# Converting the Suggested Upper Merged Ontology to Typed First-order Form

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#### 1 Introduction

Imagine a voice-enabled household robot that can pick up objects and transfer them from room to room. The owner might say "Pick up my red cart and put it in the garage." but the robot might hear either "Pick up my red car and put it in the garage." or the original sentence. A reasoning system might disambiguate the utterance or correct it by reasoning that a typical car weights 1.5 tons and that's 3000 pounds and the carrying capacity of the robot is 100 pounds and the alternate text must be what was said. A system must construct an answer to a question that has never been asked before ("Can the robot carry a car?") perform some simple computation involving unit conversions, to understand that 3000lbs >100lbs and possibly explain its answer if required ("Why didn't you move my car like I asked you to?")

Theorem proving<sup>1</sup> (or Automated Reasoning) is the process by which we can answer a question posed to a theory in a mathematic logic. It differs from search and information retrieval (such as that done with Google or Apache Lucene) in that small facts and rules are combined to synthesize an answer to a question that may never have been asked before, as opposed to matching a query to the most similar previous query and answer. Theorem proving also provides an explanation (a proof) of how the answer was derived. In this work, we use the Suggested Upper Merged Ontology (SUMO) (Niles and Pease, 2001; Pease, 2011) as the theory, or collection of statements in mathematical logic. There are many kinds of mathematical logic. One of the most popular for automated reasoning, providing a good compromise between the expressiveness and efficiency of the language, is called First Order Logic (FOL) (or predicate calculus).

<sup>1</sup>https://plato.stanford.edu/entries/reasoning-automated/

General purpose first-order theorem provers historically haven't done proofs with arithmetic, which is related to problems with Russel's Paradox (Baldwin and Lessmann, 2018). The solution to the paradox, which was discovered by Ernst Zermelo, involved separating numbers and other things into different types. The solution in automated reasoning has now been largely addressed with a new language called Typed First-order Form (TFF) (Sutcliffe et al., 2012) and implemented in several of the best modern provers, including Vampire (Kovács and Voronkov, 2013) from the University of Manchester. Writing a translator from SUMO's native formalization into TFF should open up many new opportunities for doing practical automated reasoning involving numbers and arithmetic.

To have useful and non-trivial reasoning about, for example, a robot's capabilities, or to do question answering, we need not only a language capable of arithmetic calculation (as well as first order logical reasoning) but also a non-trivial body of axioms that has the information about the real world needed to form answers to such questions. For that reason, we need to use the SUMO, which is a comprehensive and diverse set of logical statements about the world. At approximately 20,000 concepts and 80,000 logical statements, it is large enough to answer interesting questions about a wide range of topics.

In earlier work, we described (Pease and Schulz, 2014) how to translate SUMO to the strictly first order language of TPTP (Sutcliffe, 2007). SUMO has an extensive type structure and all relations have type restrictions on their arguments. Translation to TPTP involved implementing a sorted (typed) logic axiomatically in TPTP by altering all implications in SUMO to contain type restrictions on any variables that appear. Many other translation steps were needed that are described in our earlier paper and part of the translation to TFF but since they are not specific to the TFF translation they are not repeated here in detail. Note also that all the strictly higher-order content in SUMO is lost in translation to first-order, whether TPTP or TFF. Briefly however, the translation steps include:

- expanding "row variables" which allow for stating axioms without commitment to the number of arguments a relation has, similar to Lisp's @REST construct
- instantiating "predicate variables" with all possible values. This is needed for any axiom that has a variable in place of a relation.
- expanding the arity of all variable arity relations as set of relations with different names depending upon their fixed number of arguments
- renaming any relations given as arguments to other relations

The interested reader is referred to the earlier paper for more details and examples.

## 2 Typed First-order Form

Like TPTP, TFF forms are valid Prolog syntax (although obviously not the same semantics!) TFF has five disjoint *sorts*: integers, real numbers, rational numbers, booleans and everything else. These are respectively called \$int, \$real, \$rat, \$o and \$i in TFF syntax. Each variable that is used in a logical statement must be declared to be one of these sorts, or by default it will be assumed to be type \$i.

TFF has built in to the language the basic arithmetic functions and arithmetic comparison operators. Each function and operator is *polymorphic* - it is actually a set of three different operators that can handle integers, rationals and reals. Equality is also defined for \$0 and \$i.

TFF's creators have planned to include the ability to define subtypes (subsorts) but this is not yet defined in specification or implemented in any prover. An additional issue is that since all types are disjoint, and SUMO allows multiple inheritance, there is a mismatch between the two type systems. So we have to continue to implement much of SUMO's sort system axiomatically in TFF as in TPTP, but have a special treatment of integers, rationals and reals that does use the TFF type system, so we can use its arithmetic and comparison operators.

