

● 1. 考虑右边的三地址代码:

(1) 把这段代码序列划分为基本块

- (1) s = 0
- (2) i = 0
- (3) $t_1 = s \% 19$
- (4) if $t_1 == 1$ qoto (19)
- (5) j = 0
- (6) if i < 50 goto (8)
- (7) return s
- (8) if j >= 100 goto (17)
- (9) if i == j goto (15)

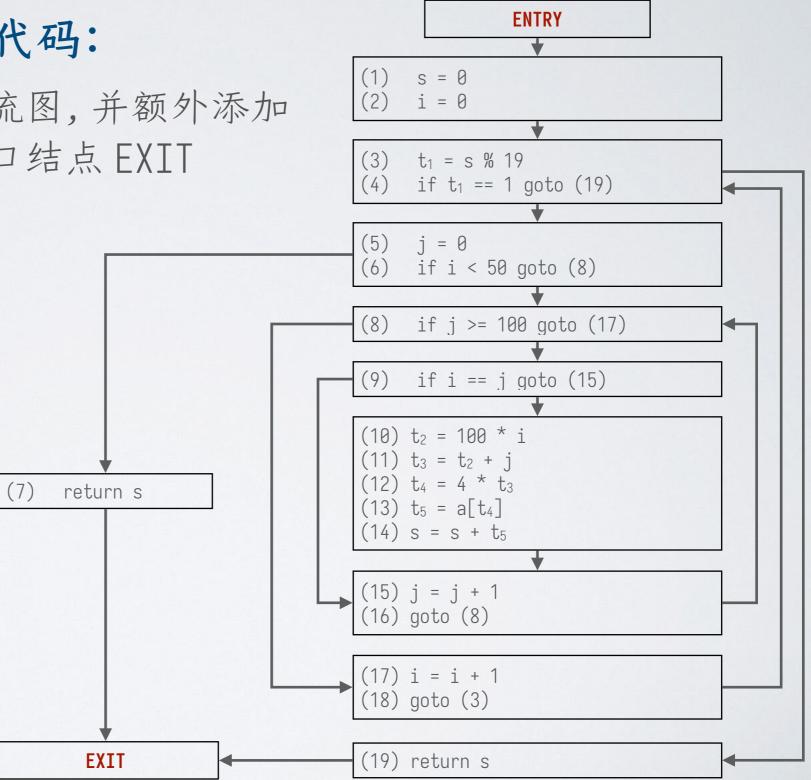
- $(10) t_2 = 100 * i$
- $(11) t_3 = t_2 + j$
- $(12) t_4 = 4 * t_3$
- $(13) t_5 = a[t_4]$
- $(14) s = s + t_5$
- (15) j = j + 1
- (16) goto (8)
- (17) i = i + 1
- (18) goto (3)
- (19) return s

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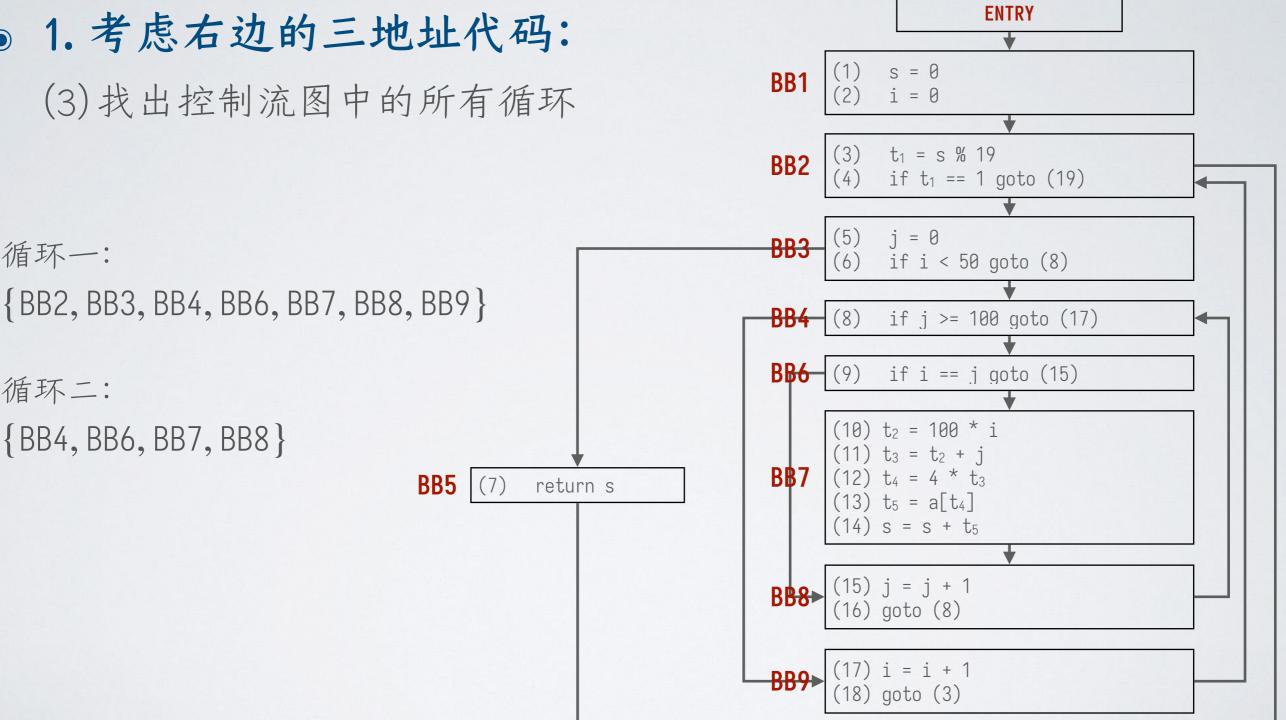
(2)为这段代码构造控制流图,并额外添加入口结点 ENTRY 和出口结点 EXIT





