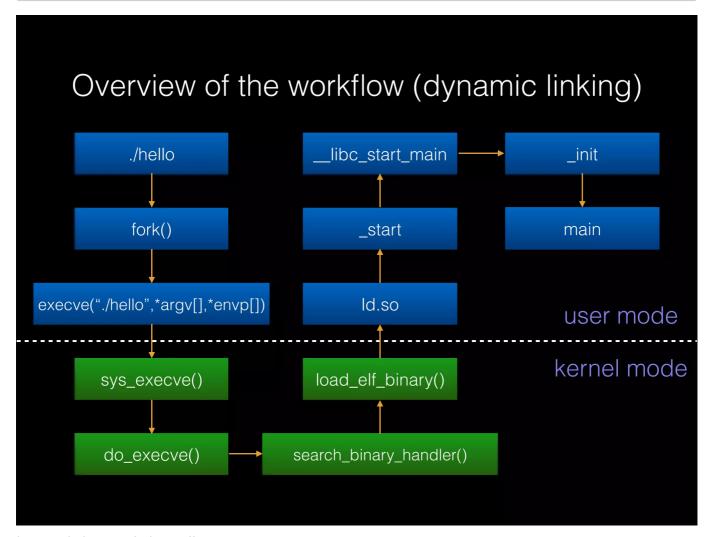
Execuation Explained

Overview

https://www.youtube.com/watch?v=Ss2e6JauS0Y

https://www.slideshare.net/AngelBoy1/execution-50215114 (recommended)

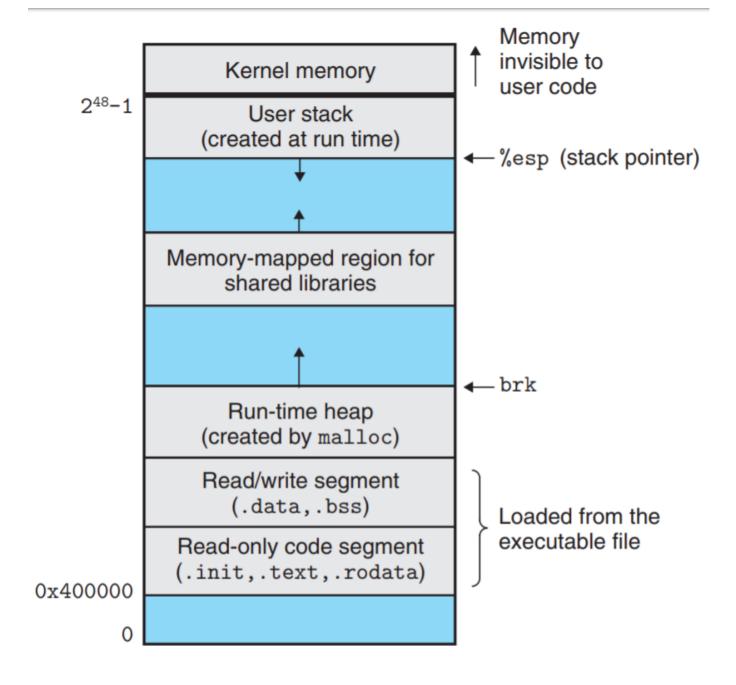


keyword: dynamic linker, rellocation, PIC, GOT, PLT

Preliminary

Unix Process Memory Layout

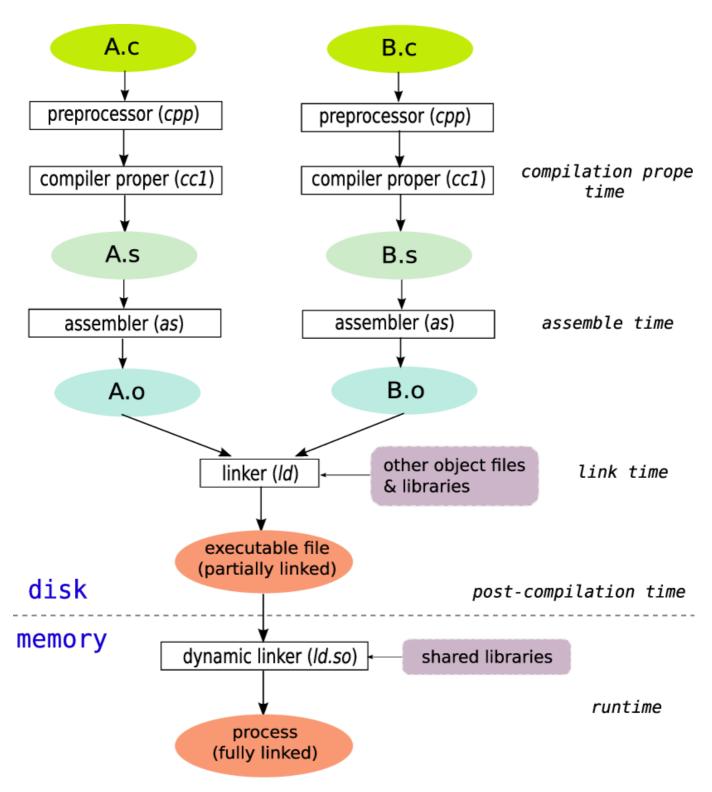
https://www.geeksforgeeks.org/memory-layout-of-c-program/