## 3 Approach

SUMO's mathematical operators take a generic Quantity type, which then has PhysicalQuantity (a number with units of measure) as well as RealNumber(s) and Integer(s) as subclasses (among others see figure 1). So we have rules like figure 2

```
Quantity
  PhysicalDimension
  Number
    RealNumber
      {\tt Rational Number}
         Integer
           {\tt EvenInteger}
           OddInteger
           PrimeNumber
           NonnegativeInteger
             {\tt PositiveInteger}
           NegativeInteger
      {\tt Irrational Number}
      NonnegativeRealNumber
         {\tt Positive Real Number}
           PositiveInteger
         NegativeRealNumber
           NegativeInteger
      BinaryNumber
    {\tt ImaginaryNumber}
    ComplexNumber
  PhysicalQuantity
    {\tt ConstantQuantity}
    TimeMeasure
      {\tt TimeDuration}
      TimePosition
         {\tt TimeInterval}
```

Figure 1: SUMO's Hierarchy of Quantity Types (somewhat simplified)

```
(=>
  (and
    (instance ?PM ParticulateMatter)
    (part ?Particle ?PM)
    (approximateDiameter ?Particle
        (MeasureFn ?Size Micrometer))
    (greaterThan 10 ?Size)
    (greaterThan ?Size 2.5))
  (instance ?PM CoarseParticulateMatter)))
```

Figure 2: A rule with an arithmetic comparison

where  ${\tt greaterThan}$  is comparing integers or reals, as well as rules like figure 3

```
(=>
  (and
    (instance ?AREA DesertClimateZone)
    (subclass ?MO Month)
    (averageTemperatureForPeriod ?AREA ?MO ?TEMP)
    (greaterThan ?TEMP
        (MeasureFn 18 CelsiusDegree)))
(instance ?AREA SubtropicalDesertClimateZone))
```

Figure 3: A rule comparing numbers with units

where  ${\tt greaterThan}$  is comparing quantities (numbers with units) like "18 degrees".

#### 3.1 Type Promotion

Since TFF doesn't have a type hierarchy, we can't have relations with numbers and quantities (numbers with units) both allowed for their arguments. So the approach is to separate relations that use numbers from those that don't.

Take every relation that has a domain specification of Quantity or a subclass of Quantity. If the domain is a subclass of Integer, then promote it to Integer and add a constraint function to the antecedent. If it's a RationalNumber leave the type as is. If it's any subclass of NonnegativeRealNumber that's also not a subclass of Integer, promote it to RealNumber. All axioms containing SUMO relations that have corresponding TFF relations will require copying

with a native TFF relation if the SUMO relation is a Quantity but not a RationalNumber, RealNumber or Integer.

#### 3.2 Type Adjustment Algorithm

Step-by-step the algorithm is

- 1. Collect the types of all variables
- 2. Constrain all types further if an equality or inequality (an instance of SUMO's RelationExtendedToQuantities) has one argument that is more specific than the stated argument type for the relation, unless it's already a RealNumber (in which case even if the other argument is an Integer, don't constrain it).
- 3. "promote" types that are specializations of Integer or RealNumber as described above
- 4. if a number is an Integer (without a decimal point) and is not a parameter or in a comparison statement with a defined Integer type, add ".0" to it so it's a real and record its type as a RealNumber
- 5. Rename the RelationExtendedToQuantities with a \_Real or \_\_Integer suffix if its arguments have been constrained to those types (or their subclasses)
- 6. If there are any variable types for the axiom at this stage that are below Quantity but not below Number in the SUMO hierarchy, create two version of the axiom one with all the original names of the predicates and one with every SUMO predicate without a TFF equivalent having a \_\_Real suffix and those that do have a TFF equivalent converted to that equivalent
- 7. Translate to TFF syntax, including translating all relations that have a corresponding native TFF relation or arithmetic operator, and have a \_\_Real or \_\_Integer suffix

#### 3.3 Examples of Relevant Axioms

Figure 2 would then become figure 4 and figure 3 becomes figure 5

Figure 4: Axiom in TFF with native comparison operator

Figure 5: Axiom in TFF with SUMO comparison operator

We already have conversion axioms like figure 6

```
(=>
  (equal ?AMOUNT
   (MeasureFn ?X Joule))
  (equal ?AMOUNT
   (MeasureFn
        (MultiplicationFn 0.0002778 ?X) Watt)))
```

Figure 6: Unit conversion axiom

So we need a new axiom to compare quantities for each equality and inequality, e.g. figure 7

```
(=>
  (and
    (equal ?Q1 (MeasureFn ?I1 ?U))
    (equal ?Q2 (MeasureFn ?I2 ?U))
    (greaterThan ?I1 ?I2))
  (greaterThan ?Q1 ?Q2))
```

Figure 7: Unit comparison axiom

but we can do this more succinctly with predicate variables as in figure 8

```
(=>
  (and
    (instance ?REL RelationExtendedToQuantities)
    (equal ?Q1 (MeasureFn ?I1 ?U))
    (equal ?Q2 (MeasureFn ?I2 ?U))
    (?REL ?I1 ?I2))
  (?REL ?Q1 ?Q2))
```

Figure 8: Unit comparison axiom for all operators

which will create the preceding axiom plus all other relevant ones (figure 9).

```
(instance equal RelationExtendedToQuantities)
(instance greaterThan RelationExtendedToQuantities)
(instance greaterThanOrEqualTo RelationExtendedToQuantities)
(instance lessThan RelationExtendedToQuantities)
(instance lessThanOrEqualTo RelationExtendedToQuantities)
```

Figure 9: All relations that take numbers and quantities

The axiom in figure 8 plus the relation statements from figure 9 will result in the axiom in figure 7 (and many others) being generated by the predicate variable instantiation process in Sigma. Then when SUMO's greaterThan appears with strictly integer or real arguments the subsequent TFF conversion step in Sigma will convert it to \$greater.

