

# CGM: An Enhanced Mechanism for Streaming Data Collection with Local Differential Privacy (LDP)

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(submitted to VLDB 2021)

# Motivation for LDP



collects data from individuals



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Daily number of visits to the merchant website



collects data from individuals



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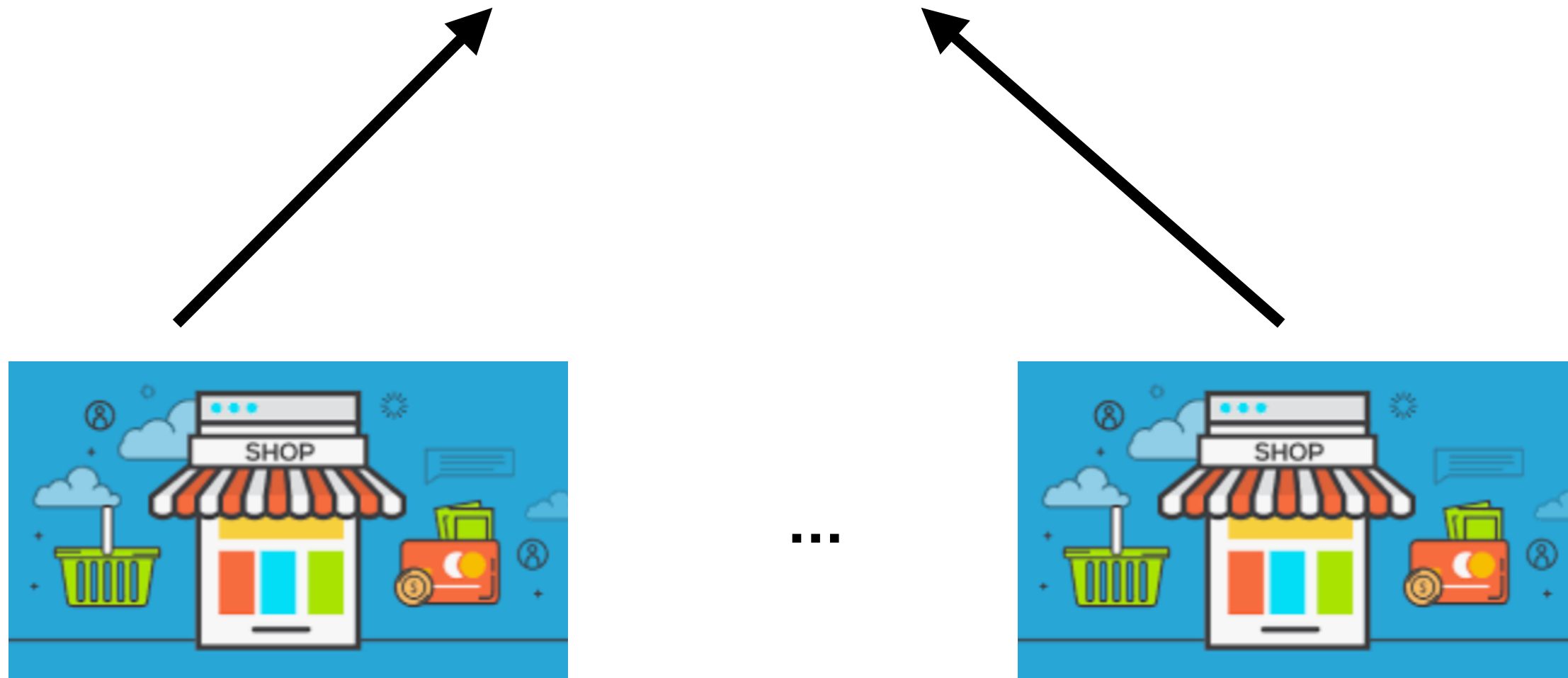


The percentage of time the taxi stays in a certain area within 30 minutes.

# Motivation for LDP



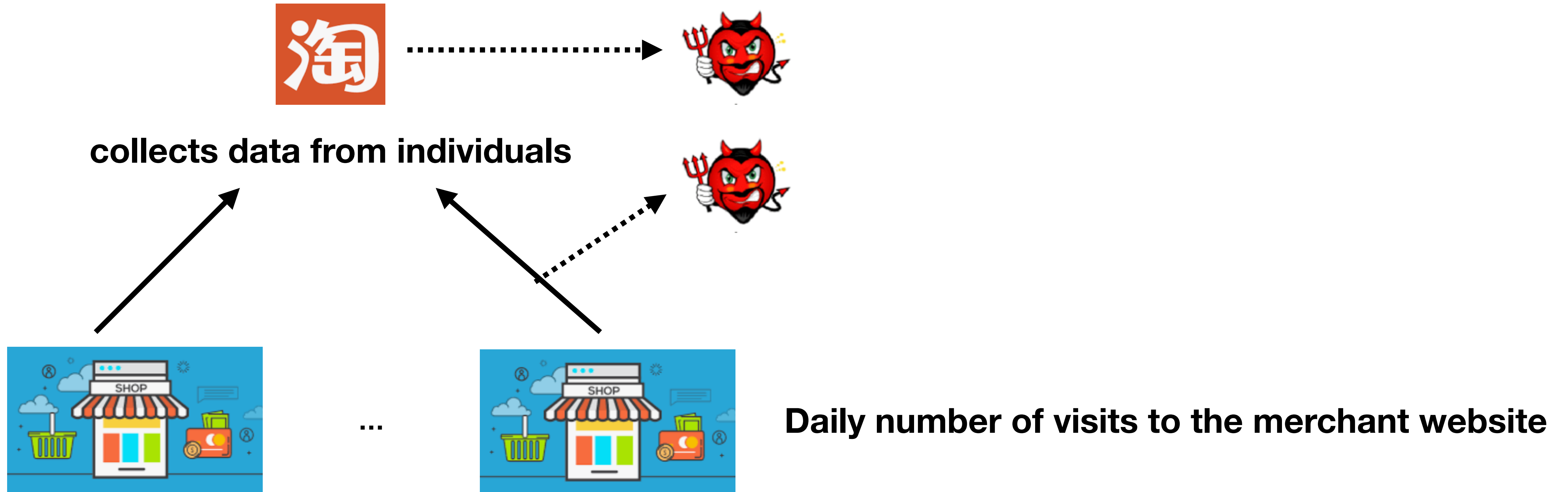
**collects data from individuals**



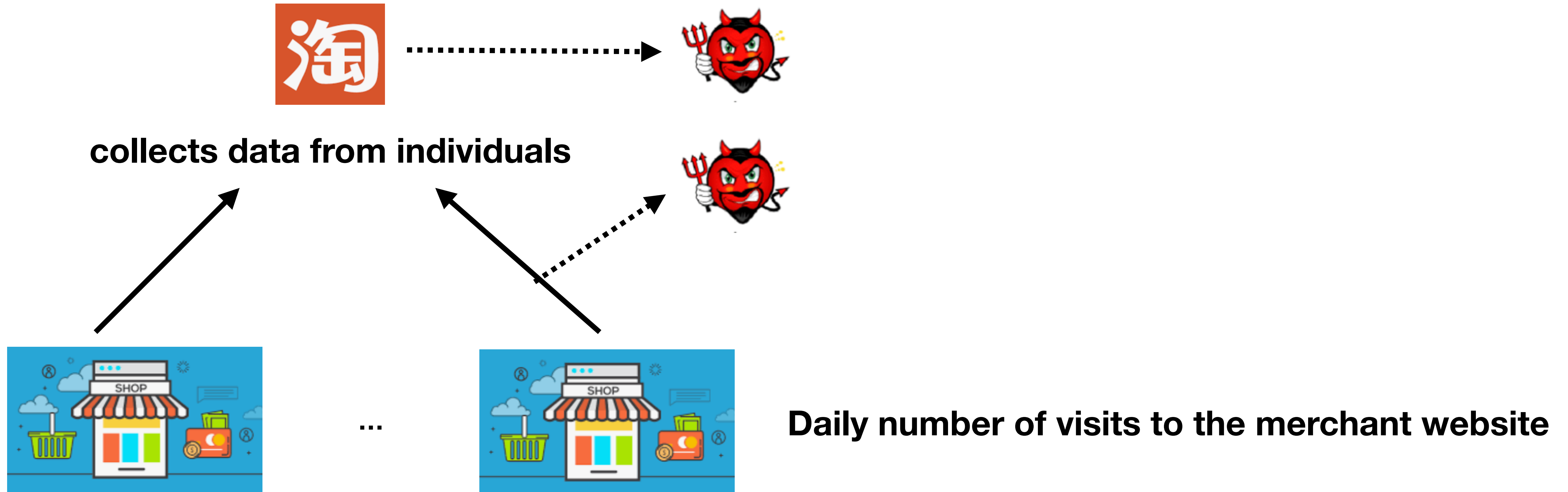
**Daily number of visits to the merchant website**

**The aggregator wants the data for e.g. traffic monitoring, setting up business plans, etc..**

# Motivation for LDP



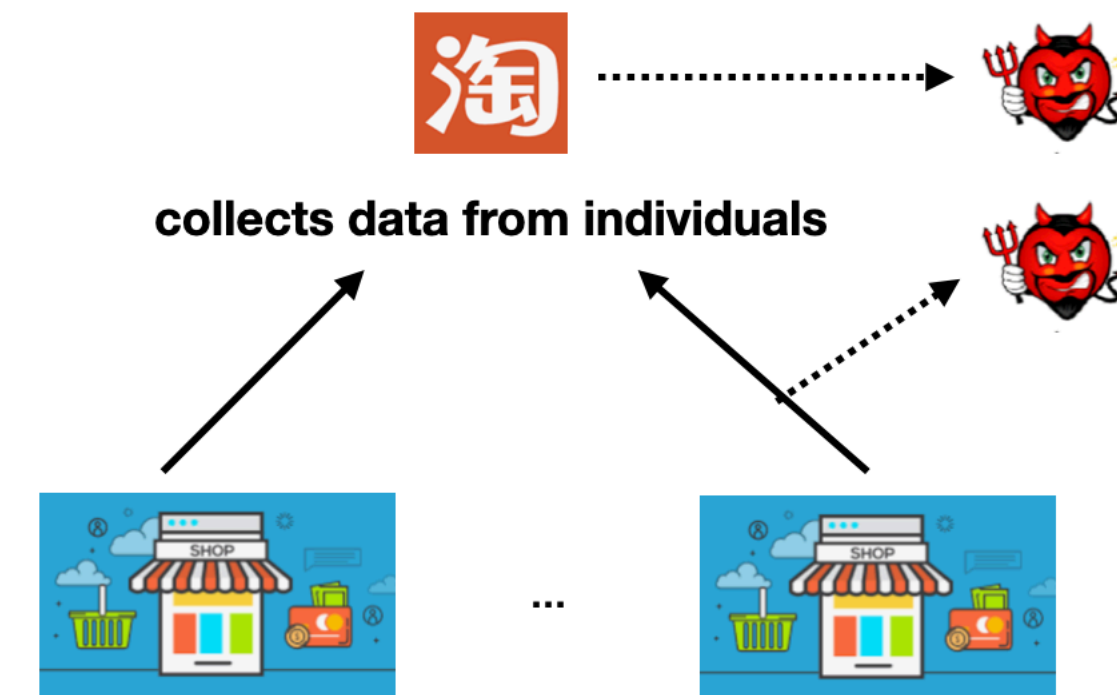
# Motivation for LDP



**The individual does not trust the communication channel, or the aggregator. Sending sensitive data directly leaks privacy. E.g., the daily number of visits is an indicator for the merchant's revenue.**

# Motivation for LDP

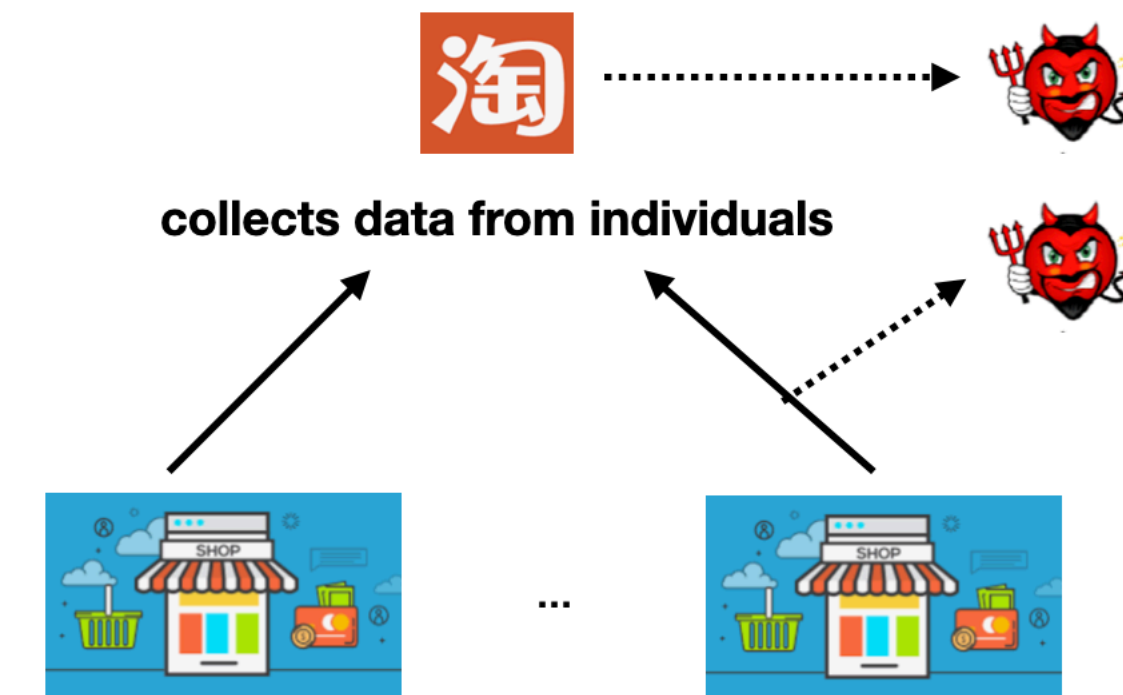
Each individual sends her data in a way that:



- A. From the data of an individual student, the aggregator can not learn the exact information.
- B. Upon the collection of all answers (e.g., hundreds of thousands of merchants), the aggregator can get a good estimate for the whole population, e.g., average daily number of visits.



