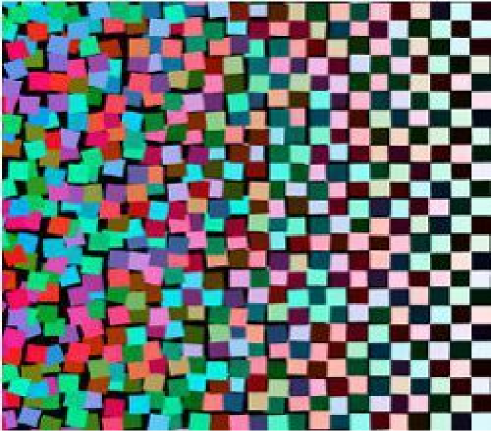
Consistency and Replication



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Tatiana Maslany in the show Orphan Black: The story of a group of clones that discover each other and the secret organization Dyad, which was responsible for their creation.

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  - Scaling in numbers
  - Scaling in geographical area
- Having multiple copies leads to the problem of consistency. When and how the copies are made consistent determines the price of replication.
- **Example**: Client caching of web pages in web browser gains performance for the client at the cost of consistency.

Replication as a Scaling Technique

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- How to keep the replicas consistent? Use global ordering using Lamport timestamps or use a coordinator. This may require a lot of communication for a large system.
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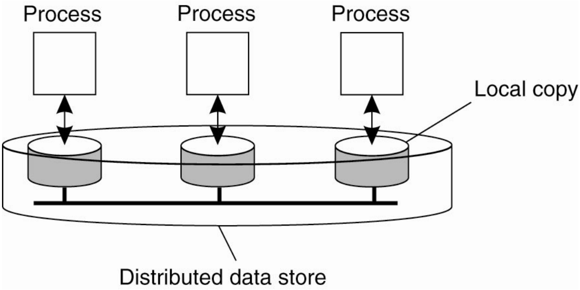
- A range of consistency models are available.

Data Store



- Examples:
  - distributed shared memory
  - distributed shared database
  - distributed file system

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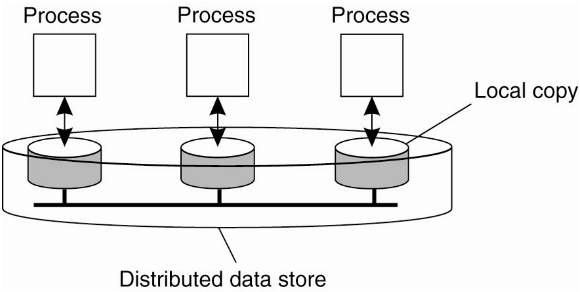

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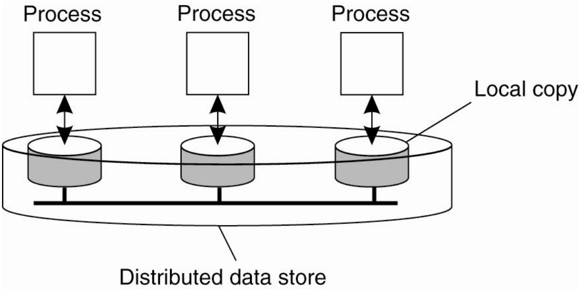

- Examples:
  - distributed shared memory
  - distributed shared database
  - distributed file system
- Each process has a local or nearby copy of the whole store or part of it.
- A data operation is classified as a write when it changes the data, and is otherwise classifies as a read operation.

What is a Consistency Model?

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- Models with minor restrictions are easy to use while models with major restrictions are more difficult to use. But then easy models don't perform as well. So we have to make trade-offs.

Consistency Models

- Data-centric consistency models
  - Continuous consistency
  - Sequential consistency
  - Causal consistency
  - Entry Consistency with Grouping operations
- Client-centric consistency models
  - Eventual consistency
  - Monotonic reads
  - Monotonic writes
  - Read your writes
  - Writes follow reads

Sequential Consistency (1)

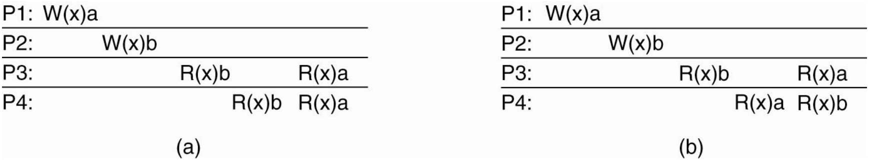
A data store is sequentially consistent when:

The result of any execution is the same as if the (read and write) operations by all processes on the data store were executed in **some sequential order** and the operations of each individual process appear in this sequence in the order specified by its program.