When caling up to processing all of SUMO to TFF, rather than a subset, we need to treat numeric type polymorphism as a "macro" - every statement that uses TFF's built-in math operators.

#### 3.4 Detailed Example of the Algorithm

Let's take the axiom in figure 10 and the argument type specification for the functions it contains as the start of the step-by-step translation

```
(equal
    (RemainderFn ?NUMBER1 ?NUMBER2)
    ?NUMBER)
    (equal
        (SignumFn ?NUMBER2)
        (SignumFn ?NUMBER)))

(domain RemainderFn 1 Quantity)
(domain RemainderFn 2 Quantity)
(range RemainderFn Quantity)

(domain SignumFn 1 RealNumber)
(range SignumFn Integer)
```

Figure 10: Axiom containing functions and their type signatures

For Step 1, the current Sigma translation will know to constrain each type to its most specific value in the axiom, which in this case means that ?NUMBER and ?NUMBER2, since they are both arguments to SignumFn, must be RealNumber, since SignumFn takes one argument that must be a RealNumber - (domain SignumFn 1 RealNumber). We'll have a list of [?NUMBER=RealNumber, ?NUMBER1=Quantity, ?NUMBER2=RealNumber]

For Step 2 we do have the RelationExtendedToQuantities RemainderFn as well as equal where we must constrain its argument to be the same lowest common type, which among [?NUMBER1=Quantity, ?NUMBER2=RealNumber] is RealNumber. So the list of variables and their types becomes [?NUMBER=RealNumber, ?NUMBER1=RealNumber, ?NUMBER2=RealNumber]. Step 3 doesn't apply since the variable list contains no specializations of Integer or RealNumber. Step 4 doesn't apply since we have no literal numbers. Since we've modified the arguments of RemainderFn we need to rename it in step 5 as RemainderFn\_Real. Then we have to look at equal. Note that we have to keep track of a modified return type for every function that returns some numerical type. In this case, although RemainderFn returns a Quantity,

its arguments have already constrained it to return a RealNumber. So we have to check that the arguments to equal are the same and constrain them if needed. In this case, ?NUMBER is already a RealNumber and doesn't need to be constrained further.

Step 6 doesn't apply since there are no variables of type Number. Now we have figure 11

```
(=>
  (equal
    (RemainderFn_Real ?NUMBER1 ?NUMBER2)
    ?NUMBER)
  (equal
    (SignumFn_Real ?NUMBER2)
    (SignumFn_Real ?NUMBER)))
```

Figure 11: Expansion of function names

The last step is to translate into TFF syntax and convert relations and functions to their native TFF names, when applicable. This results in figure 12.

Figure 12: Axiom in TFF with native comparison operator

Let's take an example axiom that maps to some TFF built in functions or relations. In the case of greaterThan, we may have SUMO axioms in which it is used with units, and cases where it is used without. The cases without units will map to TFF's built in \$greater relation, as in figure 13

```
(=>
  (and
    (instance ?X ?Y)
    (subclass ?Y PureSubstance)
    (boilingPoint ?Y
        (MeasureFn ?BOIL KelvinDegree))
    (meltingPoint ?Y
        (MeasureFn ?MELT KelvinDegree))
    (measure ?X
        (MeasureFn ?TEMP KelvinDegree))
    (greaterThan ?TEMP ?MELT)
    (lessThan ?TEMP ?BOIL))
    (attribute ?X Liquid))
```

Figure 13: Another rule with numeric comparisons

Step 1 results in the following where a '+' signifies a class rather than an instance - [?X=Quantity, ?Y=PureSubstance+, ?BOIL=RealNumber, ?TEMP=RealNumber, ?MELT=RealNumber]. For Step 2 we have greaterThan and lessThan but all are RealNumber so no change is needed. Step 3 doesn't apply. For Step 4 there are no numbers so that doesn't apply. Step 5 we have greaterThan and lessThan, both have RealNumber as arguments and so will get renamed to greaterThan\_Real and lessThan\_Real respectively, as shown in Figure 14. Step 6 doesn't apply. For Step 7 we rename greaterThan\_Real to \$greater and lessThan\_Real to \$less and translate to TFF syntax, resulting in Figure 15.

```
(=>
  (and
    (instance ?X ?Y)
    (subclass ?Y PureSubstance)
    (boilingPoint ?Y
        (MeasureFn ?BOIL KelvinDegree))
        (meltingPoint ?Y
            (MeasureFn ?MELT KelvinDegree))
        (measure ?X
            (MeasureFn ?TEMP KelvinDegree))
        (greaterThan_Real ?TEMP ?MELT)
        (lessThan_Real ?TEMP ?BOIL))
        (attribute ?X Liquid))
```

Figure 14: Rule with numeric comparisons after renaming with Real

Figure 15: Rule in TFF with numeric comparisons after renaming operators

Another case is figure 16

```
(<=>
  (and
    (instance ?LD LiquidDrop)
    (approximateDiameter ?LD
        (MeasureFn ?Size Micrometer))
    (lessThan 500 ?Size))
  (instance ?LD Droplet))
```

Figure 16: Axiom with literal numbers

For Step 1 we have [?LD=Droplet,?Size=RealNumber]. For step 2, since ?Size is already a RealNumber it doesn't need to be changed to be a RealNumber like '500.0' given that they are both arguments to lessThan. Step 3 doesn't apply. For Step 4 we promote '500' to be '500.0'. In Step 5 we change lessThan to lessThan\_Real to conform to its argument types. Step 6 doesn't apply. For step 7 we make lessThan\_Real into \$less and do the translation to TFF syntax that results in figure 17.