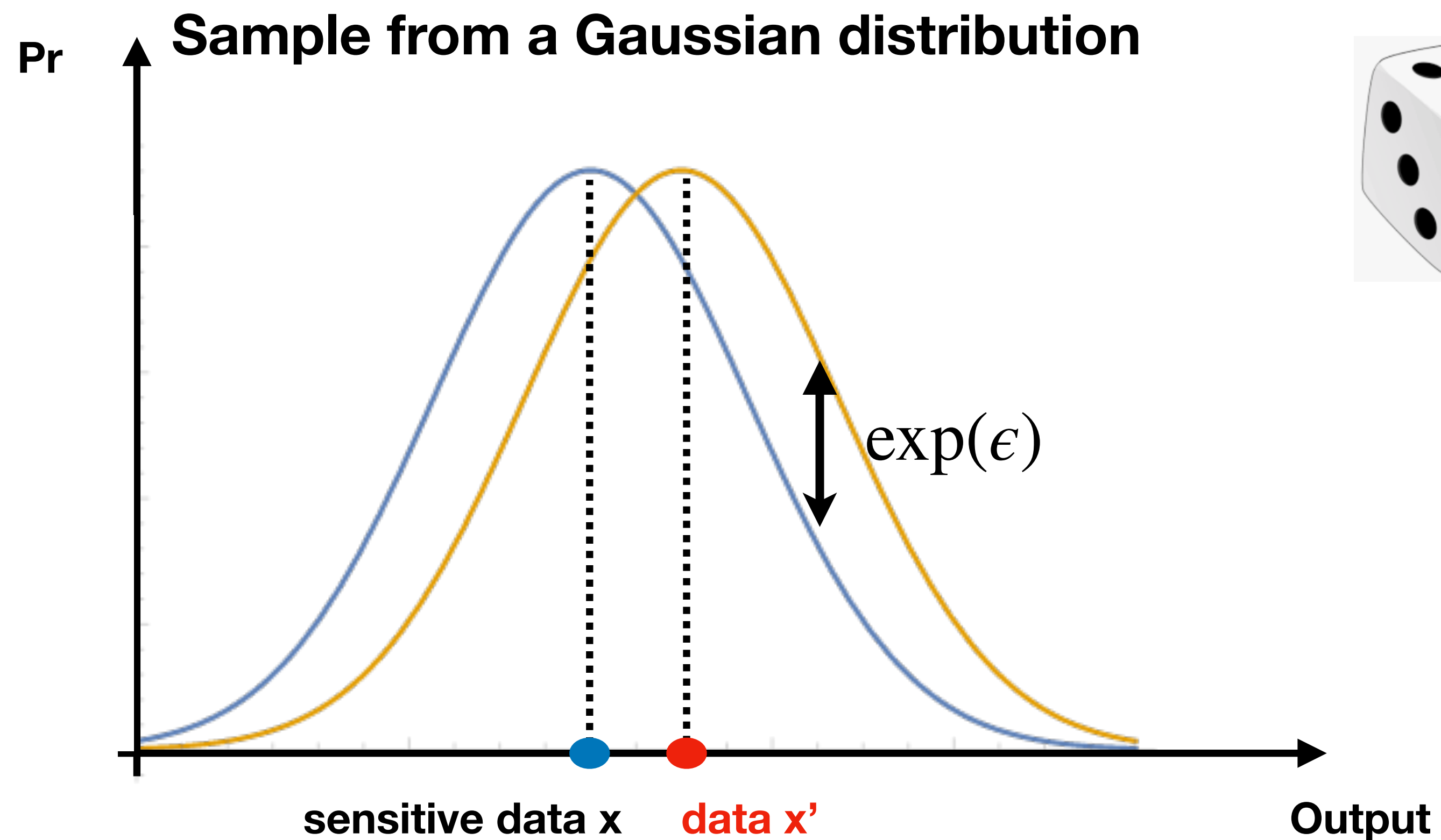
# Motivation for LDP



**Each individual sends her data in a way that:**

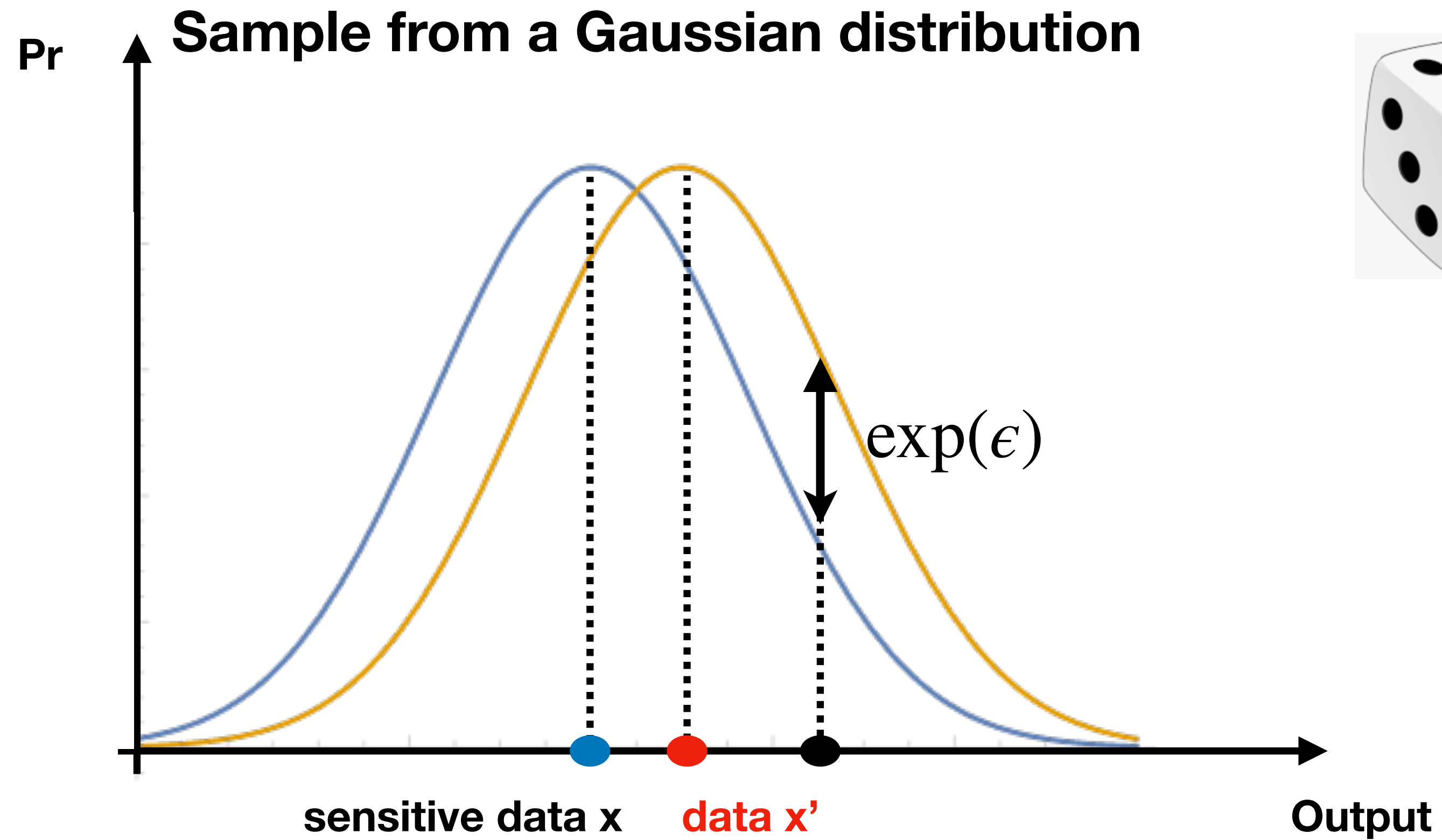
- A. From the data of an individual student, the aggregator can not learn the exact information.
  - B. Upon the collection of all answers (e.g., hundreds of thousands of merchants), the aggregator can get a good estimate for the whole population, e.g., average daily number of visits.
- A. ensures the privacy for every individual.
- B. ensures the accuracy of the statistic.

