Impact** Make the TimeLock mechanism work unexpectedly.



**Suggestion** Use \_delay instead of minDelay on line 58.

#### 2.1.2 Potential unsafe check in TimeLock

```
Status Fixed in Version 2
Introduced by Version 1
```

**Description** If timestamps[id] is accidentally to be set as 0, then the require check of the execute() function (line 70 - line 71) could be bypassed. Actually this sanity check is inconsistent with the logic isOperationReady in OpenZeppelin.

```
46 function isDoneCall(bytes32 _id) public view returns(bool) {
47    return timestamps[_id] == _DONE_TIMESTAMP;
48  }
49
50 function isReadyCall(bytes32 _id) public view returns(bool) {
51    return timestamps[_id] <= block.timestamp;
52 }</pre>
```

Listing 2.2: TimeLock.sol

```
68 function execute(address _target, bytes calldata _data, bytes32 _predecessor, bytes32 _salt,
        uint256 _delay) external onlyOwner {
69
     bytes32 id = hashOperation(_target, _data, _predecessor, _salt);
70
    require(isReadyCall(id), "TimeLock: Not ready for execution");
71
     require(!isDoneCall(id), "TimeLock: Already executed");
72
     require(_predecessor == bytes32(0) || isDoneCall(_predecessor), "TimeLock: Predecessor call
          not executed");
73
     _call(id, _target, _data);
74
     timestamps[id] = _DONE_TIMESTAMP;
75 }
```

Listing 2.3: TimeLock.sol

```
function isOperationReady(bytes32 id) public view virtual returns (bool ready) {
   uint256 timestamp = getTimestamp(id);
   return timestamp > _DONE_TIMESTAMP && timestamp <= block.timestamp;
}</pre>
```

Listing 2.4: TimelockController.sol in OpenZeppelin

**Impact** The sanity check in execute could be bypassed.

**Suggestion** Make it consistent with OpenZeppelin.

### 2.2 DeFi Security

#### 2.2.1 Insufficient check for \_token and \_simpleToken

```
Status Fixed in Version 2
Introduced by Version 1
Description
```



According to the document, \_token and \_simpleToken cannot be zero at the same time. However, there does not exist such a verification in the init() function.

```
72 function init(
73 address _apeRegistry,
74 address _token,
75 address _registry,
76 address _simpleToken,
77 address _newOwner) external {
78 require(!setup);
79
     setup = true;
80 apeRegistry = _apeRegistry;
81 if (_token != address(0))
82
     vault = VaultAPI(RegistryAPI(_registry).latestVault(_token));
83
     simpleToken = IERC20(_simpleToken);
84
85
     // Recommended to use a token with a 'Registry.latestVault(_token) != address(0)'
86
    token = IERC20(_token);
87 // Recommended to use 'v2.registry.ychad.eth'
88
    registry = RegistryAPI(_registry);
89
     _owner = _newOwner;
90
     emit OwnershipTransferred(address(0), _newOwner);
91 }
```

Listing 2.5: wrapper/beacon/ApeVault.sol

#### Impact NA

**Suggestion** Add a check in the function.

#### 2.2.2 Lack of authorization of the apeWithdrawSimpleToken function

```
Status Fixed in Version 2 Introduced by Version 1
```

#### **Description**

The public function <code>apeWithdrawSimpleToken</code> lacks a modifier so that everyone can withdraw the simple token from the contract.

```
131 function apeWithdrawSimpleToken(uint256 _amount) public {
132    simpleToken.safeTransfer(msg.sender, _amount);
133 }
```

Listing 2.6: wrapper/beacon/ApeVault.sol

**Impact** Anyone can withdraw the simple token in the contract.

**Suggestion** Add the onlyOwner modifier.

#### 2.2.3 Insufficient sanitation in uploadEpochRoot

```
Status Fixed in Version 2
Introduced by Version 1
Description
```



The uploadEpochRoot is used to uploadEpochRoot that can be used to claim the tokens in each epoch. However, we find that this function does not have enough sanitation that could be abused by attackers to claim tokens.

