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ALGEBRAIC FOUNDATIONS of COMPUTER SCIENCE
lectures by Ferucio Laurentiu Tiplea

Rings and Fields

Part III – Applications in cryptography

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Outline

Advanced Encryption Standard (AES)

AES description

AES security

Reading and exercise guide

Advanced Encryption Standard (AES)

AES – a bit of hystory

1. On January 2, 1997, the American National Institute for Standards and Technology (NIST) invited cryptographers from all over the world to develop candidates for a new standard for the protection of sensitive electronic information;
2. Twenty-one teams of cryptographers from 11 countries submitted candidates;
3. On October 2, 2000, the winner was announced: the algorithm Rijndael (pronounced "Rhine-dahl"), designed by two Flemish researchers, Joan Daemen and Vincent Rijmen. Strong points:
   - a simple and elegant design;
   - efficient and fast on modern processors, but also compact in hardware and on smartcards;
4. On November 26, 2001, Rijndael was officially published as the Advanced Encryption Standard (AES).

AES description

AES description

AES processes data blocks of $4 \times m \times 8$ bits using a key of $4 \times k \times 8$ bits, where $m, k \in \{4, 6, 8\}$, as follows:

- first, the data block is divided into groups of 8 bits each (called bytes), obtaining in this way an array of bytes which is then organized as a $4 \times m$ matrix

$$\begin{pmatrix} b_0 & b_4 & \cdots & b_{4m-4} \\ b_1 & b_5 & \cdots & b_{4m-3} \\ b_2 & b_6 & \cdots & b_{4m-2} \\ b_3 & b_7 & \cdots & b_{4m-1} \end{pmatrix}$$

- the matrix obtained as above is considered as a plaintext symbol of the cryptosystem. It is encrypted by performing a set of transformations on it. The result is a $4 \times m$ matrix of bytes as well.

AES description

In Rijndael, bytes are represented in various ways:

- as sequences of 8 bits, or
- as 8-dimensional (row) vectors over $\mathbb{Z}_2$, or
- as sequences of two hexadecimal digits.

For example, the following notations refer to the same byte:

$$00111101, \quad (0,0,1,1,1,1,0,1), \quad (3d)_h$$

("$(\cdot)_h$" stands for the hexadecimal notation).

AES description

- $\mathcal{P} = \mathcal{C} = \mathcal{M}_{4\times m}(\mathbb{Z}_2^8)$, where $m \in \{4, 6, 8\}$;
- $\mathcal{K} = \mathcal{M}_{4\times k}(\mathbb{Z}_2^8)$, where $k \in \{4, 6, 8\}$;
- For any $K \in \mathcal{K}$,

$$e_K = T^f_{K_n} \circ T_{K_{n-1}} \circ \cdots \circ T_{K_1} \circ T^i_{K_0}$$

and

$$d_K = T^{-f}_{K_0} \circ T^{-1}_{K_1} \circ \cdots \circ T^{-1}_{K_{n-1}} \circ T^i_{K_n}$$

- n denotes the number of *rounds* to be performed during the execution of the algorithm. It is dependent on the key and block length

n	$m = 4$	$m = 6$	$m = 8$
$k = 4$	10	12	14
$k = 6$	12	12	14
$k = 8$	14	14	14

AES description

- T_Z^i, T_Z, and T_Z^f are transformations given by:
  - $T_Z^i = A_Z$,
  - $T_Z = A_Z \circ Mc \circ Sh \circ S$,
  - $T_Z^f = A_Z \circ Sh \circ S$,

  for any $Z \in \mathcal{M}_{4\times m}(\mathbb{Z}_2^8)$.

- A_Z, called the AddRoundKey transformation, is just a simple bitwise XOR operation extended to matrices. That is,

$$A_Z(X)(i,j) = X(i,j) \oplus Z(i,j),$$

  for any $X \in \mathcal{M}_{4\times m}(\mathbb{Z}_2^8)$, $0 \leq i \leq 3$ and $0 \leq j \leq m-1$. We simply write $A_Z(X) = X \oplus Z$;

AES description

- S, called the SubBytes transformation, is a non-linear byte substitution that operates independently on each byte of the input matrix. It uses a substitution table that can be computed by