# LDP Mechanism [DMNS '06]



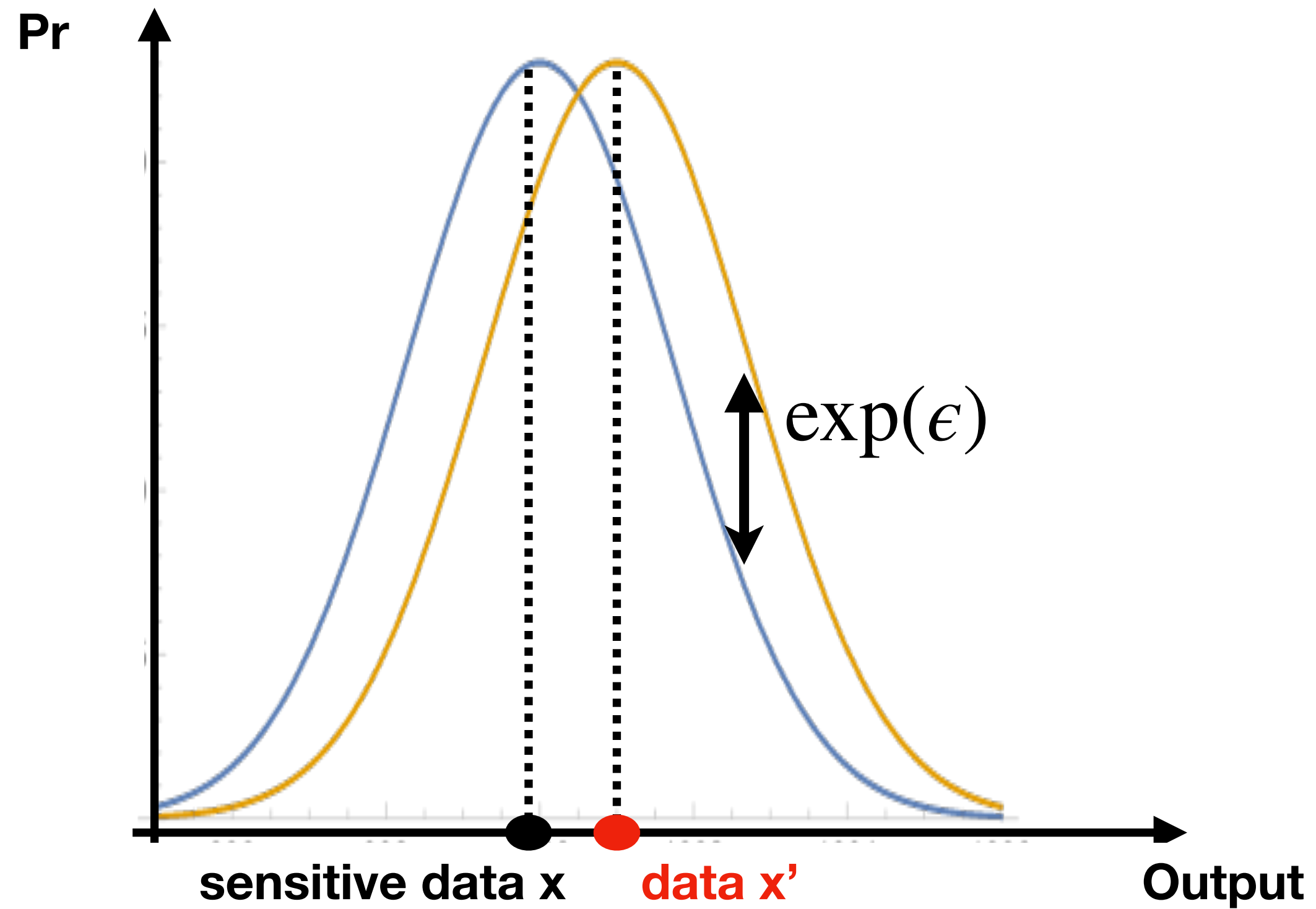


# LDP Mechanism

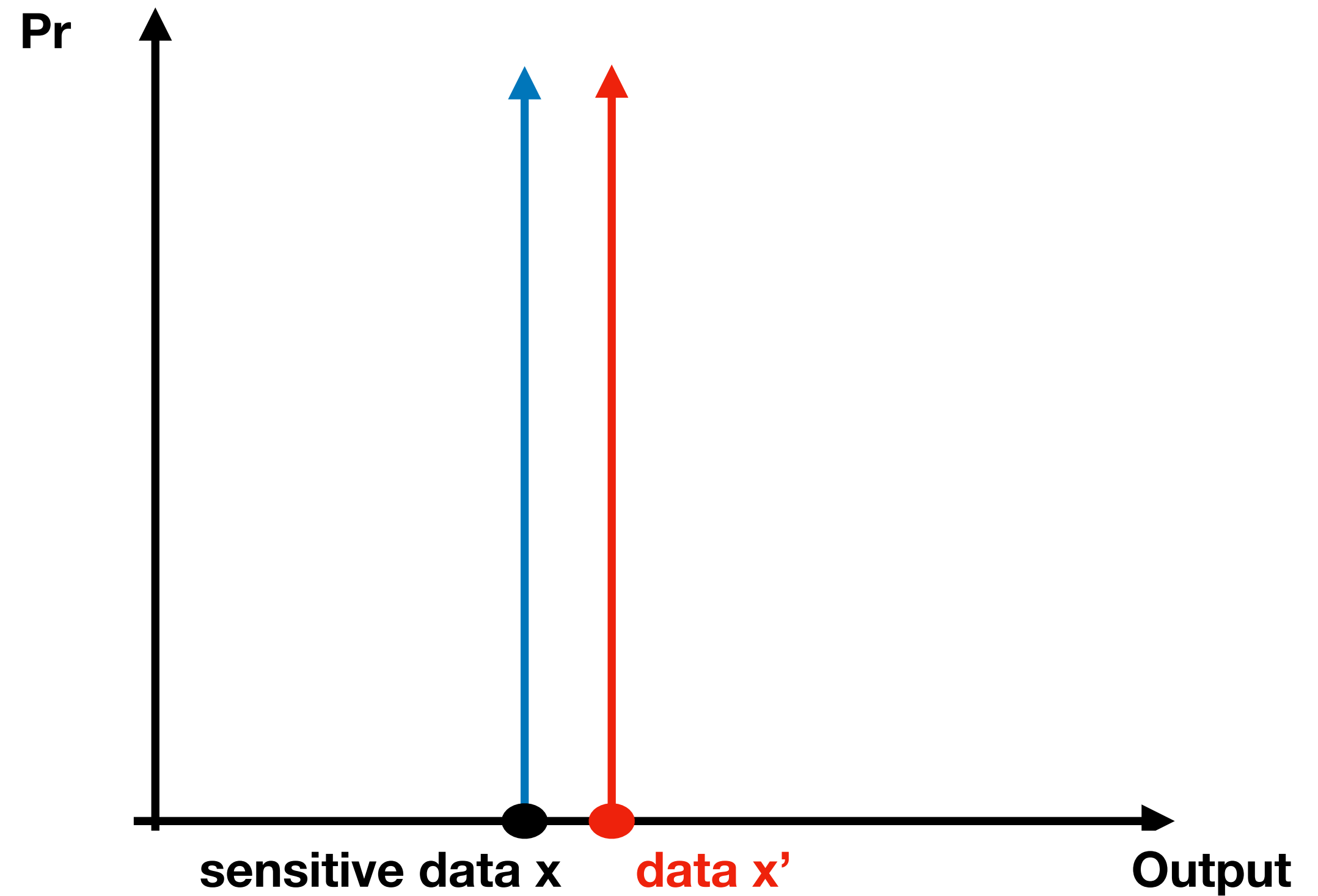


# LDP Mechanism

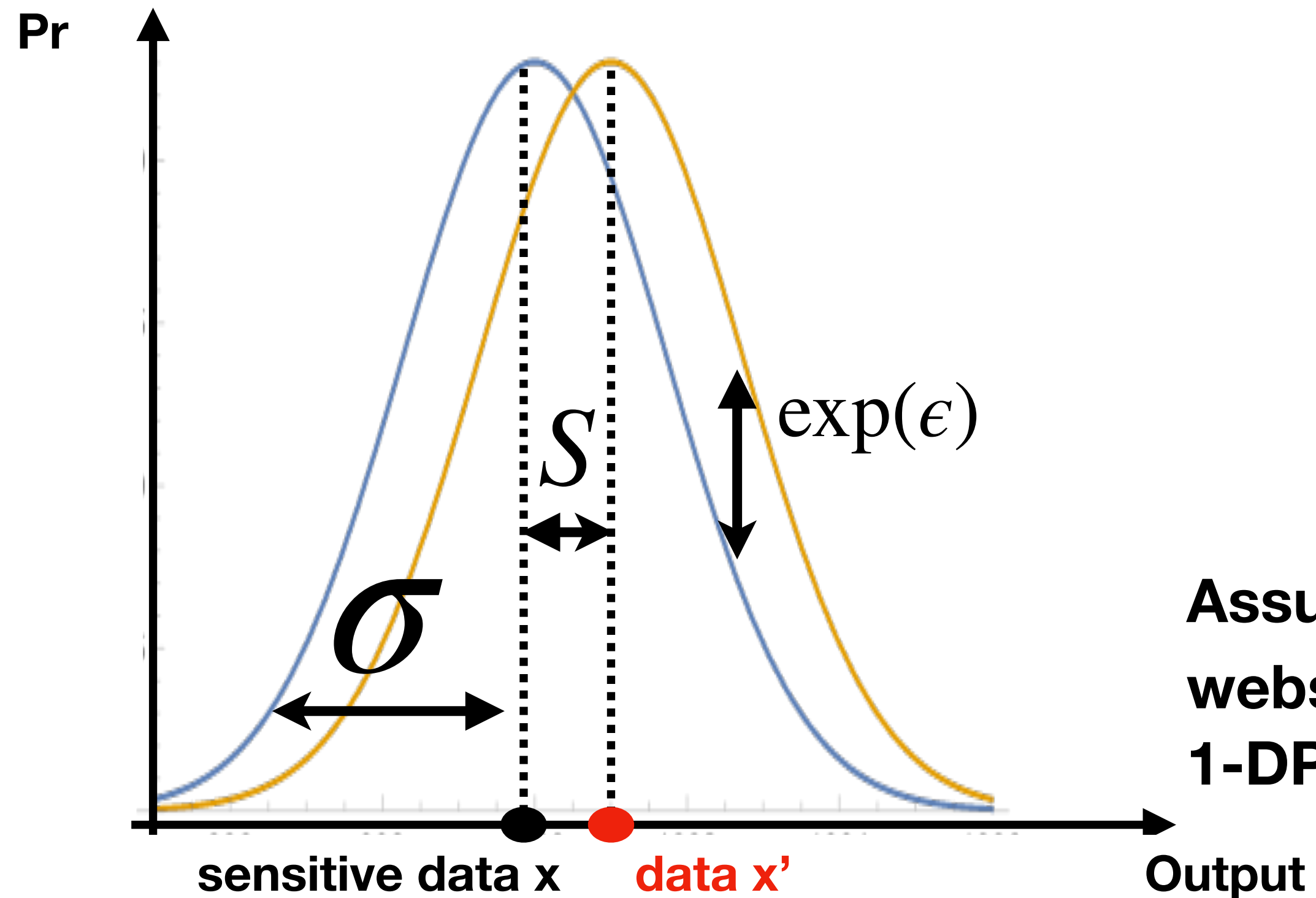
Smaller  $\epsilon$  means stronger privacy



This mechanism ensures no privacy



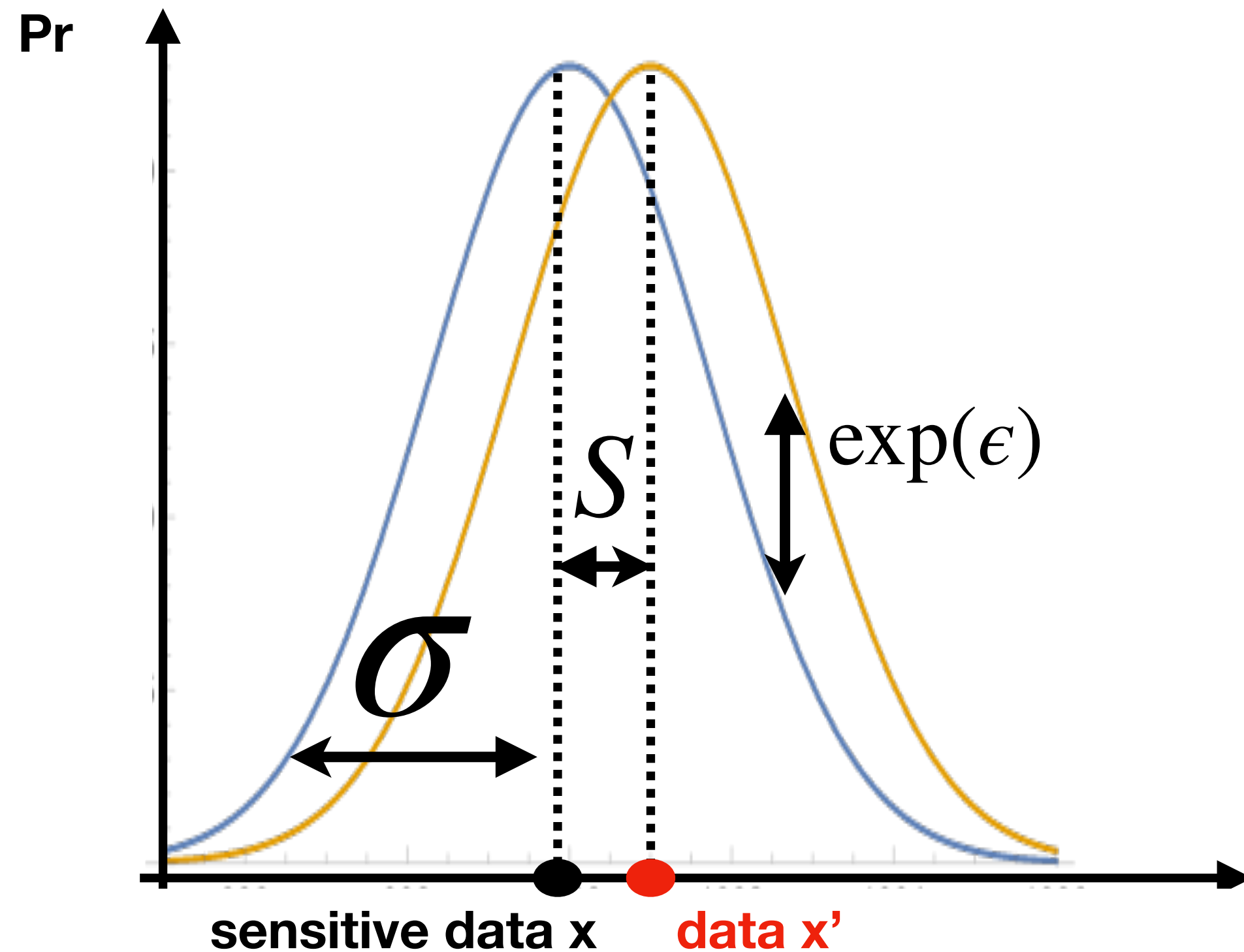
# LDP Mechanism



$$\text{Utility } \sigma \approx \frac{\text{Sensitivity } S}{\text{Privacy } \epsilon}$$

Assuming the maximum number of daily visits to a merchant website is 10000, adding Gaussian noise of scale 10000 ensures 1-DP.

