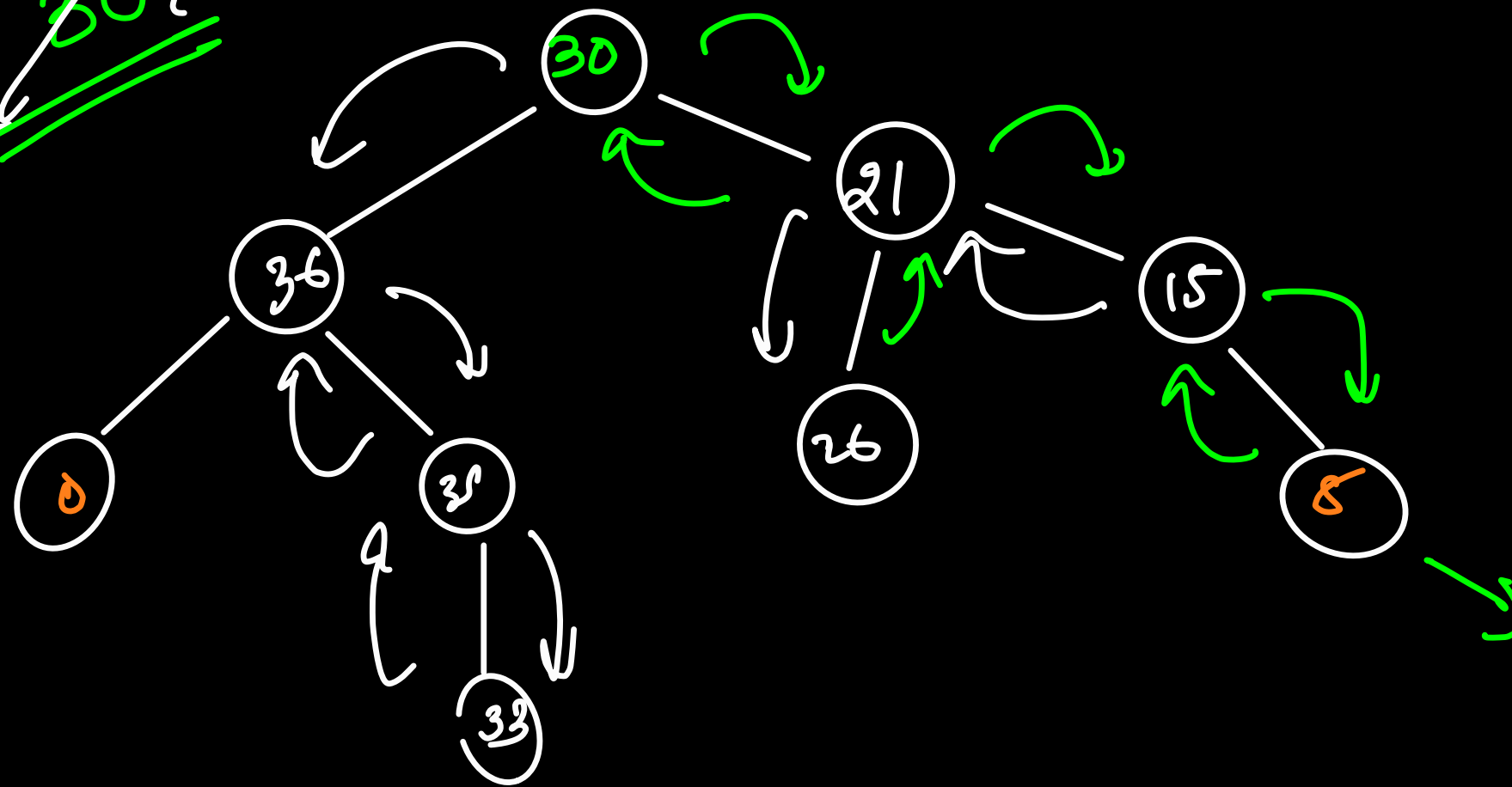


global sum = 30 ~~23~~ ~~35~~ -



Reverse In order

RST
Root
LST

$f(\text{root})$

↓

sum inc

=

$f(\text{root} \cdot \text{right})$

$\text{root.val} += \text{global sum}$

$\text{global sum} = \text{root.val}$

