

\hookrightarrow $bar[i] \rightarrow h_i \rightarrow$ we can form the biggest rectangle which covers the whole bar i^{th} bar, by using the first bar on left & right which is smaller than my current bar.

	0	1	2	3	4	5
	2	1	5	6	2	3
l si →	-1	-1	1	2	1	4
rs i →	1	6	4	4	6	6

hi(s-l-1)

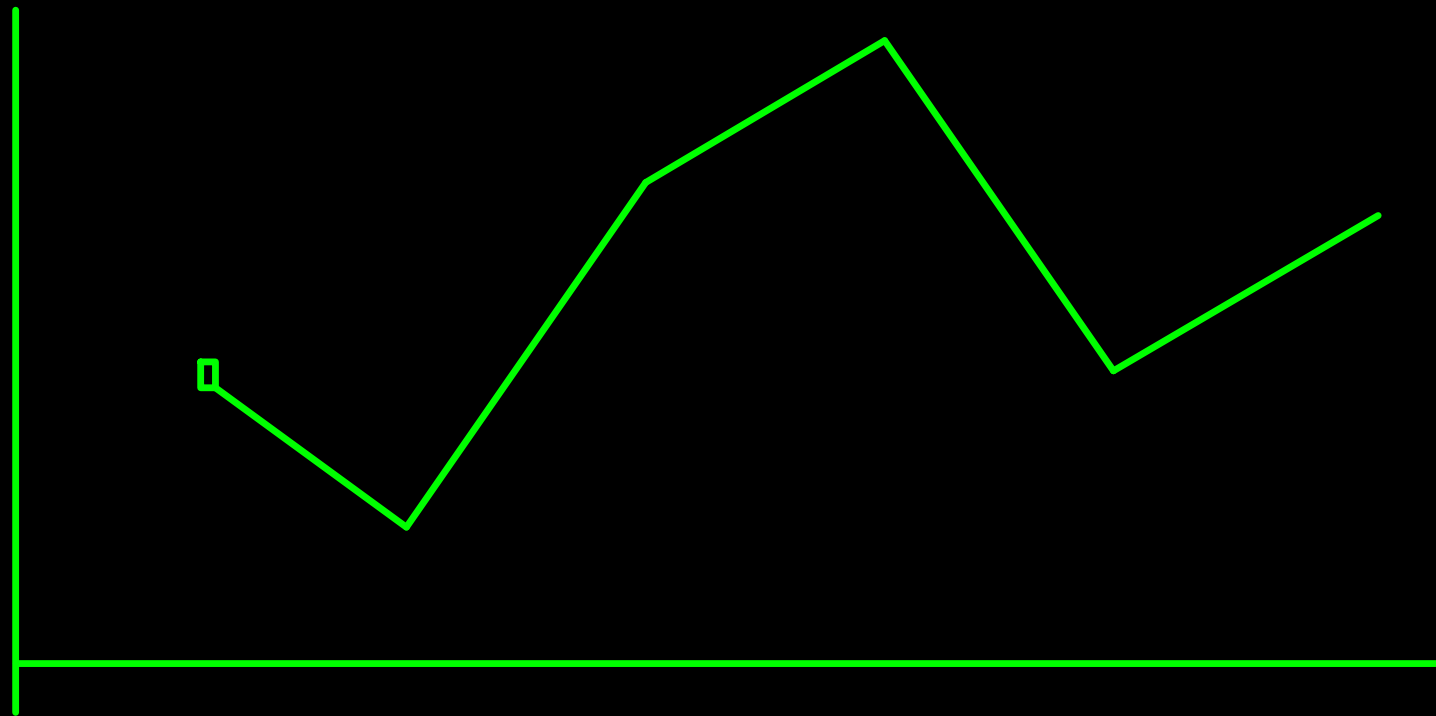
$O(n)$

$O(n)$

max (2 6 10 6 8 3) → 10

RSI

RS1 →



monitors Stack

4

1

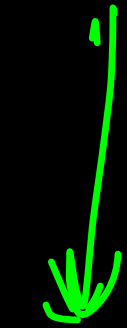
in order
client

$$u^{14} \rightarrow \begin{matrix} 10i \\ 75i \end{matrix}$$
$$x = st \cdot \log()$$

$$st \cdot \log()$$
$$\text{sf.yop}()$$
$$\text{arr}[a] \times (i - (\text{st.top}() - 1))$$
$$\gamma = 1$$
$$l = -1$$
$$2 \times (8 - 1 - 1)$$
$$2 \times (1 - \frac{1}{2}) = 1$$

ans = 610

$$arr[x] * (i - (s.empty() ? (-1) : s.top()) - 1)$$



remove elem.

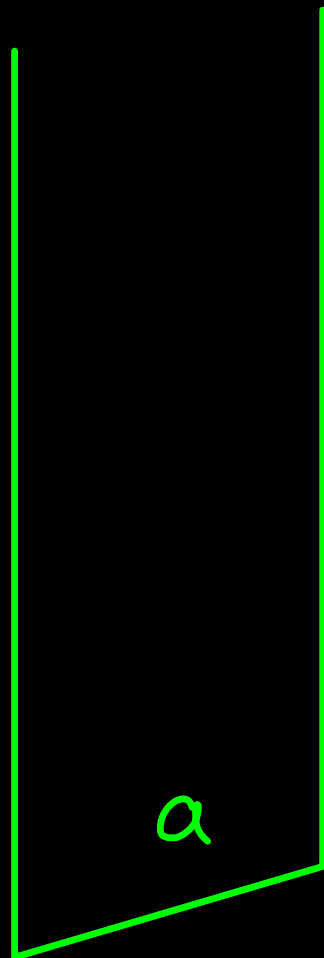
$$\rightarrow arr[x] * (arr.length - (s.empty() ? (-1) : s.top()) - 1)$$

1	0	1	0	0
1	0	1	1	1
1	1	1	1	1
1	0	0	1	0

1	0	1	0	0
2	0	2	1	1
3	1	3	2	2
4	0	0	3	0

$$\underline{\underline{\text{ans} = 4/6}}$$

$$\text{Space} \rightarrow \underline{\underline{O(n^2)}}$$



c_i c_i
 duplication

$$\frac{O(n)}{O(n)}$$

c c

b c c
 b b

a b b b b a

$[() ()]$