# LDP Mechanism



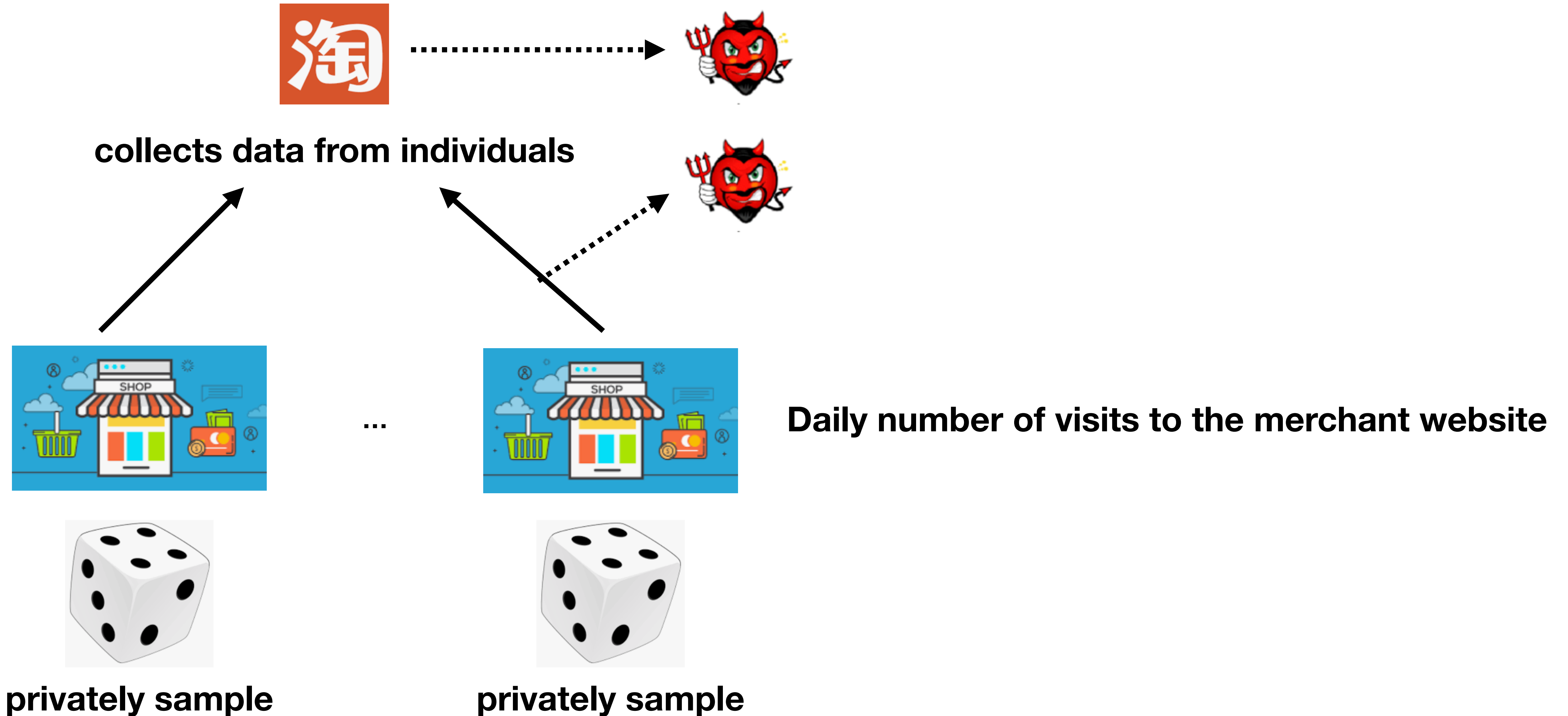
$$\text{Utility } \sigma \approx \frac{S \text{ Sensitivity}}{\epsilon \text{ Privacy}}$$

Assuming the maximum number of daily visits (i.e., the sensitivity) to a merchant website is 10000, adding Gaussian noise of scale 10000 ensures 1-DP.

Is the noise too large?

No. If there are  $10^4$  merchants, then on average the noise is only 100, introducing only 1% error to the average number of daily visits of all merchants.

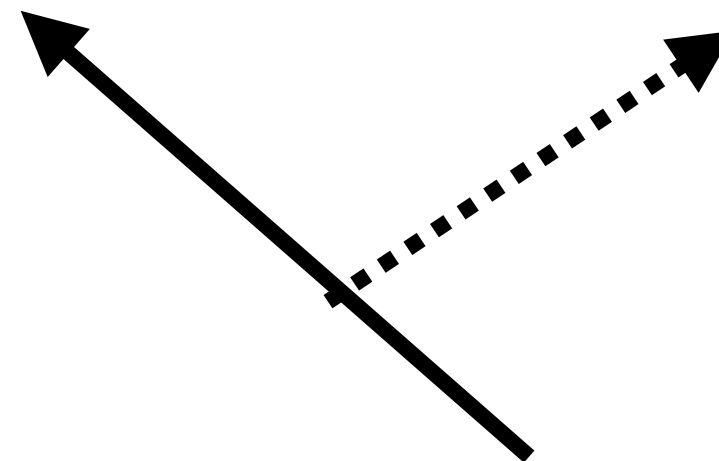
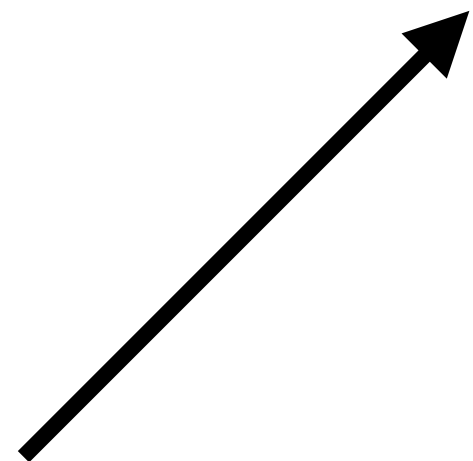
# Recap



# Streaming LDP Mechanism



collects data from individuals



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Daily number of visits to the merchant website



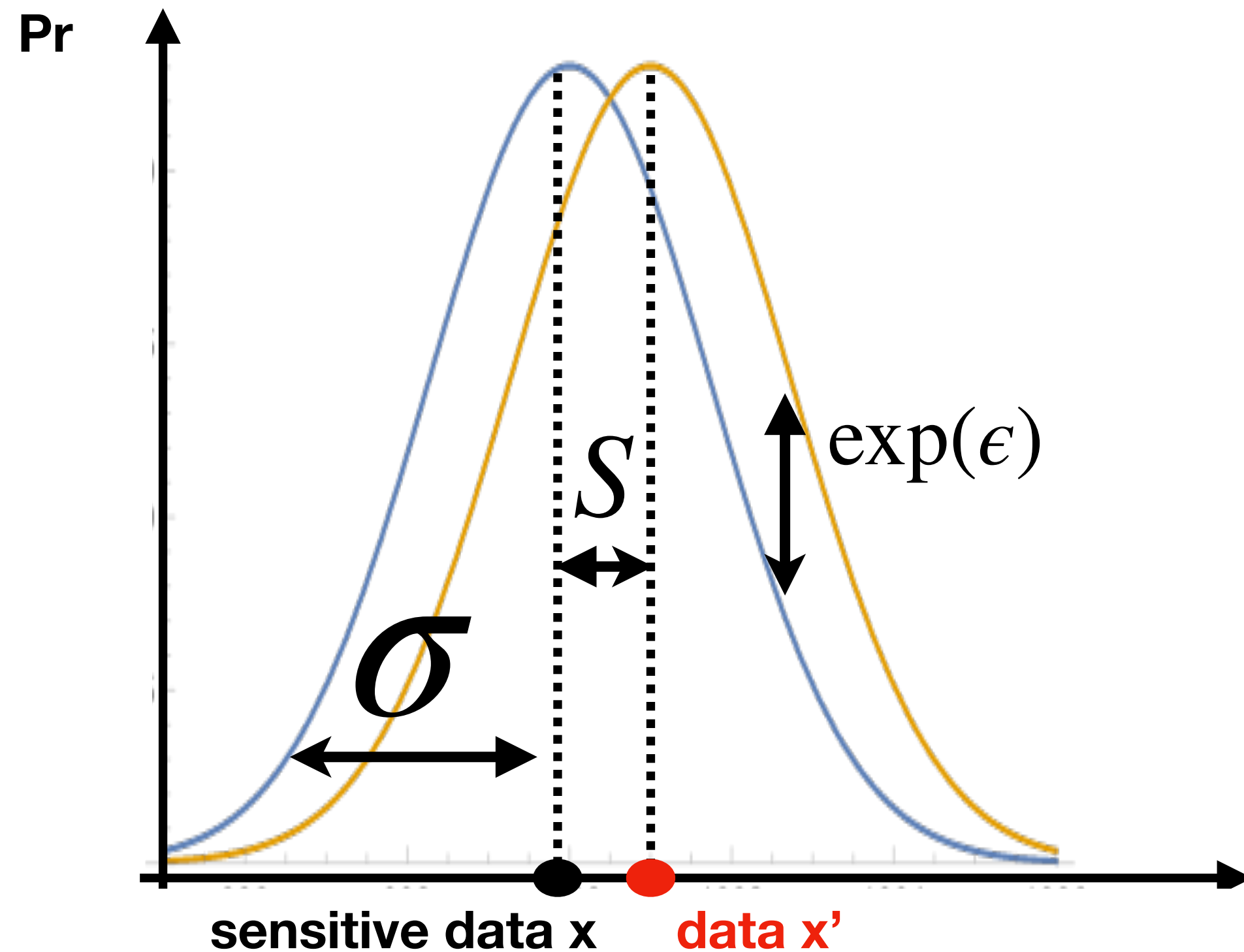
privately sample



privately sample



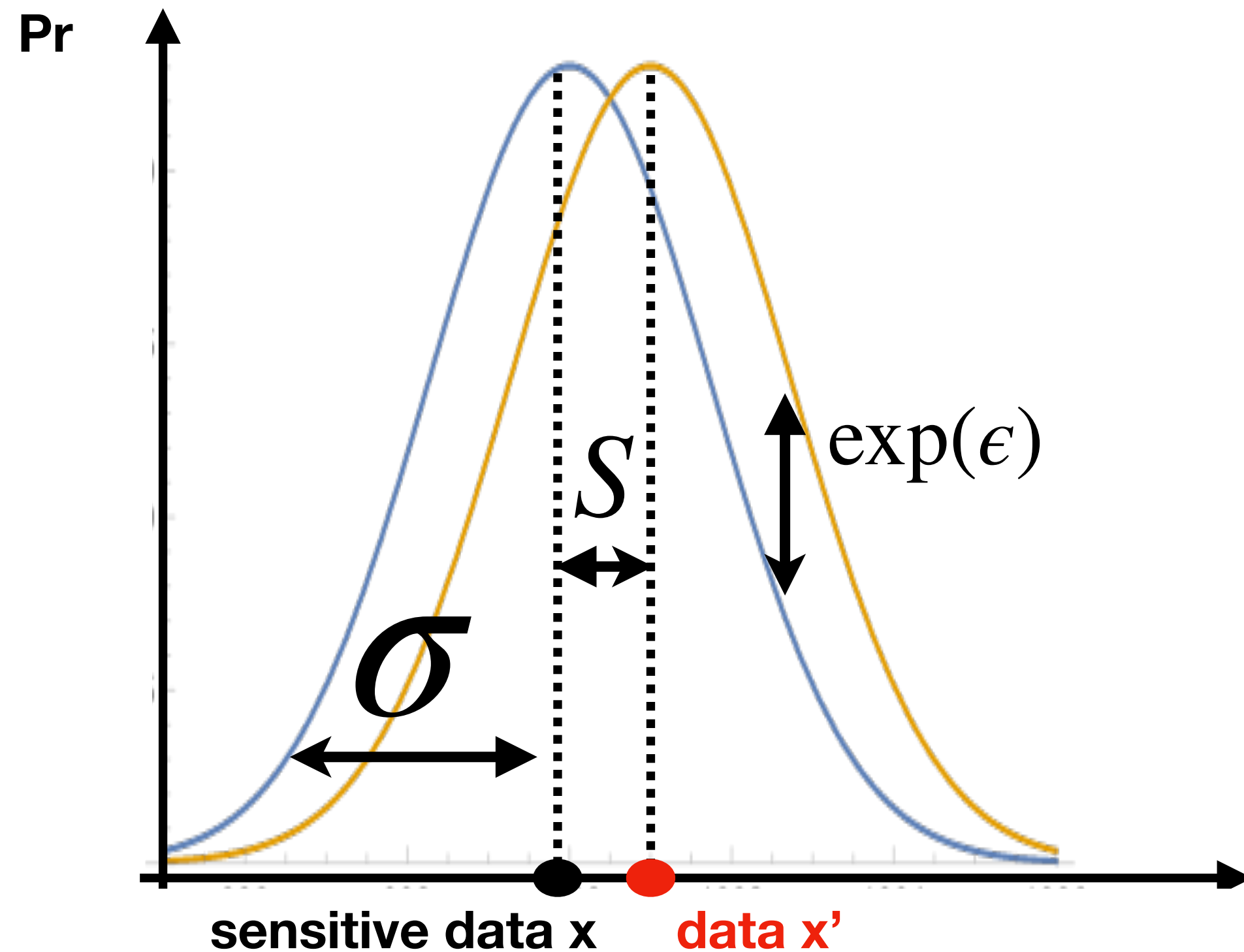
# Streaming LDP Mechanism



Repeatedly sample everyday.



