Networks and Network Programming

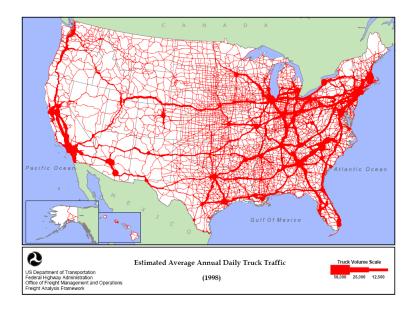
```
TOP 1220 90-36256 [PSH, ACK] Seq#122989 Ackn1 VannS4 Lenn1254 TSvalm3448276489 TSecr#120558
                                                                                                                                                 31.12.64.200
                                                                                                                                                        192,169,0,10
                                                                                                                                                     31.12.54.208
                                                                                                                                                                                                                                  TOP 66 36256-80 [ADX] SRIPLI SOCIEZAND NUTLICAN LIBRID INVALIGABLES CONTROL OF TOP 1200 00-36256 [P94, ADX] SRIPLI 27000 Adx | VirinG4 Librid 234 TSink | Adx | Adx | VirinG4 Librid | Adx | Adx
                                                                                                                                                                                                                                                                               1300 90-96256 TP9H, #CKT Sequit29159 Ackul Worlds Lenut254 TSvalu3488276574 TSecrut2655
                                                                                                                                                                                                                                                                                      1320 80-36256 [PSH, ACK] Sep+130413 Ack+1 WinvS4 Len+1254 TSHel+3446276629 TSecr+1305
                                                                                                                                                                                                                                                                                    86 36236-80 [ACK] Sept. Ack-33587 Nin-Load Len-0 TSVal-430050nse :cet-ro-excresor
1219 80-36256 [P94, ACK] Sept.23067 Ackn) Wind54 Len:1252 TSval-16446295654 TSecr-13052
66 36256-80 [ACK] Sept. Ack-320303 Nin-1444 Len-0 TSVal-430550102 TSecr-444627654
           Ethernet II, Src: ArrieGro_00:00:03 [00:00:ca:00:00:03], Det: Broadcom_dd:db:db (00:0a:f7:dd:db:db)
         Internet Protocol Version 4, Src: 31.12.64.208 (31.12.64.208), Dat: 192.168.0.10 (192.168.0.10)
                  Destination Port: 36256 (36256)
                  Sequence number: 192920 (relative sequence number)
                (Next sequence number: 134174 [relative sequence number)]
                  Acknowledgment number: 1 (relative ack number)
           + .... 0000 0001 1000 = Flags: 0x018 [P94, ACK]
                Window size value: 54
                [Calculated window size: 54]
           [Window size scaling factor: -1 (unknown)]
> Checksum: Os64bd [validation disabled]
  P CONTROLLO CONTROL (1974 de de de control canada control cont
⊕ M enp6s0: «Ive capture in progres... Packets: 1160 - Displayed: 1160 (100.0%)
```

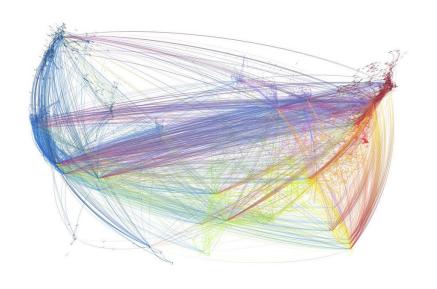
Wireshark screenshot of network traffic

Networking

- Hardware
- Protocols
- Software

The network is the computer. (John Gage)







Networking Options

Network type	maximum bandwidth	latency	
	(Mbits/second)	(microsecs)	
Fast Ethernet	100	200	
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Internet backbone uses specialized hardware and has speeds up to 500G/s. Experimental systems are available for much higher speeds.

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- ▶ If packets collide, then the machines choose a random number from the interval (0, k) and try again.
- On subsequent collisions, the value k is doubled each time, making it a lot less likely that a collision would occur again. This is an example of an exponential backoff protocol.

► The Ethernet packet has the format shown below.

Ethernet Packet Format

Preamble 10101010	Synch 11	Destination Address	Source Address	Туре	Data	Frame Check Sequence
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Use the program wireshark to watch Ethernet packets on your network live!

