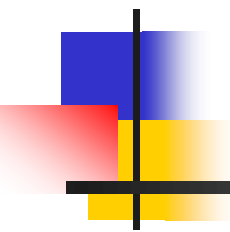


Chapter 10 Advanced topics in relational databases

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- Security and user authorization in SQL
 - **Recursion in SQL**
 - Object-relational model
 - 1. User-defined types in SQL
 - 2. Operations on object-relational data
 - Online analytic processing & data cubes



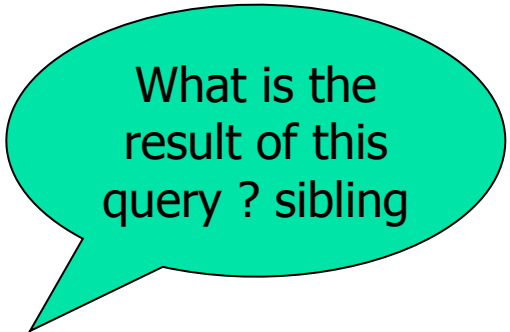
Question?

- EDB: $\text{Par}(c,p) = p$ is a parent of c .
- We want to find **generalized cousins**: people with common ancestors one or more generations back.

Select p1.c,p2.c

From Par P1,Par P2

Where $P1.p=P2.p$ and $P1.c \neq p2.c$



What is the
result of this
query ? sibling



Recursive example

$\text{Sib}(x,y) \leftarrow \text{Par}(x,p) \text{ AND } \text{Par}(y,p) \text{ AND } x \neq y$

$\text{Cousin}(x,y) \leftarrow \text{Sib}(x,y)$

$\text{Cousin}(x,y) \leftarrow \text{Par}(x,xp) \text{ AND } \text{Par}(y,yp)$
 $\text{AND } \text{Cousin}(xp,yp)$



Example: Evaluation of Cousin

- We'll proceed in rounds to infer Sib facts (**red**) and Cousin facts (**green**).
- Remember the rules:

Sib(x,y) <- Par(x,p) AND Par(y,p) AND x<>y

Cousin(x,y) <- Sib(x,y)

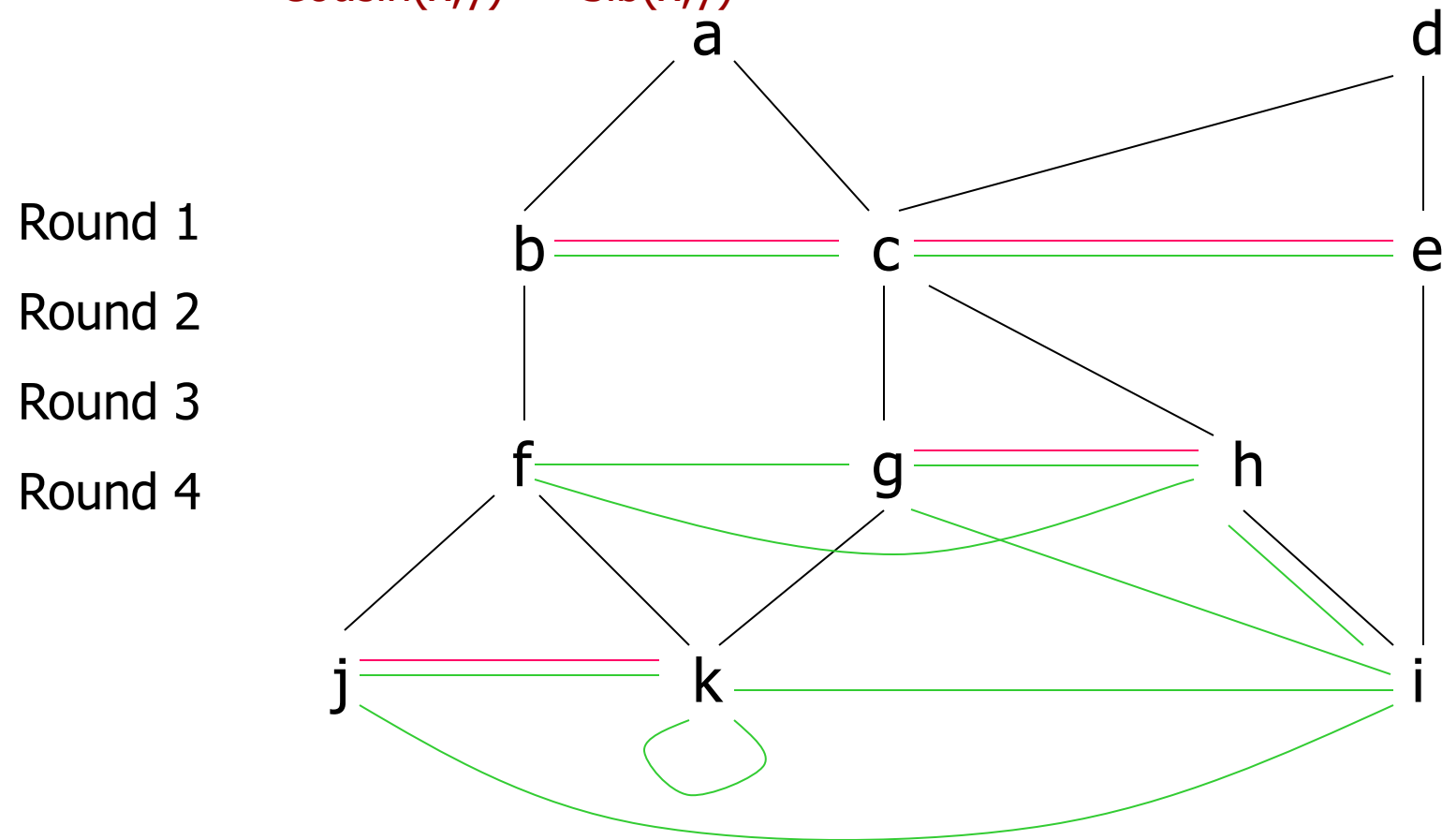
**Cousin(x,y) <- Par(x,xp) AND Par(y,yp)
AND Cousin(xp,yp)**

