

# Logic(s) for computer science - Week 11

## Natural deduction

### Tutorial

## 1 Natural Deduction Rules

$$\begin{array}{c}
 \wedge i \frac{\Gamma \vdash \varphi \quad \Gamma \vdash \varphi'}{\Gamma \vdash (\varphi \wedge \varphi')}, \quad \wedge e_1 \frac{\Gamma \vdash (\varphi \wedge \varphi')}{\Gamma \vdash \varphi}, \quad \wedge e_2 \frac{\Gamma \vdash (\varphi \wedge \varphi')}{\Gamma \vdash \varphi'}, \\
 \rightarrow e \frac{\Gamma \vdash (\varphi \rightarrow \varphi') \quad \Gamma \vdash \varphi}{\Gamma \vdash \varphi'}, \quad \rightarrow i \frac{\Gamma, \varphi \vdash \varphi'}{\Gamma \vdash (\varphi \rightarrow \varphi')}, \quad \vee i_1 \frac{\Gamma \vdash \varphi_1}{\Gamma \vdash (\varphi_1 \vee \varphi_2)}, \\
 \vee i_2 \frac{\Gamma \vdash \varphi_2}{\Gamma \vdash (\varphi_1 \vee \varphi_2)}, \quad \vee e \frac{\Gamma \vdash (\varphi_1 \vee \varphi_2) \quad \Gamma, \varphi_1 \vdash \varphi' \quad \Gamma, \varphi_2 \vdash \varphi'}{\Gamma \vdash \varphi'}, \\
 \neg e \frac{\Gamma \vdash \varphi \quad \Gamma \vdash \neg \varphi}{\Gamma \vdash \perp}, \quad \neg i \frac{\Gamma, \varphi \vdash \perp}{\Gamma \vdash \neg \varphi}, \quad \perp e \frac{\Gamma \vdash \perp}{\Gamma \vdash \varphi}, \\
 \text{IPOTEZĂ} \frac{}{\Gamma \vdash \varphi} \varphi \in \Gamma, \quad \text{EXTINDERE} \frac{\Gamma \vdash \varphi}{\Gamma, \varphi' \vdash \varphi}, \quad \neg \neg e \frac{\Gamma \vdash \neg \neg \varphi}{\Gamma \vdash \varphi}, \quad \forall e \frac{\Gamma \vdash \forall x. \varphi}{\Gamma \vdash \varphi[x \mapsto t]} \\
 \exists e \frac{\Gamma \vdash \exists x. \varphi \quad \Gamma \cup \{\varphi[x \mapsto x_0]\} \vdash \psi}{\Gamma \vdash \psi} x_0 \notin \text{vars}(\Gamma, \varphi), \quad \forall i \frac{\Gamma \vdash \varphi[x \mapsto x_0]}{\Gamma \vdash \forall x. \varphi} x_0 \notin \text{vars}(\Gamma, \varphi) \\
 \exists i \frac{\Gamma \vdash \varphi[x \mapsto t]}{\Gamma \vdash \exists x. \varphi}
 \end{array}$$

## 2 Solved exercises

1. Prove that the sequence  $\{P(a), \neg P(a)\} \vdash P(b)$  is valid.

**Proof:**

1.  $\{P(a), \neg P(a)\} \vdash P(a);$  (HYPOTHESIS)
2.  $\{P(a), \neg P(a)\} \vdash \neg P(a);$  (HYPOTHESIS)
3.  $\{P(a), \neg P(a)\} \vdash \perp;$  ( $\neg e$ , 1, 2)
4.  $\{P(a), \neg P(a)\} \vdash P(b).$  ( $\perp e$ , 3)

2. Prove that the sequence  $\{(P(a) \vee Q(a))\} \vdash (q \vee p)$  is valid.

**Proof:**

1.  $\{(P(a) \vee Q(a)), P(a)\} \vdash P(a);$  (HYPOTHESIS)
2.  $\{(P(a) \vee Q(a)), P(a)\} \vdash (Q(a) \vee P(a));$  ( $\vee i_2, 1$ )
3.  $\{(P(a) \vee Q(a)), Q(a)\} \vdash Q(a);$  (HYPOTHESIS)
4.  $\{(P(a) \vee Q(a)), Q(a)\} \vdash (Q(a) \vee P(a));$  ( $\vee i_1, 1$ )
5.  $\{(P(a) \vee Q(a))\} \vdash (P(a) \vee Q(a));$  (HYPOTHESIS)
6.  $\{(P(a) \vee Q(a))\} \vdash (Q(a) \vee P(a)).$  ( $\vee e, 5, 2, 4$ )

3. Prove that the sequence  $\{\forall x.(P(x) \rightarrow Q(x)), P(a)\} \vdash \exists x.Q(x)$  is valid.

**Proof:**

1.  $\{\forall x.(P(x) \rightarrow Q(x)), P(a)\} \vdash \forall x.(P(x) \rightarrow Q(x))$  (HYPOTHESIS)
2.  $\{\forall x.(P(x) \rightarrow Q(x)), P(a)\} \vdash P(a)$  (HYPOTHESIS)
3.  $\{\forall x.(P(x) \rightarrow Q(x)), P(a)\} \vdash (P(a) \rightarrow Q(a))$  ( $\forall e, 1, a$ )
4.  $\{\forall x.(P(x) \rightarrow Q(x)), P(a)\} \vdash Q(a)$  ( $\rightarrow e, 3, 2$ )
5.  $\{\forall x.(P(x) \rightarrow Q(x)), P(a)\} \vdash \exists x.Q(x)$  ( $\exists i, 4$ )

