CSE3666 HW2 solutions

1. Array.

a)

We only need to add one instruction to the code in the slide. See the red instruction in part b).

There are 8 instructions in the loop. The loop is executed for 100 times. There are two instructions before the loop. So the total number of executed instructions is 802.

```
b)
      for (i = 0; i < 100; i += 2) {
            B[i] = A[i] + 4;
            B[i+1] = A[i+1] + 4;
      }
            li
                  s4, 100
            li
                  s1, 0
                               # we are sure s1 < s4 after initialization
      loop:
            slli
                  t0, s1, 2
                               # t0 = i * 4
                               # compute addr of A[i]
            add
                  t2, t0, s2
                  t3, t0, s3
            add
                               # compute addr of B[i]
            lw
                  t1, 0(t2)
                               # A[i]
            addi
                  t1, t1, 4
                               # the new instruction for +4
                  t1, 0(t3)
            SW
                  t4, 4(t2)
                               # A[i+1]
            lw
                  t4, t4, 4
            addi
                  t4, 4(t3)
            SW
            1w
                  t5, 8(t2)
                               # A[i+2]
                  t5, t5, 4
            addi
            SW
                  t5, 8(t3)
                  t6, 12(t2)
                              # A[i+3]
            lw
            addi
                  t6, t6, 4
                  t6, 12(t3)
            SW
            addi s1, s1, 4
            blt
                  s1, s4, loop # 11 instructions in the loop
```

Each iteration has 17 instructions. The total number of instructions executed is 2 + 17 * 25 = 427.

2. Two-dimensional array.

blt t0, s0, LOOPI

An example implementation is shown below. Pay attention to the nested loops and the calculation of T[i][j]'s address.

Register	s9	t0	t1	s0	s1
Variable/value	address of T	i	j	16	8

```
for (i = 0; i < 16; i += 1)
    for (j = 0; j < 8; j += 1)
         T[i][j] = 256 * i + j;
# We are sure the initial value of i and j are less than 16 and 8, respectively
# So we do not jump to condition test before the loops
# Instructions that depend only on i can be placed outside of the inner loop
# Here, we just convert the loops from C to RISC-V literally.
      addi
           s0, x0, 16
      addi s1, x0, 8
      addi t0, x0, 0
                      # i = 0
      # beq x0, x0 CONDI
                          # this is not needed here
LOOPI:
      addi t1, x0, 0
                              # j = 0
LOOPJ:
                               # t5 = 256 * i
      slli t5, t0, 8
           t6, t5, t1
                               # t6 = 256 * i + j
      add
                               # 8 * i. Number of words before T[i][0]
      slli t2, t0, 3
      addi t3, t2, t1
                               # 8 * i + j. Number of words before T[i][j]
                               # offset of T[i][j]
      slli t3, t3, 2
      add
           t3, t3, s9
                               # address of T[i][j]
           t6, 0(t3)
      SW
      addi t1, t1, 1
                              # j = j + 1
      blt t1, s1, LOOPJ # continue the inner loop if j < 8
      addi t0, t0, 1
                              # i = i + 1
```

continue the outer loop if i < 16</pre>

3. Encoding.

```
Steps: Find bits in each field, and then write the machine code.
013964B3 or s1, s2, s3 (or x9,x18,x19)
0000001001110010110010010110011
R-type
opcode:
          0110011
rd:
          01001
funct3:
          110
rs1:
          10010
rs2:
          10011
funct7:
          0000000
=======
01039313 slli t1, t2, 16 (slli x6,x7,0x10)
0000001000000111001001100010011
I-type shift
shamt:
          10000
imm i:
          000000010000
opcode:
          0010011
rd:
          00110
funct3:
          001
rs1:
          00111
rs2:
          10000
funct7:
          0000000
FFF0C093 xori x1, x1, -1 (xori x1,x1,0xffffffff)
11111111111100001100000010010011
I-type
imm_i:
          111111111111
opcode:
          0010011
rd:
          00001
funct3:
          100
rs1:
          00001
rs2:
          11111
funct7:
          1111111
========
F9C1A103 lw x2, -100(x3) (lw x2,0xffffff9c(x3))
11111001110000011010000100000011
I-type
imm_i:
          111110011100
opcode:
          0000011
rd:
          00010
funct3:
          010
rs1:
          00011
rs2:
          11100
funct7:
          1111100
```

4. Decoding.

Steps: Find the bits in each field. Check opcode, identify the instruction format, find out the immediate if necessary, and then write the instruction.

```
========
FEACA823 sw x10, -16(x25)
```

11111110101011001010100000100011

opcode: 0100011 rd: 10000 funct3: 010 rs1: 11001 rs2: 01010 funct7: 1111111

S-type

111111110000 imm_s:

========

04020713 addi x14,x4,64

00000100000000100000011100010011

opcode: 0010011 rd: 01110 funct3: 000 rs1: 00100 00000 rs2: funct7: 0000010

I-type

imm i: 000001000000

========

00557BB3 and x23,x10,x5

0000000010101010111101110110011

opcode: 0110011 rd: 10111 funct3: 111 rs1: 01010 rs2: 00101 funct7: 0000000

R-type ========

414FDF13 srai x30,x31,20

0100000101001111111011111100010011

opcode: 0010011 rd: 11110 funct3: 101 rs1: 11111 rs2: 10100 funct7: 0100000

I-type shift

imm i: 0**1**0000010100

shamt: 10100