$f(\text{root} \cdot \text{left})$

LCA

$rs + ls + self$
 $0 + 0 + 1$

path from root

$O(N)$
 $O(h)$

$7 \leftarrow p, q \rightarrow 6$

path from root to p
 $\hookrightarrow \text{path}_p \rightarrow$

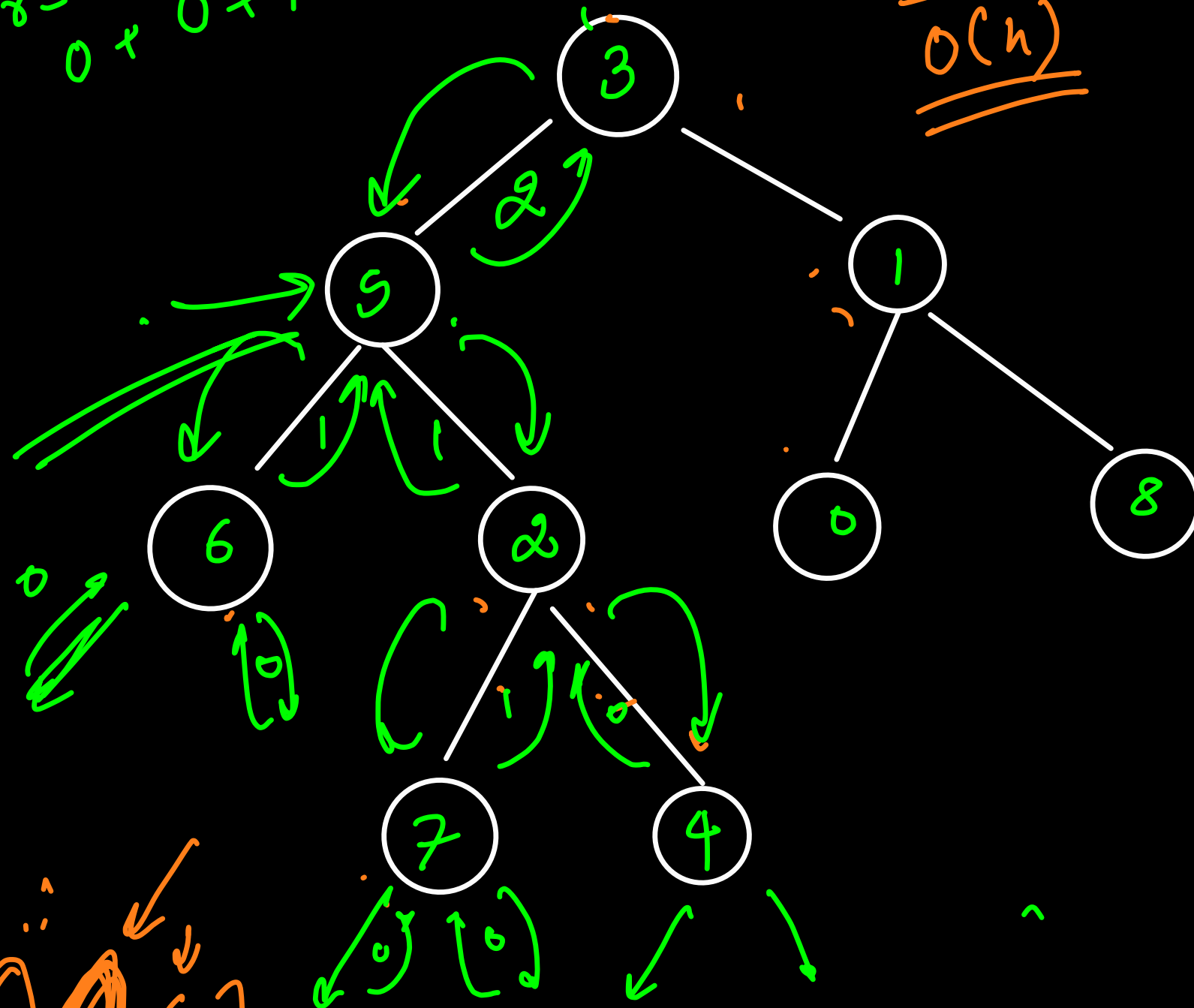
path_q =

path [7, 2, 5, 3]