4. Prove that the sequence  $\{\forall x.(P(x) \rightarrow Q(x)), \exists x.P(x)\} \vdash \exists x.Q(x)$  is valid.

**Proof:**

1.  $\{\forall x.(P(x) \rightarrow Q(x)), \exists x.P(x)\} \vdash \exists x.P(x)$  (HYPOTHESIS)
2.  $\{\forall x.(P(x) \rightarrow Q(x)), \exists x.P(x), P(x_0)\} \vdash P(x_0)$  (HYPOTHESIS)
3.  $\{\forall x.(P(x) \rightarrow Q(x)), \exists x.P(x), P(x_0)\} \vdash \forall x.(P(x) \rightarrow Q(x))$  (HYPOTHESIS)
4.  $\{\forall x.(P(x) \rightarrow Q(x)), \exists x.P(x), P(x_0)\} \vdash (P(x_0) \rightarrow Q(x_0))$  ( $\forall e, 3, x_0$ )
5.  $\{\forall x.(P(x) \rightarrow Q(x)), \exists x.P(x), P(x_0)\} \vdash Q(x_0)$  ( $\rightarrow e, 4, 2$ )
6.  $\{\forall x.(P(x) \rightarrow Q(x)), \exists x.P(x), P(x_0)\} \vdash \exists x.Q(x)$  ( $\exists i, 5$ )
7.  $\{\forall x.(P(x) \rightarrow Q(x)), \exists x.P(x)\} \vdash \exists x.Q(x)$  ( $\exists e, 1, 6$ )

5. Prove that the sequence  $\{\forall x.(P(x) \rightarrow Q(x)), P(x)\} \vdash \forall x.Q(x)$  is valid.

**Proof:**

1.  $\{\forall x.(P(x) \rightarrow Q(x)), \forall x.P(x)\} \vdash \forall x.(P(x) \rightarrow Q(x))$  (HYPOTHESIS)
2.  $\{\forall x.(P(x) \rightarrow Q(x)), \forall x.P(x)\} \vdash \forall x.P(x)$  (HYPOTHESIS)
3.  $\{\forall x.(P(x) \rightarrow Q(x)), \forall x.P(x)\} \vdash (P(x_0) \rightarrow Q(x_0))$  ( $\forall e, 1, x_0$ )
4.  $\{\forall x.(P(x) \rightarrow Q(x)), \forall x.P(x)\} \vdash P(x_0)$  ( $\forall e, 2, x_0$ )
5.  $\{\forall x.(P(x) \rightarrow Q(x)), \forall x.P(x)\} \vdash Q(x_0)$  ( $\rightarrow e, 3, 4$ )
6.  $\{\forall x.(P(x) \rightarrow Q(x)), \forall x.P(x)\} \vdash \forall x.Q(x)$  ( $\forall i, 5$ )

### 3 Proposed exercises

Are the following sequences valid?

1.  $\{((P(a) \wedge Q(a)) \wedge \forall x.P(x))\} \vdash (Q(a) \wedge \forall x.P(x));$
2.  $\{((P(a) \wedge Q(a)) \wedge \forall x.P(x)), \forall x.Q(x)\} \vdash (\forall x.Q(x) \wedge Q(a));$
3.  $\{((P(a) \wedge Q(a)) \wedge \forall x.P(x))\} \vdash (\forall x.P(x) \wedge (Q(a) \wedge P(a)));$
4.  $\{((P(a) \wedge Q(a)) \rightarrow \forall x.P(x)), P(a), Q(a)\} \vdash \forall x.P(x);$
5.  $\{(P(a) \rightarrow \forall x.P(x)), P(a), Q(a)\} \vdash (Q(a) \wedge \forall x.P(x));$
6.  $\{(P(a) \rightarrow P(b)), (Q(a) \rightarrow P(b))\} \vdash ((P(a) \vee Q(a)) \rightarrow P(b));$
7.  $\{\neg(P(a) \wedge Q(a))\} \vdash (\neg P(a) \vee \neg Q(a));$
8.  $\{\neg(\neg P(a) \vee \neg Q(a))\} \vdash (P(a) \wedge Q(a));$
9.  $\{\neg(\neg P(a) \wedge \neg Q(a))\} \vdash (P(a) \vee Q(a));$
10.  $\{\forall x.(P(x) \wedge Q(x))\} \vdash \forall x.P(x);$
11.  $\{\forall x.Q(x), P(a)\} \vdash P(a) \wedge Q(a);$
12.  $\{\forall x.P(x), \forall x.Q(x)\} \vdash \forall x.(P(x) \wedge Q(x));$
13.  $\{\exists x.\exists y.P(x, y)\} \vdash \exists y.\exists x.P(x, y);$
14.  $\{\exists x.\forall y.P(x, y)\} \vdash \forall y.\exists x.P(x, y);$  But in the other direction:  $\{\forall y.\exists x.P(x, y)\} \vdash \exists x.\forall y.P(x, y)?$
15.  $\{\neg(\exists x.P(x))\} \vdash \forall x.\neg P(x);$
16.  $\{\forall x.\neg P(x)\} \vdash \neg(\exists x.P(x));$