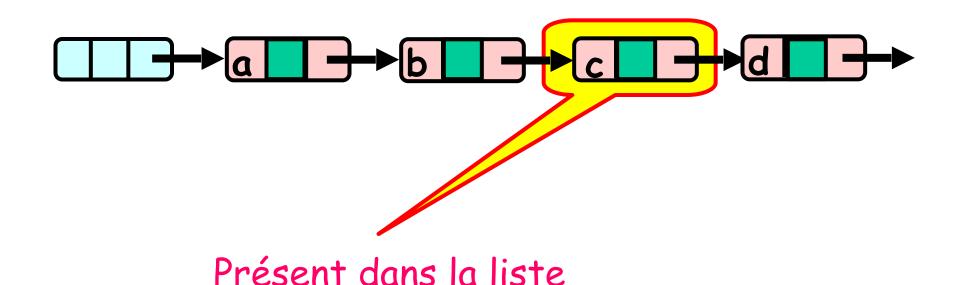
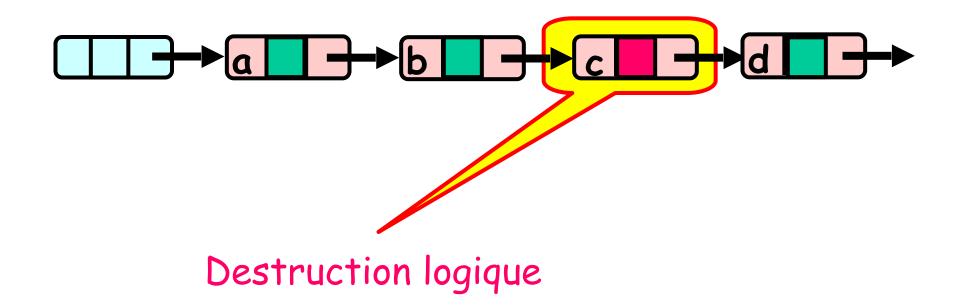
Implémentation paresseuse de Liste

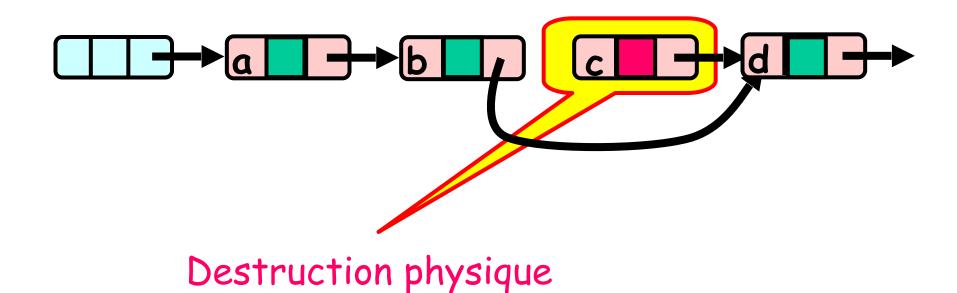
- · Comme optimiste, sauf que
 - On traverse une seule fois pour add et remove (mais restent avec des verrous)
 - contains(x) sans verrous ...
- · Comment?
 - Détruite les noeuds posent des problèmes
 - On le fait "paresseusement"

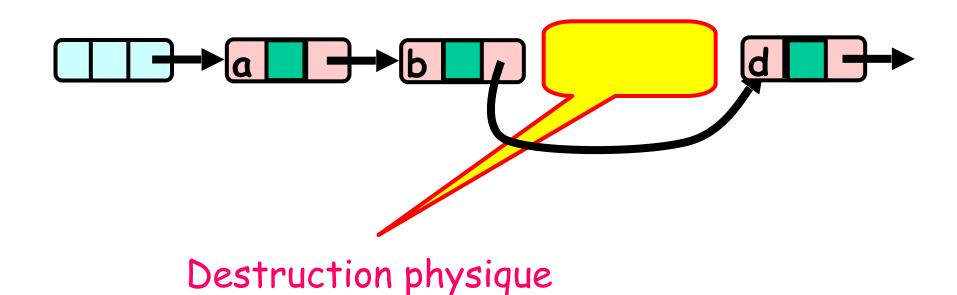
- remove()
 - Parcourir la liste (comme avant)
 - Verrouiller prédécesseur & courant (comme avant)
- · Destruction logique
 - Marqué le noeud courant comme enlevé (nouveau!)
- Destruction physique
 - Redirige le suivant de prédécesseur (comme avant)