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- Custom complicated switching technology, or
- One big switch.

Network Topology Options

Hubs and Switches.

▶ Direct wire. Two machines can be connected directly by a Ethernet cable (usually a Cat 5e cable) without needing a hub or a switch. With multiple NICs per machine, we can create networks but then we need to specify routing tables to allow packets to get through. The machines will end up doing double-duty as routers.

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- ▶ Switches. Accepts packets, interprets destination address fields and send packets down only the segment that has the destination node. Allows half the machines to communicate directly with the other half (subject to bandwidth constraints of the switch hardware). Multiple switches can be connected in a tree or sometimes other schemes. The root switch can become a bottleneck. The root switch can be a higher bandwidth switch.

Switches can be managed or unmanaged. Managed switches are more expensive but they also allow many useful configurations. Here are some examples.

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- b Stackable, High bandwidth Switches. Stackable switches with special high bandwidth interconnect in-between the switches. For example, Cisco has 24-port Gigabit stackable switches with a 32 Gbits/sec interconnect. Up to 8 such switches can be stacked together. All the stacked switches can be controlled by one switch and managed as a single switch. If the controlling switch fails, the remaining switches hold an election and a new controlling switch is elected. Baystack also has stackable switches with a 40 Gbits/sec interconnect.

Network Interface Cards

► The Ethernet card, also known as the Network Interface Controller (NIC), contains the Data Link Layer and the Physical Layer (the two lowest layers of networking). Each Ethernet card has a unique hardware address that is know as its MAC address (MAC stands for Media Access Controller). The MAC address is usually printed on the Ethernet card.

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- Another issue to consider is that having multi-processor boards may cause more load on the network cards in each node. Certain network cards have multiple network processors in them, making them better candidates for multi-processor motherboards.

Networking Models

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- ► The ARPANET Reference Model (ARM) was the network model that led to the ISO OSI seven layer standardized model.
- ▶ ISO Open System Interconnection (OSI). A reference model for networking prescribes seven layers of network protocols and strict methods of communication between them. Most systems implement simplified version of the OSI model. The ARPANET Reference Model (ARM) can be seen as a simplified OSI model.

Network Models (contd.)

ISO	ARM	4.2 BSD Layers	Example
application	process	user programs/libraries	ssh
presentation	applications		
session		sockets	sock_stream
transport	host-host	protocols	TCP/IP
network	network interface	network interface	Ethernet driver
data link			
hardware	network hardware	network hardware	interlan controller

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 - HyperText Transfer Protocol (HTTP), File Transfer Protocol (FTP)
 - Java Remote Method Invocation (RMI), Distributed Component Object Model (DCOM)

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- The portion of the address that remains fixed within a network is called the network address and the remainder is the host address.
- ▶ The address with all 0's in the host address, for example 192.168.1.0, is the network address and cannot be assigned to any machine. The address with all 1's in the host address, for example 192.168.1.255, is the network broadcast address.

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- These addresses are permanently unassigned, not forwarded by Internet backbone routers and thus do not conflict with publicly addressable IP addresses.

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Typical communication domains:

J I		
domain type	symbolic name	address format
Unix domain	AF_Unix	pathnames
Internet domain	AF_INET	Internet address and port number

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- ▶ Raw sockets. Allows access to TCP, IP or Ethernet protocol

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- ▶ close(...), shutdown(...): Close both or one end of a socket respectively.

Client-Server Setup Using Sockets

Server side	Client side
socket()	
bind()	socket()
listen()	connect()
accept()	
read/write	read/write
close/shutdown	

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 - port 37 is reserved for getting the time from a system
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 - port 21 is used by the FTP (File Transfer Protocol) client/server
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Servers

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- ► A multi-threaded server accepts connections and passes them off to their own threads for processing.
- ▶ A multi-process server forks off a copy of itself after a connection to handle the client, while the original server process goes back to accept further connections.

Client/Server Examples in C

The examples are in folder sockets-C/tcp

► Single-threaded server: timeserver.c and client timeclient.c

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- Multi-process server: tcpserver.c and client tcpclient.c
- Multi-threaded server: Left as an exercise....
- Note that read/write on sockets is slightly different than read/write on files. A read/write on a socket may return a count less than asked for. This is not an error since with sockets, the buffer in the kernel may be full. We can just keep calling read/write until the right amount of data has been read or written.

