

Support Vector Machine

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Primal and Dual Problem

Primal Problem

$$\begin{aligned} \min_x & f(x) \\ \text{s.t. } & g(x) \leq 0, h(x) = 0 \end{aligned}$$



$$\min_x \max_{\alpha \geq 0, \beta} L(x, \alpha, \beta)$$

Lagrange Dual Problem

$$\begin{aligned} \max_{\alpha > 0, \beta} & d(\alpha, \beta) \\ \text{s.t. } & \alpha > 0 \end{aligned}$$



$$\max_{\alpha \geq 0, \beta} \min_x L(x, \alpha, \beta)$$

- Weak duality theorem
 - $d(\alpha, \beta) \leq f(x^*)$ for $\forall \alpha, \forall \beta$
 - $d^* = \max_{\alpha \geq 0, \beta} \min_x L(x, \alpha, \beta) \leq \min_x \max_{\alpha \geq 0, \beta} L(x, \alpha, \beta) = p^*$
 - Maximizing the dual function provides the lower bound of $f(x^*)$
 - Duality gap = $f(x^*) - d(\alpha^*, \beta^*)$
- Strong duality
 - $d^* = \max_{\alpha \geq 0, \beta} \min_x L(x, \alpha, \beta) = \min_x \max_{\alpha \geq 0, \beta} L(x, \alpha, \beta) = p^*$
 - When Karush-Kuhn-Tucker (KKT) Conditions are satisfied

KKT Condition and Strong Duality

- Strong duality

- $$d^* = \max_{\alpha \geq 0, \beta} \min_x L(x, \alpha, \beta) = \min_x \max_{\alpha \geq 0, \beta} L(x, \alpha, \beta) = p^*$$

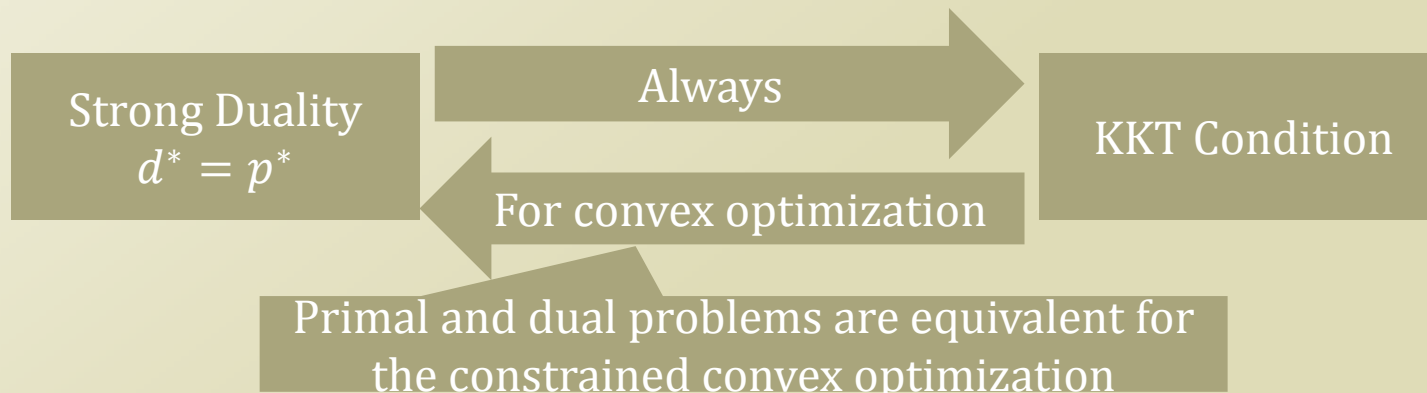
- Holds when KKT conditions are met

- $\nabla L(x^*, \alpha^*, \beta^*) = 0$
- $\alpha^* \geq 0$
- $g(x^*) \leq 0$
- $h(x^*) = 0$
- $\alpha^* g(x^*) = 0$

Active Constraint
 $\alpha^* = 0 \Rightarrow g(x^*) = 0$
 Inactive Constraint
 $g(x^*) < 0 \Rightarrow \alpha^* = 0$
 \rightarrow Complementary Slackness