The possible attack process is as following:

- 1. The attacker invokes uploadEpochRoot, with a \_vault contract that is controlled (deployed) by the attacker. The passed \_token is a valid token in the distributor contract. The passed \_amount is zero. There is no check on the passing \_vault parameter. Since \_vault is controlled by the attacker, the attacker can control the return value of the function call to this contract. Specifically, the attacker returns msg.sender to the invocation \_vault.owner(). Thus, isOwner is true on line 60. So the check on line 61 can pass. Similarly, the require statement on line 62 can be bypassed since the return value of \_vault.vault() is controlled by the attacker. Then from line 65-69, the code will update the epochRoots to the \_root passed by the attacker. The check from line 73 does not help here since the passed \_amount is zero.
- 2. Then the attacker can invoke claim function. Since the \_root has been updated in the previous step, the check on line 118 can pass and the attacker can drain the tokens in the circle.

```
52 function uploadEpochRoot(
53 address _vault,
54 bytes32 _circle,
55 address _token,
56 bytes32 _root,
57 uint256 _amount,
58 uint8 _tapType)
59 external {
60
     bool isOwner = ApeVaultWrapperImplementation(_vault).owner() == msg.sender;
61
     require(vaultApprovals[_vault][_circle] == msg.sender || isOwner, "Sender cannot upload a root
62
     require(address(ApeVaultWrapperImplementation(_vault).vault()) == _token, "Vault cannot supply
           token");
63
      if (!isOwner)
64
      _isTapAllowed(_vault, _circle, _token, _amount);
65
      uint256 epoch = epochTracking[_circle][_token];
66
      epochRoots[_circle][_token][epoch] = _root;
67
68
      epochTracking[_circle] [_token]++;
69
      circleAlloc[_circle][_token] += _amount;
70
      uint256 beforeBal = IERC20(_token).balanceOf(address(this));
71
      uint256 sharesRemoved = ApeVaultWrapperImplementation(_vault).tap(_amount, _tapType);
72
      uint256 afterBal = IERC20(_token).balanceOf(address(this));
73
      require(afterBal - beforeBal == _amount, "Did not receive correct amount of tokens");
74
      if (sharesRemoved > 0)
75
      emit apeVaultFundsTapped(_vault, address(ApeVaultWrapperImplementation(_vault).vault()),
          sharesRemoved);
76 }
```

Listing 2.7: ApeDistributor.sol



```
117
       bytes32 node = keccak256(abi.encodePacked(_index, _account, _checkpoint));
118
       require(_proof.verify(epochRoots[_circle][_token][_epoch], node), "Wrong proof");
119
       uint256 currentCheckpoint = checkpoints[_circle][_token][_account];
120
       require(_checkpoint > currentCheckpoint, "Given checkpoint not higher than current checkpoint"
           );
121
122
       uint256 claimable = _checkpoint - currentCheckpoint;
123
       require(claimable <= circleAlloc[_circle][_token], "Can't claim more than circle has to give")</pre>
           ;
124
       circleAlloc[_circle][_token] -= claimable;
125
       checkpoints[_circle][_token][_account] = _checkpoint;
126
       _setClaimed(_circle, _token, _epoch, _index);
127
       if (_redeemShares && msg.sender == _account)
128
       VaultAPI(_token).withdraw(claimable, _account);
129
130
       IERC20(_token).safeTransfer(_account, claimable);
131
       emit Claimed(_circle, _token, _epoch, _index, _account, claimable);
132 }
```

Listing 2.8: ApeDistributor.sol

**Impact** The attacker can update the EpochRoot as the vault owner.

**Suggestion** Ensure that the passing \_vault is a valid one (not a fake one deployed by the attacker.)

**Note** The project fixes this issue by using \_vault as the first index of epochRoots. In this case, even the attacker can pass a fake \_vaule, it will not cause the security impact since it only updates the epochRoot for that fake value.