Sequential Consistency (2)

- Behavior of two processes operating on the same data item. The horizontal axis is time.

Sequential Consistency (3)



(a) A sequentially consistent data store.

(b) A data store that is not sequentially consistent.

Sequential Consistency (4)

Process P1	Process P2	Process P3
$x \leftarrow 1;$	$y \leftarrow 1;$	$z \leftarrow 1;$
print(y, z);	print(x, z);	print(x, y);

- Three concurrently-executing processes.

Sequential Consistency (5)

x ← 1;	x ← 1;	y ← 1;	y ← 1;
print(y, z);	y ← 1;	z ← 1;	x ← 1;
y ← 1;	print(x, z);	print(x, y);	z ← 1;
print(x, z);	print(y, z);	print(x, z);	print(x, z);
z ← 1;	z ← 1;	x ← 1;	print(y, z);
print(x, y);	print(x, y);	print(y, z);	print(x, y);
Prints: 001011	Prints: 101011	Prints: 010111	Prints: 111111
Signature: 001011	Signature: 101011	Signature: 110101	Signature: 111111
(a)	(b)	(c)	(d)

- Four valid execution sequences for the three processes. The vertical axis is time.
- Signature is output of P1, P2, P3 concatenated.

Causal Consistency (1)

For a data store to be considered causally consistent, it is necessary that the store obeys the following condition:

- **Causal Consistency**: Writes that are potentially causally related must be seen by all processes in the same order. Concurrent writes by different processes may be seen in a different order on different machines. Write by the same process are considered to also be causally related.
- Weaker than sequential consistency.

Causal Consistency (2)

P1:	W(x)a			W(x)c		
P2:		R(x)a	W(x)b			
P3:		R(x)a			R(x)c	R(x)b
P4:		R(x)a			R(x)b	R(x)c

- This sequence is allowed with a causally-consistent store, but not with a sequentially consistent store. Why?

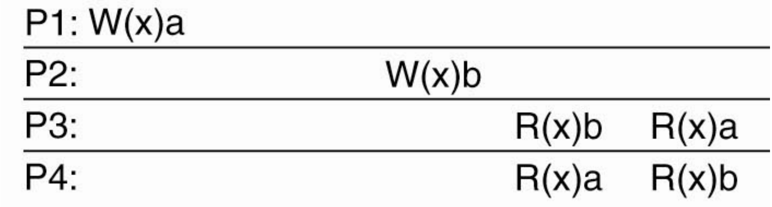
Causal Consistency (3)

P1:	W(x)a				
P2:		R(x)a	W(x)b		
P3:				R(x)b	R(x)a
P4:				R(x)a	R(x)b

(a)

- A violation of a causally-consistent store. Why?

Causal Consistency (4)



(b)

- A valid sequence of events in a causally-consistent store. Concurrent writes do not have to be globally ordered.

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- This can be done using *vector timestamps*.

Entry Consistency with Grouping Operations (1)

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- Exclusive access to a lock can succeed only if no other process has exclusive or nonexclusive access to that lock.
- Nonexclusive access to a lock is allowed only if any previous exclusive access has been completed, including updates on the lock's associated data.

This, in effect, is linearizing the usage of locks, adhering to sequential consistency.

Entry Consistency with Grouping Operations (2)

P1: L(x) W(x)a L(y) W(y)b U(x) U(y)

P2: L(x) R(x)a R(y) NIL

P3: L(y) R(y)b

- A valid event sequence for entry consistency.

Client-Centric Consistency Models

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- A special, important class of distributed data stores have the following characterization:
  - Lack of simultaneous updates (or when such updates happen, they can be easily resolved)
  - Most operations involve reading data.
- A very weak consistency model, called eventual consistency, is sufficient in such cases. This can be implemented relatively cheaply by client-centric consistency models.

