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Examples where threads are useful: Windowing systems, Web browsers, Servers and Clients

How can you be in two place at once when you're not anywhere at all?
—Firesign Theater.

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- ▶ The other way to create a thread is to declare a class that implements the `Runnable` interface. That class then implements the `run` method. An instance of the class can then be allocated, passed as an argument when creating `Thread`, and started.

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`threads/RunnableExample.java`
- ▶ Example 3: Create a thread quagmire!
`threads/MaxThreads.java`

In Java, each thread is an object!

Relevant Java Classes/Interfaces

- ▶ See documentation for basic classes: `java.lang.Thread`, `java.lang.ThreadGroup` and `java.lang.Runnable` interface.
- ▶ See the `java.lang.Object` class for synchronization methods.
- ▶ For automatic management of threads, see: `Executor` interface from `java.util.concurrent` package.

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- ▶ Example: `threads/InterruptTest.java`

A Thread's Life

A thread continues to execute until one of the following thing happens.

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What happens if the `run()` method never terminates, and the application that started the thread never calls the `stop()` method?

The thread remains alive even after the application has finished!
(so the Java interpreter has to keep on running...)

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- ▶ Code snippet:

```
class Devil extends Thread {  
    Devil() {  
        setDaemon( true);  
        start();  
    }  
    public void run() {  
        //perform evil tasks  
        ...  
    }  
}
```

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- ▶ Example of a race condition: **Account.java**, **TestAccount.java**

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// the class.
synchronized void update() { //... }

// Access to individual datum can also be synchronized.
// The object buffer can be used in several classes,
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- ▶ Every Java object has an implicit monitor associated with it to implement the synchronized keyword. Inner class has a separate monitor than the containing outer class.
- ▶ Java allows **Reentrant Synchronization**, that is, a thread can reacquire a lock it already owns. For example, a synchronized method can call another synchronized method.

Synchronization Example 1

- ▶ Example of a race condition: `Account.java`, `TestAccount.java`
- ▶ Thread safe version using `synchronized` keyword:
`SynchronizedAccount.java`

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- ▶ The method `notifyAll()` wakes up all waiting threads instead of just one waiting thread.

Example with `wait()/notify()`

```
class MyThing {
    synchronized void waiterMethod() {
        // do something
        // we need to wait for the notifier to do something
        // give up the lock, put calling thread to sleep
        wait();
        // continue where we left off
    }

    synchronized void notifierMethod() {
        // do something
        // notifier the waiter that we've done it
        notify();
        //do more things
    }

    synchronized void relatedMethod() {
        // do some related stuff
    }
}
```

Synchronization Example 2: Producer/Consumer Problem

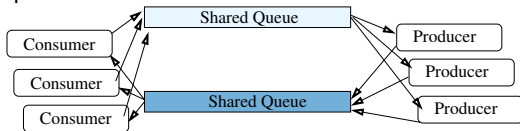
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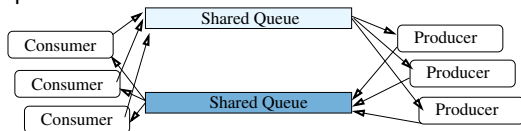
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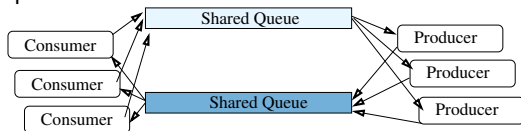
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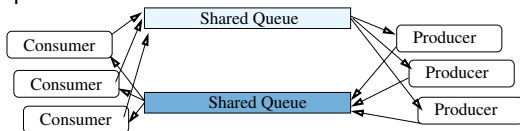
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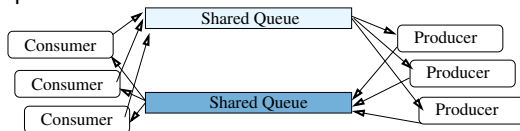
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- ▶ Example: `SharedQueue.java`, `Producer.java`, `Consumer.java`, `PC.java`
- ▶ The **Producer/Consumer** or a **Thread Pool pattern** is a widely used one for multi-threaded applications as well as in servers and clients.

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- ▶ Are the threads really simulating ping pong? We need them to exchange an object over the network!

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- ▶ Example: **ThreadGroupExample.java**

Thread Pool

- ▶ **Thread Pool**: A number of threads are created to perform a number of tasks, which are usually organized in a queue. Typically, there are many more tasks than threads.
- ▶ Java provides a thread pool via the **Executor** interface from the `java.util.concurrent` package.

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public static <T> Set<T> synchronizedSet(Set<T> s);  
public static <T> List<T> synchronizedList(List<T> list);  
public static <K,V> Map<K,V> synchronizedMap(Map<K,V> m);  
public static <T> SortedSet<T> synchronizedSortedSet(SortedSet<T> s);  
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A collection created in this fashion is every bit as thread-safe as a normally synchronized collection, such as a `Vector`.

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For more details, see:

<http://docs.oracle.com/javase/tutorial/collections/implementations/wrapper.html>

- ▶ Brian Goetz, Tim Peierls, Joshua Bloch and Joseph Bowbeer: *Java Concurrency in Practice*
- ▶ Lewis and Berg: *Multithreaded Programming with Java Technology*