

# Основные предикаты для списка

## [Head | Tail]

### I Member

- $\text{member}(X, [X|T]).$
- $\text{member}(X, [-|T]) :- \text{member}(X, T).$
- $\text{member}(X, L) :- \text{remove}(X, L, -).$
- $\text{member}(X, L) :- \text{append}(-, [X|L], L).$

### II append

- $\text{append}([ ], Y, Y).$
- $\text{append}([X|T], Y, [X|Result]) :- \text{append}(T, Y, Result).$

### III Insert

- $\text{insert}(X, L, Result) :- \text{append}(A, B, L), \text{append}(A, [X|B], Result).$
- $\text{insert}(X, L, Result) :- \text{remove}(X, L, N).$

### IV Remove.

- $\text{remove}(X, [X|T], T).$
- $\text{remove}(X, [H|T], [H|N]) :- \text{remove}(X, T, N).$

### V Subset (с корректно упорядоченными элементами)

- $\text{sub}([ ], [ ]).$
- $\text{sub}([X|T], Res) :- \text{sub}(T, Res).$
- $\text{sub}([X|T], [X|Res]) :- \text{sub}(T, Res).$

### VI subset (с неупорядоченными элементами)

- $\text{psub}(L, Res) :- \text{sub}(L, Q), \text{perm}(Q, Res).$

### VII Permutations

- $\text{perm}([X|T], Result) :- \text{perm}(T, R), \text{insert}(X, R, Result).$

VIII First element

first(F, [F|\_]).

IX Second element

• second(S, [\_|S1]).

• second(S, [H|T]) :- first(S, T).

X Last element

• last(X, L) :- append(\_, [X], L).

• last(X, [\_|T]) :- last(X, T).

XI Prefix

prefix(P, L) :- append(P, \_, L). | prefix([ ], \_).

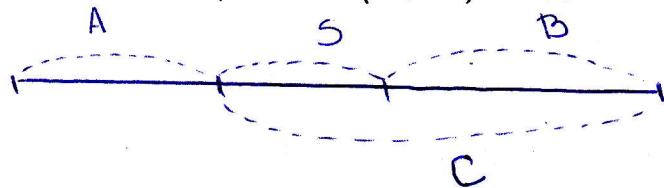
prefix([H|T], [H|Q]) :- prefix(T, Q).

XII Suffix

suffix(S, L) :- append(\_, S, L).

XIII Sublist

sublist(S, L) :- append(\_, C, L), append(S, \_, C).



XIV Is subset?

isSubset(S, L) :- not(member(X, S)), not(member(X, L)).

XV equal lists

equal(L1, L2) :- issubset(L1, L2), issubset(L2, L1).

XVI intersection

intersection(X, A, B) :- member(X, A), member(X, B).

XVII union

union(X, A, B) :- member(X, A); member(X, B)

XVIII  $\text{sub}(X, A, B) :- \text{member}(X, A), \text{not member}(X, B).$

XIX Reverse

- $\text{reverse}([J], [J])$
- $\text{reverse}(R, [H, L]) :- \text{reverse}(RL, L), \text{append}(R2, [H], R).$
- $\text{reverse}(R, L) :- \text{reverse}([J], L, R).$
- $\text{reverse}(S, [J], S).$
- $\text{reverse}(S, [H | T], R) :- \text{reverse}([H | S], T, R).$

XX length

- $\text{length}([J], 0).$
- $\text{length}([H | T], N) :- \text{length}(T, M), N \text{ is } M + 1.$

XXI Nth element

- $\text{nth}([X1 | J], 0, X).$
- $\text{nth}([ | I | J], N, X) :- \text{nth}(I, M, X), N \text{ is } M + 1.$

XXII Sorting

1. Insertion sort
  - $\text{sort}([J], [J]).$
  - $\text{sort}([H | Q, J], L) :- \text{sort}(Q, M), \text{sort}_{in}(H, M, L).$

- $\text{sort}_{in}(X, [ ], [X]).$
- $\text{sort}_{in}(X, [H | T], [X, H | T]) :- X = < H.$
- $\text{sort}_{in}(X, [H | Q], [H | L]) :- H < X, \text{sort}_{in}(X, Q, L).$

2. Selection sort

- $\text{sort}([J], [J]).$
- $\text{sort}(L, [H | Q]) :- \text{min}(H, L, M), \text{sort}(M, Q).$

- $\text{min}(X, [X], [J]).$
- $\text{min}(H, [H | T], T) :- \text{min}(X, T, -), H = < X.$
- $\text{min}(X, [H | Q], [H | L]) :- \text{min}(X, Q, L), H > X.$

- $\text{sort}([J], [J]).$
- $\text{sort}([M | S], L) :- \text{min}(M, L, N), \text{remove}(M, L, N), \text{sort}(N, S).$

- $\text{min}(M, [M]).$
- $\text{min}(M, [H | T]) :- \text{min}(N, T), \text{min}2(M, H, N).$
- $\text{min}2(A, A, B) :- A = < B.$
- $\text{min}2(B, A, B) :- \text{not}(A = < B).$

### 3. Quick Sort

quicksort ([J,J]).

quicksort ([X|XsJ, Res]) :- part(X, XsJ, L, R), quicksort(L, SL), quicksort(R, SR), append(SL, [X|SR], Res).

part( \_, [J], [J], [J]).

part(X, [H|T], [H|L], R) :- H < X, part(X, T, L, R).

part(X, [H|T], L, [H|R]) :- H >= X, part(X, T, L, R).

