

MT1003 PURE HELP SHEET

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1 Graphs

1.1 Directed Graphs

Consists of set V of *vertices* and set E of ordered pairs from V making *edges*. Written: $\Gamma = (V, E)$.

Let Γ be a directed graph:

Walk – Sequence of vertices and edges on Γ .

Path – Walk on Γ s.t. no vertex occurs twice.

Circuit – Closed walk on Γ .

Simple Circuit – Closed path on Γ .

Loop – Edge of form (v, v) with $v \in V$.

1.2 Types & Properties of Graphs

Undirected Graph – A graph s.t. $\forall (v_i, v_j) \in \Gamma \exists (v_j, v_i)$. Usually just called *graph*.

Multigraph – Graph with multiple edges between vertices.

Simple Graph – Undirected, finite, no loops, no multiple edges.

Connected Graph – \exists path between any two distinct vertices a, b .

Vertex Degree – # of edges incident on vertex.

Regular Graph – All vertices have same degree.

Complete Graph – Graph, K_n , is simple, connected with n vertices.

Null Graph – Graph, N_n , with n vertices and no edges. Also a regular graph.

Cycle Length – Graph C_n s.t. it forms a circle.

1.3 Trees

Tree – Connected, no circuits.

Spanning Tree – Subgraph of Γ which contains all vertices and is a tree. Minimal connectivity.

Theorem 8.22: \exists a unique path between any two vertices on a tree.

PROOF: Since T is connected, $\exists \geq 1$ path from $a \rightarrow b$. If 1 then \exists a circuit in $T \therefore$ not a tree.

Theorem 8.23: Let Γ be *undirected graph*. Γ connected if and only if it has spanning tree.

Theorem 8.24: \exists 3 non-isomorphic trees with 5 vertices.

Theorem 8.25: Let $T = (V, E)$ be a tree.

Then $|E| = |V| - 1$.

PROOF: [THINGS]

1.4 Eulerian & Hamiltonian Graphs

Eulerian Graph – \exists circuit in Γ passing through every vertex and edge exactly once.

Semi-Eulerian Graph – \exists walk in Γ passing through every vertex and edge exactly once.

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SSH Test)