

Computer Networks

(CISC3001)

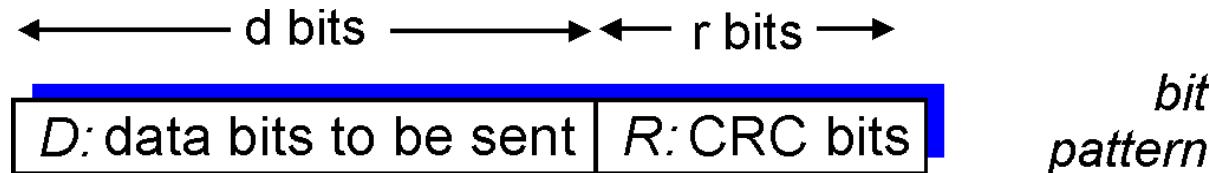
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Cyclic redundancy check

- more powerful error-detection coding
- view data bits, **D**, as a binary number
- choose $r+1$ bit pattern (generator), **G**
- goal: choose r CRC bits, **R**, such that
 - $\langle D, R \rangle$ exactly divisible by G (modulo 2)
 - receiver knows G, divides $\langle D, R \rangle$ by G. If non-zero remainder: error detected!
 - can detect all burst errors less than $r+1$ bits
- widely used in practice (Ethernet, 802.11 WiFi, ATM)



$$D * 2^r \text{ XOR } R$$

mathematical formula

Cyclic redundancy check

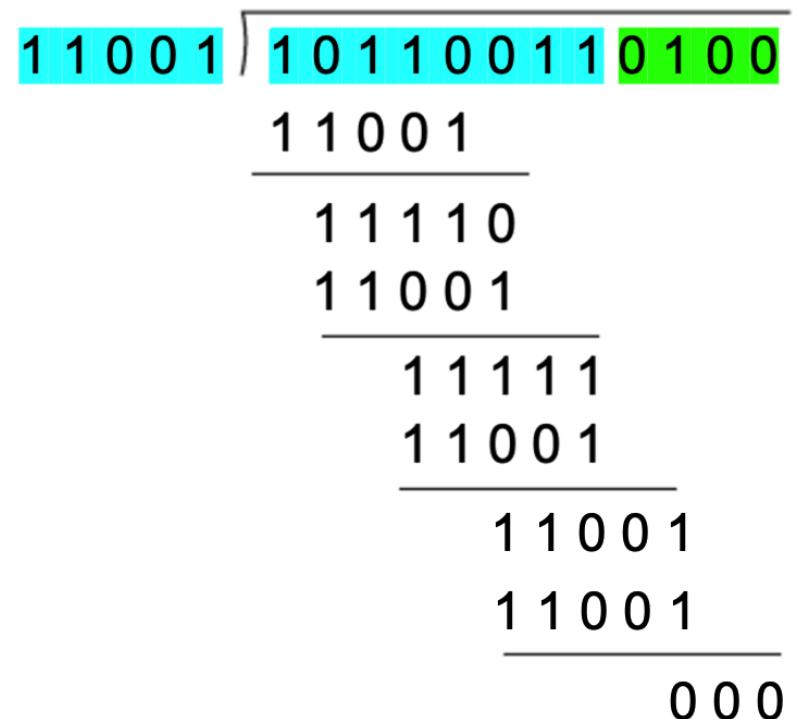
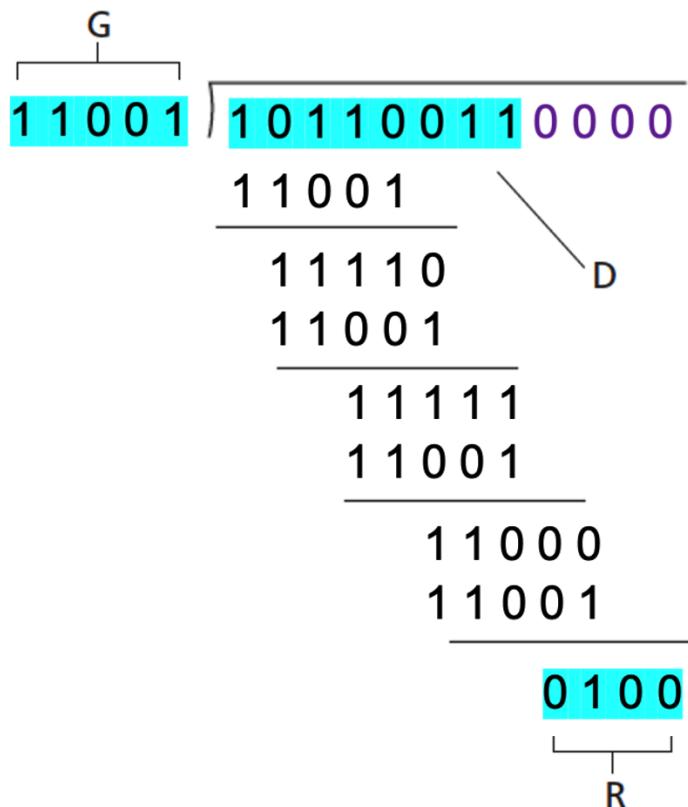
- Performance:
 - CRC can detect all **single bit** errors ($G > I$)
 - CRC can detect all **double-bit** errors provided the **G** contains at least three logic **I**'s
 - CRC can detect all burst error of **length less than $r+I$ bits**
 - CRC can detect most of the larger burst errors with a high probability
 - For example, CRC-12 detects **99.97%** of error
- Conditions for **G**:
 - The highest and lowest bits **must be I**
 - When an error occurs in any bit, the remainder should be **non-zero**
 - When errors occur in different bits, the remainders should be **different**
 - Balance the number of “I” and the number of “0”

