## CEC 15<sup>th</sup> - Past Year Paper Solution 2014-2015 Sem2 CZ1007 - Data Structures

Solver: Eric Leonardo Lim

Email Address: elim023@e.ntu.edu.sg

- 1. (a)
  - (i) The output of the code is the sum of all elements in row 0..i which are positive integers and to the left of the first 0, if any, for all rows i = 0..3.

More formally, let  $x_i$  be the first column j in row i such that  $a_{i,j} = 0$ , and  $x_i$  is equal to the number of columns (which is 4) if there doesn't exists such j. Let  $d_i$  be the sum of all  $a_{i,j}$  in row i and column j such that  $a_{i,j} > 0$  with  $0 \le j < x_i$ . Then, the output for each row i = 0..3 is  $\sum_{k=0}^i d_k$ .

Output:

3

10

12

20

(ii) Notice that f1 is a function that swaps the values of its parameters. This is because the code works like the following:

$$c = a + b$$
  
 $b' = c - b = (a + b) - b = a$ 

$$a' = c - b' = (a + b) - a = b$$

a' and b' denote the new values of a and b and they are equal to b and a respectively.

Output:

2, 1

3, 1

2, 1

(b)

```
(i) Line 5. double sum = 0;  // floating point for computing average
Line 10. while(s[i] != 0) {  // iterate through the array until a 0.
Line 18. sum /= c;  // "/=" for divide operation
Line 19. printf("%f\n",sum); // want to output average stored in sum.
```

For line 10, if operator > is used instead, it will return true if the address of the left string is larger than the address of the right string. As the strings are stored in a contiguous



memory block, it returns true if the index of the left string is larger than the index of the right string.

```
// pointer p points to a[0]
  (c)
                                   // pointer q points to b[0]
 (i) p = a;
                                   // after loop, p points to end of a
     q = b;
     while(*p) p++;
                                   // concatenate
     while(*q) *(p++) = *(q++);
                                   // add '\0' string terminator
     *p = 0;
     printf("%s\n",a);
 (i) f2(3,5)
               = 5*f2(2,3)
               = 5*5* f2(1,1)
               = 5*5*5*f2(0,-1)
               = 5*5*5*1
    Output:
    125
                           + f3(3)
 (ii) f3(5)
               = f3(4)
               = f3(3) + f3(2) + f3(2) + f3(1)
              = f3(3) + 1 + 1 + 1
              = f3(2) + f3(1) + 1 + 1 + 1
              =1+1+1+1+1=5
    Output:
   f3(): 5
(b)
              void f4(char *s, int *a, int *b)
(i) Line 2.
   Line 9. void f4(char *s, int *a, int *b) {
   Line 11. if(*s>='A' && *s<='Z') (*a)++;
   Line 12. if(*s>='a' && *s<='z') (*b)++;
```

(ii) We need to come up with how the code should work. Otherwise, debugging the code will be more difficult. The code should work in the following way.

First, find the maximum character. Then, for all the characters in its left, copy them to their immediate right. Finally, assign the maximum character to be the first character of the string.

e.g. str = abdzc. Let max\_character = z. Copy all characters in max\_character's left to its immediate right, i.e. str = aabdc. Finally, assign max\_character to be the first character of str, i.e. str = zabdc.

```
Line 4. char str[100]; // size should be initialized for input line 13. q = p+i;
```

Line 17. while(q>p){ Line 18. \*q = \*(q-1);

## CEC 15<sup>th</sup> - Past Year Paper Solution 2014-2015 Sem2 CZ1007 - Data Structures

```
Line 17. while(q>p){
       Line 18. *q = *(q-1);
   (c)
   if(ar[0] == target) return 0;
   if(rLookup(ar+1,n-1,target) >= 0)
          return rLookup(ar+1,n-1,target) + 1;
   else return -1;
3. (a)
   (i) Line 8. n * sizeof(int)
       Line 13. free(intArr)
   (ii) It may access memory space which is already allocated.
   (b)
    (i) 2,1
    (ii) A 6-operation sequence is possible if and only if:
           Both push and pop operations are done exactly three times.
           Pop operations are done only when the stack is not empty, i.e. there exists at least
           an integer in the stack. We can have O as the next character in the sequence if the
           current number of O's is less than 3 and the current number of U's is more than
           current number of O's in the current sequence.
       There are exactly 5 possible distinct such sequences:
       UOUOUO -> 1,2,3
       UOUU00 -> 1,3,2
       UUOOUO -> 2,1,3
       UUOU00 -> 2,3,1
       UUU000 -> 3,2,1
    (iii) Queue's working principle is FIFO (First-In First-Out), whereas stack's working principle is
       LIFO (Last-In Last-Out).
   (c)
   while(p->next->next != NULL){ // there are at least 4 nodes
          p = p->next;
          size++;
    }
                                         // because haven't count the last 2 nodes
    size+=2;
   tmp = head->next;
   head->next = p->next;
   head->next->next = tmp->next;
   p->next = tmp;
   tmp->next = NULL;
```



## CEC 15<sup>th</sup> - Past Year Paper Solution 2014-2015 Sem2 CZ1007 - Data Structures

- 4. (a)
  - (i) Figure Q4(1) is not a BST because 14 is located in the left sub-tree of 13 and 14>13.
  - (ii) Let L<sub>x</sub> be the set of all nodes in the left subtree of node x.
     Let R<sub>x</sub> be the set of all nodes in the right subtree of node x.
     Figure Q4(2) is a BST because for all nodes x in the tree, x has at most 2 children and ∀l ∈ L<sub>x</sub>. ∀r ∈ R<sub>x</sub>. l < x < r,</li>
  - (b)
  - (i) A ...
  - (ii) B-A-D-C
  - (c)

```
int printParent(BTNode *node, BTNode *x){
    if(node == NULL) return 0;
    int flag = 0;
    if(node->left != NULL && node->left == x){
        printf("%d\n",node->item);
        flag = 1;
    }
    if(node->right != NULL && node->right == x){
        printf("%d\n",node->item);
        flag = 1;
    }
    return flag | printParent(node->left,x) | printParent(node->right,x);
}
```

For reporting of errors and errata, please visit pypdiscuss.appspot.com
Thank you and all the best for your exams!