

## Семинар 14

задача Да се док, че  $\varphi_1 \wedge \varphi_2 \wedge \varphi_3 \Rightarrow \varphi_4$   
чрез лна  $\mathcal{F}$ , където:

$$\varphi_1 \equiv \exists z \forall x [\forall y (p(f(x), y) \vee q(z, f(x)))]$$

$$\varphi_2 \equiv \exists z \forall y [q(y, f(f(y))) \Rightarrow \forall x p(x, f(z))]$$

$$\varphi_3 \equiv \forall x \forall y [p(x, f(y)) \Rightarrow (p(f(y), x) \vee \neg \exists y \exists x p(x, y))]$$

$$\varphi_4 \equiv \exists y (\exists x p(f(x), y) \wedge \exists x p(y, x))$$

Л-вс:

"тавтология"  $\Leftrightarrow$  т е "неизпълнено"

$$\neg (\varphi_1 \wedge \varphi_2 \wedge \varphi_3 \Rightarrow \varphi_4) \equiv \varphi_1 \wedge \varphi_2 \wedge \varphi_3 \wedge \neg \varphi_4$$

step 1: Отрицания саю пред атоярни  $\mathcal{F}$ -ли (възможно най-навътре ги издвоява и се оправяме с " $\Rightarrow$ " и " $\Leftarrow$ ") / \* в стъпки от клетка

step 2: Пн $\mathcal{F}$

step 3: Сн $\mathcal{F}$

step 4: Кн $\mathcal{F}$  / \* Последната стъпка - прилагане дистрибутивния закон  $(\varphi \wedge \psi) \vee \chi \equiv (\varphi \vee \chi) \wedge (\psi \vee \chi)$

$$\varphi_1 \equiv \exists z \forall x \forall y [p(f(x), y) \vee q(z, f(x))]$$

$$\varphi_1^s \equiv \forall x \forall y [p(f(x), y) \vee q(a, f(x))]$$

$$\varphi_1^{\text{final}} \equiv \varphi_1^s$$

$$\varphi_2 \equiv \exists z \forall y \forall x [\neg q(y, f(f(y))) \vee p(x, f(z))]$$

$$\varphi_2^s \equiv \forall y \forall x [\neg q(y, f(f(y))) \vee p(x, f(b))]$$

$$\varphi_2^{\text{final}} \equiv \varphi_2^s$$



генерира  $t, v$

$$\varphi_3' \equiv \forall x \forall y [\neg p(x, f(y)) \vee p(f(y), x) \vee \forall y \forall x p(x, y)] \vdash$$

НОВАТ НАЧЕЛО

$$\vdash \forall x \forall y \forall t \forall v [\neg p(x, f(y)) \vee p(f(y), x) \vee p(v, t)]$$

$$\varphi_3^{\text{final}} \equiv \varphi_3$$

$$\varphi_4' \equiv \neg \varphi_1 \vdash \forall y (\forall x \neg p(f(x), y) \vee \forall x \neg p(y, x)) \vdash$$

$$\forall y \forall x \forall z (\neg p(f(x), y) \vee \neg p(y, z))$$

$$D_1 = \{ \neg p(f(x), y), q(a, f(x)) \}$$

$$D_2 = \{ \neg q(y, f(f(y))), p(x, f(b)) \}$$

$$D_3 = \{ \neg p(x, f(y)), p(f(y), x), \neg p(v, t) \}$$

$$D_4 = \{ \neg p(f(x), y), \neg p(y, z) \}$$

Let the game begin:

Rule 1: НОВАТ-ОДНА РЕЗОНЕНТА

$$D_1 = \{ L, \neg L \} \cup D_1'$$

$$D_2 = \{ \neg L, L \} \cup D_2'$$

НОВАТ ОДНА ПРОМЕНЛИВА!

Тогаво нопираме судетитиуна  $\sigma$ , така че

$$L\sigma = \neg L\sigma$$

$$\text{Res}_L(D_1, D_2) = D_1'\sigma \cup D_2'\sigma$$

Rule 2: КОЛАНС

$$D_1 = \{ L, \neg L \} \cup D_1'$$

\* Разн. случа, когато се види конкта ./

Нопираме судетитиуна  $\sigma$ , така че

$$L\sigma = \neg L\sigma; \text{Collapse}(D_1) = D_1'\sigma \cup \{L\sigma\}$$



① Парн  $D_3$

$$\sigma = \{x/f(x), y/f(y)\}$$

$$\text{Collapse}(D_3) = \{ \neg p(x, f(y)), p(f(y), x) \} = D_5$$

② Парн  $D_4$

$$\sigma = \{y/f(x), z/f(x)\}$$

$$\text{Collapse}(D_4) = \{ \neg p(f(x), f(x)) \} = D_6$$

③ Парн  $D_1 \cup D_2$

$$\begin{aligned} & \left| \begin{array}{l} a = y_2 \\ f(x_1) = f(f(y_2)) \\ \sigma = \{y_2/a, x_1/f(a)\} \end{array} \right. \Leftrightarrow \left| \begin{array}{l} y_2 = a \\ x_1 = f(a) \end{array} \right. \\ & \sigma = \{y_2/a, x_1/f(a)\} \end{aligned}$$