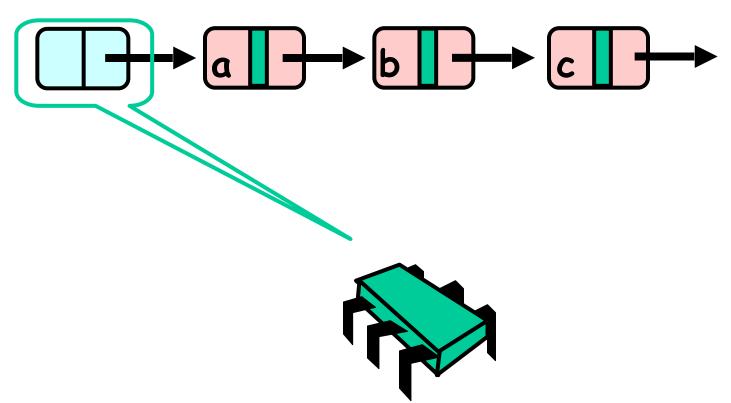


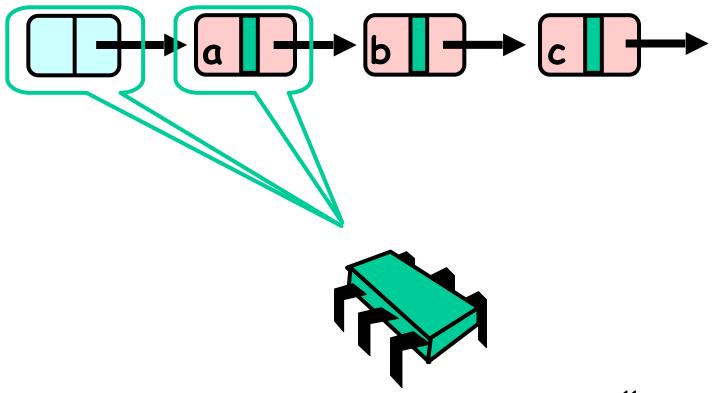


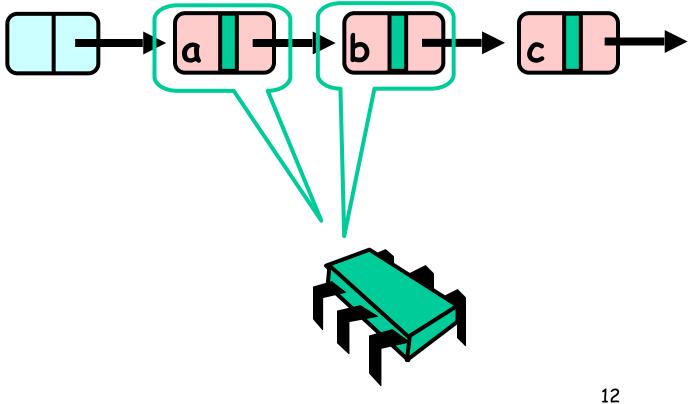
Destruction physique

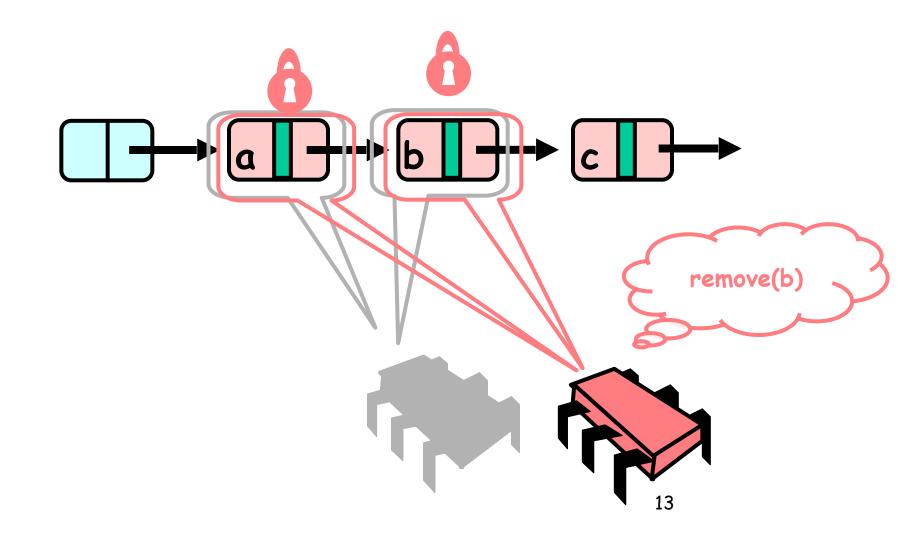
- Toutes les méthodes
 - Parcours sans verrou parmi les noeuds verrouillés et marqués
 - La destruction d'un noeud ne ralentit pas les autres appels...
- · On doit encore verrouiller pred et curr

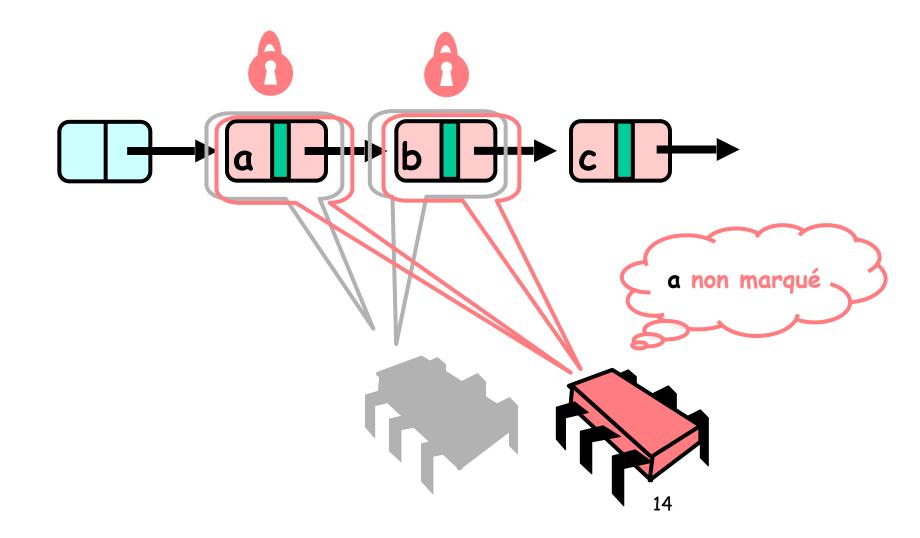
- · Pas besoin de reparcourir la liste!
- · Vérifie que pred n'est pas marqué
- · Vérifie que curr n'est pas marqué
- Vérifie que le successeur de pred est curr

