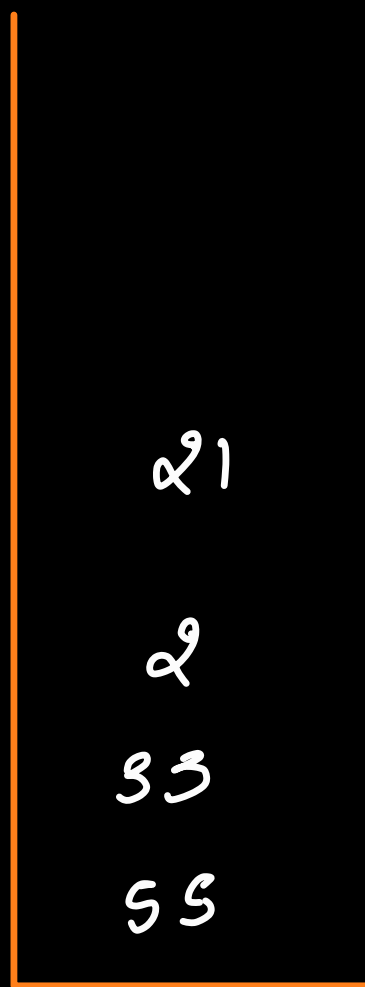


Stacks

- linear data structure
- data can be only accessed from one end of the stack
- apart from the topmost element no other element is directly accessible in the stack.
- application