Figure 17: TFF Axiom with literal numbers

```
(=>
   (instance ?X NegativeInteger)
   (greaterThan 0 ?X))
```

Figure 18: Axiom with NegativeInteger

For Figure 18 the transformation steps are, in order

- 1. [?X=NegativeInteger]
- 2. Doesn't apply
- 3. Promote NegativeInteger to Integer, so we have [?X=Integer]
- 4. Doesn't apply
- 5. We make greaterThan into greaterThan\_Integer
- 6. Doesn't apply
- 7. TFF syntax translation results in figure 19

```
! [V_X : $int] :
    (s__instance(V_X,s__NegativeInteger) =>
    $greater(0,V_X))
```

Figure 19: TFF Axiom with NegativeInteger

But this doesn't work since instance takes \$i types and not \$int which is the type of V\_X.

## 4 Modification to the Translation Algorithm

In effect we need to suppress all axioms involving types below Integer and RealNumber, but only if they are the definitions of those terms! Otherwise, we need to have the definitions become restrictions on integers and reals. So we build a table of definitional constraints for all types below Integer and RealNumber. This will consist of figure 20 and 21.

```
(=>
    (instance ?NUMBER EvenInteger)
    (equal
        (RemainderFn ?NUMBER 2) 0))
(=>
    (instance ?NUMBER OddInteger)
    (equal
        (RemainderFn ?NUMBER 2) 1))
(=>
    (instance ?PRIME PrimeNumber)
    (forall (?NUMBER)
        (=>
            (equal
                (RemainderFn ?PRIME ?NUMBER) 0)
            (or
                (equal ?NUMBER 1)
                (equal ?NUMBER ?PRIME)))))
(=>
    (instance ?X NonnegativeInteger)
    (greaterThan ?X -1))
(=>
    (instance ?X NegativeInteger)
    (greaterThan 0 ?X))
(=>
    (instance ?X PositiveInteger)
    (greaterThan ?X 0))
(<=>
    (instance ?NUMBER PositiveRealNumber)
    (and
        (greaterThan ?NUMBER 0)
        (instance ?NUMBER RealNumber)))
(=>
    (instance ?NUMBER PositiveRealNumber)
    (equal
        (SignumFn ?NUMBER) 1))
```

Figure 20: Number definitions

```
(<=>
    (instance ?NUMBER NegativeRealNumber)
    (and
        (lessThan ?NUMBER 0)
        (instance ?NUMBER RealNumber)))
(=>
    (instance ?NUMBER NegativeRealNumber)
    (equal
        (SignumFn ?NUMBER) -1))
(<=>
    (instance ?NUMBER NonnegativeRealNumber)
    (and
        (greaterThanOrEqualTo ?NUMBER 0)
        (instance ?NUMBER RealNumber)))
(=>
    (instance ?NUMBER NonnegativeRealNumber)
        (equal
            (SignumFn ?NUMBER) 1)
        (equal
            (SignumFn ?NUMBER) 0)))
```

Figure 21: Number definitions (continued)

That leaves some types of numbers that have no axioms other than definitions and which are not used as argument types: IrrationalNumber, BinaryNumber, ImaginaryNumber and ComplexNumber. We will exclude these from translation, at least for now. Note also that there are two subtypes that inherit from RealNumber and Integer. They are PositiveInteger, which inherits from PositiveRealNumber and NegativeInteger, which inherits from NegativeRealNumber. We will come back to this issue later.

We reformulate some of the axioms to be consistent, removing the bi-implications and the type assertions when they occur on both sides. This gives us a few revised axioms as shown in figure 22.

```
(=>
    (instance ?NUMBER PositiveRealNumber)
    (greaterThan ?NUMBER 0))

(=>
    (instance ?NUMBER NegativeRealNumber)
    (lessThan ?NUMBER 0))

(=>
    (instance ?NUMBER NonnegativeRealNumber)
    (greaterThanOrEqualTo ?NUMBER 0))
```

