

Stability of algorithm. (Insertion, Merge, Heap, Bubble)

Insertion Sort:

3 | 5 | 4 | 6 | 8 | 2_A | 2_B

After Insertion Sort

2_A | 2_B | 3 | 4 | 5 | 6 | 8

Insertion Sort
is stable

We just pick an element and place it in the right order, there is no swap between 2 elements, so it's stable sort.

Merge Sort

3 | 1 | 2_A | 2_B | 5 | 4

After heap sort

1 | 1_A | 2_B | 3 | 4 | 5

Merge sort
is stable

While merging 2 separated halves, there is no need to change the original order (L <= R) guarantee that we favor left hand values.

Heap Sort

6 | 4 | 8 | 2 | 2 | 3 | 1 | 5

After

1 | 2 | 2 | 3 | 4 | 5 | 6

It's not stable because operations on the heap can change the original order.

Heap sort pick the largest element and put it in the last of the list
So the element that was picked first stays first
... picked later stays last

Instable

Bubble Sort

6 | 4 | 8 | 2 | 2 | 3 | 1 | 5

After

1 | 2 | 2 | 3 | 4 | 5 | 6 | 8

Again, no swaps can be made if the two elements are equal

(it depends on the condition also,

if there is \leq or $>$ swap

then, the algorithm is no more stable.

Stable

Insertion Sort

Best case complexity : $O(n)$
Worst case complexity : $O(n^2)$

There is a slight change
in the behaviour between the 2 cases.

So, I'll assume it's adaptive.

Merge Sort

Best case complexity : $O(n \log n)$
Worst case complexity : $O(n \log n)$

There is literally no change in behaviour

⇒ Non-adaptive

Heap Sort

Best case complexity : $O(n \log n)$
Worst case complexity : $O(n \log n)$

No change

⇒ Non-adaptive

Bubble Sort:

Best case Complexity: $O(n)$
Worst case Complexity: $O(n^2)$

Change in behaviour
So, it's Adaptive