### 4. Is sorted?

sorted(X) :- not (append( \_, [A,B|\_], X), A > B).

// sorted increasingly

XXIII

Sum of elements

sum([ ], 0).

sum([H|T], N) :- sum(T, M), N is M+H.

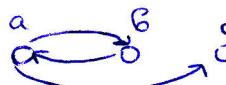
## Графы.

Неко  $G$  е представен като единица от редпор  
[[ $v_0, u_0$ ], [ $v_1, u_1$ ], ..., [ $v_n, u_n$ ]]

I Най бърз граф

- $\text{path}(G, X, Y) :- \text{member}([X, Y], G).$
- $\text{path}(G, X, Y) :- \text{member}([X, Z], G), \text{path}(G, Z, Y).$

Наподобняваме заглавките



$$G = [[a, b], [a, c], [b, a]]$$

Безупречн.

↓

- $\text{path}(G, X, Y) :- \text{path}(G, X, Y, [X], -).$

$\text{path}(G, X, X, [X], [X]).$

$\text{path}(G, X, Y, V, [X|P]) :- \text{member}([X, Z], G), X \neq Y,$   
 $\text{not}(\text{member}(Z, V)),$   
 $\text{path}(G, Z, Y, [Z|V], P).$

II Учебен бърз граф?

$\text{cycle}(G) :- \text{member}([X, Y], G), \text{path}(G, Y, X, [Y], -).$

$\text{cycle}(G, [X|T]) :- \text{member}([X, Y], G), \text{path}(G, Y, X, [Y], T).$

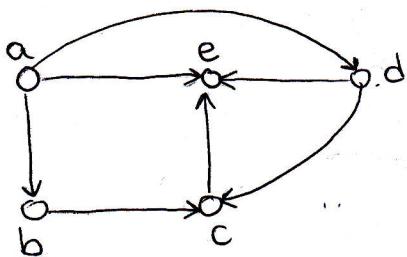
III DFS

$\text{dfs}(G, X, T) :- \text{dfs}(G, [X], [X, T]).$

$\text{next}(G, U, S, F, V) :- \text{member}([U, V], G),$   
 $\text{not}(\text{member}(V, S)),$   
 $\text{not}(\text{member}(V, F)).$

$\text{dfs}(G, [ ], -, [ ])$  visited  
 $\text{dfs}(G, [S|SS], F, [S|NR]) :-$   
список  $\text{next}(G, S, [S|SS], F, N),$   
 $\text{dfs}(G, [N|SS], F, NR).$

$\text{dfs}(G, [S|SS], F, R) :- \text{not}(\text{next}(G, S, [S|SS], F, -)),$   
 $\text{dfs}(G, SS, [S|F|R], R).$



$$G = \{[E_{a,b}], [a,e], [a,d], [b,c], [c,e], [d,g], [d,e], [e,g]\}$$

$\text{dfs}(G, a, T) : - \text{dfs}(G, [a], [ ], T)$

$\text{dfs}(G, [b, a], [ ], R) \quad R = ([a, b] | R)$

$\text{dfs}(G, [c, b, a], [ ], R) \quad R = ([b, c], [a, b] | R)$

$\text{dfs}(G, [e, c, b, a], [ ], R) \quad R = ([c, e], [b, c], [a, b] | R)$

$\text{dfs}(G, [a, b, a], [e], R)$

$\text{dfs}(G, [b, a], [c, e], R)$

$\text{dfs}(G, [a], [b, c, e], R)$

$\text{dfs}(G, [d, a], [b, c, e], R) \quad R = ([a, d], [c, e], [b, c], [a, b] | R)$

$\text{dfs}(G, [a], [d, b, c, e], R)$

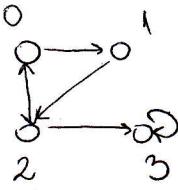
$\text{dfs}(G, [ ], [a, d, b, c, e], R)$

$T = [a, b], [b, c], [d, e], [a, d]$

# DFS

\*using stack

1. Push 0 on the stack
2. Pop the last element of the stack
3. See which elements it's connected to.
4. Put them on the stack (1, 2).
- 5 Remove 2 from the stack.
6. See which element it's connected to (0, 3).
7. Put 3 on the stack. (0 is already visited).
8. Remove 3 from the stack. // Just 1 left.
9. Remove 1 from the stack
10. Done.



Stack: 0

Visited: -

Stack: 1, 2

Visited: 0

Stack: 1, 3

Visited: 0, 2

Stack: 1

Visited: 0, 2, 3

Stack -

Visited: 0, 2, 3, 1

$\text{dfs}(G, X, T) :- \text{dfs}(G, [X], [], T)$ .

$\text{dfs}(G, [], T, T)$ .