Q1: For D:**10110011**, G:**11001**, what are the CRC bits?

Q2: Is G:**1001** a “good” generator for D:**101110**?

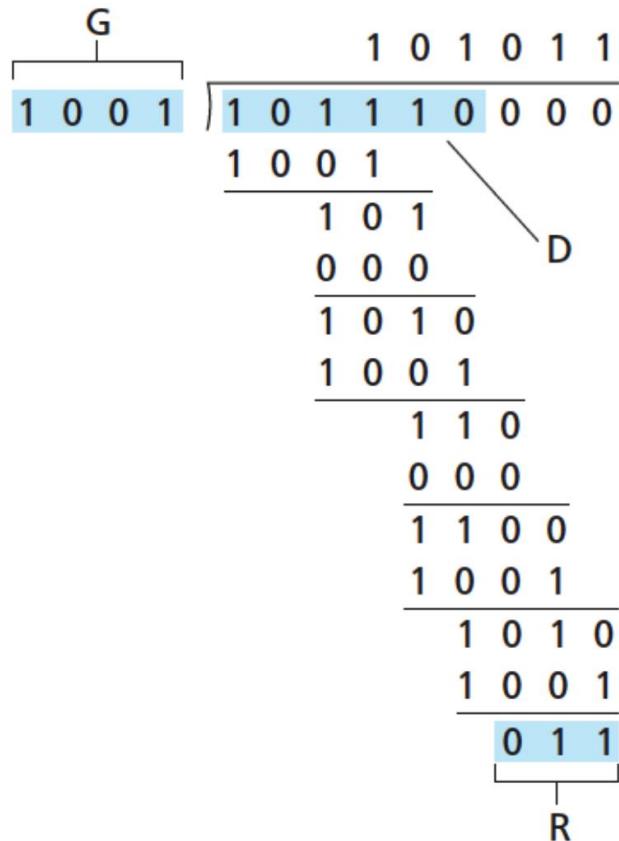
CRC example

Q1: For D:10110011, G:11001, what are the CRC bits?



CRC example

Q2: Is G:1001 a “good” generator for D:101110?



For the 2nd bit error,
the remainder:

$$\begin{array}{r} 1001 \longdiv{1\textcolor{red}{1}11110\textcolor{green}{0}11} \\ \quad 1001 \\ \hline \quad 1101 \\ \quad 1001 \\ \hline \quad 1000 \\ \quad 1001 \\ \hline \quad 1011 \\ \quad 1001 \\ \hline \quad 010 \end{array}$$

For the 5th bit error,
the remainder:

$$\begin{array}{r} 1001 \longdiv{1011\textcolor{red}{0}0011} \\ \quad 1001 \\ \hline \quad 1000 \\ \quad 1001 \\ \hline \quad 1011 \\ \quad 1001 \\ \hline \quad 010 \end{array}$$

Bit pattern: 101110011

CRC example

Q2: Is G:1001 a “good” generator for D:101110?

When errors occur in different bits, the remainders should be different!

For the second bit and the fifth bit errors,
the remainder:

$$\begin{array}{r} 1001 \overline{)111100011} \\ 1001 \\ \hline 1100 \\ 1001 \\ \hline 1010 \\ 1001 \\ \hline 1101 \\ 1001 \\ \hline 1001 \\ 1001 \\ \hline \end{array}$$

No error checked ! 0

The commonly used G

	Polynomial	G
CRC-4	$x^4 + x^3 + x^2 + 1$	11101
	$x^4 + x^2 + x + 1$	10111
	$x^4 + x + 1$	10011

In practical 802.11 protocol, normally
CRC-32 with a generator of 33 bits is used

The case of the first bit and fourth bit errors is similar

G:1001 might not be a good generator for D:101110