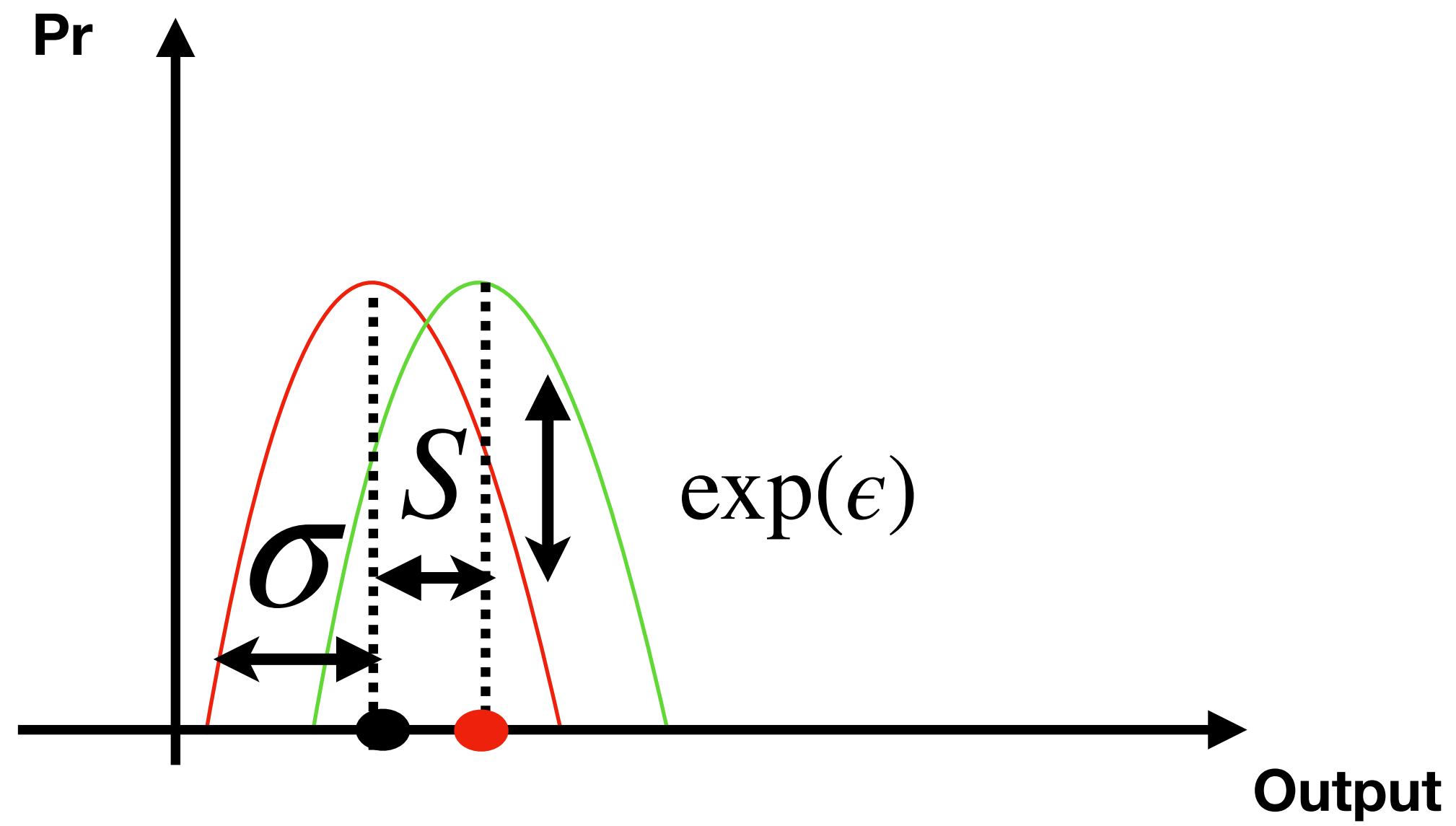
# Streaming LDP Mechanism



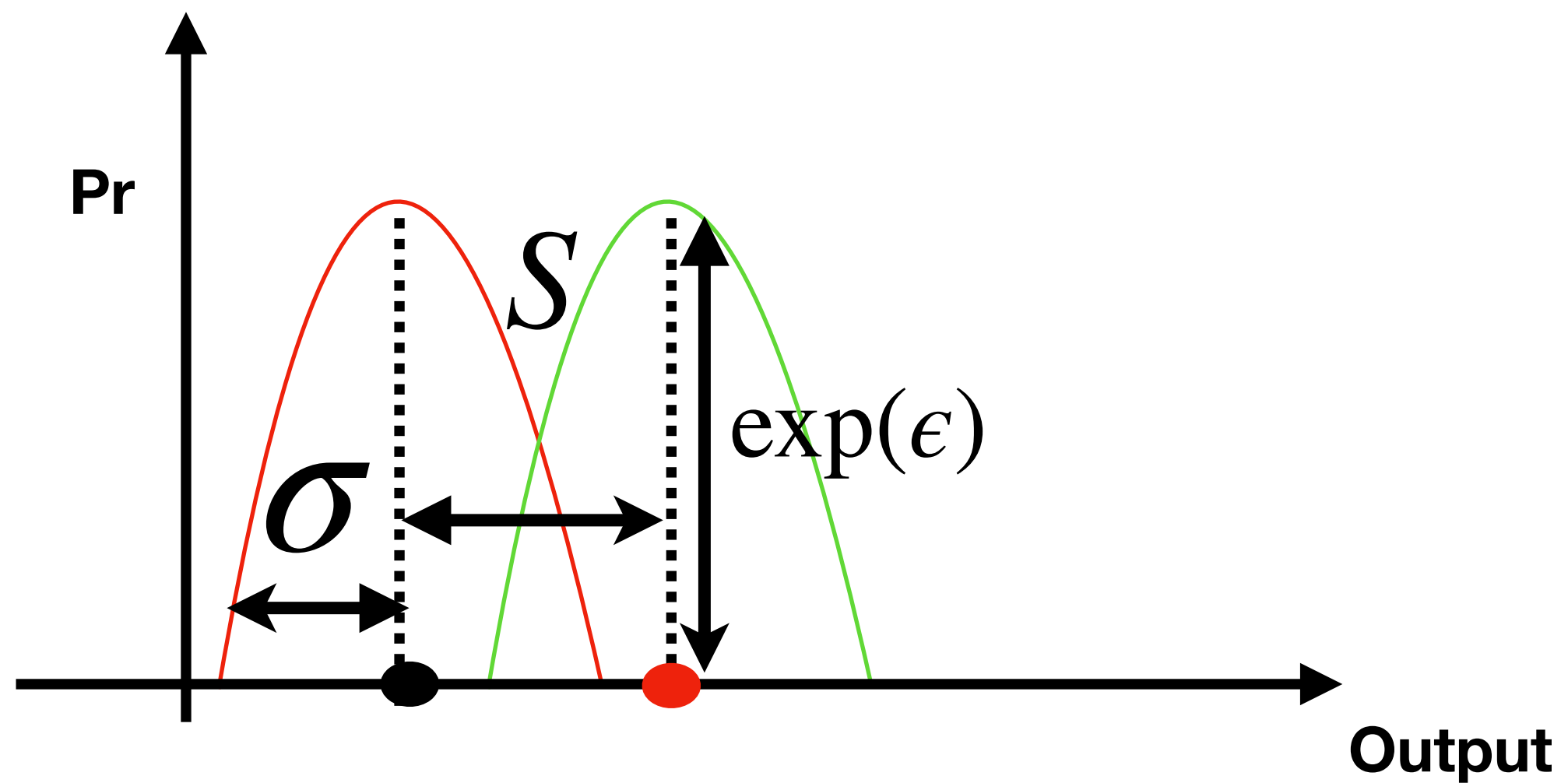
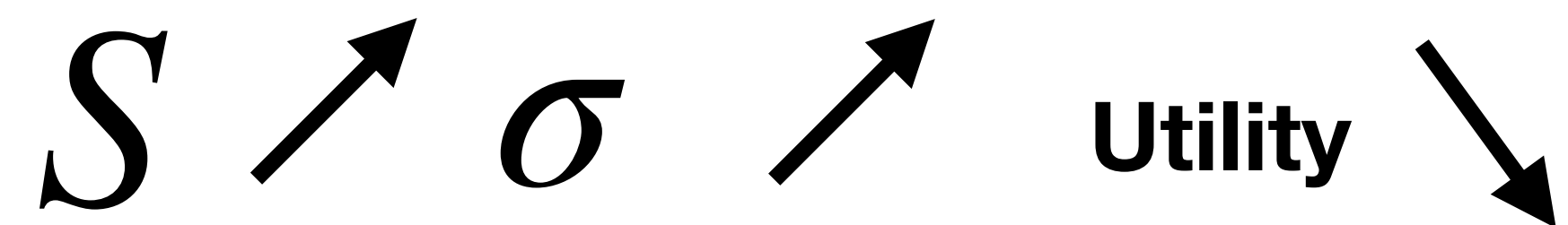
~~Repeatedly sample everyday.~~

There is a better way!

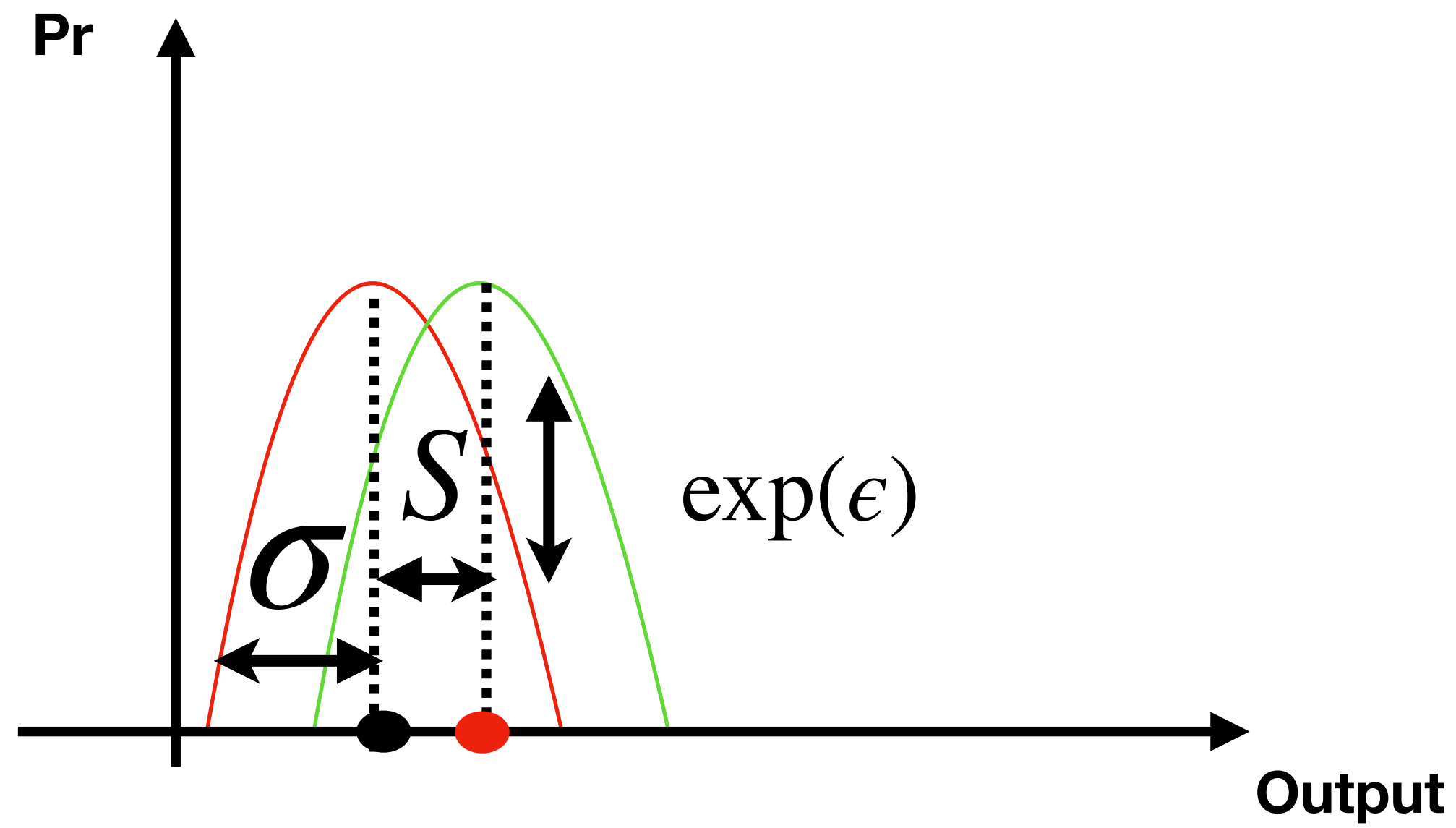
# LDP Mechanism



$$\text{Utility } \sigma \approx \frac{S}{\epsilon} \text{ Sensitivity Privacy}$$



