

level 0

level 1

level 2

level 3

level 4

1

2

3

4

5

6

7

8

9

10

11

12

12

level order traversal

FCFS

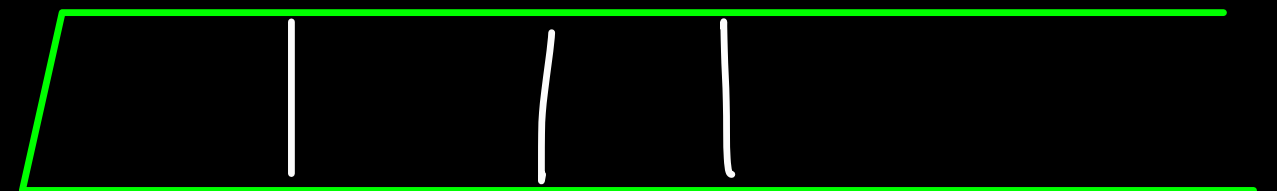
BFS

1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12

assume, (2) will be read
before (3) i.e left child
will be read before right
child

front

rear



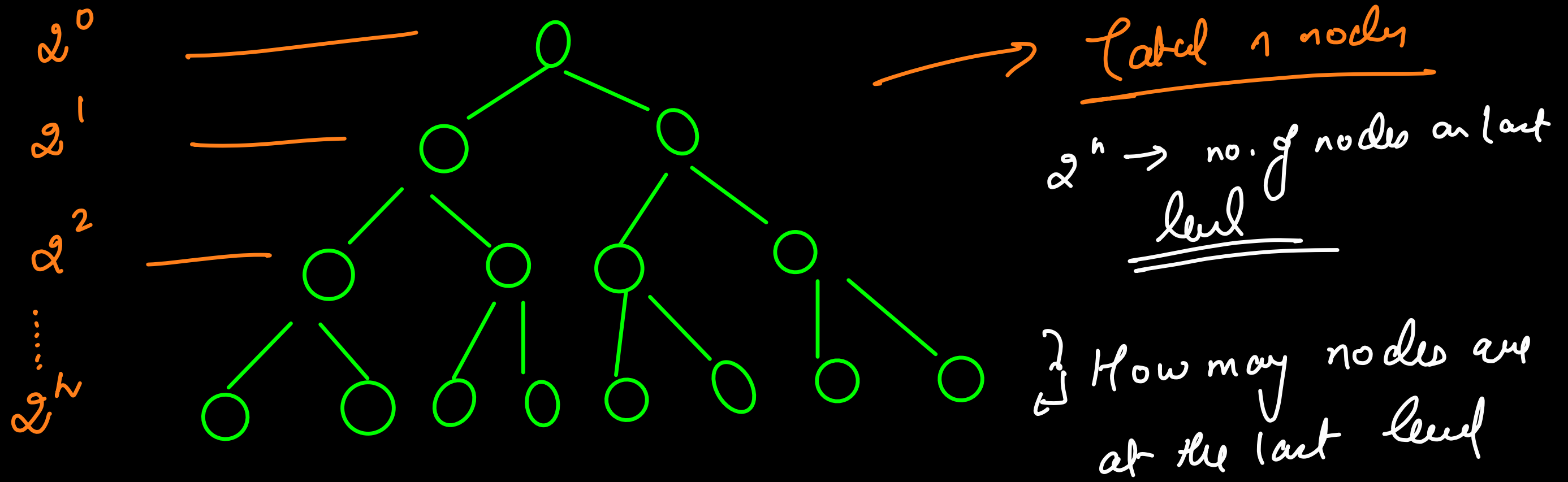
1 2 3 4 5 6 7
8 9 10 11 12

$O(n)$

$\rightarrow O(n)$ Space

```
qv.enqueue(root);  
while (not qv.empty()) {  
    front = qv.getFront();  
    qv.dequeue();  
    if (front.left != null)  
        qv.enqueue(front.left);  
    if (front.right != null)  
        qv.enqueue(front.right);  
    print(front.val);  
}
```





total no. of nodes in a perfect bt

$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{h-2} + 2^{h-1} + 2^h = n$$