$\text{dfs}(G, [Current | Rest], Visited, ([Current, Next] | R)) :-$

$\text{next}(G, Current, Next, [Current | Rest], Visited),$   
 $\text{dfs}(G, [Next, Current | Rest], Visited, R)$ .

$\text{dfs}(G, [Current | Rest], Visited, R) :-$

$\text{not}(\text{next}(G, Current, - [Current | Rest], Visited)),$   
 $\text{dfs}(G, Rest, [Current | Visited] | R)$ .

So, we need:

1) Stack

2) Visited

3) neighbours which are neither on the stack nor in the visited.

Prolog:

$\text{neighbours}(G, U, V, Stack, Visited) :-$

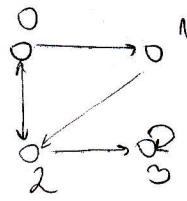
$\text{member}(U, V, G),$   
 $\text{not}(\text{member}(U, Stack)),$   
 $\text{not}(\text{member}(V, Visited))$ .

## BFS

\* uses a queue

Queue: 0

Visited: -



Queue: 1, 2

Visited: 0

Queue: 2

Visited: 0, 1

Queue: 3

Visited: 0, 1, 2

Queue: -

Visited: 0, 1, 2, 3

## Prolog:

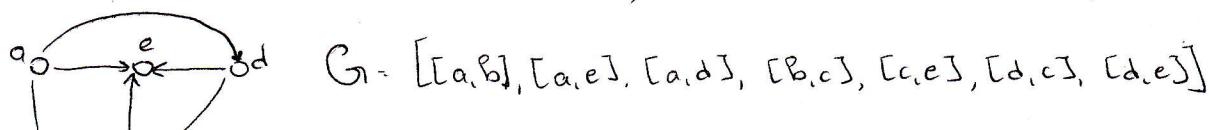
bfs(G, X, T) :- bfs(G, [X], [J, T]).

bfs(G, [ ], - , [J]).

bfs( G, [Current|Queue], Visited, ([Current|Next]|R) ) :-  
next( G, Current, Next, [Current|Queue], Visited ),  
~~append( [Current|Queue], [Next], NewQueue ),~~  
bfs( G, NewQueue, Visited, R ).

bfs( G, [Current|Queue], Visited, R ) :-

not(next( G, Current, - , [Current|Queue], Visited )),  
bfs( G, Queue, [Current|Visited], R ).



$G = [[a,b], [a,e], [a,d], [b,c], [c,e], [d,c], [d,e]]$

bfs( G, a, T ) :- bfs( G, [a], [J, T] )

bfs( G, [a,b], [J, R=[a,b]])

bfs( G, [a,b,e], [J, [a,b,[a,e]]] )

bfs( G, [a,b,e,d], [J, [[a,b],[a,e],[a,d]]] )

bfs( G, [b,e,d], [a], [[[a,b],[a,e],[a,d]]] )

bfs( G, [b,e,d,c], [a], [[[a,b],[a,e],[a,d],[b,c]]] )

bfs( G, [e,d,c], [b,a], R )

bfs( G, [d,c], [e,b,a], R )

bfs( G, [c], [d,e,b,a], R )

bfs( G, [ ], [c,d,e,b,a], R )

# Работа с числом и арифметични операции

I Генератор на едното число  
nat(0).  
nat(X) :- nat(Y), X is Y+1.

] Не същ за разпознаване

Операции за присваяване

II Най-големи същ денумери

gcd(X, Y, Z) :- X mod Y == 0.

gcd(X, Y, Z) :- R is X mod Y, gcd(Y, R, Z).

III Генератор на едното четно число  
evenNumber(0).

evenNumber(X) :- evenNumber(Y), X is Y+2.

IV Разпознаване на едното четно число  
evenNumber(0).

evenNumber(X) :- X >= 1, Y is X-2, evenNumber(Y).

V Разпознаване | Генератор на едното четно число  
evenNumber(X) :- nat(X), X mod 2 == 0.

VI Фибоначи.

Fibonacci(0, 1).

Fibonacci(A, B) :- Fibonacci(B, C), A is B+C.

Fibonacci(X) :- Fibonacci(X, -).

• Fibonacci(X, X, -).

Fibonacci(X, A, B) :- C is A+B, Fibonacci(X, C, A).

Fibonacci(X) :- Fibonacci(X, 0, 1).

VII Between

Between(A, B, A) :- A = < B.

Between(A, B, X) :- A < B, C is A+1, Between(C, B, X).

VIII Prime

isPrime(N) :- N is N//2, not(Between(2, N, X)), N mod X == 0.

prime(N) :- nat(X), isPrime(N).

IX

Общий естественный номер

pair ( $X, Y$ ) :- nat ( $N$ ), Between ( $0, N, X$ ),  $Y$  is  $N - X$ .

X

Генератор для членов ряда

•  $Z(Y)$  :- nat ( $X$ ),  $s(X, Y)$ .

$s(X, X)$ .

$s(X, Y)$  :-  $X > 0$ ,  $Y$  is  $-X$ .

