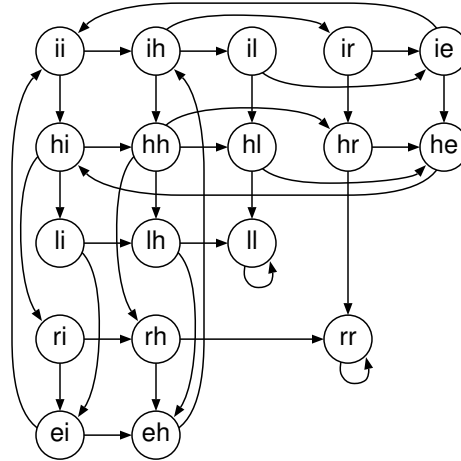


# CPRE/SE/COMS 412, COMS 512      HOMEWORK 2

Reminder: present your own work and properly cite any sources used. Solutions should be presented satisfying the *other student viewpoint*. If you need clarification, contact the instructor: [asminer@iastate.edu](mailto:asminer@iastate.edu).

## Question 1

Consider a Kripke structure model for two dining philosophers, as shown below.



Each state of the Kripke structure is denoted by the states of the two philosophers, where the philosopher states are:

**i** : Idle

**h** : Hungry, without any forks

**l** : Hungry, has left fork but not right

**r** : Hungry, has right fork but not left

**e** : Eating (has both forks)

Initially, both philosophers are *Idle*. Determine if the model satisfies the properties below (remember to show your work).

1. **(5 points)** Whenever Philosopher 1 becomes hungry, he will eventually eat:

$$\text{AG}(\text{Phil.1 hungry} \rightarrow \text{AF}(\text{Phil.1 eats}))$$

2. **(5 points)** The same property as above, using the fairness constraint  $\mathcal{C} = \{ii\}$ , i.e., only considering paths where both philosophers are idle, infinitely often:

$$\text{A}_\mathcal{C}\text{G}(\text{Phil.1 hungry} \rightarrow \text{A}_\mathcal{C}\text{F}(\text{Phil.1 eats}))$$

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## Question 2

Define a new CTL temporal operator  $\mathbf{EN}$ , that acts like ordinary “until”, except it does not require the second property to eventually hold. Formally, for any Kripke structure  $M$  and any state  $s$ ,  $M, s \models \mathbf{E} \phi_1 \mathbf{N} \phi_2$ , if and only if, there exists a path  $(p_0, p_1, p_2, \dots)$  with  $p_0 = s$ , such that either (a)  $p_i \models \phi_1, \forall i$ ; or (b)  $\exists j \geq 0 : p_j \models \phi_2, p_i \models \phi_1, \forall i < j$ .

1. **(5 points)** Prove or disprove:  $\mathbf{E} \mathbf{tt} \mathbf{N} \phi \equiv \mathbf{EF} \phi$
2. **(5 points)** Prove or disprove:  $\mathbf{E} \phi \mathbf{N} \mathbf{ff} \equiv \mathbf{EG} \phi$
3. **(5 points)** Prove:  $\mathbf{E} \phi_1 \mathbf{N} \phi_2 \equiv \phi_2 \vee (\phi_1 \wedge \mathbf{EX} \mathbf{E} \phi_1 \mathbf{N} \phi_2)$
4. **(10 points, Optional for 412)** Prove:  $\llbracket \mathbf{E} \phi_1 \mathbf{N} \phi_2 \rrbracket$  is the greatest fixed point of  $U_{\phi_1 \phi_2}$ , using the definition of  $U_{\phi_1 \phi_2}$  as given in Property 4.24.