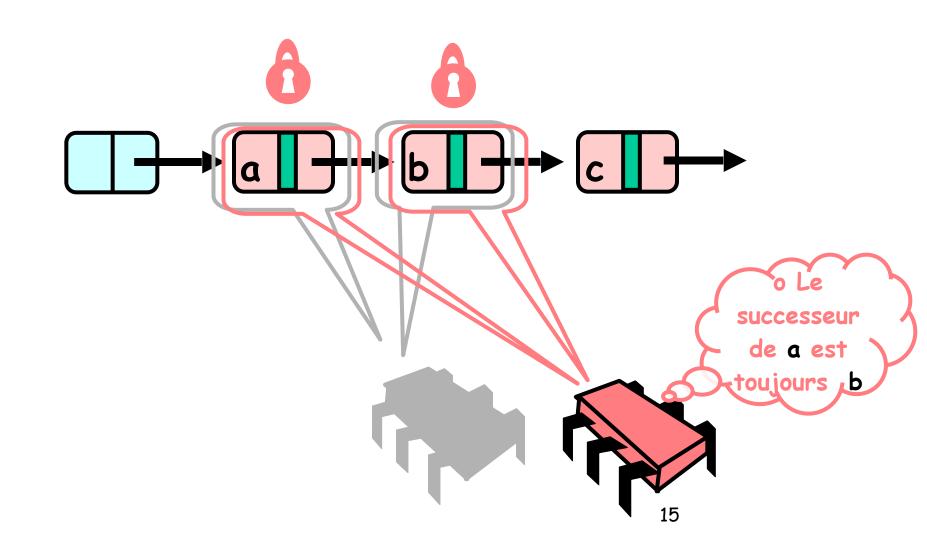


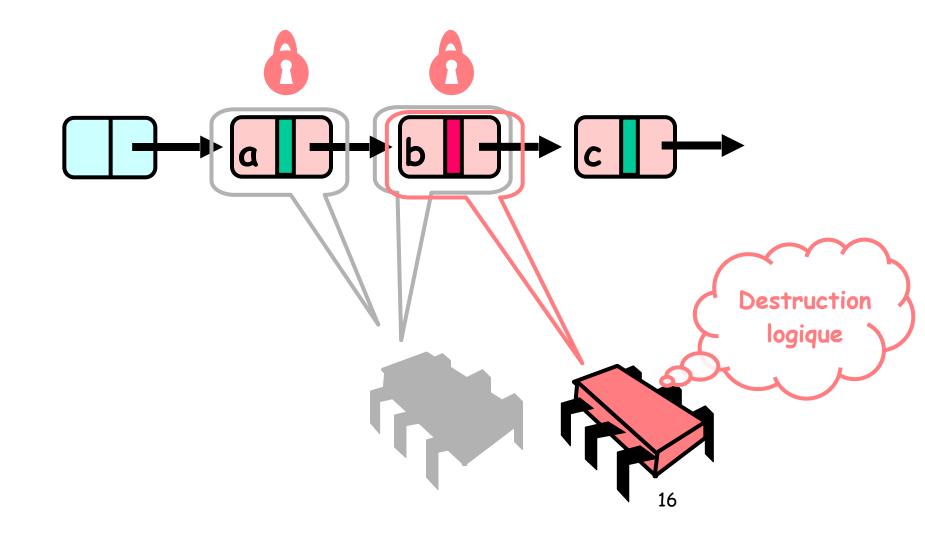


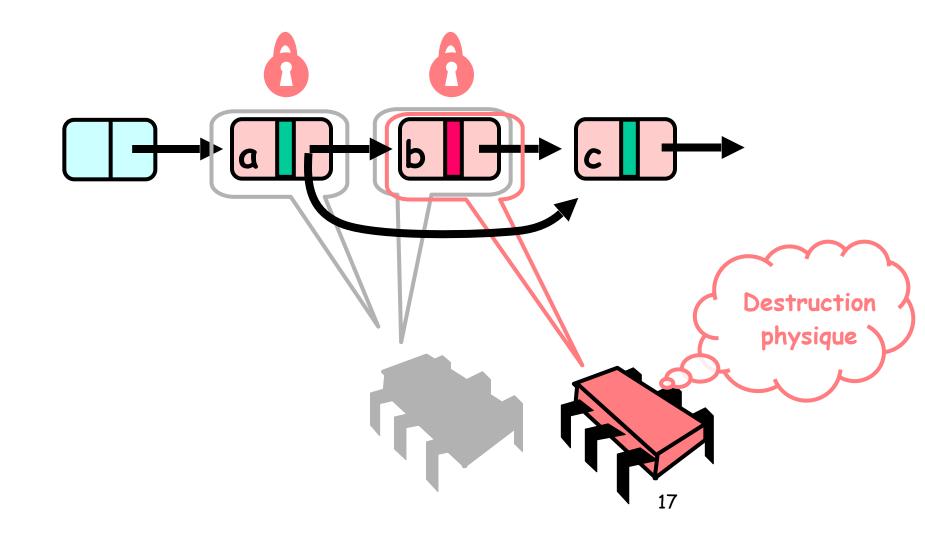


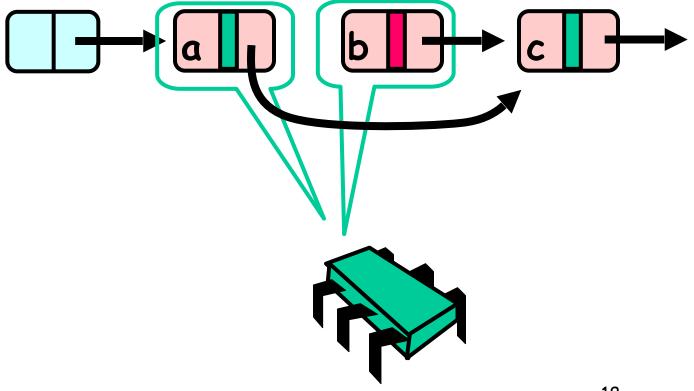












Nouvelle application

```
S(head) =
-{x | il existe un noeud a tel que
· a atteignable de head et
· a.item = x et
· a n'est pas marqué
-}
```

Invariant

- · Si non marqué l'item est dans set
- Et atteignable depuis head
- Et si la traversée n'est pas fini atteignable de pred

```
private boolean
validate(Node pred, Node curr) {
return
!pred.marked &&
!curr.marked &&
pred.next == curr);
}
```

```
private boolean
 validate(Node pred, Node curr) {
return
 !pred.marked &&
 curr marked &&
 pred.next == cur
                         Pred pas détruit
                             logiquement
```

```
private boolean
 validate(Node pred, Node curr) {
return
 !pred.marked &&
 !curr.marked &&
  med next -- curr
                                Curr pas détruit
                                   logiquement
```

```
private boolean
validate(Node pred, Node curr) {
return
!pred.marked &&
!curr.marked &&
pred.next == curr);
}
```

Le suivant de pred est toujours curr

```
public boolean remove(Item item) {
  int key = item.hashCode();
  while (true) {
    Node pred = head;
    Node curr = pred.next;
    while (curr.key < key) {
        pred = curr; curr = curr.next;}
}</pre>
```

```
try {
 pred.lock(); curr.lock();
 if (validate(pred, curr) {
  if (curr.key == key) {
   curr.marked = true;
   pred.next = curr.next;
   return true;
  } else {
   return false:
  }}} finally {
        pred.unlock();
        curr.unlock();
  }}}
```

```
try {
 pred.lock(); curr.lock();
 if (validate(pred, curr) {
  if (curr key == key) {
   curr.marked = true;
   pred.next = curr.next;
   return true:
  } else {
                                   Validate comme avant
   return false;
  }}} finally {
        pred.unlock();
        curr.unlock();
  }}}
```

```
try {
 pred.lock(); curr.lock();
 if (validate(pred,curr) {
  if (curr.key == key) {
   curr.marked = true;
   pred.next = curr.next;
   return true:
  } else {
   return false:
  }}} finally {
        pred.unlock();
                                         Key trouvé
        curr.unlock();
  }}}
```

```
try {
 pred.lock(); curr.lock();
 if (validate(pred,curr) {
  if (curr.key == key) {
   curr.marked = true;
   pred.next = curr.next;
   return true:
  } else {
   return false:
  }}} finally {
                                  Destruction logique
        pred.unlock();
        curr.unlock();
  }}}
```

```
try {
 pred.lock(); curr.lock();
 if (validate(pred,curr) {
  if (curr.key == key) {
   curr.marked = true:
   pred.next = curr.next;
  } else {
   return false:
  }}} finally {
        pred.unlock();
        curr.unlock();
                               Destruction physique
  }}}
```

```
try {
 pred.lock(); curr.lock();
 if (validate(pred, curr) {
  if (curr.key == key) {
   curr.marked = true;
   pred.next = curr.next;
   return true;
  } else {
   return false:
                          On recommence si validate faux
  }}} finally {
        pred.unlock();
        curr.unlock();
```