# LDP Mechanism

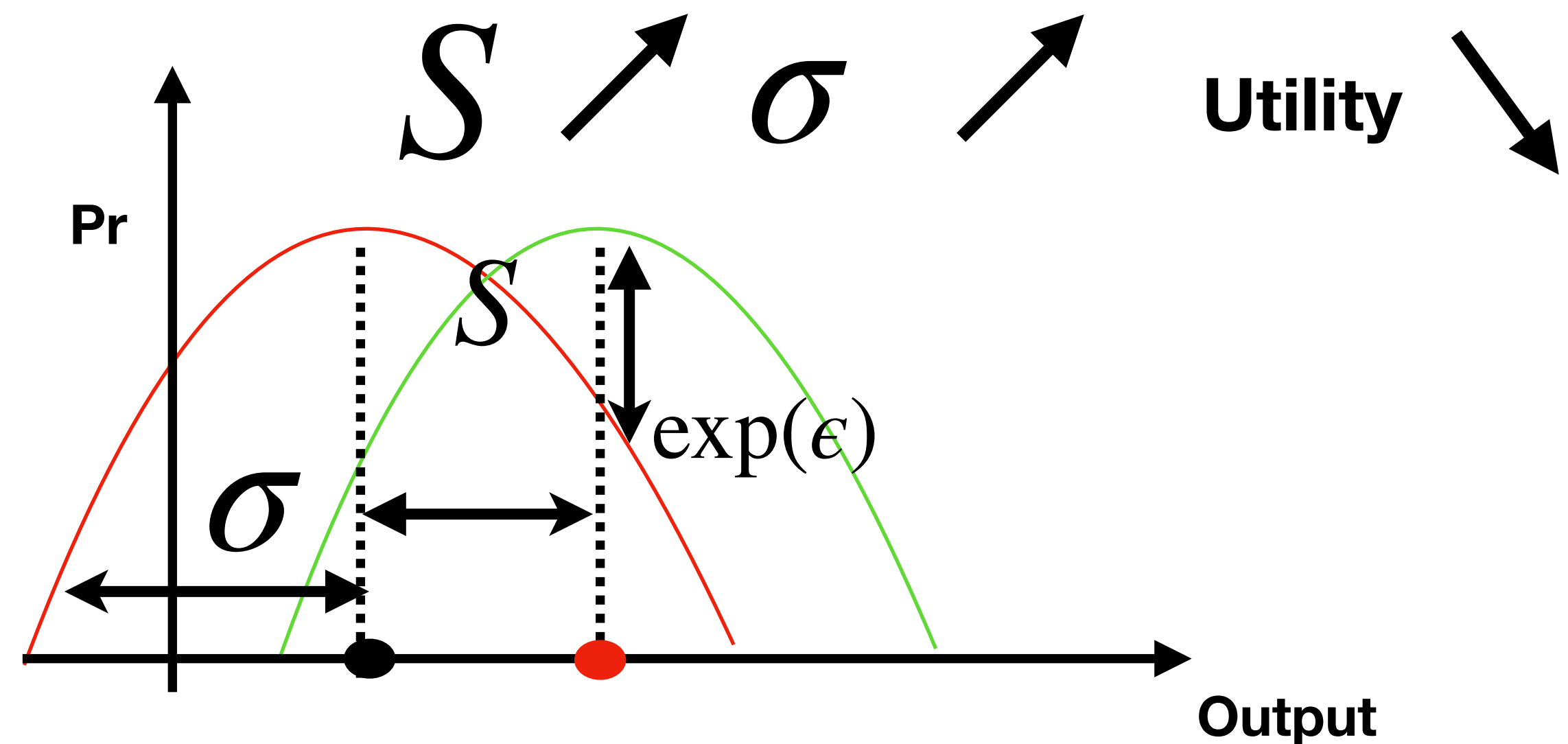
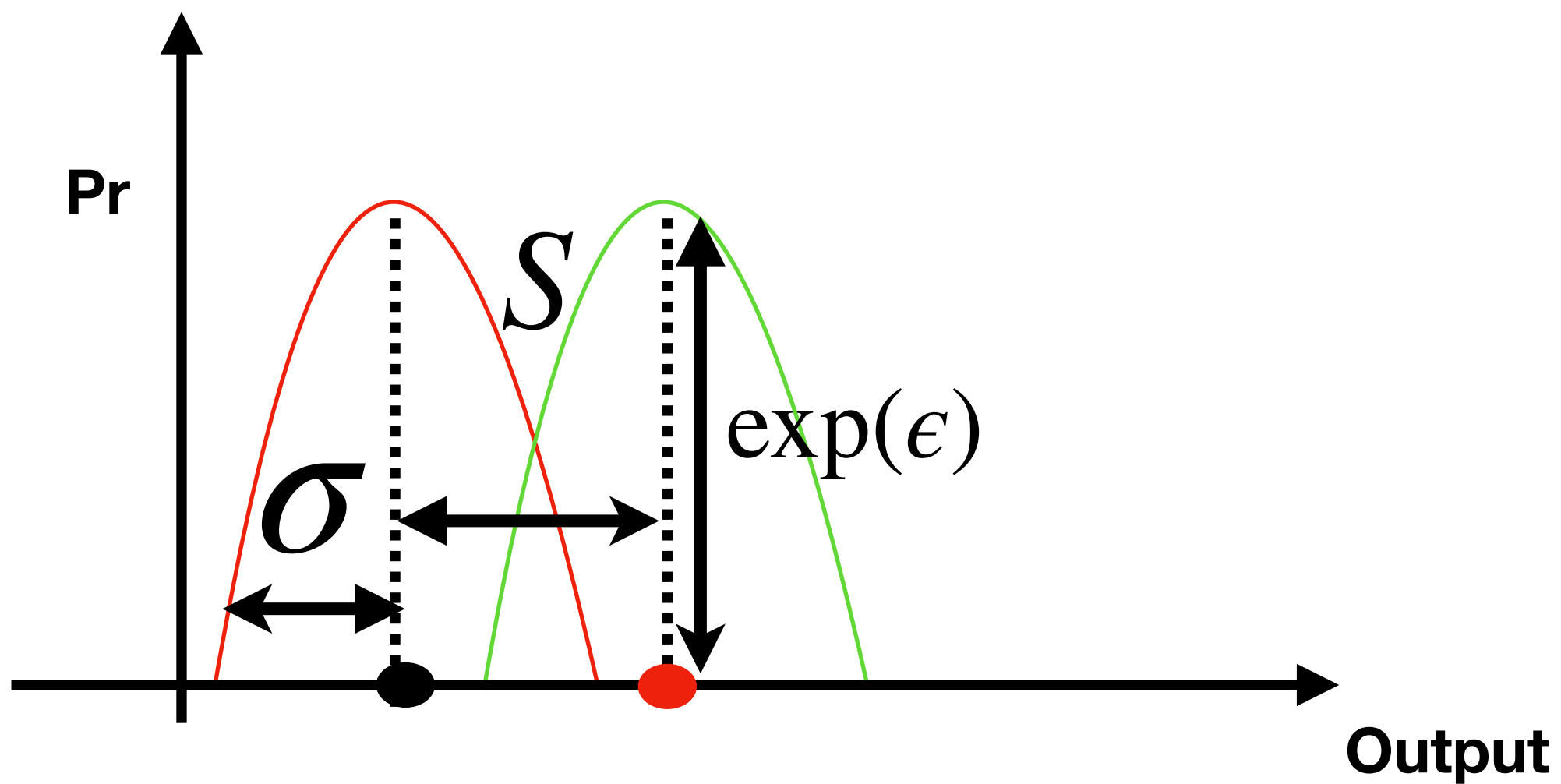


$$\sigma \approx \frac{S}{\epsilon}$$

Utility

Sensitivity

Privacy



# A better streaming LDP Mechanism

~~Repeatedly sample everyday.~~

There is a better way!

$$\sigma \approx \frac{S}{\epsilon}$$

Day 1



9000 visits

Day 2



10000 visits

+ 1000

# A better streaming LDP Mechanism

~~Repeatedly sample everyday.~~

There is a better way!

$$\sigma \approx \frac{S}{\epsilon}$$

Day 1



9000 visits

$$\pm C$$

Day 2



10000 visits

+1000

$$\underset{C}{1000} = \frac{1}{10} \times \underset{S}{10000}$$



# A better streaming LDP Mechanism

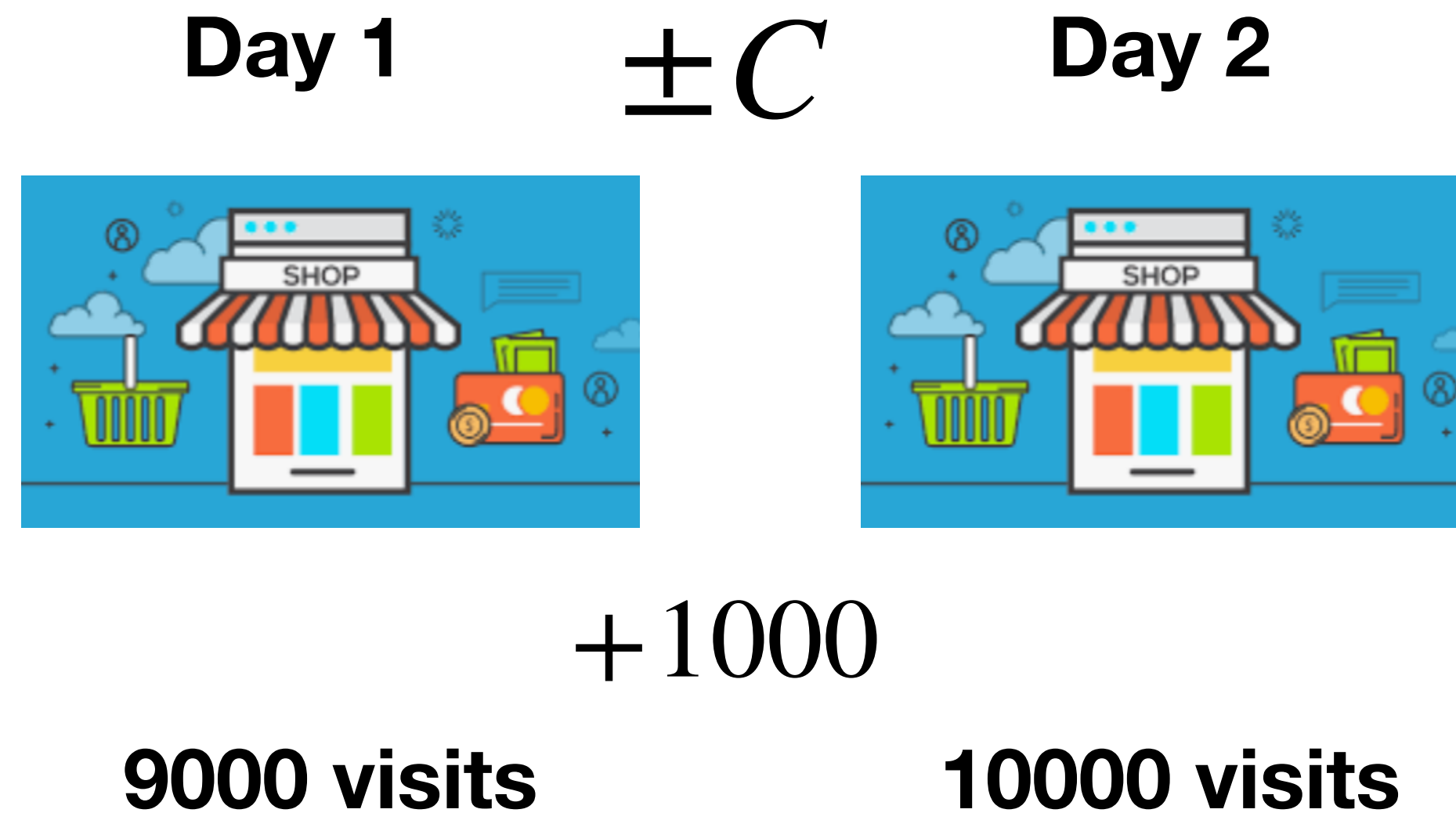
~~Repeatedly sample everyday.~~

There is a better way!