(3, 5, 2, 7)

[0, 1, 5] \rightarrow (3, 1, 5)

[3, 5, 6]
 [3, 5, 2, 7]



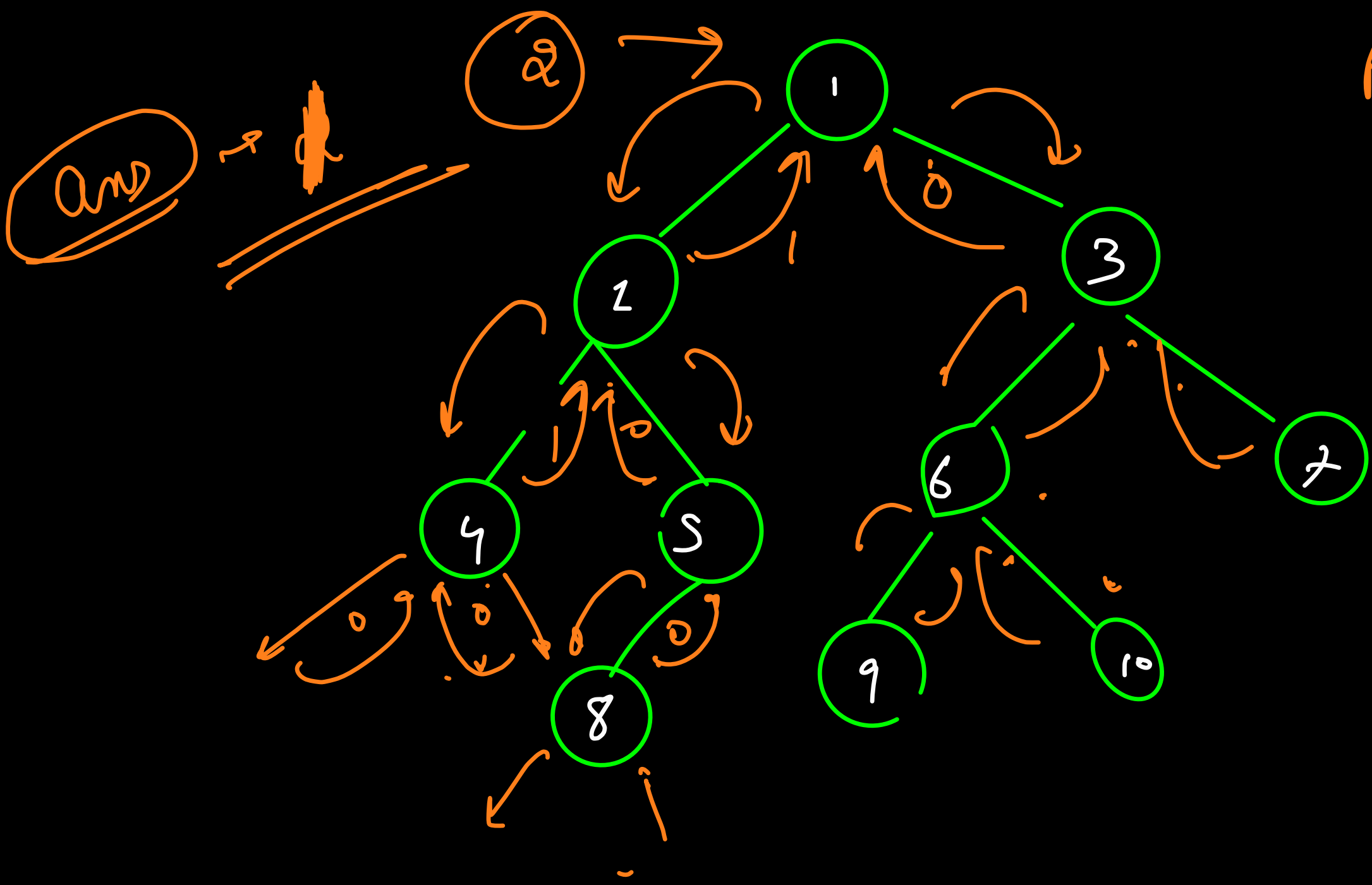
0
1
2
f (root, p, q)