$$S(X)(i,j)^t = M_1 \cdot X(i,j)' \oplus C,$$

where

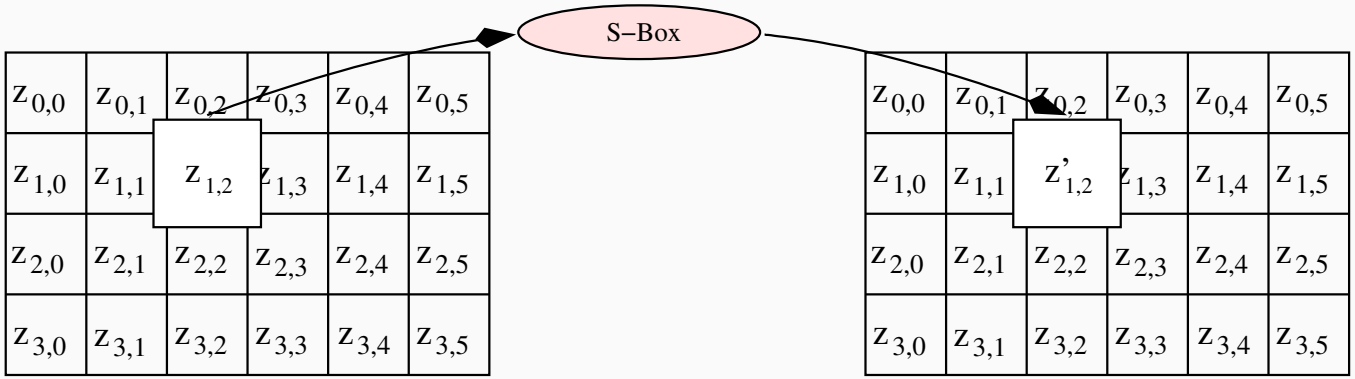
$$M_1 = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}, \quad C = \begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix},$$

AES description

and

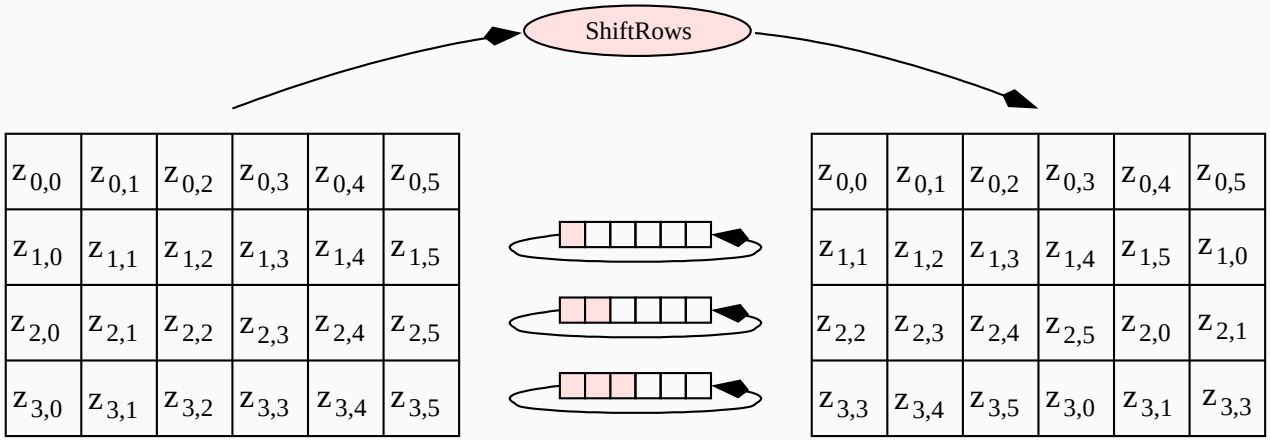
$$X(i,j)' = \begin{cases} (0,0,0,0,0,0,0,0)^t, & \text{if } X(i,j) = (00)_h \\ (X(i,j)^{-1})^t, & \text{otherwise} \end{cases}$$

(the inverse is in the finite field $GF(2^8)$ with the irreducible polynomial $x^8 + x^4 + x^3 + x + 1 \in \mathbb{Z}_2[x]$).



AES description

- *Sh*, called the ShiftRows transformation, cyclically shifts the rows of the input matrix over different numbers of positions (offsets). The *i*th row is shifted over $C_i = i$ positions (i=0,1,2,3)



Formally, we may write

$$Sh(X)(i,j) = X(i,(j + C_i) \ mod \ m),$$

for any $X \in \mathcal{M}_{4\times m}(\mathbb{Z}_2^8)$, $0 \leq i \leq 3$ and $0 \leq j \leq m-1$;

AES description

- Mc, called the MixColumns transformation, treats each column as a polynomial over $GF(2^8)$ and multiplies it modulo $x^4 + 1$ with a fixed polynomial $a(x)$ given by:

$$a(x) = (03)_h x^3 + (01)_h x^2 + (01)_h x + (02)_h.$$

This transformation can be written as a matrix multiplication in $GF(2^8)[x]$,

$$Mc(X) = M_2 \bullet X,$$

where

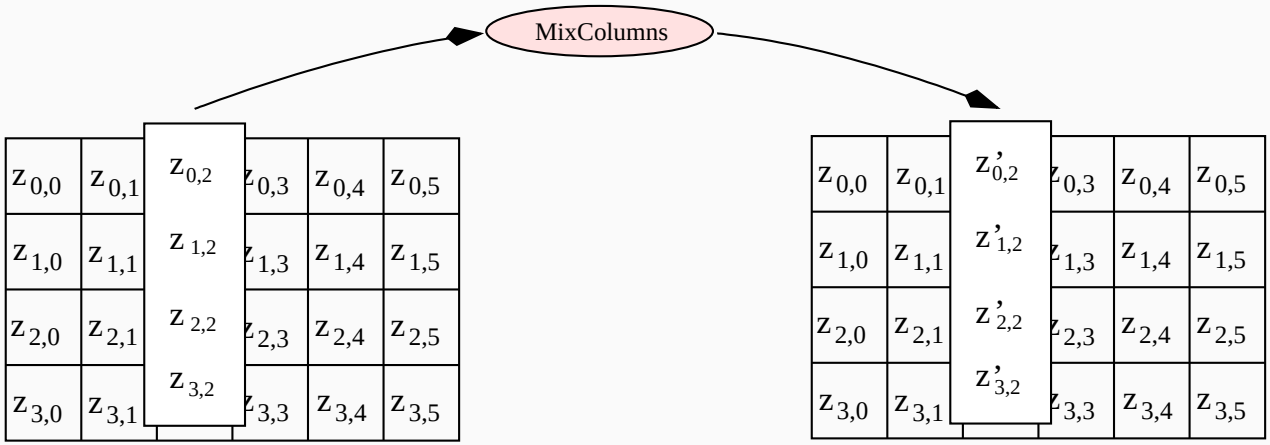
$$M_2 = \begin{pmatrix} (02)_h & (03)_h & (01)_h & (01)_h \\ (01)_h & (02)_h & (03)_h & (01)_h \\ (01)_h & (01)_h & (02)_h & (03)_h \\ (03)_h & (01)_h & (01)_h & (02)_h \end{pmatrix}$$

AES description

The matrix M_2 is invertible and its inverse is

$$M_2^{-1} = \begin{pmatrix} (0e)_h & (0b)_h & (0d)_h & (09)_h \\ (09)_h & (0e)_h & (0b)_h & (0d)_h \\ (0d)_h & (09)_h & (0e)_h & (0b)_h \\ (0b)_h & (0d)_h & (09)_h & (0e)_h \end{pmatrix}$$

The transformation is pictorially represented by



AES description

- T_Z^{-1} and T_Z^{-f} are transformations given by:
  - $T_Z^{-1} = A_{Mc^{-1}(Z)} \circ Mc^{-1} \circ Sh^{-1} \circ S^{-1}$
  - $T_Z^{-f} = A_Z \circ Sh^{-1} \circ S^{-1}$,

  for any $Z \in \mathcal{M}_{4\times m}(\mathbb{Z}_2^8)$.

  The transformations S, Sh, and Mc are invertible;

AES description

$K_0, \ldots, K_n$, called round keys, and obtained as follows:

- define first $W_0, W_1, \ldots, W_{m(n+1)-1}$ by
  - $W_i = K(-, i)$, for any $0 \leq i \leq k-1$;
  - $W_i = W_{i-k} \oplus T(W_{i-1})$, where:

$$T(W) = \begin{cases} SB(RB(W)) \oplus Rcon(i/k), & \text{if } i \ mod \ k = 0 \\ SB(W), & \text{if } k > 6 \text{ and } i \ mod \ k = 4 \\ W, & \text{otherwise} \end{cases}$$

$$RB((z_0, z_1, z_2, z_3)^t) = (z_1, z_2, z_3, z_0)^t$$

$$SB((z_0, z_1, z_0, z_3)^t) = (S(z_0), S(z_1), S(z_2), S(z_3))^t$$

$$Rcon(i) = (RC(i), (00)_h, (00)_h, (00)_h)$$

$$RC(1) = (01)_h \text{ and } RC(i) = x \bullet RC(i-1), \ \forall k \leq i \leq m(n+1)-1$$

AES security

AES security

AES Security:

- Asiacrypt 2002: Nicolas Courtois and Josef Pieprzyk showed that Rijndael can be written as an over-defined system of multivariate quadratic equations (MQ). For example authors showed that for 128-bit Rijndael, the problem of recovering the secret key from one single plaintext can be written as a system of 8000 quadratic equations with 1600 binary unknowns;
- If Shamir's XL algorithm would work for efficiently solving large systems of equations, then attacking Rijndael by such a method would require only a few known plaintexts to succeed;
- XL and XSL attacks do work in many interesting cases. Unfortunately they are heuristic, and their behavior is not well understood. There are examples where these or similar attacks do behave in practice as it is predicted, and there are examples where they do not.

Reading and exercise guide

Reading and exercise guide

Course readings:

1. Pages 343-349 from textbook [1].

References

[1] Ferucio Laurențiu Țiplea. *Algebraic Foundations of Computer Science*. "Alexandru Ioan Cuza" University Publishing House, Iași, Romania, second edition, 2021.