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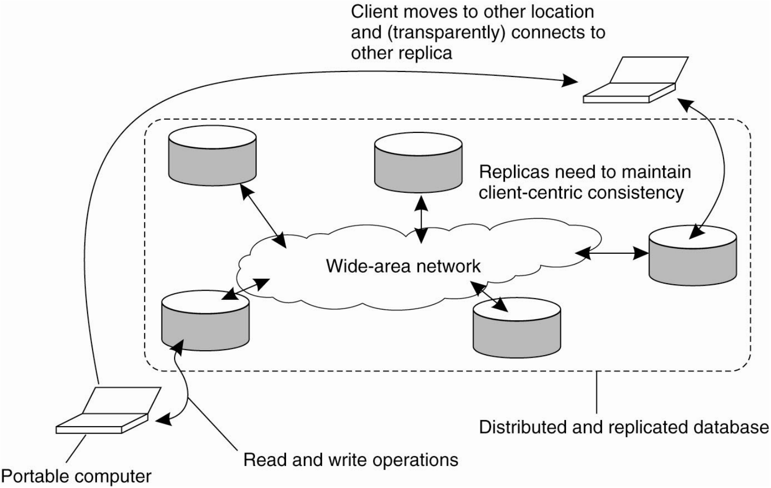
- If no updates take place, all replicas will gradually become consistent. This form of consistency is called eventual consistency.
- Eventual consistency essentially requires only that updates are guaranteed to propagate to all replicas. Write-write conflicts are easy to solve when assuming that only a small group of processes can perform updates. Eventual consistency is therefore often cheap to implement.

Eventual Consistency (2)



- The above problem can be alleviated by using *client-centric consistency*, which provides guarantees for a single client consistency in accessing the data store.

Client-centric Consistency Model Notation

- Version x_i of a data item is the result of a series of write operations that took place since initialization.
- We denote these operations by the write set $W(x_i)$
- By appending write operations to that series, we obtain another version x_j and say that x_j follows from x_i. We denote this as $W(x_i; x_j)$.
- If we do not know if x_j follows from x_i, we use the notation $W(x_i \mid x_j)$.

Monotonic Reads (1)

A data store is said to provide monotonic-read consistency if the following condition holds:

- If a process reads the value of a data item x, then any successive read operation on x by that process will always return that same value or a more recent value.
- But no guarantees on concurrent access by different clients.

Monotonic Reads (2)

L1:	$W_1(x_1)$		$R_1(x_1)$	
L2:		$W_2(x_1;x_2)$		$R_1(x_2)$

- The read operations performed by a single process P at two different local copies of the same data store. A monotonic-read consistent data store.
- L_1 and L_2 are two local data stores.
- $W_1(x_1)$: Process P_1 produces version x_1 without knowing anything about other versions.
- $W_2(x_1; x_2)$: Process P_2 produces version x_2 that follows from x_1.

Monotonic Reads (3)

L1:	$W_1(x_1)$	$R_1(x_1)$
L2:	$W_2(x_1 \| x_2)$	$R_1(x_2)$

- The read operations performed by a single process P at two different local copies of the same data store. A data store that does not provide monotonic reads.
- $W_2(x_1 \mid x_2)$ denotes that process P_2 produced version x_2 concurrently to version x_1, a potential write-write conflict.

Monotonic Writes (1)

In a monotonic-write consistent store, the following condition holds:

- A write operation by a process on a data item x is completed before any successive write operation on x by the same process.

Monotonic Writes (2)

L1: $W_1(x_1)$

L2: $W_2(x_1;x_2)$ $W_1(x_2;x_3)$

- A monotonic-write consistent data store.

Monotonic Writes (3)

L1:	$W_1(x_1)$	
L2:	$W_2(x_1 \mid x_2)$	$W_1(x_1 \mid x_3)$

- A data store that does not provide monotonic-write consistency.

Monotonic Writes (4)

L1: $W_1(x_1)$

L2: $W_2(x_1|x_2)$ $W_1(x_2;x_3)$

- A data store that does not provide monotonic-write consistency as we have both $WS(x_1 \mid x_2)$ and $WS(x_1 \mid x_3)$.

Monotonic Writes (5)

L1: $W_1(x_1)$

L2: $W_2(x_1|x_2)$ $W_1(x_1;x_3)$

- Consistent because $WS(x_1; x_3)$ although x_1 has apparently overwritten x_2.

Read Your Writes (1)

A data store is said to provide read-your-writes consistency, if the following condition holds:

- The effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same process.

Read Your Writes (2)

L1: $W_1(x_1)$

L2: $W_2(x_1;x_2)$ $R_1(x_2)$

- A data store that provides read-your-writes consistency.

Read Your Writes (3)

L1: $W_1(x_1)$

L2: $W_2(x_1|x_2)$ $R_1(x_2)$

- A data store that does not follow *Read-Your-Writes* consistency model.