Add

```
public boolean add(Item item) {
  int key = item.hashCode();
  while (true) {
    Node pred = head;
    Node curr = pred.next;
    while (curr.key < key) {
        pred = curr; curr = curr.next;}
}</pre>
```

Add

```
try {
 pred.lock(); curr.lock();
 if (validate(pred, curr) {
  if (curr.key != key) {
    Node node=new Node(item);
   Node.next=curr;
   pred.next = node;
   return true;
  } else {
   return false;
  }}} finally {
        pred.unlock();
        curr.unlock();
  }}}
```

Contains

```
public boolean contains(Item item) {
  int key = item.hashCode();
  Node curr = this.head;
  while (curr.key < key) {
    curr = curr.next;
  }
  return curr.key == key && !curr.marked;
}</pre>
```

Contains

```
public boolean contains(Item item) {
    int key = item.hashCode();
    Node curr = this.head;
    while (curr.key < key) {
        curr = curr.next;
    }
    return curr.key == key && \text{curr marked;}
}</pre>
```

On commence à head

Contains

```
public boolean contains(Item item) {
 int key = item.hashCode();
 Node curr = this.head;
 while (curr.key < key) {
   curr = curr.next;
 return curr.key == key &&!curr.marked;
```

Recherche par la clef

Contains

```
public boolean contains(Item item) {
  int key = item.hashCode();
  Node curr = this.head;
  while (curr.key < key) {
    curr = curr.next;
    return curr.key == key &&!curr.marked;
}</pre>
```

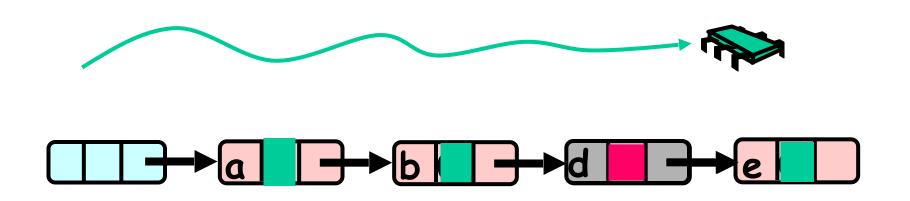
Traverse sans verrou (Les noeuds peuvent avoir été enlevés) 37

Contains

```
public boolean contains(Item item) {
 int key = item.hashCode();
 Node curr = this.head;
 while (curr.key < key) {
   curr = curr.next;
 return curr.key == key && !curr.marked;
```

Present et non enlevé?

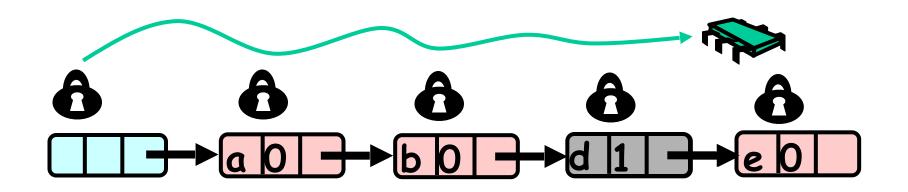
Résumé: Wait-free Contains



Utilise un bit indiquant le noeud marqué et le fait que la liste soit ordonnée

- 1. Pas marqué → dans set
- 2. Marqué ou absent → pas dans set

Paresseux



add() et remove() paresseux (blocking) + Waitfree contains()

Evaluation

· Bon:

- contains() sans lock (wait-free!)
- Bon car souvent beaucoup d'appel à contains()
- Pas de retraversée en cas de succès

Mauvais

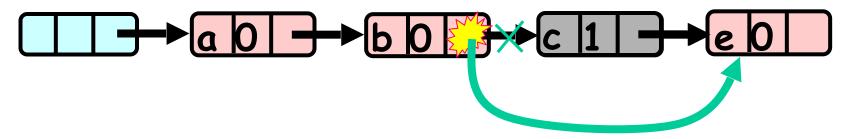
- Contention sur add() et remove(), on peut être amener à retraverser
- Une thread lente ralentit l'appel de ces méthodes

Implémentation sans verrou

- · Elimine entièrement les verrous
- contains() wait-free; add() et remove() non blocking
- Utilise compareAndSet()
- Que peut-il se passer?

Remove en utilisant CAS

Remove logique = Positionne un bit



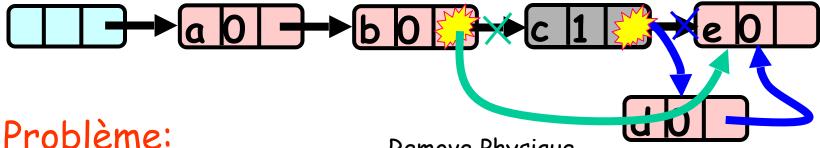
Utiliser CAS pour modifier le champ next

Remove physique modification du pointeur CAS

Insuffisant

Problème...

Remove logique = Positionne un bit



d n'a pas été ajouté à la liste...

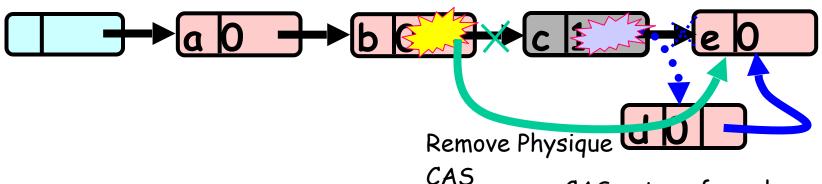
Faire attention au pointeur des noeuds enlevés

Remove Physique CAS

Noeud ajouté avant le remove physique par CAS

Solution: Combiner Bit et CAS

Remove logique = Positionne un bit



Bit de marque et le pointeur sont ensemble dans le CAS

(AtomicMarkableReference)

CAS retour faux: le noeud est ajoute après un remove logique

Solution

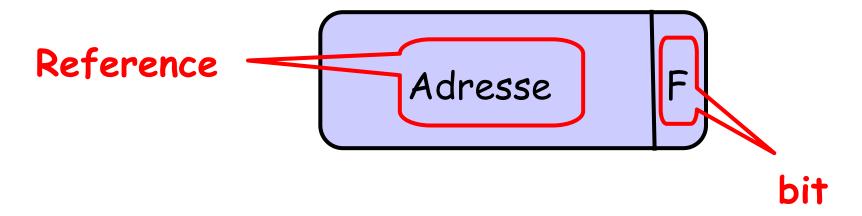
- Utilise AtomicMarkableReference
- Atomiquement
 - Modifie la référence et
 - Met à jour la marque