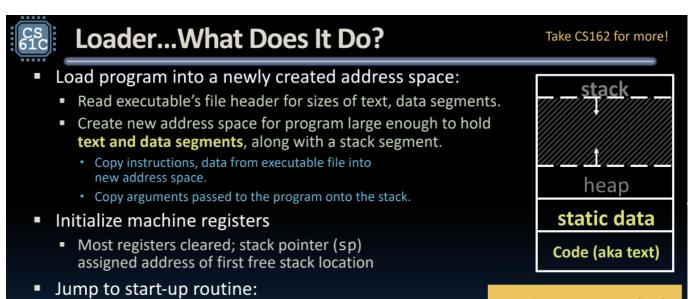
keyword: section, segment, variable, allocation.

Build Process

 ${\tt CS61C:https://drive.google.com/file/d/1Gg1pKejSDy_dUPWdwQ1zuquI0RW8UKUu/view?pli=1}$



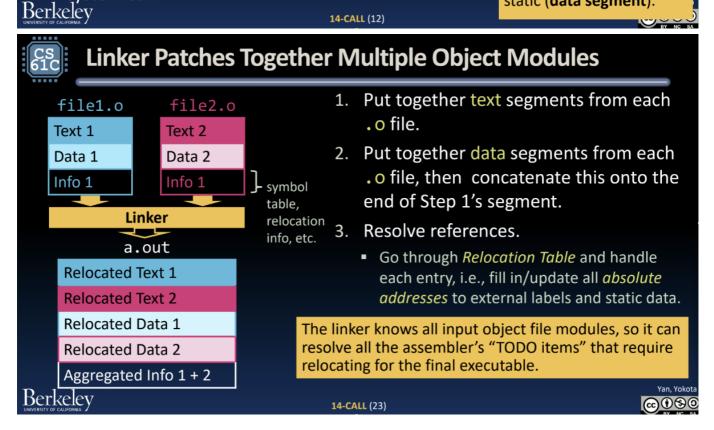
Linker



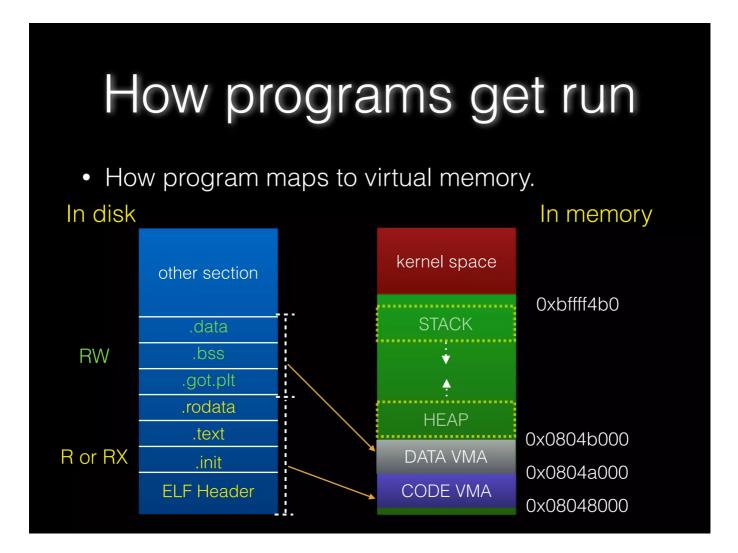
If main routine returns, terminate program with exit system call.

Copy program arguments from stack to registers, set PC

a.out must contain both the machine code itself (text segment) and any static (data segment).



Memory Mapping



Dynamic linking - Code sharing

https://bottomupcs.com/ch09s03.html
https://github.com/jserv/min-dl

https://hackmd.io/@RinHizakura/S1tNAORZv

keyword: virtual memory, shared library

Elf specification

ELF Document:

https://www.cs.cmu.edu/afs/cs/academic/class/15213-f00/docs/elf.pdf

ELF explained:

https://www.youtube.com/watch?v=nC1U1LJQL8o

Linker & Loader:

https://www.cs.cornell.edu/courses/cs3410/2013sp/lecture/15-linkers2-w.pdf

Linking View
ELF header
Program header table
optional
Section 1
Section n
Section header table

Execution View
ELF header
Program header table
Segment 1
Segment 2
• • • •
Section header table
optional

