

Logic(s) for computer science - Week 13
Normal forms for First Order Logic - Part II
Tutorialr

1. Show that the following formulae in CSNF are unsatisfiable using binary resolution:

(a) $\forall x.\forall y.\forall z.\left(\left(\neg P(x, z) \vee R(x, x, z)\right) \wedge \left(\neg R(e, x, e)\right) \wedge \left(P(e, y)\right)\right);$

(b) $\forall x.\forall y.\left(\left(\neg P(x, y) \vee Q(x) \vee Q(y)\right) \wedge \left(\neg Q(i(i(e)))\right) \wedge \left(P(i(x), i(x))\right)\right).$

2. Establish using the binary resolution that the following formulae are valid:

(a) $\left(\left(\forall x.\forall y.\forall z(P(x, y) \wedge P(y, z) \rightarrow P(x, z))\right) \wedge P(x, y) \wedge P(y, x)\right) \rightarrow P(x, x);$

(b) $(\forall x.Q(x)) \rightarrow (\exists x.Q(x));$

(c) $(\neg\forall x.Q(x)) \leftrightarrow (\exists x.\neg Q(x));$

(d) $(\neg\exists x.Q(x)) \leftrightarrow (\forall x.\neg Q(x));$

(e) $(\exists y.\forall x.P(x, y)) \rightarrow (\forall x.\exists y.P(x, y));$

(f) $(\forall x.(P(x, x) \leftrightarrow Q(x))) \rightarrow (P(e, e) \rightarrow Q(e)).$