gp

$$\frac{2^0 \times (2^h - 1)}{2 - 1} + 2^h = n$$

$$2^h - 1 + 2^h = n \Rightarrow 2 \times 2^h = n + 1$$

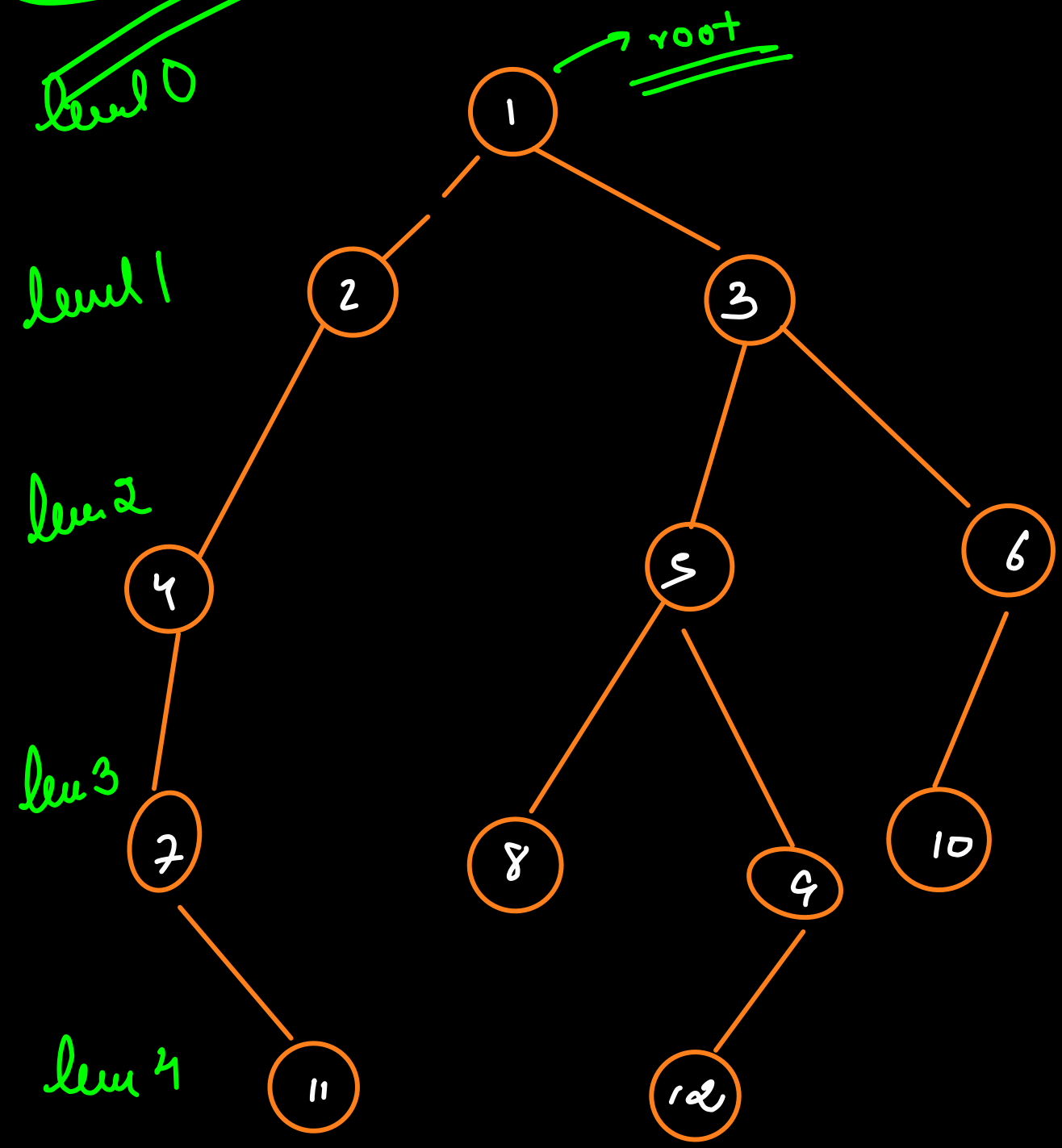
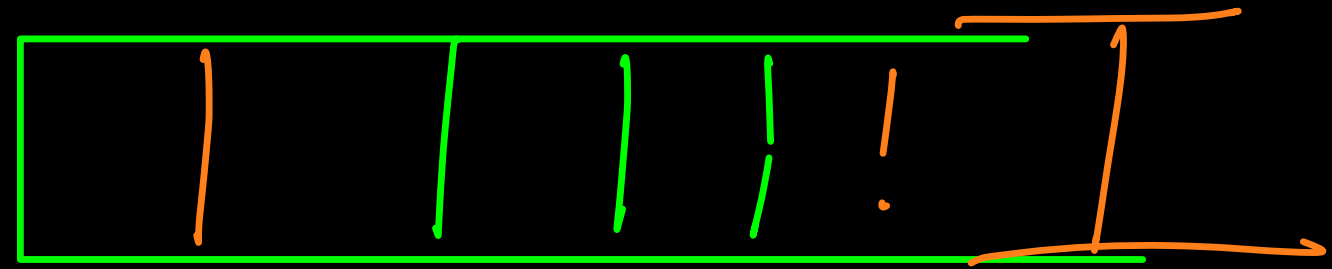
$$2^h = \frac{(n+1)}{2}$$