memory → call stack

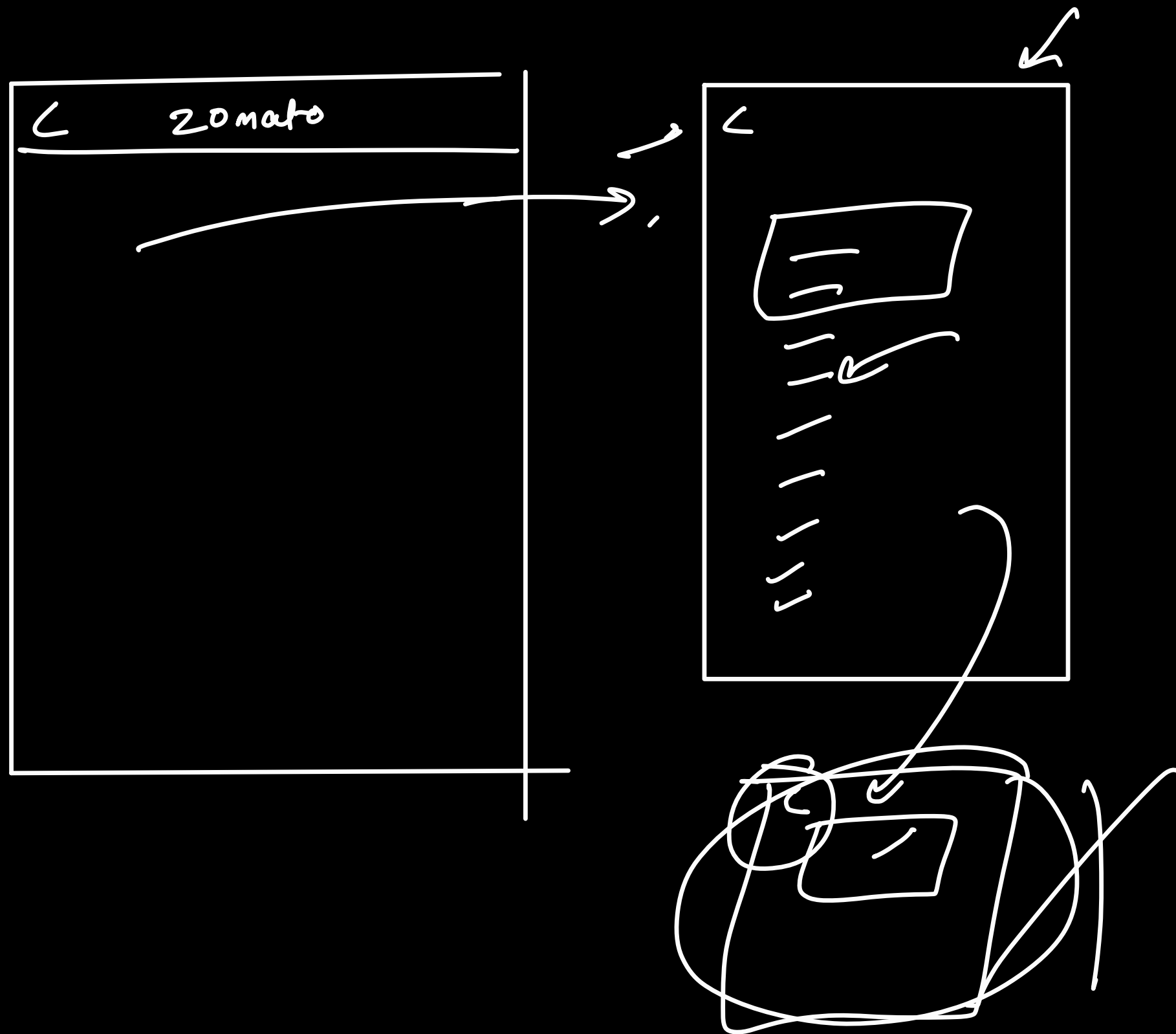
top →

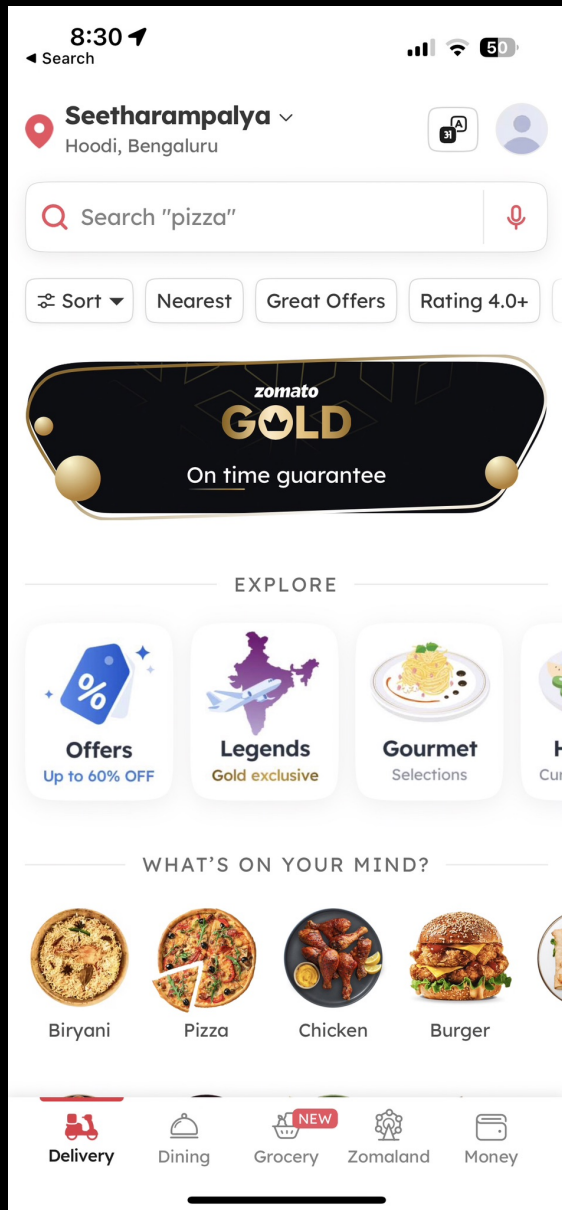


~~FI~~ LIFO

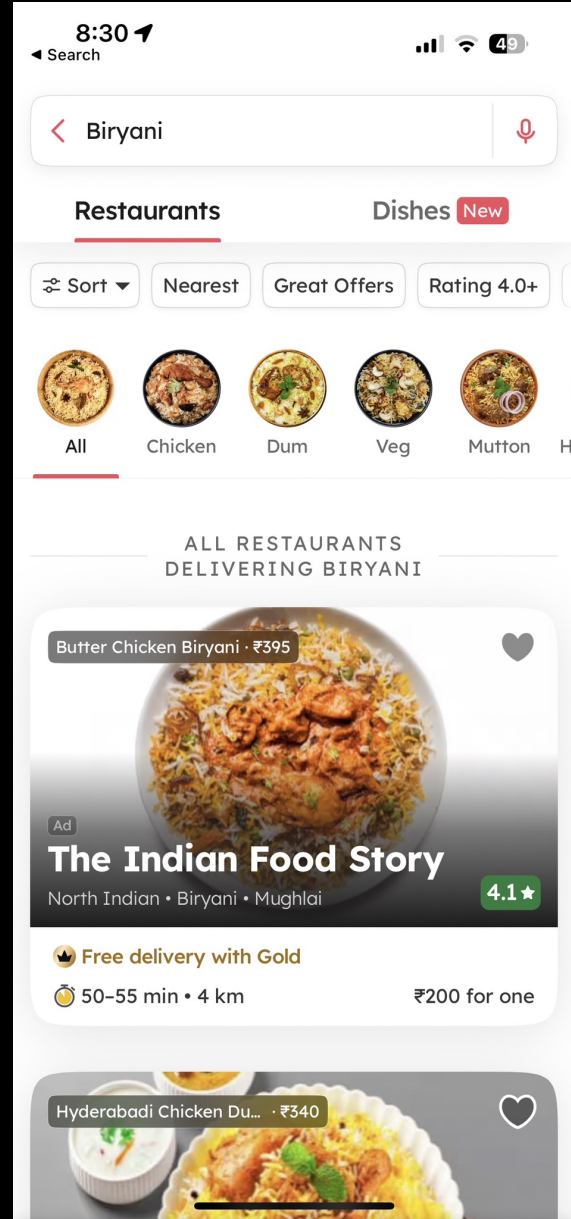
, , , , , 6

19

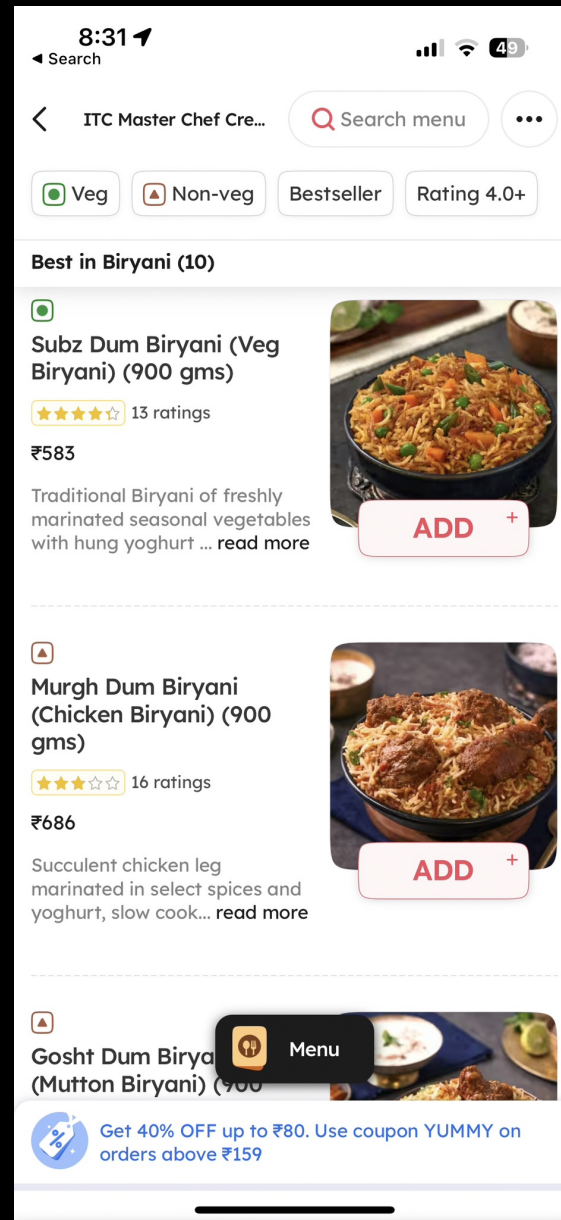




1

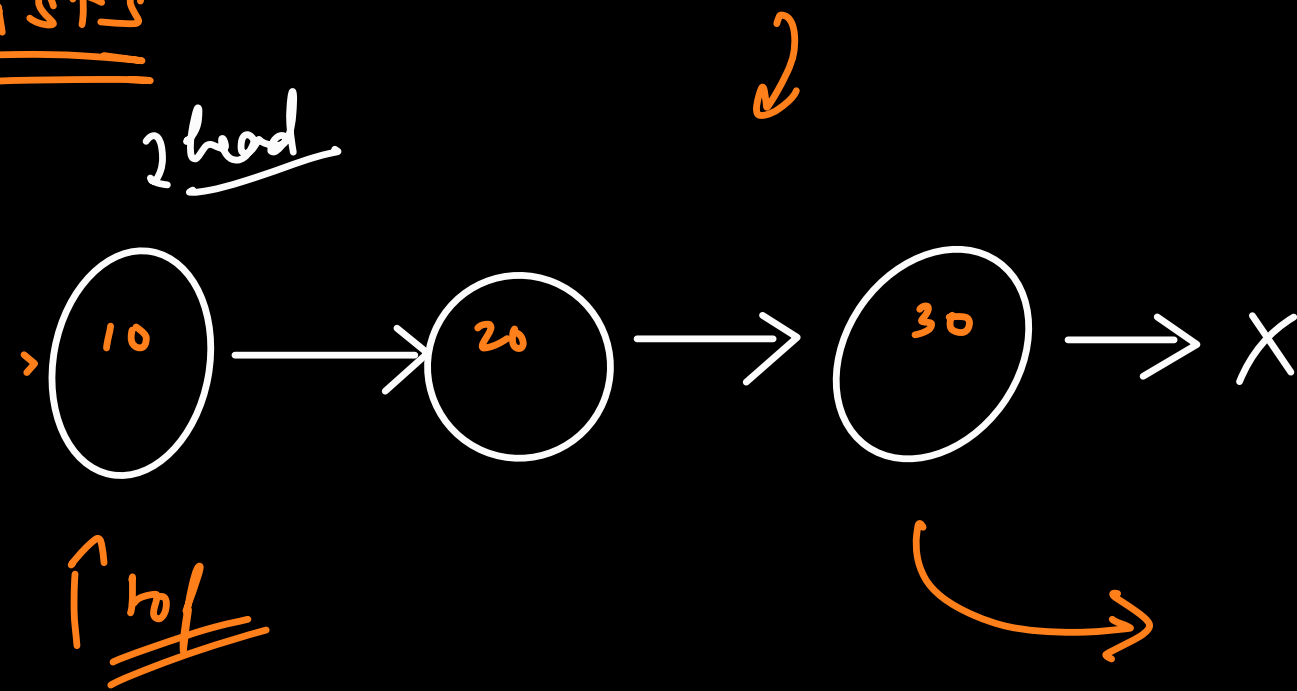


2