•  $Z(X)$  :- nat ( $Y$ ), decode ( $Y, X$ ).

decode ( $X, Y$ ) :-  $X \bmod 2 =:= 0$ ,  
 $Y$  is  $-X // 2$ .

decode ( $X, Y$ ) :-  $X \bmod 2 =:= 1$ ,  
 $Y$  is  $X // 2 + 1$

0 1 -1 2 -2 3 5  
0 1 2 3 4 5 6

# Задачи

1. Да се реализира предикат, който на зададено число, създава списък с разлагането му на прости делители (без броящия предикат).

Pred:

$\text{primeDivisor}(N, P) :- N \equiv 0 \pmod{P}, \text{isPrime}(P),$

$\text{minPrimeDivisor}(N, P) :- \text{primeDivisor}(N, P),$   
 $\forall Q \in \text{primeDivisors}(N), Q < P.$

$\text{decomposition}(0, [ ]).$

$\text{decomposition}(1, [ ]).$

$\text{decomposition}(N, [P | T]) :- \text{minPrimeDivisor}(N, P),$   
 $N \equiv 0 \pmod{P},$   
 $\text{decomposition}(N/P, T).$

2. Да се реализира предикат, който на зададено число създава списък с простите делители (без броящия предикат) и техните степени от каквато либо предикат.

Pred:

$\text{degrees}(0, [ ]).$

$\text{degrees}(1, [ ]).$

$\text{degrees}(N, D) :- \text{minPrimeDivisor}(N, P),$   
 $\text{degrees}(N/P, D),$   
 $\text{add}(\overline{P}, P, D).$

$\text{add}(\overline{P}, P, D) :- \text{append}([P, CR, CJ, E], \overline{T}), C \in CR \wedge,$   
 $\text{append}([P, CJ, E], \overline{D}).$

$\text{add}(\overline{P}, P, D) :- \text{not}(\text{member}([P, -], \overline{T})).$

3. Да се генерира равнинна фигура във вид  $0 \leq x \leq 1000, 0 \leq y \leq 1000$ , която да е симетрична по  $y = 3*(x+2)^5 + (k+3)^2$ .

Pred:

$\text{format}(X) :- \text{Between}(0, 1000, Y), \text{Between}(0, 1000, Z),$   
 $X \equiv 3*(Y+2)^5 + (Z+3)^2, X \geq 0, X \leq 1000.$

Зад 4.

Pred:

Da се генерира равнинна фигура с цели координати, които се намират във вид  $(x, y)$  предвидността на енкодинга с новите  $x_1$  и  $y_1$ .

$\text{ellipse}(X, Y) :- \text{Between}(-113, 113, X),$   
 $\text{Between}(-71, 71, Y),$   
 $(X*X)/113 + (Y*Y)/49 \leq 1.$

$11/13 \in \text{менюнирано}$  Генерации  
 $5/2 = 2,5$

5.  
Pue

Предикат, генерирующий пары чисел от  $s$ .

genKS(1, S, [S]).

genKS(K, S, [X|T]) :- K > 1, Between(0, S, X), K1 is K-1,  
S1 is S-X, genKS(K1, S1, T).

a) предикат, генерирующий пары четных чисел между  
genC(X) :- nat(N), genKS(4, N, X).

b) предикат, генерирующий пары двойки четных чисел  
pair(X, Y) :- nat(N), genKS(2, N, [X, Y]).

c) предикат, генерирующий пары кратных пяти из четных чисел

gen-all([ ]).

gen-all(X) :- nat(N), Between(1, N, K), S is N-K,  
genKS(K, S, X).

6  $\begin{array}{l} \text{hand} \rightarrow a_n = 1 \\ \quad \quad \quad a_0 = 1 \\ \quad \quad \quad a_n = 2a_{n-1} + n \end{array} \quad \begin{array}{l} \text{hand} \rightarrow b_0 = 1 \\ \quad \quad \quad b_n = 3b_{n-1} + n^2 - 1 \end{array}$

Даде наумие предикат p(N), генерирующий пары чисел от  
поди  $a_i + b_j$ .

Pue:

a(0, 1)

a(N, X) :- N > 0, N1 is N-1, a(N1, Y), X is 2\*N1+N.

b(0, 1)

b(N, X) :- N > 0, N1 is N-1, b(N1, Y), X is 3\*N1+N\*N-1.

p(N) :- nat(X), Between(0, X, I), J is X-I, a(I, A), b(J, B),  
N is A+B.

/no  
используем

s(-, [A|As], [B|Bs], Res) :- (member(X, Bs), Res is A+X);  
(member(X, [A|As]), Res is B+X).

s(N, [A|As], [B|Bs], R) :- NA is 2\*A+N,  
NB is 3\*B+N\*N-1,  
NR is N+1

s(N1, [NA, A|As], [NB, B|Bs], R).

p(N) :- s(1, [1], [1], N).

• 7. Дефиниция на L-списък:

$[I]$  е ~~списък~~

Aко  $A$  и  $B$  са L-списъци, то  $[A, [A, B]]$  също е L-списък.  
Да се покаже предикат, генериращ всички L-списъци.

$$\left[ \begin{array}{c} [I] \\ A \end{array} \right], \left[ \begin{array}{c} [I] \\ A \end{array} \right], \left[ \begin{array}{c} [I] \\ B \end{array} \right] ] ] ]$$

$$\text{rang}([I]) = 1$$

$$\text{rang}([A, [A, B]]) = \text{rang}(A) + \text{rang}(B)$$

Пример:

$r([I], N)$ .

