

# **Computer Networks**

## **(CISC3001)**

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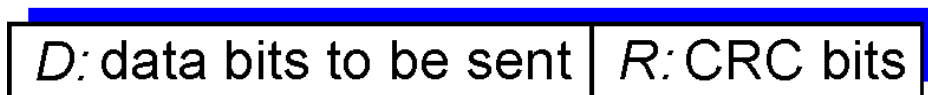
**Department of Computer and Information Science**

**SKL of Internet of Things for Smart City**

# Cyclic redundancy check

- more powerful error-detection coding
- view data bits, **D**, as a binary number
- choose  $r+1$  bit pattern (generator), **G**
- goal: choose  $r$  CRC bits, **R**, such that
  - $\langle D, R \rangle$  exactly divisible by  $G$  (modulo 2)
  - receiver knows  $G$ , divides  $\langle D, R \rangle$  by  $G$ . If non-zero remainder: error detected!
  - can detect all burst errors less than  $r+1$  bits
- widely used in practice (Ethernet, 802.11 WiFi, ATM)

← d bits → ← r bits →



*bit  
pattern*

$$D * 2^r \text{ XOR } R$$

*mathematical  
formula*

# Cyclic redundancy check

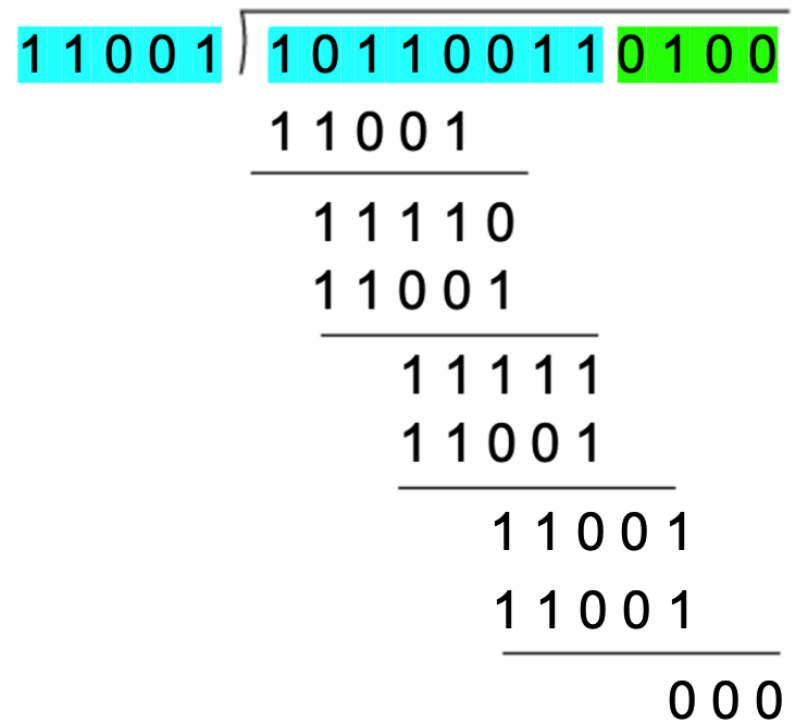
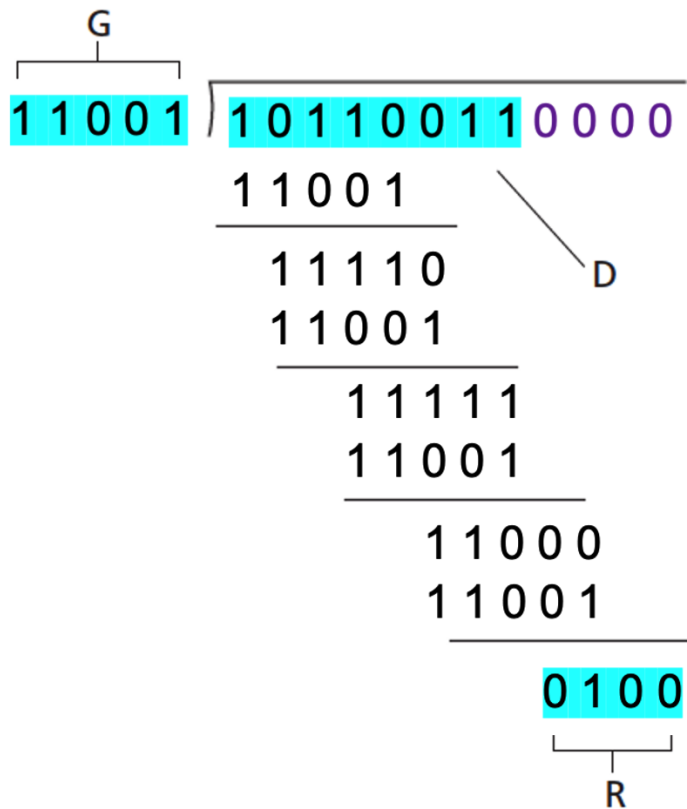
- Performance:
  - CRC can detect all **single bit** errors ( $G > 1$ )
  - CRC can detect all **double-bit** errors provided the **G** contains at least three logic 1's
  - CRC can detect all burst error of **length less than  $r+1$  bits**
  - CRC can detect most of the larger burst errors with a high probability
  - For example, CRC-12 detects 99.97% of error
- Conditions for **G**:
  - The highest and lowest bits **must be 1**
  - When an error occurs in any bit, the remainder should be **non-zero**
  - When errors occur in different bits, the remainders should be **different**
  - Balance the number of “1” and the number of “0”