$$\text{Res}(D_1, D_2) = \{ p(f(f(a)), y_1), p(x_2, f(b)) \} = D_7$$

④ Парн  $D_6 \cup D_7$

$$\sigma = \{x/b, x_2/f(b)\}$$

$$\text{Res}(D_6, D_7) = \{ p(f(f(a)), y_1) \} = D_8$$

⑤ Парн  $D_8 \cup D_8$

$$\sigma = \{x/f(a), y_1/f(f(a))\}$$

$$\text{Res}(D_8, D_8) = \emptyset$$

Заг 2 Да се пок. чрез лог. те  $\varphi_1 \Rightarrow \varphi_2$   
 $\varphi_1 \Leftarrow$  "който не получава изхвърляне"  
 $\varphi_2 \Leftarrow$  "хората не са запознати с изхвърляне"

1-130

$$\varphi_1 \Leftarrow \forall x (drink(x) \Rightarrow hangover(x))$$

$$\varphi_2 \Leftarrow \forall x (\neg hangover(x) \Rightarrow \neg drink(x))$$

- / ако някой изхвърля, то няма да  
 ни е /

$$\varphi_1 \Rightarrow \varphi_2 \quad \underline{1}, \quad (\varphi_1 \& \neg \varphi_2)$$

$$\varphi_1 \Leftarrow \forall x (\neg drink(x) \vee hangover(x))$$

$$\varphi_2 \Leftarrow \exists x (\neg hangover(x) \& drink(x))$$

$$\varphi_2 \Leftarrow (\neg hangover(a) \& drink(a))$$



$$D_1 = \{ \neg \text{drink}(x), \text{hangover}(x) \}$$

$$D_2 = \{ \neg \text{hangover}(a) \}$$

$$D_3 = \{ \text{drink}(a) \}$$

Резулт.  $D_1 \cup D_3$

$$\sigma = \{ x/a \}$$

$$\text{Res}(D_1, D_3) = \{ \text{hangover}(a) \} = D_4$$

Резулт.  $D_2 \cup D_4$

$$\text{Res}(D_2, D_4) = \emptyset$$

Заг. 3) Да се докаже, че ако  $P, Q, R$  са

$$e_1 \wedge e_2 \wedge e_3 \Rightarrow e_4, \text{ където:}$$

$$e_1 \equiv \forall x \exists y [q(y, x) \wedge \forall z (q(y, z) \Rightarrow r(z, x))]$$

$$e_2 \equiv \forall x [\exists y q(x, y) \Rightarrow \exists y (q(x, y) \wedge \neg \exists z (q(y, z) \wedge r(x, z)))]$$

$$e_3 \equiv \forall x \forall y \forall z [q(y, x) \wedge r(z, y) \Rightarrow q(z, x)]$$

$$e_4 \equiv \forall x \neg q(x, x)$$

П-до

$$e_1' \equiv \forall x \exists y [q(y, x) \wedge \forall z (\neg q(y, z) \vee r(z, x))] \vdash$$

$$\vdash \forall x (\exists y \forall z [q(y, x) \wedge (\neg q(y, z) \vee r(z, x))])$$

$$e_1^f \equiv e_1^s \equiv \forall x \forall z [\underbrace{q(f(x), x)}_{D_1} \wedge \underbrace{(\neg q(f(x), z) \vee r(z, x))}_{D_2}]$$

$$e_2' \equiv \forall x [\forall y \neg q(x, y) \vee \exists y (q(x, y) \wedge \forall z (\neg q(y, z) \vee r(x, z)))] \vdash$$

$$\vdash \forall x (\exists y \forall y \forall z [\neg q(x, y) \vee (q(x, y) \wedge (\neg q(y, z) \vee r(x, z)))]$$

$$e_2^s \vdash \forall x \forall y \forall z [\neg q(x, y) \vee (q(x, g(x)) \wedge (\neg q(g(x), z) \vee r(x, z)))]$$

$$e_2^f \vdash \forall x \forall y \forall z [\neg q(x, y) \vee q(x, g(x)) \wedge (\neg q(g(x), z) \vee r(x, z))] \vdash$$



$$e_3^f \leq \forall x \forall y \forall z (\underbrace{\neg q(y, x) \vee \neg r(z, y) \vee q(z, x)}_{D_5})$$

