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# COMS 512 Homework 0

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1. As  $p$  is a propositional variable,  $p$  can be either True or False. We denote True with  $tt$  and False with  $ff$ .

a)  $tt \rightarrow p$

If  $p = tt$  then  $tt \rightarrow p \equiv tt$ , If  $p = ff$  then  $tt \rightarrow p \equiv ff$ .

b)  $p \rightarrow tt$

If  $p = tt$  then  $p \rightarrow tt \equiv tt$ , If  $p = ff$  then  $p \rightarrow tt \equiv tt$ .

c)  $ff \rightarrow p$

If  $p = tt$  then  $ff \rightarrow p \equiv tt$ , If  $p = ff$  then  $ff \rightarrow p \equiv tt$ .

d)  $p \rightarrow ff$

If  $p = tt$  then  $p \rightarrow ff \equiv ff$ , If  $p = ff$  then  $p \rightarrow ff \equiv tt$ .

2. We can prove the given expression is a tautology by constructing the truth table of the premises.

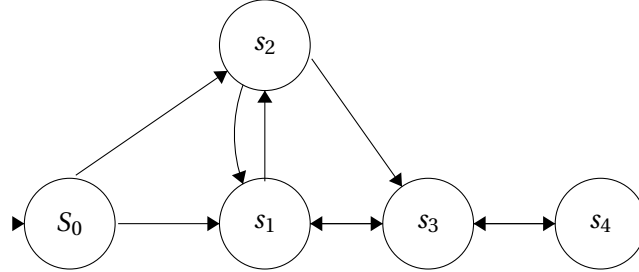
A	B	C	$A \rightarrow B$	$B \rightarrow C$	$A \rightarrow C$	$(A \rightarrow B) \wedge (B \rightarrow C)$	$[(A \rightarrow B) \wedge (B \rightarrow C)] \rightarrow (A \rightarrow C)$
tt	tt	tt	tt	tt	tt	tt	tt
tt	tt	ff	tt	ff	ff	ff	tt
tt	ff	tt	ff	tt	tt	ff	tt
tt	ff	ff	ff	tt	ff	ff	tt
ff	tt	tt	tt	tt	tt	tt	tt
ff	tt	ff	tt	ff	tt	ff	tt
ff	ff	tt	tt	tt	tt	tt	tt
ff	ff	ff	tt	tt	tt	tt	tt

First we create the truth table for  $A \rightarrow B$ ,  $B \rightarrow C$  and  $A \rightarrow C$ . Then we find the truth table for  $(A \rightarrow B) \wedge (B \rightarrow C)$ . Finally we find the truth table for the given expression. As all the

term in the given expression turn out to be True, we can conclude that the expression is a tautology.

3. Let us consider the following Kripke Structure, where  $S = \{s_0, s_1, s_2, s_3, s_4\}$ ,  $S = \{s_0\}$  is the initial state, the transition relation,

$$R = \{(s_0, s_1), (s_0, s_2), (s_1, s_2), (s_1, s_3), (s_2, s_1), (s_2, s_3), (s_3, s_1), (s_3, s_4), (s_4, s_3)\}.$$



Lets take  $X = \{s_0, s_1\}$ . The PreImage function of X will be  $PreImage(X, R) = \{s_0, s_2, s_3\}$ . Next we need to find the PostImage function of this PreImage set. The PostImage function will be  $PostImage(\{s_0, s_2, s_3\}, R) = \{s_1, s_2, s_3, s_4\}$ .

We can see that  $\{s_0, s_1\} \not\subseteq \{s_1, s_2, s_3, s_4\}$ . i.e  $X \not\subseteq PostImage(\{s_0, s_2, s_3\}, R)$ . Therefore from this counter-example, we can conclude that  $X \not\subseteq PostImage(Preimage(X, R), R)$ .

4. By the definition of Kripke structure we know that every state of the Kripke structure must have at least one outgoing edge.

Now if we consider a set of states  $X \subseteq S$  and try to find their PostImage set Y, then every state from X will have at least one state in Y that it can reach in one step.

i.e. if  $s \in X$  then  $\exists s' \in Y$  s.t  $(s, s') \in R$ .

Hence every state in Y has atleast one predecessor state in X. therefore if we construct the Preimage set of Y it will cover all those states that belong to set X. i.e.  $X \subseteq PreImage(Y, R)$ .

Therefore  $X \subseteq PreImage(PostImage(X, R), R)$ .