$$\sigma \approx \frac{S}{\epsilon}$$

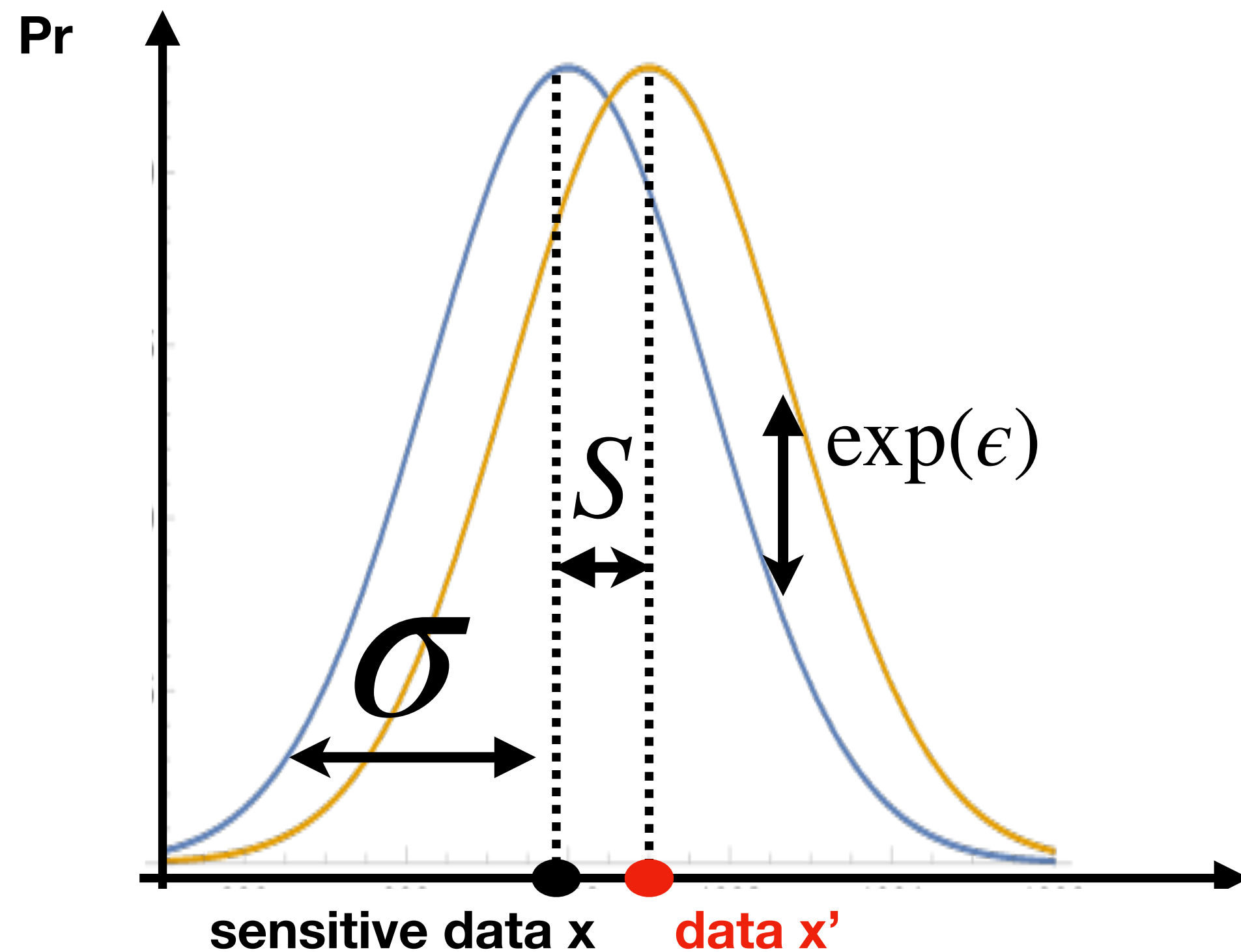
$$1000 = \frac{1}{10} \times 10000$$

$C$   $S$



$\sigma$   $\searrow$  Utility  $\nearrow$

# A better streaming LDP Mechanism



$$\text{Utility } \sigma \approx \frac{\text{Sensitivity } S}{\text{Privacy } \epsilon}$$

**CGM** helps carefully chooses the “proper” sensitivity  $S$ , in order to minimize  $\sigma$ , improving utility.

# A better streaming LDP Mechanism

Day 1

$\pm C$

Day 2



+1000

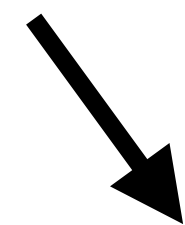
9000 visits

10000 visits

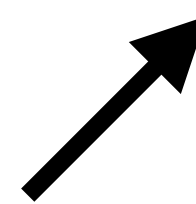
Other scenarios: daily App usage, daily phone usage, taxi locations, etc..

CGM:

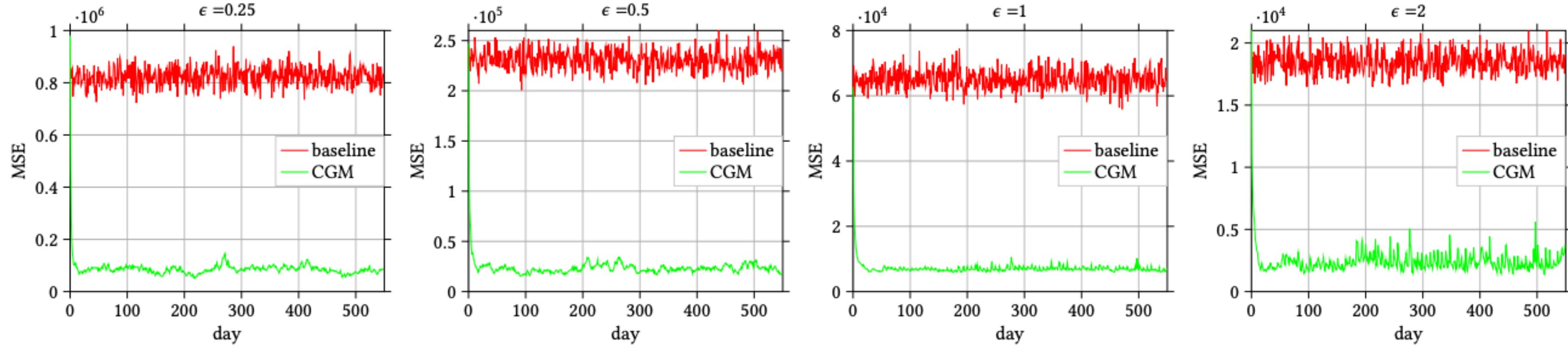
$\frac{C}{S}$



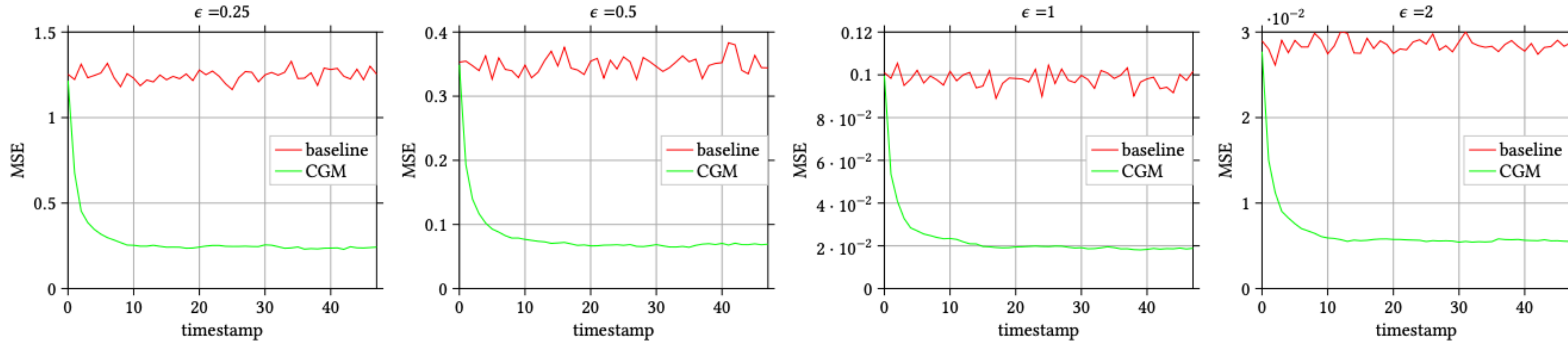
Utility improvement



# A better streaming LDP Mechanism



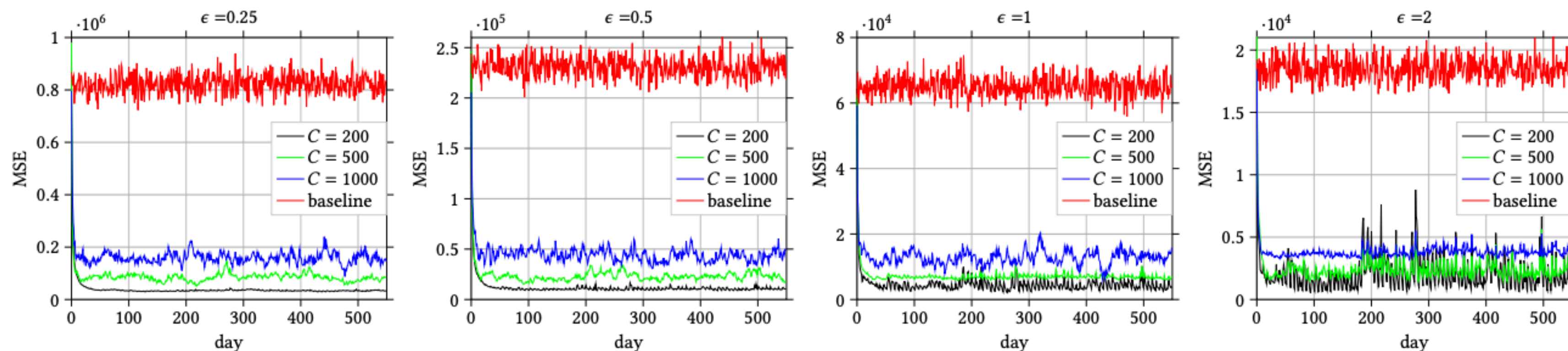
**Figure 1: Utility performances of CGM (Algorithm 3) and the baseline approach (Algorithm 1) on the Kaggle Web Traffic dataset, with varying daily privacy budget  $\epsilon \in \{0.25, 0.5, 1, 2\}$  and  $\delta = 10^{-5}$ . For CGM, the differential bound is fixed to  $C = 500$ .**



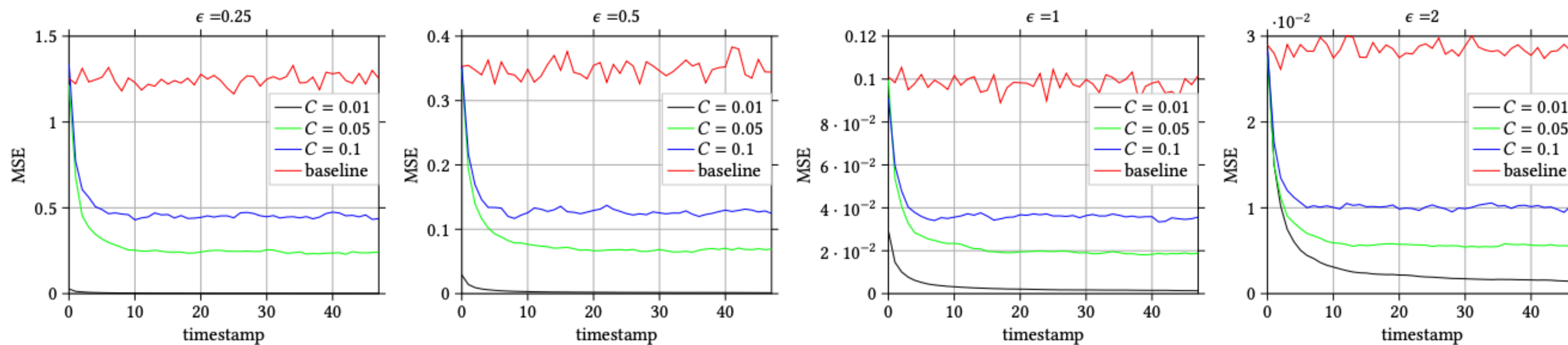
**Figure 2: Utility performances of CGM (Algorithm 3) and the baseline approach (Algorithm 1) on the Beijing Taxi dataset, with varying total privacy budget for all updates  $\epsilon \in \{0.25, 0.5, 1, 2\}$  and  $\delta = 10^{-5}$ . For CGM, the differential bound is fixed to  $C = 0.05$ . The whole space is normalized to  $[0, 1] \times [0, 1]$ , and the query region is  $[0.45, 0.55] \times [0.45, 0.55]$ .**