This funⁿ returns 0
if none of p or q are present in
Subtree of root,
it returns 1 if any one of p or q
is present in subtree, & it
returns 2. if both of them are
present in root's subtree

$f(\text{root}, p)$ =
↓
boolean if a
path is found
from root to p

```
if (root.val == p) {  
    path.push(root.val);  
    return true;  
}  
  
lans = f(root.left, p)  
if (lans == true) {  
    path.push(root.val)  
    return true;  
}  
  
rans = f(root.right, p)  
if (rans == true)  
    path.push(root.val)  
    return true;  
}  
  
return false;
```

$p = 1, q = 4$



ans = -1

$f(\text{root}, p, q) \Rightarrow$

if (ans != -1) return 0

lans = $f(\text{root}.left, p, q)$

rans = $f(\text{root}.right, p, q)$

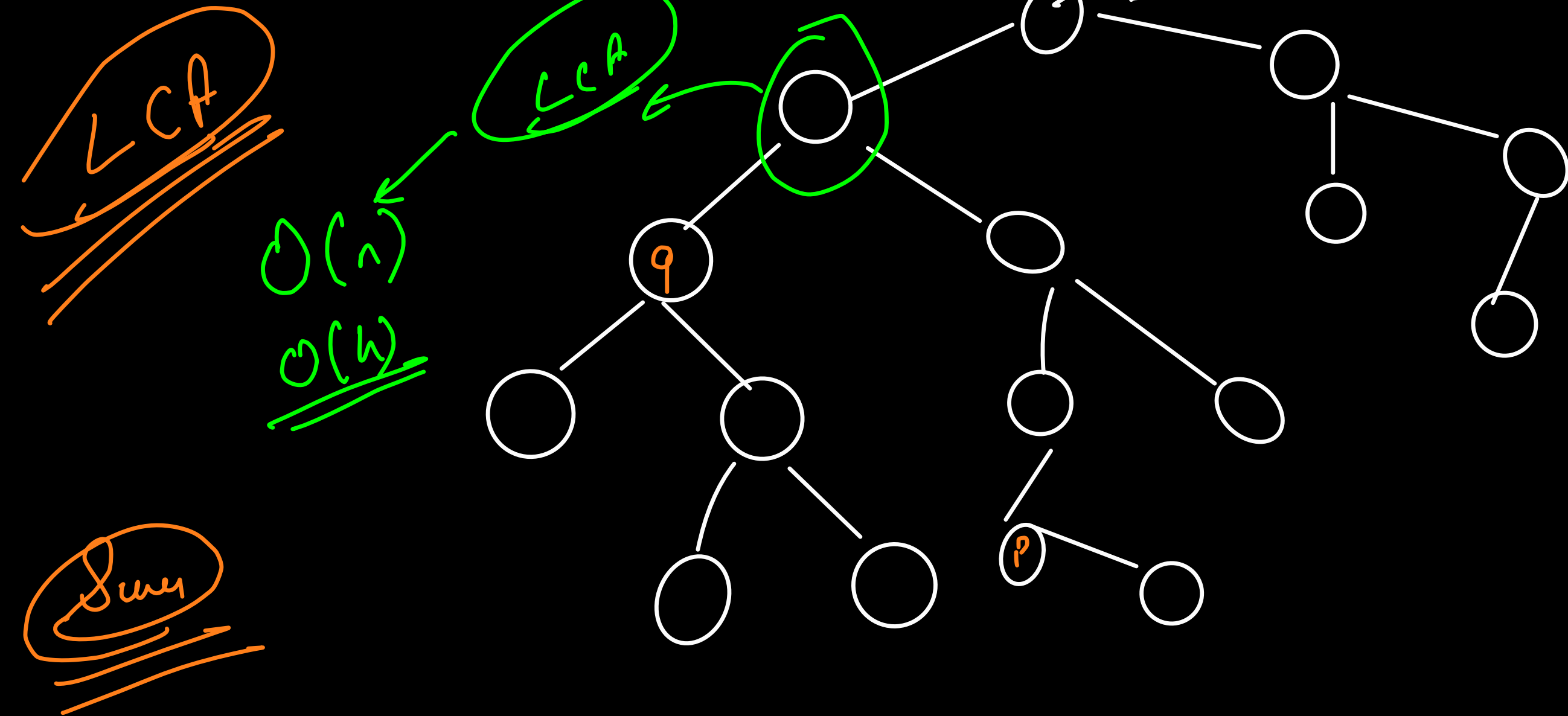
myans = ($\text{root}.val == p$ or $\text{root}.val == q$)