● 1. 考虑右边的三地址代码:

(3) 找出控制流图中的所有循环



BB10 (19) return s

循环二:

循环一:

{BB4, BB6, BB7, BB8}

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EXIT



● 2. 按照课堂上给出的数组引用翻译的 SDT 给出下面语句 S 的三地址代码:

```
a[i^2+j][a[2][i+j][k]+i][i^4+k^2] = i^j+k^a[0][4][2^k];
```

其中数组 a 的类型是 int[5][6][9]。

```
t1 = i * 2

t2 = t1 + j

t3 = i + j

t4 = 2 * 6

t5 = t4 + t3

t6 = t5 * 9

t7 = t6 + k

t8 = t7 * 4

t9 = a [ t8 ]

t10 = t9 + i

t11 = t2 * 6

t12 = t11 + t10

t13 = i * 4

t14 = k * 2

t15 = t13 + t14
```

```
t16 = t12 * 9

t17 = t16 + t15

t18 = t17 * 4

t19 = i * j

t20 = 0 * 6

t21 = t20 + 4

t22 = 2 * k

t23 = t21 * 9

t24 = t23 + t22

t25 = t24 * 4

t26 = a [ t25 ]

t27 = k * t26

t28 = t19 + t27

a [ t18 ] = t28
```

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 \odot 3. 按照课堂上给出的控制流回填翻译的 SDT 给出下面语句 S 的三地址代码和 S . nextlist (指令编号从 100 开始):

```
while ((a + b * c > x + y) && (m == n)) {
  if (x <= y) {
    while (a < b) { a = a + 10; b = c * m; }
  } else a = b + c;
}</pre>
```

```
100: t1 = b * c

101: t2 = a + t1

102: t3 = x + y

103: if t2 > t3 goto 105

104: goto _

105: if m == n goto 107

106: goto _

107: if x <= y goto 109

108: goto 117

109: if a < b goto 111
```

```
110: goto 116
111: t4 = a + 10
112: a = t4
113: t5 = c * m
114: b = t5
115: goto 109
116: goto 119
117: t6 = b + c
118: a = t6
119: goto 100
```

S. nextlist = {104, 106}

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● 4. 考虑一种基于栈的中间表示,它的所有指令都隐式地操作 一个全局的栈,填写以下生成短路求值代码的 SDT。

```
产生规则
                                                                                       语义动作
S \to ID = E;
                         { S.code = E.code || `SET_VAR(\{genvar(ID.lexeme\})`; \}
E \rightarrow E_1 + E_2
                         { E. \text{code} = E_1. \text{code} \parallel E_2. \text{code} \parallel \text{`ADD'}; }
E \rightarrow -E_1
                         { E. code = `CONST(0)` || E_1. code || `SUB`; }
E \to ID
                         { E.code = `GET_VAR(\{genvar(ID.lexeme)\})`; }
B \rightarrow \text{true}
                          { B.code = `GOTO({B.true})`; }
B \rightarrow E_1 == E_2
                         \{B. \text{code} = E_1. \text{code} \mid\mid E_2. \text{code} \mid\mid \text{`SUB'} \mid\mid \text{`GOTO}_{IF}_{ZERO}(\{B. \text{true}\})' \mid\mid \text{`GOTO}(\{B. \text{false}\})'; \}
                          { B_1.true = B.true; B_1.false = genlabel(); }
       B_1 \mid \mid
                          { B_2.true = B.true; B_2.false = B.false; }
       B_2
                          { B. \text{code} = B_1. \text{code} \parallel `\{B_1. \text{false}\}: ` \parallel B_2. \text{code}; }
S \rightarrow \text{if}
                         { B.true = genlabel(); B.false = genlabel(); }
        B)
                         { S_1.next = S.next; }
                         \{ S_2 . next = S . next; \}
       S_1 else
                         \{S. \text{code} = B. \text{code} \parallel `\{B. \text{true}\}: `\parallel S_1. \text{code} \parallel `GOTO(\{S. \text{next}\})` \parallel `\{B. \text{false}\}: `\parallel S_2. \text{code}; \}
       S_2
S \rightarrow \text{while} (
                          { B. true = genlabel(); B. false = S. next; }
                         \{ S_1. next = genlabel(); \}
       B)
                          \{ S. \text{code} = `\{S_1. \text{next}\}: ` || B. \text{code} || `\{B. \text{true}\}: ` || S_1. \text{code} || `goto \{S_1. \text{next}\}`; \}
```

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