Primal Problem

$$\begin{aligned} \min_x & f(x) \\ \text{s.t. } & g(x) \leq 0, \\ & h(x) = 0 \end{aligned}$$



Dual Problem of SVM

Primal Problem

$$\begin{aligned} \min_x & f(x) \\ \text{s.t. } & g(x) \leq 0, \\ & h(x) = 0 \end{aligned}$$

Lagrange Prime Function

$$L(x, \alpha, \beta) = f(x) + \alpha g(x) + \beta h(x)$$

Primal Problem of Linearly Separable SVM

$$\begin{aligned} \min_{w,b} & ||w|| \\ \text{s.t. } & (wx_j + b)y_j \geq 1, \forall j \end{aligned}$$

$$\begin{aligned} \min_{w,b} \max_{\alpha \geq 0, \beta} & \frac{1}{2} w \cdot w - \sum_j \alpha_j [(wx_j + b)y_j - 1] \\ \text{s.t. } & \alpha_j \geq 0, \text{ for } \forall j \end{aligned}$$

- Linearly separable case
- Lagrange Prime Function

- $L(w, b, \alpha)$

$$= \frac{1}{2} w \cdot w - \sum_j \alpha_j [(wx_j + b)y_j - 1]$$

- Lagrange Multiplier

- $\alpha_j \geq 0, \text{ for } \forall j$

Dual Problem of Linearly Separable SVM

$$\begin{aligned} \max_{\alpha \geq 0} \min_{w,b} & \frac{1}{2} w \cdot w - \sum_j \alpha_j [(wx_j + b)y_j - 1] \\ \text{s.t. } & \alpha_j \geq 0, \text{ for } \forall j \end{aligned}$$

KKT Condition to Eliminate the Duality Gap

$$\frac{\partial L(w,b,\alpha)}{\partial w} = 0, \frac{\partial L(w,b,\alpha)}{\partial b} = 0$$

$$\alpha_i \geq 0, \forall i$$

$$\alpha_i ((wx_j + b)y_j - 1) = 0, \forall i$$

Dual Representation of SVM

- $L(w, b, \alpha) = \frac{1}{2}w \cdot w - \sum_j \alpha_j [(wx_j + b)y_j - 1]$
 - $= \frac{1}{2}ww - \sum_j \alpha_j y_j wx_j - b \sum_j \alpha_j y_j + \sum_j \alpha_j$
 - $= \frac{1}{2} \sum_i \sum_j \alpha_i \alpha_j y_i y_j x_i x_j - \sum_i \sum_j \alpha_i \alpha_j y_i y_j x_i x_j - b \times 0 + \sum_j \alpha_j$
 - $= \sum_j \alpha_j - \frac{1}{2} \sum_i \sum_j \alpha_i \alpha_j y_i y_j x_i x_j$
- Again, a quadratic programming
- Once, α_j is known
 - $w = \sum_{i=1}^N \alpha_i y_i x_i$
 - $\alpha_i ((wx_j + b)y_j - 1) = 0$
- Now, we can find out the w and b again.
 - Why is this better?
 - Most of α_j are.....
 - Location of x is....
- Let's find out from the implementation...

Dual Problem of Linearly Separable SVM

$$\max_{\alpha \geq 0} \min_{w, b} \frac{1}{2}w \cdot w - \sum_j \alpha_j [(wx_j + b)y_j - 1]$$

s. t. $\alpha_j \geq 0$, for $\forall j$

KKT Condition to Eliminate the Duality Gap

$$\frac{\partial L(w, b, \alpha)}{\partial w} = 0, \frac{\partial L(w, b, \alpha)}{\partial b} = 0$$

$$\alpha_i \geq 0, \forall i$$

$$\alpha_i ((wx_j + b)y_j - 1) = 0, \forall i$$

$$w = \sum_{i=1}^N \alpha_i y_i x_i \quad \sum_{i=1}^N \alpha_i y_i = 0$$