# Effect of $\frac{C}{S}$



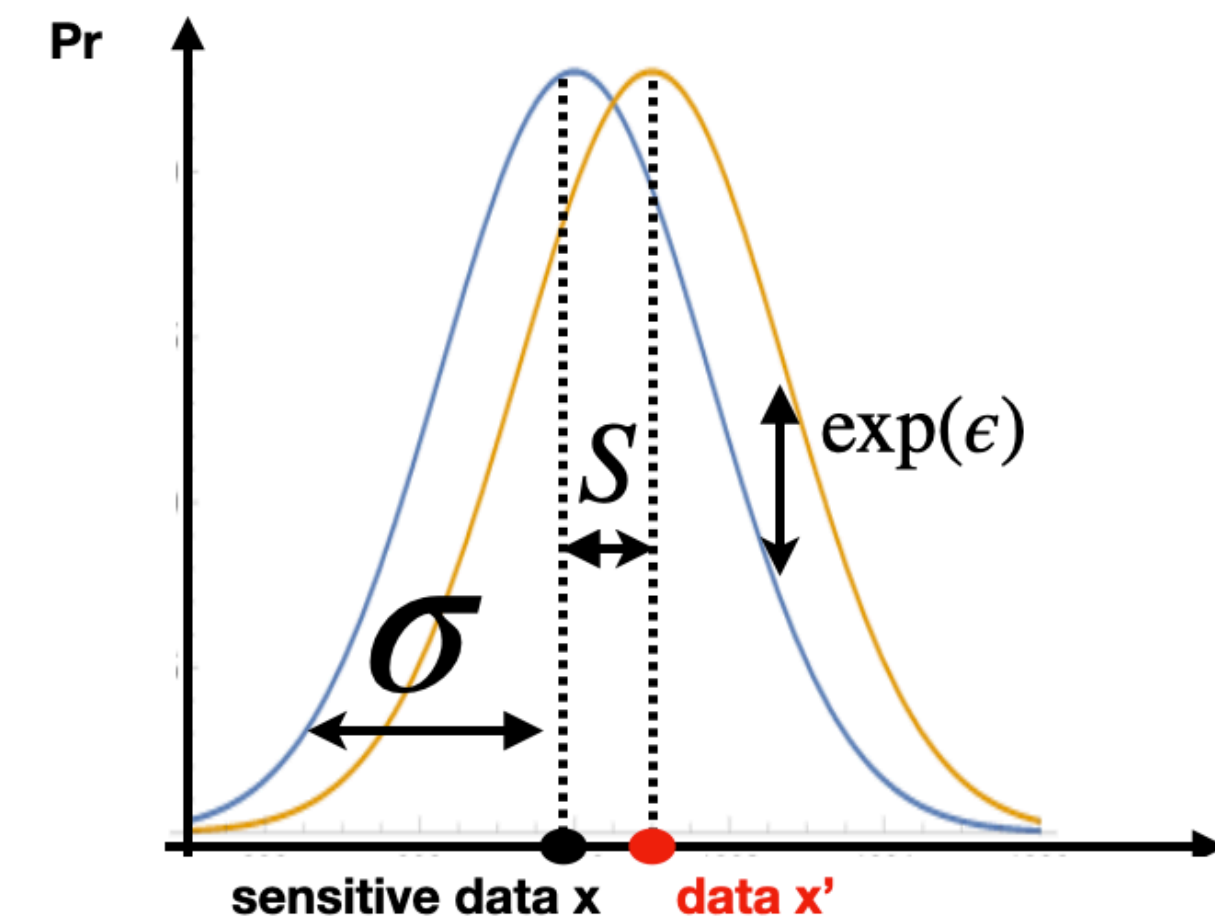
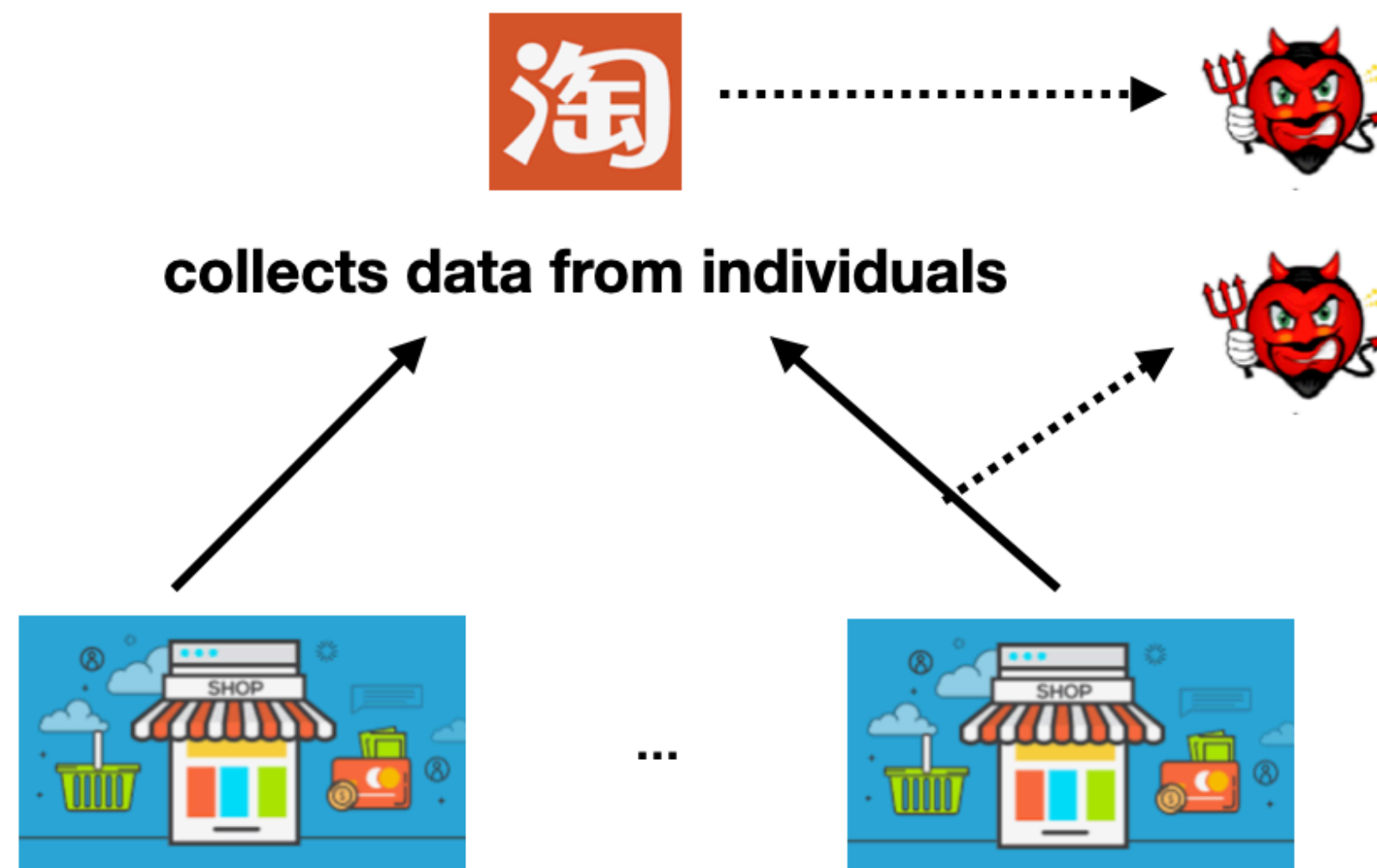
**Figure 3: Impact of varying differential bound  $C \in \{200, 500, 1000\}$  on the utility performance of CGM on the Kaggle Web Traffic dataset, where  $\epsilon \in \{0.25, 0.5, 1, 2\}$  and  $\delta = 10^{-5}$ .**



**Figure 4: Impact of varying differential bound  $C \in \{0.01, 0.05, 0.1\}$  on the utility performance of CGM on the Beijing Taxi dataset, where  $\epsilon \in \{0.25, 0.5, 1, 2\}$  and  $\delta = 10^{-5}$ . The query region is  $[0.45, 0.55] \times [0.45, 0.55]$ .**

# Summary

In this work, we study the problem of streaming data collection under local differential privacy. In this setting, an individual possesses a stream of data items, and the goal is to design a randomized mechanism for releasing the data stream without compromising the individual's privacy.

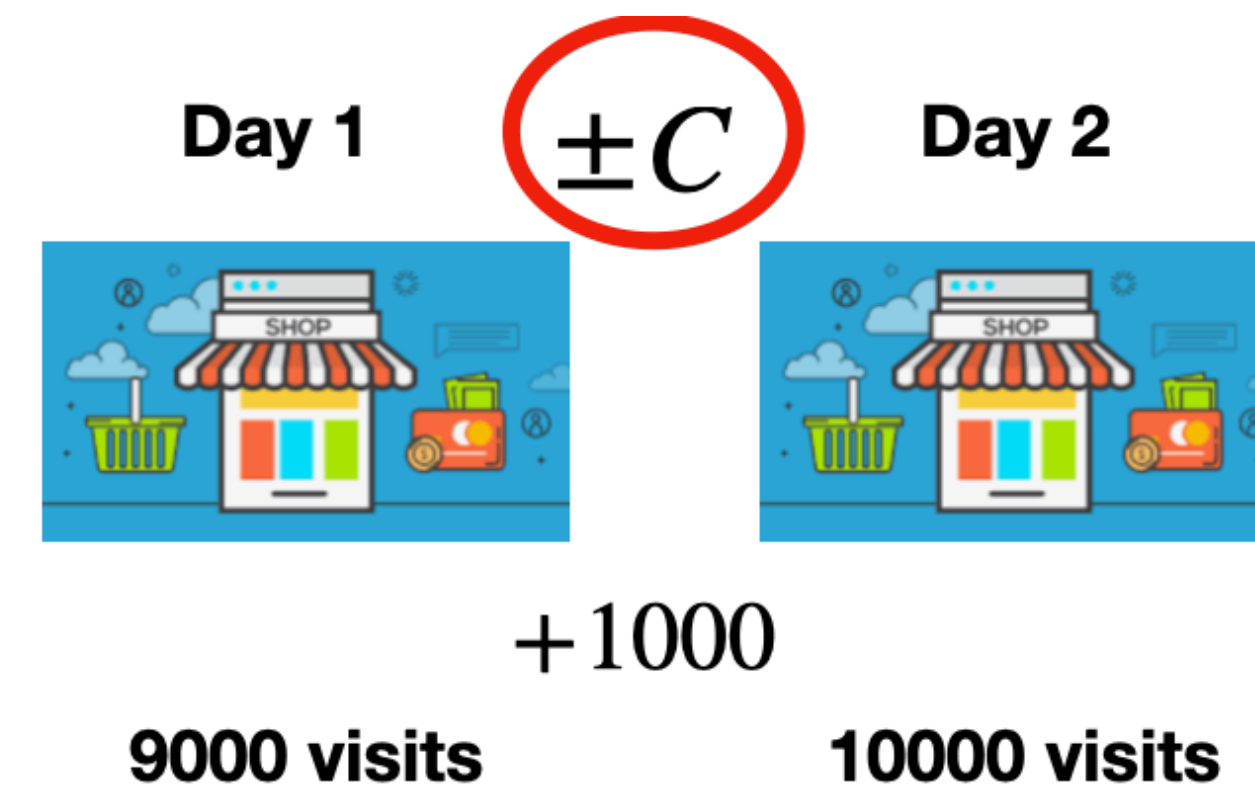




# Summary

The naive, and yet common approach, requires each user to perturb the data item independently at each timestamp, and upload the perturbed data to the untrusted aggregator. This approach leads to an excessively large amount of noise. Addressing this issue, we exploit data autocorrelations common in many real data streams, and propose a novel correlated Gaussian mechanism (CGM).

$$\text{Utility } \sigma \approx \frac{S \text{ Sensitivity}}{\epsilon \text{ Privacy}}$$



# Summary

Through both theoretical analysis and extensive experimental evaluations using real data from multiple application domains, we demonstrate that CGM consistently and significantly outperforms the baseline solution in terms of result utility.