$r([A, [A, B]], N) :- N > 1, M \text{ is } N-1, \text{ Between } (1, M, RA),$   
 $RB \text{ is } N-RA, r(A, RA), r(B, RB).$

$\text{alpha}(L) :- \text{nat}(N), r(L, N).$

$$\left[ \begin{array}{c} [I, [I, I]] \\ [I] \\ A \end{array} \right] \quad r=2$$

$$\left[ \begin{array}{c} [I] \\ [I] \\ B \end{array} \right]$$

8. Да се дефинира предикат  $s(M, L)$ , който генерира в  $M$  всички  
списъци с елементи от  $L$ .

$$L = [1, 2]$$

$$M = [1, 1, 1, 2]$$

$$M = [2, 1, 2, 1, 2]$$

Пример:

$s([I], L)$ .

$s([H | T], L) :- s(T, L), \text{member}(H, L).$

$s(M, L) :- \text{nat}(N), s(N, M, L).$

$s(0, [I], L).$

$s(N, [H | T], L) :- N > 0, \text{member}(H, L), N1 \text{ is } N-1, s(N1, T, L).$

9. Дефиниция на списък от дърво:

-  $[I]$  е дърво.

- Ако  $A$  и  $B$  са дървата,  $[A, B]$  също е дърво.

Да се покаже предикат, генериращ всички дървата.

Пример:

$t(T) :- \text{nat}(N), t(N, T).$

$t(0, [I]).$

$t(N, [A, B]) :- N > 0, NA \text{ is } N-1, \text{Between}(0, NA, NA), NB \text{ is } N-NA,$   
 $t(NA, A), t(NB, B).$

10. Да се покаже предикат  $st(G, T)$ , който по зададен граф  $G = [V, E]$ ,  
генерира покриващо дърво  $T = [V, EA]$ .

Пример:

$st([V, E], [V, EA]) :- V = [X | V1], st([V, E], [X], V1, EA).$

$st(G, V, [J, [J]]).$

$st(G, V, N, [X, Y] | R) :- \text{edge}(G, X, Y), \text{member}(X, V), \text{member}(Y, N),$   
настечки  $\downarrow$   $\downarrow$  Нестечки  $\text{remove}(Y, N, NN), st(G, [Y | V], NN, R).$

8. Нека  $G$  е представен като списък от редка. т.е  
 $G = \{[A, B]\}$  или редка от  $A$  до  $B$ .

Да се дефинира предикат  $p(G)$ , който щаден ориентиран  
граф  $G$  без изолирани върхове проверява дали върховете на  $G$   
могат така да се "съединят" в 2 цвета, че никоя три редка да  
не образуваат единувален триъгълник.

P-ue:

% vertices([E,V]) - gathers into  $V$  all vertices from list of edges  $E$   
vertices([J,J]).

vertices([[A,B]|Tail],V) :-

    vertices(Tail,TV), add(A,TV,VN), add(B,VN,V).

% add(X,L,Res) - adds  $X$  in  $L$  if it's not already stored.

add(X,L,L) :- ~~member(X,L)~~

add(X,L,[X|L]) :- not(member(X,L)).

% subset(L,S)

subset(T,[J])

subset([H|T],S) :- subset(T,S).

subset([H|T],[H|S]) :- subset(T,S).

% triangle(G,A,B,C) :- find 3 vertices forming a triangle.

triangle(G,A,B,C) :-

    member(G,[A,B]), member(G,[B,C]), member(G,[C,A]),  
    not(A=B), not(B=C), not(C=A).

% sameColour(A,B,C,Coloured) - determines if A,B,C are the same  
colour.

sameColour(A,B,C,Coloured) :- member(Coloured,A),  
                              member(Coloured,B),  
                              member(Coloured,C).

sameColour(A,B,C,Coloured) :- not(member(Coloured,A)),  
                              not(member(Coloured,B)),  
                              not(member(Coloured,C)).

% solution.

$p(G)$  :- vertices(G,V), subset(V, Colored),  
        not(triangle(G,A,B,C)), sameColor(A,B,C, Colored)).

9

Нека  $G$  е графът на изолирани върхове, представен чрез списък от ребра.

Да се дефинира предикат  $p(G)$ , който ще даде ориентирани граф  $G$  проверява дали върховете на  $G$  могат така да се уредят по дадена форма, че всяко ребро да свързва разнокраси върхове.

P-ue:

$p(G) :- \text{vertices}(G, V), \text{subset}(V, \text{Colored}),$   
 $\text{not } (\text{member}([A, B], G) \wedge \text{sameColor}(A, B, \text{Coloured})).$

% sameColor(A, B, Coloured)

sameColor(A, B, Coloured) :- member(Coloured, A),  
 $\text{member}(\text{Coloured}, B).$

sameColor(A, B, Coloured) :- not(member(Coloured, A)),  
 $\text{not}(\text{member}(\text{Coloured}, B)).$

10 ? - flatten([a, [b, [c, d], e]], X)  
 $X = [a, b, c, d, e]$

P-ue:

flatten(X, [X]) :- not(is-list(X)).

flatten([ ], [ ]).

flatten([H|T], Zs) :- flatten(H, H1), flatten(T, T1), append(H1, T1, Zs).

is-list([ ])

is-list([-|T]) :- isList(T).