Implementation of Stacks \leftarrow add \rightarrow push

\rightarrow Linked Lists



whenever we do add At Head, head contains the last element that we have added.

push \rightarrow add \rightarrow add At Head \rightarrow O(1)
pop \rightarrow remove \rightarrow remove At Head \rightarrow O(1)

10
20
30

arrays

1	2	3			
---	---	---	--	--	--

9 men

4

st = []

{ st.push() }
st.pop()

() (x
) () x

JOIN THE DARKSIDE

{ [() [()] (())] }

$s_1 + s_2 \rightarrow \underline{s_1 s_2}$ ✓

$(s_1) \rightarrow \underline{\underline{\cdot}}$

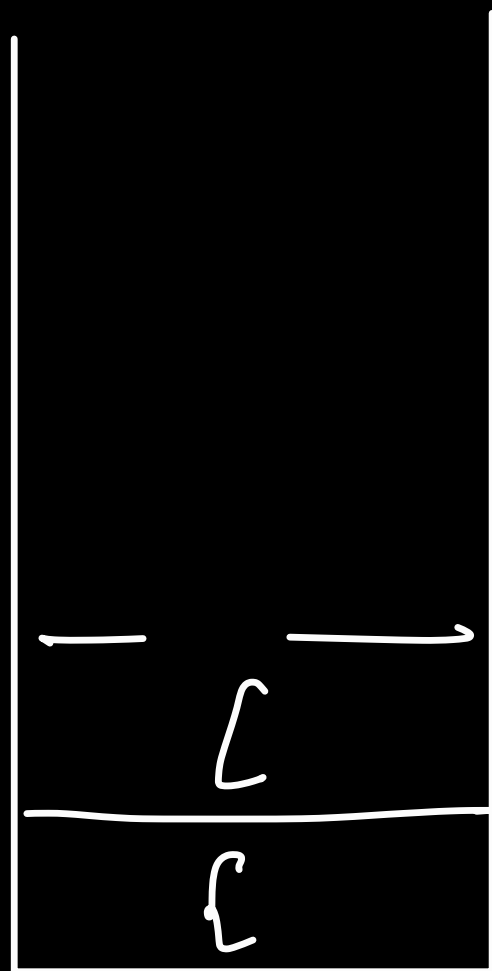
any opening bracket
↓
push it in the
stack.