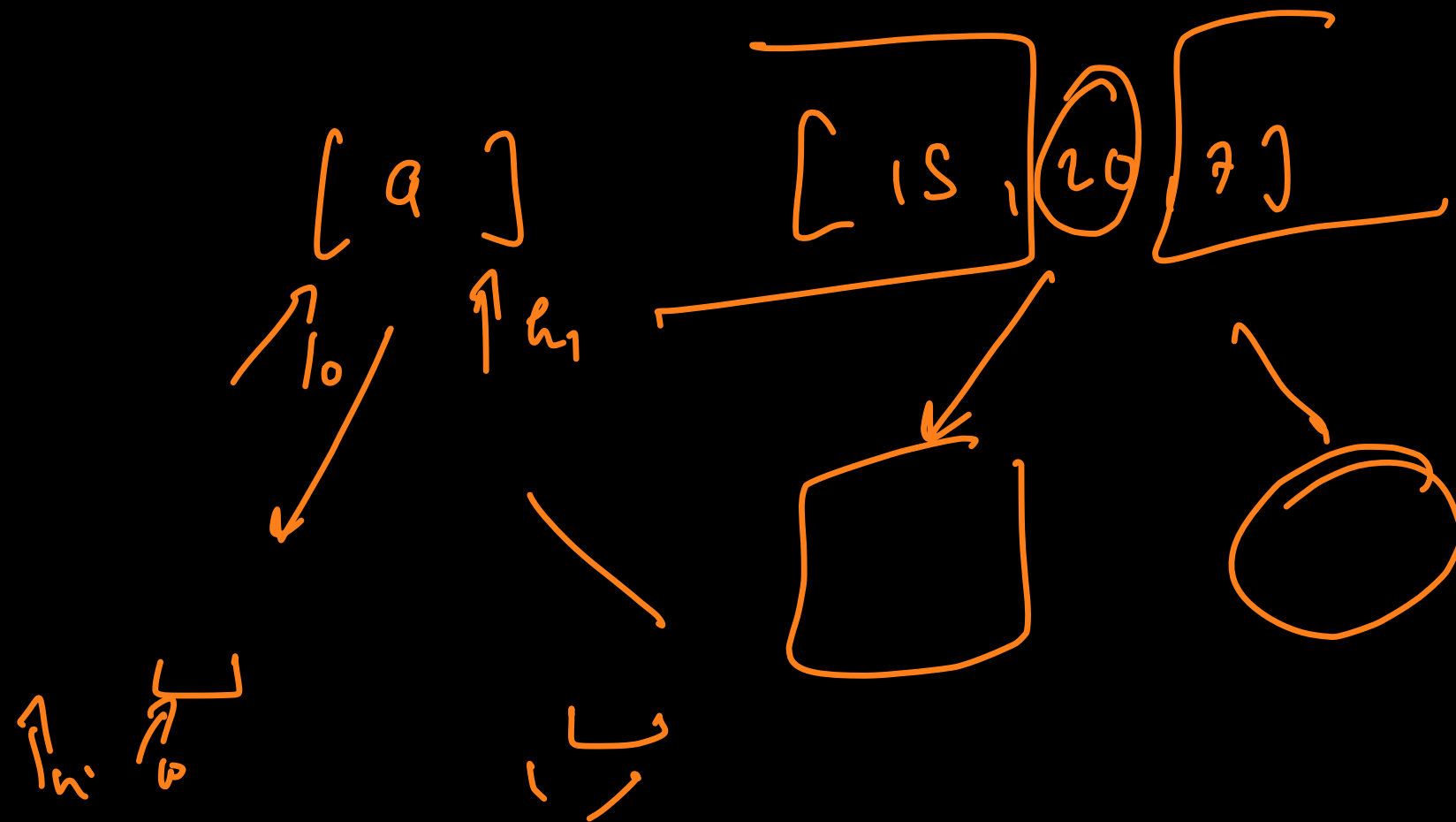
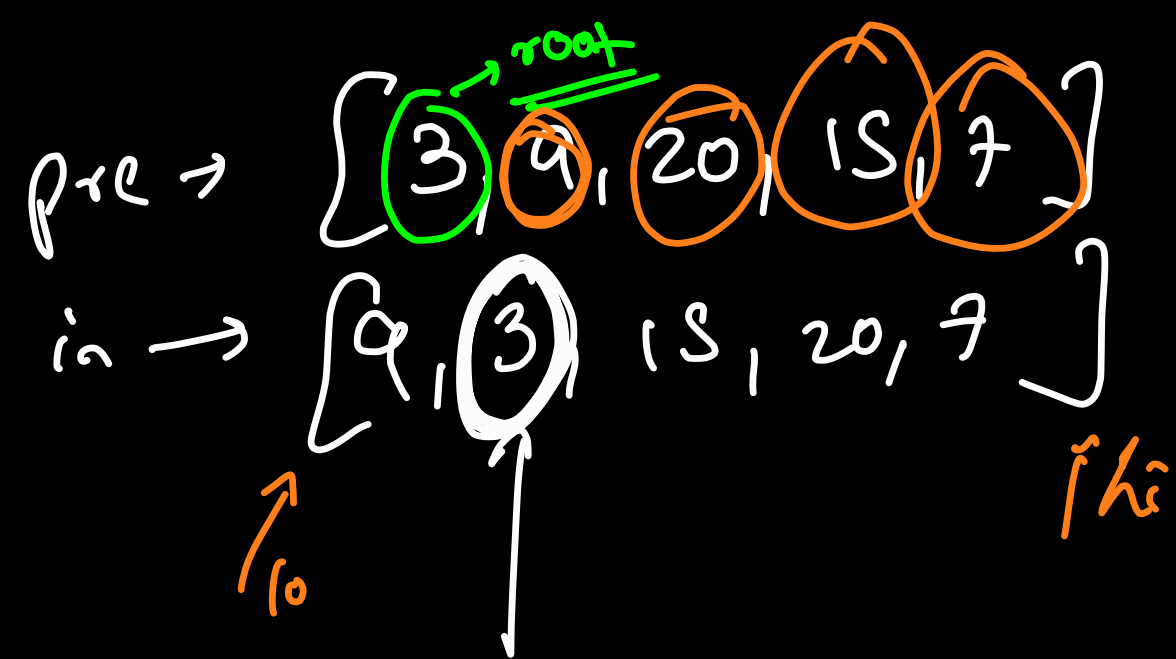
if ($\text{lans} + \text{rans} + \text{myans} > 0$)