11. Да се дефинира предикат премахващ ненедоволителните набляденияни елементи.

P-ue:

compress([ ], [ ])

compress([X], [X])

compress([X|Y|T], [X|Zs]) :- X \= Y, compress([Y|T], Zs).

compress([X|T], Zs) :- compress([T], Zs).

12. Да се напише предикат, който групира ненедоволителните наблядени елементи на листа в подлист.

? - pack([a, a, a, a, b, c, c, a, a, d, e, e, e, e], X)  
 $X = [[a, a, a, a], [b], [c, c], [a, a], [d], [e, e, e, e]]$

P-ue:

pack([ ], [ ])

pack([X|Xs], [Z|Zs]) :- transfer(X, Xs, Ys, Z), pack(Ys, Zs).

transfer(X, [ ], [X]).

transfer(X, [Y|Ys], [Y|Zs], [X]) :- X \= Y.

transfer(X, [X|Xs], Ys, [X|Zs]) :- transfer(X, Xs, Ys, Zs).

13: Run-length encoding of a list

?- encode([a,a,a,a,b,c,c,a,a,d,e,e,e,e], X)

X = ([4,a], [1,b], [2,c], [2,a], [1,d], [4,e]).

P-ue:

- encode(L<sub>1</sub>, L<sub>2</sub>): - pack(L<sub>1</sub>, L), transform(L, L<sub>2</sub>).  
transform([J, J]).
- transform([[X|Xs]|T], [[N,X]|Z]): - length([X|Xs], N), transform(T, Z).

14 Modified run-length encoding

P-ue:

- transform([J, J]).
- transform([[X|Xs]|T], [[N,X]|Z]): - length([X|Xs], N), transform(T, Z).
- transform([[X|Xs]|T], [X|Z]): - transform(T, Z).

encode-mod(L<sub>1</sub>, L<sub>2</sub>): - pack(L<sub>1</sub>, L), transform(L, L<sub>2</sub>).

- encode-mod(L<sub>1</sub>, L<sub>2</sub>): - encode(L<sub>1</sub>, L), strip(L, L<sub>2</sub>).  
strip([J, J]).
- strip([[N,X]|T], [X|Z]): - strip(T, Z).
- strip([[N,X]|T], [[N,X]|Z]): - N>1, strip(T, Z).

15 Decode a run-length encoded list.

P-ue:

- decode([J, J]).
- decode([X|Ys], [X|Zs]): - not(is-list(X)), decode(Ys, Zs).
- decode([[1,X]|Ys], [X|Zs]): - decode(Ys, Zs).
- decode([[N,X]|Ys], [X|Zs]): - N>1, N is N-1, decode([[N-1,X]|Ys], Zs).

16 Direct run-length encoding

P-ue:

- encode-direct([J, T]).
- encode-direct([X|Xs], [Z|Zs]): - count(X, Xs, Ys, K, Z),  
encode-direct(Ys, Zs).

Count(X, [J, T], N, X).

Count(X, [J, T], N, [N, X]): - N>1.

Count(X, [Y|Ys], [Y|Ys], N, X): - X1=Y.

Count(X, [Y|Ys], [Y|Ys], N, [N, X]): - N>1, X1=Y.

Count(X, [X|Xs], Ys, K, T): - X1 is K+1, count(X, Xs, Ys, K, T).

Ys is the list that remains from the list Xs when all the leading copies are removed.

T is the term [N, X], where N is K plus the number of X's that can be removed from Xs.

In the case of N=1, T is X instead of [1, X].

14. Duplicate the elements of a list.

P-ue:  $\text{duplicate}([J], [J]).$

$\text{duplicate}([H|T], [H, H|T]) :- \text{duplicate}(T, X).$

18: Duplicate the elements of a list a given number of times  
%  $\text{duplicate}(L, N, LN).$

P-ue:  $\text{duplicate}(L, N, LN) :- \text{duplicate}(L, N, LN, N).$

$\text{duplicate}([J], \_, [J], \_).$

$\text{duplicate}(\_, [Xs], N, Ys, O) :- \text{duplicate}(Xs, N, Ys, O).$

$\text{duplicate}([X|Xs], N, [X|Ys], K) :- K > 0, KN \text{ is } K-1,$

$\text{duplicate}(Xs, N, Ys, KN).$

19: Drop every n'th element of a list.