approx
leaf

level order
level wise

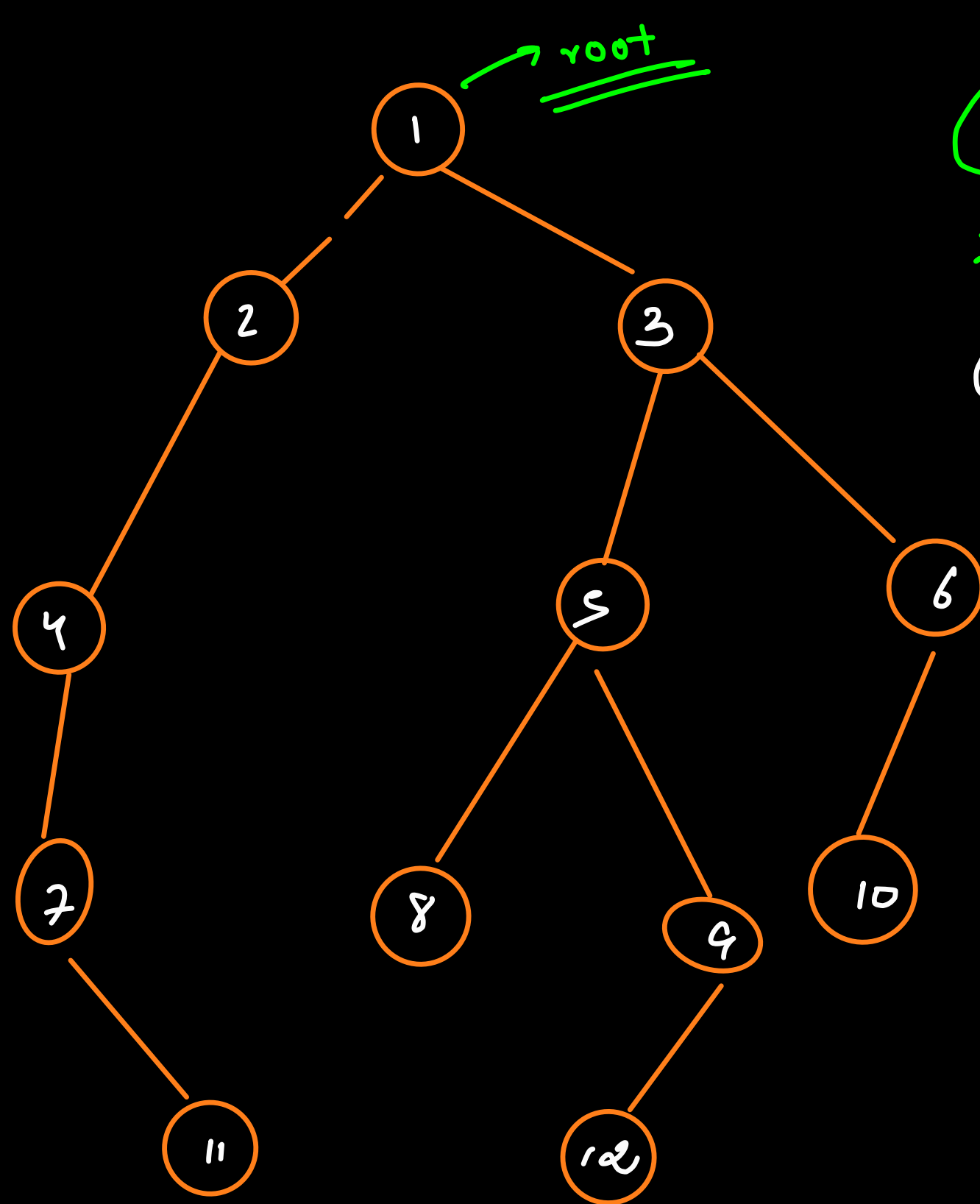
marker to represent end of a level

→ 94



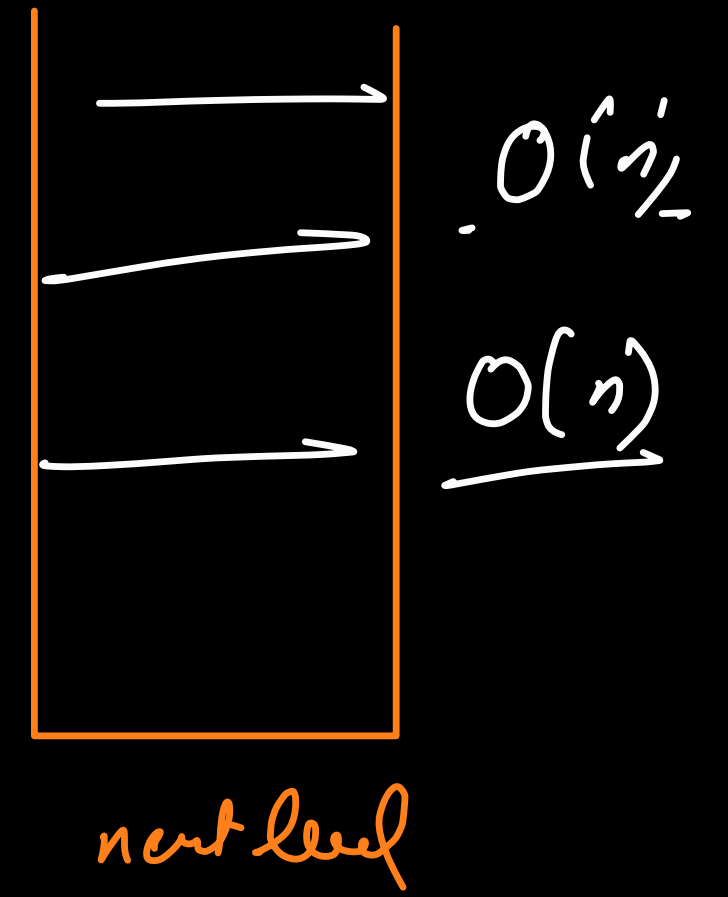
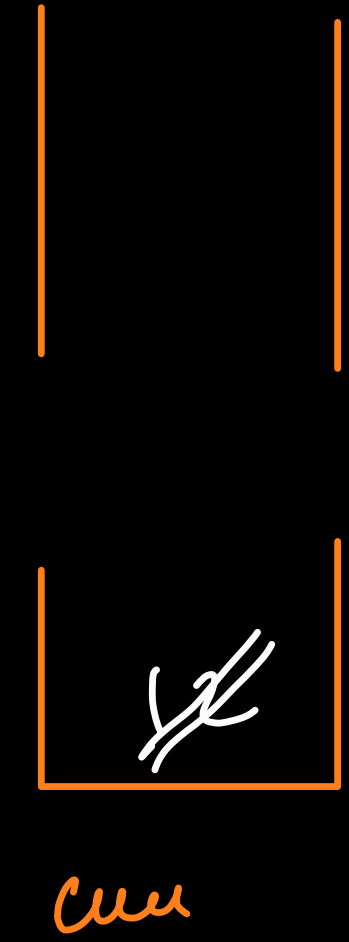
1 null 2 3 null 4 5 6 null
7 8 9 10 null 11 12 null

everytime when you remove
as null, we only have
elements of next level.