Writes Follow Reads (1)

A data store is said to provide writes-follow-reads consistency, if the following holds:

- A write operation by a process on a data item x following a previous read operation on x by the same process is guaranteed to take place on the same or a more recent value of x that was read.

Writes Follow Reads (2)

L1: $W_1(x_1)$ $R_2(x_1)$

L2: $W_3(x_1;x_2)$ $W_2(x_2;x_3)$

- A *Writes-Follow-Reads* consistent data store.

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Replica Management

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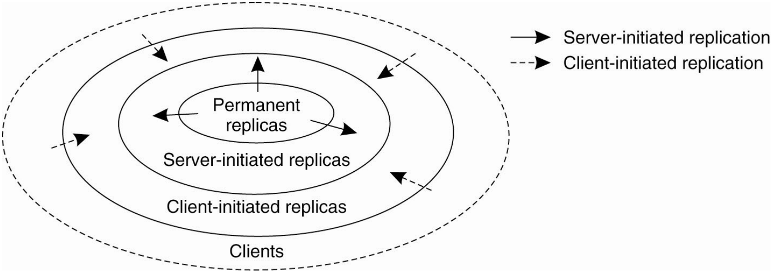
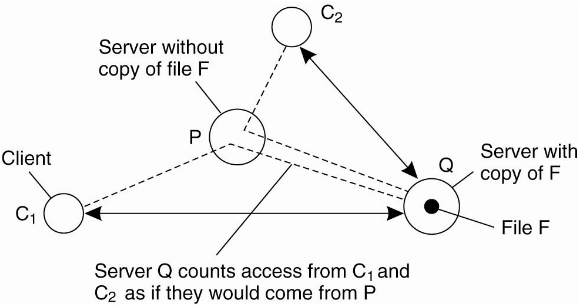
- A key issue for any distributed system that supports replication is to decide where, when and by whom replicas should be placed, and subsequently the mechanisms to use for keeping the replicas consistent. Two separate problems:
- *placing replica servers*: handled automatically by data centers
- *placing content*: permanent replicas, server-initiated, client-initiated

Content Replication and Placement



- The logical organization of different kinds of copies of a data store into three concentric rings.

Server-Initiated Replicas



- Counting access requests from different clients.
- Three possible actions: migrate, delete, replicate.

Client-Initiated Replicas

- Same as client caches.
- Used to improve access times.
- Sharing a cache between clients may or may not improve the hit rate. E.g. Web data, file servers
- Shared caches can be at departmental or organization level

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- A push-based approach is efficient when the read-to-update ratio is relatively high.
- Pull Protocol: A server or a client requests another server to send it any updates it has at the moment. A pull-based approach is efficient when the read-to-update ratio is relatively low.
- A comparison between push-based and pull-based protocols in the case of multiple-client, single-server systems.

Issue	Push-based	Pull-based
State at server	List of client replicas and caches	None
Messages sent	Update (and possibly fetch update later)	Poll and update
Response time at client	Immediate (or fetch-update time)	Fetch-update time

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- Using multicasting can be much more efficient than unicasting for the updates.

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  - Use Vector Clocks, where on server S_k, $RVC_k[i] = t_i$ implies that S_k has seen all writes that have been submitted to S_i up to time t_i.
  - Whenever a server notices that $t_k - RVC_k[i]$ is about to exceed a specified limit, it simply starts pulling in writes from S_i with a timestamp later than $RVC_k[i]$.
- Bounding ordering deviations
  - Can be bounded by specifying the maximal length of the queue of tentative writes.
  - When this is exceeded, the server no longer accepts writes, but will instead attempt to commit tentative writes by negotiating with other servers in which order its writes should be executed.

