Consider the page table shown in Figure 3.1 for a system with 12-bit virtual and physical addresses and with 256-byte pages. The list of free page frames is D, E, F (that is, D is at the head of the list, E is second, and F is last). Convert the following virtual

Page	Page Frame
0	–
1	2
2	C
3	A
4	–
5	4
6	3
7	–
8	B
9	0

Answers:

- Page Size = 256-byte = 2^8 byte => 8-bit offset

- 12-bit virtual address

- 12-bit physical address

• 9EF has offset: EF, for page 9 we find frame 0 so result is 0EF

• 111 has offset: 11, for page 1 we find frame 2 so result is 211

• 700 has offset: 00, because we can't find page 7 frame in virtual memory,so we get the first frame from free page frames list.We got D,so result is D00

• The same for 0FF,it has offset FF,we can't find page frame with page 7,so we get the second frame from free page frames list.Then we got E,so resuilt is