**Q1:** For **D:10110011**, **G:11001**, what are the CRC bits?

**Q2:** Is **G:1001** a “good” generator for **D:101110**?

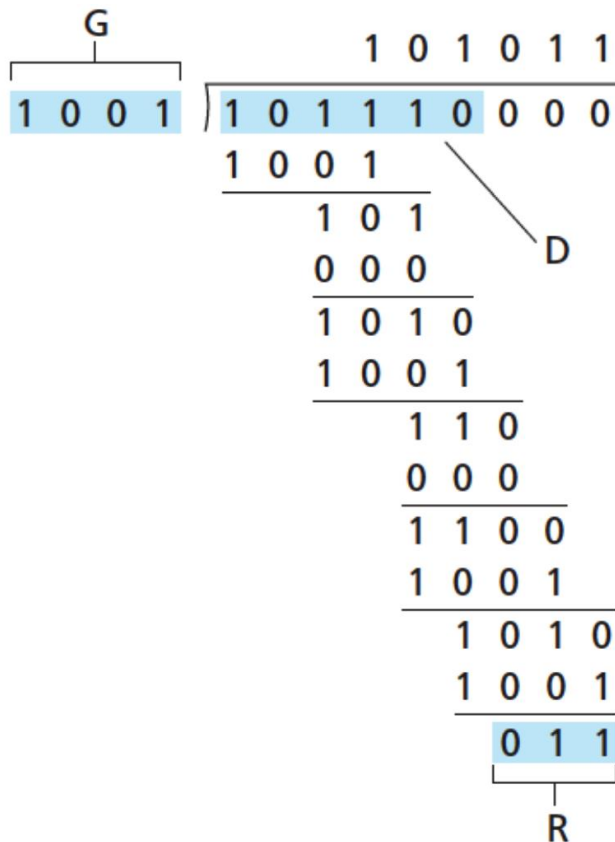
# CRC example

**Q/:** For **D:10110011**, **G:11001**, what are the CRC bits?

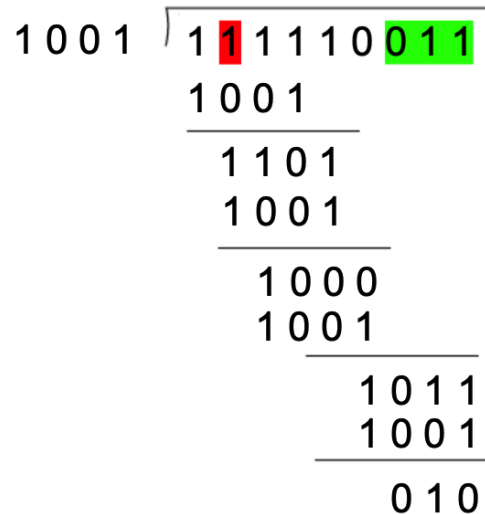


# CRC example

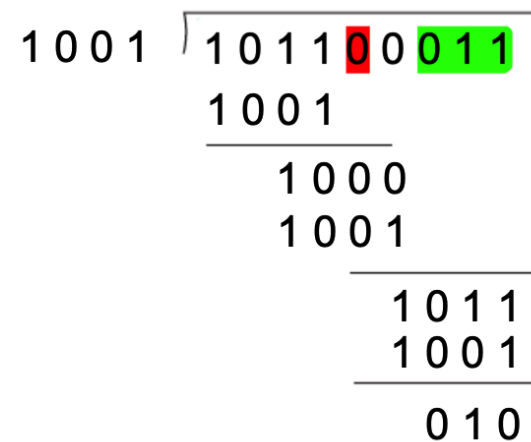
Q2: Is G:1001 a “good” generator for D:101110?



For the 2nd bit error, the remainder:



For the 5th bit error, the remainder:



Bit pattern: 1 0 1 1 1 0 0 1 1

# CRC example

**Q2:** Is **G:1001** a “good” generator for **D:101110**?

When errors occur in different bits, the remainders should be different!

For the second bit and the fifth bit errors, the remainder:

$$\begin{array}{r}
 1001 \overline{) 11110011} \\
 \underline{1001} \phantom{0000} \\
 1100 \phantom{0000} \\
 \underline{1001} \phantom{0000} \\
 1010 \phantom{0000} \\
 \underline{1001} \phantom{0000} \\
 1101 \phantom{0000} \\
 \underline{1001} \phantom{0000} \\
 1001 \phantom{0000} \\
 \underline{1001} \phantom{0000} \\
 0 \phantom{0000}
 \end{array}$$

No error checked ! **0**

The case of the first bit and fourth bit errors is similar

**G:1001** might not be a good generator for **D:101110**

The commonly used G

	Polynomial	G
CRC-4	$x^4 + x^3 + x^2 + 1$	11101
	$x^4 + x^2 + x + 1$	10111
	$x^4 + x + 1$	10011

In practical 802.11 protocol, normally  
CRC-32 with a generator of 33 bits is used