root

zigzag



cur - 1 2 3 4 5

right → left
 & then left → right

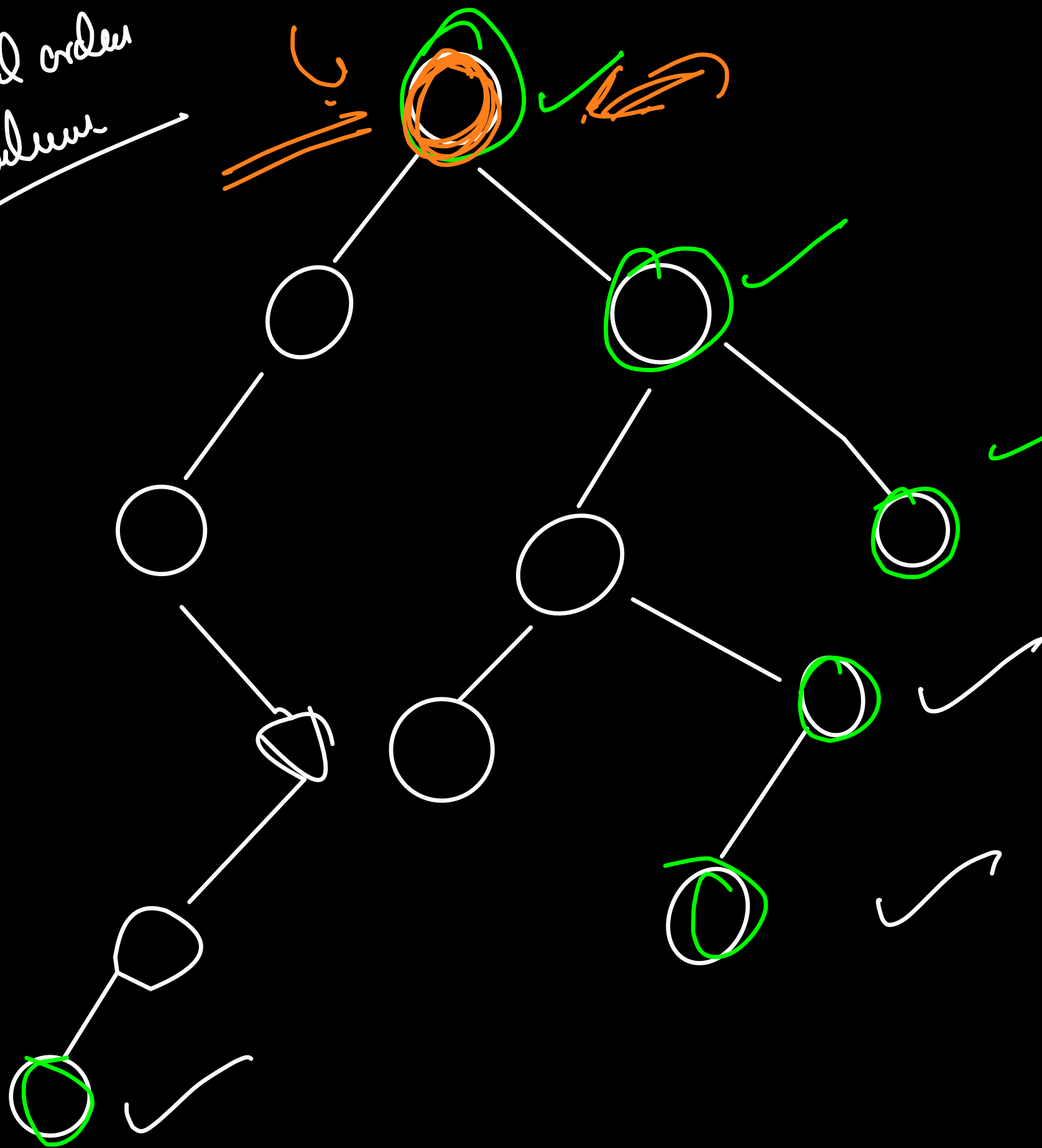
left → cur
 & then cur → 1

① ③ ② 4 ⑤ ⑥ 10 9 8 7

11 12

→

level order
level



Lead order

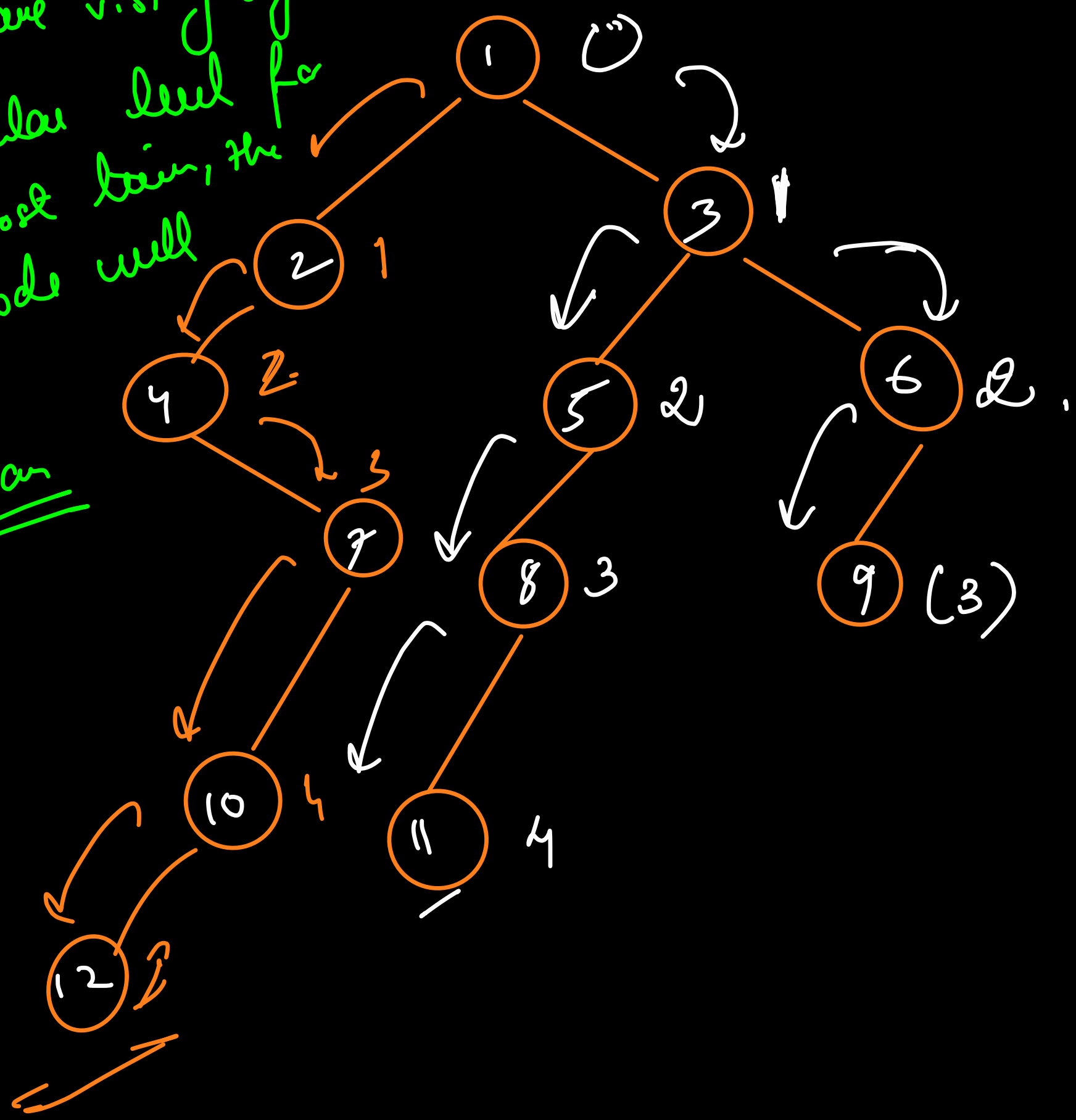
$$O(n)$$

$O(n)$

null

best

if we are visiting any particular level for the first time, the cur node will be part of an



Right cur node
 1, 3, 6, 9, 11,
12
Right
Left

max_level = 0
 cur_level = 0
 (cd, ml)

$f(\text{root}, cl) =$

- if ($cl > ml$) ?
- $ml = cl$
- $\text{root}.push(\text{root}.val)$
- $f(\text{root}.right, cl+1)$
- $f(\text{root}.left, cl+1)$