ans = $\text{root}.val$

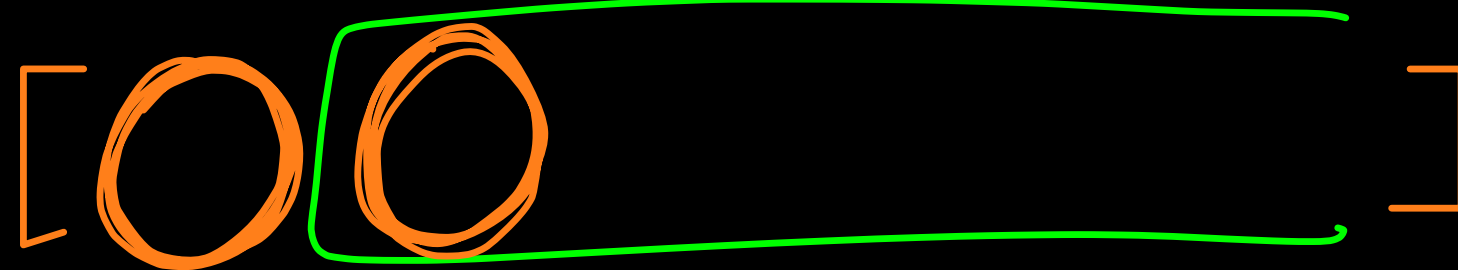
return $\text{lans} + \text{rans} + \text{myans}$;

Qⁿ Given a B.T, and 2 nodes p and q, we need to find the shortest path length between the 2 nodes. p and q.