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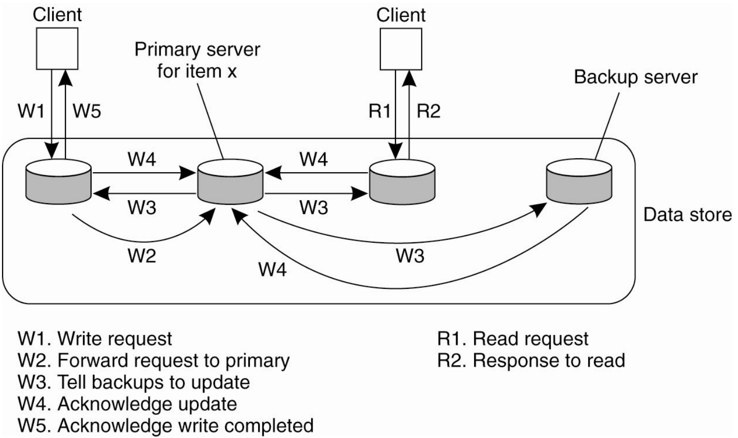
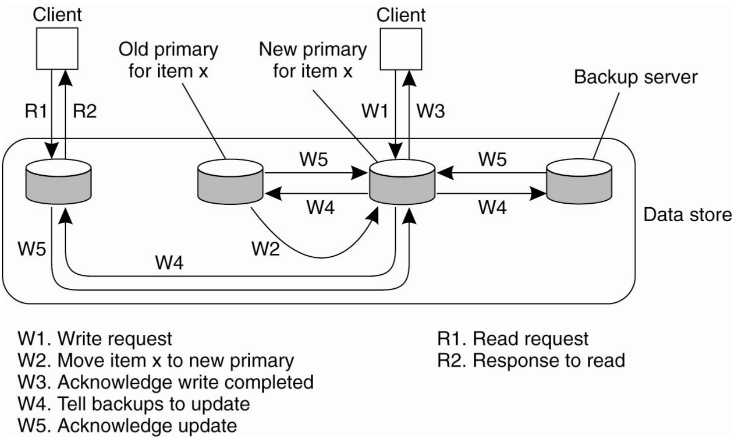
- Primary can be *fixed at a remote server* or write operations can be carried out *locally after moving the primary to the process* where the write operation was initiated.

Remote Write Protocol



- Also known as a primary-backup protocol. Implements sequential consistency.
- Non-blocking version: the primary acknowledges after updating its copy and informs backup servers afterwards

Local Write Protocol



- Primary-backup protocol in which the primary migrates to the process wanting to perform an update. Updates have to be propagated back to other replicas.

Examples of Local Write Primary-Backup Protocols

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- Quorum-based protocol: the write operation is done by majority voting.

Quorum-based Replicated-Write Protocol (1)

Quorum-based protocol

- To read a file with N replicas, a client needs to assemble a read quorum, an arbitrary collection of N_R servers. Similarly, to write to a file, a write quorum of at least N_W servers. The values of the quorums are subject to the following rules:
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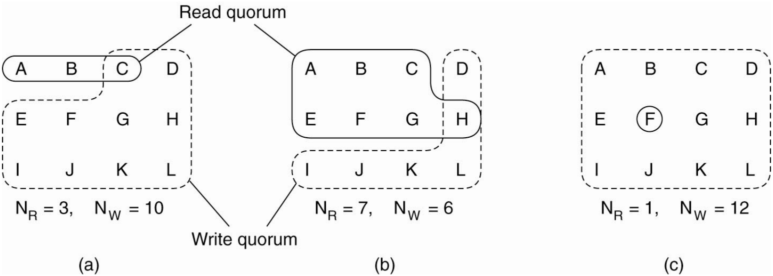
- Only one writer at a time can achieve write quorum.
- Every reader sees at least one copy of the most recent read (takes one with most recent version number or logical timestamp)

Quorum-based Replicated-Write Protocols (2)



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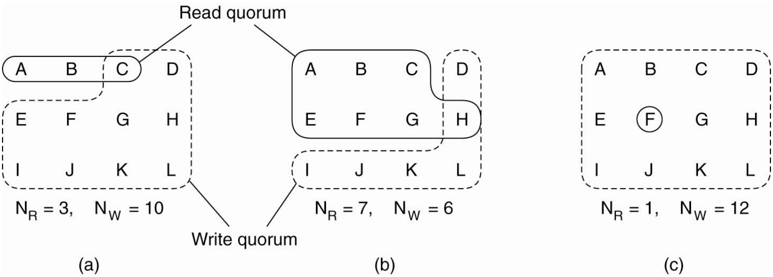

- Quorum-based Protocol. Three examples of the voting algorithm.
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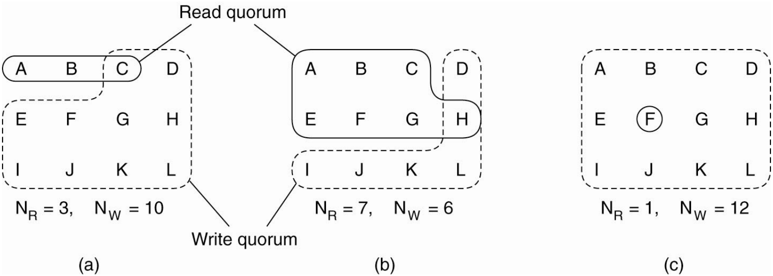

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  (c) A correct choice, known as ROWA (read one, write all).
- **In-class Exercise**. A file is replicated on 10 servers. List all combinations of read and write quorums that are permitted by the voting algorithm.

