LinkedLists

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Interface - Implementation Separation

A list interface defines the operations and expected behavior of a list, but does not specify how a list implementation will manage list elements internally.

The user of a list does not need to know anything about the implementation to interact with a list through its interface. However, there is no single, "best" implementation of a list and the internal management will affect the efficiency of interface methods and memory use.

Avoiding shifting - more organic lists

Maintaining list elements in their corresponding indexed positions in arrays makes sense and allows fast lookups by index, but imposes O(n) shifting costs for insertions and deletions.

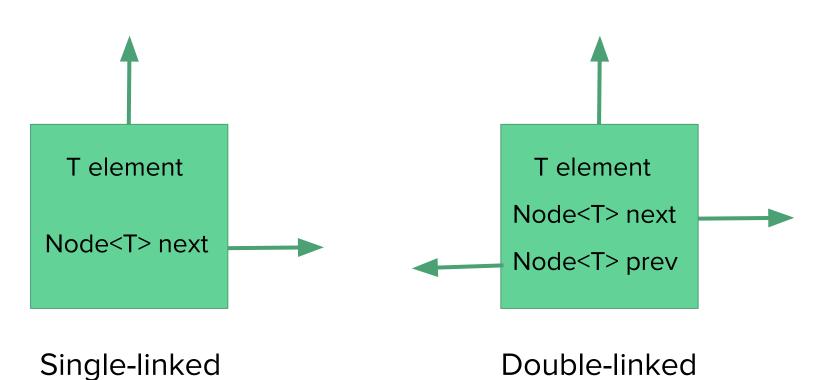
An alternative is to create an implementation that more closely matches the fluid, organic mental model of the interface, where a list is simply a sequence of elements and "space" between consecutive elements automatically expands or shrinks to accommodate list insertion or deletion modifications.

Nodes - building blocks of flexible structures

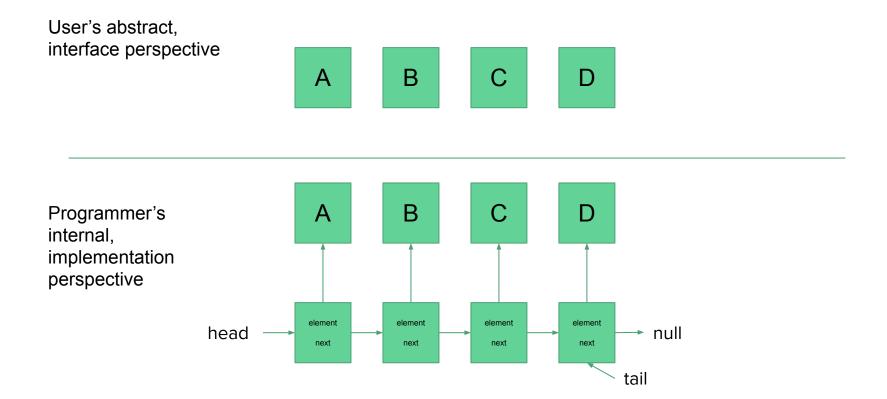
A simple, lightweight object, called a **node**, can be defined to hold the position of one element in a data structure by managing an element reference and at least one reference to another node in the structure.

For a list, a node would need a reference to the list element and a reference to the next node in the list, allowing traversal of all list elements from start to finish. Backward navigation would require an additional reference to the previous node in the list.

Node<T>

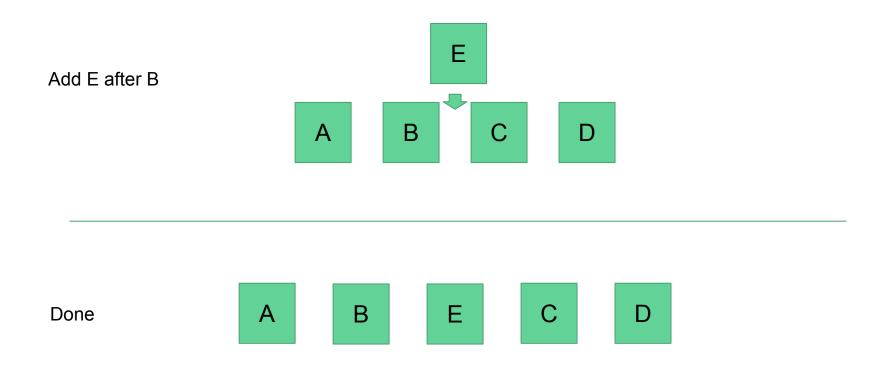


Interface vs Internal Perspective

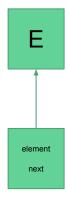


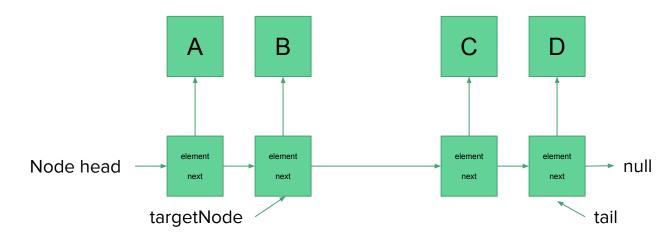
Code - Getting Started

- Write the Node<T> class (stand-alone class or private inner class of the list)
- Implement list interface
- Declare instance variables for head and tail nodes, and size
- Initialize instance variables in a constructor
- Identify useful first utility methods to implement

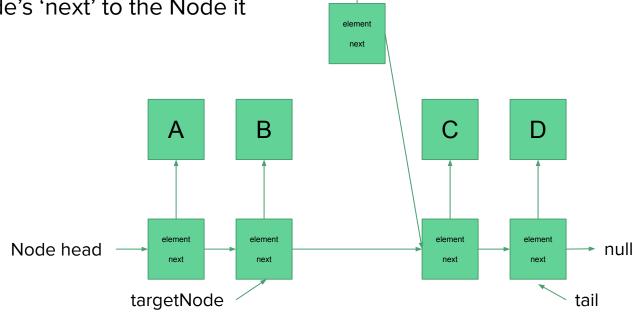


- 1. Locate the target node
- 2. Create a new Node for the new element.