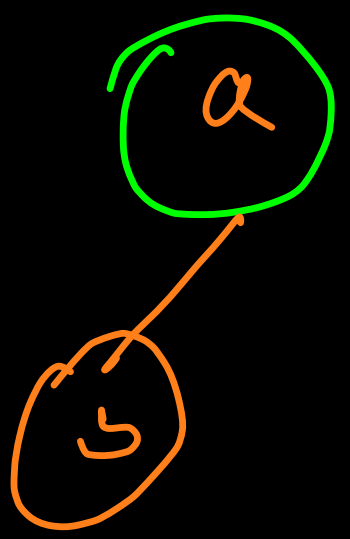




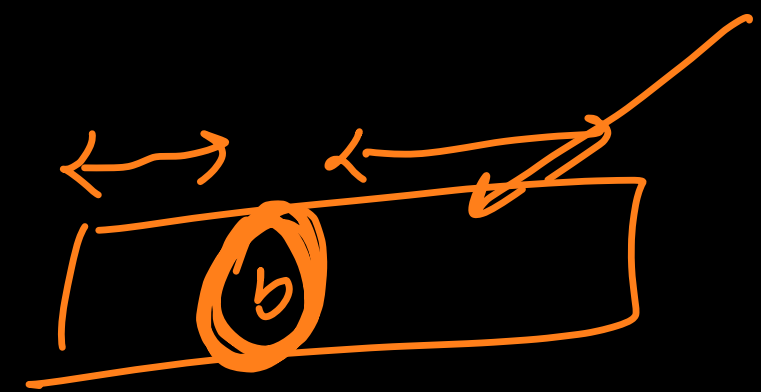
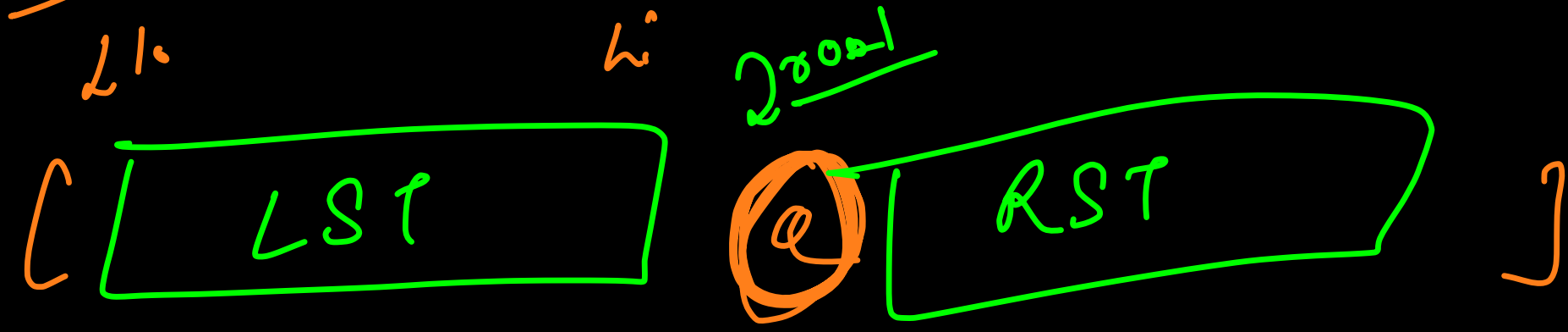
pre →



root
L1



in :