Quorum-based Replicated-Write Protocols (3)

- ROWA: $R = 1$, $W = N$
  Fast reads, slow writes
- RAWO: $R = N$, $W = 1$
  Fast writes, slow reads
- Majority: $R = W = N/2 + 1$
  Both moderately slow, but extremely high availability
- Weighted voting:
  Give more votes to "better" replicas

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- Each client tracks identifiers for two sets of writes:
  - read set: consists of writes relevant for the read operations performed by the client.
  - write set: consists of the writes performed by the client.

Implementing Client-Centric Consistency (2)

- *Implementing Monotonic Reads*: When a client performs a read operation at a server, that server is handed the client's read set to check whether all the identified writes have taken place locally. If not, it can contact other servers to be brought up to date. Or it can forward the request to another server where the write operations have already taken place.

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- After the read operation is performed, the write operations that are relevant to the read operation are added to the client's read set.
- How to determine exactly where the write operations in the client's read set have taken place? The write identifier can contain the server's identifier. Servers can be required to log their writes so they can replayed at another server. The client's can generate a globally unique identifier that is included in the write identifier.

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- *Implementing Monotonic Writes*: Whenever a client initiates a new write operation at a server, the server is handed over the client's write set. It then ensures that the identified write operations are performed first and in the correct order. After performing the new write operation, that write operation's identifier is added to the write set.

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- *Implementing Writes-Follow-Reads*: Bring the selected server up to date with the write operations in the client's read set, and then later adding the identifier for the write set to the write set, along with the identifiers of the read set (which have now become relevant for the write operation just performed).

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Implementing Client-Centric Consistency (4)

- Problems with the implementation: The read set and write set associated with each client can become very big.
- Improving efficiency: keep track of the sets only for the length of a session. This isn't sufficient. Why?
- The main problem is the representation of the read and write sets. It can be represented more efficiently by means of vector timestamps.

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- Whenever a server accepts a new write operation W, it assigns it a globally unique identifier along with a timestamp `ts`(W). That server is identified as the `origin`(W). Subsequent writes to that server get higher timestamps.

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- Whenever a client issues a request to read or write operation at a specific server, that server returns its current timestamp along with the results of the operation.

Representing Read/Write Sets Using Vector Timestamps (2)

- Read and write sets are represented by vector timestamps. For each session A, we construct a vector timestamp $\texttt{SVC}_A$ with $\texttt{SVC}_A[i]$ set equal to the maximum timestamp of all write operations in A that originate from server S_i.

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- The compactness is obtained by representing all observed writes originating from the same server through a single timestamp.

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- Once the operation is performed, the server S_i will return its current timestamp $\mathtt{WVC}_i$. At that point, $\mathtt{SVC}_A$ is adjusted to:

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Summary of Consistency Models (1)

Data-Centric Consistency Models

- **Sequential Consistency**: The result of any execution is the same as if the (read and write) operations by all processes on the data store were executed in **some sequential order** and the operations of each individual process appear in this sequence in the order specified by its program.
- **Causal Consistency**: Writes that are potentially causally related must be seen by all processes in the same order. Concurrent writes by different processes may be seen in a different order on different machines. Writes by the same process are also considered to be causally related.

Summary of Consistency Models (2)

Eventual consistency *essentially requires only that updates are guaranteed to propagate to all replicas*. **Client-Centric Consistency Models** that are used to implement eventual consistency.

- **Monotonic Reads**: If a process reads the value of a data item x, then any successive read operation on x by that process will always return that same value or a more recent value.
- **Monotonic Writes**: A write operation by a process on a data item x is completed before any successive write operation on x by the same process.
- **Read Your Writes**: The effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same process.
- **Writes Follow Reads**: A write operation by a process on a data item x following a previous read operation on x by the same process is guaranteed to take place on the same or a more recent value of x that was read.

References

- Eventually Consistent by Werner Vogels (CTO, Amazon)
- How eventual is eventual consistency? posted on Basho Blog
- 10 Lessons from 10 Years of Amazon Web Services by Werner Vogels (CTO, Amazon)