Lazy binding - Startup time reduction

https://stackoverflow.com/questions/63745016/why-use-lazy-binding-for-

 $\verb"position-independent-code-function-calls"$

 $\verb|https://rafaelchen.wordpress.com/2017/09/25/pwn\%e7\%9a\%84\%e4\%bf\%ae\%e7\%85\%89\%$

e4%b9%8b%e8%b7%af-lazy-binding/#more-1244

Position Independent Code

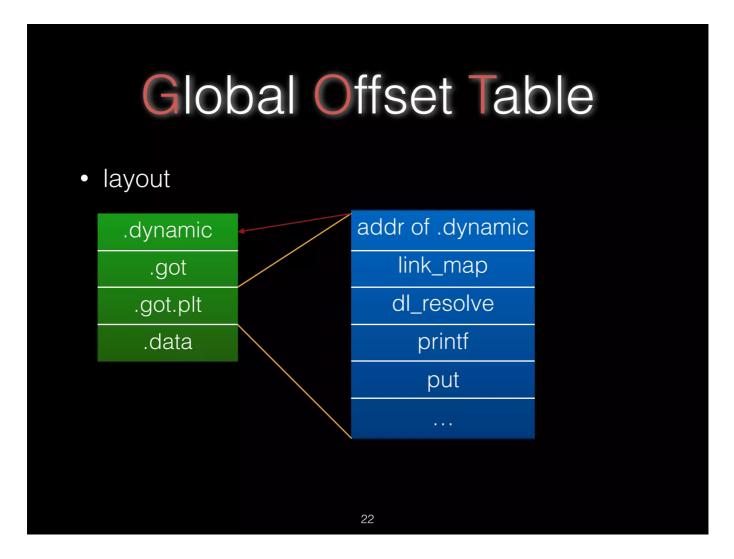
https://docs.oracle.com/cd/E26505_01/html/E26506/glmqp.html

https://bottomupcs.com/ch09s02.html#dynamic_linker_s2

In an executable file, the code and data segment is given a specified base address in virtual memory. The executable code is not shared, and each executable gets its own fresh address space. This means that the compiler knows exactly where the data section will be, and can reference it directly. Consequently all libraries must be produced with code that can execute no matter where it is put into memory, known as position independent code (or PIC for short). Note that the data section is still a fixed offset from the code section; but to actually find the address of data the offset needs to be added to the load address.

.got - Runtime symbol resolve

https://bottomupcs.com/ch09s03.html



_start - The entry point of the program

```
0000000000001050 <_start>:
    1050:
                f3 0f 1e fa
                                          endbr64
    1054:
                31 ed
                                                 %ebp,%ebp
                                          xor
                                                 %rdx,%r9
    1056:
                49 89 d1
                                          mov
    1059:
                                                 %rsi
                5e
                                          pop
    105a:
                48 89 e2
                                          mov
                                                 %rsp,%rdx
                                                 $0xfffffffffffff0,%rsp
    105d:
                48 83 e4 f0
                                          and
    1061:
                50
                                          push
                                                 %rax
    1062:
                 54
                                          push
                                                 %rsp
                                                 %r8d,%r8d
    1063:
                45 31 c0
                                          xor
    1066:
                31 c9
                                                 %ecx, %ecx
                                          xor
    1068:
                48 8d 3d d1 00 00 00
                                          lea
                                                 0xd1(%rip),%rdi
1140 <main>
    106f:
                ff 15 63 2f 00 00
                                          call
                                                 *0x2f63(%rip)
                                                                        # 3fd8
<__libc_start_main@GLIBC_2.34>
    1075:
                f4
                                          hlt
                66 2e 0f 1f 84 00 00
                                          cs nopw 0x0(%rax,%rax,1)
    1076:
    107d:
                00 00 00
```

__libc_start_main - Initialization funtion

```
https://hammertux.github.io/libc-start
https://bottomupcs.com/ch08s08.html#startup
```

Signature of __libc_start_main.

```
int __libc_start_main(int (*main)(int, char *[], char *[]), int argc, char
*argv[], void (*init)(void), void (*fini)(void), void (*rtld_fini)(void),
void *stack_end) {
    // Perform initialization tasks here
    // ...

    // Call the main function
    int result = main(argc, argv, environ);

    // Perform cleanup tasks here
    // ...
    return result;
}
```

keyword: __libc_start_main, .init, .fini, __do_global_ctors_aux

Appendix

Stack allocation padding and alignment

```
https://stackoverflow.com/questions/1061818/stack-allocation-padding-and-alignment
```

keyword: cache block size, data structure alignment, performance.

```
__Attribute__
```

In C, __attribute__ is a keyword-like syntax extension that allows you to provide additional information or directives to the compiler about various attributes of a function, variable, or type.

```
void __attribute__((constructor)) program_init(void) {
   printf("init\n");
}

void __attribute__((destructor)) program_fini(void) {
   printf("fini\n");
}

int main(void)
```

```
{
  return 0;
}

$ gcc -Wall -o test test.c

$ ./test
init
fini
```