$$e_4^1 \leq \neg e_4 \neq \underbrace{\exists x q(x, x)}_{D_6}$$

$$e_4^f \leq e_4^s \leq \underbrace{q(a, a)}_{D_6}$$

$$D_1 = \{q(f(x_1), x_1)\}$$

$$D_2 = \{\neg q(f(x_2), z_2), r(z_2, x_2)\}$$

$$D_3 = \{\neg q(x_3, y_3), q(x_3, q(x_3))\}$$

$$D_4 = \{\neg q(x_4, y_4), \neg q(g(x_4), z_4), \neg q(x_4, z_4)\}$$

$$D_5 = \{\neg q(y_5, x_5), \neg r(z_5, y_5), q(z_5, x_5)\}$$

$$D_6 = \{q(a, a)\}$$

① Резол.  $D_2 \cup D_5$

$$\sigma = \{y_2/y_5, z_2/z_5\}$$

$$\text{Res}(D_2, D_5) = \{\neg q(f(x_2), z_5), \neg q(y_5, x_5), q(z_5, x_5)\}$$

$D_7$

② Резол.  $D_3 \cup D_7$

$$\sigma = \{x_3/f(x_2), z_5/q(x_3)\}$$

$$\text{Res}(D_3, D_7) = \{\neg q(f(x_2), y_3), \neg q(y_5, x_5), q(g(f(x_3)), x_5)\} = D_8$$

③ Резол.  $D_4 \cup D_8$

$$\sigma = \{x_4/f(x_3), z_4/x_5\}$$

$$\text{Res}(D_4, D_8) = \{\neg q(f(x_3), y_4), \neg q(f(x_3), z_5), \neg q(f(x_2), y_3), \neg q(y_5, x_5)\} = D_9$$

④ Резол.  $D_9$  ;  $\sigma = \{x_1/x_3, y_4/x_5\}$



$$\text{Collapse}(D_3) = \{ \neg q(f(x_3), y_4), \neg q(y_5, x_5) \} = D_0$$

⑤ Разн.  $D_1$  и  $D_0$

$$\sigma = \{ x_3/x_1, y_4/x_1 \}$$

$$\text{Res}(D_1, D_0) = \{ \neg q(y_5, x_5) \} = D_{11}$$

⑥ Разн.  $D_0$  и  $D_{11}$

$$\sigma = \{ y_5/a, x_5/a \}$$

$$\text{Res}(D_0, D_{11}) = \square$$

### Дополнительни задачи

задач - 11-  $\varphi_1 \& \varphi_2 \Rightarrow \varphi_3$ , което:

$$\varphi_1 \leq \forall x [\forall y (s(y) \Rightarrow p(y, x)) \Rightarrow \neg r(x)]$$

$$\varphi_2 \leq \exists y \forall x [\neg s(x) \vee \neg q(x, y)]$$

$$\varphi_3 \leq \exists x \forall y [r(y) \Rightarrow \exists z (\neg p(z, y) \& \neg q(z, x))]$$

Частично решение

Трѣбва да стигнете до:

$$D_1 = \{ s(f(x)), r(x) \}$$

$$D_2 = \{ \neg p(f(x), x), \neg r(x) \}$$

$$D_3 = \{ \neg s(x), \neg q(x, a) \}$$

$$D_4 = \{ r(g(x)) \}$$

$$D_5 = \{ p(z, g(x)), q(z, x) \}$$

Примери извода

① Разн.  $D_1$  и  $D_4 \Rightarrow D_6$

② Разн.  $D_2$  и  $D_4 \Rightarrow D_7$

③ Разн.  $D_3$  и  $D_5 \Rightarrow D_8$

④ Разн.  $D_8$  и  $D_7 \Rightarrow D_9$

⑤ Разн.  $D_6$  и  $D_9 \Rightarrow \square$



задача -11-  $\varphi_1 \& \varphi_2 \Rightarrow \varphi_3$  верно

$$\varphi_1 \equiv \exists z \forall y [s(y) \vee q(z, y)]$$

$$\varphi_2 \equiv \forall y \exists z [\neg r(z) \& \forall t (p(z, t) \Rightarrow \neg q(y, t))]$$

$$\varphi_3 \equiv \exists y [\forall z (\neg s(z) \Rightarrow \neg p(y, z)) \& \neg r(y)]$$

Частичное решение:

Получаем:

$$D_1 = \{s(y), q(a, y)\}$$

$$D_2 = \{\neg r(f(y))\}$$

$$D_3 = \{\neg p(f(y), t), \neg q(y, t)\}$$

$$D_4 = \{\neg s(g(y)), r(y)\}$$

$$D_5 = \{\neg p(y, g(y)), r(y)\}$$

Примеры выбора из:

① Разрн  $D_2$  и  $D_4 \Rightarrow D_5$

② Разрн  $D_2$  и  $D_5 \Rightarrow D_5$

③ Разрн  $D_1$  и  $D_3 \Rightarrow D_5$

④ Разрн  $D_5$  и  $D_4 \Rightarrow D_5$

⑤ Разрн  $D_5$  и  $D_3 \Rightarrow D_5$

Genex ! ☺