Toute mise à jour du champ next quand la marque est posée échouera

- · Remove en deux étapes
 - Mettre la marque dans le champ next
 - Rediriger le pointeur du prédécesseur

Marquer un noeud

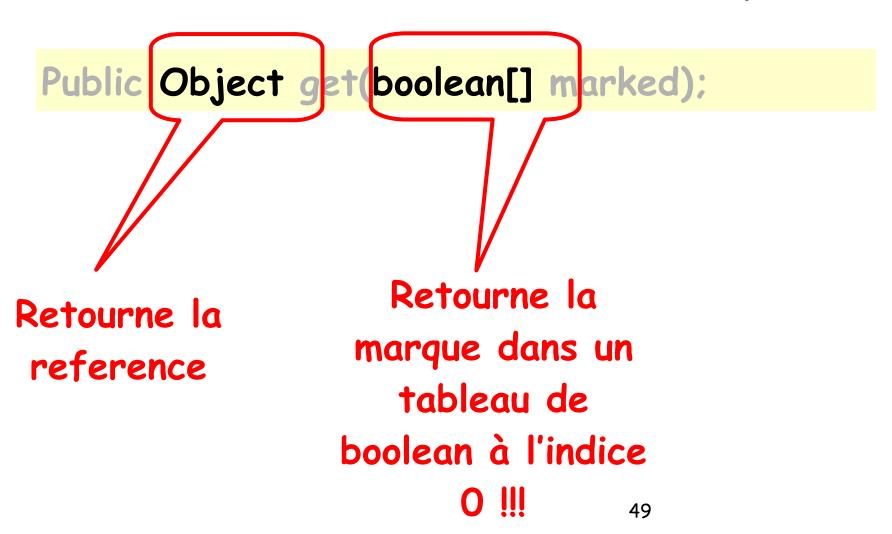
- AtomicMarkableReference class
 - Java.util.concurrent.atomic package



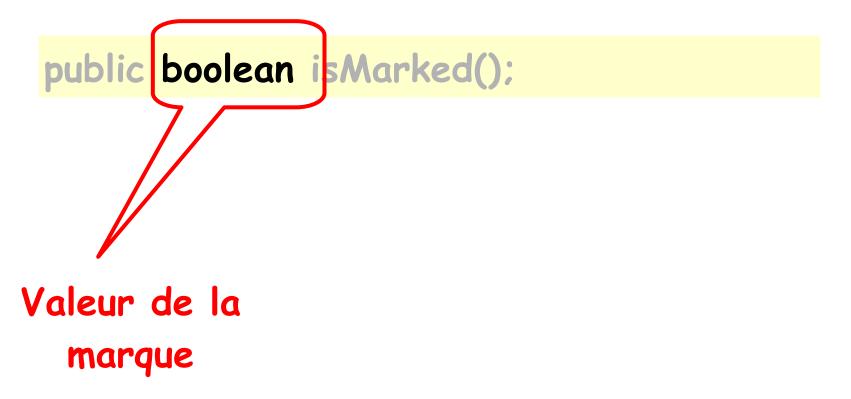
Extraire Reference & Marque

Public Object get(boolean[] marked);

Extraire Reference & Marque



Extraire Reference & Marque



```
Public boolean compareAndSet(
Object expectedRef,
Object updateRef,
boolean expectedMark,
boolean updateMark);
```

Si cette référence est la référence Public boolean compare And Scottrante Object expectedRef Object updateRef, boolean expectedMark, Et si cette marque boolean updateMa est la marque courante

...alors changer pour cette nouvelle

Public boolean compareAndSet(
Object expectedRef,
Object updateRef,
boolean expectedMark,
boolean updateMark);

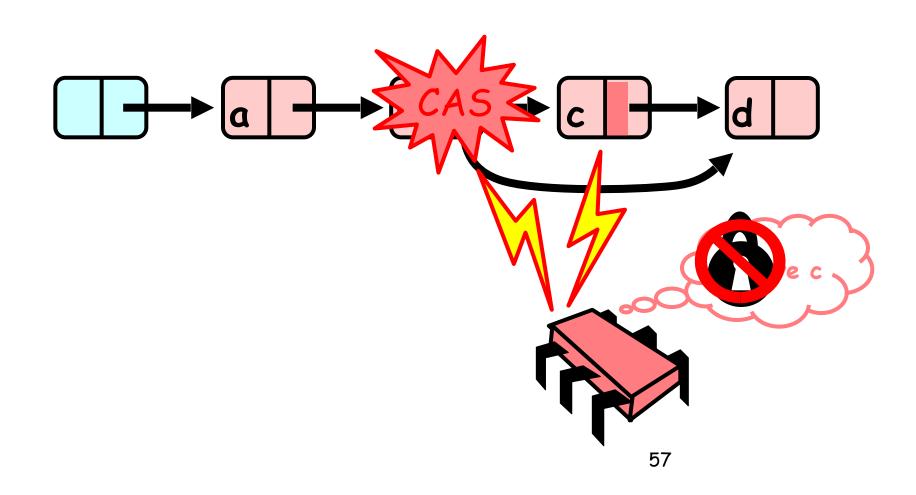
... et cette nouvelle marque

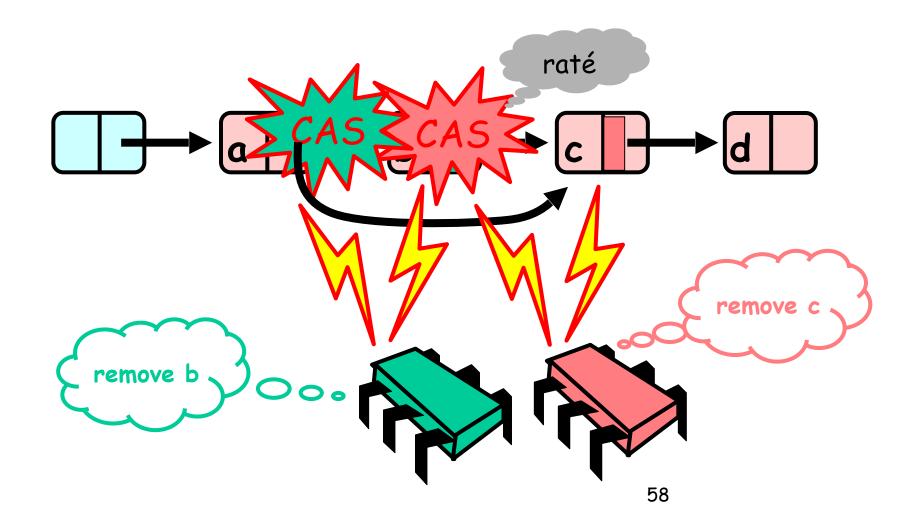
```
public boolean attemptMark(
   Object expectedRef,
   boolean updateMark);
```