- ▶ Use netstat -ni to find information on the network interfaces.
- ▶ Use netstat -rn to see the routing table.
- ► Use netstat -nap to see the processes that are using specific interfaces and ports. You need to be superuser to be able to see complete process information. Nice way of determining who has a port bound up!
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- ► Use 1sof -w -n -i tcp:<port num> gives the pid of the process using that port so you can kill a process that didn't clean up or properly release the port.

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- ► Use lsof -w -n -i tcp:<port num> gives the pid of the process using that port so you can kill a process that didn't clean up or properly release the port.
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- ▶ Use wireshark to watch network packets in real time! You will need superuser access to be able to use wireshark fully.

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- Now the server and the client can read/write to the streams associated with the sockets.
- ► Always open OutputStream before InputStream on a socket to avoid deadlock and synchronization problems.

Client Example (Java)

```
try {
    Socket server = new Socket("www.party.com",1234);
    OutputStream out = server.getOutputStream();
    InputStream in = server.getInputStream();
    out.write(42); //write a byte
    //write a newline or carriage return delimited string
    PrintWriter pout = new PrintWriter(out, true);
    pout.println("Hello!");
   //read a byte
    Byte response = in.read();
   //read a newline or carriage return delimited string
    BufferedReader bin = new BufferedReader (new
   InputStreamReader(in));
    String answer = bin.readLine();
    //send a serialized Java object
    ObjectOutputStream oout = new ObjectOutputStream(out);
    oout.writeObject(new java.util.Date());
    oout.flush():
    server.close();
 catch (IOException e) { System.err.println(e);}
```

Server Example (Java)

```
try { // meanwhile, on www.party.com...
    ServerSocket listener = new ServerSocket(1234);
    while (!finished) {
        Socket client = listener.accept();
        InputStream in = client.getInputStream();
        OutputStream out = client.getOutputStream();
        Byte someByte = in.read(); //read a byte
        //read a newline/carriage return delimited string
        BufferedReader bin =
             new BufferedReader(new InputStreamReader(in));
        String someString = bin.readLine():
        out.write(42); //write a byte
        PrintWriter pout = new PrintWriter(out, true);
        pout.println("Goodbye!");
        //read a serialized Java object
        ObjectInputStream oin = new ObjectInputStream(in);
        Date date = (Date) oin.readObject();
        client.close():
        //...
    listener.close():
 catch (IOException e) { System.err.println(e);}
```

See the folder and subfolders in sockets.

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- The SecurityManager can impose arbitrary restrictions on applications as to what hosts they may or may not talk to, and whether they can listen for connections.
- ► A server could run a proxy that lets the application communicate indirectly with anyone the server likes and allows. How would you design a proxy server?

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public class Request implements java.io.Serializable {}
public class DateRequest extends Request {}
public class WorkRequest extends Request {
    public Object execute() {return null;}
}
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- ► A sample work request class: MyCalculation.java
- ► The Server.java, Client.java that ties it all together.

Running the Object Server/Client

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- Start the client anywhere on the Internet. E.g. on restaurant.endofuniverse.net: java Client plainoldearth.net 1234
- ▶ Note that the server machine must have all the classes that the client has in order to be able to execute them on the client's behalf. That may be an unreasonable assumption since you may want to serve many kinds of clients without having to store all their classes.

Socket and ServerSocket Options

- ServerSocket and Socket classes have several useful options.
- ► For example: we can set a timeout on a socket, we can set the receive buffer sizes, etc.
- See examples ServerSocketOptions.java and ClientSocketOptions.java in the package tcp.socketoptions.

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- Example: UdpServer1.java and UdpClient1.java in package tcp.udp.