Figure 22: Revised number definitions

We preprocess all axioms looking for antecedents that specify an instance of a Number then cache the consequent in a map indexed by the type. We also extract the variable name that constrains the type in those antecedents. We wind up with the list of antecedents and variables, respectively, shown in figure 23

```
{PositiveInteger =
   (greaterThan ?X 0),
OddInteger =
   (equal (RemainderFn ?NUMBER 2) 1),
NegativeRealNumber =
   (equal (SignumFn ?NUMBER) −1),
EvenInteger =
   (equal (RemainderFn ?NUMBER 2) 0),
NonnegativeInteger =
   (greaterThan ?X -1),
NegativeInteger =
   (greaterThan 0 ?X),
PrimeNumber =
   (forall (?NUMBER)
       (equal (RemainderFn ?PRIME ?NUMBER) 0)
       (or
         (equal ?NUMBER 1)
         (equal ?NUMBER ?PRIME)))),
RationalNumber =
   (exists (?INT1 ?INT2)
     (and
       (instance ?INT1 Integer)
       (instance ?INT2 Integer)
       (equal ?NUMBER (DivisionFn ?INT1 ?INT2)))),
NonnegativeRealNumber =
   (or
     (equal (SignumFn ?NUMBER) 1)
     (equal (SignumFn ?NUMBER) 0))}
{PositiveInteger =
                         Χ,
OddInteger =
                         NUMBER,
NegativeRealNumber =
                        NUMBER,
EvenInteger =
                         NUMBER,
NonnegativeInteger =
                         Х,
NegativeInteger =
                         Х,
PrimeNumber =
                         PRIME,
RationalNumber =
                         NUMBER,
NonnegativeRealNumber = NUMBER}
```

Figure 23: All number type conditions

For example, we would have [PositiveInteger=(greaterThan ?NUMBER 0)]. The value of the index will be added as an antecedent to any axiom that has that type for a variable. For example, see figure 24 where the first axiom becomes the second, and the resulting TFF version is the third.

Figure 24: Adding a number type condition

We have a different substitution process for appearances of number subtypes in rule consequents. If there is an instance statement involving a number subtype in the consequent, then replace it with the defining condition for that sybtype and all its parents from the list in figure 23. For example, see figure 25, in which a variable is concluded to be a PositiveRealNumber and must therefore be replaced with its defining condition

Figure 25: Adding a number type in a consequent

## 5 Type Propagation from Functions

Our translation algorithm is unfortunately still incomplete, as described so far. Type changes can also propagate throughout a formula because of functions. Take Figure 26 that has an axiom with several nested functions, which in turn give the variables the types below it. But the return types of the functions should also constrain the types of the variables.

Figure 26: Nested Functions

The return types of the functions are shown in Figure 27. While only FloorFn has types more restricted than Quantity, these will actually propagate through to constrain the entire set of variables.

```
(domain RemainderFn 1 Quantity)
(domain RemainderFn 2 Quantity)
(range RemainderFn Quantity)

(domain AdditionFn 1 Quantity)
(domain AdditionFn 2 Quantity)

(domain MultiplicationFn 1 Quantity)
(domain MultiplicationFn 2 Quantity)

(domain FloorFn 1 RealNumber)
(range FloorFn Integer)

(domain DivisionFn 1 Quantity)
(domain DivisionFn 2 Quantity)
```

Figure 27: Function Argument and Return Types

FloorFn makes the return type of DivisionFn a RealNumber, which then makes ?NUMBER1 and ?NUMBER2 RealNumbers but since ?NUMBER2 is an Integer

?NUMBER1 must also be an Integer and the return type of DivisionFn becomes Integer, meaning we need to rename as DivisionFn\_Integer and then FloorFn must become FloorFn\_Integer. Since MultiplicationFn\_Integer returns an Integer, that constrains ?NUMBER also to be an Integer. RemainderFn then also becomes RemainderFn\_Integer. The result is the variable set and TFF translation shown in Figure 28.

```
{?NUMBER1=[Integer], ?NUMBER2=[Integer], ?NUMBER=[Integer]}
! [V__NUMBER1 : $int,V__NUMBER2 : $int,V__NUMBER : $int] :
  ((s__RemainderFn__OIn1In2In(V__NUMBER1, V__NUMBER2) =
  V__NUMBER =>
     $sum(
       $product(
         s__FloorFn__1In(
           $quotient_e(V__NUMBER1, V__NUMBER2)),
         V__NUMBER2),
       V__NUMBER) =
       V__NUMBER1) &
  ($sum(
   $product(
      s__FloorFn__1In(
        $quotient_e(V__NUMBER1, V__NUMBER2)) ,
      V__NUMBER2),
   V_{NUMBER} =
   V__NUMBER1 =>
      s__RemainderFn__OIn1In2In(V__NUMBER1, V__NUMBER2) =
        V__NUMBER))
```

Figure 28: Variable Type Propagation Result

We then also need to copy all the axioms on the original term over to the new term with the modified suffix.

Every relation that has an argument type between Quantity and Integer will need and additional Integer version. Those with arguments between Quantity and RealNumber will need Integer and RationalNumber versions. Those with arguments of Quantity, PhysicalDimension or Number will need all three - Integer, RealNumber and RationalNumber as well as their original version.

#### 5.1 Sort Conflict

Inequalities in TFF must be used with a single sort. However, in SUMO we have the case of figure 29

```
(=>
  (equal
    (CeilingFn ?NUMBER)
    ?INT)
  (not
    (exists (?OTHERINT)
        (and
          (instance ?OTHERINT Integer)
          (greaterThanOrEqualTo ?OTHERINT ?NUMBER)
          (lessThan ?OTHERINT ?INT)))))
```

Figure 29: Type Conflict

CeilingFn takes a RealNumber and returns the smallest Integer greater than or equal to the argument. The greaterThanOrEqualTo must compare the RealNumber and Integer but that is not allowed in TFF.