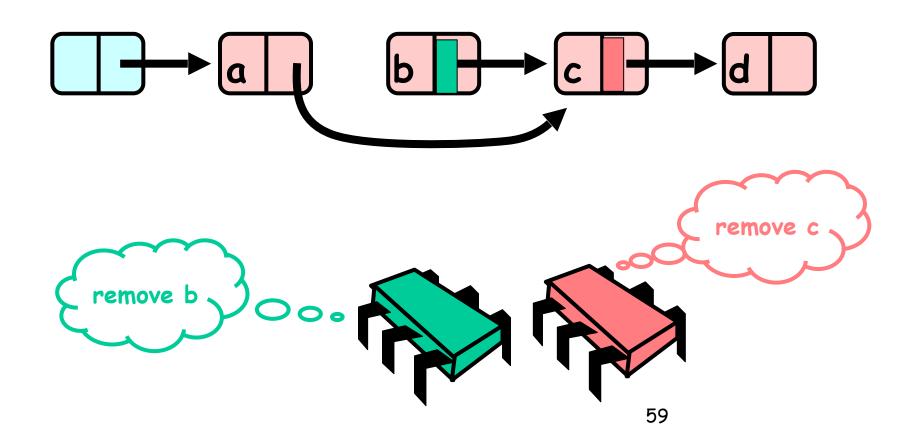
```
public boolean attemptMark(
    Object expectedRef,
    boolean updateMark);
Si cette reference est
```

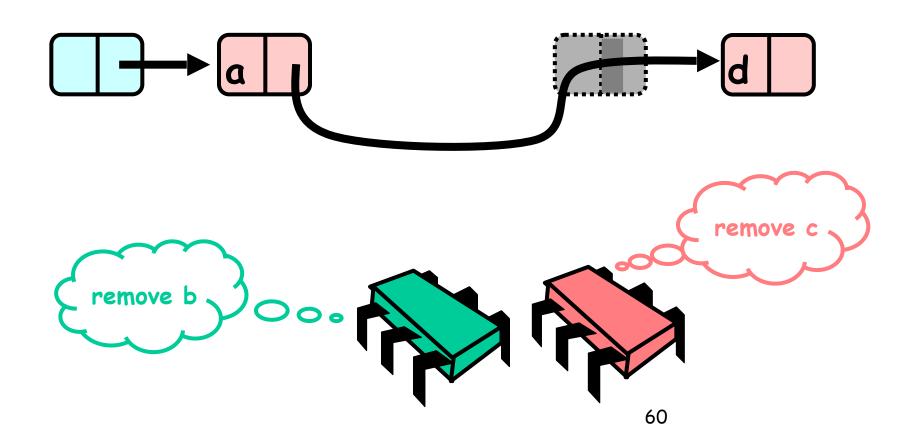
Si cette reference est la reference courante ...

```
public boolean attemptMark(
  Object expectedRef,
 boolean updateMark);
 .. alors changer
pour cette nouvelle
     marque.
```





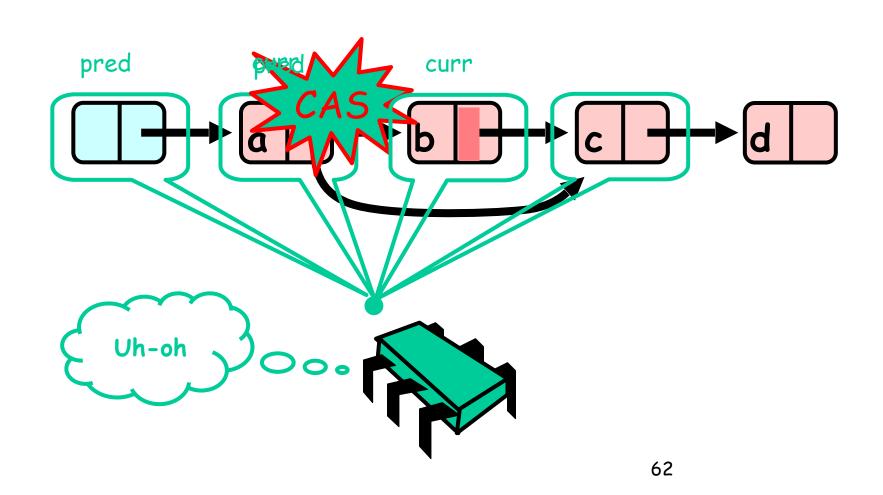




Traverser la liste

- Q: que faire quand on trouve un noeud logiquement enlevé sur le chemin?
- R: finir le travail.
 - Modifier par CAS le champ next du prédécesseur
 - Et repeter si besoin

Traversée sans verrou (seulement Add et Remove)



Window Class

```
class Window {
  public Node pred;
  public Node curr;
  Window(Node pred, Node curr) {
    this.pred = pred; this.curr = curr;
  }
}
```

Window Class

```
class Window {
  public Node pred;
  public Node curr;
  Window(Node pred, Node eurr) {
    this.pred = pred; this.curr = curr;
  }
}
```

Regroupe pred et curr

```
Window window = find(head, key);
Node pred = window.pred;
curr = window.curr;
```

```
Window window = find(head, key);

Node pred = window.pred;

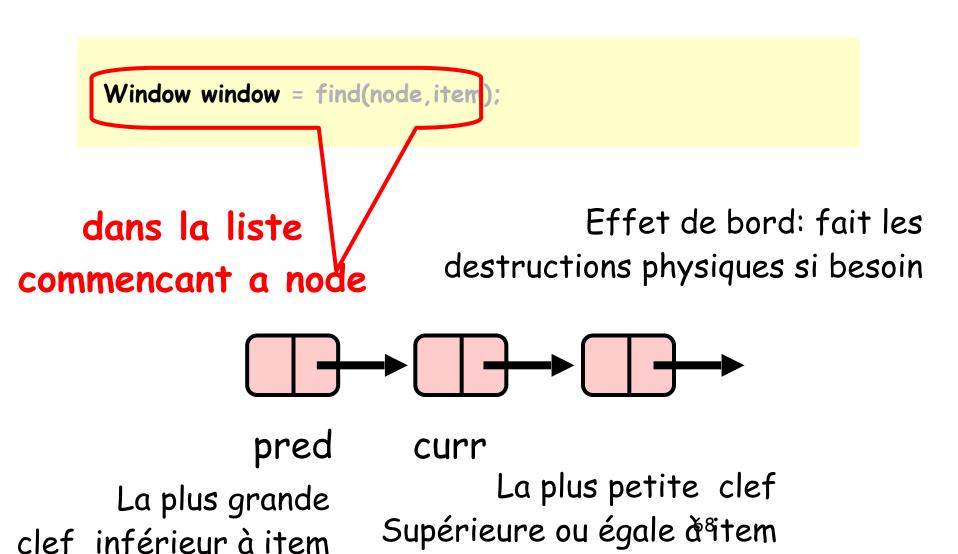
curr = window.surr;
```