$O(n^2)$

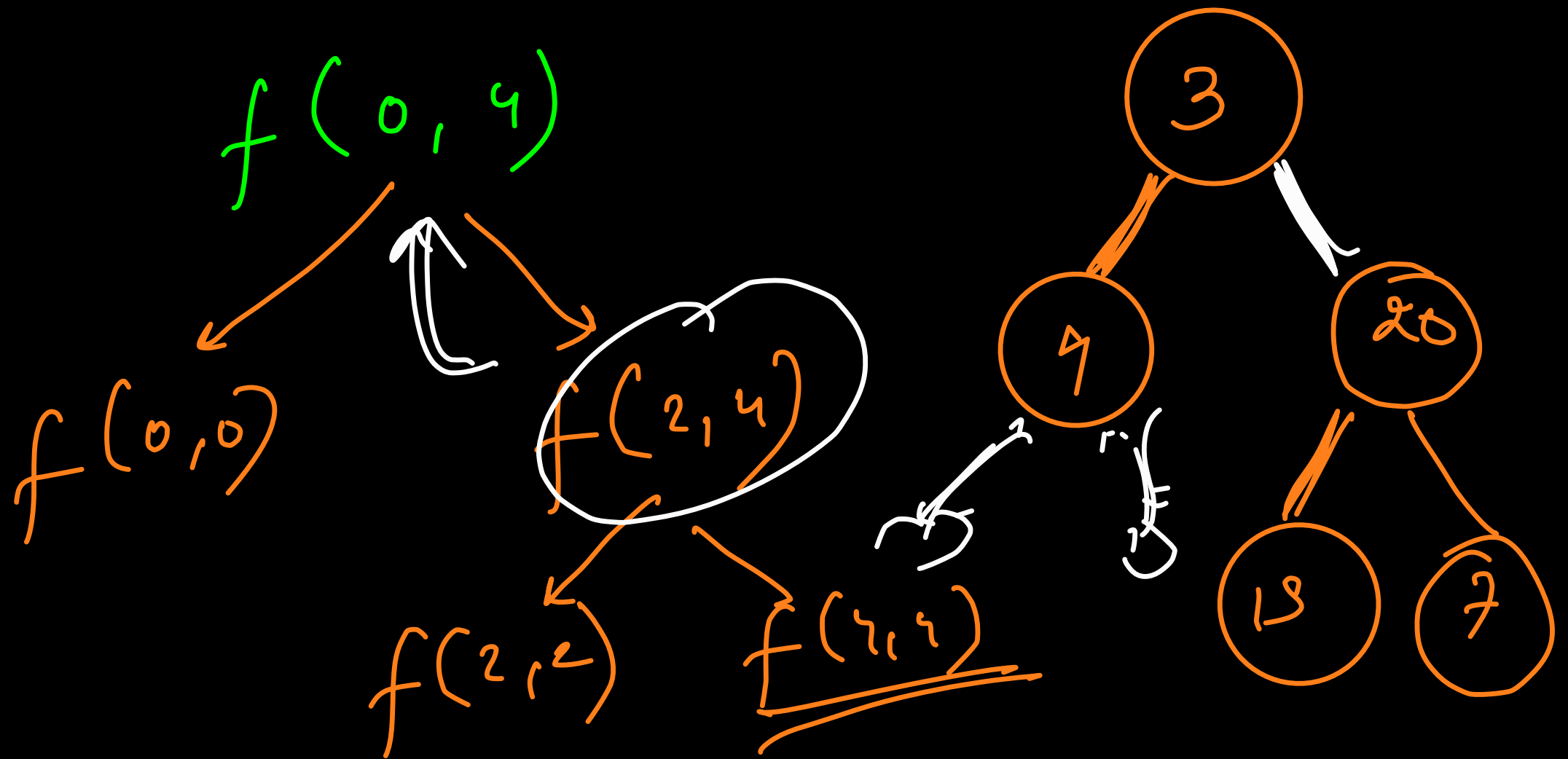
Optimal
search
el \rightarrow ind

9 \rightarrow 0
3 \rightarrow 1
15 \rightarrow 2
20 \rightarrow 3
7 \rightarrow 4

Mapper

pre \rightarrow [3 9 20 15 7]
in \rightarrow [9 3 15 20 7]

if (i > 0)
return



$f(\text{root}, l, r) \leftarrow$
 if ($l == r$)
 return new Node(pre[l])

DFS

$i++;$
 $n = \text{new Node}(\text{pre}[i])$

$m = \text{mp}[\text{pre}[i]]$

$n.\text{left} = f(\text{root}, l, m-1);$

$n.\text{right} = f(\text{root}, m+1, r);$

return root

}