We would like to be able to state the following in figure 30

Figure 30: Type ConflictTFF

but that generates an error due to the arguments to **\$greatereq** being of different types. A similar problem exists in SUMO's axiom for FloorFn. We can use TFF's type coercion function to solve this, as in figure 31

```
(! [V__NUMBER : $real,V__INT : $int] :
    (s__CeilingFn__OIn1ReFn(V__NUMBER) = V__INT =>
    ~? [V__OTHERINT:$int] :
    ($greatereq(V__OTHERINT ,$to_real(V__NUMBER)) &
    $less(V__OTHERINT ,V__INT))))).
```

Figure 31: Eliminating Type ConflictTFF with type coercion

Alternatively, for these particular axioms we could use the built-in TFF \$ceiling and \$floor functions, and then even eliminate these axioms entirely since they help define the these TFF built-in functions.

Another version of this problem occurs when mixing Quantity(s) and basic numeric types, as with

```
(=>
  (diameter ?CIRCLE ?LENGTH)
  (exists (?HALF)
    (and
        (radius ?CIRCLE ?HALF)
        (equal (MultiplicationFn ?HALF 2) ?LENGTH))))
```

Figure 32: Type ConflictTFForig

We'd like to multiply, for example, "two feet" by 2 and get "four feet". This is equivalent to the following -  $\,$ 

Figure 33: Type ConflictTFFmod

where the numeric part of the quantity with units is multiplied by the numeric argument.

#### 6 Conclusion

We hope that the extensive list of practical challenges in translating a large and comprehensive theory to TFF will help motivate future version of the TFF that may make theorem proving with arithmetic more straightforward. In particular, having a type hierarchy would greatly simplify the translation steps that are needed.

The software for this translation is implemented in Java as part of the Sigma Knowledge Engineering Environment, available with the Suggested Upper Merged Ontology (SUMO) on GitHub at https://github.com/ontologyportal. That site also has an expanded technical report providing additional translation details that space doesn't permit including in this paper.