Find retourne window

```
Window window = find(head, key);

Node pred = window.pred;
curr = window.curr;

Extraction de pred et curr
```



```
public boolean remove(T item) {
Boolean snip; int key = item.hashCode();
while (true) {
Window window = find(head, key);
Node pred = window.pred; Node curr = window.curr;
 if (curr.key != key) {
    return false:
 } else {
 Node succ = curr.next.getReference();
 snip = curr.next.attemptMark(succ, true);
 if (!snip) continue;
  pred.next.compareAndSet(curr, succ, false, false);
    return true:
```

```
public boolean remove(T item) {
        snip; int key = item.hashCode();
while (true)
Window window find(head, key);
Node pred = window.pred, curr = window.curr;
 if (curr.key != ke
    return false:
 } else {
 Node succ = curr.next.getReference();
 snip = curr.next.attemptMark(\ucc, true);
                                     On continue tant
 if (!snip) continue;
  pred.next.compareAndSet(curr, sucqu'or af pas reussi
    return true:
```

```
public boolean remove(T item) {
Boolean snip; int key = item.hashCode();
while (true) {
Window window = find(head, key);
Node pred = window.pred, curr = window.curr;
 if (curr.key != key) {
    return false:
 } else {
 Node succ = curr.next.getReference()
 snip = curr.next.attemptMark(succ, true
 if (!snip) continue;
  pred.next.compareAndSet(curr, succ, false, false);
Trouver les voisins
    return true:
```

```
public boolean remove(T item) {
Boolean snip; int key = item.hashCode();
while (true) {
Window window = find(head, key);
 Node pred = window.pred, curr = window.curr;
 if (curr.key != key) {
    return false;
 } else {
 Node succ = curr.next.getReference(); L'item n'est pas dans
 snip = curr.next.attemptMark(susc, true);
                                                  la liste
 if (!snip) continue;
  pred.next.compareAndSet(curr, succ, false, false);
    return true:
```

Remove

```
public boolean remove(T item) {
Boolean snip
wh On essaie de marquer le noeud comme
Window window = fing(heenlevé
Node pred = window.pred, curr = window.curr;
 if (curr.key != Key)
    return false;
 } else {
 Node succ = curr.next.getReference();
 snip = curr.next.attemptMark(succ, true);
 if (!snip) continue;
  pred.next.compareAndSet(curr, succ, false, false);
    return true:
```

Remove

```
public boolean remove(T item) {
Boolean snip:
whileSinca (n'a pas
 Windmarché on (head, key)
 réessaie, sinon le
   travail est
 pratiquement fait etReference();
 snip = curr.next.attemptMark(succ, true);
if (!snip) continue;
  pred.next.compareAndSet(curr, succ, false, false);
   return true;
```

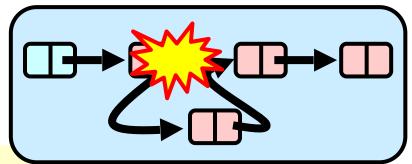
Remove

```
public boolean remove(T item) {
Boolean snip; int key = item.hashCode
while (true) {
Window window = find(head, key);
Node pred = window.pred, curr = window.curr;
 if (curr.key != key) {
                                    On essaie d'avancer la reference
    return false:
                                    (si retour faux c'est que quelqu'un
 } else {
                                    d'autre l'a fait ou le fera).
 Node succ = curr.next.getReference();
 snip = curr.next.attemptMark(succ, true);
 if (!snip) continue
  pred.next.compareAndSet(curr, succ, false, false);
    return true;
```

```
public boolean add(T item) {
int key = item.hashCode();
while (true) {
  Window window = find(head, key);
  Node pred = window.pred; Node curr = window.curr;
  if (curr.key == key) {
     return false:
  } else {
  Node node = new Node(item);
  node.next = new AtomicMarkableRef(curr, false);
  if (pred.next.compareAndSet(curr, node, false, false)) {return
true;}
```

```
public boolean add(T item) {
int key = item.hashCode();
while (true) {
  Window window = find(head, key);
  Node pred = window.pred curr = window.curr;
  f (curr.key == key) {
     return false;
  } else {
  Node node = new Node(item);
  node.next = new AtomicMarkableRef(curr, false);
  if (pred.next.compareAndSet(curr, node, false, false)) {return
true;}
                           Item déjà la.
```

```
public boolean add(T item) {
int key = item.hashCode();
while (true) {
  Window window = find(head, key);
  Node pred = window.pred, curr = window.curr;
  if (curr.key == key) {
    return false:
  } else {
  Node node = new Node(item);
  node.next = new AtomicMarkableRef(curr, false);
  if (pred.next.compareAndSet(curr, node, false, false)) {return
true:}
                   Crée le nouveau noeud
```



```
public boolean add(T item) {
int key = item.hashCode();
                                   Installe le nouveau
while (true) {
  Window window = find(head, key)
                                        noeud sinon
  Node pred = window.pred, curr
                                       recommence
  if (curr.key == key) {
     return false;
  } else {
  Node node = new Node(item);
  node.next = new AtomicMarkableRef(curr, false);
  if (pred.next.compareAndSet(curr, node, false, false)) {return
true;}
```

Wait-free Contains

```
public boolean contains(T item) {
   boolean[] marked;
   int key = item.hashCode();
   Node curr = this.head;
   while (curr.key < key)
      curr = curr.next;
   Node succ = curr.next.get(marked);
   return (curr.key == key && !marked[0])
}</pre>
```