Par Data: Parent Above Child

$\text{Sib}(x,y) \leftarrow \text{Par}(x,p) \text{ AND } \text{Par}(y,p) \text{ AND } x \neq y$

$\text{Cousin}(x,y) \leftarrow \text{Par}(x,xp) \text{ AND } \text{Par}(y,yp) \text{ AND } \text{Cousin}(xp,yp)$

$\text{Cousin}(x,y) \leftarrow \text{Sib}(x,y)$





SQL-99 Recursion

- Datalog recursion has inspired the addition of recursion to the SQL-99 standard.
- IBM DB2 does implement the SQL-99 proposal.



Form of SQL Recursive Queries

WITH

[RECURSIVE] R1 AS <Definition of R1>

[RECURSIVE] R2 AS <Definition of R2>

<a SQL query about EDB,R1,R2,...>




Example: SQL Recursion – (1)

- Find Sally's cousins, using SQL like the recursive Datalog example.
- **Par(child,parent)** is the EDB.

WITH Sib(x,y) AS

```
SELECT p1.child, p2.child
FROM Par p1, Par p2
WHERE p1.parent = p2.parent AND
      p1.child <> p2.child;
```

Like Sib(x,y) <-
Par(x,p) AND
Par(y,p) AND
x <> y



Example: SQL Recursion – (2)

WITH

RECURSIVE Cousin(x,y) AS

(SELECT * FROM Sib)

UNION

(SELECT p1.child, p2.child
FROM Par p1, Par p2, Cousin
WHERE p1.parent = Cousin.x AND
p2.parent = Cousin.y);

Required – Cousin
is recursive

Reflects Cousin(x,y) <-
Sib(x,y)

Reflects
Cousin(x,y) <-
Par(x, xp) AND
Par(y, yp) AND
Cousin(xp, yp)



Example: SQL Recursion – (3)

- With those definitions, we can add the query, which is about the virtual view Cousin(x,y):

```
SELECT y
FROM Cousin
WHERE x = 'Sally';
```



Legal SQL Recursion

- It is possible to define SQL recursions that do not have a meaning.
- The SQL standard restricts recursion so there is a meaning.



Another Example

```
create table Employee(ID int, salary int);  
create table Manager(mID int, eID int);  
create table Project(name text, mgrID int);
```

- Find total salary cost of project 'X'

Solution 1:

Employee(ID , salary)
Manager(mID, eID)
Project(name,mgrID)

with recursive

Superior as (select * from Manager

union

select S.mID, M.eID

from **Superior S**, Manager M

where S.eID = M.mID)

select sum(salary)

from Employee

where ID in

(select mgrID from Project where name = 'X'

union

select eID from Project, Superior

where Project.name = 'X' AND Project.mgrID = Superior.mID);

Solution 2:

Employee(ID , salary)
Manager(mID, eID)
Project(name,mgrID)



with recursive

Xemps(ID) as (select mgrID as ID from
Project where name = 'X'

union

select eID as ID

from Manager M, **Xemps** X

where M.mID = X.ID)

select sum(salary)

from Employee

where ID in (select ID from Xemps);