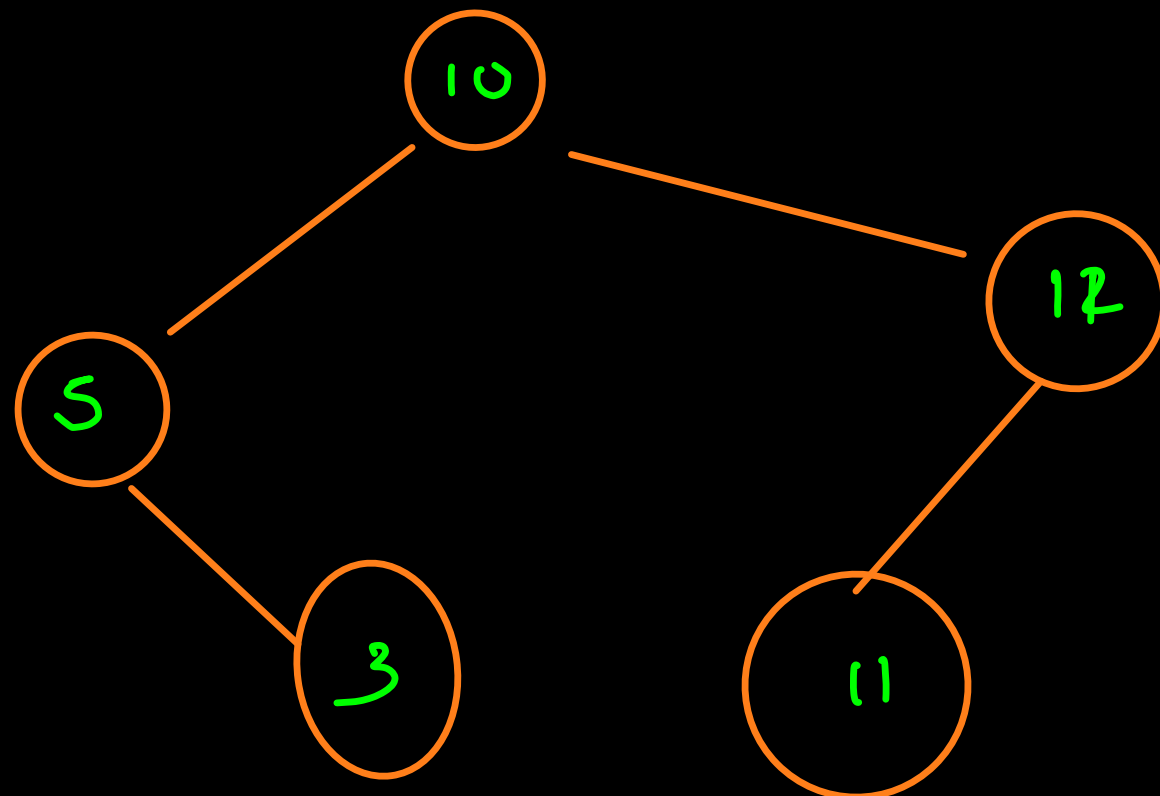
→ BST

→ Binary Tree

all the elements of $l < \text{root}$

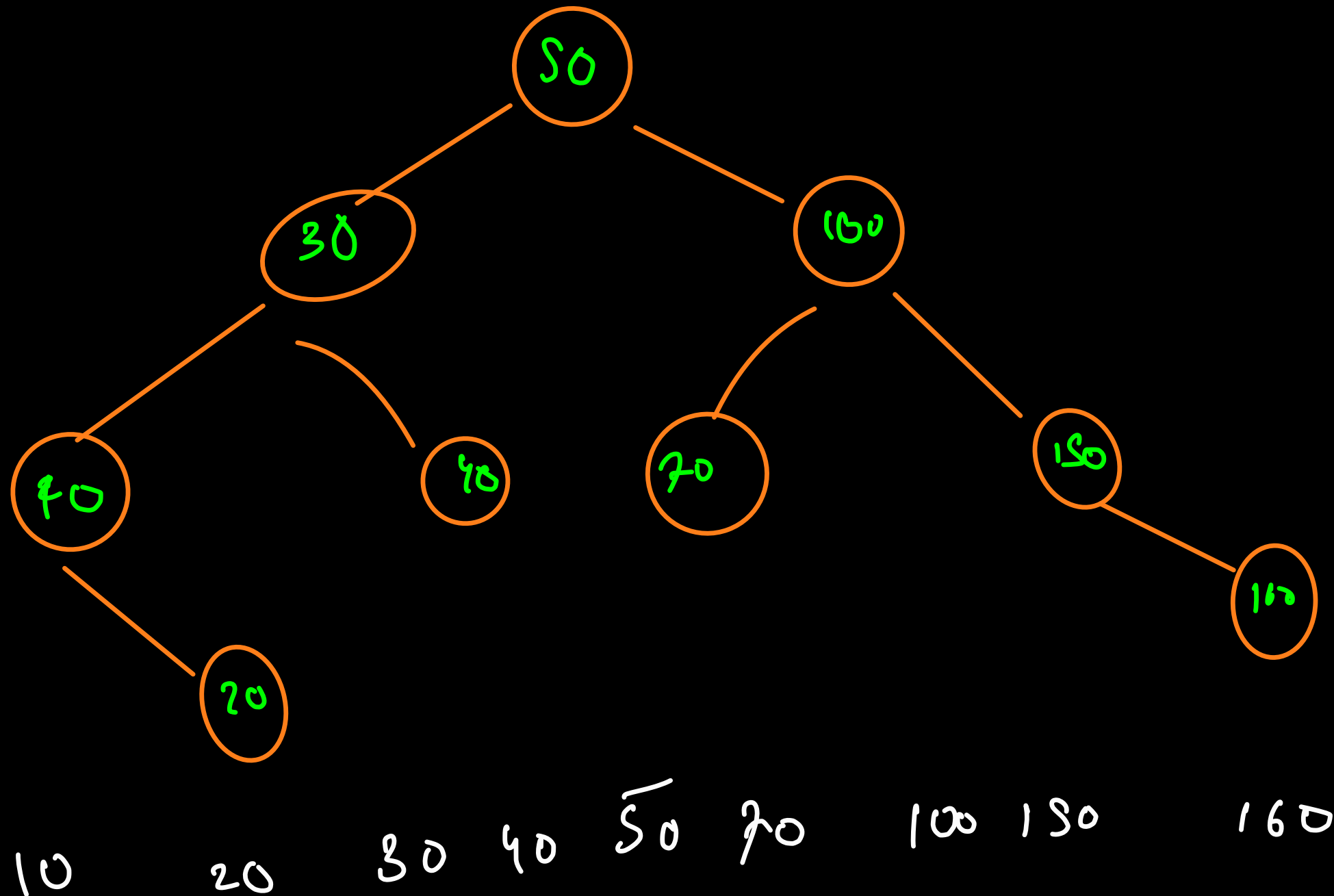
all the elements of $r > \text{root}$

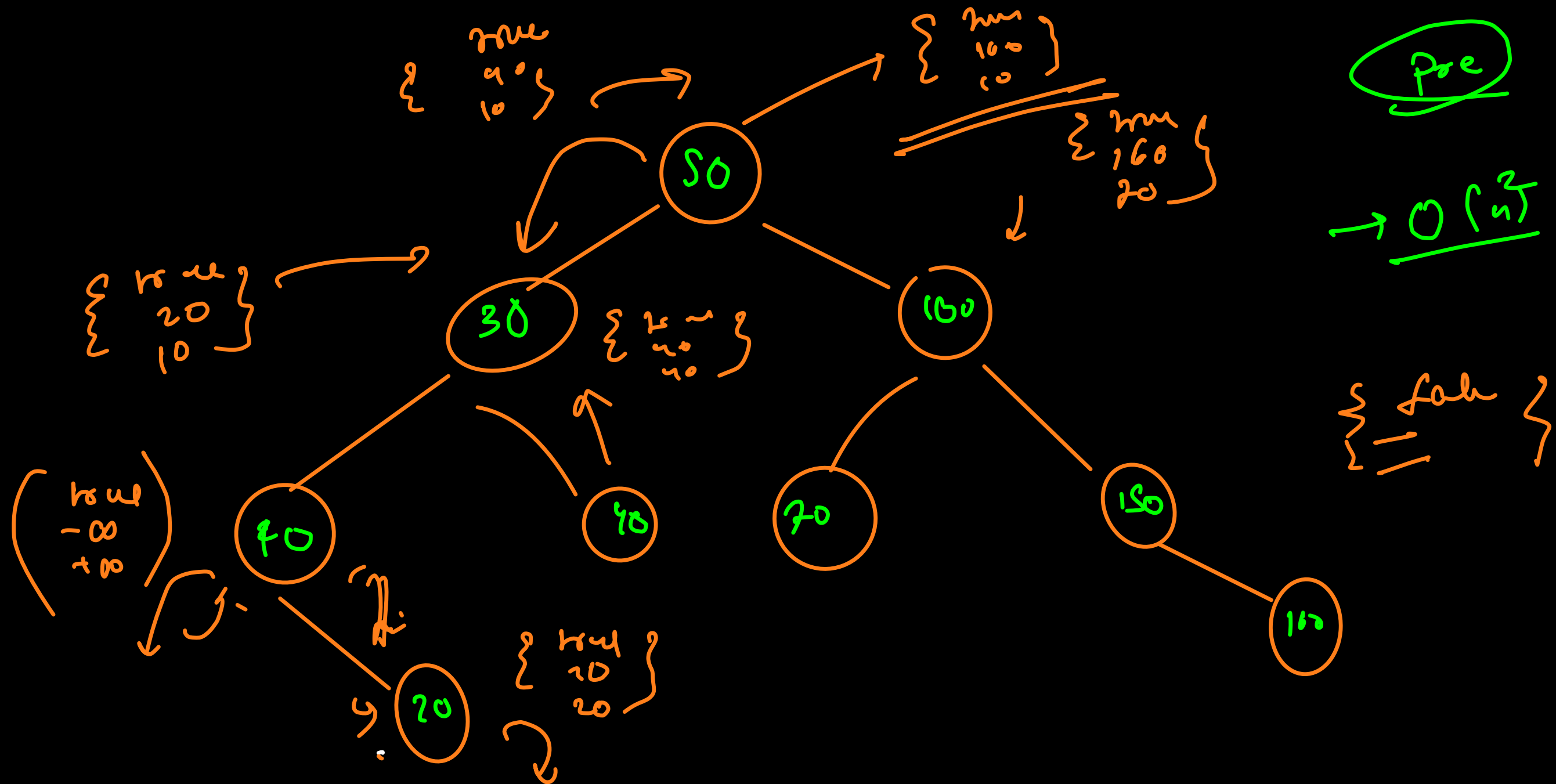
recursion



Property²