### 7 Appendix: Sample Proof

Here we show a sample proof with arithmetic of the sample problem from the introduction. We assert that a robot can only carry objects of less than 100 pounds (axiom f1) and then assert existence of an event where it carries a 1.5 ton object (axioms f2 and f4). We pose this query to the Vampire theorem prover. We see that in step f71 it applies the unit conversion axiom stated in f8 to convert 100 pounds to 0.05 tons and that 1.5 is not less than 0.05. This yields a contradiction.

```
% Refutation found. Thanks to Tanya!
% SZS status Unsatisfiable for Robot-small
% SZS output start Proof for Robot-small
tff(func_def_0, type, s__MeasureFn__1ReFn: ($real * $i) > $i).
tff(f1,axiom,(
  s__maximumPayloadCapacity(s__Robot1,
  s_MeasureFn_1ReFn(100.0,s_PoundMass))),
  file('/home/apease/.sigmakee/KBs/Robot-small.tff',kb_SUMO_3179)).
tff(f2,axiom,(
  s__measure(s__MyCar,s__MeasureFn__1ReFn(1.5,s__TonMass))),
  file('/home/apease/.sigmakee/KBs/Robot-small.tff',kb_SUMO_3180)).
tff(f3,axiom,(
  ! [X0,X1,X2,X3] : ((greaterThan(X2,X3) &
    s__maximumPayloadCapacity(X1,X3) &
    s__measure(X0,X2)) =>
  ~? [X4] : (s_instance(X4,s_Carrying) &
             s_patient(X4,X0) &
             s_{\text{instrument}}(X4,X1))),
  file('/home/apease/.sigmakee/KBs/Robot-small.tff',kb_SUMO_268)).
```

```
tff(f4,axiom,(
  s__instance(s__Carry1,s__Carrying) &
  s__patient(s__Carry1,s__MyCar) &
  s__instrument(s__Carry1,s__Robot1)),
  file('/home/apease/.sigmakee/KBs/Robot-small.tff',assert)).
tff(f8,axiom,(
  ! [X3 : $real] : s__MeasureFn__1ReFn(X3,s__PoundMass) =
    s__MeasureFn__1ReFn($quotient(X3,2000.0),s__TonMass)),
  file('/home/apease/.sigmakee/KBs/Robot-small.tff',kb_SUM0_3175b2)).
tff(f9,axiom,(
  ! [X6,X7 : $real,X8 : $real] : ($greater(X7,X8) =>
    greaterThan(s__MeasureFn__1ReFn(X7,X6),s__MeasureFn__1ReFn(X8,X6)))),
  file('/home/apease/.sigmakee/KBs/Robot-small.tff',kb_SUMO_228b)).
tff(f10,axiom,(
  ! [X6,X7 : $real,X8 : $real] :
    (greaterThan(s__MeasureFn__1ReFn(X7,X6),s__MeasureFn__1ReFn(X8,X6)) =>
    $greater(X7,X8))),
  file('/home/apease/.sigmakee/KBs/Robot-small.tff',kb_SUMO_228c)).
tff(f11,plain,(
  ! [X6,X7 : $real,X8 : $real] : ($less(X8,X7) =>
    greaterThan(s__MeasureFn__1ReFn(X7,X6),s__MeasureFn__1ReFn(X8,X6)))),
  inference(evaluation,[],[f9])).
tff(f12,plain,(
  ! [X6,X7 : $real,X8 : $real] :
    (greaterThan(s__MeasureFn__1ReFn(X7,X6),s__MeasureFn__1ReFn(X8,X6)) =>
    $less(X8,X7))),
  inference(evaluation,[],[f10])).
tff(f14,plain,(
  ! [X0 : $real] : s__MeasureFn__1ReFn(X0,s__PoundMass) =
    s__MeasureFn__1ReFn($quotient(X0,2000.0),s__TonMass)),
  inference(rectify,[],[f8])).
tff(f16,plain,(
  ! [X0,X1 : $real,X2 : $real] : ($less(X2,X1) =>
    greaterThan(s__MeasureFn__1ReFn(X1,X0),s__MeasureFn__1ReFn(X2,X0)))),
  inference(rectify,[],[f11])).
tff(f17,plain,(
  ! [X0,X1 : $real,X2 : $real] :
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(greaterThan(s__MeasureFn__1ReFn(X1,X0),s__MeasureFn__1ReFn(X2,X0)) =>
    $less(X2,X1))),
  inference(rectify,[],[f12])).
tff(f20,plain,(
  ! [X0,X1 : $real,X2 : $real] :
    (greaterThan(s__MeasureFn__1ReFn(X1,X0),s__MeasureFn__1ReFn(X2,X0)) |
    ~$less(X2,X1))),
  inference(ennf_transformation,[],[f16])).
tff(f21,plain,(
  ! [X0,X1 : $real,X2 : $real] : ($less(X2,X1) |
    ~greaterThan(s__MeasureFn__1ReFn(X1,X0),s__MeasureFn__1ReFn(X2,X0)))),
  inference(ennf_transformation,[],[f17])).
tff(f22,plain,(
  ! [X0,X1,X2,X3] : (! [X4] : (~s__instance(X4,s__Carrying) |
    "s__patient(X4,X0) | "s__instrument(X4,X1)) | ("greaterThan(X2,X3) |
    "s__maximumPayloadCapacity(X1,X3) | "s__measure(X0,X2)))),
  inference(ennf_transformation,[],[f3])).
tff(f23,plain,(
  ! [X0,X1,X2,X3] : (! [X4] : (~s__instance(X4,s__Carrying) |
"s_patient(X4,X0) | "s_instrument(X4,X1)) | "greaterThan(X2,X3) |
"s__maximumPayloadCapacity(X1,X3) | "s__measure(X0,X2))),
  inference(flattening,[],[f22])).
tff(f25,plain,(
  s_maximumPayloadCapacity(s_Robot1,s_MeasureFn_1ReFn(100.0,s_PoundMass))),
  inference(cnf_transformation,[],[f1])).
tff(f26,plain,(
  s_measure(s_MyCar,s_MeasureFn_1ReFn(1.5,s_TonMass))),
  inference(cnf_transformation,[],[f2])).
tff(f27,plain,(
  s__instrument(s__Carry1,s__Robot1)),
  inference(cnf_transformation,[],[f4])).
tff(f28,plain,(
  s__patient(s__Carry1,s__MyCar)),
  inference(cnf_transformation,[],[f4])).
tff(f29,plain,(
  s__instance(s__Carry1,s__Carrying)),
  inference(cnf_transformation,[],[f4])).
```

