编译原理·hw4

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A1

A2

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\label{eq:solution} S \to \text{repeat } \{ \text{ S}_1.\text{next} := \text{newlabel } \} \text{ S}_1 \text{ until} \\ \{ \text{ E.true} := \text{ S}.\text{next}; \text{ E.false} := \text{ S}_1.\text{next } \} \text{ E} \\ \{ \text{ S.code} := \text{gen}(\text{S}_1.\text{next }':') \mid | \text{ S}_1.\text{code } | | \text{ E.code } | | \text{ gen}('\text{goto}' \text{ S}_1.\text{next}) \} \\ \end{cases}
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A3

(a)

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P → D ; { S.inloop := 0 } S
        { P.type := if D.type = ok and S.type = ok then ok else type_error }

S → if E then { S<sub>1</sub>.inloop := S.inloop } S<sub>1</sub>
        { S.type := if E.type = bool then S<sub>1</sub>.type else type_error }

S → while E then {S<sub>1</sub>.inloop := 1 } S<sub>1</sub>
        { S.type := if E.type = bool then S<sub>1</sub>.type else type_error }

S → { S<sub>1</sub>.inloop := S.inloop; S<sub>2</sub>.inloop := S.inloop } S<sub>1</sub> ; S<sub>2</sub>
        { S.type := if S<sub>1</sub>.type = ok and S<sub>2</sub>.type = ok then ok else type_error }

S → continue when E
        { S.type := if E.type = bool and S.inloop = 1 then ok else type_error }
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\begin{array}{lll} P \rightarrow & D \;;\; \{\; S.next := newlabel;\; S.continue := newlabel\; \} \; S \; \{\; gen(S.next \; ':')\; \} \\ S \rightarrow & \text{if} \; \{\; E.true := newlabel;\; E.false := S.next\; \} \; E \; then} \\ & \; \{\; S_1.next := S.next;\; S_1.continue := S.continue\; \} \; S_1 \\ & \; \{\; S.code := E.code\; ||\; gen(E.true \; ':')\; ||\; S_1.code\; \} \\ S \rightarrow & \text{while} \; \{\; E.true := newlabel;\; E.false := S.next\; \} \; E \; do \\ & \; \{\; S_1.next := newlabel;\; S_1.continue := S_1.next\; \} \; S_1 \\ & \; \{\; S.code := gen(S_1.next \; ':')\; ||\; E.code\; ||\; gen(E.true \; ':')\; ||\; S_1.code\; ||\; gen\; ('goto'\; S_1.next)\; \} \\ S \rightarrow & \; \{\; S_1.next := newlabel;\; S_1.countinue := S.continue\; \} \; S_2 \\ & \; \{\; S_2.next := S.next;\; S_2.countinue := S.continue\; \} \; S_2 \\ & \; \{\; S.code := S_1.code\; ||\; gen(S_1.next \; ':')\; ||\; S_2.code\; \} \\ S \rightarrow & \; continue\; when\; \{\; E.true := S.continue;\; E.false := S.next\; \} \\ & \; E \; \{\; S.code := E.code\; \} \end{array}
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A4

A5

(1)

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\begin{array}{lll} A \rightarrow & A_1 + A_2 \; \{ \; A.instr := \; A_2.instr \; || \; A_1.instr \; || \; Plus \; \} \\ \\ A \rightarrow & A_1 - A_2 \; \{ \; A.instr := \; A_2.instr \; || \; A_1.instr \; || \; Minus \; \} \\ \\ A \rightarrow & (A_1) \; \{ \; A.instr := \; A_1.instr \; \} \\ \\ A \rightarrow & \underline{int} \; \{ \; A.instr := \; Push \; \underline{int}.val \; \} \end{array}
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\begin{array}{lll} E \rightarrow & E_1 \text{ if } B \ \{ \text{ E.instr} := E_1.instr \ || \ B.instr \ || \ Cond \ \} \\ A \rightarrow & \underline{id} \ \{ \text{ A.instr} := \text{Load } \underline{id}.\text{val} \} \\ B \rightarrow & A_1 > A_2 \ \{ \text{ B.instr} := A_1.instr \ || \ A_2.instr \ || \ Minus \ || \ Cmp \ || \ Push 1 \ || \ Minus \ \} \\ B \rightarrow & B_1 \ \& \ B_2 \ \{ \text{ B.instr} := B_2.instr \ || \ B_1.instr \ || \ Cond \ \} \\ B \rightarrow & \underline{true} \ \{ \text{ B.instr} := B_1.instr \ || \ Push 1 \ || \ Minus \ \} \\ B \rightarrow & \underline{true} \ \{ \text{ B.instr} := Push 1 \ \} \\ B \rightarrow & \underline{id} := E \ \{ \text{ S.instr} := E.instr \ || \ Store \ \underline{id}.val \ \} \\ S \rightarrow & \underline{id} := E \ \{ \text{ S.instr} := S_1.instr \ || \ S_2.instr \ \} \end{array}
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A6