→ Inorder traversal of a BST is always sorted.





$O(n)$

$f(\text{root})_{\text{isBST}} =$
 $l - f(\text{root}.left)$
 $r - f(\text{root}.right)$

{

max

min

isBST

if (l.isBST and r.isBST and

$\text{root.val} > l.\text{max}$ and $\text{root.val} < r.\text{min}$)

return { isBST: true

max: $\max(l.\text{max}, r.\text{max}, \text{root.val})$

min: $\min(l.\text{min}, r.\text{min}, \text{root.val})$

}

}

root = null

→

{

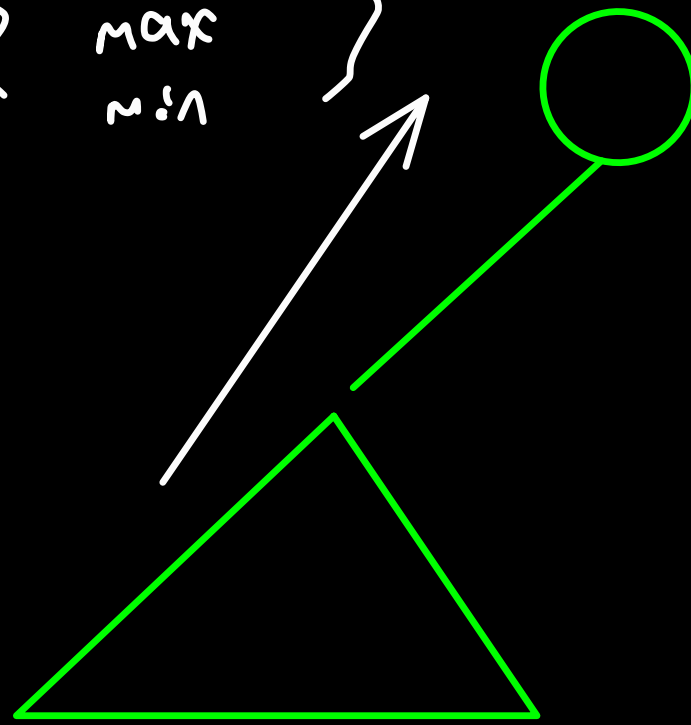
isBST → true

max → $-\infty$

min → $+\infty$

}

{ isBST
max
min }



{ isBST
max
min }