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tff(f30,plain,(
  ( ! [XO:$real] : (s_MeasureFn_1ReFn(XO,s_PoundMass) =
    s__MeasureFn__1ReFn($quotient(X0,2000.0),s__TonMass)) )),
  inference(cnf_transformation,[],[f14])).
tff(f32,plain,(
  (! [X2:$real,X0,X1:$real] :
    (greaterThan(s__MeasureFn__1ReFn(X1,X0),s__MeasureFn__1ReFn(X2,X0)) |
    ~$less(X2,X1)) )),
  inference(cnf_transformation,[],[f20])).
tff(f33,plain,(
  (! [X2:$real,X0,X1:$real] : ($less(X2,X1) |
    ~greaterThan(s__MeasureFn__1ReFn(X1,X0),s__MeasureFn__1ReFn(X2,X0))) )),
  inference(cnf_transformation,[],[f21])).
tff(f34,plain,(
  (![X4,X2,X0,X3,X1]:
    (~s__instance(X4,s__Carrying) |
    ~s__patient(X4,X0) |
    ~s__instrument(X4,X1) |
    ~greaterThan(X2,X3) |
    ~s__maximumPayloadCapacity(X1,X3) |
    ~s__measure(X0,X2)) )),
  inference(cnf_transformation,[],[f23])).
tff(f35,plain,(
  (![X4,X0,X1]:
    (~s__instance(X4,s__Carrying) |
    ~s__patient(X4,X0) |
    ~s__instrument(X4,X1) |
    sPO(X1,X0))),
  inference(cnf_transformation,[],[f35_D])).
tff(f35_D,plain,(
  (![X0,X1]:((![X4]:
    (~s__instance(X4,s__Carrying) |
     ~s__patient(X4,X0) |
     ~s__instrument(X4,X1)) ) <=>
    ~sP0(X1,X0)) )),
  introduced(general_splitting_component_introduction,[new_symbols(naming,[sP0])])).
tff(f36,plain,(
  (![X2,X0,X3,X1]:
    (~greaterThan(X2,X3) |
```

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~s__maximumPayloadCapacity(X1,X3) |
     ~s__measure(X0,X2) |
     ~sP0(X1,X0)) )),
  inference(general_splitting,[],[f34,f35_D])).
tff(f37,plain,(
  (![X2,X0,X3]:
    (sP1(X3,X0) |
     ~s__measure(X0,X2) |
     ~greaterThan(X2,X3)) )),
  inference(cnf_transformation,[],[f37_D])).
tff(f37_D,plain,(
  (![X0,X3]:((![X2]:
    (~s__measure(X0,X2) |
     ~greaterThan(X2,X3)) ) <=>
    ~sP1(X3,X0)) )),
  introduced(general_splitting_component_introduction,[new_symbols(naming,[sP1])])).
tff(f38,plain,(
  (![X0,X3,X1]:
    (~sP0(X1,X0) |
     ~s__maximumPayloadCapacity(X1,X3) |
     ~sP1(X3,X0)) )),
  inference(general_splitting,[],[f36,f37_D])).
tff(f39,plain,(
  (![X0,X1]:
    (~s__patient(s__Carry1,X0) |
     ~s__instrument(s__Carry1,X1) |
     sPO(X1,X0)) )),
  inference(resolution,[],[f35,f29])).
tff(f40,plain,(
  (! [X0:$real,X1:$real]:
    (~$less($quotient(X0,2000.0),X1) |
     greaterThan(s__MeasureFn__1ReFn(X1,s__TonMass),
                 s__MeasureFn__1ReFn(X0,s__PoundMass))) )),
  inference(superposition,[],[f32,f30])).
tff(f44,plain,(
  (! [XO] :
    (sPO(XO,s__MyCar) |
     ~s__instrument(s__Carry1,X0)) )),
  inference(resolution,[],[f39,f28])).
```

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tff(f45,plain,(
  (! [X2,X0:$real,X1:$real] :
    (greaterThan(s_MeasureFn_1ReFn(X0,s_TonMass),
                 s__MeasureFn__1ReFn(X1,s__PoundMass)) |
    ~greaterThan(s__MeasureFn__1ReFn(X0,X2),
                 s__MeasureFn__1ReFn($quotient(X1,2000.0),X2))) )),
  inference(resolution,[],[f40,f33])).
tff(f47,plain,(
  ( ! [X0,X1] :
    (~sP1(X1,s__MyCar) |
     ~s__maximumPayloadCapacity(X0,X1) |
     ~s__instrument(s__Carry1,X0)) )),
  inference(resolution,[],[f44,f38])).
tff(f50,plain,(
  (![X2,X0,X1]:
    (~s__instrument(s__Carry1,X0) |
     ~s__maximumPayloadCapacity(X0,X1) |
     ~s__measure(s__MyCar,X2) |
     ~greaterThan(X2,X1)) )),
  inference(resolution, [], [f47, f37])).
tff(f55,plain,(
  (![X0,X1]:
    (~s__maximumPayloadCapacity(s__Robot1,X0) |
     ~s__measure(s__MyCar,X1) |
     ~greaterThan(X1,X0)) )),
  inference(resolution,[],[f50,f27])).
tff(f60,plain,(
  (! [XO]:
    (~s__measure(s__MyCar,X0) |
     ~greaterThan(X0,s__MeasureFn__1ReFn(100.0,s__PoundMass))) )),
  inference(resolution,[],[f55,f25])).
tff(f65,plain,(
  ~greaterThan(s__MeasureFn__1ReFn(1.5,s__TonMass),
               s__MeasureFn__1ReFn(100.0,s__PoundMass))),
  inference(resolution,[],[f60,f26])).
tff(f71,plain,(
  (! [X1] : (~greaterThan(s__MeasureFn__1ReFn(1.5,X1),
                           s__MeasureFn__1ReFn($quotient(100.0,2000.0),X1)))),
  inference(resolution,[],[f65,f45])).
```

```
tff(f72,plain,(
 (! [X1] : (~greaterThan(s__MeasureFn__1ReFn(1.5,X1),
                        s__MeasureFn__1ReFn(0.05,X1))) )),
 inference(evaluation,[],[f71])).
tff(f78,plain,(
 ~$less(0.05,1.5)),
 inference(resolution,[],[f72,f32])).
tff(f95,plain,(
 $false),
 inference(evaluation,[],[f78])).
\% SZS output end Proof for Robot-small
% -----
\% Version: Vampire 4.2.2 (commit 6588b35 on 2018-07-19 13:39:17 +0200)
% Termination reason: Refutation
% Memory used [KB]: 9722
% Time elapsed: 0.003 s
% -----
% Success in time 99.856 s
```

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