{ [() [()] (())] }

closing bracket

$O(n)$

$\Delta p \rightarrow O(n)$

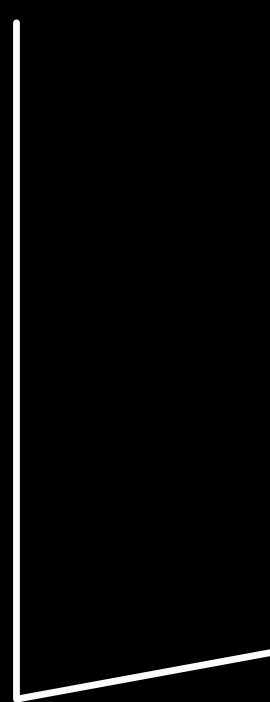
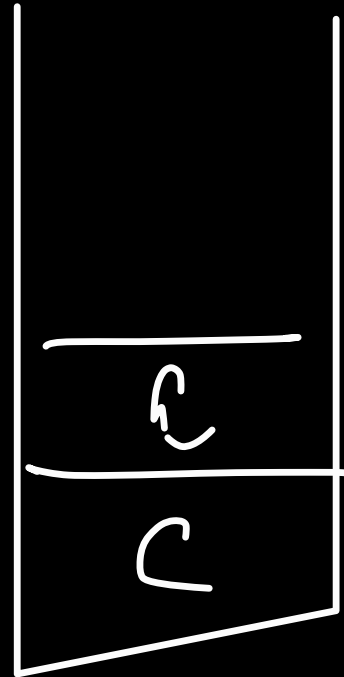


stack

→ at last stack is empty

balanced

$([])$ \rightarrow \neq \leftarrow $[]$



\sim \sim

((((((1- - - - -)))

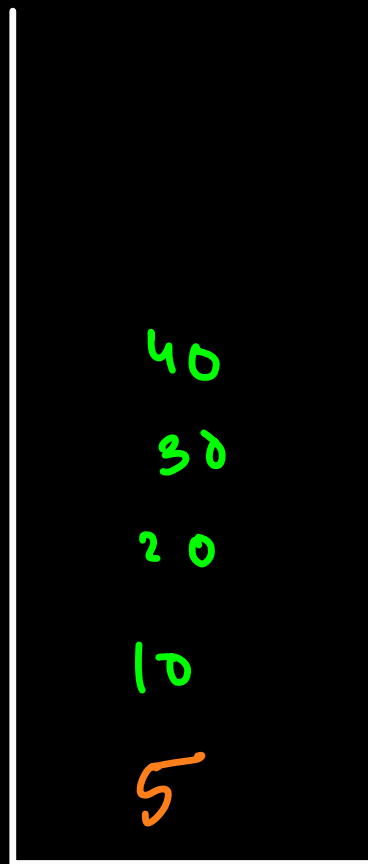
JOIN THE DARKSIDE

Q.2 Write a method that takes a stack as an input,
and an element x , and then instead of inserting
 x at top, it inserts x at bottom of stack.

(Time & Space need not be $O(1)$)

insertAtBottom (st, x)   push, pop, top

Time $\rightarrow O(n)$
Space $\rightarrow O(n)$



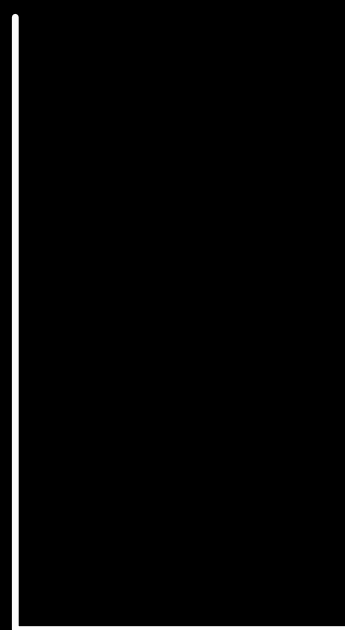
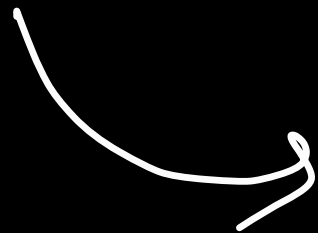
\rightarrow st

\rightarrow Remove all elements one by one
and back them up,

\rightarrow add x to the empty stack

\rightarrow insert the backed up elements
in the same order.

insertAtBottom(st, 5)