- 1. Locate the target node
- Create a new Node for the new element.
- 3. Connect new Node's 'next' to the Node it will precede.



Ε

element

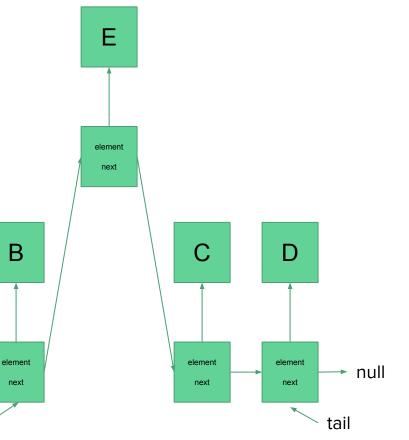
next

targetNode

- 1. Locate the target node
- Create a new Node for the new element.
- Connect new Node's 'next' to the Node it will precede.
- 4. Connect target Node's 'next' to the new

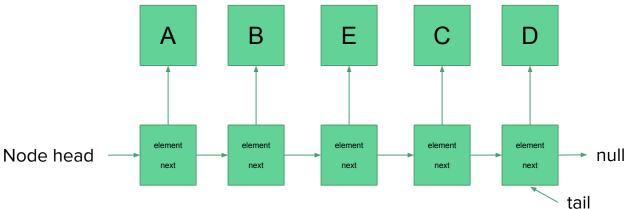
Node head

Node.



- 1. Locate the target node
- Create a new Node for the new element.
- Connect new Node's 'next' to the Node it will precede.
- Connect target Node's 'next' to the new Node.

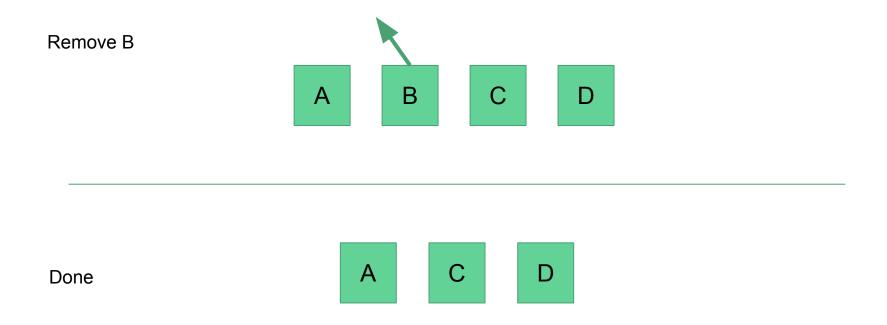
5. Increment size



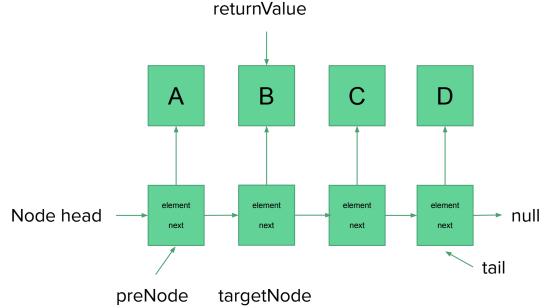
Code - addAfter(newElement, targetElement)

- Locate Node containing targetElement
 - Throw NoSuchElementException if not found
- Create a new Node with the newElement
- Connect new Node's next reference to target Node's next
- Connect target Node's next to new Node
- Update tail reference if necessary
- Increment size

Removing an element - Interface perspective



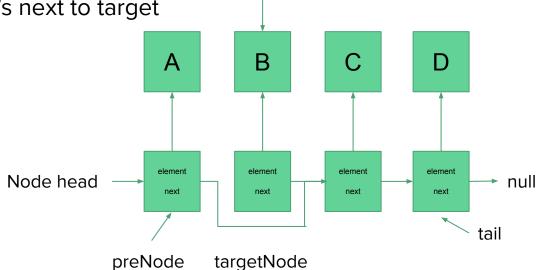
 Locate the node before the target node and remember the element in the target node for later return



 Locate the node before the target node and remember the element in the target node for later return

2. Update predecessor Node's next to target

Node's next



returnValue

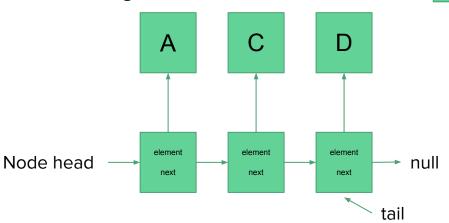
 Locate the node before the target node and remember the element in the target node for later return

2. Update predecessor Node's next to target

Node's next

3. Decrement size

4. Return removed element



returnValue

Code - remove(element)

- Handle special cases: one element or first element
- Else, locate the Node before the Node containing the element
 - Store a reference to the element being removed
 - Update predecessor Node's next to target Node's next
 - Update tail if necessary
- Decrement size
- Return removed element

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