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- ► Requests/Methods in the HTTP Protocol:

 GET <pathname> HTTP/x.y (e.g. GET /sample.html HTTP/1.0

 HEAD <pathname> (same as GET except only metadata is returned)

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HEAD <pathname> (same as GET except only metadata is returned)
POST <string>
```

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Response from server:

```
HTTP-Version status-code reason-phrase <CR><LF>
```

Selected status codes in the HTTP Protocol

Status codes:

```
200 OK
201 Created
```

```
301 Moved permanently
```

```
305 Use Proxy
```

307 Temporary redirect

400 Bad request (bad syntax)

401 Unauthorized

402 Payment required

403 Forbidden

404 Not found

500 Internal server error

501 Not implemented

503 Service unavailable

505 HTTP version not supported

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- ► Then use a web browser and point to localhost:5005 or use the provided client program.
 - java tinyhttpd.Client localhost 5005
- Warning! This web server will serve files without any protection from a system!

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catch (SecurityException e) { pout.println("403 Forbidden");}

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- ► Add the following after the catch for FileNotFoundException.

 catch (SecurityException e) { pout.println("403 Forbidden");}
- ▶ Now, recompile and run the server as follows.

```
java -Djava.security.manager
-Djava.security.policy=mysecurity.policy TinyHttpd 5005
```

Adding a Custom Security Manager to TinyHttpd

```
class TinyHttpdSecurityManager extends SecurityManager {
 public void checkAccess(Thread g) { };
 public void checkListen(int port) { };
  public void checkLink(String lib) { };
 public void checkPropertyAccess(String key) { };
 public void checkAccept(String host, int port) { };
 public void checkWrite(FileDescriptor fd) { };
 public void checkRead(FileDescriptor fd ) { };
 public void checkRead(String s) {
    if (new File(s).isAbsolute() || (s.indexOf("..") != -1))
      throw new SecurityException("Access to file: " + s
                                           + " denied.");
//add the following to the TinyHttpd at the start of the
//main method but after creating the ServerSocket
System.setSecurityManager(new TinyHttpdSecurityManager());
```

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```
66.249.69.87 - - [28/Jan/2016:23:16:04 -0700]

"GET /~amit/teaching/555/cs455-555.html HTTP/1.1" 200 7343 "-"

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- ▶ Use scalable I/O with java.nio package.

Scalable I/O with java.nio package

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Scalable I/O with java.nio package

- Nonblocking and selectable network communications are used to create services that can handle very high volumes of simultaneous client requests.
- Starting one thread per client request can consume a lot of resources. One strategy is to use nonblocking I/O operations to manage a lot of communications from a single thread. The second strategy is to use a configurable pool of threads, taking advantage of machines with many processors.
- ➤ The java.nio package provides selectable channels. A selectable channel allows for the registration of a special kind of listener called a selector that can check the readiness of the channel for operations such as reading and writing or accepting or creating network connections.

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SelectableChannel channelA = ...;
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▶ Then, we register the channels.

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- Once one or more channels are registered with the Selector, we can perform a select operations by using one of the select() methods.

```
int readyCount = selector.select(); //block until one channel is ready
int readyCount = selector.selectNow(); // returns immediately
int readyCount = selector.select(50); // timeout of 50 milliseconds
while (selector.select(50) == 0);
```

Checking for ready channels

Once select() comes back with a non-zero ready count, then we can get the set of ready channels from the Selector with the selectedKeys() method and iterate through them.

```
Set readySet = selector.selectedKeys();
for (Iterator itr = readySet.iterator(); itr.hasNext();) {
    SelectionKey key = (SelectionKey) itr.next();
    itr.remove(); // remove the key from the ready set
    // use the key in the application
}
```

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- Example: LargerHttpd.java, HttpdConnection.java, ClientQueue.java.
- A single thread executes the main loop that accepts new connections and checks the readiness of existing client connections for reading or writing.
- ▶ Whenever a client needs attention, it places the job in a queue where a thread from our thread pool waits to service it.
- ▶ Run it as follows (on one line!):

```
java -Djava.security.manager
-Djava.security.policy=mysecurity.policy
LargerHttpd <port> <maxthreads>
```

Exercises

- Modify the object server example to be single threaded.
- Modify the object server example to use the Executor interface.
- Create a EC2 virtual machine using AWS (Amazon Web Services) and run the object server on it. Connect with your laptop or a lab machine as a client.
- Convert the object server example to send JSON objects instead of Java objects.
- Write a Ping Pong programs using sockets and an object to represent the ping-pong ball. Are the two programs playing ping pong servers? Or are they clients?