Wait-free Contains

```
public boolean contains(T item) {
    boolean marked;
    int key = item.hashCode();
    Node curr = this.head;
    while (curr.key < key)
    curr = curr.next;
    iNode succ = curr.next.get(markecurr.next.get et on return (curr.key == key && !marked[0]) vérifie que
    marked[0] est vrai
```

```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null, t=null
boolean[] marked = {false}; boolean snip;
retry: while (true) {
  pred = head;
  curr = pred.next.getReference();
  while (true) {
   succ = curr.next.get(marked);
   while (marked[0]) {
   if (curr.key >= key)
        return new Window(pred, curr);
      pred = curr;
      curr = succ;
```

```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
 while (true) {
  pred = head;
                                          Si la liste a change
  curr = pred.next.getReferen
                                          pendant le parcours,
  while (true) {
                                             on recomence
   succ = curr.next.get(marked);
                                           Lock-Free car on
   while (marked[0]) {
                                             recommence si
                                           quelqu'un d'autre a
                                               progressé
   if (curr.key >= key)
        return new Window(pred, curr);
      pred = curr;
      curr = succ;
```

```
public Window find(Node head int key) (
Node pred = null, curr On commence head
boolean[] marked = {false}; boolean; snip;
retry: while (true) {
  pred = head;
  curr = pred.next.getReference();
  while (true) {
   succ = curr.next.get(marked);
   while (marked[0]) {
   if (curr.key >= key)
       return new Window(pred, curr)
      pred = curr;
      curr = succ;
```

```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
retry: while (true) { On parcours la liste
  pred = head;
  curr = pred.next.getRefer
  while (true) {
   succ = curr.next.get(marked);
   while (marked[0]) {
   if (curr.key >= key)
       return new Window(pred, curr);
      pred = curr;
      curr = succ;
```

```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
retry: while (true) {
  pred = head;
  curr = pred.next.getReference();
  while (true) {
  succ = curr.next.get(marked);
   while (marked[0]) {
   if (curr.key >= key)
       return new Window(pred, curr);
                      ref du successeur
      pred = curr;
      curr = succ;
                        et bit de marque
```

```
public Window find(Node head, int key) {
 Node pred = null, curr = null, succ = null;
 boolean[] marked = {false}; boolean snip;
 retry: while (true) {
   pred = head;
   curr = pred.next.getReference();
   while (true) {
    succ = curr.next.get(markea);
    while (marked[0]) {
   if (curr.key >= key)
On tentendienlever (des noeuds
                  » au passaage
```

```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
retry: while (true) {
  nnad - haad.
  si curr key est >= retourne
            pred et curr
   while (marked[0])
   if (curr.key >= key)
       return new Window(pred, curr);
      pred = curr;
      curr = succ:
```

```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
retry: while (true) {
  pred = head;
  curr = pred.next.getReference();
   Sinon avancer la fenêtre et
          refaire la boucle
     (curr.key >= key)
       return new Window(pred, curr);
      pred = curr;
      curr = succ;
```

```
retry: while (true) {
  while (marked[0]) {
    snip = pred.next.compareAndSet(curr, succ, false,
false);
    if (!snip) continue retry;
    curr = succ;
    succ = curr.next.get(marked);
```

On essaie d'enlever le noeud

```
retry: while (true) {
    snip = pred.next.compareAndSet(curr, succ, false,
false);
    if (!snip) continue retry;
    curr = succ;
    succ = curr.next.get(marked);
```

Si le champ suivant du predecesseur a changé on

```
retry: while (true) {
                             doit refaire tout le
                                    parcours
  while (marked[0]) {
    snip = pred.next.compare/ndSet(curr, succ, false,
false);
   if (!snip) continue retry;
    curr = succ;
    succ = curr.next.get(marked);
```

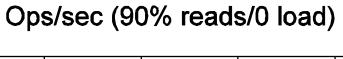
verifier si le noeud suivant retry: while (true) { est detruit while (marked[0]) { snip = pred.next.compareAndSet(curr, succ, false, false); if (!snip) continue re curr = succ; succ = curr.next.get(marked);

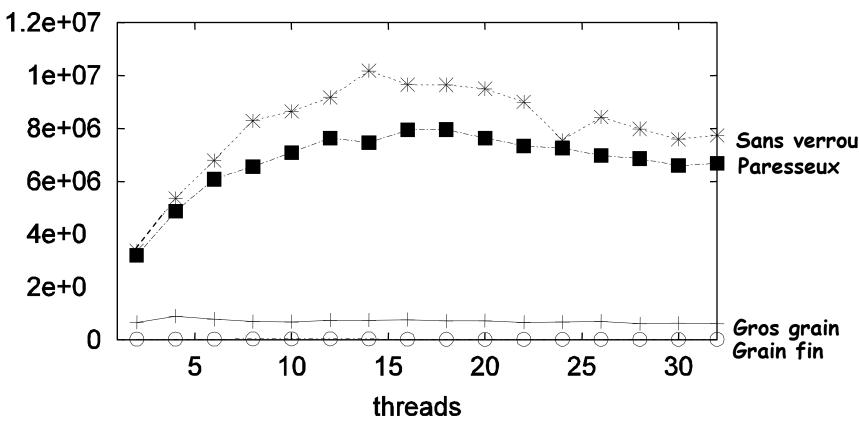
Sinon on avance pour

Implémentation sans verrou

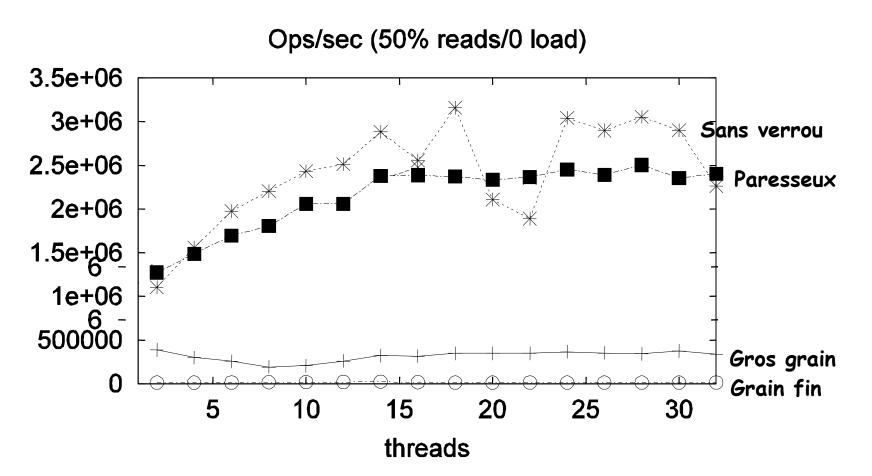
- · Elimine entièrement les verrous
- contains() wait-free; add() et remove() non blocking
- Utilise compareAndSet()

Avec un fort taux de Contains





Avec un faible taux de contains



""Avec ou sans Verrou""

- Blocking vs. Non-blocking: points de vue extrémistes des deux côtés
- Réponse: tenter le compromis, allie verrouillage et non-blocage

Exemple : la liste paresseuse combine le blocage add() et remove() et un wait-free contains() N'oubliez pas : le blocage/non-blocage est une propriété d'une méthode