P-ue:  $\text{drop}(L, N, LN) :- \text{drop}(L, N, LN, N).$

$\text{drop}(\_, \_, \_).$

$\text{drop}(\_, [Xs], N, Ys, N) :- \text{drop}(Xs, N, Ys, N).$

$\text{drop}([X|Xs], N, [X|Ys], K) :- K > 1, KN \text{ is } K-1, \text{drop}(Xs, N, Ys, KN).$

20 Split a list into two parts - the length of the first part is given.  
// Do not use any predefined predicates  
%  $\text{split}(L, N, LN, L2).$

P-ue:  $\text{split}([], 0, [], []).$

14  $\text{split}([Xs], N, [Xs], Zs) :- N > 0, NN \text{ is } N-1, \text{split}(Xs, NN, Ys, Zs).$

21 Extract a slice from a list

Given 2 indices,  $I$  and  $K$ , the slice is the list containing the elements between the  $I$ 'th and  $K$ 'th element of the original list  
Start counting the elements with 1.

?-  $\text{slice}([a,b,c,d,e,f,g,h,i,k], 3, 4, L)$   
 $L = [c, d, e, f, g]$

P-ue:  $\text{slice}([X|T], I, I, [X]).$

$\text{slice}([X|Xs], I, K, [X|T]) :- I > 1, KN \text{ is } K-1, \text{slice}(Xs, I, KN, T).$

$\text{slice}(\_, [Xs], I, K, Ys) :- I > 1, IN \text{ is } I-1, KN \text{ is } K-1,$   
 $\text{slice}(Xs, IN, KN, Ys).$

22. Rotate a list  $N$  places to the left

rotate ([a,b,c,d,e,f,g,h], 3, [d,e,f,g,h,a,b,c]).

rotate ([a,b,c,d,e,f,g,h], -2, [g,h,a,b,c,d,e,f]).

P-ue: rotate(L1, N, L2) :-  $N > 0$ , length(L1, N1),  $N2$  is  $N \bmod N1$ ,  
rotate-left(L1, N2, L2).

rotate-left(L, 0, L).

rotate-left(L, N, L2) :- split(L, N, A, B), append(B, A, L2).

23. Remove left the  $[K|T]$  element of the list  $L$  and return  $(B, A, L2)$ .

% remove-at(X, L, K, Res).

L = [a,b,c,d,e,f,g,h].

remove-at(X, [X|T], 1, T).

remove-at(X, [H|T], K, [H|Res]) :-  $K > 1$ ,  $K1$  is  $K-1$

remove-at(X, T, K1, Res).

24. Insert an element at a given position into a list.

% insert-at(X, L, K, R).

P-ue: insert-at(X, L, K, R) :- remove-at(X, R, K, L).

25. Create a list containing all integers within a given range.

% range(A, B, L) :-  $\{ i \in \mathbb{Z} : A \leq i \leq B \}$ .

P-ue: range([A, A, [A]).

range(A, B, [A|Res]) :-  $A < B$ ,  $A1$  is  $A+1$ , range(A1, B, Res).

26. Extract a given number of randomly selected elements from a list:

? - rnd-select([a,b,c,d,e,f,g,h], 3, L).

L = [e, d, a].

P-ue: rnd-select([ ], 0, [ ]).

rnd-select(Xs, N, [x1|xs]) :-  $N > 0$ , length(Xs, L),  $i$  is random(L)+1,  
remove-at(X, Xs, I, Ys),  $NA$  is  $N-1$ ,  
rnd-select(Ys, NA, Zs).

27. Draw  $N$  different numbers from the set  $1\dots M$

% lotto(N, M, L)

P-ue: lotto(N, M, L) :- range(1, M, L1), rnd-select(L1, N, L).

28. Generate a random permutation of the elements of a list.

% rnd-perm(L1, L2).

P-ue: 25  
rnd-perm(L1, L2) :- length(L1, N), rnd-select(L1, N, L2).

29. A list of prime numbers.  
%prime\_list(A,B,L)

p-list(A,B,[J]) :- A > B.

p-list(A,B,[A|L]) :- is-prime(A), next(A,A1), p-list(A1,B,L).

p-list(A,B,L) :- next(A,A1), p-list(A1,B,L).

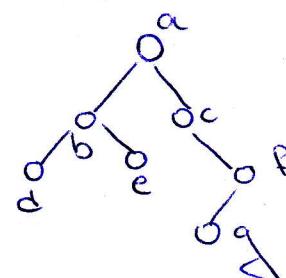
next(2,3).

next(A,A1) :- A1 is A+2.

34. -----

Binary tree

$t(x, L, R)$   
root  
left subtree      right subtree



$T1 = t(a, t(b, t(d, nil, nil), t(e, nil, nil)), t(c, nil, t(f, t(g, nil, nil), nil))),$

30 Check whether a given term represents a tree  
istree(nil).

istree( $t(-, L, R)$ ) :- is-tree(L), is-tree(R).

31 Construct completely balanced binary tree

cbal-tree([0, [J]]).

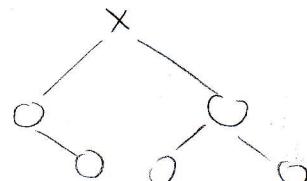
cbal-tree(N,  $t(x, L, R)$ ) :- N > 0, NO is N-1, NR is NO//2, N2 is NO-NO//2,  
distriB(NO1, NO2, NL, NR),  
cbal-tree(NO1, L), cbal-tree(NO2, R).

distriB(N, N, N, N).