\rightarrow temp

insert At Bottom (st, x) {

temp ← temp stack

while (not st.empty()) {

temp.push(st.top())

st.pop()

}

st.push(x);

while (not temp.empty()) {

st.push(temp.top())

temp.pop()

}

return st;

}

st.empty() == 0

$O(n)$

$O(n)$

st[st.top()-1]

1 pseudo

→ st

→ reverse

→

insert At Bottom

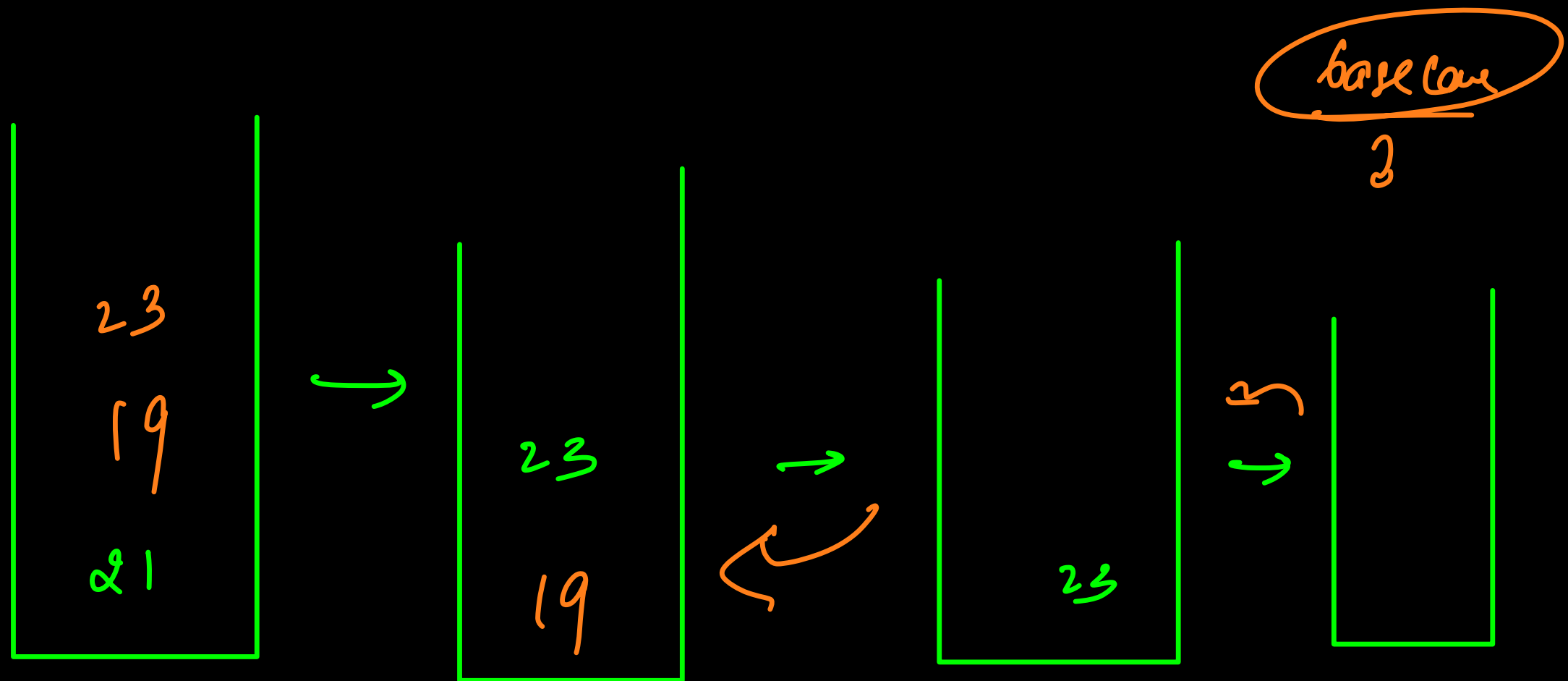
→ $O(N)$

Recursion

4
3
2
1

finally
→

1
2
3
4
5



19

↑ insert At Bottom → $O(n)$

for every element, we are calling insert At Bottom,

$O(N \times N) \rightarrow O(N^2)$

Space $\rightarrow O(n)$

```
reverse (st) {
```

```
    if (st.empty())  
        return;
```

```
    el = st.top
```

```
    st.pop()
```

```
    reverse (st);
```

```
    ↪ insertAtBottom (st, el);
```

```
}
```