#### 2.3 Additional Recommendation

#### 2.3.1 Do not use unlimited approval

**Status** Fixed in Version 2 Introduced by Version 1

#### Description

The delegateDeposit use unlimited approval, which is controversial and widely debated in recent months due to some relevant security incidents.



```
52
     IERC20(_token).safeApprove(address(vault), 0); // Avoid issues with some IERC20(_token)s
          requiring 0
53
     IERC20(_token).safeApprove(address(vault), MAX_UINT); // Vaults are trusted
54 }
55
56 uint256 beforeBal = IERC20(_token).balanceOf(address(this));
57
58 uint256 sharesMinted = vault.deposit(_amount, _apeVault);
59
60 uint256 afterBal = IERC20(_token).balanceOf(address(this));
61 deposited = beforeBal - afterBal;
62
63
64 ApeVaultWrapperImplementation(_apeVault).addFunds(deposited);
65
   emit DepositInVault(_apeVault, _token, sharesMinted);
66}
```

Listing 2.9: ApeRouter.sol

#### Impact NA.

**Suggestion** Use the actual amount for approval instead of the unlimited one.

#### 2.3.2 Use safeTransfer

Status Fixed in Version 2
Introduced by Version 1

#### **Description**

The \_tapBase use the transfer function while \_tapOnlyProfit uses safeTransfer. It's a good practice to make the invocation consistent and use safeTransfer instead.

```
200
201 function _tapOnlyProfit(uint256 _tapValue, address _recipient) internal {
202
    uint256 fee = FeeRegistry(ApeRegistry(apeRegistry).feeRegistry()).getVariableFee(_tapValue,
203
      uint256 finalTapValue = _tapValue + _tapValue * fee / TOTAL_SHARES;
204
      require(_shareValue(finalTapValue) <= profit(), "Not enough profit to cover epoch");</pre>
205
      vault.safeTransfer(_recipient, _tapValue);
206
      vault.safeTransfer(ApeRegistry(apeRegistry).treasury(), _tapValue * fee / TOTAL_SHARES);
207 }
208
209 /**
210 * @notice
211 * Used to take funds from vault by deducting a part from profits
212 * @param _tapValue Amount of funds to take
213 * Oparam _recipient recipient of funds (always distributor)
214 */
215 function _tapBase(uint256 _tapValue, address _recipient) internal {
    uint256 underlyingTapValue = _shareValue(_tapValue);
216
217
    uint256 profit_ = profit();
218
     uint256 fee = FeeRegistry(ApeRegistry(apeRegistry).feeRegistry()).getVariableFee(profit_,
          underlyingTapValue);
219
      uint256 finalTapValue = underlyingTapValue + underlyingTapValue * fee / TOTAL_SHARES;
```



```
if (finalTapValue > profit_)
underlyingValue -= finalTapValue - profit_;
vault.transfer(_recipient, _tapValue);
vault.transfer(ApeRegistry(apeRegistry).treasury(), _tapValue * fee / TOTAL_SHARES);
}
```

Listing 2.10: wrapper/beacon/ApeVault.sol

```
126
      receive() external payable {
127
          emit Deposit(msg.sender, msg.value, address(this).balance);
128
      }
129
130
      fallback() external payable {
131
          if (msg.value > 0) {
132
              emit Deposit(msg.sender, msg.value, address(this).balance);
133
          }
134
      }
```

Listing 2.11: MultiSigs.sol

#### Impact NA

**Suggestion** Use safeTransfer.

#### 2.3.3 Ensure that the same storage layout of the proxy

Status Acknowledged.

#### **Description**

In our local simulation, proxy behaves unexpectedly when the new implementation of ApeVault has a different storage layout. When updating the implementation of the ApeVault, ensure that the storage layout is same.

**Impact** Unexpected behavior.

**Suggestion** Keep the same storage layout.

**Feedback from the developer** When we will be pushing out new implementations, it will be mandatory to keep storage layout the way it is and add new storage at the end of the current layout.

#### 2.3.4 Ensure that state variables in ApeRegistry must be properly initialized

```
Status Fixed in Version 2 Introduced by Version 1
```

#### Description

The state variables (e.g., treasury) are not initialized in the constructor(), thus the corresponding setters (e.g., setTreasury()) must be properly invoked. Otherwise, it will lead to a fund loss since part of the fund will be transferred to the treasury when performing the tapping operation.

```
13 constructor(uint256 _minDelay) TimeLock(_minDelay) {}
```

Listing 2.12: ApeRegistry.sol

#### Impact NA.

**Suggestion** Add state variables to the constructor.



## 2.3.5 Remove unused variables in ApeVault.sol

Status Fixed in Version 2
Introduced by Version 1

**Description** 

There are two unused variables in ApeVault.sol, including hasAccess and allowanceModule.

Impact NA

**Suggestion** Remove unused variables.