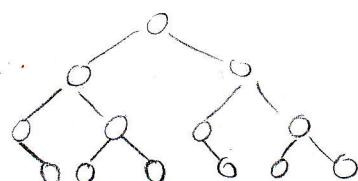
distriB(N1, N2, N1, N2) :- N1 = N2,

distriB(N1, N2, N2, N1) :- N1 ≠ N2.

cbal-tree(6,  $t(x, L, R)$ ).



cbal-tree(7,  $t(x, L, R)$ ).



32. Generate all completely balanced tree

P-ue

tree([ ], [ ]).

tree(N, [X, L, R]): - N > 0, N0 is N-1, N1 is N0//2, N2 is N0-N1  
distribute(N1, N2, NL, NR), tree(NL, L), tree(NR, R).

distribute(N, N, N, N).

distribute(N1, N2, X, Y): - N1 = N2,  
( (X is N1, Y is N2); (Y is N1, X is N2) ).

tree([X, L, R]): - nat(N), tree(N, [X, L, R]).

33 Let us call a binary tree symmetric if you can draw a vertical line through the node and then the right subtree is the mirror image of the left subtree.

Write a predicate to check if a given binary tree is symmetric.

symmetric([ ]).

symmetric([\_, L, R]): - mirror(L, R).

mirror([ ], [ ]).

mirror([T, L1, R1], [T, L2, R2]): - mirror(L1, L2), mirror(R1, R2).

34 Construct a BST from a list of integers

Bst([ ], [ ]).

Bst([H|T], T): - Bst([H|T]), insert(H, T, T).

insert(X, [ ], [X, [ ], [ ]]).

insert(X, [T, L, R], [T, LN, RN]): - X ≤ T, insert(X, L, LN).

insert(X, [T, L, R], [T, LN, RN]): - X > T, insert(X, R, RN).

35 Make a list from a BST.

ltr([ ], [ ]).

ltr([T, L, R], New): - ltr(L, LN), ltr(R, RN),  
append([H, [T|R]], New).

36. Generate all balanced symmetric trees.

syntree(T): - tree(T), symmetric(T).

- 37. Generate Height-Balanced tree

htree ([-, L, R]): - nat(N), htree (N, [-, L, R]).

htree (0, nil).

htree (N, [root, nil, nil]).

htree (N, [root, L, R]): - N > 1, N1 is N-1, N2 is N-2,  
distribute (N1, N2, NL, NR),  
htree (NL, L), htree (NR, R).

distribute (N, N, N).

distribute (N1, N2, X, Y) :- (X is N1, Y is N2); (X is N2, Y is N1).

38 Count the leaves of a binary tree.

count\_leaves (nil, 0).

count\_leaves ([-, nil, nil], 1).

count\_leaves ([-, L, R], N) :- count\_leaves (L, NL),  
count\_leaves (R, NR),  
N is NL+NR.

39 Collect the internal nodes of a tree in a list.

internals ([3, 3, 3], [3]).

internals ([X, 3, 3], [3]) :- X \= 3.

internals ([X, L, R], [X|S]) :- internals (L, XL), internals (R, XR),  
append (XL, XR, S).

40 Collect the leaves of a tree in a list.

leaves ([3, 3, 3], [3]).

leaves ([X, 3, 3], [X]) :- X \= 3.

leaves ([-, L, R], S) :- leaves (L, SL), leaves (R, SR), append (SL, SR, S).

41. Collect nodes at a given level of a binary tree.

atlevel ([3, 3, 3], 1, [3]).

atlevel ([X, 3, 3], 1, [X]).

atlevel ([-, L, R], L, S) :- D > 1, D1 is D-1, atlevel (L, D1, SL),  
atlevel (R, D1, SR), append (SL, SR, S).

42. Construct a complete binary tree.

cbt (N, T) :- cbt (N, T, 1).

cbt (N, [3, 3, 3], A) :- A > N.

cbt (N, [-, L, R], A) :- A =< N, AL is 2\*A, AR is AL+A,  
cbt (N, L, AL),  
cbt (N, R, AR).

## // Graphs

1) Methods to represent graphs in prolog.

1) Edges as facts ..

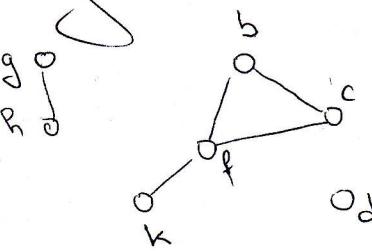
e(h,g)

e(k,f)

e(f,b).

---

- isolated nodes cannot be represented.



2)  $G = [V, E]$  where  $V = [b, c, d, f, g, h, k]$   
 $E = [[b, c], [c, d], [f, b], [k, f], [g, h]]$

3) {b-c, f-c, g-h, d, f-b, k-f, h-g}

4) Path in a graph // G - list of edges

pathl( , X, Y, , [X]).

pathl(G, X, Y, Visited, [X|R]): - X \= Y,  
member([X, Z], G),  
not(member([Z, Y], Visited)),  
pathl(G, Z, Y, [X|Visited], R)

5) Cycle in a graph

cycle(G, [X|T]): - member([X, Y], G), pathl(G, Y, X, [ ], T).