

CS 2800 HW #10

KIRILL CHERNYSHOV

Key: as. = assumption, abs. = reductio ad absurdum, e. m. = excluded middle.

Problem 1

(a) Suppose $\neg P$. Q can either evaluate to T or F :

$\neg P$	P	Q	$P \wedge Q$	$\neg(P \wedge Q)$
T	F	T	F	T
T	F	F	F	T

Therefore, $\neg P \models \neg(P \wedge Q)$.

(b)

$$\frac{\overline{\vdash (P \wedge Q) \vee \neg(P \wedge Q)}}{e. m.} \frac{\frac{\overline{P \wedge Q \vdash P \wedge Q}}{as.} \wedge elim \quad \frac{\overline{\neg P \vdash \neg P}}{as.}}{\neg P, P \wedge Q \vdash \neg(P \wedge Q)} abs. \quad \frac{\overline{\neg(P \wedge Q) \vdash \neg(P \wedge Q)}}{as.} \vee elim$$

Problem 2

Inductively assume $P(\phi)$ and $P(\psi)$. $I \models \phi \rightarrow \psi$ means that $\text{eval}(\phi \rightarrow \psi, I) = T$, that is, $I \models \neg\phi$ or $I \models \psi$. If $I \models \neg\phi$ then using our assumption, $I \vdash \neg\phi$, and so $I \vdash \phi \rightarrow \psi$. This can be seen from the following proof tree:

$$\frac{\frac{\overline{\phi \vdash \phi}}{as.} \quad \frac{\overline{\neg\phi \vdash \neg\phi}}{phi} as.}{\neg\phi, \phi \vdash \psi} abs. \quad \frac{\neg\phi \vdash \phi \rightarrow \psi}{\rightarrow intro}$$

If $I \models \psi$, then using our assumption, $I \vdash \psi$, and once again we have $I \vdash \phi \rightarrow \psi$:

$$\frac{\overline{\phi, \psi \vdash \psi}}{as.} \rightarrow intro$$

Conversely, if $I \not\models \phi \rightarrow \psi$, then $I \models \phi \wedge \neg\psi$. Using the inductive assumption, we have $I \vdash \phi$ and $I \vdash \neg\psi$. Therefore:

$$\frac{\overline{\vdash (\phi \rightarrow \psi) \vee \neg(\phi \rightarrow \psi)}}{e. m.} \frac{\frac{\overline{\phi \vdash \phi}}{as.} \quad \frac{\overline{\phi \rightarrow \psi \vdash \phi \rightarrow \psi}}{as.} \rightarrow elim \quad \frac{\overline{\neg\psi \vdash \neg\psi}}{as.}}{\phi, \phi \rightarrow \psi \vdash \psi} abs. \quad \frac{\overline{\neg(\phi \rightarrow \psi) \vdash \neg(\phi \rightarrow \psi)}}{as.} \vee elim$$